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## Foreword

- 1 ISO (the International Organization for Standardization) and IEC (the International Electrotechnical Commission) form the specialized system for worldwide standardization. National bodies that are members of ISO or IEC participate in the development of International Standards through technical committees established by the respective organization to deal with particular fields of technical activity. ISO and IEC technical committees collaborate in fields of mutual interest. Other international organizations, governmental and nongovernmental, in liaison with ISO and IEC, also take part in the work. In the field of information technology, ISO and IEC have established a joint technical committee, ISO/IEC JTC 1.
- 2 International Standards are drafted in accordance with the rules given in the ISO/IEC Directives, Part 2.
- 3 The main task of the joint technical committee is to prepare International Standards. Draft International Standards adopted by the joint technical committee are circulated to national bodies for voting. Publication as an International Standard requires approval by at least 75 % of the national bodies casting a vote.
- 4 Attention is drawn to the possibility that some of the elements of this document may be the subject of patent rights. ISO and IEC shall not be held responsible for identifying any or all such patent rights.
- 5 ISO/IEC 1539-1 was prepared by Joint Technical Committee ISO/IEC JTC 1, *Information technology*, Subcommittee SC 22, *Programming languages, their environments and system software interfaces*.
- 6 This fourth edition cancels and replaces the third edition (ISO/IEC 1539-1:2010), which has been technically revised. It also incorporates the Technical Corrigenda ISO/IEC 1539-1:2010/Cor. 1:2012, ISO/IEC 1539-1:2010/Cor. 2:2013, ISO/IEC 1539-1:2010/Cor. 3:2014, and ISO/IEC 1539-1:2010/Cor. 4:2016, and the Technical Specifications ISO/IEC TS 29113:2012 and ISO/IEC TS 18508:2015.
- 7 ISO/IEC 1539 consists of the following parts, under the general title *Information technology — Programming languages — Fortran*:
  - 8 — *Part 1: Base language*
  - 9 — *Part 2: Varying length character strings*

## Introduction

- 1 This document comprises the specification of the base Fortran language, informally known as Fortran 2015. With the limitations noted in 4.3.3, the syntax and semantics of Fortran 2008 are contained entirely within Fortran 2015. Therefore, any standard-conforming Fortran 2008 program not affected by such limitations is a standard-conforming Fortran 2015 program. New features of Fortran 2015 can be compatibly incorporated into such Fortran 2008 programs, with any exceptions indicated in the text of this document.
- 2 Fortran 2015 contains several extensions to Fortran 2008; these are listed below.
  - Data declaration:

Constant properties of an object declared in its *entity-decl* can be used in its *initialization*. The **EQUIVALENCE** and **COMMON** statements and the **block data program unit** have been redundant since Fortran 90 and are now specified to be **obsolescent**. Diagnosis of the appearance of a **PROTECTED TARGET** variable accessed by **use association** as a *data-target* in a structure constructor is required.
  - Data usage and computation:

The **declared type** of the value supplied for a **polymorphic allocatable component** in a **structure constructor** is no longer required to be the same as the **declared type** of the **component**. **FORALL** is now specified to be **obsolescent**. The type and kind of an implied DO variable in an **array constructor** or **DATA statement** can be specified within the constructor or statement. The **SELECT RANK construct** provides structured access to the elements of an **assumed-rank array**. Completing execution of a **BLOCK construct** can cause the **association status of a pointer** with the **PROTECTED attribute** to become undefined. The standard intrinsic operations **<**, **<=**, **>**, and **>=** (also known as **.LT.**, **.LE.**, **.GT.**, and **.GE.**) on IEEE numbers provide **compareSignaling{relation}** operations; the **=** and **/=** operations (also known as **.EQ.** and **.NE.**) provide **compareQuiet{relation}** operations. Finalization of an allocatable subobject during intrinsic assignment has been clarified.
  - Input/output:

The **SIZE= specifier** can be used with advancing input. It is no longer prohibited to open a file on more than one unit. The value assigned by the **RECL= specifier** in an **INQUIRE statement** has been standardized. The values assigned by the **POS=** and **SIZE=** specifiers in an **INQUIRE statement** for a unit that has pending asynchronous operations have been standardized. The **G0.d edit descriptor** can be used for list items of type **Integer**, **Logical**, and **Character**. The **D**, **E**, **EN**, and **ES edit descriptors** can have a field width of zero, analogous to the **F edit descriptor**. The exponent width *e* in a **data edit descriptor** can be zero, analogous to a field width of zero. Floating-point formatted input accepts hexadecimal-significand numbers that conform to ISO/IEC/IEEE 60559:2011. The **EX edit descriptor** provides hexadecimal-significand formatted output conforming to ISO/IEC/IEEE 60559:2011. An error condition occurs if unacceptable characters are presented for logical or numeric editing during execution of a formatted input statement.
  - Execution control:

The **arithmetic IF statement** has been deleted. Labeled **DO loops** have been redundant since Fortran 90 and are now specified to be **obsolescent**. The **nonblock DO construct** has been deleted. The locality of a variable used in a **DO CONCURRENT construct** can be explicitly specified. The stop code in a **STOP** or **ERROR STOP** statement can be nonconstant. Output of the stop code and exception summary from the **STOP** and **ERROR STOP** statements can be controlled.
  - Intrinsic procedures and modules:

In references to the intrinsic functions **ALL**, **ANY**, **FINDLOC**, **IALL**, **IANY**, **IPARITY**, **MAXLOC**, **MAXVAL**, **MINLOC**, **MINVAL**, **NORM2**, **PARITY**, **PRODUCT**, **SUM**, and **THIS\_IMAGE**, the **actual argument** for **DIM** can be a present optional **dummy argument**. In a reference to the intrinsic function **CMPLX** with an actual argument of type complex, no keyword is needed for a **KIND** argument. The new intrinsic function **COSHAPE** returns the coshape of a **coarray**. The new intrinsic function **OUT\_OF\_RANGE** tests whether a numeric value can be safely converted to a different type or kind. The new intrinsic subroutine **RANDOM\_INIT** establishes the initial state of the pseudorandom number generator used by **RANDOM\_NUMBER**. The new intrinsic function **REDUCE** performs user-specified array reductions. A processor is required to report use of a **nonstandard intrinsic** procedure, use of a **nonstandard intrinsic** module, and use of a non-standard procedure from a **standard intrinsic** module. Integer and logical arguments to intrinsic procedures and intrinsic module procedures that were previously required to be of default kind no longer have that

requirement, except for `RANDOM_SEED`. Specific names for intrinsic functions are now deemed *obsolescent*. All standard procedures in the intrinsic module `ISO_C_BINDING`, other than `C_F_POINTER`, are now *pure*. The arguments to the intrinsic function `SIGN` can be of different kind. Nonpolymorphic pointer arguments to the intrinsic functions `EXTENDS_TYPE_OF` and `SAME_TYPE_AS` need not have defined pointer association status. The effects of invoking the intrinsic procedures `COMMAND_ARGUMENT_COUNT`, `GET_COMMAND`, and `GET_COMMAND_ARGUMENT`, on *images* other than *image* one, are no longer processor dependent. Access to error messages from the intrinsic subroutines `GET_COMMAND`, `GET_COMMAND_ARGUMENT`, and `GET_ENVIRONMENT_VARIABLE` is provided by an optional `ERRMSG` argument. The result of `NORM2` for a zero-sized array argument has been clarified.

- Program units and procedures:
 

The `IMPORT` statement can appear in a contained subprogram or `BLOCK` construct, and can restrict access via host association; diagnosis of violation of the `IMPORT` restrictions is required. The `GENERIC` statement can be used to declare generic interfaces. The number of procedure arguments is used in *generic resolution*. In a module, the *default accessibility* of entities accessed from another module can be controlled separately from the default accessibility of entities declared in the using module. An `IMPLICIT NONE` statement can require explicit declaration of the `EXTERNAL` attribute throughout a *scoping unit* and its contained *scoping units*. A *defined operation* need not specify `INTENT (IN)` for a *dummy argument* with the `VALUE` attribute. A *defined assignment* need not specify `INTENT (IN)` for the second *dummy argument* if it has the `VALUE` attribute. Procedures, including elemental procedures, can be invoked recursively by default; the `RECURSIVE` keyword is advisory only. The `NON_RECURSIVE` keyword specifies that a procedure is not recursive. The `ERROR STOP` statement can appear in a *pure subprogram*. A *dummy argument* of a *pure function* is permitted in a variable definition context, if it has the `VALUE` attribute. A *coarray dummy argument* can be referenced or defined by another *image*.
- Features previously described by ISO/IEC TS 29113:2012:
 

A *dummy data object* can *assume its rank* from its *effective argument*. A *dummy data object* can *assume the type* from its *effective argument*, without having the ability to perform type selection. An *interoperable* procedure can have *dummy arguments* that are *assumed-type* and/or *assumed-rank*. An *interoperable* procedure can have *dummy data objects* that are *allocatable*, *assumed-shape*, *optional*, or *pointers*. The *character length* of a *dummy data object* of an *interoperable* procedure can be assumed. The argument to `C_LOC` can be a noninteroperable array. The `FPTR` argument to `C_F_POINTER` can be a noninteroperable array pointer. The argument to `C_FUNLOC` can be a noninteroperable procedure. The `FPTR` argument to `C_F_PROCPONTER` can be a noninteroperable procedure pointer. Additionally to ISO/IEC TS 29113:2012, an *assumed-rank dummy data object* that is not associated with an *assumed-size array* can be used as the argument to the function `C_SIZEOF` from the intrinsic module `ISO_C_BINDING`.
- Changes to the intrinsic modules `IEEE_ARITHMETIC`, `IEEE_EXCEPTIONS`, and `IEEE_FEATURES` for conformance with ISO/IEC/IEEE 60559:2011:
 

There is a new, optional, rounding mode `IEEE_AWAY`. The new type `IEEE_MODES_TYPE` encapsulates all floating-point modes. Features associated with subnormal numbers can be accessed with functions and types named `...SUBNORMAL...` (the old `...DENORMAL...` names remain). The new function `IEEE_FMA` performs fused multiply-add operations. The function `IEEE_INT` performs rounded conversions to integer type. The new functions `IEEE_MAX_NUM`, `IEEE_MAX_NUM_MAG`, `IEEE_MIN_NUM`, and `IEEE_MIN_NUM_MAG` calculate maximum and minimum numeric values. The new functions `IEEE_NEXT_DOWN` and `IEEE_NEXT_UP` return the adjacent machine numbers. The new functions `IEEE_QUIET_EQ`, `IEEE_QUIET_GE`, `IEEE_QUIET_GT`, `IEEE_QUIET_LE`, `IEEE_QUIET_LT`, and `IEEE_QUIET_NE` perform quiet comparisons. The new functions `IEEE_SIGNALING_EQ`, `IEEE_SIGNALING_GE`, `IEEE_SIGNALING_GT`, `IEEE_SIGNALING_LE`, `IEEE_SIGNALING_LT`, and `IEEE_SIGNALING_NE` perform signaling comparisons. The decimal rounding mode can be inquired and set independently of the binary rounding mode, using the `RADIX` argument to `IEEE_GET_ROUNDING_MODE` and `IEEE_SET_ROUNDING_MODE`. The new function `IEEE_REAL` performs rounded conversions to real type. The function `IEEE_REM` now requires its arguments to have the same radix. The function `IEEE_RINT` now has a `ROUND` argument to perform specific rounding. The new function `IEEE_SIGNBIT` tests the sign bit of an IEEE number.
- Features previously described by ISO/IEC TS 18508:2015:
 

The `CRITICAL` statement has optional `ERRMSG=` and `STAT=` specifiers. The intrinsic subroutines

`ATOMIC_DEFINE` and `ATOMIC_REF` have an optional `STAT` argument. The new intrinsic subroutines `ATOMIC_ADD`, `ATOMIC_AND`, `ATOMIC_CAS`, `ATOMIC_FETCH_ADD`, `ATOMIC_FETCH_AND`, `ATOMIC_FETCH_OR`, `ATOMIC_FETCH_XOR`, `ATOMIC_OR`, and `ATOMIC_XOR` perform atomic operations. The new intrinsic functions `FAILED_IMAGES` and `STOPPED_IMAGES` return indices of images known to have `failed` or `stopped` respectively. The new intrinsic function `IMAGE_STATUS` returns the image execution status of an image. The intrinsic subroutine `MOVE_ALLOC` has optional `ERRMSG` and `STAT` arguments. The intrinsic functions `IMAGE_INDEX` and `NUM_IMAGES` have additional forms with a `TEAM` or `TEAM_NUMBER` argument. The intrinsic function `THIS_IMAGE` has an optional `TEAM` argument. The `EVENT POST` and `EVENT WAIT` statements, the intrinsic subroutine `EVENT_QUERY`, and the type `EVENT_TYPE` provide an event facility for one-sided segment ordering. The `CHANGE TEAM` construct, derived type `TEAM_TYPE`, `FORM TEAM` and `SYNC TEAM` statements, intrinsic functions `GET_TEAM` and `TEAM_NUMBER`, and the `TEAM=` and `TEAM_NUMBER=` specifiers on image selectors, provide a team facility for a subset of the program's images to act in concert as if it were the set of all images. This team facility allows an allocatable `coarray` to be allocated or deallocated on a subset of images. The new intrinsic subroutines `CO_BROADCAST`, `CO_MAX`, `CO_MIN`, `CO_REDUCE`, and `CO_SUM` perform collective reduction operations on the images of the `current team`. The concept of `failed images`, the `FAIL IMAGE` statement, the `STAT=` specifier on image selectors, and the named constant `STAT_FAILED_IMAGE` from the intrinsic module `ISO_FORTRAN_ENV` provide support for fault-tolerant parallel execution.

- Changes to features previously described by ISO/IEC TS 18508:2015:
  - The `CHANGE TEAM` and `SYNC TEAM` statements, and the `TEAM=` specifier on image selectors, permit the team to be specified by an expression. The intrinsic functions `FAILED_IMAGES` and `STOPPED_IMAGES` have no restriction on the kind of their result. The name of the function argument to the intrinsic function `CO_REDUCE` is `OPERATION` instead of `OPERATOR`; this argument is not required to be commutative. The named constant `STAT_UNLOCKED_FAILED_IMAGE` from the intrinsic module `ISO_FORTRAN_ENV` indicates that a `lock variable` was locked by an image that `failed`. The `team number` for the `initial team` can be used in image selectors, and in the intrinsic functions `NUM_IMAGES` and `IMAGE_INDEX`. A `team variable` that appears in a `CHANGE TEAM` statement can no longer be defined or become undefined during execution of the `CHANGE TEAM` construct. All `images` of the `current team` are no longer required to execute the same `CHANGE TEAM` statement.

- 3 This document is organized in 19 clauses, dealing with 8 conceptual areas. These 8 areas, and the clauses in which they are treated, are:

High/low level concepts	Clauses 4, 5, 6
Data concepts	Clauses 7, 8, 9
Computations	Clauses 10, 16, 17
Execution control	Clause 11
Input/output	Clauses 12, 13
Program units	Clauses 14, 15
Interoperability with C	Clause 18
Scoping and association rules	Clause 19

- 4 It also contains the following nonnormative material:

Processor dependencies	Annex A
Deleted and obsolescent features	Annex B
Extended notes	Annex C
Index	Index

# Information technology — Programming languages — Fortran —

## Part 1: Base language

### 1 Scope

1 This document specifies the form and establishes the interpretation of programs expressed in the base Fortran  
2 language. The purpose of this document is to promote portability, reliability, maintainability, and efficient  
3 execution of Fortran programs for use on a variety of computing systems.

4 This document specifies

- 5 • the forms that a program written in the Fortran language may take,
- 6 • the rules for interpreting the meaning of a program and its data,
- 7 • the form of the input data to be processed by such a program, and
- 8 • the form of the output data resulting from the use of such a program.

9 Except where stated otherwise, requirements and prohibitions specified by this document apply to programs  
10 rather than processors.

11 This document does not specify

- 12 • the mechanism by which programs are transformed for use on computing systems,
- 13 • the operations required for setup and control of the use of programs on computing systems,
- 14 • the method of transcription of programs or their input or output data to or from a storage medium,
- 15 • the program and processor behavior when this document fails to establish an interpretation except for the  
16 processor detection and reporting requirements in items (2) to (10) of 4.2,
- 17 • the maximum number of [images](#), or the size or complexity of a program and its data that will exceed the  
18 capacity of any particular computing system or the capability of a particular processor,
- 19 • the mechanism for determining the number of [images](#) of a program,
- 20 • the physical properties of an [image](#) or the relationship between [images](#) and the computational elements of  
21 a computing system,
- 22 • the physical properties of the representation of quantities and the method of rounding, approximating, or  
23 computing numeric values on a particular processor, except by reference to ISO/IEC/IEEE 60559:2011  
24 under conditions specified in Clause 17,
- 25 • the physical properties of input/output records, files, and units, or
- 26 • the physical properties and implementation of storage.

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ISO/IEC JTC 1/SC 22/WG5/N2137

## 1 2 Normative references

2 1 The following documents are referred to in the text in such a way that some or all of their content constitutes  
3 requirements of this document. For dated references, only the edition cited applies. For undated references, the  
4 latest edition of the referenced document (including any amendments) applies.

5 ISO/IEC 646:1991 (International Reference Version), *Information technology—ISO 7-bit coded character set for*  
6 *information interchange*

7 ISO/IEC 9899:2011, *Programming languages—C*

8 ISO/IEC 10646, *Information technology—Universal Multiple-Octet Coded Character Set (UCS)*

9 ISO/IEC/IEEE 60559:2011, *Information technology — Microprocessor Systems — Floating-Point arithmetic*

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## 3 Terms and definitions

1 For the purposes of this document, the following terms and definitions apply. ISO and IEC maintain terminological databases for use in standardization at the following addresses:

- IEC Electropedia: available at <http://www.electropedia.org/>
- ISO Online browsing platform: available at <http://www.iso.org/obp>

### 3.1

#### **actual argument**

entity (R1524) that appears in a [procedure reference](#)

### 3.2

#### **allocatable**

having the [ALLOCATABLE attribute](#) (8.5.3)

### 3.3

#### **array**

set of scalar data, all of the same type and [type parameters](#), whose individual elements are arranged in a rectangular pattern (8.5.8, 9.5)

#### **3.3.1**

##### **array element**

scalar individual element of an array

#### **3.3.2**

##### **array pointer**

array with the [POINTER attribute](#) (8.5.14)

#### **3.3.3**

##### **array section**

array [subobject](#) designated by [array-section](#), and which is itself an array (9.5.3.3)

#### **3.3.4**

##### **assumed-shape array**

nonallocatable nonpointer [dummy argument](#) array that takes its shape from its [effective argument](#) (8.5.8.3)

#### **3.3.5**

##### **assumed-size array**

[dummy argument](#) array whose size is assumed from that of its [effective argument](#) (8.5.8.5)

#### **3.3.6**

##### **deferred-shape array**

[allocatable](#) array or array pointer, declared with a [deferred-shape-spec-list](#) (8.5.8.4)

#### **3.3.7**

##### **explicit-shape array**

array declared with an [explicit-shape-spec-list](#), which specifies explicit values for the [bounds](#) in each dimension of the array (8.5.8.2)

### 3.4

#### **ASCII character**

character whose representation method corresponds to ISO/IEC 646:1991 (International Reference Version)

- 1     **3.5**  
2     **associate name**  
3     name of [construct entity](#) associated with a selector of an [ASSOCIATE](#), [CHANGE TEAM](#), [SELECT RANK](#), or  
4     [SELECT TYPE](#) construct ([11.1.3](#), [11.1.5](#), [11.1.10](#), [11.1.11](#))
- 5     **3.6**  
6     **associating entity**  
7     ⟨in a dynamically-established association⟩ the entity that did not exist prior to the establishment of the association  
8     ([19.5.5](#))
- 9     **3.7**  
10    **association**  
11    [inheritance association](#), [name association](#), [pointer association](#), or [storage association](#).
- 12    **3.7.1**  
13    **argument association**  
14    association between an [effective argument](#) and a [dummy argument](#) ([15.5.2](#))
- 15    **3.7.2**  
16    **construct association**  
17    association between a selector and an [associate name](#) in an [ASSOCIATE](#), [CHANGE TEAM](#), [SELECT RANK](#),  
18    or [SELECT TYPE](#) construct([11.1.3](#), [11.1.5](#), [11.1.10](#), [11.1.11](#), [19.5.1.6](#))
- 19    **3.7.3**  
20    **host association**  
21    name association, other than argument association, between entities in a submodule or contained [scoping unit](#)  
22    and entities in its host ([19.5.1.4](#))
- 23    **3.7.4**  
24    **inheritance association**  
25    association between the [inherited](#) components of an [extended type](#) and the components of its [parent component](#)  
26    ([19.5.4](#))
- 27    **3.7.5**  
28    **linkage association**  
29    association between a variable or common block with the [BIND attribute](#) and a C global variable ([18.9](#), [19.5.1.5](#))
- 30    **3.7.6**  
31    **name association**  
32    [argument association](#), [construct association](#), [host association](#), [linkage association](#), or [use association](#) ([19.5.1](#))
- 33    **3.7.7**  
34    **pointer association**  
35    association between a [pointer](#) and an entity with the [TARGET attribute](#) ([19.5.2](#))
- 36    **3.7.8**  
37    **storage association**  
38    association between storage sequences ([19.5.3](#))
- 39    **3.7.9**  
40    **use association**  
41    association between entities in a module and entities in a [scoping unit](#) or construct that references that module,  
42    as specified by a [USE statement](#) ([14.2.2](#))
- 43    **3.8**  
44    **assumed-rank dummy data object**  
45    [dummy data object](#) that assumes the rank, shape, and size of its [effective argument](#) ([8.5.8.7](#))

- 1     **3.9**  
2     **assumed-type**  
3     declared with a TYPE(\*) type specifier (7.3.2)
- 4     **3.10**  
5     **attribute**  
6     property of an entity that determines its uses (8.1)
- 7     **3.11**  
8     **automatic data object**  
9     nondummy [data object](#) with a [type parameter](#) or array [bound](#) that depends on the value of a [specification-expr](#)  
10    that is not a [constant expression](#) (8.3)
- 11    **3.12**  
12    **base object**  
13    <[data-ref](#)> object designated by the leftmost [part-name](#) (9.4.2)
- 14    **3.13**  
15    **binding**  
16    [type-bound procedure](#) or [final subroutine](#) (7.5.5)
- 17    **3.14**  
18    **binding name**  
19    name given to a specific or generic [type-bound procedure](#) in the type definition (7.5.5)
- 20    **3.15**  
21    **binding label**  
22    default character value specifying the name by which a global entity with the [BIND attribute](#) is known to the  
23    [companion processor](#) (18.10.2, 18.9.2)
- 24    **3.16**  
25    **block**  
26    sequence of executable constructs formed by the syntactic class [block](#) and which is treated as a unit by the  
27    executable constructs described in 11.1
- 28    **3.17**  
29    **bound**  
30    **array bound**  
31    limit of a dimension of an [array](#) (8.5.8)
- 32    **3.18**  
33    **branch target statement**  
34    [action-stmt](#), [associate-stmt](#), [end-associate-stmt](#), [if-then-stmt](#), [end-if-stmt](#), [select-case-stmt](#), [end-select-stmt](#), [select-](#)  
35    [rank-stmt](#), [end-select-rank-stmt](#), [select-type-stmt](#), [end-select-type-stmt](#), [do-stmt](#), [end-do-stmt](#), [block-stmt](#), [end-](#)  
36    [block-stmt](#), [critical-stmt](#), [end-critical-stmt](#), [forall-construct-stmt](#), [where-construct-stmt](#), [end-function-stmt](#), [end-mp-](#)  
37    [subprogram-stmt](#), [end-program-stmt](#), or [end-subroutine-stmt](#), whose statement label appears as a [label](#) in a GO TO  
38    statement, computed GO TO statement, [alt-return-spec](#), END= specifier, EOR= specifier, or ERR= specifier (11.2.1)
- 39    **3.19**  
40    **C address**  
41    value identifying the location of a [data object](#) or procedure either defined by the [companion processor](#) or which  
42    might be accessible to the [companion processor](#)

**NOTE 3.1**

This is the concept that ISO/IEC 9899:2011 calls the address.

1 **3.20**2 **C descriptor**3 C structure of type [CFI\\_cdesc\\_t](#) defined in the source file [ISO\\_Fortran\\_binding.h](#) ([18.4](#), [18.5](#))4 **3.21**5 **character context**6 within a character literal constant ([7.4.4](#)) or within a character string edit descriptor ([13.3.2](#))7 **3.22**8 **characteristics**9 [⟨dummy argument⟩](#) being a [dummy data object](#), [dummy procedure](#), or an asterisk (alternate return indicator)10 **3.23**11 **characteristics**12 [⟨dummy data object⟩](#) properties listed in [15.3.2.2](#)13 **3.24**14 **characteristics**15 [⟨dummy procedure or dummy procedure pointer⟩](#) properties listed in [15.3.2.3](#)16 **3.25**17 **characteristics**18 [⟨function result⟩](#) properties listed in [15.3.3](#)19 **3.26**20 **characteristics**21 [⟨procedure⟩](#) properties listed in [15.3.1](#)22 **3.27**23 **coarray**24 [data entity](#) that has nonzero corank ([5.4.7](#))25 **3.27.1**26 **established coarray**27 [coarray](#) that is accessible using an [image-selector](#) ([5.4.8](#))28 **3.28**29 **cobound**30 bound (limit) of a [codimension](#) ([8.5.6](#))31 **3.29**32 **codimension**33 dimension of the pattern formed by a set of corresponding [coarrays](#) ([8.5.6](#))34 **3.30**35 **coindexed object**36 [data object](#) whose [designator](#) includes an [image-selector](#) ([R924](#), [9.6](#))37 **3.31**38 **collating sequence**39 one-to-one mapping from a character set into the nonnegative integers ([7.4.4.4](#))40 **3.32**41 **common block**42 block of physical storage specified by a [COMMON statement](#) ([8.10.2](#))43 **3.32.1**44 **blank common**

45 unnamed common block

1 **3.33**2 **companion processor**

3 processor-dependent mechanism by which global data and procedures may be referenced or defined (5.5.7)

4 **3.34**5 **component**

6 part of a derived type, or of an object of derived type, defined by a *component-def-stmt* (7.5.4)

7 **3.34.1**8 **direct component**

9 one of the components, or one of the direct components of a nonpointer nonallocatable component (7.5.1)

10 **3.34.2**11 **parent component**

12 component of an *extended type* whose type is that of the *parent type* and whose components are *inheritance*  
13 *associated* with the *inherited* components of the *parent type* (7.5.7.2)

14 **3.34.3**15 **potential subobject component**

16 nonpointer component, or potential subobject component of a nonpointer component (7.5.1)

17 **3.34.4**18 **subcomponent**

19 *<structure>* *direct component* that is a *subobject* of the *structure* (9.4.2)

20 **3.34.5**21 **ultimate component**

22 component that is of *intrinsic type*, a pointer, or allocatable; or an ultimate component of a nonpointer nonal-  
23 locatable component of derived type

24 **3.35**25 **component order**

26 ordering of the nonparent components of a derived type that is used for *intrinsic* formatted input/output and  
27 *structure constructors* (where component keywords are not used) (7.5.4.7)

28 **3.36**29 **conformable**

30 *<of two data entities>* having the same shape, or one being an array and the other being scalar

31 **3.37**32 **connected**

33 relationship between a *unit* and a file: each is connected if and only if the *unit* refers to the file (12.5.4)

34 **3.38**35 **constant**

36 *data object* that has a value and which cannot be defined, redefined, or become undefined during execution of a  
37 program (6.2.3, 9.3)

38 **3.38.1**39 **literal constant**

40 constant that does not have a name (R605, 7.4)

41 **3.38.2**42 **named constant**

43 named *data object* with the *PARAMETER attribute* (8.5.13)

1 **3.39**2 **construct entity**

3 entity whose identifier has the scope of a construct (19.1, 19.4)

4 **3.40**5 **constant expression**

6 expression satisfying the requirements specified in 10.1.12, thus ensuring that its value is constant

7 **3.41**8 **contiguous**

9 ⟨array⟩ having array elements in order that are not separated by other data objects, as specified in 8.5.7

10 **3.42**11 **contiguous**

12 ⟨multi-part data object⟩ that the parts in order are not separated by other data objects

13 **3.43**14 **corank**15 number of [codimensions](#) of a [coarray](#) (zero for objects that are not coarrays) (8.5.6)16 **3.44**17 **cosubscript**18 (R925) scalar integer expression in an *image-selector* (R924)19 **3.45**20 **data entity**21 [data object](#), result of the evaluation of an expression, or the result of the execution of a function reference22 **3.46**23 **data object**24 **object**25 constant (7.1.4), [variable](#) (9), or [subobject](#) of a constant (5.4.3.2.3)26 **3.47**27 **decimal symbol**28 character that separates the whole and fractional parts in the decimal representation of a real number in a file  
29 (13.6)**3.48****declaration**

specification of attributes for various program entities

**NOTE 3.2**

Often this involves specifying the type of a named [data object](#) or specifying the shape of a named array [object](#).

30 **3.49**31 **default initialization**32 mechanism for automatically initializing pointer components to have a defined pointer association status, and  
33 nonpointer components to have a particular value (7.5.4.6)34 **3.50**35 **default-initialized**36 ⟨[subcomponent](#)⟩ subject to a [default initialization](#) specified in the type definition for that [component](#) (7.5.4.6)37 **3.51**38 **definable**39 capable of [definition](#) and permitted to become [defined](#)

1 **3.52**2 **defined**3 [⟨data object⟩](#) has a valid value4 **3.53**5 **defined**6 [⟨pointer⟩](#) has a pointer association status of associated or [disassociated](#) (19.5.2.2)7 **3.54**8 **defined assignment**

9 assignment defined by a procedure (10.2.1.4, 15.4.3.4.3)

10 **3.55**11 **defined input/output**12 input/output defined by a procedure and accessed via a [defined-io-generic-spec](#) (R1509, 12.6.4.8)13 **3.56**14 **defined operation**

15 operation defined by a procedure (10.1.6.1, 15.4.3.4.2)

16 **3.57**17 **definition**18 [⟨data object⟩](#) process by which the data object becomes defined (19.6.5)19 **3.58**20 **definition**21 [⟨derived type](#) (7.5.2), enumeration (7.6), or [procedure](#) (15.6)⟩ specification of the type, enumeration, or procedure22 **3.59**23 **descendant**24 [⟨module or submodule⟩](#) submodule that extends that module or submodule or that extends another descendant  
25 thereof (14.2.3)26 **3.60**27 **designator**28 name followed by zero or more component selectors, complex part selectors, array section selectors, array element  
29 selectors, image selectors, and substring selectors (9.1)30 **3.60.1**31 **complex part designator**32 designator that designates the real or imaginary part of a complex [data object](#), independently of the other part  
33 (9.4.4)**3.60.2****object designator****data object designator**[designator](#) for a [data object](#)**NOTE 3.3**

An object name is a special case of an object designator.

34 **3.60.3**35 **procedure designator**36 [designator](#) for a procedure

1 **3.61**  
2 **disassociated**  
3 [⟨pointer association⟩ pointer association](#) status of not being associated with any [target](#) and not being [undefined](#)  
4 ([19.5.2.2](#))

5 **3.62**  
6 **disassociated**  
7 [⟨pointer⟩](#) has a [pointer association](#) status of disassociated

8 **3.63**  
9 **dummy argument**  
10 entity whose identifier appears in a dummy argument list ([R1536](#)) in a [FUNCTION](#), [SUBROUTINE](#), [ENTRY](#), or  
11 [statement function](#) statement, or whose name can be used as an [argument keyword](#) in a reference to an [intrinsic](#)  
12 procedure or a procedure in an [intrinsic](#) module

13 **3.63.1**  
14 **dummy data object**  
15 [dummy argument](#) that is a data object

16 **3.63.2**  
17 **dummy function**  
18 [dummy procedure](#) that is a function

19 **3.64**  
20 **effective argument**  
21 entity that is argument-associated with a dummy argument ([15.5.2.3](#))

22 **3.65**  
23 **effective item**  
24 scalar object resulting from the application of the rules in [12.6.3](#) to an input/output list

25 **3.66**  
26 **elemental**  
27 independent scalar application of an action or operation to elements of an array or corresponding elements of a  
28 set of conformable arrays and scalars, or possessing the capability of elemental operation

#### NOTE 3.4

Combination of scalar and array operands or arguments combine the scalar operand(s) with each element of the array operand(s).
--

29 **3.66.1**  
30 **elemental assignment**  
31 assignment that operates elementally

32 **3.66.2**  
33 **elemental operation**  
34 operation that operates elementally

35 **3.66.3**  
36 **elemental operator**  
37 operator in an elemental operation

38 **3.66.4**  
39 **elemental procedure**  
40 elemental [intrinsic](#) procedure or procedure defined by an elemental subprogram ([15.8](#))



- 1     **3.66.5**  
2     **elemental reference**  
3     reference to an elemental procedure with at least one array actual argument
- 4     **3.66.6**  
5     **elemental subprogram**  
6     subprogram with the [ELEMENTAL](#) prefix ([15.8.1](#))
- 7     **3.67**  
8     **END statement**  
9     *[end-block-data-stmt](#), [end-function-stmt](#), [end-module-stmt](#), [end-mp-subprogram-stmt](#), [end-program-stmt](#),*  
10    *[end-submodule-stmt](#), or [end-subroutine-stmt](#)*
- 11    **3.68**  
12    **explicit initialization**  
13    initialization of a data object by a specification statement ([8.4](#), [8.6.7](#))
- 14    **3.69**  
15    **extent**  
16    number of elements in a single dimension of an [array](#)
- 17    **3.70**  
18    **external file**  
19    file that exists in a medium external to the program ([12.3](#))
- 20    **3.71**  
21    **external unit**  
22    **external input/output unit**  
23    entity that can be [connected](#) to an [external file](#) ([12.5.3](#), [12.5.4](#))
- 24    **3.72**  
25    **file storage unit**  
26    unit of storage in a [stream file](#) or an unformatted [record file](#) ([12.3.5](#))
- 27    **3.73**  
28    **final subroutine**  
29    subroutine whose name appears in a [FINAL statement](#) ([7.5.6](#)) in a type definition, and which can be automatically  
30    invoked by the processor when an object of that type is finalized ([7.5.6.2](#))
- 31    **3.74**  
32    **finalizable**  
33    ⟨type⟩ has a final subroutine or a nonpointer nonallocatable component of finalizable type
- 34    **3.75**  
35    **finalizable**  
36    ⟨nonpointer data entity⟩ of finalizable type
- 37    **3.76**  
38    **finalization**  
39    process of calling [final subroutines](#) when one of the events listed in [7.5.6.3](#) occurs
- 40    **3.77**  
41    **function**  
42    procedure that is invoked by an expression
- 43    **3.78**  
44    **function result**  
45    entity that returns the value of a function ([15.6.2.2](#))

- 1 **3.79**  
2 **generic identifier**  
3 [lexical token](#) that identifies a generic set of procedures, [intrinsic](#) operations, and/or [intrinsic](#) assignments (15.4.3.4.1)
- 4 **3.80**  
5 **host instance**  
6 [internal procedure](#), or [dummy procedure](#) or [procedure pointer](#) associated with an [internal procedure](#) instance  
7 of the host procedure that supplies the host environment of the [internal procedure](#) (15.6.2.4)
- 8 **3.81**  
9 **host scoping unit**  
10 **host**  
11 [scoping unit](#) immediately surrounding another [scoping unit](#), or the [scoping unit](#) extended by a submodule
- 12 **3.82**  
13 **IEEE infinity**  
14 ISO/IEC/IEEE 60559:2011 conformant infinite floating-point value
- 15 **3.83**  
16 **IEEE NaN**  
17 ISO/IEC/IEEE 60559:2011 conformant floating-point datum that does not represent a number
- 18 **3.84**  
19 **image**  
20 instance of a Fortran program (5.3.4)
- 21 **3.84.1**  
22 **active image**  
23 [image](#) that has not failed or stopped (5.3.6)
- 24 **3.84.2**  
25 **failed image**  
26 [image](#) that has not initiated termination but which has ceased to participate in program execution (5.3.6)
- 27 **3.84.3**  
28 **stopped image**  
29 [image](#) that has initiated [normal termination](#) (5.3.6)
- 30 **3.85**  
31 **image index**  
32 integer value identifying an [image](#) within a [team](#)
- 33 **3.86**  
34 **image control statement**  
35 statement that affects the execution ordering between [images](#) (11.6)
- 36 **3.87**  
37 **inclusive scope**  
38 nonblock [scoping unit](#) plus every [block scoping unit](#) whose [host](#) is that [scoping unit](#) or that is nested within such  
39 a [block scoping unit](#)

**NOTE 3.5**

That is, inclusive scope is the scope as if **BLOCK** constructs were not [scoping units](#).

- 40 **3.88**  
41 **inherit**  
42 [extended type](#) acquire entities ([components](#), [type-bound procedures](#), and [type parameters](#)) through type exten-  
43 sion from the parent type (7.5.7.2)

- 1 **3.89**  
2 **inquiry function**  
3 [intrinsic](#) function, or function in an [intrinsic](#) module, whose result depends on the properties of one or more of  
4 its arguments instead of their values
- 5 **3.90**  
6 **interface**  
7 ⟨procedure⟩ name, procedure characteristics, dummy argument names, binding label, and generic identifiers  
8 ([15.4.1](#))
- 9 **3.90.1**  
10 **abstract interface**  
11 set of procedure characteristics with dummy argument names ([15.4.1](#))
- 12 **3.90.2**  
13 **explicit interface**  
14 interface of a procedure that includes all the characteristics of the procedure and names for its dummy arguments  
15 except for asterisk dummy arguments ([15.4.2](#))
- 16 **3.90.3**  
17 **generic interface**  
18 set of procedure interfaces identified by a [generic identifier](#)
- 19 **3.90.4**  
20 **implicit interface**  
21 interface of a procedure that is not an explicit interface ([15.4.2](#), [15.4.3.8](#))
- 22 **3.90.5**  
23 **specific interface**  
24 [interface](#) identified by a nongeneric name
- 25 **3.91**  
26 **interface block**  
27 [abstract interface block](#), [generic interface block](#), or [specific interface block](#) ([15.4.3.2](#))
- 28 **3.91.1**  
29 **abstract interface block**  
30 interface block with the [ABSTRACT](#) keyword; collection of [interface bodies](#) that specify named [abstract interfaces](#)
- 31 **3.91.2**  
32 **generic interface block**  
33 interface block with a [generic-spec](#); collection of [interface bodies](#) and procedure statements that are to be given  
34 that generic identifier
- 35 **3.91.3**  
36 **specific interface block**  
37 interface block with no [generic-spec](#) or [ABSTRACT](#) keyword; collection of [interface bodies](#) that specify the  
38 interfaces of procedures
- 39 **3.92**  
40 **interoperable**  
41 ⟨Fortran entity⟩ equivalent to an entity defined by or definable by the [companion processor](#) ([18.3](#))
- 42 **3.93**  
43 **intrinsic**  
44 type, procedure, module, assignment, operator, or input/output operation defined in this document and accessible  
45 without further definition or specification, or a procedure or module provided by a processor but not defined in  
46 this document

- 1 **3.93.1**  
2 **standard intrinsic**  
3 ⟨procedure or module⟩ defined in this document (16)
- 4 **3.93.2**  
5 **nonstandard intrinsic**  
6 ⟨procedure or module⟩ provided by a processor but not defined in this document
- 7 **3.94**  
8 **internal file**  
9 character variable that is [connected](#) to an [internal unit](#) (12.4)
- 10 **3.95**  
11 **internal unit**  
12 [input/output unit](#) that is [connected](#) to an [internal file](#) (12.5.4)
- 13 **3.96**  
14 **ISO 10646 character**  
15 character whose representation method corresponds to UCS-4 in ISO/IEC 10646
- 16 **3.97**  
17 **keyword**  
18 statement keyword, argument keyword, type parameter keyword, or component keyword
- 19 **3.97.1**  
20 **argument keyword**  
21 word that identifies the corresponding [dummy argument](#) in an [actual argument](#) list (15.5.2.1)
- 22 **3.97.2**  
23 **component keyword**  
24 word that identifies a [component](#) in a [structure constructor](#) (7.5.10)
- 25 **3.97.3**  
26 **statement keyword**  
27 word that is part of the syntax of a statement (5.5.2)
- 28 **3.97.4**  
29 **type parameter keyword**  
30 word that identifies a [type parameter](#) in a type parameter list
- 31 **3.98**  
32 **lexical token**  
33 keyword, name, literal constant other than a complex literal constant, operator, label, delimiter, comma, =, =>, ;, ::, :, or % (6.2)
- 35 **3.99**  
36 **line**  
37 sequence of zero or more characters
- 38 **3.100**  
39 **main program**  
40 [program unit](#) that is not a [subprogram](#), [module](#), [submodule](#), or [block data program unit](#) (14.1)
- 41 **3.101**  
42 **masked array assignment**  
43 [assignment statement](#) in a [WHERE statement](#) or [WHERE construct](#) (10.2.3)

- 1 **3.102**  
2 **module**  
3 [program unit](#) containing (or accessing from other modules) definitions that are to be made accessible to other  
4 [program units](#) (14.2)
- 5 **3.103**  
6 **name**  
7 identifier of a program constituent, formed according to the rules given in [6.2.2](#)
- 8 **3.104**  
9 **NaN**  
10 Not a Number, a symbolic floating-point datum (ISO/IEC/IEEE 60559:2011)
- 11 **3.105**  
12 **operand**  
13 data value that is the subject of an operator
- 14 **3.106**  
15 **operator**  
16 *[intrinsic-operator](#)*, *[defined-unary-op](#)*, or *[defined-binary-op](#)* (R608, R1003, R1023)
- 17 **3.107**  
18 **passed-object dummy argument**  
19 dummy argument of a [type-bound procedure](#) or procedure pointer component that becomes associated with the  
20 object through which the procedure is invoked ([7.5.4.5](#))
- 21 **3.108**  
22 **pointer**  
23 [data pointer](#) or [procedure pointer](#)
- 24 **3.108.1**  
25 **data pointer**  
26 [data entity](#) with the [POINTER](#) attribute ([8.5.14](#))
- 27 **3.108.2**  
28 **procedure pointer**  
29 procedure with the [EXTERNAL](#) and [POINTER](#) attributes ([8.5.9](#), [8.5.14](#))
- 30 **3.108.3**  
31 **local procedure pointer**  
32 [procedure pointer](#) that is part of a [local variable](#), or a named [procedure pointer](#) that is not a [dummy argument](#)  
33 or accessed by use or host association
- 34 **3.109**  
35 **pointer assignment**  
36 association of a pointer with a target, by execution of a [pointer assignment statement](#) ([10.2.2](#)) or an [intrinsic](#)  
37 [assignment statement](#) ([10.2.1.2](#)) for a derived-type object that has the pointer as a subobject
- 38 **3.110**  
39 **polymorphic**  
40 [data entity](#) able to be of differing [dynamic types](#) during program execution ([7.3.2.3](#))
- 41 **3.111**  
42 **preconnected**  
43 [file](#) or [unit](#) [connected](#) at the beginning of execution of the program ([12.5.5](#))

- 1 **3.112**  
2 **procedure**  
3 entity encapsulating an arbitrary sequence of actions that can be invoked directly during program execution
- 4 **3.112.1**  
5 **dummy procedure**  
6 procedure that is a [dummy argument](#) (15.2.2.3)
- 7 **3.112.2**  
8 **external procedure**  
9 procedure defined by an external subprogram (R503) or by means other than Fortran (15.6.3)
- 10 **3.112.3**  
11 **internal procedure**  
12 procedure defined by an internal subprogram (R512)
- 13 **3.112.4**  
14 **module procedure**  
15 procedure defined by a module subprogram, or a procedure provided by an [intrinsic](#) module (R1408)
- 16 **3.112.5**  
17 **pure procedure**  
18 procedure declared or defined to be pure (15.7)
- 19 **3.112.6**  
20 **type-bound procedure**  
21 procedure that is bound to a derived type and referenced via an object of that type (7.5.5)
- 22 **3.113**  
23 **processor**  
24 combination of a computing system and mechanism by which programs are transformed for use on that computing  
25 system
- 26 **3.114**  
27 **processor dependent**  
28 not completely specified in this document, having methods and semantics determined by the processor
- 29 **3.115**  
30 **program**  
31 set of Fortran [program units](#) and entities defined by means other than Fortran that includes exactly one [main](#)  
32 [program](#)
- 33 **3.116**  
34 **program unit**  
35 [main program](#), [external subprogram](#), [module](#), [submodule](#), or [block data program unit](#) (5.2.1)
- 36 **3.117**  
37 **rank**  
38 number of array dimensions of a [data entity](#) (zero for a scalar entity)
- 39 **3.118**  
40 **record**  
41 sequence of values or characters in a file (12.2)
- 42 **3.119**  
43 **record file**  
44 file composed of a sequence of records (12.1)

- 1 **3.120**  
2 **reference**  
3 [data object](#) reference, [procedure](#) reference, or [module](#) reference
- 4 **3.120.1**  
5 **data object reference**  
6 appearance of a [data object designator](#) (9.1) in a context requiring its value at that point during execution
- 7 **3.120.2**  
8 **function reference**  
9 appearance of the [procedure designator](#) for a function, or operator symbol for a [defined operation](#), in a context  
10 requiring execution of the function during expression evaluation (15.5.3)
- 11 **3.120.3**  
12 **module reference**  
13 appearance of a module name in a [USE statement](#) (14.2.2)
- 14 **3.120.4**  
15 **procedure reference**  
16 appearance of a [procedure designator](#), operator symbol, or assignment symbol in a context requiring execution  
17 of the procedure at that point during execution; or occurrence of defined input/output (13.7.6) or derived-type  
18 [finalization](#) (7.5.6.2)
- 19 **3.121**  
20 **saved**  
21 having the [SAVE attribute](#) (8.5.16)
- 22 **3.122**  
23 **scalar**  
24 [data entity](#) that can be represented by a single value of the type and that is not an array (9.5)
- 25 **3.123**  
26 **scoping unit**  
27 [BLOCK construct](#), derived-type definition, [interface body](#), [program unit](#), or subprogram, excluding all nested  
28 scoping units in it
- 29 **3.123.1**  
30 **block scoping unit**  
31 scoping unit of a [BLOCK construct](#)
- 32 **3.124**  
33 **sequence**  
34 set of elements ordered by a one-to-one correspondence with the numbers 1, 2, to  $n$
- 35 **3.125**  
36 **sequence structure**  
37 [scalar data object](#) of a [sequence type](#) (7.5.2.3)
- 38 **3.126**  
39 **sequence type**  
40 derived type with the [SEQUENCE attribute](#) (7.5.2.3)
- 41 **3.126.1**  
42 **character sequence type**  
43 sequence type with no [allocatable](#) or [pointer components](#), and whose [components](#) are all default character or of  
44 another character sequence type

1 **3.126.2**  
2 **numeric sequence type**  
3 sequence type with no [allocatable](#) or [pointer components](#), and whose [components](#) are all default complex, default  
4 integer, default logical, default real, double precision real, or of another numeric sequence type

5 **3.127**  
6 **shape**  
7 array dimensionality of a data entity, represented as a rank-one array whose size is the [rank](#) of the data entity  
8 and whose elements are the extents of the data entity

**NOTE 3.6**

Thus the shape of a scalar data entity is an array with rank one and size zero.

