# Advanced Mac OS X Rootkits

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#### Overview



- Mac OS X and Mach
- Why use Mach for rootkits?
- User-mode Mach rootkit techniques
- Kernel Mach rootkit techniques

## Why Mach Rootkits?



- Traditional Unix rootkit techniques are well understood
- Mach functionality is more obscure
- Rootkits using obscure functionality are less likely to be detected or noticed
- Mach is fun to program

#### Introduction to Mach

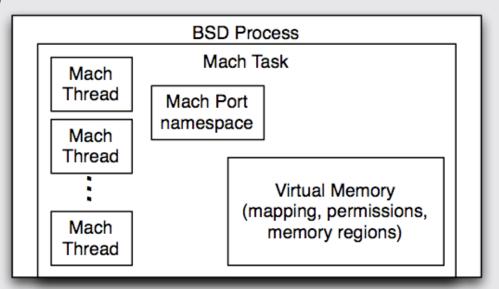


- Mac OS X kernel (xnu) is a hybrid between Mach 3.0 and FreeBSD
  - FreeBSD kernel top-half runs on Mach kernel bottom-half
  - Multiple system call interfaces: BSD (positive numbers), Mach (negative)
  - BSD sysctls, ioctls
  - Mach in-kernel RPC servers, IOKit user clients, etc.
- Mach inter-process communication (IPC)
  - Communicates over uni-directional ports, access controlled via rights
  - Multiple tasks may hold port send rights, only one may hold receive rights

## Tasks and Processes



- Mach Tasks own Threads, Ports, and Virtual Memory
- BSD Processes own file descriptors, etc.
- BSD Processes <=> Mach Task
  - task\_for\_pid(), pid\_for\_task()
- POSIX Thread != Mach Thread
  - Library functions use TLS



## Mach Task/Thread System Calls



- task\_create(parent\_task, ledgers, ledgers\_count, inherit\_memory, \*child\_task)
- thread\_create(parent\_task, \*child\_activation)
- vm\_allocate(task, \*address, size, flags)
- vm deallocate(task, address, size)
- vm\_read(task, address, size, \*data)
- vm\_write(task, address, data, data\_count)

#### **User-mode Mach Rootkits**



- Not as "sexy" as kernel mode rootkits
- Can be just as effective and harder to detect
- Are typically application/process -specific
- Based on thread injection or executable infection
- Would you notice an extra bundle and thread in your web browser?

## Injecting Mach Threads



- Get access to another task's task port
  - task\_for\_pid() or by exploiting a local privilege escalation vulnerability
- Allocate memory in remote process for thread stack and code trampoline
- Create new mach thread in remote process
  - Execute trampoline with previously allocated thread stack segment
  - Trampoline code promotes Mach Thread to POSIX Thread
    - Call \_pthread\_set\_self(pthread\_t) and cthread\_set\_self(pthread\_t)

## Mach Exceptions



- Tasks and Threads generate exceptions on memory errors
- Another thread (possibly in another task) may register as the exception handler for another thread or task
- Exception handling process:
  - 1. A Thread causes a runtime error, generates an exception
  - 2. Exception is delivered to thread exception handler (if exists)
  - 3. Exception is delivered to task's exception handler (if exists)
  - 4. Exception converted to Unix signal and delivered to BSD Process

## Injecting Mach Bundles



- Inject threads to call functions in the remote process
  - Remote thread calls injected trampoline code and then target function
  - Function returns to chosen bad address, generates an exception
  - Injector handles exception, retrieves function return value
- Call dlopen(), dlsym(), dlclose() to load bundle from disk
- Inject memory, call NSCreateObjectFileImageFromMemory(), NSLinkModule()
- Injected bundle can hook library functions, Objective-C methods

## inject-bundle



### inject-bundle

- —Inject a bundle from disk into a running process
- –Usage: inject\_bundle path\_to\_bundle [ pid ]

#### Sample bundles

- —test: Print output on load/run/unload
- -isight: Take a picture using iSight camera
- –sslspy: Log SSL traffic sent through SecureTransport
- -ichat: Log IMs from within iChat

# Hooking and Swizzling



- Hooking C functions is basically the same as on any other platform
  - -see Rentzsch's mach\_override
- Objective-C runtime has hooking built-in:
  - -method\_exchangeImplementations()
  - -or just switch the method pointers manually
  - -all due to Obj-C's dynamic runtime
  - –use JRSwizzle for portability



# DEMO

## Rootkitting the Web Browser



- What client system doesn't have the web browser open at all times?
- Will be allowed to connect to \*:80 and \*:443 by host-based firewalls (i.e. Little Snitch)
- Background thread can poll a known site for command and control instructions or look for instructions in HTML content from any site
- Injected bundles do not invalidate dynamic code signatures (used by Keychain, etc)

#### Kernel Mach Rootkits



- Mach system calls allow Mach RPC to in-kernel servers which perform task, thread, and VM operations
- RPC routines are stored in the mig\_buckets hash table by subsystem id + subroutine id
- Analogous to sysent table for Unix system calls
- Incoming Mach messages sent to a kernel-owned port are dispatched through mig\_buckets
- We can interpose on these function calls or inject new RPC servers by modifying this hash table