9 **3.128**  
10 **simply contiguous**  
11  $\langle$ array designator or variable $\rangle$  satisfying the conditions specified in [9.5.4](#)

**NOTE 3.7**

These conditions are simple ones which make it clear that the designator or variable designates a [contiguous](#) array.

12 **3.129**  
13 **size**  
14  $\langle$ array $\rangle$  total number of elements in the array

15 **3.130**  
16 **specification expression**  
17 expression satisfying the requirements specified in [10.1.11](#), thus being suitable for use in specifications

18 **3.131**  
19 **specific name**  
20 name that is not a generic name

21 **3.132**  
22 **standard-conforming program**  
23 program that uses only those forms and relationships described in, and has an interpretation according to, this  
24 document

25 **3.133**  
26 **statement**  
27 sequence of one or more complete or partial lines satisfying a syntax rule that ends in *-stmt* ([6.3](#))

28 **3.133.1**  
29 **executable statement**  
30 *end-function-stmt*, *end-mp-subprogram-stmt*, *end-program-stmt*, *end-subroutine-stmt*, or statement that is a mem-  
31 ber of the syntactic class *executable-construct*, excluding those in the *block-specification-part* of a **BLOCK con-**  
32 **struct**

33 **3.133.2**  
34 **nonexecutable statement**  
35 statement that is not an [executable statement](#)

36 **3.134**  
37 **statement entity**  
38 entity whose identifier has the scope of a statement or part of a statement ([19.1](#), [19.4](#))



- 1 **3.135**  
2 **statement label**  
3 **label**  
4 unsigned positive number of up to five digits that refers to an individual statement (6.2.5)
- 5 **3.136**  
6 **storage sequence**  
7 contiguous sequence of [storage units](#) (19.5.3.2)
- 8 **3.137**  
9 **storage unit**  
10 [character storage unit](#), [numeric storage unit](#), [file storage unit](#), or [unspecified storage unit](#) (19.5.3.2)
- 11 **3.137.1**  
12 **character storage unit**  
13 unit of storage that holds a default character value (19.5.3.2)
- 14 **3.137.2**  
15 **numeric storage unit**  
16 unit of storage that holds a default real, default integer, or default logical value (19.5.3.2)
- 17 **3.137.3**  
18 **unspecified storage unit**  
19 unit of storage that holds a value that is not default character, default real, double precision real, default logical,  
20 or default complex (19.5.3.2)
- 21 **3.138**  
22 **stream file**  
23 file composed of a sequence of file storage units (12.1)
- 24 **3.139**  
25 **structure**  
26 [scalar data object](#) of [derived type](#) (7.5)
- 27 **3.139.1**  
28 **structure component**  
29 [component](#) of a structure
- 30 **3.139.2**  
31 **structure constructor**  
32 syntax ([structure-constructor](#), 7.5.10) that specifies a structure value or creates such a value
- 33 **3.140**  
34 **submodule**  
35 [program unit](#) that extends a [module](#) or another [submodule](#) (14.2.3)
- 36 **3.141**  
37 **subobject**  
38 portion of [data object](#) that can be referenced, and if it is a [variable](#) defined, independently of any other portion
- 39 **3.142**  
40 **subprogram**  
41 *function-subprogram* (R1529) or *subroutine-subprogram* (R1534)
- 42 **3.142.1**  
43 **external subprogram**  
44 subprogram that is not contained in a [main program](#), [module](#), [submodule](#), or another subprogram

- 1 **3.142.2**  
2 **internal subprogram**  
3 subprogram that is contained in a [main program](#) or another subprogram
- 4 **3.142.3**  
5 **module subprogram**  
6 subprogram that is contained in a [module](#) or [submodule](#) but is not an internal subprogram
- 7 **3.143**  
8 **subroutine**  
9 procedure invoked by a [CALL statement](#), by [defined assignment](#), or by some operations on derived-type entities
- 10 **3.143.1**  
11 **atomic subroutine**  
12 [intrinsic](#) subroutine that performs an action on its ATOM argument atomically
- 13 **3.143.2**  
14 **collective subroutine**  
15 [intrinsic](#) subroutine that performs a calculation on a [team](#) of images without requiring synchronization
- 16 **3.144**  
17 **target**  
18 entity that is [pointer associated](#) with a [pointer](#) (19.5.2.2), entity on the right-hand-side of a [pointer assignment](#)  
19 [statement](#) (R1033), or entity with the [TARGET](#) attribute (8.5.17)
- 20 **3.145**  
21 **team**  
22 ordered set of [images](#) created by execution of a [FORM TEAM statement](#), or the initial ordered set of all [images](#)
- 23 **3.145.1**  
24 **current team**  
25 team specified by the most recently executed [CHANGE TEAM statement](#) of a [CHANGE TEAM construct](#) that  
26 has not completed execution (11.1.5), or initial team if no [CHANGE TEAM construct](#) is being executed
- 27 **3.145.2**  
28 **initial team**  
29 team existing at the beginning of program execution, consisting of all images
- 30 **3.145.3**  
31 **parent team**  
32 ⟨team except for initial team⟩ current team at time of execution of the [FORM TEAM statement](#) that created the  
33 team (11.6.9)
- 34 **3.145.4**  
35 **team number**  
36 –1 which identifies the initial team, or positive integer that identifies a team within its [parent team](#)
- 37 **3.146**  
38 **transformational function**  
39 [intrinsic](#) function, or function in an [intrinsic](#) module, that is neither [elemental](#) nor an [inquiry function](#)
- 40 **3.147**  
41 **type**  
42 **data type**  
43 named category of data characterized by a set of values, a syntax for denoting these values, and a set of operations  
44 that interpret and manipulate the values (7.1)

- 1 **3.147.1**  
2 **abstract type**  
3 type with the [ABSTRACT attribute](#) (7.5.7.1)
- 4 **3.147.2**  
5 **declared type**  
6 type that a data entity is declared to have, either explicitly or implicitly (7.3.2, 10.1.9)
- 7 **3.147.3**  
8 **derived type**  
9 type defined by a type definition (7.5) or by an [intrinsic](#) module
- 10 **3.147.4**  
11 **dynamic type**  
12 type of a data entity at a particular point during execution of a program (7.3.2.3, 10.1.9)
- 13 **3.147.5**  
14 **extended type**  
15 type with the [EXTENDS attribute](#) (7.5.7.1)
- 16 **3.147.6**  
17 **extensible type**  
18 type that may be extended using the EXTENDS clause (7.5.7.1)
- 19 **3.147.7**  
20 **extension type**  
21 ⟨of one type with respect to another⟩ is the same type or is an extended type whose parent type is an extension  
22 type of the other type
- 23 **3.147.8**  
24 **intrinsic type**  
25 type defined by this document that is always accessible (7.4)
- 26 **3.147.9**  
27 **numeric type**  
28 one of the types integer, real, and complex
- 29 **3.147.10**  
30 **parent type**  
31 ⟨[extended type](#)⟩ type named in the EXTENDS clause
- 32 **3.147.11**  
33 **type compatible**  
34 compatibility of the type of one entity with respect to another for purposes such as [argument association](#), [pointer](#)  
35 [association](#), and allocation (7.3.2)
- 36 **3.147.12**  
37 **type parameter**  
38 value used to parameterize a type (7.2)
- 39 **3.147.12.1**  
40 **assumed type parameter**  
41 [length type parameter](#) that assumes the type parameter value from another entity

**NOTE 3.8**

The other entity is

- the selector for an [associate name](#),
- the [constant-expr](#) for a [named constant](#) of type character, or

## NOTE 3.8 (cont.)

- the [effective argument](#) for a [dummy argument](#).

- 1 **3.147.12.2**  
2 **deferred type parameter**  
3 [length type parameter](#) whose value can change during execution of a program and whose *type-param-value* is a  
4 colon
- 5 **3.147.12.3**  
6 **kind type parameter**  
7 type parameter whose value is required to be defaulted or given by a [constant expression](#)
- 8 **3.147.12.4**  
9 **length type parameter**  
10 type parameter whose value is permitted to be [assumed](#), [deferred](#), or given by a [specification expression](#)
- 11 **3.147.12.5**  
12 **type parameter inquiry**  
13 syntax (*type-param-inquiry*) that is used to inquire the value of a type parameter of a data object (9.4.5)
- 14 **3.147.12.6**  
15 **type parameter order**  
16 ordering of the type parameters of a type (7.5.3.2) used for derived-type specifiers (*derived-type-spec*, 7.5.9)
- 17 **3.148**  
18 **ultimate argument**  
19 nondummy entity with which a [dummy argument](#) is associated via a chain of argument associations (15.5.2.3)
- 20 **3.149**  
21 **undefined**  
22 [data object](#) does not have a valid value
- 23 **3.150**  
24 **undefined**  
25 [pointer](#) does not have a pointer association status of associated or [disassociated](#) (19.5.2.2)
- 26 **3.151**  
27 **unit**  
28 **input/output unit**  
29 means, specified by an *io-unit*, for referring to a file (12.5.1)
- 30 **3.152**  
31 **unlimited polymorphic**  
32 able to have any [dynamic type](#) during program execution (7.3.2.3)
- 33 **3.153**  
34 **unsaved**  
35 not having the [SAVE attribute](#) (8.5.16)
- 36 **3.154**  
37 **variable**  
38 [data entity](#) that can be [defined](#) and redefined during execution of a program
- 39 **3.154.1**  
40 **event variable**  
41 scalar variable of type [EVENT\\_TYPE](#) (16.10.2.10) from the intrinsic module [ISO\\_FORTRAN\\_ENV](#)

- 1     **3.154.2**  
2     **local variable**  
3     variable in a [scoping unit](#) that is not a [dummy argument](#) or part thereof, is not a global entity or part thereof,  
4     and is not an entity or part of an entity that is accessible outside that [scoping unit](#)
- 5     **3.154.3**  
6     **lock variable**  
7     scalar variable of type [LOCK\\_TYPE](#) (16.10.2.19) from the intrinsic module [ISO\\_FORTRAN\\_ENV](#)
- 8     **3.154.4**  
9     **team variable**  
10    scalar variable of type [TEAM\\_TYPE](#) (16.10.2.32) from the intrinsic module [ISO\\_FORTRAN\\_ENV](#)
- 11    **3.155**  
12    **vector subscript**  
13    *section-subscript* that is an array (9.5.3.3.2)
- 14    **3.156**  
15    **whole array**  
16    array component or array name without further qualification (9.5.2)

1

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## 4 Notation, conformance, and compatibility

### 4.1 Notation, symbols and abbreviated terms

#### 4.1.1 Syntax rules

1 Syntax rules describe the forms that Fortran [lexical tokens](#), statements, and constructs may take. These syntax rules are expressed in a variation of Backus-Naur form (BNF) with the following conventions.

- 2 • Characters from the Fortran character set ([6.1](#)) are interpreted literally as shown, except where otherwise noted.
- 3 • Lower-case italicized letters and words (often hyphenated and abbreviated) represent general syntactic classes for which particular syntactic entities shall be substituted in actual statements.

4 Common abbreviations used in syntactic terms are:

<i>arg</i>	for	argument	<i>attr</i>	for	attribute
<i>decl</i>	for	declaration	<i>def</i>	for	definition
<i>desc</i>	for	descriptor	<i>expr</i>	for	expression
<i>int</i>	for	integer	<i>op</i>	for	operator
<i>spec</i>	for	specifier	<i>stmt</i>	for	statement

- 5 • The syntactic metasympols used are:

<b>is</b>	introduces a syntactic class definition
<b>or</b>	introduces a syntactic class alternative
[ ]	encloses an optional item
[ ] ...	encloses an optionally repeated item that may occur zero or more times
■	continues a syntax rule

- 6 • Each syntax rule is given a unique identifying number of the form R<sub>snn</sub>, where s is a one- or two-digit clause number and nn is a two-digit sequence number within that clause. The syntax rules are distributed as appropriate throughout the text, and are referenced by number as needed. Some rules in [Clauses 5](#) and [6](#) are more fully described in later clauses; in such cases, the clause number s is the number of the later clause where the rule is repeated.
- 7 • The syntax rules are not a complete and accurate syntax description of Fortran, and cannot be used to generate a Fortran parser automatically; where a syntax rule is incomplete, it is restricted by corresponding constraints and text.

#### NOTE 4.1

An example of the use of the syntax rules is:

*digit-string*                      **is** *digit* [ *digit* ] ...

The following are examples of forms for a digit string allowed by the above rule:

*digit*  
*digit digit*  
*digit digit digit digit*  
*digit digit digit digit digit digit digit digit*

If particular entities are substituted for *digit*, actual digit strings might be:

**NOTE 4.1 (cont.)**

4  
67  
1999  
10243852

**4.1.2 Constraints**

- 1  
2 1 Each constraint is given a unique identifying number of the form Csmn, where s is a one- or two-digit clause  
3 number and nn is a two- or three-digit sequence number within that clause.
- 4 2 Often a constraint is associated with a particular syntax rule. Where that is the case, the constraint is annotated  
5 with the syntax rule number in parentheses. A constraint that is associated with a syntax rule constitutes part of  
6 the definition of the syntax term defined by the rule. It thus applies in all places where the syntax term appears.
- 7 3 Some constraints are not associated with particular syntax rules. The effect of such a constraint is similar to  
8 that of a restriction stated in the text, except that a processor is required to have the capability to detect and  
9 report violations of constraints (4.2). In some cases, a broad requirement is stated in text and a subset of the  
10 same requirement is also stated as a constraint. This indicates that a standard-conforming program is required to  
11 adhere to the broad requirement, but that a standard-conforming processor is required only to have the capability  
12 of diagnosing violations of the constraint.

**4.1.3 Assumed syntax rules**

- 13  
14 1 In order to minimize the number of additional syntax rules and convey appropriate constraint information, the  
15 following rules are assumed.
- 16 R401 *xyz-list* is *xyz* [ , *xyz* ] ...
- 17 R402 *xyz-name* is *name*
- 18 R403 *scalar-xyz* is *xyz*
- 19 C401 (R403) *scalar-xyz* shall be scalar.
- 20 2 The letters “*xyz*” stand for any syntactic class phrase. An explicit syntax rule for a term overrides an assumed  
21 rule.

**4.1.4 Syntax conventions and characteristics**

- 22  
23 1 Any syntactic class name ending in “-*stmt*” follows the source form statement rules: it shall be delimited by  
24 end-of-line or semicolon, and may be labeled unless it forms part of another statement (such as an **IF** or **WHERE**  
25 statement). Conversely, everything considered to be a source form statement is given a “-*stmt*” ending in the  
26 syntax rules.
- 27 2 The rules on statement ordering are described rigorously in the definition of *program-unit* (R502). Expression  
28 hierarchy is described rigorously in the definition of *expr* (R1022).
- 29 3 The suffix “-*spec*” is used consistently for specifiers, such as input/output statement specifiers. It also is used for  
30 type declaration attribute specifications (for example, “*array-spec*” in R815), and in a few other cases.
- 31 4 Where reference is made to a type parameter, including the surrounding parentheses, the suffix “-*selector*” is  
32 used. See, for example, “*kind-selector*” (R706) and “*length-selector*” (R722).



### 4.1.5 Text conventions

- 1 In descriptive text, an equivalent English word is frequently used in place of a syntactic term. Particular statements and attributes are identified in the text by an upper-case keyword, e.g., “END statement”. The descriptions of obsolescent features appear in a smaller type size.

#### NOTE 4.2

This sentence is an example of the type size used for obsolescent features.

## 4.2 Conformance

- 1 A **program** (5.2.2) is a **standard-conforming program** if it uses only those forms and relationships described herein and if the program has an interpretation according to this document. A **program unit** (5.2.1) conforms to this document if it can be included in a **program** in a manner that allows the **program** to be standard conforming.

- 2 A **processor** conforms to this document if:

- (1) it executes any **standard-conforming program** in a manner that fulfills the interpretations herein, subject to any limits that the **processor** may impose on the size and complexity of the **program**;
- (2) it contains the capability to detect and report the use within a submitted **program unit** of a form designated herein as obsolescent, insofar as such use can be detected by reference to the numbered syntax rules and constraints;
- (3) it contains the capability to detect and report the use within a submitted **program unit** of a form or relationship that is not permitted by the numbered syntax rules or constraints, including the deleted features described in Annex B;
- (4) it contains the capability to detect and report the use within a submitted **program unit** of an intrinsic type with a **kind type parameter** value not supported by the processor (7.4);
- (5) it contains the capability to detect and report the use within a submitted **program unit** of source form or characters not permitted by Clause 6;
- (6) it contains the capability to detect and report the use within a submitted **program** of name usage not consistent with the scope rules for names, labels, operators, and assignment symbols in Clause 19;
- (7) it contains the capability to detect and report the use within a submitted **program unit** of a **non-standard intrinsic** procedure (including one with the same name as a **standard intrinsic** procedure but with different requirements);
- (8) it contains the capability to detect and report the use within a submitted **program unit** of a **non-standard intrinsic** module;
- (9) it contains the capability to detect and report the use within a submitted **program unit** of a procedure from a **standard intrinsic** module, if the procedure is not defined by this document or the procedure has different requirements from those specified by this document; and
- (10) it contains the capability to detect and report the reason for rejecting a submitted **program**.

- 3 However, in a format specification that is not part of a **FORMAT statement** (13.2.1), a processor need not detect or report the use of deleted or obsolescent features, or the use of additional forms or relationships.

- 4 A standard-conforming processor may allow additional forms and relationships provided that such additions do not conflict with the standard forms and relationships. However, a standard-conforming processor may allow additional intrinsic procedures even though this could cause a conflict with the name of a procedure in a standard-conforming program. If such a conflict occurs and involves the name of an **external procedure**, the processor is permitted to use the intrinsic procedure unless the name has the **EXTERNAL attribute** (8.5.9) where it is used. A standard-conforming program shall not use **nonstandard intrinsic** procedures or modules that have been added by the processor.

- 5 Because a standard-conforming program may place demands on a processor that are not within the scope of this document or may include standard items that are not portable, such as **external procedures** defined by means

- 1 other than Fortran, conformance to this document does not ensure that a program will execute consistently on  
2 all or any standard-conforming processors.
- 3 6 The semantics of facilities that are identified as [processor dependent](#) are not completely specified in this document.  
4 They shall be provided, with methods or semantics determined by the processor.
- 5 7 The [processor](#) should be accompanied by documentation that specifies the limits it imposes on the size and com-  
6 plexity of a [program](#) and the means of reporting when these limits are exceeded, that defines the additional forms  
7 and relationships it allows, and that defines the means of reporting the use of additional forms and relationships  
8 and the use of deleted or obsolescent forms. In this context, the use of a deleted form is the use of an additional  
9 form.
- 10 8 The [processor](#) should be accompanied by documentation that specifies the methods or semantics of processor-  
11 dependent facilities.

## 12 4.3 Compatibility

### 13 4.3.1 Previous Fortran standards

- 14 1 Table [4.3](#) lists the previous editions of the Fortran International Standard, along with their informal names.

Table 4.3: Previous editions of the Fortran International Standard

Official designation	Informal name
ISO R 1539-1972	FORTRAN 66
ISO 1539-1980	FORTRAN 77
ISO/IEC 1539:1991	Fortran 90
ISO/IEC 1539-1:1997	Fortran 95
ISO/IEC 1539-1:2004	Fortran 2003
ISO/IEC 1539-1:2010	Fortran 2008

### 15 4.3.2 New intrinsic procedures

- 16 1 Each Fortran International Standard since ISO 1539:1980 (FORTRAN 77), defines more intrinsic procedures than  
17 the previous one. Therefore, a Fortran program conforming to an older standard might have a different inter-  
18 pretation under a newer standard if it invokes an external procedure having the same name as one of the new  
19 standard intrinsic procedures, unless that procedure is specified to have the [EXTERNAL attribute](#).

### 20 4.3.3 Fortran 2008 compatibility

- 21 1 Except as identified in this subclause, and except for the deleted features noted in Annex [B.2](#), this document is  
22 an upward compatible extension to the preceding Fortran International Standard, ISO/IEC 1539-1:2010 (Fortran  
23 2008). Any standard-conforming Fortran 2008 program that does not use any deleted features, and does not use  
24 any feature identified in this subclause as being no longer permitted, remains standard-conforming under this  
25 document.
- 26 2 Fortran 2008 specifies that the [IOSTAT=](#) variable shall be set to a processor-dependent negative value if the flush  
27 operation is not supported for the unit specified. This document specifies that the processor-dependent negative  
28 integer value shall be different from the [named constants](#) [IOSTAT\\_EOR](#) or [IOSTAT\\_END](#) from the intrinsic  
29 module [ISO\\_FORTRAN\\_ENV](#).
- 30 3 Fortran 2008 permitted a noncontiguous array that was supplied as an [actual argument](#) corresponding to a  
31 [contiguous INTENT \(INOUT\) dummy argument](#) in one iteration of a [DO CONCURRENT construct](#), without  
32 being previously defined in that iteration, to be defined in another iteration; this document does not permit this.

- 1 4 Fortran 2008 permitted a [pure](#) statement function to reference a [volatile](#) variable, and permitted a [local variable](#)  
2 of a [pure](#) subprogram or of a [BLOCK construct](#) within a [pure](#) subprogram to be [volatile](#) (provided it was not  
3 used); this document does not permit that.
- 4 5 Fortran 2008 permitted a [pure](#) function to have a result that has a [polymorphic](#) allocatable [ultimate component](#);  
5 this document does not permit that.
- 6 6 Fortran 2008 permitted a [PROTECTED TARGET](#) variable accessed by [use association](#) to be used as an [initial-  
7 data-target](#); this document does not permit that.
- 8 7 Fortran 2008 permitted a named constant to have declared type [LOCK\\_TYPE](#), or have a noncoarray [potential  
9 subobject component](#) with declared type [LOCK\\_TYPE](#); this document does not permit that.
- 10 8 Fortran 2008 permitted a [polymorphic](#) object to be finalized within a [DO CONCURRENT construct](#); this docu-  
11 ment does not permit that.

#### 12 4.3.4 Fortran 2003 compatibility

- 13 1 Except as identified in this subclause, this document is an upward compatible extension to ISO/IEC 1539-1:2004  
14 (Fortran 2003). Except as identified in this subclause, any standard-conforming Fortran 2003 program remains  
15 standard-conforming under this document.
- 16 2 Fortran 2003 permitted a [sequence type](#) to have [type parameters](#); that is not permitted by this document.
- 17 3 Fortran 2003 specified that array constructors and structure constructors of [finalizable](#) type are finalized. This  
18 document specifies that these constructors are not finalized.
- 19 4 The form produced by the G edit descriptor for some values and some input/output rounding modes differs from  
20 that specified by Fortran 2003.
- 21 5 Fortran 2003 required an [explicit interface](#) only for a procedure that was actually referenced in the scope, not  
22 merely passed as an [actual argument](#). This document requires an [explicit interface](#) for a procedure under the  
23 conditions listed in [15.4.2.2](#), regardless of whether the procedure is [referenced](#) in the scope.
- 24 6 Fortran 2003 permitted the [function result](#) of a [pure](#) function to be a [polymorphic](#) allocatable variable, to have  
25 a [polymorphic](#) allocatable [ultimate component](#), or to be [finalizable](#) by an impure [final subroutine](#). These are not  
26 permitted by this document.
- 27 7 Fortran 2003 permitted an [INTENT \(OUT\)](#) argument of a [pure](#) subroutine to be [polymorphic](#); that is not  
28 permitted by this document.
- 29 8 Fortran 2003 interpreted assignment to an allocatable variable from a nonconformable array as [intrinsic assign-  
30 ment](#), even when an elemental [defined assignment](#) was in scope; this document does not permit assignment from  
31 a nonconformable array in this context.
- 32 9 Fortran 2003 permitted a statement function to be of parameterized derived type; this document does not permit  
33 that.
- 34 10 Fortran 2003 permitted a [pure](#) statement function to reference a [volatile](#) variable, and permitted a [local variable](#)  
35 of a [pure](#) subprogram to be [volatile](#) (provided it was not used); this document does not permit that.

#### 36 4.3.5 Fortran 95 compatibility

- 37 1 Except as identified in this subclause, this document is an upward compatible extension to ISO/IEC 1539-1:1997  
38 (Fortran 95). Except as identified in this subclause, any standard-conforming Fortran 95 program remains  
39 standard-conforming under this document.
- 40 2 Fortran 95 permitted defined assignment between character strings of the same rank and different kinds. This  
41 document does not permit that if both of the different kinds are ASCII, ISO 10646, or default kind.

- 1 3 The following Fortran 95 features might have different interpretations in this document.
- 2 • Earlier Fortran standards had the concept of printing, meaning that column one of formatted output had  
3 special meaning for a processor-dependent (possibly empty) set of [external files](#). This could be neither  
4 detected nor specified by a standard-specified means. The interpretation of the first column is not specified  
5 by this document.
  - 6 • This document specifies a different output format for real zero values in list-directed and namelist output.
  - 7 • If the processor distinguishes between positive and negative real zero, this document requires different  
8 returned values for [ATAN2\(Y,X\)](#) when  $X < 0$  and  $Y$  is negative real zero and for [LOG\(X\)](#) and [SQRT\(X\)](#)  
9 when  $X$  is complex with  $X\%RE < 0$  and  $X\%IM$  is negative real zero.
  - 10 • This document has fewer restrictions on [constant expressions](#) than Fortran 95; this might affect whether a  
11 variable is considered to be an [automatic data object](#).
  - 12 • The form produced by the G edit descriptor with  $d$  equal to zero differs from that specified by Fortran 95  
13 for some values.

#### 14 4.3.6 Fortran 90 compatibility

- 15 1 Except for the deleted features noted in Annex [B.1](#), and except as identified in this subclause, this document  
16 is an upward compatible extension to ISO/IEC 1539:1991 (Fortran 90). Any standard-conforming Fortran 90  
17 program that does not use one of the deleted features remains standard-conforming under this document.
- 18 2 The [PAD= specifier](#) in the [INQUIRE statement](#) in this document returns the value UNDEFINED if there is no  
19 connection or the connection is for unformatted input/output. Fortran 90 specified YES.
- 20 3 Fortran 90 specified that if the second argument to [MOD](#) or [MODULO](#) was zero, the result was processor  
21 dependent. This document specifies that the second argument shall not be zero.
- 22 4 Fortran 90 permitted defined assignment between character strings of the same rank and different kinds. This  
23 document does not permit that if both of the different kinds are ASCII, ISO 10646, or default kind.
- 24 5 The following Fortran 90 features have different interpretations in this document:
- 25 • if the processor distinguishes between positive and negative real zero, the result value of the intrinsic function  
26 [SIGN](#) when the second argument is a negative real zero;
  - 27 • formatted output of negative real values (when the output value is zero);
  - 28 • whether an expression is a [constant expression](#) (thus whether a variable is considered to be an [automatic  
29 data object](#));
  - 30 • the G edit descriptor with  $d$  equal to zero for some values.

#### 31 4.3.7 FORTRAN 77 compatibility

- 32 1 Except for the deleted features noted in Annex [B.1](#), and except as identified in this subclause, this document is an  
33 upward compatible extension to ISO 1539:1980 (FORTRAN 77). Any standard-conforming FORTRAN 77 program  
34 that does not use one of the deleted features noted in Annex [B.1](#) and that does not depend on the differences  
35 specified here remains standard-conforming under this document. This document restricts the behavior for some  
36 features that were processor dependent in FORTRAN 77. Therefore, a standard-conforming FORTRAN 77 program  
37 that uses one of these processor-dependent features might have a different interpretation under this document, yet  
38 remain a standard-conforming program. The following FORTRAN 77 features might have different interpretations  
39 in this document.
- 40 • FORTRAN 77 permitted a processor to supply more precision derived from a default real constant than can  
41 be represented in a default real datum when the constant is used to initialize a double precision real data  
42 object in a [DATA statement](#). This document does not permit a processor this option.
  - 43 • If a named variable that was not in a [common block](#) was initialized in a [DATA statement](#) and did not have  
44 the [SAVE attribute](#) specified, FORTRAN 77 left its [SAVE attribute](#) processor dependent. This document  
45 specifies ([8.6.7](#)) that this named variable has the [SAVE attribute](#).

- 1 • FORTRAN 77 specified that the number of characters required by the input list was to be less than or equal  
2 to the number of characters in the record during formatted input. This document specifies (12.6.4.5.3) that  
3 the input record is logically padded with blanks if there are not enough characters in the record, unless the  
4 **PAD= specifier** with the value 'NO' is specified in an appropriate **OPEN** or **READ** statement.
- 5 • A value of 0 for a list item in a formatted **output statement** will be formatted in a different form for some  
6 G edit descriptors. In addition, this document specifies how rounding of values will affect the output field  
7 form, but FORTRAN 77 did not address this issue. Therefore, the form produced for certain combinations  
8 of values and G edit descriptors might differ from that produced by some FORTRAN 77 processors.
- 9 • FORTRAN 77 did not permit a processor to distinguish between positive and negative real zero; if the  
10 processor does so distinguish, the result will differ for the intrinsic function **SIGN** when the second argument  
11 is negative real zero, and formatted output of negative real zero will be different.

## 12 4.4 Deleted and obsolescent features

### 13 4.4.1 General

- 14 1 This document protects the users' investment in existing software by including all but six of the language elements  
15 of Fortran 90 that are not processor dependent. This document identifies two categories of outmoded features.  
16 The first category, deleted features, consists of features considered to have been redundant in FORTRAN 77 and  
17 largely unused in Fortran 90. Those in the second category, obsolescent features, are considered to have been  
18 redundant in Fortran 90 and Fortran 95, but are still frequently used.

### 19 4.4.2 Nature of deleted features

- 20 1 There are two groups of deleted features. The first group contains features for which better methods existed in  
21 FORTRAN 77; these features were not included in Fortran 95, Fortran 2003, or Fortran 2008, and are not included  
22 in this document. The second group contains features for which better methods existed in Fortran 90; these  
23 features were included in Fortran 2008, but are not included in this document.

### 24 4.4.3 Nature of obsolescent features

- 25 1 Better methods existed in Fortran 90 and Fortran 95 for each obsolescent feature. It is recommended that  
26 programmers use these better methods in new programs and convert existing code to these methods.
- 27 2 The obsolescent features are identified in the text of this document by a distinguishing type font (4.1.5).
- 28 3 A future revision of this document might delete an obsolescent feature if its use has become insignificant.

1

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## 5 Fortran concepts

### 5.1 High level syntax

- 1 This subclause introduces the syntax associated with [program units](#) and other Fortran concepts above the construct, statement, and expression levels and illustrates their relationships.

#### NOTE 5.1

Constraints and other information related to the rules that do not begin with R5 appear in the appropriate clause.

5	R501	<i>program</i>	<b>is</b>	<i>program-unit</i>
6				[ <i>program-unit</i> ] ...
7	R502	<i>program-unit</i>	<b>is</b>	<i>main-program</i>
8			<b>or</b>	<i>external-subprogram</i>
9			<b>or</b>	<i>module</i>
10			<b>or</b>	<i>submodule</i>
11			<b>or</b>	<i>block-data</i>
12	R1401	<i>main-program</i>	<b>is</b>	[ <i>program-stmt</i> ]
13				[ <i>specification-part</i> ]
14				[ <i>execution-part</i> ]
15				[ <i>internal-subprogram-part</i> ]
16				<i>end-program-stmt</i>
17	R503	<i>external-subprogram</i>	<b>is</b>	<i>function-subprogram</i>
18			<b>or</b>	<i>subroutine-subprogram</i>
19	R1529	<i>function-subprogram</i>	<b>is</b>	<i>function-stmt</i>
20				[ <i>specification-part</i> ]
21				[ <i>execution-part</i> ]
22				[ <i>internal-subprogram-part</i> ]
23				<i>end-function-stmt</i>
24	R1534	<i>subroutine-subprogram</i>	<b>is</b>	<i>subroutine-stmt</i>
25				[ <i>specification-part</i> ]
26				[ <i>execution-part</i> ]
27				[ <i>internal-subprogram-part</i> ]
28				<i>end-subroutine-stmt</i>
29	R1404	<i>module</i>	<b>is</b>	<i>module-stmt</i>
30				[ <i>specification-part</i> ]
31				[ <i>module-subprogram-part</i> ]
32				<i>end-module-stmt</i>
33	R1416	<i>submodule</i>	<b>is</b>	<i>submodule-stmt</i>
34				[ <i>specification-part</i> ]
35				[ <i>module-subprogram-part</i> ]
36				<i>end-submodule-stmt</i>
37	R1420	<i>block-data</i>	<b>is</b>	<i>block-data-stmt</i>
38				[ <i>specification-part</i> ]

1			<i>end-block-data-stmt</i>
2	R504	<i>specification-part</i>	<b>is</b> [ <i>use-stmt</i> ] ...
3			[ <i>import-stmt</i> ] ...
4			[ <i>implicit-part</i> ]
5			[ <i>declaration-construct</i> ] ...
6	R505	<i>implicit-part</i>	<b>is</b> [ <i>implicit-part-stmt</i> ] ...
7			<i>implicit-stmt</i>
8	R506	<i>implicit-part-stmt</i>	<b>is</b> <i>implicit-stmt</i>
9			<b>or</b> <i>parameter-stmt</i>
10			<b>or</b> <i>format-stmt</i>
11			<b>or</b> <i>entry-stmt</i>
12	R507	<i>declaration-construct</i>	<b>is</b> <i>specification-construct</i>
13			<b>or</b> <i>data-stmt</i>
14			<b>or</b> <i>format-stmt</i>
15			<b>or</b> <i>entry-stmt</i>
16			<b>or</b> <i>stmt-function-stmt</i>
17	R508	<i>specification-construct</i>	<b>is</b> <i>derived-type-def</i>
18			<b>or</b> <i>enum-def</i>
19			<b>or</b> <i>generic-stmt</i>
20			<b>or</b> <i>interface-block</i>
21			<b>or</b> <i>parameter-stmt</i>
22			<b>or</b> <i>procedure-declaration-stmt</i>
23			<b>or</b> <i>other-specification-stmt</i>
24			<b>or</b> <i>type-declaration-stmt</i>
25	R509	<i>execution-part</i>	<b>is</b> <i>executable-construct</i>
26			[ <i>execution-part-construct</i> ] ...
27	R510	<i>execution-part-construct</i>	<b>is</b> <i>executable-construct</i>
28			<b>or</b> <i>format-stmt</i>
29			<b>or</b> <i>entry-stmt</i>
30			<b>or</b> <i>data-stmt</i>
31	R511	<i>internal-subprogram-part</i>	<b>is</b> <i>contains-stmt</i>
32			[ <i>internal-subprogram</i> ] ...
33	R512	<i>internal-subprogram</i>	<b>is</b> <i>function-subprogram</i>
34			<b>or</b> <i>subroutine-subprogram</i>
35	R1407	<i>module-subprogram-part</i>	<b>is</b> <i>contains-stmt</i>
36			[ <i>module-subprogram</i> ] ...
37	R1408	<i>module-subprogram</i>	<b>is</b> <i>function-subprogram</i>
38			<b>or</b> <i>subroutine-subprogram</i>
39			<b>or</b> <i>separate-module-subprogram</i>
40	R1538	<i>separate-module-subprogram</i>	<b>is</b> <i>mp-subprogram-stmt</i>
41			[ <i>specification-part</i> ]
42			[ <i>execution-part</i> ]
43			[ <i>internal-subprogram-part</i> ]
44			<i>end-mp-subprogram-stmt</i>
45	R513	<i>other-specification-stmt</i>	<b>is</b> <i>access-stmt</i>
46			<b>or</b> <i>allocatable-stmt</i>



1		<b>or</b>	<i>asynchronous-stmt</i>
2		<b>or</b>	<i>bind-stmt</i>
3		<b>or</b>	<i>codimension-stmt</i>
4		<b>or</b>	<i>contiguous-stmt</i>
5		<b>or</b>	<i>dimension-stmt</i>
6		<b>or</b>	<i>external-stmt</i>
7		<b>or</b>	<i>intent-stmt</i>
8		<b>or</b>	<i>intrinsic-stmt</i>
9		<b>or</b>	<i>namelist-stmt</i>
10		<b>or</b>	<i>optional-stmt</i>
11		<b>or</b>	<i>pointer-stmt</i>
12		<b>or</b>	<i>protected-stmt</i>
13		<b>or</b>	<i>save-stmt</i>
14		<b>or</b>	<i>target-stmt</i>
15		<b>or</b>	<i>volatile-stmt</i>
16		<b>or</b>	<i>value-stmt</i>
17		<b>or</b>	<i>common-stmt</i>
18		<b>or</b>	<i>equivalence-stmt</i>
19	R514	<i>executable-construct</i>	<b>is</b> <i>action-stmt</i>
20			<b>or</b> <i>associate-construct</i>
21			<b>or</b> <i>block-construct</i>
22			<b>or</b> <i>case-construct</i>
23			<b>or</b> <i>change-team-construct</i>
24			<b>or</b> <i>critical-construct</i>
25			<b>or</b> <i>do-construct</i>
26			<b>or</b> <i>if-construct</i>
27			<b>or</b> <i>select-rank-construct</i>
28			<b>or</b> <i>select-type-construct</i>
29			<b>or</b> <i>where-construct</i>
30			<b>or</b> <i>forall-construct</i>
31	R515	<i>action-stmt</i>	<b>is</b> <i>allocate-stmt</i>
32			<b>or</b> <i>assignment-stmt</i>
33			<b>or</b> <i>backspace-stmt</i>
34			<b>or</b> <i>call-stmt</i>
35			<b>or</b> <i>close-stmt</i>
36			<b>or</b> <i>continue-stmt</i>
37			<b>or</b> <i>cycle-stmt</i>
38			<b>or</b> <i>deallocate-stmt</i>
39			<b>or</b> <i>endfile-stmt</i>
40			<b>or</b> <i>error-stop-stmt</i>
41			<b>or</b> <i>event-post-stmt</i>
42			<b>or</b> <i>event-wait-stmt</i>
43			<b>or</b> <i>exit-stmt</i>
44			<b>or</b> <i>fail-image-stmt</i>
45			<b>or</b> <i>flush-stmt</i>
46			<b>or</b> <i>form-team-stmt</i>
47			<b>or</b> <i>goto-stmt</i>
48			<b>or</b> <i>if-stmt</i>
49			<b>or</b> <i>inquire-stmt</i>
50			<b>or</b> <i>lock-stmt</i>
51			<b>or</b> <i>nullify-stmt</i>
52			<b>or</b> <i>open-stmt</i>
53			<b>or</b> <i>pointer-assignment-stmt</i>
54			<b>or</b> <i>print-stmt</i>

1 or *read-stmt*  
2 or *return-stmt*  
3 or *rewind-stmt*  
4 or *stop-stmt*  
5 or *sync-all-stmt*  
6 or *sync-images-stmt*  
7 or *sync-memory-stmt*  
8 or *sync-team-stmt*  
9 or *unlock-stmt*  
10 or *wait-stmt*  
11 or *where-stmt*  
12 or *write-stmt*  
13 or *computed-goto-stmt*  
14 or *forall-stmt*

## 15 5.2 Program unit concepts

### 16 5.2.1 Program units and scoping units

- 17 1 Program units are the fundamental components of a Fortran program. A program unit is a main program, an  
18 external subprogram, a module, a submodule, or a block data program unit.
- 19 2 A subprogram is a function subprogram or a subroutine subprogram. A module contains definitions that can be  
20 made accessible to other program units. A submodule is an extension of a module; it may contain the definitions  
21 of procedures declared in a module or another submodule. A block data program unit is used to specify initial values for  
22 data objects in named common blocks.
- 23 3 Each type of program unit is described in Clause 14 or 15.
- 24 4 A program unit consists of a set of nonoverlapping scoping units.

#### NOTE 5.2

The module or submodule containing a module subprogram is the host scoping unit of the module subprogram. The containing main program or subprogram is the host scoping unit of an internal subprogram.

An internal procedure is local to its host in the sense that its name is accessible within the host scoping unit and all its other internal procedures but is not accessible elsewhere.

### 25 5.2.2 Program

- 26 1 A program shall consist of exactly one main program, any number (including zero) of other kinds of program units,  
27 any number (including zero) of external procedures, and any number (including zero) of other entities defined by  
28 means other than Fortran. The main program shall be defined by a Fortran *main-program program-unit* or by  
29 means other than Fortran, but not both.

### 30 5.2.3 Procedure

- 31 1 A procedure is either a function or a subroutine. Invocation of a function in an expression causes a value to be  
32 computed which is then used in evaluating the expression.
- 33 2 A procedure that is not pure might change the program state by changing the value of accessible data objects or  
34 procedure pointers.
- 35 3 Procedures are described further in Clause 15.

## 5.2.4 Module

- 1 A **module** contains (or accesses from other modules) definitions that can be made accessible to other **program units**. These definitions include **data object declarations**, type definitions, procedure definitions, and **interface blocks**. Modules are further described in Clause 14.

## 5.2.5 Submodule

- 1 A **submodule** extends a **module** or another **submodule**.
- 2 It may provide definitions (15.6) for procedures whose **interfaces** are declared (15.4.3.2) in an ancestor module or submodule. It may also contain declarations and definitions of other entities, which are accessible in its **descendants**. An entity declared in a submodule is not accessible by **use association** unless it is a module procedure whose **interface** is declared in the ancestor module. Submodules are further described in Clause 14.

### NOTE 5.3

A submodule has access to entities in its parent module or submodule by **host association**.

## 5.3 Execution concepts

### 5.3.1 Statement classification

- 1 Each Fortran statement is classified as either an **executable statement** or a **nonexecutable statement**.
- 2 An **executable statement** is an instruction to perform or control an action. Thus, the executable statements of a **program unit** determine the behavior of the **program unit**.
- 3 **Nonexecutable statements** are used to configure the program environment in which actions take place.

### 5.3.2 Statement order

Table 5.1: Requirements on statement ordering

PROGRAM, FUNCTION, SUBROUTINE, MODULE, SUBMODULE, or BLOCK DATA statement		
USE statements		
IMPORT statements		
FORMAT and ENTRY statements	IMPLICIT NONE	
	PARAMETER statements	IMPLICIT statements
	PARAMETER and DATA statements	Derived-type definitions, interface blocks, type declaration statements, enumeration definitions, procedure declarations, specification statements, and statement function statements
	DATA statements	Executable constructs
<b>CONTAINS statement</b>		
Internal subprograms or module subprograms		
END statement		

1 The syntax rules of subclause 5.1 specify the statement order within **program units** and **subprograms**. These  
 2 rules are illustrated in Table 5.1 and Table 5.2. Table 5.1 shows the ordering rules for statements and applies  
 3 to all **program units**, **subprograms**, and **interface bodies**. Vertical lines delineate varieties of statements that  
 4 may be interspersed and horizontal lines delineate varieties of statements that shall not be interspersed. **Internal**  
 5 or **module subprograms** shall follow a **CONTAINS statement**. Between **USE** and **CONTAINS** statements in a  
 6 **subprogram**, nonexecutable statements generally precede executable statements, although the **ENTRY statement**,  
 7 **FORMAT statement**, and **DATA statement** may appear among the executable statements. Table 5.2 shows which  
 8 statements are allowed in some kinds of **scoping units**.

Table 5.2: Statements allowed in scoping units

Statement type	Kind of scoping unit						
	Main program	Module or submodule	Block data	External subprogram	Module subprogram	Internal subprogram	Interface body
<b>USE</b>	Yes	Yes	Yes	Yes	Yes	Yes	Yes
<b>IMPORT</b>	No	Submodule	No	No	Yes	Yes	Yes
<b>ENTRY</b>	No	No	No	Yes	Yes	No	No
<b>FORMAT</b>	Yes	No	No	Yes	Yes	Yes	No
Misc. decl.s <sup>1</sup>	Yes	Yes	Yes	Yes	Yes	Yes	Yes
<b>DATA</b>	Yes	Yes	Yes	Yes	Yes	Yes	No
Derived-type	Yes	Yes	Yes	Yes	Yes	Yes	Yes
Interface	Yes	Yes	No	Yes	Yes	Yes	Yes
Executable	Yes	No	No	Yes	Yes	Yes	No
<b>CONTAINS</b>	Yes	Yes	No	Yes	Yes	No	No
Statement function	Yes	No	No	Yes	Yes	Yes	No

(1) Miscellaneous declarations are **PARAMETER statements**, **IMPLICIT statements**, **type declaration statements**, enumeration definitions, **procedure declaration statements**, and specification statements.

### 9 5.3.3 The END statement

10 1 Each **program unit**, **module subprogram**, and **internal subprogram** shall have exactly one **END statement**. The  
 11 *end-program-stmt*, *end-function-stmt*, *end-subroutine-stmt*, and *end-mp-subprogram-stmt* statements are execut-  
 12 able, and may be **branch target statements** (11.2). Executing an *end-program-stmt* initiates **normal termination**.  
 13 Executing an *end-function-stmt*, *end-subroutine-stmt*, or *end-mp-subprogram-stmt* is equivalent to executing a  
 14 *return-stmt* with no *scalar-int-expr*.

15 2 The *end-module-stmt*, *end-submodule-stmt*, and *end-block-data-stmt* statements are nonexecutable.

### 16 5.3.4 Program execution

17 1 Execution of a program consists of the asynchronous execution of a fixed number (which may be one) of its **images**.  
 18 Each **image** has its own execution state, floating-point status (17.7), and set of **data objects**, **input/output units**,  
 19 and procedure pointers. The **image index** that identifies an **image** is an integer value in the range one to the  
 20 number of **images** in a **team**.

21 2 A team is an ordered set of **images** that is either the initial team, consisting of all images, or a subset of a parent  
 22 team formed by execution of a **FORM TEAM statement**. The initial team has no parent; every other team has a  
 23 unique parent team. Within its parent team, each team is identified by its team number; this is the integer value  
 24 that was specified in the **FORM TEAM statement**.

25 3 During execution, each image has a current team, which is only changed by execution of **CHANGE TEAM**  
 26 and **END TEAM** statements. Image indices, and thus coindexing of variable names with an *image-selector*, are

1 relative to the current team unless a different team is specified. Initially, the current team is the initial team.

#### NOTE 5.4

Fortran control constructs (11.1, 11.2) control the progress of execution in each *image*. *Image control statements* (11.6.1) affect the relative progress of execution between *images*. *Coarrays* (5.4.7) provide a mechanism for accessing data on one *image* from another *image*.

#### NOTE 5.5

A processor might allow the number of *images* to be chosen at compile time, link time, or run time. It might be the same as the number of CPUs but this is not required. Compiling for a single *image* might permit the optimizer to eliminate overhead associated with parallel execution. A program that makes assumptions about the number of *images* is unlikely to be portable.

### 2 5.3.5 Execution sequence

3 1 Following the creation of a fixed number of *images*, execution begins on each *image*. *Image* execution is a  
 4 sequence, in time, of actions. Actions take place during execution of the statement that performs them (except  
 5 when explicitly stated otherwise). Segments (11.6.2) executed by a single *image* are totally ordered, and segments  
 6 executed by separate *images* are partially ordered by *image control statements* (11.6.1).

7 2 If the program contains a Fortran *main program*, each *image* begins execution with the first executable construct  
 8 of the *main program*. The execution of a *main program* or *subprogram* involves execution of the executable  
 9 constructs within its *scoping unit*. When a Fortran procedure is invoked, the *specification expressions* within  
 10 the *specification-part* of the invoked procedure, if any, are evaluated in a processor dependent order. Thereafter,  
 11 execution proceeds to the first executable construct appearing within the *scoping unit* of the procedure after the  
 12 invoked entry point. With the following exceptions, the effect of execution is as if the executable constructs are  
 13 executed in the order in which they appear in the *main program* or *subprogram* until a *STOP*, *ERROR STOP*,  
 14 *RETURN*, or *END* statement is executed.

- 15 • Execution of a branching statement (11.2) changes the execution sequence. These statements explicitly  
 16 specify a new starting place for the execution sequence.
- 17 • *DO* constructs, *IF* constructs, *SELECT CASE* constructs, *SELECT RANK* constructs, and *SELECT TYPE*  
 18 *constructs* contain an internal statement structure and execution of these constructs involves implicit in-  
 19 ternal transfer of control. See Clause 11 for the detailed semantics of each of these constructs.
- 20 • A *BLOCK* construct may contain *specification expressions*; see 11.1.4 for detailed semantics of this construct.
- 21 • An *END=*, *ERR=*, or *EOR=* specifier (12.11) might result in a branch.
- 22 • An alternate return might result in a branch.

### 23 5.3.6 Image execution states

24 1 There are three image execution states: active, stopped, and failed. An *image* that has initiated *normal termin-*  
 25 *ation* of execution is a *stopped image*. An *image* that has ceased participating in program execution but has not  
 26 initiated termination is a *failed image*. All other *images* are *active images*.

27 2 A *failed image* remains failed for the remainder of the execution of the program. The conditions that cause an  
 28 *image* to fail are processor dependent.

29 3 Defining a *coindexed object* on a *failed image* has no effect other than defining the *stat-variable*, if one appears,  
 30 with the value *STAT\_FAILED\_IMAGE* (16.10.2.26). The value of a reference to a *coindexed object* on a *failed*  
 31 *image* is processor dependent. Execution continues after such a reference.

32 4 When an image fails during the execution of a segment, a data object on a nonfailed image becomes undefined  
 33 if it is not a *lock variable* or *event variable* and it might be defined or become undefined by execution of a  
 34 statement of the segment other than an invocation of an *atomic subroutine* with the object as an *actual argument*  
 35 corresponding to the *ATOM dummy argument*.

### 5.3.7 Termination of execution

- 1 Termination of execution of a program is either normal termination or error termination. Normal termination occurs only when all **images** initiate normal termination and occurs in three steps: initiation, synchronization, and completion. In this case, all **images** synchronize execution at the second step so that no image starts the completion step until all **images** have finished the initiation step. Error termination occurs when any **image** initiates error termination. Once error termination has been initiated on an **image**, error termination is initiated on all **images** that have not already initiated error termination. Termination of execution of the program occurs when all **images** have terminated execution.
- 2 Normal termination of execution of an **image** is initiated when a **STOP statement** or *end-program-stmt* is executed. Normal termination of execution of an **image** also may be initiated during execution of a procedure defined by a **companion processor** (ISO/IEC 9899:2011 5.1.2.2.3 and 7.22.4.4). If normal termination of execution is initiated within a Fortran **program unit** and the program incorporates procedures defined by a **companion processor**, the process of execution termination shall include the effect of executing the C `exit()` function (ISO/IEC 9899:2011 7.22.4.4) during the completion step.
- 3 Error termination of execution of an **image** is initiated if an **ERROR STOP statement** is executed or as specified elsewhere in this document. When error termination on an **image** has been initiated, the processor should initiate error termination on other **images** as quickly as possible.
- 4 If the processor supports the concept of a process exit status, it is recommended that error termination initiated other than by an **ERROR STOP statement** supplies a processor-dependent nonzero value as the process exit status.

#### NOTE 5.6

As well as in the circumstances specified in this document, error termination might be initiated by means other than Fortran.

#### NOTE 5.7

If an **image** has initiated normal termination, its data remain available for possible reference or definition by other **images** that are still executing.

## 5.4 Data concepts

### 5.4.1 Type

- 1 A **type** is a named categorization of data that, together with its **type parameters**, determines the set of values, syntax for denoting these values, and the set of operations that interpret and manipulate the values. This central concept is described in 7.1.
- 2 A **type** is either an **intrinsic type** or a **derived type**.

#### 5.4.1.1 Intrinsic type

- 1 The **intrinsic types** are integer, real, complex, character, and logical. The properties of **intrinsic types** are described in 7.4.
- 2 All **intrinsic types** have a **kind type parameter** called KIND, which determines the representation method for the specified type. The **intrinsic type** character also has a **length type parameter** called LEN, which determines the length of the character string.

#### 5.4.1.2 Derived type

- 1 **Derived types** may be parameterized. A scalar **object** of **derived type** is a **structure**; assignment of structures is defined intrinsically (10.2.1.3), but there are no **intrinsic** operations for structures. For each **derived type**, a

1 [structure constructor](#) is available to create values (7.5.10). In addition, [objects](#) of [derived type](#) may be used as  
2 procedure arguments and function results, and may appear in input/output lists. If additional operations are  
3 needed for a [derived type](#), they shall be defined by procedures (10.1.6).

4 2 [Derived types](#) are described further in 7.5.

## 5 5.4.2 Data value

6 1 Each [intrinsic type](#) has associated with it a set of values that a datum of that type may take, depending on the  
7 values of the type parameters. The values for each [intrinsic type](#) are described in 7.4. The values that [objects](#) of  
8 a [derived type](#) may assume are determined by the type definition, type parameter values, and the sets of values  
9 of its components.

## 10 5.4.3 Data entity

### 11 5.4.3.1 General

12 1 A [data entity](#) has a type and type parameters; it might have a data value (an exception is an undefined variable).  
13 Every [data entity](#) has a [rank](#) and is thus either a scalar or an array.

14 2 A [data entity](#) that is the result of the execution of a [function reference](#) is called the function result.

### 15 5.4.3.2 Data object

16 1 A [data object](#) is either a constant, variable, or a [subobject](#) of a constant. The type and type parameters of a  
17 named [data object](#) may be specified explicitly (8.2) or implicitly (8.7).

18 2 [Subobjects](#) are portions of [data objects](#) that may be referenced and defined (variables only) independently of the  
19 other portions.

20 3 These include portions of arrays (array elements and array sections), portions of character strings (substrings),  
21 portions of complex [objects](#) (real and imaginary parts), and portions of [structures](#) (components). [Subobjects](#)  
22 are themselves [data objects](#), but [subobjects](#) are referenced only by [object designators](#) or [intrinsic](#) functions. A  
23 [subobject](#) of a variable is a variable. [Subobjects](#) are described in Clause 9.

24 4 The following [objects](#) are referenced by a name:

- 25 • a named scalar (a scalar object);
- a named array (an array object).

26 5 The following [subobjects](#) are referenced by an [object designator](#):

- an array element (a scalar subobject);
- an array section (an array subobject);
- 27 • a [complex part designator](#) (the real or imaginary part of a complex object);
- a structure component (a scalar or an array subobject);
- a substring (a scalar subobject).

### 28 5.4.3.2.1 Variable

29 1 A [variable](#) can have a value or be undefined; during execution of a program it can be defined, redefined, or become  
30 undefined.

31 2 A [local variable](#) of a [module](#), [submodule](#), [main program](#), [subprogram](#), or [BLOCK construct](#) is accessible only in  
32 that [scoping unit](#) or construct and in any contained [scoping units](#) and constructs.

#### NOTE 5.8

A [subobject](#) of a [local variable](#) is also a [local variable](#).

**NOTE 5.8 (cont.)**

A [local variable](#) cannot be in [COMMON](#) or have the [BIND attribute](#), because [common blocks](#) and [variables](#) with the [BIND attribute](#) are global entities.

**1 5.4.3.2.2 Constant**

2 1 A constant is either a [named constant](#) or a [literal constant](#).

3 2 Named constants are defined using the [PARAMETER attribute](#) ([8.5.13](#), [8.6.11](#)). The syntax of literal constants  
4 is described in [7.4](#).

**5 5.4.3.2.3 Subobject of a constant**

6 1 A [subobject](#) of a [constant](#) is a portion of a [constant](#).

7 2 In an [object designator](#) for a [subobject](#) of a [constant](#), the portion referenced may depend on the value of a [variable](#).

**NOTE 5.9**

For example, given:

```
CHARACTER (LEN = 10), PARAMETER :: DIGITS = '0123456789'  
CHARACTER (LEN = 1)           :: DIGIT  
INTEGER :: I  
...  
DIGIT = DIGITS (I:I)
```

DIGITS is a [named constant](#) and DIGITS (I:I) designates a [subobject](#) of the [constant](#) DIGITS.

**8 5.4.3.3 Expression**

9 1 An expression ([10.1](#)) produces a [data entity](#) when evaluated. An expression represents either a [data object](#)  
10 [reference](#) or a computation; it is formed from operands, operators, and parentheses. The type, type parameters,  
11 value, and [rank](#) of an expression result are determined by the rules in [Clause 10](#).

**12 5.4.3.4 Function reference**

13 1 A [function reference](#) produces a [data entity](#) when the function is executed during expression evaluation. The  
14 type, type parameters, and [rank](#) of a function result are determined by the [interface](#) of the function ([15.3.3](#)). The  
15 value of a function result is determined by execution of the function.

**16 5.4.4 Definition of objects and pointers**

17 1 When an [object](#) is given a valid value during program execution, it becomes [defined](#). This is often accomplished  
18 by execution of an [assignment](#) or [input](#) statement. When a variable does not have a predictable value, it is  
19 [undefined](#).

20 2 Similarly, when a pointer is associated with a [target](#) or nullified, its [pointer association](#) status becomes [defined](#).  
21 When the association status of a pointer is not predictable, its [pointer association](#) status is [undefined](#).

22 3 [Clause 19](#) describes the ways in which variables become [defined](#) and [undefined](#) and the association status of  
23 pointers becomes [defined](#) and [undefined](#).

**24 5.4.5 Reference**

25 1 A [data object](#) is [referenced](#) when its value is required during execution. A procedure is [referenced](#) when it is  
26 executed.



1 2 The appearance of a [data object designator](#) or [procedure designator](#) as an [actual argument](#) does not constitute  
2 a [reference](#) to that [data object](#) or procedure unless such a [reference](#) is necessary to complete the specification of  
3 the [actual argument](#).

#### 4 5.4.6 Array

5 1 An array may have up to fifteen dimensions minus its [corank](#), and any [extent](#) in any dimension. The [size](#) of an  
6 array is the total number of elements, which is equal to the product of the [extents](#). An array may have zero  
7 size. The [shape](#) of an array is determined by its [rank](#) and its [extent](#) in each dimension, and is represented as  
8 a rank-one array whose elements are the extents. All named arrays shall be declared, and the [rank](#) of a named  
9 array is specified in its declaration. Except for an [assumed-rank](#) array, the [rank](#) of a named array, once declared,  
10 is constant; the extents may be constant or may vary during execution.

11 2 Any [intrinsic](#) operation defined for scalar [objects](#) may be applied to [conformable objects](#). Such operations are  
12 performed [elementally](#) to produce a resultant array [conformable](#) with the array operands. If an [elemental](#) operation  
13 is intrinsically pure or is implemented by a pure [elemental function](#) (15.8), the element operations may be  
14 performed simultaneously or in any order.

15 3 A rank-one array may be constructed from scalars and other arrays and may be reshaped into any allowable array  
16 shape (7.8).

17 4 Arrays may be of any type and are described further in 9.5.

#### 18 5.4.7 Coarray

19 1 A [coarray](#) is a [data entity](#) that has nonzero [corank](#); it can be directly referenced or defined by other [images](#). It  
20 may be a scalar or an array.

21 2 For each [coarray](#) on an [image](#), there is a corresponding [coarray](#) with the same type, type parameters, and [bounds](#)  
22 on every other [image](#) that is in the same [team](#).

23 3 The set of corresponding [coarrays](#) on all [images](#) in a [team](#) is arranged in a rectangular pattern. The dimensions of  
24 this pattern are the [codimensions](#); the number of [codimensions](#) is the [corank](#). The bounds for each [codimension](#)  
25 are the [cobounds](#).

##### NOTE 5.10

If the total number of [images](#) is not a multiple of the product of the sizes of each but the rightmost of the [codimensions](#), the rectangular pattern will be incomplete.

26 4 A [coarray](#) on any [image](#) can be accessed directly by using [cosubscripts](#). On its own [image](#), a [coarray](#) can also be  
27 accessed without use of [cosubscripts](#).

28 5 A [subobject](#) of a [coarray](#) is a [coarray](#) if it does not have any [cosubscripts](#), [vector subscripts](#), [allocatable](#) component  
29 selection, or pointer component selection.

30 6 For a [coindexed object](#), its [cosubscript](#) list determines the [image index](#) (9.6) in the same way that a subscript list  
31 determines the subscript order value for an [array element](#) (9.5.3.2).

32 7 [Intrinsic](#) procedures are provided for mapping between an [image index](#) and a list of [cosubscripts](#).

##### NOTE 5.11

The mechanism for an [image](#) to reference and define a [coarray](#) on another [image](#) might vary according to the hardware. On a shared-memory machine, a [coarray](#) on an [image](#) and the corresponding [coarrays](#) on other [images](#) could be implemented as a sequence of arrays with evenly spaced starting addresses. On a distributed-memory machine with separate physical memory for each [image](#), a processor might store a [coarray](#) at the same virtual address in each physical memory.

**NOTE 5.12**

Except in contexts where [coindexed objects](#) are disallowed, accessing a [coarray](#) on its own [image](#) by using a set of [cosubscripts](#) that specify that [image](#) has the same effect as accessing it without [cosubscripts](#). In particular, the segment ordering rules ([11.6.2](#)) apply whether or not [cosubscripts](#) are used to access the [coarray](#).

**5.4.8 Established coarrays**

- 1 A nonallocatable [coarray](#) with the [SAVE](#) attribute is [established](#) in the [initial team](#).
- 2 An allocated allocatable [coarray](#) is [established](#) in the [team](#) in which it was allocated. An unallocated allocatable [coarray](#) is not established.
- 3 A [coarray](#) that is an [associating entity](#) in a [coarray-association](#) of a [CHANGE TEAM](#) statement is [established](#) in the [team](#) of its [CHANGE TEAM](#) construct.
- 4 A nonallocatable [coarray](#) that is an [associating entity](#) in an [ASSOCIATE](#), [SELECT RANK](#), or [SELECT TYPE](#) construct is [established](#) in the team in which the [ASSOCIATE](#), [SELECT RANK](#), or [SELECT TYPE](#) statement is executed.
- 5 A nonallocatable [coarray](#) that is a dummy argument or host associated with a dummy argument is [established](#) in the [team](#) in which the procedure was invoked. A [coarray dummy argument](#) is not established in any [ancestor team](#) even if the corresponding [actual argument](#) is [established](#) in one or more of them.

**5.4.9 Pointer**

- 1 A [pointer](#) has an association status which is either associated, [disassociated](#), or undefined ([19.5.2.2](#)).
- 2 A [pointer](#) that is not associated shall not be referenced or defined.
- 3 If a [data pointer](#) is an [array](#), the [rank](#) is declared, but the [bounds](#) are determined when it is associated with a [target](#).

**5.4.10 Allocatable variables**

- 1 The allocation status of an [allocatable](#) variable is either allocated or unallocated. An [allocatable](#) variable becomes allocated as described in [9.7.1.3](#). It becomes unallocated as described in [9.7.3.2](#).
- 2 An unallocated [allocatable](#) variable shall not be [referenced](#) or [defined](#).
- 3 If an [allocatable](#) variable is an array, the [rank](#) is declared, but the [bounds](#) are determined when it is allocated. If an [allocatable](#) variable is a [coarray](#), the [corank](#) is declared, but the [cobounds](#) are determined when it is allocated.

**5.4.11 Storage**

- 1 Many of the facilities of this document make no assumptions about the physical storage characteristics of [data objects](#). However, [program units](#) that include [storage association](#) dependent features shall observe the storage restrictions described in [19.5.3](#).

**5.5 Fundamental concepts****5.5.1 Names and designators**

- 1 A [name](#) is used to identify a program constituent, such as a [program unit](#), named [variable](#), named [constant](#), [dummy argument](#), or [derived type](#).

1 2 A [designator](#) is used to identify a program constituent or a part thereof.

## 2 5.5.2 Statement keyword

3 1 A [statement keyword](#) is not a reserved word; that is, a name with the same spelling is allowed. In the syntax  
4 rules, such keywords appear literally. In descriptive text, this meaning is denoted by the term “keyword” without  
5 any modifier. Examples of [statement keywords](#) are IF, READ, UNIT, KIND, and INTEGER.

## 6 5.5.3 Other keywords

7 1 Other keywords denote names that identify items in a list. In this case, items are identified by a preceding  
8 [keyword](#)= rather than their position within the list.

9 2 An [argument keyword](#) is the name of a [dummy argument](#) in the [interface](#) for the procedure being referenced, and  
10 may appear in an [actual argument](#) list. A [type parameter keyword](#) is the name of a type parameter in the type  
11 being specified, and may appear in a type parameter list. A [component keyword](#) is the name of a component in  
12 a [structure constructor](#).

13 R516 *keyword* is *name*

### NOTE 5.13

Use of keywords rather than position to identify items in a list can make such lists more readable and allows them to be reordered. This facilitates specification of a list in cases where optional items are omitted.

## 14 5.5.4 Association

15 1 [Name association](#) (19.5.1) permits an entity to be identified by different names in the same [scoping unit](#) or by  
16 the same name or different names in different [scoping units](#).

17 2 [Pointer association](#) (19.5.2) between a pointer and a target allows the target to be denoted by the pointer.

18 3 [Storage association](#) (19.5.3) causes different entities to use the same storage.

19 4 [Inheritance association](#) (19.5.4) occurs between components of the parent component and components inherited  
20 by type extension.

## 21 5.5.5 Intrinsic

22 1 All [intrinsic](#) types, procedures, assignments, and operators may be used in any [scoping unit](#) without further  
23 definition or specification. [Intrinsic](#) modules (16.10, 17, 18.2) may be accessed by [use association](#).

## 24 5.5.6 Operator

25 1 This document specifies a number of [intrinsic](#) operators (e.g., the arithmetic operators +, −, \*, /, and \*\* with  
26 numeric operands and the logical operators .AND., .OR., etc. with logical operands). Additional operators may  
27 be defined within a program (7.5.5, 15.4.3.4).

## 28 5.5.7 Companion processors

29 1 A [processor](#) has one or more [companion processors](#). A [companion processor](#) may be a mechanism that references  
30 and defines such entities by a means other than Fortran (15.6.3), it may be the [Fortran processor](#) itself, or it may  
31 be another [Fortran processor](#). If there is more than one [companion processor](#), the means by which the [Fortran](#)  
32 [processor](#) selects among them are [processor dependent](#).

33 2 If a [procedure](#) is defined by means of a [companion processor](#) that is not the [Fortran processor](#) itself, this document  
34 refers to the C function that defines the [procedure](#), although the [procedure](#) need not be defined by means of the

1 C programming language.

**NOTE 5.14**

A [companion processor](#) might or might not be a mechanism that conforms to the requirements of ISO/IEC 9899:2011. If it does, [5.3.7](#) states that a program unit that is defined by means other than Fortran and that initiates normal termination is required to include the effect of executing the C `exit()` function.

For example, a [processor](#) might allow a [procedure](#) defined by some language other than Fortran or C to be invoked if it can be described by a C prototype as defined in 6.7.6.3 of ISO/IEC 9899:2011.

## 6 Lexical tokens and source form

### 6.1 Processor character set

#### 6.1.1 Characters

1 The processor character set is processor dependent. Each character in a processor character set is either a control character or a graphic character. The set of graphic characters is further divided into letters (6.1.2), digits (6.1.3), underscore (6.1.4), special characters (6.1.5), and other characters (6.1.6).

2 The letters, digits, underscore, and special characters make up the Fortran character set. Together, the set of letters, digits, and underscore define the syntax class *alphanumeric-character*.

```
R601  alphanumeric-character    is  letter
                                     or  digit
                                     or  underscore
```

3 Except for the currency symbol, the graphics used for the characters shall be as given in 6.1.2, 6.1.3, 6.1.4, and 6.1.5. However, the style of any graphic is not specified.

#### 6.1.2 Letters

1 The twenty-six letters are:

```
2      A B C D E F G H I J K L M N O P Q R S T U V W X Y Z
```

3 The set of letters defines the syntactic class *letter*. The processor character set shall include lower-case and upper-case letters. A lower-case letter is equivalent to the corresponding upper-case letter in *program units* except in a *character context* (3.21).

##### NOTE 6.1

The following statements are equivalent:

```
CALL BIG_COMPLEX_OPERATION (NDATE)
call big_complex_operation (ndate)
Call Big_Complex_Operation (NDate)
```

#### 6.1.3 Digits

1 The ten digits are:

```
2      0 1 2 3 4 5 6 7 8 9
```

3 The ten digits define the syntactic class *digit*.

#### 6.1.4 Underscore

```
R602  underscore                is  _
```

#### 6.1.5 Special characters

1 The special characters are shown in Table 6.1.

Table 6.1: Special characters

Character	Name of character	Character	Name of character
	Blank	;	Semicolon
=	Equals	!	Exclamation point
+	Plus	"	Quotation mark or quote
-	Minus	%	Percent
*	Asterisk	&	Ampersand
/	Slash	~	Tilde
\	Backslash	<	Less than
(	Left parenthesis	>	Greater than
)	Right parenthesis	?	Question mark
[	Left square bracket	'	Apostrophe
]	Right square bracket	`	Grave accent
{	Left curly bracket	^	Circumflex accent
}	Right curly bracket		Vertical line
,	Comma	\$	Currency symbol
.	Decimal point or period	#	Number sign
:	Colon	@	Commercial at

1 2 Some of the special characters are used for operator symbols, bracketing, and various forms of separating and  
2 delimiting other lexical tokens.

### 3 6.1.6 Other characters

4 1 Additional characters may be representable in the processor, but may appear only in comments (6.3.2.3, 6.3.3.2),  
5 character constants (7.4.4), input/output records (12.2.2), and character string edit descriptors (13.3.2).

## 6 6.2 Low-level syntax

### 7 6.2.1 Tokens

8 1 The low-level syntax describes the fundamental lexical tokens of a [program unit](#). A [lexical token](#) is a keyword,  
9 name, literal constant other than a complex literal constant, operator, statement label, delimiter, comma, =, =>,  
10 :, ::, :, or %.

### 11 6.2.2 Names

12 1 [Names](#) are used for various entities such as [variables](#), [program units](#), [dummy arguments](#), [named constants](#), and  
13 derived types.

14 R603 *name* is *letter* [ *alphanumeric-character* ] ...

15 C601 (R603) The maximum length of a *name* is 63 characters.

#### NOTE 6.2

Examples of names:

```
A1
NAME_LENGTH      (single underscore)
S_P_R_E_A_D__O_U_T (two consecutive underscores)
TRAILER_         (trailing underscore)
```

**NOTE 6.3**

The word “name” always denotes this particular syntactic form. The word “identifier” is used where entities can be identified by other syntactic forms or by values; its particular meaning depends on the context in which it is used.

1 **6.2.3 Constants**

2	R604	<i>constant</i>	<b>is</b>	<i>literal-constant</i>
3			<b>or</b>	<i>named-constant</i>
4	R605	<i>literal-constant</i>	<b>is</b>	<i>int-literal-constant</i>
5			<b>or</b>	<i>real-literal-constant</i>
6			<b>or</b>	<i>complex-literal-constant</i>
7			<b>or</b>	<i>logical-literal-constant</i>
8			<b>or</b>	<i>char-literal-constant</i>
9			<b>or</b>	<i>boz-literal-constant</i>
10	R606	<i>named-constant</i>	<b>is</b>	<i>name</i>
11	R607	<i>int-constant</i>	<b>is</b>	<i>constant</i>
12	C602	(R607) <i>int-constant</i> shall be of type integer.		

13 **6.2.4 Operators**

14	R608	<i>intrinsic-operator</i>	<b>is</b>	<i>power-op</i>
15			<b>or</b>	<i>mult-op</i>
16			<b>or</b>	<i>add-op</i>
17			<b>or</b>	<i>concat-op</i>
18			<b>or</b>	<i>rel-op</i>
19			<b>or</b>	<i>not-op</i>
20			<b>or</b>	<i>and-op</i>
21			<b>or</b>	<i>or-op</i>
22			<b>or</b>	<i>equiv-op</i>
23	R1007	<i>power-op</i>	<b>is</b>	**
24	R1008	<i>mult-op</i>	<b>is</b>	*
25			<b>or</b>	/
26	R1009	<i>add-op</i>	<b>is</b>	+
27			<b>or</b>	-
28	R1011	<i>concat-op</i>	<b>is</b>	//
29	R1013	<i>rel-op</i>	<b>is</b>	.EQ.
30			<b>or</b>	.NE.
31			<b>or</b>	.LT.
32			<b>or</b>	.LE.
33			<b>or</b>	.GT.
34			<b>or</b>	.GE.
35			<b>or</b>	==
36			<b>or</b>	/=
37			<b>or</b>	<
38			<b>or</b>	<=
39			<b>or</b>	>
40			<b>or</b>	>=

1	R1018	<i>not-op</i>	is	.NOT.
2	R1019	<i>and-op</i>	is	.AND.
3	R1020	<i>or-op</i>	is	.OR.
4	R1021	<i>equiv-op</i>	is	.EQV.
5			or	.NEQV.
6	R609	<i>defined-operator</i>	is	<i>defined-unary-op</i>
7			or	<i>defined-binary-op</i>
8			or	<i>extended-intrinsic-op</i>
9	R1003	<i>defined-unary-op</i>	is	. letter [ letter ] ... .
10	R1023	<i>defined-binary-op</i>	is	. letter [ letter ] ... .
11	R610	<i>extended-intrinsic-op</i>	is	<i>intrinsic-operator</i>

## 6.2.5 Statement labels

13 1 A [statement label](#) provides a means of referring to an individual statement.

14 R611 *label* is *digit [ digit [ digit [ digit [ digit ] ] ] ]*

15 C603 (R611) At least one digit in a *label* shall be nonzero.

16 2 If a statement is labeled, the statement shall contain a nonblank character. The same statement label shall not  
 17 be given to more than one statement in its scope. Leading zeros are not significant in distinguishing between  
 18 statement labels. There are 99999 possible unique statement labels and a processor shall accept any of them as  
 19 a statement label. However, a processor may have a limit on the total number of unique statement labels in one  
 20 [program unit](#).

### NOTE 6.4

For example:

99999

10

010

are all statement labels. The last two are equivalent.

21 3 Any statement that is not part of another statement, and that is not preceded by a semicolon in fixed form, may begin  
 22 with a statement label, but the labels are used only in the following ways.

- 23 • The label on a [branch target statement](#) (11.2) is used to identify that statement as the possible destination  
 24 of a branch.
- 25 • The label on a [FORMAT statement](#) (13.2.1) is used to identify that statement as the format specification  
 26 for a [data transfer statement](#) (12.6).
- 27 • In some forms of the [DO construct](#) (11.1.7), the terminal statement of the construct is identified by a label.

## 6.2.6 Delimiters

29 1 A [lexical token](#) that is a delimiter is a (, ), /, [, ], (/, or /).



## 6.3 Source form

### 6.3.1 Program units, statements, and lines

- 1 A Fortran **program unit** is a sequence of one or more **lines**, organized as Fortran statements, comments, and **INCLUDE lines**. A **line** is a sequence of zero or more characters. **Lines** following a **program unit END statement** are not part of that **program unit**. A Fortran **statement** is a sequence of one or more complete or partial **lines**.
- 2 A comment may contain any character that may occur in any **character context**.
- 3 There are two source forms. Subclause 6.3.2 applies only to free form source. Subclause 6.3.3 applies only to fixed source form. Free form and fixed form shall not be mixed in the same program unit. The means for specifying the source form of a **program unit** are processor dependent.

### 6.3.2 Free source form

#### 6.3.2.1 Free form line length

- 1 In free source form there are no restrictions on where a statement (or portion of a statement) may appear within a **line**. A **line** may contain zero characters. If a **line** consists entirely of characters of default kind (7.4.4), it may contain at most 132 characters. If a **line** contains any character that is not of default kind, the maximum number of characters allowed on the **line** is processor dependent.

#### 6.3.2.2 Blank characters in free form

- 1 In free source form blank characters shall not appear within lexical tokens other than in a **character context** or in a format specification. Blanks may be inserted freely between tokens to improve readability; for example, blanks may occur between the tokens that form a complex literal constant. A sequence of blank characters outside of a **character context** is equivalent to a single blank character.
- 2 A blank shall be used to separate names, constants, or labels from adjacent keywords, names, constants, or labels.

#### NOTE 6.5

For example, the blanks after REAL, READ, 30, and DO are required in the following:

```
REAL X
READ 10
30 DO K=1,3
```

- 3 One or more blanks shall be used to separate adjacent keywords except in the following cases, where blanks are optional:

Table 6.2: Adjacent keywords where separating blanks are optional

BLOCK DATA	END FILE	END SUBROUTINE
DOUBLE PRECISION	END FORALL	END TEAM
ELSE IF	END FUNCTION	END TYPE
ELSE WHERE	END IF	END WHERE
END ASSOCIATE	END INTERFACE	GO TO
END BLOCK	END MODULE	IN OUT
END BLOCK DATA	END PROCEDURE	SELECT CASE
END CRITICAL	END PROGRAM	SELECT TYPE
END DO	END SELECT	
END ENUM	END SUBMODULE	

### 1 6.3.2.3 Free form commentary

2 1 The character “!” initiates a comment except where it appears within a [character context](#). The comment extends  
3 to the end of the [line](#). If the first nonblank character on a [line](#) is an “!”, the [line](#) is a comment line. [Lines](#)  
4 containing only blanks or containing no characters are also comment lines. Comments may appear anywhere in  
5 a [program unit](#) and may precede the first statement of a [program unit](#) or follow the last statement of a [program](#)  
6 [unit](#). Comments have no effect on the interpretation of the [program unit](#).

#### NOTE 6.6

This document does not restrict the number of consecutive comment lines.

### 7 6.3.2.4 Free form statement continuation

8 1 The character “&” is used to indicate that the statement is continued on the next [line](#) that is not a comment  
9 line. Comment lines cannot be continued; an “&” in a comment has no effect. Comments may occur within a  
10 continued statement. When used for continuation, the “&” is not part of the statement. No [line](#) shall contain  
11 a single “&” as the only nonblank character or as the only nonblank character before an “!” that initiates a  
12 comment.

13 2 If a non[character context](#) is to be continued, an “&” shall be the last nonblank character on the [line](#), or the last  
14 nonblank character before an “!”. There shall be a later [line](#) that is not a comment; the statement is continued  
15 on the next such [line](#). If the first nonblank character on that [line](#) is an “&”, the statement continues at the next  
16 character position following that “&”; otherwise, it continues with the first character position of that [line](#).

17 3 If a lexical token is split across the end of a [line](#), the first nonblank character on the first following noncomment  
18 [line](#) shall be an “&” immediately followed by the successive characters of the split token.

19 4 If a [character context](#) is to be continued, an “&” shall be the last nonblank character on the [line](#). There shall be  
20 a later [line](#) that is not a comment; an “&” shall be the first nonblank character on the next such [line](#) and the  
21 statement continues with the next character following that “&”.

### 22 6.3.2.5 Free form statement termination

23 1 If a statement is not continued, a comment or the end of the [line](#) terminates the statement.

24 2 A statement may alternatively be terminated by a “;” character that appears other than in a [character context](#)  
25 or in a comment. The “;” is not part of the statement. After a “;” terminator, another statement may appear  
26 on the same [line](#), or begin on that [line](#) and be continued. A sequence consisting only of zero or more blanks and  
27 one or more “;” terminators, in any order, is equivalent to a single “;” terminator.

### 28 6.3.2.6 Free form statements

29 1 A label may precede any statement not forming part of another statement.

#### NOTE 6.7

No Fortran statement begins with a digit.

30 2 A statement shall not have more than 255 continuation [lines](#).

## 31 6.3.3 Fixed source form

### 32 6.3.3.1 General

33 1 In fixed source form, there are restrictions on where a statement may appear within a [line](#). If a source [line](#) contains only characters  
34 of default kind, it shall contain exactly 72 characters; otherwise, its maximum number of characters is processor dependent.

35 2 Except in a [character context](#), blanks are insignificant and may be used freely throughout the program.

### 1 6.3.3.2 Fixed form commentary

2 1 The character “!” initiates a comment except where it appears within a [character context](#) or in character position 6. The comment  
3 extends to the end of the [line](#). If the first nonblank character on a [line](#) is an “!” in any character position other than character  
4 position 6, the [line](#) is a comment line. [Lines](#) beginning with a “C” or “\*” in character position 1 and lines containing only blanks are  
5 also comment lines. Comments may appear anywhere in a [program unit](#) and may precede the first statement of the [program unit](#) or  
6 follow the last statement of a [program unit](#). Comments have no effect on the interpretation of the [program unit](#).

#### NOTE 6.8

This document does not restrict the number of consecutive comment lines.

### 7 6.3.3.3 Fixed form statement continuation

8 1 Except within commentary, character position 6 is used to indicate continuation. If character position 6 contains a blank or zero, the  
9 [line](#) is the initial [line](#) of a new statement, which begins in character position 7. If character position 6 contains any character other  
10 than blank or zero, character positions 7–72 of the [line](#) constitute a continuation of the preceding noncomment [line](#).

#### NOTE 6.9

An “!” or “;” in character position 6 is interpreted as a continuation indicator unless it appears within commentary indicated  
by a “C” or “\*” in character position 1 or by an “!” in character positions 1–5.

11 2 Comment lines cannot be continued. Comment lines may occur within a continued statement.

### 12 6.3.3.4 Fixed form statement termination

13 1 If a statement is not continued, a comment or the end of the [line](#) terminates the statement.

14 2 A statement may alternatively be terminated by a “;” character that appears other than in a [character context](#), in a comment, or in  
15 character position 6. The “;” is not part of the statement. After a “;” terminator, another statement may begin on the same [line](#), or  
16 begin on that [line](#) and be continued. A “;” shall not appear as the first nonblank character on an initial [line](#). A sequence consisting  
17 only of zero or more blanks and one or more “;” terminators, in any order, is equivalent to a single “;” terminator.

### 18 6.3.3.5 Fixed form statements

19 1 A label, if it appears, shall occur in character positions 1 through 5 of the first [line](#) of a statement; otherwise, positions 1 through  
20 5 shall be blank. Blanks may appear anywhere within a label. A statement following a “;” on the same [line](#) shall not be labeled.  
21 Character positions 1 through 5 of any continuation [lines](#) shall be blank. A statement shall not have more than 255 continuation  
22 [lines](#). The [program unit END statement](#) shall not be continued. A statement whose initial [line](#) appears to be a [program unit END](#)  
23 [statement](#) shall not be continued.

## 24 6.4 Including source text

25 1 Additional text may be incorporated into the source text of a [program unit](#) during processing. This is accomplished  
26 with the INCLUDE line, which has the form

27 INCLUDE *char-literal-constant*

28 2 The *char-literal-constant* shall not have a kind type parameter value that is a *named-constant*.

29 3 An INCLUDE line is not a Fortran statement.

30 4 An INCLUDE line shall appear on a single source [line](#) where a statement may appear; it shall be the only  
31 nonblank text on this [line](#) other than an optional trailing comment. Thus, a statement label is not allowed.

32 5 The effect of the INCLUDE line is as if the referenced source text physically replaced the INCLUDE line prior  
33 to program processing. Included text may contain any source text, including additional INCLUDE lines; such  
34 nested INCLUDE lines are similarly replaced with the specified source text. The maximum depth of nesting of  
35 any nested INCLUDE lines is processor dependent. Inclusion of the source text referenced by an INCLUDE line  
36 shall not, at any level of nesting, result in inclusion of the same source text.

- 1 6 When an INCLUDE line is resolved, the first included statement *line* shall not be a continuation *line* and the last  
2 included statement *line* shall not be continued.
- 3 7 The interpretation of *char-literal-constant* is processor dependent. An example of a possible valid interpretation  
4 is that *char-literal-constant* is the name of a file that contains the source text to be included.

**NOTE 6.10**

In some circumstances, for example where source code is maintained in an INCLUDE file for use in programs whose source form might be either fixed or free, observing the following rules allows the code to be used with either source form.

- Confine statement labels to character positions 1 to 5 and statements to character positions 7 to 72.
- Treat blanks as being significant.
- Use only the exclamation mark (!) to indicate a comment, but do not start the comment in character position 6.
- For continued statements, place an ampersand (&) in both character position 73 of a continued *line* and character position 6 of a continuation *line*.

## 7 Types

### 7.1 Characteristics of types

#### 7.1.1 The concept of type

1 Fortran provides an abstract means whereby data can be categorized without relying on a particular physical representation. This abstract means is the concept of type.

2 A type has a name, a set of valid values, a means to denote such values (constants), and a set of operations to manipulate the values.

#### 7.1.2 Type classification

1 A type is either an [intrinsic type](#) or a [derived type](#).

2 This document defines five intrinsic types: integer, real, complex, character, and logical.

3 A derived type is one that is defined by a derived-type definition ([7.5.2](#)) or by an intrinsic module. It shall be used only where it is accessible ([7.5.2.2](#)). An intrinsic type is always accessible.

#### 7.1.3 Set of values

1 For each type, there is a set of valid values. The set of valid values for logical is completely determined by this document. The sets of valid values for integer, character, and real are processor dependent. The set of valid values for complex consists of the set of all the combinations of the values of the real and imaginary parts. The set of valid values for a derived type is as defined in [7.5.8](#).

#### 7.1.4 Constants

1 The syntax for denoting a value indicates the type, type parameters, and the particular value.

2 The syntax for literal constants of each intrinsic type is specified in [7.4](#).

3 A [structure constructor](#) ([7.5.10](#)) that is a [constant expression](#) ([10.1.12](#)) denotes a scalar constant value of derived type. An array constructor ([7.8](#)) that is a [constant expression](#) denotes a constant array value of intrinsic or derived type.

4 A constant value can be named ([8.5.13](#), [8.6.11](#)).

#### 7.1.5 Operations

1 For each of the intrinsic types, a set of operations and corresponding operators is defined intrinsically. These are described in [Clause 10](#). The intrinsic set can be augmented with operations and operators defined by functions with the [OPERATOR](#) interface ([15.4.3.2](#)). Operator definitions are described in [Clauses 10](#) and [15](#).

2 For derived types, there are no intrinsic operations. Operations on derived types can be defined by the program ([7.5.11](#)).

## 7.2 Type parameters

- 1 A type might be parameterized. In this case, the set of values, the syntax for denoting the values, and the set of operations on the values of the type depend on the values of the parameters.
- 2 A [type parameter](#) is either a [kind type parameter](#) or a [length type parameter](#). All type parameters are of type integer. A [kind type parameter](#) participates in generic resolution (15.5.5.2), but a [length type parameter](#) does not.
- 3 Each intrinsic type has a [kind type parameter](#) named KIND. The intrinsic character type has a [length type parameter](#) named LEN. A derived type can have type parameters.
- 4 A type parameter value may be specified by a type specification (7.4, 7.5.9).
- R701 *type-param-value*                    **is** *scalar-int-expr*  
    **or** \*  
    **or** :
- C701 (R701) The *type-param-value* for a kind type parameter shall be a [constant expression](#).
- C702 (R701) A colon shall not be used as a *type-param-value* except in the declaration of an entity that has the [POINTER](#) or [ALLOCATABLE](#) attribute.
- 5 A colon as a *type-param-value* specifies a [deferred type parameter](#).
- 6 The values of the [deferred type parameters](#) of an object are determined by successful execution of an [ALLOCATE statement](#) (9.7.1), execution of an [intrinsic assignment statement](#) (10.2.1.3), execution of a [pointer assignment statement](#) (10.2.2), or by [argument association](#) (15.5.2).

### NOTE 7.1

[Deferred type parameters](#) of functions, including function procedure pointers, have no values. Instead, they indicate that those [type parameters](#) of the function result will be determined by execution of the function, if it returns an allocated [allocatable](#) result or an associated pointer result.

- 7 An asterisk as a *type-param-value* specifies that a [length type parameter](#) is an [assumed type parameter](#). It is used for a [dummy argument](#) to assume the type parameter value from the [effective argument](#), for an [associate name](#) in a [SELECT TYPE construct](#) to assume the type parameter value from the corresponding selector, and for a [named constant](#) of type character to assume the character length from the *constant-expr*.

### NOTE 7.2

The value of a [kind type parameter](#) is always known at compile time. Some parameterizations that involve multiple representation forms need to be distinguished at compile time for practical implementation and performance. Examples include the multiple precisions of the intrinsic real type and the possible multiple character sets of the intrinsic character type.

The adjective “length” is used for type parameters other than kind type parameters because they often specify a length, as for intrinsic character type. However, they can be used for other purposes. The important difference from kind type parameters is that their values need not be known at compile time and might change during execution.

## 7.3 Types, type specifiers, and values

### 7.3.1 Relationship of types and values to objects

- 1 The name of a type serves as a type specifier and may be used to declare objects of that type. A declaration specifies the type of a named object. A data object may be declared explicitly or implicitly. A data object has

1 **attributes** in addition to its type. Clause 8 describes the way in which a data object is declared and how its type  
2 and other attributes are specified.

3 2 Scalar data of any intrinsic or derived type may be shaped in a rectangular pattern to compose an array of the  
4 same type and type parameters. An array object has a type and type parameters just as a scalar object does.

5 3 A variable is a data object. The type and type parameters of a variable determine which values that variable  
6 may take. Assignment (10.2) provides one means of changing the value of a variable of any type.

7 4 The type of a variable determines the operations that may be used to manipulate the variable.

## 8 **7.3.2 Type specifiers and type compatibility**

### 9 **7.3.2.1 Type specifier syntax**

10 1 A type specifier specifies a type and type parameter values. It is either a *type-spec* or a *declaration-type-spec*.

11 R702 *type-spec* is *intrinsic-type-spec*  
12 or *derived-type-spec*

13 C703 (R702) The *derived-type-spec* shall not specify an **abstract type** (7.5.7).

14 R703 *declaration-type-spec* is *intrinsic-type-spec*  
15 or TYPE ( *intrinsic-type-spec* )  
16 or TYPE ( *derived-type-spec* )  
17 or CLASS ( *derived-type-spec* )  
18 or CLASS ( \* )  
19 or TYPE ( \* )

20 C704 (R703) In a *declaration-type-spec*, every *type-param-value* that is not a colon or an asterisk shall be a  
21 *specification-expr*.

22 C705 (R703) In a *declaration-type-spec* that uses the **CLASS** keyword, *derived-type-spec* shall specify an **ex-**  
23 **ensible type** (7.5.7).

24 C706 (R703) TYPE(*derived-type-spec*) shall not specify an **abstract type** (7.5.7).

25 C707 (R702) In TYPE(*intrinsic-type-spec*) the *intrinsic-type-spec* shall not end with a comma.

26 C708 An entity declared with the **CLASS** keyword shall be a **dummy argument** or have the **ALLOCATABLE**  
27 or **POINTER** attribute.

28 2 An *intrinsic-type-spec* specifies the named intrinsic type and its type parameter values. A *derived-type-spec*  
29 specifies the named derived type and its type parameter values.

#### **NOTE 7.3**

A *type-spec* is used in an array constructor, a **SELECT TYPE** construct, or an **ALLOCATE** statement. An *integer-type-spec* is used in a **DO CONCURRENT** or **FORALL** statement. Elsewhere, a *declaration-type-spec* is used.

### 30 **7.3.2.2 TYPE type specifier**

31 1 A **TYPE** type specifier is used to declare entities that are **assumed-type**, or of an intrinsic or derived type.