## Example: inject\_subsystem



```
int inject_subsystem(const struct mig_subsystem * mig)
        mach_msg_id_t h, i, r;
        // Insert each subroutine into mig_buckets hash table
        for (i = mig \rightarrow start; i < mig \rightarrow end; i++) {
                mig_hash_t* bucket;
                h = MIG_HASH(i);
                do { bucket = &mig_buckets[h % MAX_MIG_ENTRIES];
                } while (mig_buckets[h++ % MAX_MIG_ENTRIES].num != 0 &&
                                  h < MIG_HASH(i) + MAX_MIG_ENTRIES);</pre>
                if (bucket->num == 0) { // We found a free spot
                         r = mig -> start - i;
                         bucket->num = i;
                         bucket->routine = mig->routine[r].stub_routine;
                         if (mig->routine[r].max_reply_msq)
                                 bucket->size = mig->routine[r].max_reply_msg;
                         else
                                 bucket->size = mig->maxsize;
                         return 0;
        return -1;
```

## Mach Kernel RPC servers



- In-kernel Mach RPC subsystems are enumerated in the mig\_e table and interfaces are in /usr/ include/mach/subsystem.defs
  - —mach\_vm, mach\_port, mach\_host, host\_priv, host\_security, clock, clock\_priv, processor, processor\_set, is\_iokit, memory\_object\_name, lock\_set, ledger, semaphore, task, thread\_act, vm\_map, UNDReply, default\_pager\_object, security

#### Machiavelli



- Mach RPC provides high-level remote control
   —vm\_alloc(), vm\_write(), thread\_create() on kernel or any task
- Want to still use MiG generated client RPC stubs
- Machiavelli Proxy runs as background thread in control utilities on attacker's system
- Machiavelli Agents run on the remote compromised host as user-mode process or in kernel

## NetMessage and NetName servers



- Network transparency of IPC was a design goal
- Old Mach releases included the NetMessage Server
  - Mach servers could register themselves on the local NetName server
  - -Clients could lookup named servers on remote hosts
  - Local NetMessage server would act as a proxy,
     transmitting Mach IPC messages over the network
- These features no longer exist in Mac OS X

#### Machiavelli Architecture



#### Machiavelli Proxy

- Runs as background thread of a Machiavelli utility
- Receives messages on proxy ports and sends to remote Agent
- Replaces port names in messages received from Agent with proxy ports

#### Machiavelli Agent

- Receives messages over network from Proxy, sends to real destination
- Receives and transmits reply message if a reply is expected

#### Machiavelli Utilities

Run on control host, use Proxy to control compromised host

## Mach messages



- Mach messages are structured and unidirectional
- Header:

```
typedef struct
{
mach_msg_bits_t msgh_bits;
mach_msg_size_t msgh_size;
mach_port_t msgh_remote_port;
mach_port_t msgh_local_port;
mach_msg_size_t msgh_reserved;
mach_msg_id_t msgh_id;
} mach_msg_header_t;
```

Body consists of typed data items

## Complex Mach Messages



- "Complex" Mach messages contain out-of-line data and may transfer port rights and/or memory pages to other tasks
- In the message body, descriptors describe the port rights and memory pages to be transferred
- Kernel grants port rights to the receiving process
- Kernel maps transferred pages to receiving process, sometimes at message-specified address

# Proxying Mach Messages



- Proxy maintains a Mach port set
  - —A port set has the same interface as a single port and can be used identically in mach\_msg()
  - Each proxy port in the set corresponds to the real destination port name in the remote Agent
  - Port names can be arbitrary 32-bit values, so port set names are pointers to real destination port name values
- Received messages must be translated (local <=> remote ports and descriptor bits)
- Messages are serialized to byte buffers and then sent

# Serializing Mach Messages



- Serializing "simple" messages is simple as they don't contain any out-of-line data
- Out-of-line data is appended to the serialized buffer in order of the descriptors in the body
- Port names are translated during deserialization
  - —Translating to an intermediate "virtual port name" might be cleaner

## Deserializing Mach Messages



- Port names in the mach message must be replaced with local port names
- On Agent, this is done to receive the reply
- On Proxy, this is done to replace transferred port names with proxy port names
  - Ensures that only the initial port must be manually obtained from the proxy, the rest are handled automatically
- OOL memory is mapped+copied into address space

## Machiavelli example



```
int main(int argc, char* argv[])
{
    kern_return_t kr;
    mach_port_t port;
    vm_size_t page_size;

    machiavelli_t m = machiavelli_init();
    machiavelli_connect_tcp(m, "192.168.13.37", "31337");
    port = machiavelli_get_port(m, HOST_PORT);

    if ((kr = _host_page_size(port, &page_size)) != KERN_SUCCESS) {
        errx(EXIT_FAILURE, "_host_page_size: %s", mach_error_string(kr));
    }

    printf("Host page size: %d\n", page_size);
    return 0;
}
```



# DEMO

## Miscellaneous Agent services



- Agent must provide initial Mach ports:
  - -host port
  - -task\_for\_pid() (if pid == 0 => returns kernel task port)
- As OS X is a Mach/Unix hybrid, just controlling Mach is not enough
  - -i.e. How to list processes?
- Instead of implementing Unix functionality in Agent, inject Mach RPC server code into pid 1 (launchd)

# Network Kernel Extensions (NKEs)



- NKEs can extend or modify kernel networking functionality via:
  - Socket filters
  - IP filters
  - Interface filters
  - Network interfaces
  - Protocol plumbers

#### Conclusion



- Mach is a whole lot of fun
- Mach IPC can be made network transparent and provides a good abstraction for remote host control
- I wish my desktop was as secure as my iPhone
- For updated slides and tools go to:
  - -http://trailofbits.com/