32 2 A *derived-type-spec* in a **TYPE** type specifier in a **type declaration statement** shall specify a previously defined  
33 derived type. If the data entity is a function result, the derived type may be specified in the **FUNCTION**  
34 **statement** provided the derived type is defined within the body of the function or is accessible there by use or  
35 **host** association. If the derived type is specified in the **FUNCTION statement** and is defined within the body  
36 of the function, it is as if the **function result** were declared with that derived type immediately following the

1 *derived-type-def* of the specified derived type.

2 3 An entity that is declared using the TYPE(\*) type specifier is **assumed-type** and is an **unlimited polymorphic**  
 3 entity. It is not declared to have a type, and is not considered to have the same **declared type** as any other entity,  
 4 including another **unlimited polymorphic** entity. Its **dynamic type** and **type parameters** are assumed from its  
 5 **effective argument**.

6 C709 An **assumed-type** entity shall be a **dummy data object** that does not have the **ALLOCATABLE**, **CODI-**  
 7 **MENSION**, **INTENT (OUT)**, **POINTER**, or **VALUE** attribute and is not an **explicit-shape array**.

8 C710 An **assumed-type** variable name shall not appear in a designator or expression except as an **actual**  
 9 **argument** corresponding to a **dummy argument** that is **assumed-type**, or as the first argument to the  
 10 **intrinsic** function **IS\_CONTIGUOUS**, **LBOUND**, **PRESENT**, **RANK**, **SHAPE**, **SIZE**, or **UBOUND**, or  
 11 the function **C\_LOC** from the intrinsic module **ISO\_C\_BINDING**.

12 C711 An **assumed-type actual argument** that corresponds to an **assumed-rank dummy argument** shall be  
 13 **assumed-shape** or **assumed-rank**.

### 14 7.3.2.3 CLASS type specifier

15 1 The **CLASS** type specifier is used to declare **polymorphic** entities. A **polymorphic** entity is a data entity that is  
 16 able to be of differing **dynamic types** during program execution.

17 2 A *derived-type-spec* in a CLASS type specifier in a **type declaration statement** shall specify a previously defined  
 18 derived type. If the data entity is a function result, the derived type may be specified in the **FUNCTION**  
 19 **statement** provided the derived type is defined within the body of the function or is accessible there by **use** or  
 20 **host** association. If the derived type is specified in the **FUNCTION statement** and is defined within the body  
 21 of the function, it is as if the function result were declared with that derived type immediately following its  
 22 *derived-type-def*.

23 3 The **declared type** of a **polymorphic** entity is the specified type if the **CLASS** type specifier contains a type name.

24 4 An entity declared with the **CLASS(\*)** specifier is an **unlimited polymorphic** entity. It is not declared to have  
 25 a type, and is not considered to have the same **declared type** as any other entity, including another **unlimited**  
 26 **polymorphic** entity.

27 5 A nonpolymorphic entity is **type compatible** only with entities of the same **declared type**. A **polymorphic** entity  
 28 that is not an **unlimited polymorphic** entity is **type compatible** with entities of the same **declared type** or any of  
 29 its **extensions**. Even though an **unlimited polymorphic** entity is not considered to have a **declared type**, it is **type**  
 30 **compatible** with all entities. An entity is **type compatible** with a type if it is **type compatible** with entities of that  
 31 type.

#### NOTE 7.4

Given

```
TYPE TROOT
...
TYPE,EXTENDS(TROOT) :: TEXTENDED
...
CLASS(TROOT) A
CLASS(TEXTENDED) B
...
```

A is **type compatible** with B but B is not **type compatible** with A.

32 6 A **polymorphic allocatable** object may be allocated to be of any type with which it is **type compatible**. A  
 33 **polymorphic** pointer or **dummy argument** may, during program execution, be associated with objects with which  
 34 it is **type compatible**.



1 7 The **dynamic type** of an allocated **allocatable polymorphic** object is the type with which it was allocated. The  
 2 **dynamic type** of an associated **polymorphic** pointer is the **dynamic type** of its **target**. The **dynamic type** of a  
 3 nonallocatable nonpointer **polymorphic dummy argument** is the **dynamic type** of its **effective argument**. The  
 4 **dynamic type** of an unallocated **allocatable** object or a **disassociated** pointer is the same as its **declared type**. The  
 5 **dynamic type** of an entity identified by an **associate name** (11.1.3) is the **dynamic type** of the selector with which  
 6 it is associated. The **dynamic type** of an object that is not **polymorphic** is its **declared type**.

## 7 7.4 Intrinsic types

### 8 7.4.1 Classification and specification

9 1 Each intrinsic type is classified as a **numeric type** or a nonnumeric type. The **numeric types** are integer, real, and  
 10 complex. The nonnumeric intrinsic types are character and logical.

11 2 Each intrinsic type has a **kind type parameter** named KIND; this **type parameter** is of type integer with default  
 12 kind.

13 R704 *intrinsic-type-spec* is *integer-type-spec*  
 14 or REAL [ *kind-selector* ]  
 15 or DOUBLE PRECISION  
 16 or COMPLEX [ *kind-selector* ]  
 17 or CHARACTER [ *char-selector* ]  
 18 or LOGICAL [ *kind-selector* ]

19 R705 *integer-type-spec* is INTEGER [ *kind-selector* ]

20 R706 *kind-selector* is ( [ KIND = ] *scalar-int-constant-expr* )

21 C712 (R706) The value of *scalar-int-constant-expr* shall be nonnegative and shall specify a representation  
 22 method that exists on the processor.

### 23 7.4.2 Intrinsic operations on intrinsic types

24 1 Intrinsic numeric operations are defined as specified in 10.1.5.2.1 for the numeric intrinsic types. Relational  
 25 intrinsic operations are defined as specified in 10.1.5.5 for numeric and character intrinsic types. The intrinsic  
 26 concatenation operation is defined as specified in 10.1.5.3 for the character type. Logical intrinsic operations are  
 27 defined as specified in 10.1.5.4 for the logical type.

### 28 7.4.3 Numeric intrinsic types

#### 29 7.4.3.1 Integer type

30 1 The set of values for the integer type is a subset of the mathematical integers. The processor shall provide one or  
 31 more representation methods that define sets of values for data of type integer. Each such method is characterized  
 32 by a value for the **kind type parameter** KIND. The **kind type parameter** of a representation method is returned  
 33 by the intrinsic function **KIND** (16.9.108). The decimal exponent range of a representation method is returned  
 34 by the intrinsic function **RANGE** (16.9.158). The intrinsic function **SELECTED\_INT\_KIND** (16.9.169) returns  
 35 a kind value based on a specified decimal exponent range requirement. The integer type includes a zero value,  
 36 which is considered to be neither negative nor positive. The value of a signed integer zero is the same as the  
 37 value of an unsigned integer zero.

38 2 The processor shall provide at least one representation method with a decimal exponent range greater than or  
 39 equal to 18.

40 3 The type specifier for the integer type uses the keyword **INTEGER**.

- 1 4 The keyword INTEGER with no *kind-selector* specifies type integer with default kind; the *kind type parameter*  
 2 value is equal to `KIND (0)`. The decimal exponent range of default integer shall be at least 5.
- 3 5 Any integer value may be represented as a *signed-int-literal-constant*.
- 4 R707 *signed-int-literal-constant* is [ *sign* ] *int-literal-constant*
- 5 R708 *int-literal-constant* is *digit-string* [ *\_ kind-param* ]
- 6 R709 *kind-param* is *digit-string*  
 7 or *scalar-int-constant-name*
- 8 R710 *signed-digit-string* is [ *sign* ] *digit-string*
- 9 R711 *digit-string* is *digit* [ *digit* ] ...
- 10 R712 *sign* is +  
 11 or -
- 12 C713 (R709) A *scalar-int-constant-name* shall be a *named constant* of type integer.
- 13 C714 (R709) The value of *kind-param* shall be nonnegative.
- 14 C715 (R708) The value of *kind-param* shall specify a representation method that exists on the processor.
- 15 6 The optional *kind type parameter* following *digit-string* specifies the *kind type parameter* of the integer constant;  
 16 if it does not appear, the constant is default integer.
- 17 7 An integer constant is interpreted as a decimal value.

**NOTE 7.5**

Examples of signed integer literal constants are:

```
473
+56
-101
21_2
21_SHORT
1976354279568241_8
```

where SHORT is a scalar integer *named constant*. A program that uses a *digit-string* as a *kind-param* is unlikely to be portable.

18 **7.4.3.2 Real type**

- 19 1 The real type has values that approximate the mathematical real numbers. The processor shall provide two  
 20 or more approximation methods that define sets of values for data of type real. Each such method has a  
 21 representation method and is characterized by a value for the *kind type parameter* `KIND`. The *kind type parameter*  
 22 of an approximation method is returned by the intrinsic function `KIND (16.9.108)`.
- 23 2 The decimal precision, decimal exponent range, and radix of an approximation method are returned by the  
 24 intrinsic functions `PRECISION (16.9.151)`, `RANGE (16.9.158)`, and `RADIX (16.9.154)`. The intrinsic function  
 25 `SELECTED_REAL_KIND (16.9.170)` returns a kind value based on specified precision, range, and radix re-  
 26 quirements.

**NOTE 7.6**

See C.2.1 for remarks concerning selection of approximation methods.

1 3 The real type includes a zero value. Processors that distinguish between positive and negative zeros shall treat  
2 them as mathematically equivalent

- 3 • in all intrinsic relational operations, and
- 4 • as [actual arguments](#) to intrinsic procedures other than those for which it is explicitly specified that negative  
5 zero is distinguished.

#### NOTE 7.7

On a processor that distinguishes between 0.0 and  $-0.0$ ,

( X >= 0.0 )

evaluates to true if X = 0.0 or if X =  $-0.0$ ,

( X < 0.0 )

evaluates to false for X =  $-0.0$ .

In order to distinguish between 0.0 and  $-0.0$ , a program can use the intrinsic function [SIGN](#). [SIGN](#) (1.0, X) will return  $-1.0$  if X < 0.0 or if the processor distinguishes between 0.0 and  $-0.0$  and X has the value  $-0.0$ .

6 4 The type specifier for the real type uses the keyword REAL. The keyword DOUBLE PRECISION is an alternative  
7 specifier for one kind of real type.

8 5 If the type keyword REAL is used without a [kind type parameter](#), the real type with default real kind is specified  
9 and the kind value is [KIND](#) (0.0). The type specifier DOUBLE PRECISION specifies type real with double  
10 precision kind; the kind value is [KIND](#) (0.0D0). The decimal precision of the double precision real approximation  
11 method shall be greater than that of the default real method.

12 6 The decimal precision of double precision real shall be at least 10, and its decimal exponent range shall be at  
13 least 37. It is recommended that the decimal precision of default real be at least 6, and that its decimal exponent  
14 range be at least 37.

15 R713 *signed-real-literal-constant* is [ *sign* ] *real-literal-constant*

16 R714 *real-literal-constant* is *significand* [ *exponent-letter exponent* ] [ *kind-param* ]  
17 or *digit-string exponent-letter exponent* [ *kind-param* ]

18 R715 *significand* is *digit-string* . [ *digit-string* ]  
19 or . *digit-string*

20 R716 *exponent-letter* is E  
21 or D

22 R717 *exponent* is *signed-digit-string*

23 C716 (R714) If both *kind-param* and *exponent-letter* appear, *exponent-letter* shall be E.

24 C717 (R714) The value of *kind-param* shall specify an approximation method that exists on the processor.

25 7 A real literal constant without a [kind type parameter](#) is a default real constant if it is without an exponent part  
26 or has exponent letter E, and is a double precision real constant if it has exponent letter D. A real literal constant  
27 written with a [kind type parameter](#) is a real constant with the specified [kind type parameter](#).

28 8 The exponent represents the power of ten scaling to be applied to the significand or digit string. The meaning of  
29 these constants is as in decimal scientific notation.

30 9 The significand may be written with more digits than a processor will use to approximate the value of the constant.

1

**NOTE 7.8**

Examples of signed real literal constants are:

```
-12.78
+1.6E3
2.1
-16.E4_8
0.45D-4
10.93E7_QUAD
.123
3E4
```

where QUAD is a scalar integer [named constant](#).

2 **7.4.3.3 Complex type**

3 1 The complex type has values that approximate the mathematical complex numbers. The values of a complex  
4 type are ordered pairs of real values. The first real value is called the real part, and the second real value is called  
5 the imaginary part.

6 2 Each approximation method used to represent data entities of type real shall be available for both the real and  
7 imaginary parts of a data entity of type complex. The (default integer) [kind type parameter](#) KIND for a complex  
8 entity specifies for both parts the real approximation method characterized by this kind type parameter value.  
9 The [kind type parameter](#) of an approximation method is returned by the intrinsic function [KIND](#) (16.9.108).

10 3 The type specifier for the complex type uses the keyword COMPLEX. There is no keyword for double precision  
11 complex. If the type keyword COMPLEX is used without a [kind type parameter](#), the complex type with default  
12 complex kind is specified, the kind value is [KIND](#) (0.0), and both parts are default real.

13 R718 *complex-literal-constant* is ( *real-part* , *imag-part* )

14 R719 *real-part* is [signed-int-literal-constant](#)  
15 or [signed-real-literal-constant](#)  
16 or [named-constant](#)

17 R720 *imag-part* is [signed-int-literal-constant](#)  
18 or [signed-real-literal-constant](#)  
19 or [named-constant](#)

20 C718 (R718) Each [named constant](#) in a complex literal constant shall be of type integer or real.

21 4 If the real part and the imaginary part of a complex literal constant are both real, the [kind type parameter](#) value  
22 of the complex literal constant is the [kind type parameter](#) value of the part with the greater decimal precision; if  
23 the precisions are the same, it is the [kind type parameter](#) value of one of the parts as determined by the processor.  
24 If a part has a [kind type parameter](#) value different from that of the complex literal constant, the part is converted  
25 to the approximation method of the complex literal constant.

26 5 If both the real and imaginary parts are integer, they are converted to the default real approximation method  
27 and the constant is default complex. If only one of the parts is an integer, it is converted to the approximation  
28 method selected for the part that is real and the [kind type parameter](#) value of the complex literal constant is  
29 that of the part that is real.

**NOTE 7.9**

Examples of complex literal constants are:

## NOTE 7.9 (cont.)

```
(1.0, -1.0)
(3, 3.1E6)
(4.0_4, 3.6E7_8)
( 0., PI)      ! where PI is a previously declared named real constant.
```

1 **7.4.4 Character type**2 **7.4.4.1 Character sets**

- 3 1 The character type has a set of values composed of character strings. A character string is a sequence of characters,  
4 numbered from left to right 1, 2, 3, ... up to the number of characters in the string. The number of characters in  
5 the string is called the length of the string. The length is a type parameter; its kind is processor dependent and  
6 its value is greater than or equal to zero.
- 7 2 The processor shall provide one or more representation methods that define sets of values for data of type  
8 character. Each such method is characterized by a value for the (default integer) **kind type parameter** KIND.  
9 The **kind type parameter** of a representation method is returned by the intrinsic function **KIND** (16.9.108). The  
10 intrinsic function **SELECTED\_CHAR\_KIND** (16.9.168) returns a kind value based on the name of a character  
11 type. Any character of a particular representation method representable in the processor may occur in a character  
12 string of that representation method.
- 13 3 The character set specified in ISO/IEC 646:1991 (International Reference Version) is referred to as the **ASCII**  
14 **character** set and its corresponding representation method is **ASCII character** kind. The character set UCS-4 as  
15 specified in ISO/IEC 10646 is referred to as the **ISO 10646 character** set and its corresponding representation  
16 method is the **ISO 10646 character** kind.

17 **7.4.4.2 Character type specifier**

- 18 1 The type specifier for the character type uses the keyword CHARACTER.
- 19 2 If the type keyword CHARACTER is used without a **kind type parameter**, the character type with default  
20 character kind is specified and the kind value is **KIND** ('A').
- 21 3 The default character kind shall support a character set that includes the characters in the Fortran character  
22 set (6.1). By supplying nondefault character kinds, the processor may support additional character sets. The  
23 characters available in nondefault character kinds are not specified by this document, except that one character  
24 in each nondefault character set shall be designated as a blank character to be used as a padding character.

25 R721 *char-selector*                    **is** *length-selector*  
26    **or** ( LEN = *type-param-value* , ■  
27    ■ KIND = *scalar-int-constant-expr* )  
28    **or** ( *type-param-value* , ■  
29    ■ [ KIND = ] *scalar-int-constant-expr* )  
30    **or** ( KIND = *scalar-int-constant-expr* ■  
31    ■ [ , LEN = *type-param-value* ] )

32 R722 *length-selector*                **is** ( [ LEN = ] *type-param-value* )  
33    **or** \* *char-length* [ , ]

34 R723 *char-length*                    **is** ( *type-param-value* )  
35    **or** *int-literal-constant*

36 C719 (R721) The value of *scalar-int-constant-expr* shall be nonnegative and shall specify a representation  
37 method that exists on the processor.

38 C720 (R723) The *int-literal-constant* shall not include a *kind-param*.

- 1 C721 (R723) A *type-param-value* in a *char-length* shall be a colon, asterisk, or *specification-expr*.
- 2 C722 (R721 R722 R723) A *type-param-value* of \* shall be used only
- 3       • to declare a *dummy argument*,
- 4       • to declare a *named constant*,
- 5       • in the *type-spec* of an **ALLOCATE** statement wherein each *allocate-object* is a *dummy argument* of
- 6       type CHARACTER with an assumed character length,
- 7       • in the *type-spec* or *derived-type-spec* of a *type guard statement* (11.1.11), or
- 8       • in an external function, to declare the character length parameter of the function result.
- 9 C723 A function name shall not be declared with an asterisk *type-param-value* unless it is of type CHARACTER
- 10 and is the name of a *dummy function* or the name of the result of an external function.
- 11 C724 A function name declared with an asterisk *type-param-value* shall not be an array, a pointer, *elemental*, recursive, or pure.
- 12 C725 (R722) The optional comma in a *length-selector* is permitted only in a *declaration-type-spec* in a *type-declaration-stmt*.
- 13 C726 (R722) The optional comma in a *length-selector* is permitted only if no double-colon separator appears in the *type-*
- 14 *declaration-stmt*.
- 15 C727 (R721) The length specified for a character *statement function* or for a *statement function dummy argument* of type
- 16 character shall be a *constant expression*.
- 17 4 The *char-selector* in a CHARACTER *intrinsic-type-spec* and the \* *char-length* in an *entity-decl* or in a *component-*
- 18 *decl* of a type definition specify character length. The \* *char-length* in an *entity-decl* or a *component-decl* specifies
- 19 an individual length and overrides the length specified in the *char-selector*, if any. If a \* *char-length* is not specified
- 20 in an *entity-decl* or a *component-decl*, the *length-selector* or *type-param-value* specified in the *char-selector* is the
- 21 character length. If the length is not specified in a *char-selector* or a \* *char-length*, the length is 1.
- 22 5 If the character length parameter value evaluates to a negative value, the length of character entities declared
- 23 is zero. A character length parameter value of : indicates a *deferred type parameter* (7.2). A *char-length* type
- 24 parameter value of \* has the following meanings.
- 25       • If used to declare a *dummy argument* of a procedure, the *dummy argument* assumes its length from its
- 26 *effective argument*.
- 27       • If used to declare a *named constant*, the length is that of the constant value.
- 28       • If used in the *type-spec* of an **ALLOCATE** statement, each *allocate-object* assumes its length from its *effective*
- 29 *argument*.
- 30       • If used in the *type-spec* of a *type guard statement*, the *associating entity* assumes its length from the selector.
- 31       • If used to specify the character length parameter of a function result, any *scoping unit* invoking the function or passing it as
- 32 an actual argument shall declare the function name with a character length parameter value other than \* or access such a
- 33 definition by *argument*, *host*, or *use* association. When the function is invoked, the length of the *function result* is assumed
- 34 from the value of this type parameter.

#### 35 7.4.4.3 Character literal constant

- 36 1 The syntax of a character literal constant is given by R724.
- 37 R724 *char-literal-constant*            **is** [ *kind-param* \_ ] ' [ *rep-char* ] ... '
- 38   **or** [ *kind-param* \_ ] " [ *rep-char* ] ... "
- 39 C728 (R724) The value of *kind-param* shall specify a representation method that exists on the processor.
- 40 2 The optional *kind type parameter* preceding the leading delimiter specifies the *kind type parameter* of the char-
- 41 acter constant; if it does not appear, the constant is default character.

- 1 3 For the type character with kind *kind-param*, if it appears, and for default character otherwise, a representable  
2 character, *rep-char*, is defined as follows.
- 3 • In free source form, it is any graphic character in the processor-dependent character set.
  - 4 • In fixed source form, it is any character in the processor-dependent character set. A processor may restrict the occurrence of  
5 some or all of the control characters.
- 6 4 The delimiting apostrophes or quotation marks are not part of the value of the character literal constant.
- 7 5 An apostrophe character within a character constant delimited by apostrophes is represented by two consecutive  
8 apostrophes (without intervening blanks); in this case, the two apostrophes are counted as one character. Sim-  
9 ilarly, a quotation mark character within a character constant delimited by quotation marks is represented by  
10 two consecutive quotation marks (without intervening blanks) and the two quotation marks are counted as one  
11 character.
- 12 6 A zero-length character literal constant is represented by two consecutive apostrophes (without intervening blanks)  
13 or two consecutive quotation marks (without intervening blanks) outside of a *character context*.

**NOTE 7.10**

Examples of character literal constants are:

```
"DON'T"  
'DON' 'T'
```

both of which have the value DON'T and

```
''
```

which has the zero-length character string as its value.

**NOTE 7.11**

An example of a nondefault character literal constant, where the processor supports the corresponding character set, is:

```
NIHONGO_ '彼女なしでは何もできない。'
```

where NIHONGO is a *named constant* whose value is the kind type parameter for Nihongo (Japanese) characters. This means “Without her, nothing is possible”.

14 **7.4.4.4 Collating sequence**

- 15 1 The processor defines a *collating sequence* for the character set of each kind of character. The *collating sequence*  
16 is an isomorphism between the character set and the set of integers  $\{I : 0 \leq I < N\}$ , where  $N$  is the number of  
17 characters in the set. The intrinsic functions *CHAR* (16.9.44) and *ICHAR* (16.9.95) provide conversions between  
18 the characters and the integers according to this mapping.

**NOTE 7.12**

For example:

```
ICHAR ( 'X' )
```

returns the integer value of the character 'X' according to the *collating sequence* of the processor.

- 19 2 The *collating sequence* of the default character kind shall satisfy the following constraints.
- 20 • *ICHAR* ('A') < *ICHAR* ('B') < ... < *ICHAR* ('Z') for the twenty-six upper-case letters.





- 1 4 The **direct components** of a derived type are the components of that type, plus the **direct components** of the  
2 components that are of derived type and have neither the **ALLOCATABLE** nor **POINTER** attribute.
- 3 5 The **potential subobject components** of a derived type are the nonpointer components of that type together with  
4 the **potential subobject components** of the nonpointer components that are of derived type. This includes all the  
5 components that could be a subobject of an object of the type (9.4.2).
- 6 6 The **components**, **direct components**, **potential subobject components**, and **ultimate components** of an object of  
7 derived type are the **components**, **direct components**, **potential subobject components**, and **ultimate components**  
8 of its type, respectively.
- 9 7 By default, no **storage sequence** is implied by the order of the component definitions. However, a storage sequence  
10 is implied for a **sequence type** (7.5.2.3). If the derived type has the **BIND** attribute, the **storage sequence** is that  
11 required by the **companion processor** (5.5.7, 18.3.4).
- 12 8 A scalar entity of derived type is a **structure**. If a derived type has the **SEQUENCE** attribute, a scalar entity of  
13 the type is a **sequence structure**.

**NOTE 7.14**

The **ultimate components** of an object of the derived type `kids` defined below are `oldest_child%name`, `oldest_child%age`, and `other_kids`. The **direct components** of such an object are `oldest_child%name`, `oldest_child%age`, `other_kids`, and `oldest_child`.

```

type :: person
  character(len=20) :: name
  integer :: age
end type person

type :: kids
  type(person) :: oldest_child
  type(person), allocatable, dimension(:) :: other_kids
end type kids

```

14 **7.5.2 Derived-type definition**15 **7.5.2.1 Syntax of a derived-type definition**

- 16 R726 *derived-type-def* is *derived-type-stmt*  
17 [ *type-param-def-stmt* ] ...  
18 [ *private-or-sequence* ] ...  
19 [ *component-part* ]  
20 [ *type-bound-procedure-part* ]  
21 *end-type-stmt*
- 22 R727 *derived-type-stmt* is TYPE [ [ , *type-attr-spec-list* ] :: ] *type-name* ■  
23 ■ [ ( *type-param-name-list* ) ]
- 24 R728 *type-attr-spec* is **ABSTRACT**  
25 **or** *access-spec*  
26 **or** BIND (C)  
27 **or** EXTENDS ( *parent-type-name* )
- 28 C730 (R727) A derived type *type-name* shall not be **DOUBLEPRECISION** or the same as the name of any  
29 intrinsic type defined in this document.
- 30 C731 (R727) The same *type-attr-spec* shall not appear more than once in a given *derived-type-stmt*.
- 31 C732 The same *type-param-name* shall not appear more than once in a given *derived-type-stmt*.



**NOTE 7.16 (cont.)**

Such a type would be accessible only within the module in which it is defined, and within its [descendants](#).

1 **7.5.2.3 Sequence type**2 R731 *sequence-stmt* is SEQUENCE3 C741 (R726) If SEQUENCE appears, the type shall have at least one component, each data component shall  
4 be declared to be of an intrinsic type or of a [sequence type](#), the derived type shall not have any [type](#)  
5 [parameter](#), and a [type-bound-procedure-part](#) shall not appear.6 1 If the SEQUENCE statement appears, the type has the [SEQUENCE attribute](#) and is a [sequence type](#). The order  
7 of the component definitions in a [sequence type](#) specifies a [storage sequence](#) for objects of that type. The type  
8 is a [numeric sequence type](#) if there are no [pointer](#) or [allocatable components](#), and each [component](#) is default  
9 integer, default real, double precision real, default complex, default logical, or of [numeric sequence type](#). The  
10 type is a [character sequence type](#) if there are no [pointer](#) or [allocatable components](#), and each component is default  
11 character or of [character sequence type](#).**NOTE 7.17**

An example of a [numeric sequence type](#) is:

```
TYPE NUMERIC_SEQ
  SEQUENCE
  INTEGER :: INT_VAL
  REAL    :: REAL_VAL
  LOGICAL :: LOG_VAL
END TYPE NUMERIC_SEQ
```

**NOTE 7.18**

A structure resolves into a sequence of components. Unless the structure includes a SEQUENCE statement, the use of this terminology in no way implies that these components are stored in this, or any other, order. Nor is there any requirement that [contiguous](#) storage be used. The sequence merely refers to the fact that in writing the definitions there will necessarily be an order in which the components appear, and this will define a sequence of components. This order is of limited significance because a component of an object of derived type will always be accessed by a component name except in the following contexts: the sequence of expressions in a derived-type value constructor, intrinsic assignment, the data values in namelist input data, and the inclusion of the structure in an input/output list of a formatted data transfer, where it is expanded to this sequence of components. Provided the processor adheres to the defined order in these cases, it is otherwise free to organize the storage of the components for any nonsequence structure in memory as best suited to the particular architecture.

12 **7.5.2.4 Determination of derived types**13 1 Derived-type definitions with the same type name may appear in different [scoping units](#), in which case they might  
14 be independent and describe different derived types or they might describe the same type.15 2 Two data entities have the same type if they are declared with reference to the same derived-type definition. Data  
16 entities also have the same type if they are declared with reference to different derived-type definitions that specify  
17 the same type name, all have the [SEQUENCE attribute](#) or all have the [BIND attribute](#), have no components  
18 with [PRIVATE](#) accessibility, and have components that agree in order, name, and attributes. Otherwise, they  
19 are of different derived types. A data entity declared using a type with the [SEQUENCE attribute](#) or with the  
20 [BIND attribute](#) is not of the same type as an entity of a type that has any components that are [PRIVATE](#).

**NOTE 7.19**

An example of declaring two entities with reference to the same derived-type definition is:

```

TYPE POINT
  REAL X, Y
END TYPE POINT
TYPE (POINT) :: X1
CALL SUB (X1)
...
CONTAINS
  SUBROUTINE SUB (A)
    TYPE (POINT) :: A
    ...
  END SUBROUTINE SUB

```

The definition of derived type POINT is known in subroutine SUB by host association. Because the declarations of X1 and A both reference the same derived-type definition, X1 and A have the same type. X1 and A also would have the same type if the derived-type definition were in a module and both SUB and its containing [program unit](#) accessed that derived type from the module.

**NOTE 7.20**

An example of data entities in different [scoping units](#) having the same type is:

```

PROGRAM PGM
  TYPE EMPLOYEE
    SEQUENCE
    INTEGER ID_NUMBER
    CHARACTER (50) NAME
  END TYPE EMPLOYEE
  TYPE (EMPLOYEE) PROGRAMMER
  CALL SUB (PROGRAMMER)
  ...
END PROGRAM PGM
SUBROUTINE SUB (POSITION)
  TYPE EMPLOYEE
    SEQUENCE
    INTEGER ID_NUMBER
    CHARACTER (50) NAME
  END TYPE EMPLOYEE
  TYPE (EMPLOYEE) POSITION
  ...
END SUBROUTINE SUB

```

The [actual argument](#) PROGRAMMER and the [dummy argument](#) POSITION have the same type because they are declared with reference to a derived-type definition with the same name, the [SEQUENCE attribute](#), and components that agree in order, name, and [attributes](#).

Suppose the component name ID\_NUMBER was ID\_NUM in the subroutine. Because all the component names are not identical to the component names in derived type EMPLOYEE in the main program, the [actual argument](#) PROGRAMMER would not be of the same type as the [dummy argument](#) POSITION. Thus, the program would not be standard-conforming.

**NOTE 7.21**

The requirement that the two types have the same name applies to the *type-names* in the respective derived

## NOTE 7.21 (cont.)

type definitions, not to local names introduced via renaming in <a href="#">USE statements</a> .
--

1 **7.5.3 Derived-type parameters**2 **7.5.3.1 Type parameter definition statement**

3 R732 *type-param-def-stmt* is *integer-type-spec*, *type-param-attr-spec* :: ■  
 4 ■ *type-param-decl-list*

5 R733 *type-param-decl* is *type-param-name* [ = *scalar-int-constant-expr* ]

6 C742 (R732) A *type-param-name* in a *type-param-def-stmt* in a *derived-type-def* shall be one of the *type-param-*  
 7 *names* in the *derived-type-stmt* of that *derived-type-def*.

8 C743 (R732) Each *type-param-name* in the *derived-type-stmt* in a *derived-type-def* shall appear exactly once as  
 9 a *type-param-name* in a *type-param-def-stmt* in that *derived-type-def*.

10 R734 *type-param-attr-spec* is KIND  
 11 or LEN

12 1 The derived type is parameterized if the *derived-type-stmt* has any *type-param-names*.

13 2 Each type parameter is itself of type integer. If its kind selector is omitted, the [kind type parameter](#) is default  
 14 integer.

15 3 The *type-param-attr-spec* explicitly specifies whether a type parameter is a kind parameter or a length parameter.

16 4 If a *type-param-decl* has a *scalar-int-constant-expr*, the type parameter has a default value which is specified by  
 17 the expression. If necessary, the value is converted according to the rules of intrinsic assignment (10.2.1.3) to a  
 18 value of the same kind as the type parameter.

19 5 A [type parameter](#) may be used as a primary in a [specification expression](#) (10.1.11) in the *derived-type-def*. A  
 20 [kind type parameter](#) may also be used as a primary in a [constant expression](#) (10.1.12) in the *derived-type-def*.

## NOTE 7.22

The following example uses derived-type parameters.

```
TYPE humongous_matrix(k, d)
  INTEGER, KIND :: k = KIND (0.0)
  INTEGER (SELECTED_INT_KIND (12)), LEN :: d
  !-- Specify a potentially nondefault kind for d.
  REAL (k) :: element (d, d)
END TYPE
```

In the following example, *dim* is declared to be a kind parameter, allowing generic overloading of procedures distinguished only by *dim*.

```
TYPE general_point(dim)
  INTEGER, KIND :: dim
  REAL :: coordinates(dim)
END TYPE
```

21 **7.5.3.2 Type parameter order**

22 1 [Type parameter order](#) is an ordering of the type parameters of a derived type; it is used for derived-type specifiers.

- 1 2 The **type parameter order** of a nonextended type is the order of the type parameter list in the derived-type  
 2 definition. The **type parameter order** of an **extended type** (7.5.7) consists of the **type parameter order** of its  
 3 **parent type** followed by any additional type parameters in the order of the type parameter list in the derived-type  
 4 definition.

**NOTE 7.23**

Given

```

TYPE :: t1 (k1, k2)
  INTEGER, KIND :: k1, k2
  REAL (k1) a (k2)
END TYPE
TYPE, EXTENDS(t1) :: t2 (k3)
  INTEGER, KIND :: k3
  LOGICAL (k3) flag
END TYPE

```

the type parameter order for type t1 is k1 then k2, and the type parameter order for type t2 is k1 then k2 then k3.

5 **7.5.4 Components**6 **7.5.4.1 Component definition statement**7 R735 *component-part* is [ *component-def-stmt* ] ...8 R736 *component-def-stmt* is *data-component-def-stmt*  
9 or *proc-component-def-stmt*10 R737 *data-component-def-stmt* is *declaration-type-spec* [ [ , *component-attr-spec-list* ] :: ] ■  
11 ■ *component-decl-list*12 R738 *component-attr-spec* is *access-spec*  
13 or ALLOCATABLE  
14 or CODIMENSION *lbracket coarray-spec rbracket*  
15 or CONTIGUOUS  
16 or DIMENSION ( *component-array-spec* )  
17 or POINTER18 R739 *component-decl* is *component-name* [ ( *component-array-spec* ) ] ■  
19 ■ [ *lbracket coarray-spec rbracket* ] ■  
20 ■ [ \* *char-length* ] [ *component-initialization* ]21 R740 *component-array-spec* is *explicit-shape-spec-list*  
22 or *deferred-shape-spec-list*  
2324 C744 (R737) No *component-attr-spec* shall appear more than once in a given *component-def-stmt*.25 C745 (R737) If neither the **POINTER** nor the **ALLOCATABLE** attribute is specified, the *declaration-type-spec*  
26 in the *component-def-stmt* shall specify an intrinsic type or a previously defined derived type.27 C746 (R737) If the **POINTER** or **ALLOCATABLE** attribute is specified, each *component-array-spec* shall be  
28 a *deferred-shape-spec-list*.29 C747 (R737) If a *coarray-spec* appears, it shall be a *deferred-coshape-spec-list* and the component shall have  
30 the **ALLOCATABLE** attribute.

- 1 C748 (R737) If a *coarray-spec* appears, the component shall not be of type `C_PTR` or `C_FUNPTR` from the  
2 intrinsic module `ISO_C_BINDING` (18.3.3).
- 3 C749 A data component whose type has a *coarray ultimate component* shall be a nonpointer nonallocatable  
4 scalar and shall not be a *coarray*.
- 5 C750 (R737) If neither the `POINTER` nor the `ALLOCATABLE` attribute is specified, each *component-array-*  
6 *spec* shall be an *explicit-shape-spec-list*.
- 7 C751 (R740) Each *bound* in the *explicit-shape-spec* shall be a *specification expression* in which there are no ref-  
8 erences to specification functions or the intrinsic functions `ALLOCATED`, `ASSOCIATED`, `EXTENDS_`  
9 `TYPE_OF`, `PRESENT`, or `SAME_TYPE_AS`, every specification inquiry reference is a *constant ex-*  
10 *pression*, and the value does not depend on the value of a variable.
- 11 C752 (R737) A component shall not have both the `ALLOCATABLE` and `POINTER` attributes.
- 12 C753 (R737) If the `CONTIGUOUS` attribute is specified, the component shall be an array with the `POINTER`  
13 attribute.
- 14 C754 (R739) The \* *char-length* option is permitted only if the component is of type character.
- 15 C755 (R736) Each *type-param-value* within a *component-def-stmt* shall be a colon or a *specification expres-*  
16 *sion* in which there are no references to specification functions or the intrinsic functions `ALLOCATED`,  
17 `ASSOCIATED`, `EXTENDS_TYPE_OF`, `PRESENT`, or `SAME_TYPE_AS`, every specification inquiry  
18 reference is a *constant expression*, and the value does not depend on the value of a variable.

**NOTE 7.24**

Because a type parameter is not an object, a *type-param-value* or a *bound* in an *explicit-shape-spec* can contain a *type-param-name*.

- 19 R741 *proc-component-def-stmt* is PROCEDURE ( [ *proc-interface* ] ) , ■  
20 ■ *proc-component-attr-spec-list* :: *proc-decl-list*

**NOTE 7.25**

See 15.4.3.6 for definitions of *proc-interface* and *proc-decl*.

- 21 R742 *proc-component-attr-spec* is *access-spec*  
22 or NOPASS  
23 or PASS [ ( *arg-name* ) ]  
24 or POINTER
- 25 C756 (R741) The same *proc-component-attr-spec* shall not appear more than once in a given *proc-component-*  
26 *def-stmt*.
- 27 C757 (R741) `POINTER` shall appear in each *proc-component-attr-spec-list*.
- 28 C758 (R741) If the procedure pointer component has an *implicit interface* or has no arguments, `NOPASS` shall  
29 be specified.
- 30 C759 (R741) If `PASS` (*arg-name*) appears, the *interface* of the procedure pointer component shall have a *dummy*  
31 *argument* named *arg-name*.
- 32 C760 (R741) `PASS` and `NOPASS` shall not both appear in the same *proc-component-attr-spec-list*.
- 33 1 The *declaration-type-spec* in the *data-component-def-stmt* specifies the type and type parameters of the com-  
34 ponents in the *component-decl-list*, except that the character length parameter may be specified or overridden  
35 for a component by the appearance of \* *char-length* in its *entity-decl*. The *component-attr-spec-list* in the *data-*  
36 *component-def-stmt* specifies the attributes whose keywords appear for the components in the *component-decl-list*,

1 except that the [DIMENSION attribute](#) may be specified or overridden for a component by the appearance of a  
 2 [component-array-spec](#) in its [component-decl](#), and the [CODIMENSION attribute](#) may be specified or overridden  
 3 for a component by the appearance of a [coarray-spec](#) in its [component-decl](#).

#### 4 7.5.4.2 Array components

5 1 A data component is an array if its [component-decl](#) contains a [component-array-spec](#) or its [data-component-def-stmt](#)  
 6 contains a DIMENSION clause. If the [component-decl](#) contains a [component-array-spec](#), it specifies the  
 7 array [rank](#), and if the array is explicit shape ([8.5.8.2](#)), the [array bounds](#); otherwise, the [component-array-spec](#) in  
 8 the DIMENSION clause specifies the array [rank](#), and if the array is explicit shape, the [array bounds](#).

##### NOTE 7.26

An example of a derived type definition with an array component is:

```
TYPE LINE
  REAL, DIMENSION (2, 2) :: COORD      !
                                     ! COORD(:,1) has the value of [X1, Y1]
                                     ! COORD(:,2) has the value of [X2, Y2]
  REAL                          :: WIDTH ! Line width in centimeters
  INTEGER                        :: PATTERN ! 1 for solid, 2 for dash, 3 for dot
END TYPE LINE
```

An example of declaring a variable LINE\_SEGMENT to be of the type LINE is:

```
TYPE (LINE)          :: LINE_SEGMENT
```

The scalar variable LINE\_SEGMENT has a component that is an array. In this case, the array is a subobject of a scalar. The double colon in the definition for COORD is required; the double colon in the definition for WIDTH and PATTERN is optional.

##### NOTE 7.27

An example of a derived type definition with an [allocatable](#) component is:

```
TYPE STACK
  INTEGER          :: INDEX
  INTEGER, ALLOCATABLE :: CONTENTS (:)
END TYPE STACK
```

For each scalar variable of type STACK, the shape of the component CONTENTS is determined by execution of an [ALLOCATE statement](#) or [assignment statement](#), or by [argument association](#).

##### NOTE 7.28

[Default initialization](#) of an [explicit-shape array](#) component can be specified by a [constant expression](#) consisting of an array constructor ([7.8](#)), or of a single scalar that becomes the value of each array element.

#### 9 7.5.4.3 Coarray components

10 1 A data component is a [coarray](#) if its [component-decl](#) contains a [coarray-spec](#) or its [data-component-def-stmt](#)  
 11 contains a CODIMENSION clause. If the [component-decl](#) contains a [coarray-spec](#) it specifies the [corank](#); otherwise,  
 12 the [coarray-spec](#) in the CODIMENSION clause specifies the [corank](#).

##### NOTE 7.29

An example of a derived type definition with a coarray component is:

```
TYPE GRID_TYPE
```



**NOTE 7.29 (cont.)**

```
REAL, ALLOCATABLE, CODIMENSION [:, :, :] :: GRID (:, :, :)  
END TYPE GRID_TYPE
```

An object of type `grid_type` cannot be an array, an [allocatable](#) object, a [coarray](#), or a [pointer](#).

1 **7.5.4.4 Pointer components**

- 2 1 A data component is a pointer ([5.4.9](#)) if its *component-attr-spec-list* contains the **POINTER** keyword. A procedure  
3 pointer component has the **POINTER** keyword in its *proc-component-attr-spec-list*.

**NOTE 7.30**

An example of a derived type definition with a pointer component is:

```
TYPE REFERENCE  
  INTEGER                               :: VOLUME, YEAR, PAGE  
  CHARACTER (LEN = 50)                  :: TITLE  
  PROCEDURE (printer_interface), POINTER :: PRINT => NULL()  
  CHARACTER, DIMENSION (:), POINTER    :: SYNOPSIS  
END TYPE REFERENCE
```

Any object of type REFERENCE will have the four nonpointer components VOLUME, YEAR, PAGE, and TITLE, the procedure pointer PRINT, which has an [explicit interface](#) the same as `printer_interface`, plus a pointer to an array of characters holding SYNOPSIS. The size of this [target](#) array will be determined by the length of the synopsis. The space for the [target](#) could be allocated ([9.7.1](#)) or the pointer component could be associated with a target by a [pointer assignment statement](#) ([10.2.2](#)).

4 **7.5.4.5 The passed-object dummy argument**

- 5 1 A [passed-object dummy argument](#) is a distinguished [dummy argument](#) of a procedure pointer component or  
6 [type-bound procedure](#) ([7.5.5](#)). It affects procedure overriding ([7.5.7.3](#)) and [argument association](#) ([15.5.2.2](#)).
- 7 2 If **NOPASS** is specified, the procedure pointer component or [type-bound procedure](#) has no [passed-object dummy](#)  
8 [argument](#).
- 9 3 If neither **PASS** nor **NOPASS** is specified or **PASS** is specified without *arg-name*, the first [dummy argument](#) of a  
10 procedure pointer component or [type-bound procedure](#) is its [passed-object dummy argument](#).
- 11 4 If **PASS** (*arg-name*) is specified, the [dummy argument](#) named *arg-name* is the [passed-object dummy argument](#) of  
12 the procedure pointer component or named [type-bound procedure](#).
- 13 C761 The [passed-object dummy argument](#) shall be a scalar, nonpointer, nonallocatable [dummy data object](#)  
14 with the same [declared type](#) as the type being defined; all of its [length type parameters](#) shall be assumed;  
15 it shall be [polymorphic](#) ([7.3.2.3](#)) if and only if the type being defined is [extensible](#) ([7.5.7](#)). It shall not  
16 have the **VALUE** attribute.

**NOTE 7.31**

If a procedure is bound to several types as a [type-bound procedure](#), different [dummy arguments](#) might be the [passed-object dummy argument](#) in different contexts.

17 **7.5.4.6 Default initialization for components**

- 18 1 [Default initialization](#) provides a means of automatically initializing pointer components to be [disassociated](#) or  
19 associated with specific [targets](#), and nonpointer nonallocatable components to have a particular value. [Allocatable](#)  
20 components are always initialized to unallocated.

- 1 2 A pointer variable or component is data-pointer-initialization compatible with a **target** if the pointer is **type**  
 2 **compatible** with the **target**, they have the same **rank**, all nondeferred **type parameters** of the pointer have the  
 3 same values as the corresponding **type parameters** of the **target**, and the **target** is **contiguous** if the pointer has  
 4 the **CONTIGUOUS** attribute.
- 5 R743 *component-initialization* **is** = *constant-expr*  
 6 **or** => *null-init*  
 7 **or** => *initial-data-target*
- 8 R744 *initial-data-target* **is** *designator*
- 9 C762 (R737) If *component-initialization* appears, a double-colon separator shall appear before the *component-*  
 10 *decl-list*.
- 11 C763 (R737) If *component-initialization* appears, every type parameter and **array bound** of the component  
 12 shall be a colon or **constant expression**.
- 13 C764 (R737) If => appears in *component-initialization*, **POINTER** shall appear in the *component-attr-spec-*  
 14 *list*. If = appears in *component-initialization*, neither **POINTER** nor **ALLOCATABLE** shall appear in  
 15 the *component-attr-spec-list*.
- 16 C765 If *initial-data-target* appears in a *component-initialization* in a *component-decl*, *component-name* shall be  
 17 data-pointer-initialization compatible with it.
- 18 C766 A *designator* that is an *initial-data-target* shall designate a nonallocatable, **noncoindexed** variable that  
 19 has the **TARGET** and **SAVE** attributes and does not have a **vector subscript**. Every subscript, sec-  
 20 tion subscript, substring starting point, and substring ending point in *designator* shall be a **constant**  
 21 **expression**.
- 22 3 If *null-init* appears for a pointer component, that component in any object of the type has an initial association  
 23 status of **disassociated** (3.61) or becomes disassociated as specified in 19.5.2.4.
- 24 4 If *initial-data-target* appears for a data pointer component, that component in any object of the type is initially  
 25 associated with the **target** or becomes associated with the **target** as specified in 19.5.2.3.
- 26 5 If *initial-proc-target* (15.4.3.6) appears in *proc-decl* for a procedure pointer component, that component in any  
 27 object of the type is initially associated with the **target** or becomes associated with the **target** as specified in  
 28 19.5.2.3.
- 29 6 If *constant-expr* appears for a nonpointer component, that component in any object of the type is initially defined  
 30 (19.6.3) or becomes defined as specified in 19.6.5 with the value determined from *constant-expr*. If necessary,  
 31 the value is converted according to the rules of intrinsic assignment (10.2.1.3) to a value that agrees in type,  
 32 type parameters, and shape with the component. If the component is of a type for which **default initialization** is  
 33 specified for a component, the **default initialization** specified by *constant-expr* overrides the **default initialization**  
 34 specified for that component. When one **initialization** overrides another it is as if only the overriding **initialization**  
 35 were specified (see Note 7.33). **Explicit initialization** in a **type declaration statement** (8.2) overrides **default**  
 36 **initialization** (see Note 7.32). Unlike **explicit initialization**, **default initialization** does not imply that the object  
 37 has the **SAVE** attribute.
- 38 7 A **subcomponent** (9.4.2) is **default-initialized** if the type of the object of which it is a component specifies **default**  
 39 **initialization** for that component, and the **subcomponent** is not a subobject of an object that is **default-initialized**  
 40 or **explicitly initialized**.
- 41 8 A type has **default initialization** if *component-initialization* is specified for any **direct component** of the type. An  
 42 object has **default initialization** if it is of a type that has **default initialization**.

**NOTE 7.32**

It is not required that [initialization](#) be specified for each component of a derived type. For example:

```
TYPE DATE
  INTEGER DAY
  CHARACTER (LEN = 5) MONTH
  INTEGER :: YEAR = 2008      ! Partial default initialization
END TYPE DATE
```

In the following example, the default initial value for the YEAR component of TODAY is overridden by [explicit initialization](#) in the [type declaration statement](#):

```
TYPE (DATE), PARAMETER :: TODAY = DATE (21, "Feb.", 2009)
```

**NOTE 7.33**

The default initial value of a component of derived type can be overridden by [default initialization](#) specified in the definition of the type. Continuing the example of Note 7.32:

```
TYPE SINGLE_SCORE
  TYPE (DATE) :: PLAY_DAY = TODAY
  INTEGER SCORE
  TYPE (SINGLE_SCORE), POINTER :: NEXT => NULL ( )
END TYPE SINGLE_SCORE
TYPE (SINGLE_SCORE) SETUP
```

The PLAY\_DAY component of SETUP receives its initial value from TODAY, overriding the [initialization](#) for the YEAR component.

**NOTE 7.34**

Arrays of structures can be declared with elements that are partially or totally initialized by default. Continuing the example of Note 7.33 :

```
TYPE MEMBER (NAME_LEN)
  INTEGER, LEN :: NAME_LEN
  CHARACTER (LEN = NAME_LEN) :: NAME = ''
  INTEGER :: TEAM_NO, HANDICAP = 0
  TYPE (SINGLE_SCORE), POINTER :: HISTORY => NULL ( )
END TYPE MEMBER
TYPE (MEMBER(9)) LEAGUE (36)      ! Array of partially initialized elements
TYPE (MEMBER(9)) :: ORGANIZER = MEMBER (9) ("I. Manage",1,5,NULL ( ))
```

ORGANIZER is [explicitly initialized](#), overriding the [default initialization](#) for an object of type MEMBER.

Allocated objects can also be initialized partially or totally. For example:

```
ALLOCATE (ORGANIZER % HISTORY)   ! A partially initialized object of type
                                ! SINGLE_SCORE is created.
```

**NOTE 7.35**

A pointer component of a derived type can have as its [target](#) an object of that derived type. The type definition can specify that in objects declared to be of this type, such a pointer is default initialized to [disassociated](#). For example:

```
TYPE NODE
```

**NOTE 7.35 (cont.)**

```

INTEGER                :: VALUE = 0
TYPE (NODE), POINTER  :: NEXT_NODE => NULL ( )
END TYPE

```

A type such as this can be used to construct linked lists of objects of type NODE. Linked lists can also be constructed using [allocatable](#) components.

**NOTE 7.36**

A pointer component of a derived type can be default initialized to have an initial [target](#).

```

TYPE NODE
  INTEGER                :: VALUE = 0
  TYPE (NODE), POINTER  :: NEXT_NODE => SENTINEL
END TYPE

TYPE(NODE), SAVE, TARGET :: SENTINEL

```

1 **7.5.4.7 Component order**

- 2 1 [Component order](#) is an ordering of the nonparent components of a derived type; it is used for intrinsic formatted input/output and [structure constructors](#) where [component keywords](#) are not used. [Parent components](#) are excluded from the [component order](#) of an [extended type](#) (7.5.7).
- 3  
4
- 5 2 The [component order](#) of a nonextended type is the order of the declarations of the components in the derived-type definition. The component order of an [extended type](#) consists of the [component order](#) of its [parent type](#) followed by any additional components in the order of their declarations in the extended derived-type definition.
- 6  
7

**NOTE 7.37**

Given the same type definitions as in Note 7.23, the [component order](#) of type T1 is just A (there is only one component), and the [component order](#) of type T2 is A then FLAG. The [parent component](#) (T1) does not participate in the [component order](#).

8 **7.5.4.8 Component accessibility**

9 R745 *private-components-stmt* is PRIVATE

10 C767 (R745) A *private-components-stmt* is permitted only if the type definition is within the specification part  
11 of a module.

- 12 1 The default accessibility for the components that are declared in a type's *component-part* is private if the type definition contains a *private-components-stmt*, and public otherwise. The accessibility of a component may be explicitly declared by an *access-spec*; otherwise its accessibility is the default for the type definition in which it is declared.
- 13  
14  
15
- 16 2 If a component is private, that component name is accessible only within the module containing the definition, and within its [descendants](#).
- 17

**NOTE 7.38**

Type parameters are not components. They are effectively always public.

**NOTE 7.39**

The accessibility of the components of a type is independent of the accessibility of the type name. It is possible to have all four combinations: a public type name with a public component, a private type name

**NOTE 7.39 (cont.)**

with a private component, a public type name with a private component, and a private type name with a public component.

**NOTE 7.40**

An example of a public type with private components is:

```
TYPE, PUBLIC :: POINT
  PRIVATE
  REAL :: X, Y
END TYPE POINT
```

Such a type definition can be accessed by use association; however, the components X and Y are accessible only within the module and its [descendants](#).

**NOTE 7.41**

The following example illustrates the use of an individual component [access-spec](#) to override the default accessibility:

```
TYPE MIXED
  PRIVATE
  INTEGER :: I
  INTEGER, PUBLIC :: J
END TYPE MIXED

TYPE (MIXED) :: M
```

The component M%J is accessible in any [scoping unit](#) where M is accessible; M%I is accessible only within the module containing the TYPE MIXED definition, and within its [descendants](#).

**7.5.5 Type-bound procedures**

- 1
- 2 R746 *type-bound-procedure-part* is *contains-stmt*
- 3 [ *binding-private-stmt* ]
- 4 [ *type-bound-proc-binding* ] ...
- 5 R747 *binding-private-stmt* is PRIVATE
- 6 C768 (R746) A *binding-private-stmt* is permitted only if the type definition is within the specification part of
- 7 a module.
- 8 R748 *type-bound-proc-binding* is *type-bound-procedure-stmt*
- 9 or *type-bound-generic-stmt*
- 10 or *final-procedure-stmt*
- 11 R749 *type-bound-procedure-stmt* is PROCEDURE [ [ , *binding-attr-list* ] :: ] *type-bound-proc-decl-list*
- 12 or PROCEDURE (*interface-name*), *binding-attr-list* :: *binding-name-list*
- 13 R750 *type-bound-proc-decl* is *binding-name* [ => *procedure-name* ]
- 14 C769 (R749) If => *procedure-name* appears in a *type-bound-proc-decl*, the double-colon separator shall appear.
- 15 C770 (R750) The *procedure-name* shall be the name of an accessible [module procedure](#) or an [external procedure](#)
- 16 that has an [explicit interface](#).
- 17 C771 A *binding-name* in a *type-bound-proc-decl* in a derived type definition shall not be the same as any other

- 1 *binding-name* within that derived type definition.
- 2 1 If => *procedure-name* does not appear in a *type-bound-proc-decl*, it is as though => *procedure-name* had appeared  
3 with a procedure name the same as the *binding name*.
- 4 R751 *type-bound-generic-stmt* is GENERIC [ , *access-spec* ] :: *generic-spec* => *binding-name-list*
- 5 C772 (R751) Within the *specification-part* of a module, each *type-bound-generic-stmt* shall specify, either im-  
6 plicitly or explicitly, the same accessibility as every other *type-bound-generic-stmt* with that *generic-spec*  
7 in the same derived type.
- 8 C773 (R751) Each *binding-name* in *binding-name-list* shall be the name of a specific *binding* of the type.
- 9 C774 A *binding-name* in a type-bound GENERIC statement shall not specify a specific *binding* that was  
10 *inherited* or specified previously for the same generic identifier in that derived type definition.
- 11 C775 (R751) If *generic-spec* is not *generic-name*, each of its specific *bindings* shall have a *passed-object dummy*  
12 *argument* (7.5.4.5).
- 13 C776 (R751) If *generic-spec* is OPERATOR ( *defined-operator* ), the *interface* of each *binding* shall be as  
14 specified in 15.4.3.4.2.
- 15 C777 (R751) If *generic-spec* is ASSIGNMENT ( = ), the *interface* of each *binding* shall be as specified in  
16 15.4.3.4.3.
- 17 C778 (R751) If *generic-spec* is *defined-io-generic-spec*, the *interface* of each *binding* shall be as specified in  
18 12.6.4.8. The type of the *dtv* argument shall be *type-name*.
- 19 R752 *binding-attr* is *access-spec*  
20 or DEFERRED  
21 or NON\_OVERRIDABLE  
22 or NOPASS  
23 or PASS [ ( *arg-name* ) ]
- 24 C779 (R752) The same *binding-attr* shall not appear more than once in a given *binding-attr-list*.
- 25 C780 (R749) If the *interface* of the *binding* has no *dummy argument* of the type being defined, NOPASS shall  
26 appear.
- 27 C781 (R749) If PASS ( *arg-name* ) appears, the *interface* of the *binding* shall have a *dummy argument* named  
28 *arg-name*.
- 29 C782 (R752) PASS and NOPASS shall not both appear in the same *binding-attr-list*.
- 30 C783 (R752) NON\_OVERRIDABLE and DEFERRED shall not both appear in the same *binding-attr-list*.
- 31 C784 (R752) DEFERRED shall appear if and only if *interface-name* appears.
- 32 C785 (R749) An overriding *binding* (7.5.7.3) shall have the DEFERRED attribute only if the *binding* it over-  
33 rides is deferred.
- 34 C786 (R749) A *binding* shall not override an *inherited binding* (7.5.7.2) that has the NON\_OVERRIDABLE  
35 attribute.
- 36 2 A type-bound procedure statement declares one or more specific *type-bound procedures*. A specific *type-bound*  
37 *procedure* can have a *passed-object dummy argument* (7.5.4.5). A *type-bound procedure* with the DEFERRED  
38 *attribute* is a deferred type-bound procedure. The DEFERRED keyword shall appear only in the definition of  
39 an *abstract type*.
- 40 3 A GENERIC statement declares a generic *type-bound procedure*, which is a type-bound *generic interface* for its

- 1 specific [type-bound procedures](#).
- 2 4 A [binding](#) of a type is a [type-bound procedure](#) (specific or generic), a [generic type-bound interface](#), or a [final](#)  
3 [subroutine](#). These are referred to as specific bindings, generic bindings, and final bindings respectively.
- 4 5 A [type-bound procedure](#) may be identified by a [binding name](#) in the scope of the type definition. This name is the  
5 *binding-name* for a specific [type-bound procedure](#), and the *generic-name* for a generic binding whose *generic-spec*  
6 is *generic-name*. A final binding, or a generic binding whose *generic-spec* is not *generic-name*, has no [binding](#)  
7 [name](#).
- 8 6 The [interface](#) of a specific [type-bound procedure](#) is that of the procedure specified by *procedure-name* or the  
9 [interface](#) specified by *interface-name*.

**NOTE 7.42**

An example of a type and a [type-bound procedure](#) is:

```
TYPE POINT
  REAL :: X, Y
CONTAINS
  PROCEDURE, PASS :: LENGTH => POINT_LENGTH
END TYPE POINT
...
```

and in the *module-subprogram-part* of the same module:

```
REAL FUNCTION POINT_LENGTH (A, B)
  CLASS (POINT), INTENT (IN) :: A, B
  POINT_LENGTH = SQRT ( (A%X - B%X)**2 + (A%Y - B%Y)**2 )
END FUNCTION POINT_LENGTH
```

- 10 7 The same *generic-spec* may be used in several GENERIC statements within a single derived-type definition. Each  
11 additional GENERIC statement with the same *generic-spec* extends the [generic interface](#).

**NOTE 7.43**

Unlike the situation with generic procedure names, a generic [type-bound procedure](#) name is not permitted to be the same as a specific [type-bound procedure](#) name in the same type (19.3).

- 12 8 The default accessibility for the [type-bound procedures](#) of a type is private if the type definition contains a *binding-*  
13 *private-stmt*, and public otherwise. The accessibility of a [type-bound procedure](#) may be explicitly declared by an  
14 *access-spec*; otherwise its accessibility is the default for the type definition in which it is declared.
- 15 9 A public [type-bound procedure](#) is accessible via any accessible object of the type. A private [type-bound procedure](#)  
16 is accessible only within the module containing the type definition, and within its [descendants](#).

**NOTE 7.44**

The accessibility of a [type-bound procedure](#) is not affected by a PRIVATE statement in the *component-part*; the accessibility of a data component is not affected by a PRIVATE statement in the *type-bound-procedure-part*.

## 17 7.5.6 Final subroutines

### 18 7.5.6.1 FINAL statement

19 R753 *final-procedure-stmt* is FINAL [ :: ] *final-subroutine-name-list*

20 C787 (R753) A *final-subroutine-name* shall be the name of a module procedure with exactly one [dummy](#)

1            **argument**. That argument shall be nonoptional and shall be a noncoarray, nonpointer, nonallocatable,  
2            nonpolymorphic variable of the derived type being defined. All length type parameters of the **dummy**  
3            **argument** shall be assumed. The **dummy argument** shall not have the **INTENT (OUT)** or **VALUE**  
4            attribute.

5            C788 (R753) A *final-subroutine-name* shall not be one previously specified as a **final subroutine** for that type.

6            C789 (R753) A **final subroutine** shall not have a **dummy argument** with the same **kind type parameters** and  
7            **rank** as the **dummy argument** of another **final subroutine** of that type.

8            C790 (R753) If a **final subroutine** has an **assumed-rank dummy argument**, no other **final subroutine** of that  
9            type shall have a **dummy argument** with the same **kind type parameters**.

10          1 The FINAL statement specifies that each procedure it names is a **final subroutine**. A **final subroutine** might be  
11          executed when a data entity of that type is finalized (7.5.6.2).

12          2 A derived type is **finalizable** if and only if it has a **final subroutine** or a nonpointer, nonallocatable component of  
13          **finalizable** type. A nonpointer data entity is **finalizable** if and only if it is of **finalizable** type. No other entity is  
14          **finalizable**.

#### NOTE 7.45

**Final subroutines** are effectively always “accessible”. They are called for entity **finalization** regardless of the accessibility of the type, its other **type-bound procedures**, or the subroutine name itself.

#### NOTE 7.46

**Final subroutines** are not **inherited** through type extension and cannot be overridden. The **final subroutines** of the **parent type** are called after any additional **final subroutines** of an **extended type** are called.

### 15          7.5.6.2 The finalization process

16          1 Only **finalizable** entities are finalized. When an entity is **finalized**, the following steps are carried out in sequence.

17            (1) If the **dynamic type** of the entity has a **final subroutine** whose **dummy argument** has the same **kind**  
18            **type parameters** and **rank** as the entity being finalized, it is called with the entity as an **actual**  
19            **argument**. Otherwise, if there is an **elemental final subroutine** whose **dummy argument** has the same  
20            **kind type parameters** as the entity being finalized, or a **final subroutine** whose **dummy argument** is  
21            **assumed-rank** with the same **kind type parameters** as the entity being finalized, it is called with the  
22            entity as an **actual argument**. Otherwise, no subroutine is called at this point.

23            (2) All **finalizable** components that appear in the type definition are finalized in a processor-dependent  
24            order. If the entity being finalized is an array, each **finalizable** component of each element of that  
25            entity is finalized separately.

26            (3) If the entity is of **extended type** and the **parent type** is **finalizable**, the **parent component** is finalized.

27          2 If several entities are to be finalized as a consequence of an event specified in 7.5.6.3, the order in which they  
28          are finalized is processor dependent. During this process, execution of a **final subroutine** for one of these entities  
29          shall not reference or define any of the other entities that have already been finalized.

#### NOTE 7.47

An implementation might need to ensure that when an event causes more than one **coarray** to be deallocated, they are deallocated in the same order on all **images** in the **current team**.

### 30          7.5.6.3 When finalization occurs

31          1 When an **intrinsic assignment statement** is executed (10.2.1.3), if the variable is not an unallocated allocatable  
32          variable, it is finalized after evaluation of *expr* and before the definition of the variable. If the variable is an  
33          allocated allocatable variable, or has an allocated allocatable subobject, that would be deallocated by intrinsic  
34          assignment, the finalization occurs before the deallocation.



- 1 2 When a pointer is deallocated its **target** is finalized. When an **allocatable** entity is deallocated, it is finalized  
2 unless it is the variable in an **intrinsic assignment statement** or a subobject thereof. If an error condition occurs  
3 during deallocation, it is processor dependent whether finalization occurs.
- 4 3 A nonpointer, nonallocatable object that is not a **dummy argument** or function result is finalized immediately  
5 before it would become undefined due to execution of a RETURN or **END statement** (19.6.6, item (3)).
- 6 4 A nonpointer nonallocatable local variable of a **BLOCK construct** is finalized immediately before it would become  
7 undefined due to termination of the **BLOCK construct** (19.6.6, item (23)).
- 8 5 If an executable construct references a nonpointer function, the result is finalized after execution of the innermost  
9 executable construct containing the reference.
- 10 6 If a **specification expression** in a **scoping unit** references a function, the result is finalized before execution of the  
11 executable constructs in the **scoping unit**.
- 12 7 When a procedure is invoked, a nonpointer, nonallocatable, **INTENT (OUT) dummy argument** of that procedure  
13 is finalized before it becomes undefined. The finalization caused by **INTENT (OUT)** is considered to occur within  
14 the invoked procedure; so for elemental procedures, an **INTENT (OUT)** argument will be finalized only if a scalar  
15 or elemental **final subroutine** is available, regardless of the rank of the **actual argument**.
- 16 8 If an object is allocated via pointer allocation and later becomes unreachable due to all pointers associated with  
17 that object having their **pointer association** status changed, it is processor dependent whether it is finalized. If it  
18 is finalized, it is processor dependent as to when the **final subroutines** are called.

**NOTE 7.48**

If **finalization** is used for storage management, it often needs to be combined with defined assignment.

**7.5.6.4 Entities that are not finalized**

- 19  
20 1 If image execution is **terminated**, either by an error (e.g. an allocation failure) or by execution of a *stop-stmt*,  
21 *error-stop-stmt*, or *end-program-stmt*, entities existing immediately prior to **termination** are not finalized.

**NOTE 7.49**

A nonpointer, nonallocatable object that has the **SAVE attribute** is never finalized as a direct consequence of the execution of a **RETURN** or **END statement**.

**7.5.7 Type extension****7.5.7.1 Extensible, extended, and abstract types**

- 24 1 A derived type, other than the type **C\_PTR** or **C\_FUNPTR** from the intrinsic module **ISO\_C\_BINDING**, that  
25 does not have the **BIND attribute** or the **SEQUENCE attribute** is an **extensible type**.
- 26 2 A type with the **EXTENDS attribute** is an **extended type**; its **parent type** is the type named in the **EXTENDS**  
27 *type-attr-spec*.

**NOTE 7.50**

The name of the **parent type** might be a local name introduced via renaming in a USE statement.

- 28 3 An **extensible type** that does not have the **EXTENDS attribute** is an **extension type** of itself only. An **extended**  
29 **type** is an **extension** of itself and of all types for which its **parent type** is an **extension**.
- 30 4 An **abstract type** is a type that has the **ABSTRACT attribute**.

**NOTE 7.51**

The **DEFERRED** attribute (7.5.5) defers the implementation of a **type-bound procedure** to **extensions** of the type; it can appear only in an **abstract type**. The **dynamic type** of an object cannot be **abstract**; therefore, a deferred **type-bound procedure** cannot be invoked. An **extension** of an **abstract type** need not be **abstract** if it has no deferred **type-bound procedures**. A short example of an **abstract type** is:

```
TYPE, ABSTRACT :: FILE_HANDLE
CONTAINS
    PROCEDURE(OPEN_FILE), DEFERRED, PASS(HANDLE) :: OPEN
    ...
END TYPE
```

For a more elaborate example see C.2.4.

1 **7.5.7.2 Inheritance**

- 2 1 An **extended type** includes all of the type parameters, all of the components, and the nonoverridden (7.5.7.3)  
 3 **type-bound procedures** of its **parent type**. These are **inherited** by the **extended type** from the **parent type**. They  
 4 retain all of the attributes that they had in the **parent type**. Additional type parameters, components, and  
 5 procedure **bindings** may be declared in the derived-type definition of the **extended type**.

**NOTE 7.52**

Inaccessible components and **bindings** of the **parent type** are also **inherited**, but they remain inaccessible in the **extended type**. Inaccessible entities occur if the type being extended is accessed via **use association** and has a private entity.

**NOTE 7.53**

An extensible derived type is not required to have any components, **bindings**, or parameters; an **extended type** is not required to have more components, **bindings**, or parameters than its **parent type**.

- 6 2 An **extended type** has a scalar, nonpointer, nonallocatable, **parent component** with the type and type parameters  
 7 of the **parent type**. The name of this component is the **parent type** name. If the **extended type** is defined in a  
 8 module, the **parent component** has the accessibility of the **parent type** in the module in which the **parent type**  
 9 was defined. Components of the **parent component** are **inheritance associated** (19.5.4) with the corresponding  
 10 components **inherited** from the **parent type**. An ancestor component of a type is the **parent component** of the  
 11 type or an ancestor component of the **parent component**.
- 12 3 If a generic **binding** specified in a type definition has the same *generic-spec* as an **inherited binding**, it extends  
 13 the **generic interface** and shall satisfy the requirements specified in 15.4.3.4.5.

**NOTE 7.54**

A component or type parameter declared in an **extended type** cannot have the same name as any accessible component or type parameter of its **parent type**.

**NOTE 7.55**

Examples:

```
TYPE POINT                                ! A base type
    REAL :: X, Y
END TYPE POINT

TYPE, EXTENDS(POINT) :: COLOR_POINT      ! An extension of TYPE(POINT)
    ! Components X and Y, and component name POINT, inherited from parent
```

**NOTE 7.55 (cont.)**

```

INTEGER :: COLOR
END TYPE COLOR_POINT

```

**1 7.5.7.3 Type-bound procedure overriding**

2 1 If a specific [type-bound procedure](#) specified in a type definition has the same [binding name](#) as an accessible  
3 [type-bound procedure](#) from the [parent type](#) then the [binding](#) specified in the type definition overrides the one  
4 from the [parent type](#).

5 2 The overriding and overridden [type-bound procedures](#) shall satisfy the following conditions.

- 6 • Either both shall have a [passed-object dummy argument](#) or neither shall.
- 7 • If the overridden [type-bound procedure](#) is pure then the overriding one shall also be pure.
- 8 • Either both shall be [elemental](#) or neither shall.
- 9 • They shall have the same number of [dummy arguments](#).
- 10 • [Passed-object dummy arguments](#), if any, shall correspond by name and position.
- 11 • [Dummy arguments](#) that correspond by position shall have the same names and [characteristics](#), except for  
12 the type of the [passed-object dummy arguments](#).
- 13 • Either both shall be subroutines or both shall be functions having the same result [characteristics](#) (15.3.3).
- 14 • If the overridden [type-bound procedure](#) is [PUBLIC](#) then the overriding one shall not be [PRIVATE](#).

15 3 A [binding](#) of a type and a [binding](#) of an [extension](#) of that type correspond if the latter [binding](#) is the same [binding](#)  
16 as the former, overrides a corresponding [binding](#), or is an [inherited](#) corresponding [binding](#).

**NOTE 7.56**

The following is an example of procedure overriding, expanding on the example in Note 7.42.

```

TYPE, EXTENDS (POINT) :: POINT_3D
  REAL :: Z
CONTAINS
  PROCEDURE, PASS :: LENGTH => POINT_3D_LENGTH
END TYPE POINT_3D
...

```

and in the *module-subprogram-part* of the same module:

```

REAL FUNCTION POINT_3D_LENGTH ( A, B )
  CLASS (POINT_3D), INTENT (IN) :: A
  CLASS (POINT), INTENT (IN) :: B
  SELECT TYPE(B)
    CLASS IS(POINT_3D)
      POINT_3D_LENGTH = SQRT( (A%X-B%X)**2 + (A%Y-B%Y)**2 + (A%Z-B%Z)**2 )
    RETURN
  END SELECT
  PRINT *, 'In POINT_3D_LENGTH, dynamic type of argument is incorrect.'
  STOP
END FUNCTION POINT_3D_LENGTH

```

**17 7.5.8 Derived-type values**

18 1 The component value of

- 19 • a pointer component is its [pointer association](#),

- 1 • an **allocatable** component is its allocation status and, if it is allocated, its **dynamic type** and type parameters,
- 2 **bounds** and value, and
- 3 • a nonpointer nonallocatable component is its value.

4 2 The set of values of a particular derived type consists of all possible sequences of the component values of its  
5 components.

## 6 7.5.9 Derived-type specifier

7 1 A derived-type specifier is used in several contexts to specify a particular derived type and type parameters.

8 R754 *derived-type-spec* is *type-name* [ ( *type-param-spec-list* ) ]

9 R755 *type-param-spec* is [ *keyword* = ] *type-param-value*

10 C791 (R754) *type-name* shall be the name of an accessible derived type.

11 C792 (R754) *type-param-spec-list* shall appear only if the type is parameterized.

12 C793 (R754) There shall be at most one *type-param-spec* corresponding to each parameter of the type. If a  
13 type parameter does not have a default value, there shall be a *type-param-spec* corresponding to that  
14 type parameter.

15 C794 (R755) The *keyword*= may be omitted from a *type-param-spec* only if the *keyword*= has been omitted  
16 from each preceding *type-param-spec* in the *type-param-spec-list*.

17 C795 (R755) Each *keyword* shall be the name of a parameter of the type.

18 C796 (R755) An asterisk may be used as a *type-param-value* in a *type-param-spec* only in the declaration of a  
19 **dummy argument** or **associate name** or in the allocation of a **dummy argument**.

20 2 Type parameter values that do not have **type parameter keywords** specified correspond to type parameters in type  
21 parameter order (7.5.3.2). If a **type parameter keyword** appears, the value corresponds to the type parameter  
22 named by the keyword. If necessary, the value is converted according to the rules of intrinsic assignment (10.2.1.3)  
23 to a value of the same kind as the type parameter.

24 3 The value of a type parameter for which no *type-param-value* has been specified is its default value.

## 25 7.5.10 Construction of derived-type values

26 1 A derived-type definition implicitly defines a corresponding **structure constructor** that allows construction of  
27 scalar values of that derived type. The type and type parameters of a constructed value are specified by a derived  
28 type specifier.

29 R756 *structure-constructor* is *derived-type-spec* ( [ *component-spec-list* ] )

30 R757 *component-spec* is [ *keyword* = ] *component-data-source*

31 R758 *component-data-source* is *expr*  
32 **or** *data-target*  
33 **or** *proc-target*

34 C797 (R756) The *derived-type-spec* shall not specify an **abstract type** (7.5.7).

35 C798 (R756) At most one *component-spec* shall be provided for a component.

36 C799 (R756) If a *component-spec* is provided for an ancestor component, a *component-spec* shall not be provided  
37 for any component that is **inheritance associated** with a **subcomponent** of that ancestor component.

- 1 C7100 (R756) A *component-spec* shall be provided for a nonallocatable component unless it has *default initialization* or is *inheritance associated* with a *subcomponent* of another component for which a *component-spec* is provided.
- 2  
3
- 4 C7101 (R757) The *keyword=* may be omitted from a *component-spec* only if the *keyword=* has been omitted from each preceding *component-spec* in the constructor.
- 5
- 6 C7102 (R757) Each *keyword* shall be the name of a component of the type.
- 7 C7103 (R756) The type name and all components of the type for which a *component-spec* appears shall be accessible in the *scoping unit* containing the *structure constructor*.
- 8
- 9 C7104 (R756) If *derived-type-spec* is a type name that is the same as a generic name, the *component-spec-list* shall not be a valid *actual-arg-spec-list* for a function reference that is resolvable as a generic reference to that name (15.5.5.2).
- 10  
11
- 12 C7105 (R758) A *data-target* shall correspond to a data pointer component; a *proc-target* shall correspond to a procedure pointer component.
- 13
- 14 C7106 (R758) A *data-target* shall have the same *rank* as its corresponding component.

**NOTE 7.57**

The form 'name(...)' is interpreted as a generic *function-reference* if possible; it is interpreted as a *structure-constructor* only if it cannot be interpreted as a generic *function-reference*.

- 15 2 In the absence of a *component keyword*, each *component-data-source* is assigned to the corresponding *component* in *component order* (7.5.4.7). If a *component keyword* appears, the *expr* is assigned to the *component* named by the keyword. For a nonpointer *component*, the *declared type* and *type parameters* of the *component* and *expr* shall conform in the same way as for a *variable* and *expr* in an *intrinsic assignment statement* (10.2.1.2). If necessary, each value of intrinsic type is converted according to the rules of intrinsic assignment (10.2.1.3) to a value that agrees in type and type parameters with the corresponding component of the derived type. For a nonpointer nonallocatable component, the shape of the expression shall conform with the shape of the component.
- 16  
17  
18  
19  
20  
21
- 22 3 If a component with *default initialization* has no corresponding *component-data-source*, then the *default initialization* is applied to that component. If an *allocatable* component has no corresponding *component-data-source*, then that component has an allocation status of unallocated.
- 23  
24

**NOTE 7.58**

Because no *parent components* appear in the defined *component ordering*, a value for a *parent component* can be specified only with a *component keyword*. Examples of equivalent values using types defined in Note 7.55:

```
! Create values with components x = 1.0, y = 2.0, color = 3.
TYPE(POINT) :: PV = POINT(1.0, 2.0)      ! Assume components of TYPE(POINT)
                                           ! are accessible here.
...
COLOR_POINT( point=point(1,2), color=3)  ! Value for parent component
COLOR_POINT( point=PV, color=3)          ! Available even if TYPE(point)
                                           ! has private components
COLOR_POINT( 1, 2, 3)                    ! All components of TYPE(point)
                                           ! need to be accessible.
```

- 25 4 A *structure constructor* shall not appear before the referenced type is defined.

**NOTE 7.59**

This example illustrates a derived-type constant expression using a derived type defined in Note 7.15:

**NOTE 7.59 (cont.)**

```
PERSON (21, 'JOHN SMITH')
```

This could also be written as

```
PERSON (NAME = 'JOHN SMITH', AGE = 21)
```

**NOTE 7.60**

An example constructor using the derived type GENERAL\_POINT defined in Note 7.22 is

```
general_point(dim=3) ( [ 1., 2., 3. ] )
```

- 1 5 For a pointer component, the corresponding *component-data-source* shall be an allowable *data-target* or *proc-target* for such a pointer in a *pointer assignment statement* (10.2.2). If the component data source is a pointer, the association of the component is that of the pointer; otherwise, the component is *pointer associated* with the component data source.
- 2  
3  
4

**NOTE 7.61**

For example, if the variable TEXT were declared (8.2) to be

```
CHARACTER, DIMENSION (1:400), TARGET :: TEXT
```

and BIBLIO were declared using the derived-type definition REFERENCE in Note 7.30

```
TYPE (REFERENCE) :: BIBLIO
```

the statement

```
BIBLIO = REFERENCE (1, 1987, 1, "This is the title of the referenced &
&paper", SYNOPSIS=TEXT)
```

is valid and associates the pointer component SYNOPSIS of the object BIBLIO with the *target* object TEXT. The keyword SYNOPSIS is required because the fifth component of the type REFERENCE is a procedure pointer component, not a data pointer component of type character. It is not necessary to specify a *proc-target* for the procedure pointer component because it has *default initialization*.

- 5 6 If a component of a derived type is *allocatable*, the corresponding constructor expression shall be a reference to the intrinsic function NULL with no arguments, an *allocatable* entity of the same *rank*, or shall evaluate to an entity of the same *rank*. If the expression is a reference to the intrinsic function NULL, the corresponding component of the constructor has a status of unallocated.
- 6  
7  
8
- 9 7 If the component is *allocatable* and the expression is an *allocatable* entity, the corresponding component of the constructor has the same allocation status as that *allocatable* entity. If it is allocated, it has the same *bounds*; if a length parameter of the component is *deferred*, its value is the same as the corresponding parameter of the expression. If the component is *polymorphic*, it has the same *dynamic type* and value; otherwise, it has the value converted, if necessary, to the *declared type* of the component.
- 10  
11  
12  
13
- 14 8 If the component is *allocatable* and the expression is not an *allocatable* entity, the component has an allocation status of allocated and the same *bounds* as the expression; if a length parameter of the component is *deferred*, its value is the same as the corresponding parameter of the expression. If the component is *polymorphic*, it has the same *dynamic type* and value; otherwise, it has the value converted, if necessary, to the *declared type* of the component.
- 15  
16  
17  
18

## 7.5.11 Derived-type operations and assignment

1 Intrinsic assignment of derived-type entities is described in 10.2.1. This document does not specify any intrinsic operations on derived-type entities. Any operation on derived-type entities or defined assignment (10.2.1.4) for derived-type entities shall be defined explicitly by a function or a subroutine, and a generic interface (7.5.5, 15.4.3.2).

## 7.6 Enumerations and enumerators

1 An enumeration is a set of enumerators. An enumerator is a named integer constant. An enumeration definition specifies the enumeration and its set of enumerators of the corresponding integer kind.

R759 *enum-def* is *enum-def-stmt*  
*enumerator-def-stmt*  
 [ *enumerator-def-stmt* ] ...  
*end-enum-stmt*

R760 *enum-def-stmt* is ENUM, BIND(C)

R761 *enumerator-def-stmt* is ENUMERATOR [ :: ] *enumerator-list*

R762 *enumerator* is *named-constant* [ = *scalar-int-constant-expr* ]

R763 *end-enum-stmt* is END ENUM

C7107 (R761) If = appears in an *enumerator*, a double-colon separator shall appear before the *enumerator-list*.

2 For an enumeration, the kind is selected such that an integer type with that kind is interoperable (18.3.2) with the corresponding C enumeration type. The corresponding C enumeration type is the type that would be declared by a C enumeration specifier (6.7.2.2 of ISO/IEC 9899:2011) that specified C enumeration constants with the same values as those specified by the *enum-def*, in the same order as specified by the *enum-def*.

3 The companion processor (5.5.7) shall be one that uses the same representation for the types declared by all C enumeration specifiers that specify the same values in the same order.

### NOTE 7.62

If a companion processor uses an unsigned type to represent a given enumeration type, the Fortran processor will use the signed integer type of the same width for the enumeration, even though some of the values of the C enumerators might not be representable in this signed integer type. The types of any such enumerators will be interoperable with the type declared in the C enumeration.

### NOTE 7.63

ISO/IEC 9899:2011 guarantees the enumeration constants fit in a C int (6.7.2.2 of ISO/IEC 9899:2011). Therefore, the Fortran processor can evaluate all enumerator values using the integer type with kind parameter C\_INT, and then determine the kind parameter of the integer type that is interoperable with the corresponding C enumerated type.

### NOTE 7.64

ISO/IEC 9899:2011 specifies that two enumeration types are compatible only if they specify enumeration constants with the same names and same values in the same order. This document further requires that a C processor that is to be a companion processor of a Fortran processor use the same representation for two enumeration types if they both specify enumeration constants with the same values in the same order, even if the names are different.

4 An enumerator is treated as if it were explicitly declared with the PARAMETER attribute. The enumerator is defined in accordance with the rules of intrinsic assignment (10.2) with the value determined as follows.

- 1 (1) If *scalar-int-constant-expr* is specified, the value of the enumerator is the result of *scalar-int-constant-*  
 2 *expr*.  
 3 (2) If *scalar-int-constant-expr* is not specified and the enumerator is the first enumerator in *enum-def*,  
 4 the enumerator has the value 0.  
 5 (3) If *scalar-int-constant-expr* is not specified and the enumerator is not the first enumerator in *enum-*  
 6 *def*, its value is the result of adding 1 to the value of the enumerator that immediately precedes it  
 7 in the *enum-def*.

**NOTE 7.65**

Example of an enumeration definition:

```
ENUM, BIND(C)
  ENUMERATOR :: RED = 4, BLUE = 9
  ENUMERATOR YELLOW
END ENUM
```

The [kind type parameter](#) for this enumeration is processor dependent, but the processor is required to select a kind sufficient to represent the values 4, 9, and 10, which are the values of its enumerators. The following declaration might be equivalent to the above enumeration definition.

```
INTEGER (SELECTED_INT_KIND (2)), PARAMETER :: RED = 4, BLUE = 9, YELLOW = 10
```

An entity of the same [kind type parameter](#) value can be declared using the intrinsic function [KIND](#) with one of the enumerators as its argument, for example

```
INTEGER (KIND (RED)) :: X
```

**NOTE 7.66**

There is no difference in the effect of declaring the enumerators in multiple ENUMERATOR statements or in a single ENUMERATOR statement. The order in which the enumerators in an enumeration definition are declared is significant, but the number of ENUMERATOR statements is not.

## 8 7.7 Binary, octal, and hexadecimal literal constants

- 9 1 A binary, octal, or hexadecimal constant (*boz-literal-constant*) is a sequence of digits that represents an ordered  
 10 sequence of bits. Such a constant has no type.

11 R764 *boz-literal-constant* is *binary-constant*  
 12 or *octal-constant*  
 13 or *hex-constant*

14 R765 *binary-constant* is B ' *digit* [ *digit* ] ... '  
 15 or B " *digit* [ *digit* ] ... "

16 C7108 (R765) *digit* shall have one of the values 0 or 1.

17 R766 *octal-constant* is O ' *digit* [ *digit* ] ... '  
 18 or O " *digit* [ *digit* ] ... "

19 C7109 (R766) *digit* shall have one of the values 0 through 7.

20 R767 *hex-constant* is Z ' *hex-digit* [ *hex-digit* ] ... '  
 21 or Z " *hex-digit* [ *hex-digit* ] ... "

22 R768 *hex-digit* is *digit*  
 23 or A



1                   or B  
 2                   or C  
 3                   or D  
 4                   or E  
 5                   or F

- 6 2 The *hex-digits* A through F represent the numbers ten through fifteen, respectively; they may be represented  
 7 by their lower-case equivalents. Each digit of a *boz-literal-constant* represents a sequence of bits, according to  
 8 its numerical interpretation, using the model of 16.3, with  $z$  equal to one for binary constants, three for octal  
 9 constants or four for hexadecimal constants. A *boz-literal-constant* represents a sequence of bits that consists of  
 10 the concatenation of the sequences of bits represented by its digits, in the order the digits are specified. The  
 11 positions of bits in the sequence are numbered from right to left, with the position of the rightmost bit being zero.  
 12 The length of a sequence of bits is the number of bits in the sequence. The processor shall allow the position  
 13 of the leftmost nonzero bit to be at least  $z - 1$ , where  $z$  is the maximum value that could result from invoking  
 14 the intrinsic function `STORAGE_SIZE` (16.9.184) with an argument that is a real or integer scalar of any kind  
 15 supported by the processor.

16 C7110 (R764) A *boz-literal-constant* shall appear only as a *data-stmt-constant* in a *DATA statement*, or where  
 17 explicitly allowed in subclause 16.9 as an actual argument of an *intrinsic* procedure.

## 18 7.8 Construction of array values

- 19 1 An array constructor constructs a rank-one array value from a sequence of scalar values, array values, and implied  
 20 DO loops.

21 R769 *array-constructor*           is *( / ac-spec / )*  
 22                                       or *lbracket ac-spec rbracket*

23 R770 *ac-spec*                       is *type-spec ::*  
 24                                       or *[ type-spec :: ] ac-value-list*

25 R771 *lbracket*                     is *[*

26 R772 *rbracket*                     is *]*

27 R773 *ac-value*                    is *expr*  
 28                                       or *ac-IMPLIED-DO*

29 R774 *ac-IMPLIED-DO*               is *( ac-value-list , ac-IMPLIED-DO-control )*

30 R775 *ac-IMPLIED-DO-control*       is *[ integer-type-spec :: ] ac-do-variable = scalar-int-expr , ■*  
 31                                       ■ *scalar-int-expr [ , scalar-int-expr ]*

32 R776 *ac-do-variable*             is *do-variable*

33 C7111 (R770) If *type-spec* is omitted, each *ac-value* expression in the *array-constructor* shall have the same  
 34 declared type and kind type parameters.

35 C7112 (R770) If *type-spec* specifies an intrinsic type, each *ac-value* expression in the *array-constructor* shall be  
 36 of an intrinsic type that is in type conformance with a variable of type *type-spec* as specified in Table  
 37 10.8.

38 C7113 (R770) If *type-spec* specifies a derived type, the declared type of each *ac-value* expression in the *array-*  
 39 *constructor* shall be that derived type and shall have the same kind type parameter values as specified  
 40 by *type-spec*.

41 C7114 (R773) An *ac-value* shall not be unlimited polymorphic.

- 1 C7115 (R773) The **declared type** of an *ac-value* shall not be abstract.
- 2 C7116 (R774) The *ac-do-variable* of an *ac-implied-do* that is in another *ac-implied-do* shall not appear as the  
3 *ac-do-variable* of the containing *ac-implied-do*.
- 4 2 If *type-spec* is omitted, corresponding **length type parameters** of the **declared type** of each *ac-value* expression  
5 shall have the same value; in this case, the **declared type** and **type parameters** of the array constructor are those  
6 of the *ac-value* expressions.
- 7 3 If *type-spec* appears, it specifies the **declared type** and type parameters of the array constructor. Each *ac-value*  
8 expression in the *array-constructor* shall be compatible with intrinsic assignment to a variable of this type and  
9 type parameters. Each value is converted to the type and **type parameters** of the *array-constructor* in accordance  
10 with the rules of intrinsic assignment (10.2.1.3).
- 11 4 The **dynamic type** of an array constructor is the same as its **declared type**.
- 12 5 The character length of an *ac-value* in an *ac-implied-do* whose iteration count is zero shall not depend on the  
13 value of the *ac-do-variable* and shall not depend on the value of an expression that is not a **constant expression**.
- 14 6 If an *ac-value* is a scalar expression, its value specifies an element of the array constructor. If an *ac-value* is  
15 an array expression, the values of the elements of the expression, in array element order (9.5.3.2), specify the  
16 corresponding sequence of elements of the array constructor. If an *ac-value* is an *ac-implied-do*, it is expanded to  
17 form a sequence of elements under the control of the *ac-do-variable*, as in the **DO construct** (11.1.7.4).
- 18 7 For an *ac-implied-do*, the loop initialization and execution is the same as for a **DO construct**.
- 19 8 An empty sequence forms a zero-sized array.

**NOTE 7.67**

A one-dimensional array can be reshaped into any allowable array shape using the intrinsic function **RESHAPE** (16.9.163). An example is:

```
X = (/ 3.2, 4.01, 6.5 /)
Y = RESHAPE (SOURCE = [ 2.0, [ 4.5, 4.5 ], X ], SHAPE = [ 3, 2 ])
```

This results in Y having the 3 × 2 array of values:

```
2.0  3.2
4.5  4.01
4.5  6.5
```

**NOTE 7.68**

Examples of array constructors containing an implied DO are:

```
(/ (I, I = 1, 1075) /)
```

and

```
[ 3.6, (3.6 / I, I = 1, N) ]
```

**NOTE 7.69**

Using the type definition for PERSON in Note 7.15, an example of the construction of a derived-type array value is:

```
[ PERSON (40, 'SMITH'), PERSON (20, 'JONES') ]
```

**NOTE 7.70**

Using the type definition for LINE in Note 7.26, an example of the construction of a derived-type scalar value with a rank-2 array component is:

```
LINE (RESHAPE ( [ 0.0, 0.0, 1.0, 2.0 ], [ 2, 2 ] ), 0.1, 1)
```

The intrinsic function RESHAPE is used to construct a value that represents a solid line from (0, 0) to (1, 2) of width 0.1 centimeters.

**NOTE 7.71**

Examples of zero-size array constructors are:

```
[ INTEGER :: ]  
[ ( I, I = 1, 0 ) ]
```

**NOTE 7.72**

An example of an array constructor that specifies a length type parameter:

```
[ CHARACTER(LEN=7) :: 'Takata', 'Tanaka', 'Hayashi' ]
```

In this constructor, without the type specification, it would have been necessary to specify all of the constants with the same character length.

1

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## 8 Attribute declarations and specifications

### 8.1 Attributes of procedures and data objects

- 1 Every data object has a type and **rank** and may have **type parameters** and other properties that determine the uses of the object. Collectively, these properties are the **attributes** of the object. The **declared type** of a named data object is either specified explicitly in a **type declaration statement** or determined implicitly by the first letter of its name (8.7). All of its **attributes** may be specified in a **type declaration statement** or individually in separate specification statements.
- 2 A function has a type and **rank** and may have type parameters and other **attributes** that determine the uses of the function. The type, **rank**, and type parameters are the same as those of the **function result**.
- 3 A subroutine does not have a type, **rank**, or type parameters, but may have other **attributes** that determine the uses of the subroutine.

### 8.2 Type declaration statement

R801 *type-declaration-stmt*            **is** *declaration-type-spec* [ [ , *attr-spec* ] ... :: ] *entity-decl-list*

- 1 The type declaration statement specifies the **declared type** of the entities in the entity declaration list. The type and type parameters are those specified by *declaration-type-spec*, except that the character length type parameter may be overridden for an entity by the appearance of \* *char-length* in its *entity-decl*.

R802 *attr-spec*                            **is** *access-spec*  
**or** ALLOCATABLE  
**or** ASYNCHRONOUS  
**or** CODIMENSION *lbracket coarray-spec rbracket*  
**or** CONTIGUOUS  
**or** DIMENSION ( *array-spec* )  
**or** EXTERNAL  
**or** INTENT ( *intent-spec* )  
**or** INTRINSIC  
**or** *language-binding-spec*  
**or** OPTIONAL  
**or** PARAMETER  
**or** POINTER  
**or** PROTECTED  
**or** SAVE  
**or** TARGET  
**or** VALUE  
**or** VOLATILE

C801 (R801) The same *attr-spec* shall not appear more than once in a given *type-declaration-stmt*.

C802 (R801) If a *language-binding-spec* with a NAME= specifier appears, the *entity-decl-list* shall consist of a single *entity-decl*.

C803 (R801) If a *language-binding-spec* is specified, the *entity-decl-list* shall not contain any procedure names.

- 2 The type declaration statement also specifies the **attributes** whose keywords appear in the *attr-spec*, except that

1 the **DIMENSION** attribute may be specified or overridden for an entity by the appearance of *array-spec* in its  
 2 *entity-decl*, and the **CODIMENSION** attribute may be specified or overridden for an entity by the appearance of  
 3 *coarray-spec* in its *entity-decl*.

4 R803 *entity-decl* is *object-name* [ ( *array-spec* ) ] ■  
 5 ■ [ *lbracket coarray-spec rbracket* ] ■  
 6 ■ [ \* *char-length* ] [ *initialization* ]  
 7 or *function-name* [ \* *char-length* ]

8 C804 (R803) If the entity is not of type character, \* *char-length* shall not appear.

9 C805 (R801) If *initialization* appears, a double-colon separator shall appear before the *entity-decl-list*.

10 C806 (R801) If the PARAMETER keyword appears, *initialization* shall appear in each *entity-decl*.

11 C807 (R803) An *initialization* shall not appear if *object-name* is a **dummy argument**, a function result, an object  
 12 in a named **common block** unless the type declaration is in a block data program unit, an object in **blank common**, an  
 13 **allocatable** variable, or an **automatic data object**.

14 C808 (R803) The *function-name* shall be the name of an external function, an intrinsic function, a **dummy**  
 15 **function**, a **procedure pointer**, or a statement function.

16 R804 *object-name* is *name*

17 C809 (R804) The *object-name* shall be the name of a data object.

18 R805 *initialization* is = *constant-expr*  
 19 or => *null-init*  
 20 or => *initial-data-target*

21 R806 *null-init* is *function-reference*

22 C810 (R803) If => appears in *initialization*, the entity shall have the **POINTER** attribute. If = appears in  
 23 *initialization*, the entity shall not have the **POINTER** attribute.

24 C811 (R803) If *initial-data-target* appears, *object-name* shall be data-pointer-initialization compatible with it  
 25 (7.5.4.6).

26 C812 (R806) The *function-reference* shall be a reference to the intrinsic function **NULL** with no arguments.

27 3 A name that identifies a specific intrinsic function has a type as specified in 16.8. An explicit type declaration statement is not  
 28 required; however, it is permitted. Specifying a type for a generic intrinsic function name in a type declaration  
 29 statement has no effect.

30 4 If *initialization* appears for a nonpointer entity,  
 31 • its type and **type parameters** shall conform as specified for intrinsic assignment (10.2.1.2);  
 32 • if the entity has implied shape, the rank of *initialization* shall be the same as the rank of the entity;  
 33 • if the entity does not have implied shape, *initialization* shall either be scalar or have the same shape as the  
 34 entity.

#### NOTE 8.1

Examples of type declaration statements:

```
REAL A (10)
LOGICAL, DIMENSION (5, 5) :: MASK1, MASK2
COMPLEX :: CUBE_ROOT = (-0.5, 0.866)
INTEGER, PARAMETER :: SHORT = SELECTED_INT_KIND (4)
INTEGER (SHORT) K      ! Range at least -9999 to 9999.
```

## NOTE 8.1 (cont.)

```

REAL (KIND (0.0D0)) B1
REAL (KIND = 2) B2
COMPLEX (KIND = KIND (0.0D0)) :: C
CHARACTER (LEN = 10, KIND = 2) TEXT2
CHARACTER CHAR, STRING *20
TYPE (PERSON) :: CHAIRMAN
TYPE(NODE), POINTER :: HEAD => NULL ( )
TYPE (humongous_matrix (k=8, d=1000)) :: MAT

```

(The last line above uses a type definition from Note 7.22.)

## 8.3 Automatic data objects

1 An *automatic data object* is a *nondummy data object* with a *type parameter* or *array bound* that depends on the value of a *specification-expr* that is not a *constant expression*.

C813 An *automatic data object* shall not have the *SAVE attribute*.

2 If a *type parameter* in a *declaration-type-spec* or in a *char-length* in an *entity-decl* for a *local variable* of a subprogram or *BLOCK construct* is defined by an expression that is not a *constant expression*, the *type parameter value* is established on entry to a procedure defined by the subprogram, or on execution of the *BLOCK statement*, and is not affected by any redefinition or undefinition of the variables in the expression during execution of the procedure or *BLOCK construct*.

## 8.4 Initialization

1 The appearance of *initialization* in an *entity-decl* for an entity without the *PARAMETER attribute* specifies that the entity is a variable with *explicit initialization*. *Explicit initialization* alternatively may be specified in a *DATA statement* unless the variable is of a derived type for which *default initialization* is specified. If *initialization* is = *constant-expr*, the variable is initially defined with the value specified by the *constant-expr*; if necessary, the value is converted according to the rules of intrinsic assignment (10.2.1.3) to a value that agrees in type, type parameters, and shape with the variable. A variable, or part of a variable, shall not be *explicitly initialized* more than once in a program. If the variable is an array, it shall have its shape specified in either the type declaration statement or a previous attribute specification statement in the same *scoping unit*.

2 If *null-init* appears, the initial association status of the object is *disassociated*. If *initial-data-target* appears, the object is initially associated with the *target*.

3 *Explicit initialization* of a variable that is not in a *common block* implies the *SAVE attribute*, which may be confirmed by explicit specification.

## 8.5 Attributes

### 8.5.1 Attribute specification

1 An *attribute* may be explicitly specified by an *attr-spec* in a *type declaration statement* or by an attribute specification statement (8.6). The following constraints apply to *attributes*.

C814 An entity shall not be explicitly given any *attribute* more than once in a *scoping unit*.

C815 An *array-spec* for a nonallocatable nonpointer function result shall be an *explicit-shape-spec-list*.

## 8.5.2 Accessibility attribute

1 The accessibility attribute specifies the accessibility of an entity via a particular identifier.

R807 *access-spec* is PUBLIC  
or PRIVATE

C816 An *access-spec* shall appear only in the *specification-part* of a module.

2 An *access-spec* in a type declaration statement specifies the accessibility of the names of all the entities declared by that statement. An *access-spec* in a *derived-type-stmt* specifies the accessibility of the derived type name. Accessibility can also be specified by an *access-stmt*.

3 An identifier that is specified in a module or is accessible in a module by use association has either the **PUBLIC attribute** or **PRIVATE attribute**. An identifier whose accessibility is not explicitly specified has default accessibility (8.6.1).

4 The default accessibility attribute for a module is PUBLIC unless it has been changed by a PRIVATE statement. Only an identifier that has the PUBLIC attribute in that module is available to be accessed from that module by use association.

### NOTE 8.2

An identifier can only be accessed by **use association** if it has the **PUBLIC attribute** in the module from which it is accessed. It can nonetheless have the **PRIVATE attribute** in a module in which it is accessed by **use association**, and therefore not be available by **use association** from that module.

### NOTE 8.3

An example of an accessibility specification is:

```
REAL, PRIVATE :: X, Y, Z
```

## 8.5.3 ALLOCATABLE attribute

1 A variable with the **ALLOCATABLE attribute** is a variable for which space is allocated during execution.

### NOTE 8.4

Only variables and components can have the **ALLOCATABLE attribute**. The result of referencing a function whose result variable has the **ALLOCATABLE attribute** is a value that does not itself have the **ALLOCATABLE attribute**.

## 8.5.4 ASYNCHRONOUS attribute

1 An entity with the **ASYNCHRONOUS attribute** is a variable that may be subject to asynchronous input/output or asynchronous communication.

2 The **base object** of a variable shall have the **ASYNCHRONOUS attribute** in a **scoping unit** if

- the variable is a **dummy argument** or appears in an executable statement or **specification expression** in that **scoping unit**, and
- any statement of the **scoping unit** is executed while the variable is a pending input/output storage sequence affector (12.6.2.5) or a pending communication affector (18.10.4).

3 Use of a variable in an asynchronous data transfer statement can imply the **ASYNCHRONOUS attribute**; see subclause 12.6.2.5.

4 An object with the **ASYNCHRONOUS attribute** may be associated with an object that does not have the **ASYNCHRONOUS attribute**, including by **use** (14.2.2) or **host** association (19.5.1.4). If an object that is not a



1 **local variable** of a **BLOCK construct** is specified to have the **ASYNCHRONOUS attribute** in the *specification-*  
 2 *part* of the construct, the object has the attribute within the construct even if it does not have the attribute  
 3 outside the construct. If an object has the **ASYNCHRONOUS attribute**, then all of its subobjects also have the  
 4 **ASYNCHRONOUS attribute**.

#### NOTE 8.5

The **ASYNCHRONOUS attribute** specifies the variables that might be associated with a pending input/output storage sequence (the actual memory locations on which asynchronous input/output is being performed) while the **scoping unit** is in execution. This information could be used by the compiler to disable certain code motion optimizations.

### 5 8.5.5 BIND attribute for data entities

6 1 The **BIND attribute** for a variable or **common block** specifies that it is capable of interoperating with a C variable  
 7 whose name has external linkage (18.9).

8 R808 *language-binding-spec* is BIND (C [ , NAME = *scalar-default-char-constant-expr* ])

9 C817 An entity with the **BIND attribute** shall be a **common block**, variable, type, or procedure.

10 C818 A variable with the **BIND attribute** shall be declared in the specification part of a module.

11 C819 A variable with the **BIND attribute** shall be **interoperable** (18.3).

12 C820 Each variable of a **common block** with the **BIND attribute** shall be **interoperable**.

13 2 If the value of the *scalar-default-char-constant-expr* after discarding leading and trailing blanks has nonzero  
 14 length, it shall be valid as an identifier on the **companion processor**.

#### NOTE 8.6

ISO/IEC 9899:2011 provides a facility for creating C identifiers whose characters are not restricted to the C basic character set. Such a C identifier is referred to as a universal character name (6.4.3 of ISO/IEC 9899:2011). The name of such a C identifier might include characters that are not part of the representation method used by the processor for default character. If so, the C entity cannot be referenced from Fortran.

15 3 The **BIND attribute** for a **common block** implies the **SAVE attribute**, which may be confirmed by explicit specification.

### 16 8.5.6 CODIMENSION attribute

#### 17 8.5.6.1 General

18 1 The **CODIMENSION attribute** specifies that an entity is a **coarray**. The *coarray-spec* specifies its **corank** or  
 19 **corank** and **cobounds**.

20 R809 *coarray-spec* is *deferred-coshape-spec-list*  
 21 or *explicit-coshape-spec*

22 C821 The sum of the **rank** and **corank** of an entity shall not exceed fifteen.

23 C822 A **coarray** shall be a component or a variable that is not a function result.

24 C823 A **coarray** shall not be of type **C\_PTR** or **C\_FUNPTR** from the intrinsic module **ISO\_C\_BINDING**  
 25 (18.3.3).

26 C824 An entity whose type has a **coarray ultimate component** shall be a nonpointer nonallocatable scalar, shall  
 27 not be a **coarray**, and shall not be a function result.

28 C825 A **coarray** or an object with a **coarray ultimate component** shall be an **associate name**, a **dummy argument**,

1 or have the [ALLOCATABLE](#) or [SAVE](#) attribute.

#### NOTE 8.7

A [coarray](#) is permitted to be of a derived type with pointer or [allocatable](#) components. The [target](#) of such a pointer component is always on the same [image](#) as the pointer.

#### NOTE 8.8

This requirement for the [SAVE attribute](#) has the effect that [automatic coarrays](#) are not permitted; for example, the [coarray](#) WORK in the following code fragment is not valid.

```
SUBROUTINE SOLVE3(N,A,B)
INTEGER :: N
REAL    :: A(N) [*], B(N)
REAL    :: WORK(N) [*]    ! Not permitted
```

If this were permitted, it would require an implicit synchronization on entry to the procedure.

[Explicit-shape coarrays](#) that are declared in a subprogram and are not [dummy arguments](#) are required to have the [SAVE attribute](#) because otherwise they might be implemented as if they were [automatic coarrays](#).

#### NOTE 8.9

Examples of CODIMENSION attribute specifications are:

```
REAL W(100,100) [0:2,*]           ! Explicit-shape coarray
REAL, CODIMENSION[*] :: X         ! Scalar coarray
REAL, CODIMENSION[3,*] :: Y(:)   ! Assumed-shape coarray
REAL, CODIMENSION[:,],ALLOCATABLE :: Z(:,,:) ! Allocatable coarray
```

### 2 8.5.6.2 Allocatable coarray

3 1 A [coarray](#) with the [ALLOCATABLE attribute](#) has a specified [corank](#), but its [cobounds](#) are determined by  
4 allocation or [argument association](#).

5 R810 *deferred-coshape-spec* is :

6 C826 A [coarray](#) with the [ALLOCATABLE attribute](#) shall have a *coarray-spec* that is a *deferred-coshape-spec-*  
7 *list*.

8 2 The [corank](#) of an [allocatable coarray](#) is equal to the number of colons in its *deferred-coshape-spec-list*.

9 3 The [cobounds](#) of an unallocated [allocatable coarray](#) are undefined. No part of such a [coarray](#) shall be referenced  
10 or defined; however, the [coarray](#) may appear as an argument to an intrinsic [inquiry function](#) as specified in 16.1.

11 4 The [cobounds](#) of an allocated [allocatable coarray](#) are those specified when the [coarray](#) is allocated.

12 5 The [cobounds](#) of an [allocatable coarray](#) are unaffected by any subsequent redefinition or undefinition of the  
13 variables on which the [cobounds](#)' expressions depend.

### 14 8.5.6.3 Explicit-coshape coarray

15 1 An explicit-coshape coarray is a named [coarray](#) that has its [corank](#) and [cobounds](#) declared by an *explicit-coshape-*  
16 *spec*.

17 R811 *explicit-coshape-spec* is [ [ *lower-cobound* : ] *upper-cobound*, ]... ■  
18 ■ [ *lower-cobound* : ] \*

19 C827 A nonallocatable [coarray](#) shall have a *coarray-spec* that is an *explicit-coshape-spec*.

- 1 2 The **corank** is equal to one plus the number of *upper-cobounds*.
- 2 R812 *lower-cobound* is *specification-expr*
- 3 R813 *upper-cobound* is *specification-expr*
- 4 C828 (R811) A *lower-cobound* or *upper-cobound* that is not a **constant expression** shall appear only in a sub-  
5 program, **BLOCK construct**, or **interface body**.
- 6 3 If an explicit-coshape **coarray** is a **local variable** of a subprogram or **BLOCK construct** and has **cobounds** that are  
7 not **constant expressions**, the **cobounds** are determined on entry to a procedure defined by the subprogram, or  
8 on execution of the **BLOCK statement**, by evaluating the **cobounds** expressions. The **cobounds** of such a **coarray**  
9 are unaffected by the redefinition or undefinition of any variable during execution of the procedure or **BLOCK**  
10 **construct**.
- 11 4 The values of each *lower-cobound* and *upper-cobound* determine the **cobounds** of the **coarray** along a particular  
12 **codimension**. The **cosubscript** range of the **coarray** in that **codimension** is the set of integer values between and  
13 including the lower and upper **cobounds**. If the lower **cobound** is omitted, the default value is 1. The upper  
14 **cobound** shall not be less than the lower **cobound**.

### 15 8.5.7 CONTIGUOUS attribute

16 C829 An entity with the **CONTIGUOUS** attribute shall be an **array pointer**, an **assumed-shape array**, or an  
17 **assumed-rank dummy data object**.

- 18 1 The **CONTIGUOUS** attribute specifies that an **assumed-shape array** is contiguous, that an **array pointer** can  
19 only be **pointer associated** with a **contiguous target**, or that an **assumed-rank dummy data object** is contiguous.
- 20 2 An object is **contiguous** if it is
- 21 (1) an object with the **CONTIGUOUS** attribute,  
22 (2) a nonpointer whole array that is not **assumed-shape**,  
23 (3) an **assumed-shape array** that is argument associated with an array that is contiguous,  
24 (4) an **assumed-rank dummy data object** whose **effective argument** is contiguous,  
25 (5) an array allocated by an **ALLOCATE** statement,  
26 (6) a **pointer associated** with a contiguous **target**, or  
27 (7) a nonzero-sized **array section** (9.5.3) provided that
- 28 (a) its **base object** is **contiguous**,  
29 (b) it does not have a **vector subscript**,  
30 (c) the array element ordering of the elements of the section is the same as the array element  
31 ordering of those elements of the **base object**,  
32 (d) in the array element ordering of the **base object**, every element of the **base object** that is not  
33 an element of the section either precedes every element of the section or follows every element  
34 of the section,  
35 (e) if the array is of type character and a *substring-range* appears, the *substring-range* specifies all  
36 of the characters of the *parent-string* (9.4.1),  
37 (f) only its final *part-ref* has nonzero **rank**, and  
38 (g) it is not the real or imaginary part (9.4.4) of an array of type complex.
- 39 3 An object is not contiguous if it is an array subobject, and
- 40 • the object has two or more elements,  
41 • the elements of the object in array element order are not consecutive in the elements of the **base object**,  
42 • the object is not of type character with length zero, and  
43 • the object is not of a derived type that has no **ultimate components** other than zero-sized arrays and  
44 characters with length zero.

- 1 4 It is processor dependent whether any other object is contiguous.

**NOTE 8.10**

If a derived type has only one component that is not zero-sized, it is processor dependent whether a structure component of a contiguous array of that type is contiguous. That is, the derived type might contain padding on some processors.

**NOTE 8.11**

The **CONTIGUOUS attribute** makes it easier for a processor to enable optimizations that depend on the memory layout of the object occupying a contiguous block of memory. Examples of **CONTIGUOUS attribute** specifications are:

```
REAL, POINTER, CONTIGUOUS      :: SPTR(:)
REAL, CONTIGUOUS, DIMENSION(:, :) :: D
```

2 **8.5.8 DIMENSION attribute**

3 **8.5.8.1 General**

- 4 1 The **DIMENSION attribute** specifies that an entity is **assumed-rank** or an array. An **assumed-rank dummy data**  
5 **object** has the rank, shape, and size of its **effective argument**; otherwise, the **rank** or **rank** and shape is specified  
6 by its **array-spec**.

7	R814	<i>dimension-spec</i>	is	DIMENSION ( <i>array-spec</i> )
8	R815	<i>array-spec</i>	is	<i>explicit-shape-spec-list</i>
9			or	<i>assumed-shape-spec-list</i>
10			or	<i>deferred-shape-spec-list</i>
11			or	<i>assumed-size-spec</i>
12			or	<i>implied-shape-spec</i>
13			or	<i>implied-shape-or-assumed-size-spec</i>
14			or	<i>assumed-rank-spec</i>

**NOTE 8.12**

The maximum **rank** of an entity is fifteen minus the **corank**.

**NOTE 8.13**

Examples of **DIMENSION attribute** specifications are:

```
SUBROUTINE EX (N, A, B)
  REAL, DIMENSION (N, 10) :: W      ! Automatic explicit-shape array
  REAL A (:), B (0:)              ! Assumed-shape arrays
  REAL, POINTER :: D (:, :)       ! Array pointer
  REAL, DIMENSION (:), POINTER :: P ! Array pointer
  REAL, ALLOCATABLE, DIMENSION (:) :: E ! Allocatable array
  REAL, PARAMETER :: V(0:*) = [0.1, 1.1] ! Implied-shape array
```

15 **8.5.8.2 Explicit-shape array**

16	R816	<i>explicit-shape-spec</i>	is	[ <i>lower-bound</i> : ] <i>upper-bound</i>
17	R817	<i>lower-bound</i>	is	<i>specification-expr</i>
18	R818	<i>upper-bound</i>	is	<i>specification-expr</i>

1 C830 (R816) An *explicit-shape-spec* whose **bounds** are not **constant expressions** shall appear only in a subpro-  
 2 gram, derived type definition, **BLOCK construct**, or **interface body**.

3 1 An **explicit-shape array** is an array whose shape is explicitly declared by an *explicit-shape-spec-list*. The **rank** is  
 4 equal to the number of *explicit-shape-specs*.

5 2 An **explicit-shape array** that is a named **local variable** of a subprogram or **BLOCK construct** may have **bounds**  
 6 that are not **constant expressions**. The **bounds**, and hence shape, are determined on entry to a procedure defined  
 7 by the subprogram, or on execution of the **BLOCK statement**, by evaluating the bounds' expressions. The  
 8 **bounds** of such an array are unaffected by the redefinition or undefinition of any variable during execution of the  
 9 procedure or **BLOCK construct**.

10 3 The values of each *lower-bound* and *upper-bound* determine the bounds of the array along a particular dimension  
 11 and hence the extent of the array in that dimension. If *lower-bound* appears it specifies the lower bound; otherwise  
 12 the lower bound is 1. The value of a lower bound or an upper bound may be positive, negative, or zero. The  
 13 subscript range of the array in that dimension is the set of integer values between and including the lower and  
 14 upper bounds, provided the upper bound is not less than the lower bound. If the upper bound is less than the  
 15 lower bound, the range is empty, the extent in that dimension is zero, and the array is of zero size.

### 16 8.5.8.3 Assumed-shape array

17 1 An **assumed-shape array** is a nonallocatable nonpointer **dummy argument** array that takes its shape from its  
 18 **effective argument**.

19 R819 *assumed-shape-spec* is [ *lower-bound* ] :

20 2 The **rank** is equal to the number of colons in the *assumed-shape-spec-list*.

21 3 The extent of a dimension of an **assumed-shape array dummy argument** is the extent of the corresponding  
 22 dimension of its **effective argument**. If the lower bound value is  $d$  and the extent of the corresponding dimension  
 23 of its **effective argument** is  $s$ , then the value of the upper bound is  $s + d - 1$ . If *lower-bound* appears it specifies  
 24 the lower bound; otherwise the lower bound is 1.

### 25 8.5.8.4 Deferred-shape array

26 1 A **deferred-shape array** is an **allocatable** array or an **array pointer**. (An **allocatable** array has the **ALLOCATABLE**  
 27 **attribute**; an **array pointer** has the **POINTER attribute**.)

28 R820 *deferred-shape-spec* is :

29 C831 An array with the **POINTER** or **ALLOCATABLE attribute** shall have an *array-spec* that is a *deferred-*  
 30 *shape-spec-list*.

31 2 The **rank** is equal to the number of colons in the *deferred-shape-spec-list*.

32 3 The size, bounds, and shape of an unallocated **allocatable** array or a **disassociated array pointer** are undefined.  
 33 No part of such an array shall be referenced or defined; however, the array may appear as an argument to an  
 34 intrinsic **inquiry function** as specified in 16.1.

35 4 The bounds of each dimension of an allocated **allocatable** array are those specified when the array is allocated  
 36 or, if it is a **dummy argument**, when it is argument associated with an allocated **effective argument**.

37 5 The bounds of each dimension of an associated **array pointer**, and hence its shape, may be specified

- 38 • in an **ALLOCATE statement** (9.7.1) when the **target** is allocated,
- 39 • by pointer assignment (10.2.2), or
- 40 • if it is a **dummy argument**, by **argument association** with a nonpointer **actual argument** or an associated  
 41 pointer **effective argument**.

1 6 The bounds of an [array pointer](#) or [allocatable](#) array are unaffected by any subsequent redefinition or undefinition  
2 of variables on which the bounds' expressions depend.

### 3 8.5.8.5 Assumed-size array

4 1 An [assumed-size array](#) is a [dummy argument](#) array whose size is assumed from that of its [effective argument](#), or  
5 the associate name of a [RANK \( \\* \)](#) block in a [SELECT RANK construct](#). The [rank](#) and extents may differ for  
6 the [effective](#) and [dummy](#) arguments; only the size of the [effective argument](#) is assumed by the [dummy argument](#).  
7 A [dummy argument](#) is declared to be an [assumed-size array](#) by an [assumed-size-spec](#) or an [implied-shape-or-](#)  
8 [assumed-size-spec](#).

9 R821 *assumed-implied-spec* is [ *lower-bound* : ] \*

10 R822 *assumed-size-spec* is *explicit-shape-spec-list*, *assumed-implied-spec*

11 C832 An object whose array bounds are specified by an [assumed-size-spec](#) shall be a [dummy data object](#).

12 C833 An [assumed-size array](#) with the [INTENT \(OUT\)](#) attribute shall not be [polymorphic](#), [finalizable](#), of a  
13 type with an [allocatable ultimate component](#), or of a type for which [default initialization](#) is specified.

14 R823 *implied-shape-or-assumed-size-spec* is *assumed-implied-spec*

15 C834 An object whose array bounds are specified by an [implied-shape-or-assumed-size-spec](#) shall be a [dummy](#)  
16 [data object](#) or a [named constant](#).

17 2 The size of an [assumed-size array](#) is determined as follows.

- 18 • If the [effective argument](#) associated with the [assumed-size](#) dummy array is an array of any type other than  
19 default character, the size is that of the [effective argument](#).
- 20 • If the [actual argument](#) corresponding to the [assumed-size](#) dummy array is an array element of any type  
21 other than default character with a subscript order value of  $r$  (9.5.3.2) in an array of size  $x$ , the size of the  
22 dummy array is  $x - r + 1$ .
- 23 • If the [actual argument](#) is a default character array, default character array element, or a default character  
24 array element substring (9.4.1), and if it begins at [character storage unit](#)  $t$  of an array with  $c$  [character](#)  
25 [storage units](#), the size of the dummy array is  $\text{MAX}(\text{INT}((c - t + 1)/e), 0)$ , where  $e$  is the length of an  
26 element in the dummy character array.
- 27 • If the [actual argument](#) is a default character scalar that is not an array element or array element substring  
28 [designator](#), the size of the dummy array is  $\text{MAX}(\text{INT}(l/e), 0)$ , where  $e$  is the length of an element in the  
29 dummy character array and  $l$  is the length of the [actual argument](#).

30 3 The [rank](#) is equal to one plus the number of [explicit-shape-specs](#).

31 4 An [assumed-size array](#) has no upper bound in its last dimension and therefore has no extent in its last dimension  
32 and no shape. An [assumed-size array](#) shall not appear in a context that requires its shape.

33 5 If a list of [explicit-shape-specs](#) appears, it specifies the bounds of the first [rank](#)–1 dimensions. If [lower-bound](#)  
34 appears it specifies the lower bound of the last dimension; otherwise that lower bound is 1. An [assumed-size](#)  
35 [array](#) may be subscripted or sectioned (9.5.3.3). The upper bound shall not be omitted from a subscript triplet  
36 in the last dimension.

37 6 If an [assumed-size array](#) has bounds that are not [constant expressions](#), the bounds are determined on entry to  
38 the procedure. The bounds of such an array are unaffected by the redefinition or undefinition of any variable  
39 during execution of the procedure.

### 40 8.5.8.6 Implied-shape array

41 1 An implied-shape array is a [named constant](#) that takes its shape from the [constant-expr](#) in its declaration. A  
42 named constant is declared to be an implied-shape array with an [array-spec](#) that is an [implied-shape-or-assumed-](#)

1 *size-spec* or an *implied-shape-spec*.

2 R824 *implied-shape-spec* is *assumed-implied-spec*, *assumed-implied-spec-list*

3 C835 An implied-shape array shall be a **named constant**.

4 2 The **rank** of an implied-shape array is the number of *assumed-implied-specs* in its *array-spec*.

5 3 The extent of each dimension of an implied-shape array is the same as the extent of the corresponding dimension  
6 of the *constant-expr*. The lower bound of each dimension is *lower-bound*, if it appears, and 1 otherwise; the upper  
7 bound is one less than the sum of the lower bound and the extent.

### 8 8.5.8.7 Assumed-rank entity

9 1 An assumed-rank entity is a **dummy data object** whose **rank** is assumed from its **effective argument**, or the associate  
10 name of a **RANK DEFAULT** block in a **SELECT RANK construct**; this rank can be zero. An assumed-rank  
11 entity is declared with an *array-spec* that is an *assumed-rank-spec*.

12 R825 *assumed-rank-spec* is ..

13 C836 An assumed-rank entity shall be a **dummy data object** that does not have the **CODIMENSION** or **VALUE**  
14 attribute.

15 C837 An assumed-rank variable name shall not appear in a **designator** or expression except as an **actual**  
16 **argument** that corresponds to a **dummy argument** that is assumed-rank, the argument of the function  
17 **C\_LOC** or **C\_SIZEOF** from the intrinsic module **ISO\_C\_BINDING** (18.2), the first **dummy argument**  
18 of an **intrinsic inquiry function**, or the selector of a **SELECT RANK statement**.

19 C838 If an **assumed-size** or nonallocatable nonpointer assumed-rank array is an **actual argument** that corres-  
20 ponds to a **dummy argument** that is an **INTENT (OUT)** assumed-rank array, it shall not be **polymorphic**,  
21 **finalizable**, of a type with an **allocatable ultimate component**, or of a type for which **default initialization**  
22 is specified.

### 23 8.5.9 EXTERNAL attribute

24 1 The **EXTERNAL attribute** specifies that an entity is an **external procedure**, **dummy procedure**, **procedure pointer**,  
25 or block data subprogram.

26 C839 An entity shall not have both the **EXTERNAL attribute** and the **INTRINSIC attribute**.

27 C840 In an external subprogram, the **EXTERNAL attribute** shall not be specified for a procedure defined by  
28 the subprogram.

29 C841 In an interface body, the **EXTERNAL attribute** shall not be specified for the procedure declared by the  
30 interface body.

31 2 If an **external procedure** or **dummy procedure** is used as an **actual argument** or is the **target** of a **procedure pointer**  
32 **assignment**, it shall be declared to have the **EXTERNAL attribute**.

33 3 A procedure that has both the **EXTERNAL** and **POINTER** attributes is a **procedure pointer**.

#### NOTE 8.14

The **EXTERNAL attribute** can be specified in a **type declaration statement**, by an **interface body** (15.4.3.2),  
by an **EXTERNAL statement** (15.4.3.5), or by a **procedure declaration statement** (15.4.3.6).

### 34 8.5.10 INTENT attribute

35 1 The **INTENT attribute** specifies the intended use of a **dummy argument**. An **INTENT (IN)** **dummy argument**  
36 is suitable for receiving data from the invoking **scoping unit**, an **INTENT (OUT)** **dummy argument** is suitable

1 for returning data to the invoking [scoping unit](#), and an INTENT (INOUT) [dummy argument](#) is suitable for use  
2 both to receive data from and to return data to the invoking [scoping unit](#).

3 R826 *intent-spec* is IN  
4 or OUT  
5 or INOUT

6 C842 An entity with the [INTENT attribute](#) shall be a [dummy data object](#) or a dummy procedure pointer.

7 C843 (R826) A nonpointer object with the [INTENT \(IN\) attribute](#) shall not appear in a variable definition  
8 context ([19.6.7](#)).

9 C844 A pointer with the [INTENT \(IN\) attribute](#) shall not appear in a pointer association context ([19.6.8](#)).

10 C845 An [INTENT \(OUT\) dummy argument](#) of a nonintrinsic procedure shall not be an [allocatable coarray](#) or  
11 have a [subobject](#) that is an [allocatable coarray](#).

12 C846 An entity with the [INTENT \(OUT\) attribute](#) shall not be of, or have a [subcomponent](#) of, type [EVENT\\_-](#)  
13 [TYPE](#) ([16.10.2.10](#)) or [LOCK\\_TYPE](#) ([16.10.2.19](#)) from the intrinsic module [ISO\\_FORTRAN\\_ENV](#).

14 2 The [INTENT \(IN\) attribute](#) for a nonpointer [dummy argument](#) specifies that it shall neither be defined nor  
15 become undefined during the invocation and execution of the procedure. The [INTENT \(IN\) attribute](#) for a  
16 pointer [dummy argument](#) specifies that during the invocation and execution of the procedure its association shall  
17 not be changed except that it may become undefined if the [target](#) is deallocated other than through the pointer  
18 ([19.5.2.5](#)).

19 3 The [INTENT \(OUT\) attribute](#) for a nonpointer [dummy argument](#) specifies that the [dummy argument](#) becomes  
20 undefined on invocation of the procedure, except for any subcomponents that are [default-initialized](#) ([7.5.4.6](#)). Any  
21 [actual argument](#) that corresponds to such a [dummy argument](#) shall be [definable](#). The [INTENT \(OUT\) attribute](#)  
22 for a pointer [dummy argument](#) specifies that on invocation of the procedure the [pointer association](#) status of  
23 the [dummy argument](#) becomes undefined. Any [actual argument](#) that corresponds to such a pointer dummy shall  
24 be a pointer variable or a [procedure pointer](#) that is not the result of a function reference. Any undefinition or  
25 definition implied by association of an [actual argument](#) with an [INTENT \(OUT\) dummy argument](#) shall not  
26 affect any other entity within the statement that invokes the procedure.

27 4 The [INTENT \(INOUT\) attribute](#) for a nonpointer [dummy argument](#) specifies that any [actual argument](#) that  
28 corresponds to the [dummy argument](#) shall be [definable](#). The [INTENT \(INOUT\) attribute](#) for a pointer [dummy](#)  
29 [argument](#) specifies that any [actual argument](#) that corresponds to the [dummy argument](#) shall be a pointer variable  
30 or a [procedure pointer](#) that is not the result of a function reference.

#### NOTE 8.15

The [INTENT attribute](#) for an [allocatable dummy argument](#) applies to both the allocation status and the definition status. An [actual argument](#) that corresponds to an [INTENT \(OUT\) allocatable dummy argument](#) is deallocated on procedure invocation ([9.7.3.2](#)). To avoid this deallocation for coarrays, [INTENT \(OUT\)](#) is not allowed for a [dummy argument](#) that is an [allocatable coarray](#) or has a [subobject](#) that is an [allocatable coarray](#).

31 5 If no [INTENT attribute](#) is specified for a [dummy argument](#), its use is subject to the limitations of its [effective](#)  
32 [argument](#) ([15.5.2](#)).

33 6 If an object has an [INTENT attribute](#), then all of its subobjects have the same [INTENT attribute](#).

#### NOTE 8.16

An example of INTENT specification is:

```
SUBROUTINE MOVE (FROM, TO)
  USE PERSON_MODULE
```



**NOTE 8.16 (cont.)**

TYPE (PERSON), INTENT (IN) :: FROM TYPE (PERSON), INTENT (OUT) :: TO
---

**NOTE 8.17**

If a [dummy argument](#) is a derived-type object with a pointer component, then the pointer as a pointer is a subobject of the [dummy argument](#), but the [target](#) of the pointer is not. Therefore, the restrictions on subobjects of the [dummy argument](#) apply to the pointer in contexts where it is used as a pointer, but not in contexts where it is dereferenced to indicate its [target](#). For example, if X is a [dummy argument](#) of derived type with an integer pointer component P, and X is [INTENT \(IN\)](#), then the statement

```
X%P => NEW_TARGET
```

is prohibited, but

```
X%P = 0
```

is allowed (provided that X%P is associated with a [definable target](#)).

Similarly, the INTENT restrictions on pointer [dummy arguments](#) apply only to the association of the [dummy argument](#); they do not restrict the operations allowed on its [target](#).

**NOTE 8.18**

Argument intent specifications serve several purposes in addition to documenting the intended use of [dummy arguments](#). A processor can check whether an [INTENT \(IN\) dummy argument](#) is used in a way that could redefine it. A slightly more sophisticated processor could check to see whether an [INTENT \(OUT\) dummy argument](#) could possibly be referenced before it is defined. If the procedure's [interface](#) is explicit, the processor can also verify that [actual arguments](#) corresponding to [INTENT \(OUT\)](#) or [INTENT \(INOUT\) dummy arguments](#) are [definable](#). A more sophisticated processor could use this information to optimize the translation of the referencing [scoping unit](#) by taking advantage of the fact that [actual arguments](#) corresponding to [INTENT \(IN\) dummy arguments](#) will not be changed and that any prior value of an [actual argument](#) corresponding to an [INTENT \(OUT\) dummy argument](#) will not be referenced and could thus be discarded.

[INTENT \(OUT\)](#) means that the value of the argument after invoking the procedure is entirely the result of executing that procedure. If an argument might not be redefined and it is desired to have the argument retain its value in that case, [INTENT \(OUT\)](#) cannot be used because it would cause the argument to become undefined; however, [INTENT \(INOUT\)](#) can be used, even if there is no explicit reference to the value of the dummy argument.

[INTENT \(INOUT\)](#) is not equivalent to omitting the [INTENT attribute](#). The [actual argument](#) corresponding to an [INTENT \(INOUT\) dummy argument](#) is always required to be [definable](#), while an [actual argument](#) corresponding to a [dummy argument](#) without an [INTENT attribute](#) need be [definable](#) only if the [dummy argument](#) is actually redefined.

## 1 **8.5.11 INTRINSIC attribute**

2 1 The [INTRINSIC attribute](#) specifies that the entity is an intrinsic procedure. The procedure name may be a  
3 generic name (16.7), a [specific name](#) (16.8), or both.

4 2 If the [specific name](#) of an intrinsic procedure (16.8) is used as an [actual argument](#), the name shall be explicitly specified to have the  
5 [INTRINSIC attribute](#). Note that a specific intrinsic procedure listed in Table 16.3 is not permitted to be used as an [actual argument](#)  
6 (C1534).

7 C847 If the generic name of an intrinsic procedure is explicitly declared to have the [INTRINSIC attribute](#),

1 and it is also the generic name of one or more [generic interfaces \(15.4.3.2\)](#) accessible in the same [scoping](#)  
 2 [unit](#), the procedures in the interfaces and the generic intrinsic procedure shall all be functions or all be  
 3 subroutines.

#### 4 **8.5.12 OPTIONAL attribute**

5 1 The [OPTIONAL attribute](#) specifies that the [dummy argument](#) need not have a corresponding [actual argument](#)  
 6 in a [reference](#) to the procedure ([15.5.2.12](#)).

7 C848 An entity with the [OPTIONAL attribute](#) shall be a [dummy argument](#).

##### NOTE 8.19

The intrinsic function [PRESENT \(16.9.152\)](#) can be used to determine whether an optional [dummy argument](#) has a corresponding [actual argument](#).

#### 8 **8.5.13 PARAMETER attribute**

9 1 The [PARAMETER attribute](#) specifies that an entity is a [named constant](#). The entity has the value specified by  
 10 its [constant-expr](#), converted, if necessary, to the type, type parameters and shape of the entity.

11 C849 An entity with the [PARAMETER attribute](#) shall not be a [variable](#), a [coarray](#), or a [procedure](#).

12 C850 An expression that specifies a [length type parameter](#) or array bound of a named constant shall be a  
 13 constant expression.

14 2 A [named constant](#) shall not be referenced unless it has been defined previously; it may be defined previously in  
 15 the same statement.

##### NOTE 8.20

Examples of declarations with a [PARAMETER attribute](#) are:

```
REAL, PARAMETER :: ONE = 1.0, Y = 4.1 / 3.0
INTEGER, DIMENSION (3), PARAMETER :: ORDER = (/ 1, 2, 3 /)
TYPE(NODE), PARAMETER :: DEFAULT = NODE(0, NULL ( ))
```

#### 16 **8.5.14 POINTER attribute**

17 1 Entities with the [POINTER attribute](#) can be associated with different data objects or procedures during execution  
 18 of a program. A pointer is either a data pointer or a procedure pointer.

19 C851 An entity with the [POINTER attribute](#) shall not have the [ALLOCATABLE](#), [INTRINSIC](#), or [TARGET](#)  
 20 attribute, and shall not be a [coarray](#).

21 C852 A procedure with the [POINTER attribute](#) shall have the [EXTERNAL attribute](#).

22 2 A data pointer shall not be referenced unless it is [pointer associated](#) with a [target](#) object that is defined. A data  
 23 pointer shall not be defined unless it is [pointer associated](#) with a [target](#) object that is [definable](#).

24 3 If a data pointer is associated, the values of its [deferred type parameters](#) are the same as the values of the  
 25 corresponding type parameters of its [target](#).

26 4 A [procedure pointer](#) shall not be referenced unless it is [pointer associated](#) with a [target](#) procedure.

##### NOTE 8.21

Examples of [POINTER attribute](#) specifications are:

```
TYPE (NODE), POINTER :: CURRENT, TAIL
```

**NOTE 8.21 (cont.)**

REAL, DIMENSION (:, :), POINTER :: IN, OUT, SWAP
--

**8.5.15 PROTECTED attribute**

1 The **PROTECTED attribute** imposes limitations on the usage of module entities.

C853 The **PROTECTED attribute** shall be specified only in the specification part of a module.

C854 An entity with the **PROTECTED attribute** shall be a procedure pointer or variable.

C855 An entity with the **PROTECTED attribute** shall not be in a **common block**.

C856 A nonpointer object that has the **PROTECTED attribute** and is accessed by **use association** shall not appear in a variable definition context (19.6.7) or as a *data-target* or *initial-data-target*.

C857 A pointer that has the **PROTECTED attribute** and is accessed by **use association** shall not appear in a pointer association context (19.6.8).

2 Other than within the module in which an entity is given the **PROTECTED attribute**, or within any of its **descendants**,

- if it is a nonpointer object, it is not **definable**, and
- if it is a pointer, its association status shall not be changed except that it may become undefined if its **target** is deallocated other than through the pointer (19.5.2.5), or if its **target** becomes undefined by completing execution of a **BLOCK construct** or by execution of a **RETURN** or **END** statement.

3 If an object has the **PROTECTED attribute**, all of its subobjects have the **PROTECTED attribute**.

**NOTE 8.22**

An example of the **PROTECTED attribute**:

```

MODULE temperature
  REAL, PROTECTED :: temp_c, temp_f
CONTAINS
  SUBROUTINE set_temperature_c(c)
    REAL, INTENT(IN) :: c
    temp_c = c
    temp_f = temp_c*(9.0/5.0) + 32
  END SUBROUTINE
END MODULE

```

The **PROTECTED attribute** ensures that the variables `temp_c` and `temp_f` cannot be modified other than via the `set_temperature_c` procedure, thus keeping them consistent with each other.

**8.5.16 SAVE attribute**

1 The **SAVE attribute** specifies that a **local variable** of a **program unit** or subprogram retains its association status, allocation status, definition status, and value after execution of a **RETURN** or **END** statement unless it is a pointer and its **target** becomes undefined (19.5.2.5(6)). If it is a **local variable** of a subprogram it is shared by all instances (15.6.2.4) of the subprogram.

2 The **SAVE attribute** specifies that a **local variable** of a **BLOCK construct** retains its association status, allocation status, definition status, and value after termination of the construct unless it is a pointer and its **target** becomes undefined (19.5.2.5(7)). If the **BLOCK construct** is within a subprogram the variable is shared by all instances (15.6.2.4) of the subprogram.

1 3 Giving a **common block** the **SAVE attribute** confers the attribute on all entities in the **common block**.

2 C858 An entity with the **SAVE attribute** shall be a **common block**, variable, or procedure pointer.

3 C859 The **SAVE attribute** shall not be specified for a **dummy argument**, a function result, an **automatic data**  
4 **object**, or an object that is in a **common block**.

5 4 A variable, **common block**, or procedure pointer declared in the **scoping unit** of a main program, module, or  
6 submodule implicitly has the **SAVE attribute**, which may be confirmed by explicit specification. If a **common block**  
7 has the **SAVE attribute** in any other kind of **scoping unit**, it shall have the **SAVE attribute** in every **scoping unit** that is not of a  
8 main program, module, or submodule.

### 9 8.5.17 TARGET attribute

10 1 The **TARGET attribute** specifies that a data object may have a **pointer associated** with it (10.2.2). An object  
11 without the **TARGET attribute** shall not have a **pointer associated** with it.

12 C860 An entity with the **TARGET attribute** shall be a variable.

13 C861 An entity with the **TARGET attribute** shall not have the **POINTER attribute**.

14 2 If an object has the **TARGET attribute**, then all of its nonpointer subobjects also have the **TARGET attribute**.

#### NOTE 8.23

In addition to variables explicitly declared to have the **TARGET attribute**, the objects created by allocation of pointers (9.7.1.4) have the **TARGET attribute**.

#### NOTE 8.24

Examples of **TARGET attribute** specifications are:

```
TYPE (NODE), TARGET :: HEAD
REAL, DIMENSION (1000, 1000), TARGET :: A, B
```

#### NOTE 8.25

Every **object designator** that starts from an object with the **TARGET attribute** will have either the **TARGET** or **POINTER** attribute. If pointers are involved, the **designator** might not necessarily be a subobject of the original object, but because a pointer can point only to an entity with the **TARGET attribute**, there is no way to end up at a nonpointer that does not have the **TARGET attribute**.

### 15 8.5.18 VALUE attribute

16 1 The **VALUE attribute** specifies a type of **argument association** (15.5.2.4) for a **dummy argument**.

17 C862 An entity with the **VALUE attribute** shall be a **dummy data object**. It shall not be an **assumed-size**  
18 **array**, a **coarray**, or a variable with a **coarray ultimate component**.

19 C863 An entity with the **VALUE attribute** shall not have the **ALLOCATABLE**, **INTENT (INOUT)**, **INTENT**  
20 **(OUT)**, **POINTER**, or **VOLATILE** attributes.

21 C864 A **dummy argument** of a procedure with the **BIND attribute** shall not have both the **OPTIONAL** and  
22 **VALUE** attributes.

### 23 8.5.19 VOLATILE attribute

24 1 The **VOLATILE attribute** specifies that an object may be referenced, defined, or become undefined, by means  
25 not specified by the program. A pointer with the **VOLATILE attribute** may additionally have its association

- 1 status, [dynamic type](#) and type parameters, and array bounds changed by means not specified by the program.  
 2 An [allocatable](#) object with the [VOLATILE attribute](#) may additionally have its allocation status, [dynamic type](#)  
 3 and type parameters, and array bounds changed by means not specified by the program.
- 4 C865 An entity with the [VOLATILE attribute](#) shall be a variable that is not an [INTENT \(IN\) dummy argu-](#)  
 5 [ment](#).
- 6 C866 The [VOLATILE attribute](#) shall not be specified for a [coarray](#), or a variable with a [coarray ultimate](#)  
 7 [component](#), that is accessed by [use \(14.2.2\)](#) or [host \(19.5.1.4\)](#) association.
- 8 C867 Within a [BLOCK construct \(11.1.4\)](#), the [VOLATILE attribute](#) shall not be specified for a [coarray](#), or a  
 9 variable with a [coarray ultimate component](#), that is not a [construct entity \(19.4\)](#) of that construct.
- 10 2 A noncoarray object that has the [VOLATILE attribute](#) may be associated with an object that does not have  
 11 the [VOLATILE attribute](#), including by [use \(14.2.2\)](#) or [host](#) association ([19.5.1.4](#)). If an object that is not a  
 12 [local variable](#) of a [BLOCK construct](#) is specified to have the [VOLATILE attribute](#) in the [specification-part](#) of  
 13 the construct, the object has the attribute within the construct even if it does not have the attribute outside the  
 14 construct. The relationship between [coarrays](#), the [VOLATILE attribute](#), and [argument association](#) is described  
 15 in [15.5.2.8](#). The relationship between between [coarrays](#), the [VOLATILE attribute](#), and [pointer association](#) is  
 16 described in [10.2.2.3](#).
- 17 3 A pointer should have the [VOLATILE attribute](#) if its target has the [VOLATILE attribute](#). If, by means not  
 18 specified by the program, the target is referenced, defined, or becomes undefined, the pointer shall have the  
 19 [VOLATILE attribute](#). All members of an EQUIVALENCE group should have the [VOLATILE attribute](#) if any member has the  
 20 [VOLATILE attribute](#).
- 21 4 If an object has the [VOLATILE attribute](#), then all of its subobjects also have the [VOLATILE attribute](#).
- 22 5 The Fortran processor should use the most recent definition of a volatile object each time its value is required.  
 23 When a volatile object is defined by means of Fortran, it should make that definition available to the non-Fortran  
 24 parts of the program as soon as possible.

## 25 8.6 Attribute specification statements

### 26 8.6.1 Accessibility statement

27 R827 *access-stmt* is *access-spec* [ [ *::* ] *access-id-list* ]

28 R828 *access-id* is *access-name*  
 29 or *generic-spec*

30 C868 (R827) An *access-stmt* shall appear only in the *specification-part* of a module. Only one accessibility  
 31 statement with an omitted *access-id-list* is permitted in the *specification-part* of a module.

32 C869 (R828) Each *access-name* shall be the name of a module, variable, procedure, derived type, [named](#)  
 33 [constant](#), or namelist group.

34 C870 A module whose name appears in an *access-stmt* shall be referenced by a [USE statement](#) in the [scoping](#)  
 35 [unit](#) that contains the *access-stmt*.

36 C871 The name of a module shall appear at most once in all of the *access-stmts* in a module.

37 1 An *access-stmt* with an *access-id-list* specifies the [accessibility attribute](#), [PUBLIC](#) or [PRIVATE](#), of each *access-id*  
 38 in the list that is not a module name. An *access-stmt* without an *access-id* list specifies the default accessibility  
 39 of the identifiers of entities declared in the module, and of entities accessed from a module whose name does  
 40 not appear in any *access-stmt* in the module. If an identifier is accessed from another module and also declared  
 41 locally, it has the default accessibility of a locally declared identifier. The statement

42 PUBLIC

- 1 specifies a default of public accessibility. The statement  
 2 PRIVATE  
 3 specifies a default of private accessibility. If no such statement appears in a module, the default is public  
 4 accessibility.
- 5 2 If an identifier is accessed by [use association](#) and not declared in the module, and the name of every module  
 6 from which it is accessed appears in an *access-stmt* in the scoping unit, its default accessibility is PRIVATE if  
 7 the *access-spec* in every such *access-stmt* is PRIVATE, or PUBLIC if the *access-spec* in any such *access-stmt* is  
 8 PUBLIC.

**NOTE 8.26**

Examples of accessibility statements are:

```
MODULE EX
  PRIVATE
  PUBLIC :: A, B, C, ASSIGNMENT (=), OPERATOR (+)
```

**NOTE 8.27**

The following is an example of using an accessibility statement on a module name.

```
MODULE m2
  USE m1
  ! We want to use the types and procedures in m1, but we only want to
  ! re-export m_type from m1, and export our own procedures.
  PRIVATE m1
  PUBLIC m_type
  ... definitions for our own entities and module procedures.
END MODULE
```

9 **8.6.2 ALLOCATABLE statement**

- 10 R829 *allocatable-stmt* is ALLOCATABLE [ :: ] *allocatable-decl-list*
- 11 R830 *allocatable-decl* is *object-name* [ ( *array-spec* ) ] ■  
 12 ■ [ *lbracket coarray-spec rbracket* ]

- 13 1 The ALLOCATABLE statement specifies the [ALLOCATABLE attribute](#) (8.5.3) for a list of objects.

**NOTE 8.28**

An example of an ALLOCATABLE statement is:

```
REAL A, B (:), SCALAR
ALLOCATABLE :: A (:, :), B, SCALAR
```

14 **8.6.3 ASYNCHRONOUS statement**

- 15 R831 *asynchronous-stmt* is ASYNCHRONOUS [ :: ] *object-name-list*

- 16 1 The ASYNCHRONOUS statement specifies the [ASYNCHRONOUS attribute](#) (8.5.4) for a list of objects.

17 **8.6.4 BIND statement**

- 18 R832 *bind-stmt* is *language-binding-spec* [ :: ] *bind-entity-list*
- 19 R833 *bind-entity* is *entity-name*  
 20 or / *common-block-name* /

1 C872 (R832) If the *language-binding-spec* has a **NAME=** specifier, the *bind-entity-list* shall consist of a single  
2 *bind-entity*.

3 1 The BIND statement specifies the **BIND attribute** for a list of variables and **common blocks**.

#### 4 8.6.5 CODIMENSION statement

5 R834 *codimension-stmt* is CODIMENSION [ :: ] *codimension-decl-list*

6 R835 *codimension-decl* is *coarray-name lbracket coarray-spec rbracket*

7 1 The CODIMENSION statement specifies the **CODIMENSION attribute** (8.5.6) for a list of objects.

#### NOTE 8.29

An example of a CODIMENSION statement is:

```
CODIMENSION a[*], b[3,*], c[:]
```

#### 8 8.6.6 CONTIGUOUS statement

9 R836 *contiguous-stmt* is CONTIGUOUS [ :: ] *object-name-list*

10 1 The CONTIGUOUS statement specifies the **CONTIGUOUS attribute** (8.5.7) for a list of objects.

#### 11 8.6.7 DATA statement

12 R837 *data-stmt* is DATA *data-stmt-set* [ [ , ] *data-stmt-set* ] ...

13 1 The DATA statement specifies **explicit initialization** (8.4).

14 2 If a nonpointer variable has **default initialization**, it shall not appear in a *data-stmt-object-list*.

15 3 A variable that appears in a DATA statement and has not been typed previously may appear in a subsequent type  
16 declaration only if that declaration confirms the implicit typing. An array name, **array section**, or array element  
17 that appears in a DATA statement shall have had its array properties established by a previous specification  
18 statement.

19 4 Except for variables in named **common blocks**, a named variable has the **SAVE attribute** if any part of it is initialized  
20 in a DATA statement, and this may be confirmed by explicit specification.

21 R838 *data-stmt-set* is *data-stmt-object-list* / *data-stmt-value-list* /

22 R839 *data-stmt-object* is *variable*  
23 or *data-implied-do*

24 R840 *data-implied-do* is ( *data-i-do-object-list* , [ *integer-type-spec* :: ] *data-i-do-variable* = ■  
25 ■ *scalar-int-constant-expr* , ■  
26 ■ *scalar-int-constant-expr* ■  
27 ■ [ , *scalar-int-constant-expr* ] )

28 R841 *data-i-do-object* is *array-element*  
29 or *scalar-structure-component*  
30 or *data-implied-do*

31 R842 *data-i-do-variable* is *do-variable*

32 C873 A *data-stmt-object* or *data-i-do-object* shall not be a **coindexed** variable.

- 1 C874 (R839) A *data-stmt-object* that is a *variable* shall be a *designator*. Each subscript, section subscript,  
2 substring starting point, and substring ending point in the variable shall be a *constant expression*.
- 3 C875 (R839) A variable whose *designator* appears as a *data-stmt-object* or a *data-i-do-object* shall not be a  
4 *dummy argument*, accessed by use or *host* association, in a named *common block* unless the DATA statement is  
5 in a block data program unit, in *blank common*, a function name, a function result name, an *automatic data*  
6 *object*, or an *allocatable variable*.
- 7 C876 (R839) A *data-i-do-object* or a *variable* that appears as a *data-stmt-object* shall not be an *object designator*  
8 in which a pointer appears other than as the entire rightmost *part-ref*.
- 9 C877 (R841) The *array-element* shall be a variable.
- 10 C878 (R841) The *scalar-structure-component* shall be a variable.
- 11 C879 (R841) The *scalar-structure-component* shall contain at least one *part-ref* that contains a *subscript-list*.
- 12 C880 (R841) In an *array-element* or *scalar-structure-component* that is a *data-i-do-object*, any subscript shall  
13 be a *constant expression*, and any primary within that subscript that is a *data-i-do-variable* shall be a  
14 DO variable of this *data-implied-do* or of a containing *data-implied-do*.
- 15 R843 *data-stmt-value*                    **is** [ *data-stmt-repeat* \* ] *data-stmt-constant*
- 16 R844 *data-stmt-repeat*                 **is** *scalar-int-constant*  
17    **or** *scalar-int-constant-subobject*
- 18 C881 (R844) The *data-stmt-repeat* shall be positive or zero. If the *data-stmt-repeat* is a *named constant*, it  
19 shall have been defined previously.
- 20 R845 *data-stmt-constant*               **is** *scalar-constant*  
21    **or** *scalar-constant-subobject*  
22    **or** *signed-int-literal-constant*  
23    **or** *signed-real-literal-constant*  
24    **or** *null-init*  
25    **or** *initial-data-target*  
26    **or** *structure-constructor*
- 27 C882 (R845) If a DATA statement constant value is a *named constant* or a *structure constructor*, the *named*  
28 *constant* or derived type shall have been defined previously.
- 29 C883 (R845) If a *data-stmt-constant* is a *structure-constructor*, it shall be a *constant expression*.
- 30 R846 *int-constant-subobject*           **is** *constant-subobject*
- 31 C884 (R846) *int-constant-subobject* shall be of type integer.
- 32 R847 *constant-subobject*               **is** *designator*
- 33 C885 (R847) *constant-subobject* shall be a subobject of a constant.
- 34 C886 (R847) Any subscript, substring starting point, or substring ending point shall be a *constant expression*.
- 35 5 The *data-stmt-object-list* is expanded to form a sequence of pointers and scalar variables, referred to as “sequence  
36 of variables” in subsequent text. A nonpointer array whose unqualified name appears as a *data-stmt-object* or  
37 *data-i-do-object* is equivalent to a complete sequence of its array elements in array element order (9.5.3.2). An  
38 *array section* is equivalent to the sequence of its array elements in array element order. A *data-implied-do* is  
39 expanded to form a sequence of array elements and *structure components*, under the control of the *data-i-do-*  
40 *variable*, as in the DO construct (11.1.7.4).
- 41 6 The *data-stmt-value-list* is expanded to form a sequence of *data-stmt-constants*. A *data-stmt-repeat* indicates the



- 1 number of times the following *data-stmt-constant* is to be included in the sequence; omission of a *data-stmt-repeat*  
2 has the effect of a repeat factor of 1.
- 3 7 A zero-sized array or a *data-implied-do* with an iteration count of zero contributes no variables to the expanded  
4 sequence of variables, but a zero-length scalar character variable does contribute a variable to the expanded  
5 sequence. A *data-stmt-constant* with a repeat factor of zero contributes no *data-stmt-constants* to the expanded  
6 sequence of scalar *data-stmt-constants*.
- 7 8 The expanded sequences of variables and *data-stmt-constants* are in one-to-one correspondence. Each *data-stmt-*  
8 *constant* specifies the initial value, initial data *target*, or *null-init* for the corresponding variable. The lengths of  
9 the two expanded sequences shall be the same.
- 10 9 A *data-stmt-constant* shall be *null-init* or *initial-data-target* if and only if the corresponding *data-stmt-object* has  
11 the **POINTER** attribute. If *data-stmt-constant* is *null-init*, the initial association status of the corresponding data  
12 statement object is **disassociated**. If *data-stmt-constant* is *initial-data-target* the corresponding data statement  
13 object shall be data-pointer-initialization compatible (7.5.4.6) with the initial data *target*; the data statement  
14 object is initially associated with the *target*.
- 15 10 A *data-stmt-constant* other than *boz-literal-constant*, *null-init*, or *initial-data-target* shall be compatible with its  
16 corresponding variable according to the rules of intrinsic assignment (10.2.1.2). The variable is initially defined  
17 with the value specified by the *data-stmt-constant*; if necessary, the value is converted according to the rules of  
18 intrinsic assignment (10.2.1.3) to a value that agrees in type, type parameters, and shape with the variable.
- 19 11 If a *data-stmt-constant* is a *boz-literal-constant*, the corresponding variable shall be of type integer. The *boz-*  
20 *literal-constant* is treated as if it were converted by the intrinsic function **INT** (16.9.100) to type integer with the  
21 kind type parameter of the variable.

**NOTE 8.30**

Examples of DATA statements are:

```
CHARACTER (LEN = 10) NAME
INTEGER, DIMENSION (0:9) :: MILES
REAL, DIMENSION (100, 100) :: SKEW
TYPE (NODE), POINTER :: HEAD_OF_LIST
TYPE (PERSON) MYNAME, YOURNAME
DATA NAME / 'JOHN DOE' /, MILES / 10 * 0 /
DATA ((SKEW (K, J), J = 1, K), K = 1, 100) / 5050 * 0.0 /
DATA ((SKEW (K, J), J = K + 1, 100), K = 1, 99) / 4950 * 1.0 /
DATA HEAD_OF_LIST / NULL() /
DATA MYNAME / PERSON (21, 'JOHN SMITH') /
DATA YOURNAME % AGE, YOURNAME % NAME / 35, 'FRED BROWN' /
```

The character variable NAME is initialized with the value JOHN DOE with padding on the right because the length of the constant is less than the length of the variable. All ten elements of the integer array MILES are initialized to zero. The two-dimensional array SKEW is initialized so that the lower triangle of SKEW is zero and the strict upper triangle is one. The structures MYNAME and YOURNAME are declared using the derived type PERSON from Note 7.15. The pointer HEAD\_OF\_LIST is declared using the derived type NODE from Note 7.35; it is initially **disassociated**. MYNAME is initialized by a **structure constructor**. YOURNAME is initialized by supplying a separate value for each component.

**8.6.8 DIMENSION statement**

22 R848 *dimension-stmt* is DIMENSION [ :: ] *array-name* ( *array-spec* ) ■  
23 ■ [ , *array-name* ( *array-spec* ) ] ...  
24

- 25 1 The DIMENSION statement specifies the **DIMENSION attribute** (8.5.8) for a list of objects.

**NOTE 8.31**

An example of a DIMENSION statement is:

```
DIMENSION A (10), B (10, 70), C (:)
```

**8.6.9 INTENT statement**

R849 *intent-stmt* is INTENT ( *intent-spec* ) [ :: ] *dummy-arg-name-list*

- 1 The INTENT statement specifies the **INTENT attribute** (8.5.10) for the **dummy arguments** in the list.

**NOTE 8.32**

An example of an INTENT statement is:

```
SUBROUTINE EX (A, B)
  INTENT (INOUT) :: A, B
```

**8.6.10 OPTIONAL statement**

R850 *optional-stmt* is OPTIONAL [ :: ] *dummy-arg-name-list*

- 1 The OPTIONAL statement specifies the **OPTIONAL attribute** (8.5.12) for the **dummy arguments** in the list.

**NOTE 8.33**

An example of an OPTIONAL statement is:

```
SUBROUTINE EX (A, B)
  OPTIONAL :: B
```

**8.6.11 PARAMETER statement**

- 1 The PARAMETER statement specifies the **PARAMETER attribute** (8.5.13) and the values for the **named constants** in the list.

R851 *parameter-stmt* is PARAMETER ( *named-constant-def-list* )

R852 *named-constant-def* is *named-constant* = *constant-expr*

- 2 If a **named constant** is defined by a PARAMETER statement, it shall not be subsequently declared to have a type or type parameter value that differs from the type and type parameters it would have if declared implicitly (8.7). A named array constant defined by a PARAMETER statement shall have its **rank** specified in a prior specification statement.
- 3 The constant expression that corresponds to a named constant shall have type and **type parameters** that conform with the named constant as specified for intrinsic assignment (10.2.1.2). If the named constant has implied shape, the expression shall have the same rank as the named constant; otherwise, the expression shall either be scalar or have the same shape as the named constant.
- 4 The value of each **named constant** is that specified by the corresponding **constant expression**; if necessary, the value is converted according to the rules of intrinsic assignment (10.2.1.3) to a value that agrees in type, type parameters, and shape with the **named constant**.

**NOTE 8.34**

An example of a PARAMETER statement is:

```
PARAMETER (MODULUS = MOD (28, 3), NUMBER_OF_SENATORS = 100)
```

### 8.6.12 POINTER statement

R853 *pointer-stmt* is POINTER [ :: ] *pointer-decl-list*

R854 *pointer-decl* is *object-name* [ ( *deferred-shape-spec-list* ) ]  
or *proc-entity-name*

C887 A *proc-entity-name* shall have the EXTERNAL attribute.

- 1 The POINTER statement specifies the POINTER attribute (8.5.14) for a list of entities.

#### NOTE 8.35

An example of a POINTER statement is:

```
TYPE (NODE) :: CURRENT
POINTER :: CURRENT, A (:, :)
```

### 8.6.13 PROTECTED statement

R855 *protected-stmt* is PROTECTED [ :: ] *entity-name-list*

- 1 The PROTECTED statement specifies the PROTECTED attribute (8.5.15) for a list of entities.

### 8.6.14 SAVE statement

R856 *save-stmt* is SAVE [ [ :: ] *saved-entity-list* ]

R857 *saved-entity* is *object-name*  
or *proc-pointer-name*  
or / *common-block-name* /

R858 *proc-pointer-name* is *name*

C888 (R856) If a SAVE statement with an omitted saved entity list appears in a scoping unit, no other appearance of the SAVE *attr-spec* or SAVE statement is permitted in that scoping unit.

C889 A *proc-pointer-name* shall be the name of a procedure pointer.

- 1 A SAVE statement with a saved entity list specifies the SAVE attribute (8.5.16) for a list of entities. A SAVE statement without a saved entity list is treated as though it contained the names of all allowed items in the same scoping unit.

#### NOTE 8.36

An example of a SAVE statement is:

```
SAVE A, B, C, / BLOCKA /, D
```

### 8.6.15 TARGET statement

R859 *target-stmt* is TARGET [ :: ] *target-decl-list*

R860 *target-decl* is *object-name* [ ( *array-spec* ) ] ■  
■ [ *lbracket coarray-spec rbracket* ]

- 1 The TARGET statement specifies the TARGET attribute (8.5.17) for a list of objects.

**NOTE 8.37**

An example of a TARGET statement is:

```
TARGET :: A (1000, 1000), B
```

**1 8.6.16 VALUE statement**

2 R861 *value-stmt* is VALUE [ :: ] *dummy-arg-name-list*

3 1 The VALUE statement specifies the **VALUE attribute** (8.5.18) for a list of **dummy arguments**.

**4 8.6.17 VOLATILE statement**

5 R862 *volatile-stmt* is VOLATILE [ :: ] *object-name-list*

6 1 The VOLATILE statement specifies the **VOLATILE attribute** (8.5.19) for a list of objects.

**7 8.7 IMPLICIT statement**

8 1 In a **scoping unit**, an IMPLICIT statement specifies a type, and possibly type parameters, for all implicitly  
9 typed data entities whose names begin with one of the letters specified in the statement. An IMPLICIT NONE  
10 statement can indicate that no implicit typing rules are to apply in a particular **scoping unit**, or that **external**  
11 and **dummy** procedures need to be explicitly given the **EXTERNAL attribute**.

12 R863 *implicit-stmt* is IMPLICIT *implicit-spec-list*  
13 or IMPLICIT NONE [ ( [ *implicit-none-spec-list* ] ) ]

14 R864 *implicit-spec* is *declaration-type-spec* ( *letter-spec-list* )

15 R865 *letter-spec* is *letter* [ – *letter* ]

16 R866 *implicit-none-spec* is EXTERNAL  
17 or TYPE

18 C890 (R863) If an IMPLICIT NONE statement appears in a **scoping unit**, it shall precede any **PARAMETER**  
19 **statements** that appear in the **scoping unit**. No more than one IMPLICIT NONE statement shall appear  
20 in a **scoping unit**.

21 C891 If an IMPLICIT NONE statement in a scoping unit has an *implicit-none-spec* of TYPE or has no *implicit-*  
22 *none-spec-list*, there shall be no other IMPLICIT statements in the scoping unit.

23 C892 (R865) If the minus and second *letter* appear, the second letter shall follow the first letter alphabetically.

24 C893 If IMPLICIT NONE with an *implicit-none-spec* of EXTERNAL appears within a **scoping unit**, the  
25 name of an **external** or **dummy** procedure in that **scoping unit** or in a contained **subprogram** or **BLOCK**  
26 **construct** shall have an **explicit interface** or be explicitly declared to have the **EXTERNAL attribute**.

27 2 A *letter-spec* consisting of two *letters* separated by a minus is equivalent to writing a list containing all of the letters  
28 in alphabetical order in the alphabetic sequence from the first letter through the second letter. For example, A–C  
29 is equivalent to A, B, C. The same letter shall not appear as a single letter, or be included in a range of letters,  
30 more than once in all of the IMPLICIT statements in a **scoping unit**.

31 3 In each **scoping unit**, there is a mapping, which may be null, between each of the letters A, B, ..., Z and a  
32 type (and type parameters). An IMPLICIT statement specifies the mapping for the letters in its *letter-spec-*  
33 *list*. IMPLICIT NONE with an *implicit-none-spec* of TYPE or with no *implicit-none-spec-list* specifies the null  
34 mapping for all the letters. If a mapping is not specified for a letter, the default for a **program unit** or an **interface**  
35 **body** is default integer if the letter is I, J, ..., or N and default real otherwise, and the default for a **BLOCK**

- 1 [construct](#), [internal subprogram](#), or [module subprogram](#) is the mapping in the [host scoping unit](#).
- 2 4 Any data entity that is not explicitly declared by a [type declaration statement](#), is not an intrinsic function, is  
 3 not a [component](#), and is not accessed by [use](#) or [host](#) association is declared implicitly to be of the type (and  
 4 type parameters) mapped from the first letter of its name, provided the mapping is not null. The mapping for  
 5 the first letter of the data entity shall either have been established by a prior IMPLICIT statement or be the  
 6 default mapping for the letter. The data entity is treated as if it were declared in an explicit type declaration; if  
 7 the outermost [inclusive scope](#) in which it appears is not a type definition, it is declared in that scope, otherwise  
 8 it is declared in the host of that scope. An explicit type specification in a [FUNCTION statement](#) overrides an  
 9 IMPLICIT statement for the [result](#) of that function.

**NOTE 8.38**

The following are examples of the use of IMPLICIT statements:

```

MODULE EXAMPLE_MODULE
  IMPLICIT NONE
  ...
  INTERFACE
    FUNCTION FUN (I)      ! Not all data entities need to
      INTEGER FUN        ! be declared explicitly
    END FUNCTION FUN
  END INTERFACE
CONTAINS
  FUNCTION JFUN (J)      ! All data entities need to
    INTEGER JFUN, J     ! be declared explicitly.
    ...
  END FUNCTION JFUN
END MODULE EXAMPLE_MODULE
SUBROUTINE SUB
  IMPLICIT COMPLEX (C)
  C = (3.0, 2.0)        ! C is implicitly declared COMPLEX
  ...
CONTAINS
  SUBROUTINE SUB1
    IMPLICIT INTEGER (A, C)
    C = (0.0, 0.0)      ! C is host associated and of
                       ! type complex
    Z = 1.0             ! Z is implicitly declared REAL
    A = 2               ! A is implicitly declared INTEGER
    CC = 1             ! CC is implicitly declared INTEGER
    ...
  END SUBROUTINE SUB1
  SUBROUTINE SUB2
    Z = 2.0             ! Z is implicitly declared REAL and
                       ! is different from the variable of
                       ! the same name in SUB1
    ...
  END SUBROUTINE SUB2
  SUBROUTINE SUB3
    USE EXAMPLE_MODULE ! Accesses integer function FUN
                       ! by use association
    Q = FUN (K)        ! Q is implicitly declared REAL and
    ...               ! K is implicitly declared INTEGER
  END SUBROUTINE SUB3
END SUBROUTINE SUB

```

**NOTE 8.39**

The following is an example of a mapping to a derived type that is inaccessible in the local scope:

```
PROGRAM MAIN
  IMPLICIT TYPE(BLOB) (A)
  TYPE BLOB
    INTEGER :: I
  END TYPE BLOB
  TYPE(BLOB) :: B
  CALL STEVE
CONTAINS
  SUBROUTINE STEVE
    INTEGER :: BLOB
    ..
    AA = B
    ..
  END SUBROUTINE STEVE
END PROGRAM MAIN
```

In the subroutine STEVE, it is not possible to explicitly declare a variable to be of type BLOB because BLOB has been given a different meaning, but implicit mapping for the letter A still maps to type BLOB, so AA is of type BLOB.

**NOTE 8.40**

Implicit typing is not affected by **BLOCK constructs**. For example, in

```
SUBROUTINE S(N)
  ...
  IF (N>0) THEN
    BLOCK
      NSQP = CEILING (SQRT (DBLE (N)))
    END BLOCK
  END IF
  ...
  IF (N>0) THEN
    BLOCK
      PRINT *,NSQP
    END BLOCK
  END IF
END SUBROUTINE
```

even if the only two appearances of NSQP are within the **BLOCK constructs**, the scope of NSQP is the whole subroutine S.

**NOTE 8.41**

In the subprogram

```
SUBROUTINE EXAMPLE (X, Y)
  IMPLICIT NONE (EXTERNAL)
  REAL, EXTERNAL :: G
  REAL :: X, Y
  X = F (Y)           ! Invalid: F lacks the EXTERNAL attribute.
  X = G (Y)           ! Valid: G has the EXTERNAL attribute.
END SUBROUTINE
```

**NOTE 8.41 (cont.)**

the referenced function F needs to have the **EXTERNAL** attribute (8.5.9).

**8.8 IMPORT statement**

R867 *import-stmt* is IMPORT [[ :: ] *import-name-list* ]  
 or IMPORT, ONLY : *import-name-list*  
 or IMPORT, NONE  
 or IMPORT, ALL

C894 (R867) An IMPORT statement shall not appear in the **scoping unit** of a *main-program*, *external-subprogram*, *module*, or *block-data*.

C895 (R867) Each *import-name* shall be the name of an entity in the **host scoping unit**.

C896 If any IMPORT statement in a **scoping unit** has an **ONLY** specifier, all IMPORT statements in that **scoping unit** shall have an **ONLY** specifier.

C897 IMPORT, NONE shall not appear in the **scoping unit** of a submodule.

C898 If an IMPORT, NONE or IMPORT, ALL statement appears in a **scoping unit**, no other IMPORT statement shall appear in that **scoping unit**.

C899 Within an interface body, an entity that is accessed by host association shall be accessible by host or use association within the **host scoping unit**, or explicitly declared prior to the interface body.

C8100 An entity whose name appears as an *import-name* or which is made accessible by an IMPORT, ALL statement shall not appear in any context described in 19.5.1.4 that would cause the host entity of that name to be inaccessible.

1 If the **ONLY** specifier appears on an IMPORT statement in a **scoping unit**, an entity is only accessible by host association if its name appears as an *import-name* in that **scoping unit**.

2 An IMPORT, NONE statement in a **scoping unit** specifies that no entities in the **host scoping unit** are accessible by host association in that **scoping unit**. This is the default for an interface body for an **external** or **dummy** procedure.

3 An IMPORT, ALL statement in a **scoping unit** specifies that all entities from the **host scoping unit** are accessible by host association in that **scoping unit**.

4 If an IMPORT statement with no specifier and no *import-name-list* appears in a **scoping unit**, every entity in the **host scoping unit** is accessible unless its name appears in a context described in 19.5.1.4 that causes it to be inaccessible. This is the default for a nested scoping unit other than an interface body for an **external** or **dummy** procedure.

5 If an IMPORT statement with an *import-name-list* appears in a **scoping unit**, each named entity from the **host scoping unit** is accessible by host association.

**NOTE 8.42**

The IMPORT, NONE statement can be used to prevent accidental host association:

```
SUBROUTINE s(x,n)
  IMPLICIT NONE
  IMPORT, NONE
  ...
  DO i=1,n ! Forces I to be locally declared.
```

**NOTE 8.43**

The IMPORT, ALL statement can be used to confirm the default rules and prevent accidental “shadowing” of host entities:

```

SUBROUTINE outer
  REAL x
  ...
CONTAINS
  SUBROUTINE inner
    IMPORT, ALL
    ...
    x = x + 1 ! X was prevented from being locally declared...
              ! so this is the host X.
  
```

**NOTE 8.44**

The IMPORT, ONLY statement can be used to document deliberate access via host association whilst blocking accidental access:

```

SUBROUTINE sub
  IMPORT, ONLY :: x, y
  ...
  x = y + z ! Only X and Y are imported, so Z is local.

```

**NOTE 8.45**

The IMPORT statement can be used to allow [module procedures](#) to have [dummy arguments](#) that are procedures with [assumed-shape](#) arguments of an opaque type. For example:

```

MODULE M
  TYPE T
    PRIVATE ! T is an opaque type
    ...
  END TYPE
CONTAINS
  SUBROUTINE PROCESS(X,Y,RESULT,MONITOR)
    TYPE(T), INTENT(IN) :: X(:, :), Y(:, :)
    TYPE(T), INTENT(OUT) :: RESULT(:, :)
    INTERFACE
      SUBROUTINE MONITOR(ITERATION_NUMBER, CURRENT_ESTIMATE)
        IMPORT T
        INTEGER, INTENT(IN) :: ITERATION_NUMBER
        TYPE(T), INTENT(IN) :: CURRENT_ESTIMATE(:, :)
      END SUBROUTINE
    END INTERFACE
    ...
  END SUBROUTINE
END MODULE

```

The MONITOR [dummy procedure](#) requires an [explicit interface](#) because it has an [assumed-shape](#) array argument, but TYPE(T) would not be available inside the interface body without the IMPORT statement.



## 8.9 NAMELIST statement

1 A NAMELIST statement specifies a group of named data objects, which may be referred to by a single name for the purpose of data transfer (12.6, 13.11).

R868 *namelist-stmt* is NAMELIST ■  
 ■ / *namelist-group-name* / *namelist-group-object-list* ■  
 ■ [ [ , ] / *namelist-group-name* / ■  
 ■ *namelist-group-object-list* ] ...

C8101 (R868) The *namelist-group-name* shall not be a name accessed by use association.

R869 *namelist-group-object* is *variable-name*

C8102 (R869) A *namelist-group-object* shall not be an assumed-size array.

C8103 (R868) A *namelist-group-object* shall not have the PRIVATE attribute if the *namelist-group-name* has the PUBLIC attribute.

2 The order in which the values appear on output is the same as the order of the *namelist-group-objects* in the namelist group object list; if a variable appears more than once as a *namelist-group-object* for the same namelist group, its value appears once for each occurrence.

3 Any *namelist-group-name* may occur more than once in the NAMELIST statements in a scoping unit. The *namelist-group-object-list* following each successive appearance of the same *namelist-group-name* in a scoping unit is treated as a continuation of the list for that *namelist-group-name*.

4 A namelist group object may be a member of more than one namelist group.

5 A namelist group object shall either be accessed by use or host association or shall have its declared type, kind type parameters of the declared type, and rank specified by previous specification statements or the procedure heading in the same scoping unit or by the implicit typing rules in effect for the scoping unit. If a namelist group object is typed by the implicit typing rules, its appearance in any subsequent type declaration statement shall confirm the implied type and type parameters.

### NOTE 8.46

An example of a NAMELIST statement is:

```
NAMELIST /NLIST/ A, B, C
```

## 8.10 Storage association of data objects

### 8.10.1 EQUIVALENCE statement

#### 8.10.1.1 General

1 An EQUIVALENCE statement is used to specify the sharing of storage units by two or more objects in a scoping unit. This causes storage association (19.5.3) of the objects that share the storage units.

2 If the equivalenced objects have differing type or type parameters, the EQUIVALENCE statement does not cause type conversion or imply mathematical equivalence. If a scalar and an array are equivalenced, the scalar does not have array properties and the array does not have the properties of a scalar.

R870 *equivalence-stmt* is EQUIVALENCE *equivalence-set-list*

R871 *equivalence-set* is ( *equivalence-object* , *equivalence-object-list* )

R872 *equivalence-object* is *variable-name*  
 or *array-element*  
 or *substring*

- 1 C8104 (R872) An *equivalence-object* shall not be a designator with a base object that is a dummy argument, a function result,  
2 a pointer, an allocatable variable, a derived-type object that has an allocatable or pointer ultimate component, an object  
3 of a nonsequence derived type, an automatic data object, a coarray, a variable with the BIND attribute, a variable in a  
4 common block that has the BIND attribute, or a named constant.
- 5 C8105 (R872) An *equivalence-object* shall not be a designator that has more than one *part-ref*.
- 6 C8106 (R872) An *equivalence-object* shall not have the TARGET attribute.
- 7 C8107 (R872) Each subscript or substring range expression in an *equivalence-object* shall be an integer constant expression  
8 (10.1.12).
- 9 C8108 (R871) If an *equivalence-object* is default integer, default real, double precision real, default complex, default logical, or of  
10 numeric sequence type, all of the objects in the equivalence set shall be of these types and kinds.
- 11 C8109 (R871) If an *equivalence-object* is default character or of character sequence type, all of the objects in the equivalence set  
12 shall be of these types and kinds.
- 13 C8110 (R871) If an *equivalence-object* is of a sequence type that is not a numeric sequence or character sequence type, all of the  
14 objects in the equivalence set shall be of that type.
- 15 C8111 (R871) If an *equivalence-object* is of an intrinsic type but is not default integer, default real, double precision real, default  
16 complex, default logical, or default character, all of the objects in the equivalence set shall be of the same type with the  
17 same kind type parameter value.
- 18 C8112 (R872) If an *equivalence-object* has the PROTECTED attribute, all of the objects in the equivalence set shall have the  
19 PROTECTED attribute.
- 20 C8113 (R872) The name of an *equivalence-object* shall not be a name made accessible by use association.
- 21 C8114 (R872) A *substring* shall not have length zero.

**NOTE 8.47**

The EQUIVALENCE statement allows the equivalencing of sequence structures and the equivalencing of objects of intrinsic type with nondefault type parameters, but there are strict rules regarding the appearance of these objects in an EQUIVALENCE statement.

In addition to the above constraints, further rules on the interaction of EQUIVALENCE statements and default initialization are given in 19.5.3.4.

### 22 8.10.1.2 Equivalence association

- 23 1 An EQUIVALENCE statement specifies that the storage sequences (19.5.3.2) of the data objects specified in an *equivalence-set* are  
24 storage associated. All of the nonzero-sized sequences in the *equivalence-set*, if any, have the same first storage unit, and all of  
25 the zero-sized sequences in the *equivalence-set*, if any, are storage associated with one another and with the first storage unit of  
26 any nonzero-sized sequences. This causes the storage association of the data objects in the *equivalence-set* and may cause storage  
27 association of other data objects.
- 28 2 If any data object in an *equivalence-set* has the SAVE attribute, all other objects in the *equivalence-set* have the SAVE attribute;  
29 this may be confirmed by explicit specification.

### 30 8.10.1.3 Equivalence of default character objects

- 31 1 A default character data object shall not be equivalenced to an object that is not default character and not of a character sequence  
32 type. The lengths of equivalenced default character objects need not be the same.
- 33 2 An EQUIVALENCE statement specifies that the storage sequences of all the default character data objects specified in an *equivalence-*  
34 *set* are storage associated. All of the nonzero-sized sequences in the *equivalence-set*, if any, have the same first character storage unit,  
35 and all of the zero-sized sequences in the *equivalence-set*, if any, are storage associated with one another and with the first character  
36 storage unit of any nonzero-sized sequences. This causes the storage association of the data objects in the *equivalence-set* and may  
37 cause storage association of other data objects.

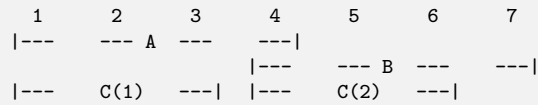
**NOTE 8.48**

For example, using the declarations:

```
CHARACTER (LEN = 4) :: A, B
CHARACTER (LEN = 3) :: C (2)
EQUIVALENCE (A, C (1)), (B, C (2))
```

**NOTE 8.48 (cont.)**

the association of A, B, and C can be illustrated graphically as:

**8.10.1.4 Array names and array element designators**

- 1 For a nonzero-sized array, the use of the array name unqualified by a subscript list as an *equivalence-object* has the same effect as using an array element *designator* that identifies the first element of the array.

**8.10.1.5 Restrictions on EQUIVALENCE statements**

- 1 An EQUIVALENCE statement shall not specify that the same *storage unit* is to occur more than once in a *storage sequence*.  
 2 An EQUIVALENCE statement shall not specify that consecutive *storage units* are to be nonconsecutive.

**8.10.2 COMMON statement****8.10.2.1 General**

- 1 The COMMON statement specifies blocks of physical storage, called *common blocks*, that can be accessed by any of the *scoping units* in a program. Thus, the COMMON statement provides a global data facility based on *storage association* (19.5.3).  
 2 A *common block* that does not have a name is called *blank common*.

R873 *common-stmt* is COMMON ■  
 ■ [ / [ *common-block-name* ] / ] *common-block-object-list* ■  
 ■ [ [ , ] / [ *common-block-name* ] / ■  
 ■ *common-block-object-list* ] ...

R874 *common-block-object* is *variable-name* [ ( *array-spec* ) ]

C8115 (R874) An *array-spec* in a *common-block-object* shall be an *explicit-shape-spec-list*.

C8116 (R874) Only one appearance of a given *variable-name* is permitted in all *common-block-object-lists* within a *scoping unit*.

C8117 (R874) A *common-block-object* shall not be a *dummy argument*, a *function result*, an *allocatable* variable, a derived-type object with an *ultimate component* that is *allocatable*, a *procedure pointer*, an *automatic data object*, a variable with the *BIND* attribute, an *unlimited polymorphic pointer*, or a *coarray*.

C8118 (R874) If a *common-block-object* is of a derived type, the type shall have the *BIND* attribute or the *SEQUENCE* attribute and it shall have no *default initialization*.

C8119 (R874) A *variable-name* shall not be a name made accessible by *use association*.

- 3 In each COMMON statement, the data objects whose names appear in a common block object list following a *common block* name are declared to be in that *common block*. If the first *common block* name is omitted, all data objects whose names appear in the first common block object list are specified to be in *blank common*. Alternatively, the appearance of two slashes with no *common block* name between them declares the data objects whose names appear in the common block object list that follows to be in *blank common*.

- 4 Any *common block* name or an omitted *common block* name for *blank common* may occur more than once in one or more COMMON statements in a *scoping unit*. The common block list following each successive appearance of the same common block name in a *scoping unit* is treated as a continuation of the list for that common block name. Similarly, each blank common block object list in a *scoping unit* is treated as a continuation of *blank common*.

- 5 The form *variable-name* (*array-spec*) specifies the *DIMENSION* attribute for that variable.

- 6 If derived-type objects of *numeric sequence type* or *character sequence type* (7.5.2.3) appear in *common*, it is as if the individual components were enumerated directly in the common list.

**8.10.2.2 Common block storage sequence**

- 1 For each *common block* in a *scoping unit*, a common block *storage sequence* is formed as follows:

- 1 (1) A **storage sequence** is formed consisting of the sequence of **storage units** in the **storage sequences** (19.5.3.2) of all data  
 2 objects in the common block object lists for the **common block**. The order of the **storage sequences** is the same as the  
 3 order of the appearance of the common block object lists in the **scoping unit**.
- 4 (2) The **storage sequence** formed in (1) is extended to include all **storage units** of any **storage sequence** associated with it  
 5 by equivalence association. The **sequence** shall be extended only by adding **storage units** beyond the last **storage unit**.  
 6 Data objects associated with an entity in a **common block** are considered to be in that **common block**.
- 7 2 Only COMMON statements and EQUIVALENCE statements appearing in the **scoping unit** contribute to common block **storage**  
 8 **sequences** formed in that **scoping unit**.

### 9 8.10.2.3 Size of a common block

- 10 1 The size of a common block is the size of its common block **storage sequence**, including any extensions of the **sequence** resulting from  
 11 equivalence association.

### 12 8.10.2.4 Common association

- 13 1 Within a program, the common block **storage sequences** of all nonzero-sized **common blocks** with the same name have the same first  
 14 **storage unit**, and the common block **storage sequences** of all zero-sized **common blocks** with the same name are storage associated  
 15 with one another. Within a program, the common block **storage sequences** of all nonzero-sized **blank common** blocks have the same  
 16 first **storage unit** and the **storage sequences** of all zero-sized **blank common** blocks are associated with one another and with the first  
 17 **storage unit** of any nonzero-sized **blank common** blocks. This results in the association of objects in different **scoping units**. Use or  
 18 **host** association may cause these associated objects to be accessible in the same **scoping unit**.
- 19 2 A nonpointer object that is default integer, default real, double precision real, default complex, default logical, or of **numeric sequence**  
 20 **type** shall be associated only with nonpointer objects of these types and kinds.
- 21 3 A nonpointer object that is default character or of **character sequence type** shall be associated only with nonpointer objects of these  
 22 types and kinds.
- 23 4 A nonpointer object of a derived type that is not a **numeric sequence** or **character sequence type** shall be associated only with  
 24 nonpointer objects of the same type.
- 25 5 A nonpointer object of intrinsic type but which is not default integer, default real, double precision real, default complex, default  
 26 logical, or default character shall be associated only with nonpointer objects of the same type and type parameters.
- 27 6 A data pointer shall be storage associated only with data pointers of the same type and **rank**. Data pointers that are storage  
 28 associated shall have **deferred** the same type parameters; corresponding nondeferred **type parameters** shall have the same value.
- 29 7 An object with the **TARGET attribute** shall be storage associated only with another object that has the **TARGET attribute** and the  
 30 same type and type parameters.

#### NOTE 8.49

A **common block** is permitted to contain sequences of different **storage units**, provided each **scoping unit** that accesses the **common block** specifies an identical sequence of **storage units** for the **common block**. For example, this allows a single **common block** to contain both **numeric** and **character storage units**.

Association in different **scoping units** between objects of default type, objects of double precision real type, and sequence structures is permitted according to the rules for equivalence objects (8.10.1).

### 31 8.10.2.5 Differences between named common and blank common

- 32 1 A **blank common** block has the same properties as a named **common block**, except for the following.
- 33 • Execution of a **RETURN** or **END** statement might cause data objects in a named **common block** to become undefined unless  
 34 the **common block** has the **SAVE attribute**, but never causes nonpointer data objects in **blank common** to become undefined  
 35 (19.6.6).
  - 36 • Named **common blocks** of the same name shall be of the same size in all **scoping units** of a program in which they appear, but  
 37 **blank common** blocks may be of different sizes.
  - 38 • A data object in a named **common block** may be initially defined by means of a **DATA statement** or **type declaration statement**  
 39 in a **block data program unit** (14.3), but objects in **blank common** shall not be initially defined.

### 40 8.10.3 Restrictions on common and equivalence

- 41 1 An **EQUIVALENCE statement** shall not cause the **storage sequences** of two different **common blocks** to be associated.
- 42 2 Equivalence association shall not cause a derived-type object with **default initialization** to be associated with an object in a **common**  
 43 **block**.

- 1 3 Equivalence association shall not cause a common block [storage sequence](#) to be extended by adding [storage units](#) preceding the first  
2 [storage unit](#) of the first object specified in a COMMON statement for the [common block](#).

1

(Blank page)

## 9 Use of data objects

### 9.1 Designator

R901 *designator* is *object-name*  
 or *array-element*  
 or *array-section*  
 or *coindexed-named-object*  
 or *complex-part-designator*  
 or *structure-component*  
 or *substring*

1 The appearance of a *data object designator* in a context that requires its value is termed a reference.

### 9.2 Variable

R902 *variable* is *designator*  
 or *function-reference*

C901 (R902) *designator* shall not be a constant or a subobject of a constant.

C902 (R902) *function-reference* shall have a data pointer result.

1 A variable is either the data object denoted by *designator* or the *target* of the pointer resulting from the evaluation of *function-reference*; this pointer shall be associated.

2 A reference is permitted only if the variable is defined. A reference to a data pointer is permitted only if the pointer is associated with a *target* object that is defined. A variable becomes defined with a value when events described in 19.6.5 occur.

R903 *variable-name* is *name*

C903 (R903) *variable-name* shall be the name of a variable.

R904 *logical-variable* is *variable*

C904 (R904) *logical-variable* shall be of type logical.

R905 *char-variable* is *variable*

C905 (R905) *char-variable* shall be of type character.

R906 *default-char-variable* is *variable*

C906 (R906) *default-char-variable* shall be default character.

R907 *int-variable* is *variable*

C907 (R907) *int-variable* shall be of type integer.

#### NOTE 9.1

For example, given the declarations:

```
CHARACTER (10) A, B (10)
```

**NOTE 9.1 (cont.)**

TYPE (PERSON) P ! See Note 7.15

then A, B, B (1), B (1:5), P % AGE, and A (1:1) are all variables.

**9.3 Constants**

- 1 A constant (6.2.3) is a literal constant or a **named constant**. A literal constant is a scalar denoted by a syntactic form, which indicates its type, type parameters, and value. A **named constant** is a constant that has a name; the name has the **PARAMETER attribute** (8.5.13, 8.6.11). A reference to a constant is always permitted; redefinition of a constant is never permitted.

**9.4 Scalars****9.4.1 Substrings**

- 1 A substring is a **contiguous** portion of a character string (7.4.4).

R908 *substring* is *parent-string* ( *substring-range* )

R909 *parent-string* is *scalar-variable-name*  
 or *array-element*  
 or *coindexed-named-object*  
 or *scalar-structure-component*  
 or *scalar-constant*

R910 *substring-range* is [ *scalar-int-expr* ] : [ *scalar-int-expr* ]

C908 (R909) *parent-string* shall be of type character.

- 2 The value of the first *scalar-int-expr* in *substring-range* is the starting point of the substring and the value of the second one is the ending point of the substring. The length of a substring is the number of characters in the substring and is  $\text{MAX}(l - f + 1, 0)$ , where  $f$  and  $l$  are the starting and ending points, respectively.
- 3 Let the characters in the parent string be numbered 1, 2, 3, ...,  $n$ , where  $n$  is the length of the parent string. Then the characters in the substring are those from the parent string from the starting point and proceeding in sequence up to and including the ending point. If the starting point is greater than the ending point, the substring has length zero; otherwise, both the starting point and the ending point shall be within the range 1, 2, ...,  $n$ . If the starting point is not specified, the default value is 1. If the ending point is not specified, the default value is  $n$ .

**NOTE 9.2**

Examples of character substrings are:

B(1)(1:5)	array element as parent string
P%NAME(1:1)	structure component as parent string
ID(4:9)	scalar variable name as parent string
'0123456789'(N:N)	character constant as parent string

**9.4.2 Structure components**

- 1 A **structure component** is part of an object of derived type; it may be referenced by an **object designator**. A **structure component** may be a scalar or an array.

R911 *data-ref* is *part-ref* [ % *part-ref* ] ...



- 1 R912 *part-ref* is *part-name* [ ( *section-subscript-list* ) ] [ *image-selector* ]
- 2 C909 (R911) Each *part-name* except the rightmost shall be of derived type.
- 3 C910 (R911) Each *part-name* except the leftmost shall be the name of a component of the **declared type** of the  
4 preceding *part-name*.
- 5 C911 (R911) If the rightmost *part-name* is of **abstract type**, *data-ref* shall be **polymorphic**.
- 6 C912 (R911) The leftmost *part-name* shall be the name of a data object.
- 7 C913 (R912) If a *section-subscript-list* appears, the number of *section-subscripts* shall equal the **rank** of *part-*  
8 *name*.
- 9 C914 (R912) If *image-selector* appears, the number of *cosubscripts* shall be equal to the **corank** of *part-name*.
- 10 C915 A *data-ref* shall not be of type **C\_PTR** or **C\_FUNPTR** from the intrinsic module **ISO\_C\_BINDING** (18.2) if one of its *part-refs* has an *image-selector*.
- 11
- 12 C916 (R912) If *image-selector* appears and *part-name* is an array, *section-subscript-list* shall appear.
- 13 C917 (R911) Except as an **actual argument** to an intrinsic **inquiry function** or as the *designator* in a type  
14 parameter inquiry, a *data-ref* shall not be a **polymorphic subobject** of a **coindexed object** and shall not  
15 be a **coindexed object** that has a **polymorphic allocatable potential subobject component**.
- 16 2 The **rank** of a *part-ref* of the form *part-name* is the **rank** of *part-name*. The **rank** of a *part-ref* that has a section  
17 subscript list is the number of subscript triplets and **vector subscripts** in the list.
- 18 C918 (R911) There shall not be more than one *part-ref* with nonzero **rank**. A *part-name* to the right of a  
19 *part-ref* with nonzero **rank** shall not have the **ALLOCATABLE** or **POINTER** attribute.
- 20 3 The **rank** of a *data-ref* is the **rank** of the *part-ref* with nonzero **rank**, if any; otherwise, the **rank** is zero. The **base**  
21 **object** of a *data-ref* is the data object whose name is the leftmost part name.
- 22 4 The type and type parameters, if any, of a *data-ref* are those of the rightmost part name.
- 23 5 A *data-ref* with more than one *part-ref* is a subobject of its **base object** if none of the *part-names*, except for  
24 possibly the rightmost, is a pointer. If the rightmost *part-name* is the only pointer, then the *data-ref* is a subobject  
25 of its **base object** in contexts that pertain to its **pointer association** status but not in any other contexts.

**NOTE 9.3**

If X is an object of derived type with a pointer component P, then the pointer X%P is a subobject of X when considered as a pointer – that is in contexts where it is not dereferenced.

However the **target** of X%P is not a subobject of X. Thus, in contexts where X%P is dereferenced to refer to the **target**, it is not a subobject of X.

- 26 R913 *structure-component* is *data-ref*
- 27 C919 (R913) There shall be more than one *part-ref* and the rightmost *part-ref* shall not have a *section-subscript-*  
28 *list*.
- 29 6 A **structure component** shall be neither referenced nor defined before the declaration of the **base object**. A  
30 **structure component** is a pointer only if the rightmost part name has the **POINTER** attribute.

**NOTE 9.4**

Examples of structure components are:

SCALAR\_PARENT%SCALAR\_FIELD scalar component of scalar parent

**NOTE 9.4 (cont.)**

ARRAY_PARENT(J)%SCALAR_FIELD	component of array element parent
ARRAY_PARENT(1:N)%SCALAR_FIELD	component of array section parent

For a more elaborate example see C.4.1.

**NOTE 9.5**

The syntax rules are structured such that a *data-ref* that ends in a component name without a following subscript list is a structure component, even when other component names in the *data-ref* are followed by a subscript list. A *data-ref* that ends in a component name with a following subscript list is either an array element or an *array section*. A *data-ref* of nonzero *rank* that ends with a *substring-range* is an *array section*. A *data-ref* of zero *rank* that ends with a *substring-range* is a substring.

**9.4.3 Coindexed named objects**

1 A *coindexed-named-object* is a named scalar *coarray* variable followed by an image selector.

2 R914 *coindexed-named-object* is *data-ref*

3 C920 (R914) The *data-ref* shall contain exactly one *part-ref*. The *part-ref* shall contain an *image-selector*.  
4 The *part-name* shall be the name of a scalar *coarray*.

**9.4.4 Complex parts**

5 R915 *complex-part-designator* is *designator* % RE  
6 or *designator* % IM

7 C921 (R915) The *designator* shall be of complex type.

8 1 If *complex-part-designator* is *designator*%RE it designates the real part of *designator*. If it is *designator*%IM  
9 it designates the imaginary part of *designator*. The type of a *complex-part-designator* is real, and its kind and  
10 shape are those of the *designator*, which can be an *array* or *scalar*.

**NOTE 9.6**

The following are examples of *complex part designators*:

impedance%re	!-- Same value as REAL(impedance)
fft%im	!-- Same value as AIMAG(fft)
x%im = 0.0	!-- Sets the imaginary part of X to zero

**9.4.5 Type parameter inquiry**

13 1 A *type parameter inquiry* is used to inquire about a type parameter of a data object. It applies to both intrinsic  
14 and derived types.

15 R916 *type-param-inquiry* is *designator* % *type-param-name*

16 C922 (R916) The *type-param-name* shall be the name of a type parameter of the *declared type* of the object  
17 designated by the *designator*.

18 2 A *deferred type parameter* of a pointer that is not associated or of an unallocated *allocatable* variable shall not  
19 be inquired about.

**NOTE 9.7**

A *type-param-inquiry* has a syntax like that of a *structure component* reference, but it does not have the same semantics. It is not a variable and thus can never be assigned to. It can be used only as a primary in an expression. It is scalar even if *designator* is an array.

The intrinsic type parameters can also be inquired about by using the intrinsic functions *KIND* and *LEN*.

**NOTE 9.8**

The following are examples of type parameter inquiries:

```

a%kind      !-- A is real.  Same value as KIND(a).
s%len       !-- S is character.  Same value as LEN(s).
b(10)%kind  !-- Inquiry about an array element.
p%dim       !-- P is of the derived type general_point.

```

See Note 7.22 for the definition of the *general\_point* type used in the last example above.

**9.5 Arrays****9.5.1 Order of reference**

- 1 No order of reference to the elements of an array is indicated by the appearance of the array *designator*, except where array element ordering (9.5.3.2) is specified.

**9.5.2 Whole arrays**

- 1 A *whole array* is a named array or a structure component whose final *part-ref* is an array component name; no subscript list is appended.
- 2 The appearance of a *whole array* variable in an executable construct specifies all the elements of the array (5.4.6). The appearance of a *whole array designator* in a nonexecutable statement specifies the entire array except for the appearance of a *whole array* designator in an equivalence set (8.10.1.4). An *assumed-size array* (8.5.8.5) is permitted to appear as a *whole array* in an executable construct or *specification expression* only as an *actual argument* in a *procedure reference* that does not require the shape.

**9.5.3 Array elements and array sections****9.5.3.1 Syntax**

R917 *array-element* is *data-ref*

C923 (R917) Every *part-ref* shall have *rank* zero and the last *part-ref* shall contain a *subscript-list*.

R918 *array-section* is *data-ref* [ ( *substring-range* ) ]  
or *complex-part-designator*

C924 (R918) Exactly one *part-ref* shall have nonzero *rank*, and either the final *part-ref* shall have a *section-subscript-list* with nonzero *rank*, another *part-ref* shall have nonzero *rank*, or the *complex-part-designator* shall be an array.

C925 (R918) If a *substring-range* appears, the rightmost *part-name* shall be of type character.

R919 *subscript* is *scalar-int-expr*

R920 *section-subscript* is *subscript*  
or *subscript-triplet*

or *vector-subscript*

R921 *subscript-triplet* is [ *subscript* ] : [ *subscript* ] [ : *stride* ]

R922 *stride* is *scalar-int-expr*

R923 *vector-subscript* is *int-expr*

C926 (R923) A *vector-subscript* shall be an integer array expression of rank one.

C927 (R921) The second subscript shall not be omitted from a *subscript-triplet* in the last dimension of an assumed-size array.

1 An array element is a scalar. An array section is an array. If a *substring-range* appears in an *array-section*, each element is the designated substring of the corresponding element of the array section.

2 The value of a subscript in an array element shall be within the bounds for its dimension.

**NOTE 9.9**

For example, with the declarations:

```
REAL A (10, 10)
CHARACTER (LEN = 10) B (5, 5, 5)
```

A (1, 2) is an array element, A (1:N:2, M) is a rank-one array section, and B (:, :, :) (2:3) is an array of shape (5, 5, 5) whose elements are substrings of length 2 of the corresponding elements of B.

**NOTE 9.10**

Unless otherwise specified, an array element or array section does not have an attribute of the whole array. In particular, an array element or an array section does not have the POINTER or ALLOCATABLE attribute.

**NOTE 9.11**

Examples of array elements and array sections are:

```
ARRAY_A(1:N:2)%ARRAY_B(I, J)%STRING(K)(:)    array section
SCALAR_PARENT%ARRAY_FIELD(J)                array element
SCALAR_PARENT%ARRAY_FIELD(1:N)              array section
SCALAR_PARENT%ARRAY_FIELD(1:N)%SCALAR_FIELD array section
```

**9.5.3.2 Array element order**

1 The elements of an array form a sequence known as the array element order. The position of an array element in this sequence is determined by the subscript order value of the subscript list designating the element. The subscript order value is computed from the formulas in Table 9.1.

Table 9.1: Subscript order value

Rank	Subscript bounds	Subscript list	Subscript order value
1	$j_1:k_1$	$s_1$	$1 + (s_1 - j_1)$
2	$j_1:k_1, j_2:k_2$	$s_1, s_2$	$1 + (s_1 - j_1) + (s_2 - j_2) \times d_1$
3	$j_1:k_1, j_2:k_2, j_3:k_3$	$s_1, s_2, s_3$	$1 + (s_1 - j_1) + (s_2 - j_2) \times d_1 + (s_3 - j_3) \times d_2 \times d_1$
⋮	⋮	⋮	⋮
⋮	⋮	⋮	⋮

**Subscript order value** (cont.)

Rank	Subscript bounds	Subscript list	Subscript order value
.	.	.	.
15	$j_1:k_1, \dots, j_{15}:k_{15}$	$s_1, \dots, s_{15}$	$1 + (s_1 - j_1)$ $+ (s_2 - j_2) \times d_1$ $+ (s_3 - j_3) \times d_2 \times d_1$ $+ \dots$ $+ (s_{15} - j_{15}) \times d_{14}$ $\times d_{13} \times \dots \times d_1$
Notes for Table 9.1: 1) $d_i = \max(k_i - j_i + 1, 0)$ is the size of the $i$ th dimension. 2) If the size of the array is nonzero, $j_i \leq s_i \leq k_i$ for all $i = 1, 2, \dots, 15$ .			

1 **9.5.3.3 Array sections**

- 2 1 In an *array-section* having a *section-subscript-list*, each *subscript-triplet* and *vector-subscript* in the section sub-
- 3 script list indicates a sequence of subscripts, which may be empty. Each subscript in such a sequence shall be
- 4 within the bounds for its dimension unless the sequence is empty. The *array section* is the set of elements from
- 5 the array determined by all possible subscript lists obtainable from the single subscripts or sequences of subscripts
- 6 specified by each section subscript.
- 7 2 In an *array-section* with no *section-subscript-list*, the *rank* and shape of the array is the *rank* and shape of the
- 8 *part-ref* with nonzero *rank*; otherwise, the *rank* of the *array section* is the number of subscript triplets and *vector*
- 9 *subscripts* in the section subscript list. The shape is the rank-one array whose  $i$ th element is the number of
- 10 integer values in the sequence indicated by the  $i$ th subscript triplet or *vector subscript*. If any of these sequences
- 11 is empty, the *array section* has size zero. The subscript order of the elements of an *array section* is that of the
- 12 array data object that the *array section* represents.

13 **9.5.3.3.1 Subscript triplet**

- 14 1 A subscript triplet designates a regular sequence of subscripts consisting of zero or more subscript values. The
- 15 *stride* in the subscript triplet specifies the increment between the subscript values. The subscripts and stride of a
- 16 subscript triplet are optional. An omitted first subscript in a subscript triplet is equivalent to a subscript whose
- 17 value is the lower bound for the array and an omitted second subscript is equivalent to the upper bound. An
- 18 omitted stride is equivalent to a stride of 1.
- 19 2 The stride shall not be zero.
- 20 3 When the stride is positive, the subscripts specified by a triplet form a regularly spaced sequence of integers
- 21 beginning with the first subscript and proceeding in increments of the stride to the largest such integer not
- 22 greater than the second subscript; the sequence is empty if the first subscript is greater than the second.

**NOTE 9.12**

For example, suppose an array is declared as A (5, 4, 3). The section A (3 : 5, 2, 1 : 2) is the array of shape (3, 2):

A (3, 2, 1)	A (3, 2, 2)
A (4, 2, 1)	A (4, 2, 2)
A (5, 2, 1)	A (5, 2, 2)

- 23 4 When the stride is negative, the sequence begins with the first subscript and proceeds in increments of the stride
- 24 down to the smallest such integer equal to or greater than the second subscript; the sequence is empty if the
- 25 second subscript is greater than the first.

**NOTE 9.13**

For example, if an array is declared B (10), the section B (9 : 1 : -2) is the array of shape (5) whose elements are B (9), B (7), B (5), B (3), and B (1), in that order.

**NOTE 9.14**

A subscript in a subscript triplet need not be within the declared bounds for that dimension if all values used in selecting the array elements are within the declared bounds.

For example, if an array is declared as B (10), the **array section** B (3 : 11 : 7) is the array of shape (2) consisting of the elements B (3) and B (10), in that order.

1 **9.5.3.3.2 Vector subscript**

- 2 1 A **vector subscript** designates a sequence of subscripts corresponding to the values of the elements of the expression.  
 3 Each element of the expression shall be defined.
- 4 2 An **array section** with a **vector subscript** shall not be **finalized** by a nonelemental **final subroutine**.
- 5 3 If a **vector subscript** has two or more elements with the same value, an **array section** with that **vector subscript**  
 6 is not **definable** and shall not be **defined** or become **undefined**.

**NOTE 9.15**

For example, suppose Z is a two-dimensional array of shape [5, 7] and U and V are one-dimensional arrays of shape (3) and (4), respectively. Assume the values of U and V are:

U = [ 1, 3, 2 ]

V = [ 2, 1, 1, 3 ]

Then Z (3, V) consists of elements from the third row of Z in the order:

Z (3, 2) Z (3, 1) Z (3, 1) Z (3, 3)

and Z (U, 2) consists of the column elements:

Z (1, 2) Z (3, 2) Z (2, 2)

and Z (U, V) consists of the elements:

Z (1, 2) Z (1, 1) Z (1, 1) Z (1, 3)

Z (3, 2) Z (3, 1) Z (3, 1) Z (3, 3)

Z (2, 2) Z (2, 1) Z (2, 1) Z (2, 3)

Because Z (3, V) and Z (U, V) contain duplicate elements from Z, the sections Z (3, V) and Z (U, V) shall not be redefined as sections.

7 **9.5.4 Simply contiguous array designators**

- 8 1 A **section-subscript-list** specifies a **simply contiguous** section if and only if it does not have a **vector subscript** and  
 9 • all but the last **subscript-triplet** is a colon,  
 10 • the last **subscript-triplet** does not have a **stride**, and  
 11 • no **subscript-triplet** is preceded by a **section-subscript** that is a **subscript**.
- 12 2 An array designator is **simply contiguous** if and only if it is  
 13 • an **object-name** that has the **CONTIGUOUS** attribute,

- 1 • an *object-name* that is not a *pointer*, not *assumed-shape*, and not *assumed-rank*,
  - 2 • a *structure-component* whose final *part-name* is an array and that either has the *CONTIGUOUS* attribute
  - 3 or is not a pointer, or
  - 4 • an *array section*
    - 5 – that is not a *complex-part-designator*,
    - 6 – that does not have a *substring-range*,
    - 7 – whose final *part-ref* has nonzero *rank*,
    - 8 – whose rightmost *part-name* has the *CONTIGUOUS* attribute or is neither *assumed-shape* nor a pointer,
    - 9 and
    - 10 – which either does not have a *section-subscript-list*, or has a *section-subscript-list* which specifies a
    - 11 *simply contiguous* section.
- 12 3 An array *variable* is *simply contiguous* if and only if it is a *simply contiguous* array designator or a reference to
- 13 a function that returns a pointer with the *CONTIGUOUS* attribute.

**NOTE 9.16**

Array sections that are *simply contiguous* include column, plane, cube, and hypercube subobjects of a *simply contiguous* base object, for example:

```

ARRAY1 (10:20, 3)    ! passes part of the third column of ARRAY1.
X3D (:, i:j, 2)     ! passes part of the second plane of X3D (or the whole
                   ! plane if i==LBOUND(X3D,2) and j==UBOUND(X3D,2)).
Y5D (:, :, :, :, 7) ! passes the seventh hypercube of Y5D.

```

All *simply contiguous* designators designate *contiguous* objects.

## 14 9.6 Image selectors

- 15 1 An image selector determines the *image index* for a *coindexed* object.
- 16 R924 *image-selector* is *lbracket cosubscript-list* [ , *image-selector-spec-list* ] *rbracket*
- 17 R925 *cosubscript* is *scalar-int-expr*
- 18 R926 *image-selector-spec* is STAT = *stat-variable*
- 19 or TEAM = *team-value*
- 20 or TEAM\_NUMBER = *scalar-int-expr*
- 21 C928 No specifier shall appear more than once in a given *image-selector-spec-list*.
- 22 C929 TEAM and TEAM\_NUMBER shall not both appear in the same *image-selector-spec-list*.
- 23 C930 A *stat-variable* in an *image-selector* shall not be a *coindexed* object.
- 24 2 The number of *cosubscripts* shall be equal to the *corank* of the object. The value of a *cosubscript* in an image
- 25 selector shall be within the *cobounds* for its *codimension*. Taking account of the *cobounds*, the *cosubscript* list in
- 26 an image selector determines the *image index* in the same way that a subscript list in an array element determines
- 27 the subscript order value (9.5.3.2), taking account of the bounds.
- 28 3 If a *TEAM= specifier* appears in an *image-selector*, the *team* of the image selector is specified by *team-value*,
- 29 which shall identify the *current* or an *ancestor team*; the object shall be an *established coarray* in that *team* or
- 30 an *ancestor* thereof. If a *TEAM\_NUMBER= specifier* appears in an *image-selector* and the current team is not
- 31 the *initial team*, the value of the *scalar-int-expr* shall be equal to the value of a *team number* for one of the
- 32 teams that were formed by execution of the *FORM TEAM* statement for the *current team* and the team of the
- 33 image selector is that team; the object shall be an *established coarray* in an *ancestor* of the *current team*. If a

- 1 *TEAM\_NUMBER= specifier* appears in an *image-selector* and the *current team* is the *initial team*, the value of  
 2 *scalar-int-expr* shall be the *team number* for the *initial team*; the object shall be an *established coarray* in the  
 3 *initial team*. Otherwise, the *team* of the image selector is the *current team*.
- 4 4 An image selector shall specify an *image index* value that is not greater than the number of *images* in the *team*  
 5 of the image selector, and identifies the image with that index in that *team*.
- 6 5 Execution of a statement containing an *image-selector* with a *STAT= specifier* causes the *stat-variable* to become  
 7 defined. If the designator is part of an operand that is evaluated or is a variable that is being defined or partly  
 8 defined, and the object designated is on a *failed image*, the *stat-variable* is defined with the value *STAT\_-*  
 9 *FAILED\_IMAGE* (16.10.2.26) in the intrinsic module *ISO\_FORTRAN\_ENV*; otherwise, it is defined with the  
 10 value zero.
- 11 6 The denotation of a *stat-variable* in an *image-selector* shall not depend on the evaluation of any entity in the  
 12 same statement. The value of an expression shall not depend on the value of any *stat-variable* that appears in  
 13 the same statement. The value of a *stat-variable* in an *image-selector* shall not be affected by the execution of  
 14 any part of the statement, other than by whether the image specified by the *image-selector* has failed.

**NOTE 9.17**

For example, if there are 16 *images* and the *coarray* A is declared

```
REAL :: A(10)[5,*]
```

A(:)[1,4] is valid because it specifies *image* 16, but A(:)[2,4] is invalid because it specifies *image* 17.

**9.7 Dynamic association****9.7.1 ALLOCATE statement****9.7.1.1 Form of the ALLOCATE statement**

- 1 The ALLOCATE statement dynamically creates pointer *targets* and *allocatable* variables.

19	R927	<i>allocate-stmt</i>	is	ALLOCATE ( [ <i>type-spec</i> :: ] <i>allocation-list</i> ■ 20 ■ [ , <i>alloc-opt-list</i> ] )
21	R928	<i>alloc-opt</i>	is	ERRMSG = <i>errmsg-variable</i> 22 or MOLD = <i>source-expr</i> 23 or SOURCE = <i>source-expr</i> 24 or STAT = <i>stat-variable</i>
25	R929	<i>stat-variable</i>	is	<i>scalar-int-variable</i>
26	R930	<i>errmsg-variable</i>	is	<i>scalar-default-char-variable</i>
27	R931	<i>source-expr</i>	is	<i>expr</i>
28	R932	<i>allocation</i>	is	<i>allocate-object</i> [ ( <i>allocate-shape-spec-list</i> ) ] ■ 29 ■ [ <i>lbracket allocate-coarray-spec rbracket</i> ]
30	R933	<i>allocate-object</i>	is	<i>variable-name</i> 31 or <i>structure-component</i>
32	R934	<i>allocate-shape-spec</i>	is	[ <i>lower-bound-expr</i> : ] <i>upper-bound-expr</i>
33	R935	<i>lower-bound-expr</i>	is	<i>scalar-int-expr</i>
34	R936	<i>upper-bound-expr</i>	is	<i>scalar-int-expr</i>



- 1 R937 *allocate-coarray-spec* is [ *allocate-coshape-spec-list* , ] [ *lower-bound-expr* : ] \*
- 2 R938 *allocate-coshape-spec* is [ *lower-bound-expr* : ] *upper-bound-expr*
- 3 C931 (R933) Each *allocate-object* shall be a data pointer or an *allocatable* variable.
- 4 C932 (R927) If any *allocate-object* has a *deferred type* parameter, is *unlimited polymorphic*, or is of *abstract*  
5 *type*, either *type-spec* or *source-expr* shall appear.
- 6 C933 (R927) If *type-spec* appears, it shall specify a type with which each *allocate-object* is *type compatible*.
- 7 C934 (R927) A *type-param-value* in a *type-spec* shall be an asterisk if and only if each *allocate-object* is a *dummy*  
8 *argument* for which the corresponding type parameter is assumed.
- 9 C935 (R927) If *type-spec* appears, the kind type parameter values of each *allocate-object* shall be the same as  
10 the corresponding type parameter values of the *type-spec*.
- 11 C936 (R927) *type-spec* shall not specify a type that has a *coarray ultimate component*.
- 12 C937 (R927) If an *allocate-object* is a *coarray*, *type-spec* shall not specify the type *C\_PTR* or *C\_FUNPTR*  
13 from the intrinsic module *ISO\_C\_BINDING*.
- 14 C938 (R927) If an *allocate-object* is an array, either *allocate-shape-spec-list* shall appear in its *allocation*, or  
15 *source-expr* shall appear in the *ALLOCATE* statement and have the same *rank* as the *allocate-object*.
- 16 C939 (R932) If *allocate-object* is scalar, *allocate-shape-spec-list* shall not appear.
- 17 C940 (R932) An *allocate-coarray-spec* shall appear if and only if the *allocate-object* is a *coarray*.
- 18 C941 (R932) The number of *allocate-shape-specs* in an *allocate-shape-spec-list* shall be the same as the *rank*  
19 of the *allocate-object*. The number of *allocate-coshape-specs* in an *allocate-coarray-spec* shall be one less  
20 than the *corank* of the *allocate-object*.
- 21 C942 (R928) No *alloc-opt* shall appear more than once in a given *alloc-opt-list*.
- 22 C943 (R927) At most one of *source-expr* and *type-spec* shall appear.
- 23 C944 (R927) Each *allocate-object* shall be *type compatible* (7.3.2.3) with *source-expr*. If *SOURCE=* appears,  
24 *source-expr* shall be a scalar or have the same *rank* as each *allocate-object*.
- 25 C945 (R927) If *source-expr* appears, the *kind type parameters* of each *allocate-object* shall have the same values  
26 as the corresponding *type parameters* of *source-expr*.
- 27 C946 (R927) The *declared type* of *source-expr* shall not be *C\_PTR* or *C\_FUNPTR* from the intrinsic module  
28 *ISO\_C\_BINDING* if an *allocate-object* is a *coarray*.
- 29 C947 (R927) If *SOURCE=* appears, the *declared type* of *source-expr* shall not be *EVENT\_TYPE* or *LOCK\_-*  
30 *TYPE* from the intrinsic module *ISO\_FORTRAN\_ENV*, or have a *potential subobject component* of  
31 type *EVENT\_TYPE* or *LOCK\_TYPE*.
- 32 C948 (R931) The *declared type* of *source-expr* shall not have a *coarray ultimate component*.
- 33 C949 (R933) An *allocate-object* shall not be a *coindexed object*.

**NOTE 9.18**

If a *coarray* is of a derived type that has an *allocatable* component, the component shall be allocated by its own *image*:

```
TYPE(SOMETHING), ALLOCATABLE :: T[:]
...
```

## NOTE 9.18 (cont.)

ALLOCATE(T[*])	! Allowed - implies synchronization
ALLOCATE(T% <i>AAC</i> (N))	! Allowed - allocated by its own <i>image</i>
ALLOCATE(T[Q]% <i>AAC</i> (N))	! Not allowed, because it is not necessarily executed on <i>image</i> Q.

- 1 2 An *allocate-object* or a bound or type parameter of an *allocate-object* shall not depend on the value of *stat-variable*,  
 2 the value of *errmsg-variable*, or on the value, bounds, *length type parameters*, allocation status, or association  
 3 status of any *allocate-object* in the same ALLOCATE statement.
- 4 3 *source-expr* shall not be allocated within the ALLOCATE statement in which it appears; nor shall it depend on  
 5 the value, *bounds*, *deferred type parameters*, allocation status, or association status of any *allocate-object* in that  
 6 statement.
- 7 4 If an ALLOCATE statement has a SOURCE= specifier and an *allocate-object* that is a *coarray*, *source-expr*  
 8 shall not have a *dynamic type* of C\_PTR or C\_FUNPTR from the intrinsic module ISO\_C\_BINDING, or  
 9 EVENT\_TYPE or LOCK\_TYPE from the intrinsic module ISO\_FORTRAN\_ENV, or have a *subcomponent*  
 10 whose *dynamic type* is EVENT\_TYPE or LOCK\_TYPE.
- 11 5 If *type-spec* is specified, each *allocate-object* is allocated with the specified *dynamic type* and type parameter  
 12 values; if *source-expr* is specified, each *allocate-object* is allocated with the *dynamic type* and type parameter  
 13 values of *source-expr*; otherwise, each *allocate-object* is allocated with its *dynamic type* the same as its *declared*  
 14 *type*.
- 15 6 If a *type-param-value* in a *type-spec* in an ALLOCATE statement is an asterisk, it denotes the current value of  
 16 that assumed type parameter. If it is an expression, subsequent redefinition or undefinition of any entity in the  
 17 expression does not affect the type parameter value.

## NOTE 9.19

An example of an ALLOCATE statement is:

```
ALLOCATE (X (N), B (-3 : M, 0:9), STAT = IERR_ALLOC)
```

## 18 9.7.1.2 Execution of an ALLOCATE statement

- 19 1 When an ALLOCATE statement is executed for an array for which *allocate-shape-spec-list* is specified, the values  
 20 of the lower bound and upper bound expressions determine the bounds of the array. Subsequent redefinition  
 21 or undefinition of any entities in the bound expressions do not affect the array bounds. If the lower bound is  
 22 omitted, the default value is 1. If the upper bound is less than the lower bound, the extent in that dimension is  
 23 zero and the array has zero size.
- 24 2 When an ALLOCATE statement is executed for a *coarray*, the values of the lower *cobound* and upper *cobound*  
 25 expressions determine the *cobounds* of the *coarray*. Subsequent redefinition or undefinition of any entities in the  
 26 *cobound* expressions do not affect the *cobounds*. If the lower *cobound* is omitted, the default value is 1. The  
 27 upper *cobound* shall not be less than the lower *cobound*.
- 28 3 If an *allocation* specifies a *coarray*, its *dynamic type* and the values of corresponding type parameters shall be the  
 29 same on every *active image* in the *current team*. The values of corresponding bounds and corresponding *cobounds*  
 30 shall be the same on those *images*. If the *coarray* is a *dummy argument*, its *ultimate argument* (15.5.2.3) shall  
 31 be the same *coarray* on those *images*.
- 32 4 When an ALLOCATE statement is executed for which an *allocate-object* is a *coarray*, there is an implicit syn-  
 33 chronization of all *active images* in the *current team*. If the *current team* contains a *stopped* or *failed* image,  
 34 an error condition occurs. If no other error condition occurs, execution on the *active images* of the segment  
 35 (11.6.2) following the statement is delayed until all other *active images* in the *current team* have executed the  
 36 same statement the same number of times in this team. The *coarray* shall not become allocated on an *image*

1 unless it is successfully allocated on all **active images** in this team.

#### NOTE 9.20

When an **image** executes an ALLOCATE statement, communication is not necessarily involved apart from any required for synchronization. The **image** allocates its **coarray** and records how the corresponding **coarrays** on other **images** are to be addressed. The processor is not required to detect violations of the rule that the bounds are the same on all **images** of the **current team**, nor is it responsible for detecting or resolving deadlock problems (such as two **images** waiting on different **ALLOCATE statements**).

- 2 5 If *source-expr* is a pointer, it shall be associated with a **target**. If *source-expr* is **allocatable**, it shall be allocated.
- 3 6 When an ALLOCATE statement is executed for an array with no *allocate-shape-spec-list*, the bounds of *source-expr* determine the bounds of the array. Subsequent changes to the bounds of *source-expr* do not affect the array bounds.
- 4 5
- 6 7 If SOURCE= appears, *source-expr* shall be **conformable** with *allocation*. If an *allocate-object* is not **polymorphic** and the *source-expr* is **polymorphic** with a **dynamic type** that differs from its **declared type**, the value provided for that *allocate-object* is the ancestor component of the *source-expr* that has the type of the *allocate-object*; otherwise the value provided is the value of the *source-expr*. On successful allocation, if *allocate-object* and *source-expr* have the same **rank** the value of *allocate-object* becomes the value provided, otherwise the value of each element of *allocate-object* becomes the value provided. The *source-expr* is evaluated exactly once for each execution of an **ALLOCATE statement**.
- 7 8
- 8 8 If MOLD= appears and *source-expr* is a variable, its value need not be defined.
- 9 9 If *type-spec* appears and the value of a **length type parameter** it specifies differs from the value of the corresponding nondeferred type parameter specified in the declaration of any *allocate-object*, an error condition occurs. If the value of a nondeferred **length type parameter** of an *allocate-object* differs from the value of the corresponding type parameter of *source-expr*, an error condition occurs.
- 10 10 The set of error conditions for an ALLOCATE statement is processor dependent. If an error condition occurs during execution of an ALLOCATE statement that does not contain the **STAT= specifier**, **error termination** is initiated. The **STAT= specifier** is described in 9.7.4. The **ERRMSG= specifier** is described in 9.7.5.
- 11 12

#### 21 9.7.1.3 Allocation of allocatable variables

- 22 1 The allocation status of an **allocatable** entity is one of the following at any time.
- 23
- 24 • The status of an **allocatable** variable becomes “allocated” if it is allocated by an ALLOCATE statement, if it is allocated during assignment, or if it is given that status by the intrinsic subroutine **MOVE\_ALLOC** (16.9.137). An **allocatable** variable with this status may be referenced, defined, or deallocated; allocating it causes an error condition in the ALLOCATE statement. The result of the intrinsic function **ALLOCATED** (16.9.11) is true for such a variable.
  - 25
  - 26 • An **allocatable** variable has a status of “unallocated” if it is not allocated. The status of an **allocatable** variable becomes unallocated if it is deallocated (9.7.3) or if it is given that status by the intrinsic subroutine **MOVE\_ALLOC**. An **allocatable** variable with this status shall not be referenced or defined. It shall not be supplied as an **actual argument** corresponding to a nonallocatable nonoptional **dummy argument**, except to certain intrinsic **inquiry functions**. It may be allocated with the ALLOCATE statement. Deallocating it causes an error condition in the DEALLOCATE statement. The result of the intrinsic function **ALLOCATED** (16.9.11) is false for such a variable.
  - 27
  - 28
  - 29
  - 30
  - 31
  - 32
  - 33
  - 34
- 35 2 At the beginning of execution of a program, **allocatable** variables are unallocated.
- 36 3 When the allocation status of an **allocatable** variable changes, the allocation status of any associated **allocatable** variable changes accordingly. Allocation of an **allocatable** variable establishes values for the **deferred type parameters** of all associated **allocatable** variables.
- 37
- 38

- 1 4 An **unsaved allocatable local variable** of a procedure has a status of unallocated at the beginning of each invocation  
 2 of the procedure. An **unsaved allocatable local variable** of a construct has a status of unallocated at the beginning  
 3 of each execution of the construct.
- 4 5 When an object of derived type is created by an ALLOCATE statement, any **allocatable ultimate components**  
 5 have an allocation status of unallocated unless the **SOURCE= specifier** appears and the corresponding component  
 6 of the *source-expr* is allocated.
- 7 6 If the evaluation of a function would change the allocation status of a variable and if a reference to the function  
 8 appears in an expression in which the value of the function is not needed to determine the value of the expression,  
 9 the allocation status of the variable after evaluation of the expression is processor dependent.

#### 10 9.7.1.4 Allocation of pointer targets

- 11 1 Allocation of a **pointer** creates an object that implicitly has the **TARGET attribute**. Following successful execution  
 12 of an ALLOCATE statement for a **pointer**, the **pointer** is associated with the **target** and may be used to reference  
 13 or define the **target**. Additional **pointers** may become associated with the **pointer target** or a part of the **pointer**  
 14 **target** by **pointer assignment**. It is not an error to allocate a **pointer** that is already associated with a **target**.  
 15 In this case, a new **pointer target** is created as required by the attributes of the **pointer** and any array bounds,  
 16 type, and type parameters specified by the ALLOCATE statement. The **pointer** is then associated with this  
 17 new **target**. Any previous association of the **pointer** with a **target** is broken. If the previous **target** had been  
 18 created by allocation, it becomes inaccessible unless other **pointers** are associated with it. The intrinsic function  
 19 **ASSOCIATED** (16.9.16) may be used to determine whether a **pointer** that does not have undefined association  
 20 status is associated.
- 21 2 At the beginning of execution of a function whose result is a **pointer**, the association status of the result **pointer**  
 22 is undefined. Before such a function returns, it shall either associate a **target** with this **pointer** or cause the  
 23 association status of this **pointer** to become **disassociated**.

#### 24 9.7.2 NULLIFY statement

25 R939 *nullify-stmt* is NULLIFY ( *pointer-object-list* )

26 R940 *pointer-object* is *variable-name*

27 or *structure-component*

28 or *proc-pointer-name*

29 C950 (R940) Each *pointer-object* shall have the **POINTER attribute**.

- 30 1 A *pointer-object* shall not depend on the value, bounds, or association status of another *pointer-object* in the  
 31 same NULLIFY statement.
- 32 2 Execution of a NULLIFY statement causes each *pointer-object* to become **disassociated**.

##### NOTE 9.21

When a NULLIFY statement is applied to a **polymorphic pointer** (7.3.2.3), its **dynamic type** becomes the same as its **declared type**.

#### 33 9.7.3 DEALLOCATE statement

##### 34 9.7.3.1 Form of the DEALLOCATE statement

- 35 1 The DEALLOCATE statement causes **allocatable** variables to be deallocated; it causes pointer **targets** to be  
 36 deallocated and the pointers to be **disassociated**.

37 R941 *deallocate-stmt* is DEALLOCATE ( *allocate-object-list* [ , *dealloc-opt-list* ] )

- 1 R942 *dealloc-opt* is STAT = *stat-variable*  
 2 or ERRMSG = *errmsg-variable*
- 3 C951 (R942) No *dealloc-opt* shall appear more than once in a given *dealloc-opt-list*.
- 4 2 An *allocate-object* shall not depend on the value, bounds, allocation status, or association status of another  
 5 *allocate-object* in the same DEALLOCATE statement; it also shall not depend on the value of the *stat-variable*  
 6 or *errmsg-variable* in the same DEALLOCATE statement.
- 7 3 The set of error conditions for a DEALLOCATE statement is processor dependent. If an error condition occurs  
 8 during execution of a DEALLOCATE statement that does not contain the STAT= specifier, error termination is  
 9 initiated. The STAT= specifier is described in 9.7.4. The ERRMSG= specifier is described in 9.7.5.
- 10 4 When more than one allocated object is deallocated by execution of a DEALLOCATE statement, the order of  
 11 deallocation is processor dependent.

**NOTE 9.22**

An example of a DEALLOCATE statement is:

```
DEALLOCATE (X, B)
```

**9.7.3.2 Deallocation of allocatable variables**

- 12
- 13 1 Deallocating an unallocated *allocatable* variable causes an error condition in the DEALLOCATE statement.  
 14 Deallocating an *allocatable* variable with the TARGET attribute causes the *pointer association* status of any  
 15 *pointer* associated with it to become *undefined*. An *allocatable* variable shall not be deallocated if it or any  
 16 subobject of it is *argument associated* with a *dummy argument* or *construct associated* with an *associate name*.
- 17 2 When the execution of a procedure is terminated by execution of a RETURN or END statement, an *unsaved*  
 18 *allocatable local variable* of the procedure retains its allocation and definition status if it is a *function result* or a  
 19 subobject thereof; otherwise, if it is allocated it will be deallocated.
- 20 3 When a BLOCK construct terminates, any *unsaved* allocated *allocatable local variable* of the construct is deal-  
 21 located.
- 22 4 If an executable construct references a function whose result is *allocatable* or has an *allocatable* subobject, and  
 23 the function reference is executed, an *allocatable* result and any allocated *allocatable* subobject of the result is  
 24 deallocated after execution of the innermost executable construct containing the reference.
- 25 5 If a function whose result is *allocatable* or has an *allocatable* subobject is referenced in the specification part of a  
 26 *scoping unit*, and the function reference is executed, an *allocatable* result and any allocated *allocatable* subobject  
 27 of the result is deallocated before execution of the executable constructs of the *scoping unit*.
- 28 6 When a procedure is invoked, any allocated *allocatable* object that is an *actual argument* corresponding to an  
 29 INTENT (OUT) *allocatable dummy argument* is deallocated; any allocated *allocatable* object that is a subobject  
 30 of an *actual argument* corresponding to an INTENT (OUT) *dummy argument* is deallocated. If a Fortran proced-  
 31 ure that has an INTENT (OUT) *allocatable dummy argument* is invoked by a C function and the corresponding  
 32 argument in the C function call is a C descriptor that describes an allocated *allocatable* variable, the variable  
 33 is deallocated on entry to the Fortran procedure. If a C function is invoked from a Fortran procedure via an  
 34 interface with an INTENT (OUT) *allocatable dummy argument* and the corresponding *actual argument* in the  
 35 reference to the C function is an allocated *allocatable* variable, the variable is deallocated on invocation (before  
 36 execution of the C function begins).
- 37 7 When an *intrinsic assignment statement* (10.2.1.3) is executed, any noncoarray allocated *allocatable* subobject of  
 38 the variable is deallocated before the assignment takes place.
- 39 8 When a variable of derived type is deallocated, any allocated *allocatable* subobject is deallocated. If an error  
 40 condition occurs during deallocation, it is *processor dependent* whether an allocated *allocatable subobject* is

- 1 deallocated.
- 2 9 If an [allocatable](#) component is a subobject of a [finalizable](#) object, that object is finalized before the component is  
3 automatically deallocated.
- 4 10 When a statement that deallocates a [coarray](#) is executed, there is an implicit synchronization of all [active images](#)  
5 in the [current team](#). If the [current team](#) contains a [stopped](#) or [failed](#) image, an error condition occurs. If no other  
6 error condition occurs, execution on the [active images](#) of the segment (11.6.2) following the statement is delayed  
7 until all other [active images](#) in the [current team](#) have executed the same statement the same number of times in  
8 this [team](#). A [coarray](#) shall not become deallocated on an [image](#) unless it is successfully deallocated on all [active](#)  
9 [images](#) in this [team](#).
- 10 11 If an [allocate-object](#) is a [coarray dummy](#) argument, its [ultimate argument](#) (15.5.2.3) shall be the same [coarray](#) on  
11 those [images](#).
- 12 12 The effect of automatic deallocation is the same as that of a DEALLOCATE statement without a [dealloc-opt-list](#).

**NOTE 9.23**

In the following example:

```
SUBROUTINE PROCESS
  REAL, ALLOCATABLE :: TEMP(:)
  REAL, ALLOCATABLE, SAVE :: X(:)
  ...
END SUBROUTINE PROCESS
```

on return from subroutine PROCESS, the allocation status of X is preserved because X has the [SAVE attribute](#). TEMP does not have the [SAVE attribute](#), so it will be deallocated if it was allocated. On the next invocation of PROCESS, TEMP will have an allocation status of unallocated.

**NOTE 9.24**

For example, executing a [RETURN](#), [END](#), or [END BLOCK](#) statement, or deallocating an object that has an allocatable subobject, can cause deallocation of a [coarray](#), and thus an implicit synchronization of all [active images](#) in the [current team](#).

**9.7.3.3 Deallocation of pointer targets**

- 13
- 14 1 If a pointer appears in a DEALLOCATE statement, its association status shall be defined. Deallocating a pointer  
15 that is [disassociated](#) or whose [target](#) was not created by an ALLOCATE statement causes an error condition  
16 in the DEALLOCATE statement. If a pointer is associated with an [allocatable](#) entity, the pointer shall not be  
17 deallocated. A [pointer](#) shall not be deallocated if its [target](#) or any [subobject](#) thereof is [argument associated](#) with  
18 a [dummy argument](#) or [construct associated](#) with an [associate name](#).
- 19 2 If a pointer appears in a DEALLOCATE statement, it shall be associated with the whole of an object that was  
20 created by allocation. The pointer shall have the same [dynamic type](#) and type parameters as the allocated object,  
21 and if the allocated object is an array the pointer shall be an array whose elements are the same as those of the  
22 allocated object in array element order. Deallocating a pointer [target](#) causes the pointer association status of any  
23 other pointer that is associated with the [target](#) or a portion of the [target](#) to become undefined.

**9.7.4 STAT= specifier**

- 24
- 25 1 This subclause applies where an [alloc-opt](#) or [dealloc-opt](#) that is a STAT= specifier appears in an [ALLOCATE](#) or  
26 [DEALLOCATE](#) statement.
- 27 2 The [stat-variable](#) shall not be allocated or deallocated within the [ALLOCATE](#) or [DEALLOCATE](#) statement  
28 in which it appears; nor shall it depend on the value, [bounds](#), [deferred type parameters](#), allocation status, or

- 1 association status of any *allocate-object* in that statement. The *stat-variable* shall not depend on the value of the  
2 *errmsg-variable*.
- 3 3 Successful execution of the **ALLOCATE** or **DEALLOCATE** statement causes the *stat-variable* to become defined  
4 with a value of zero.
- 5 4 If an **ALLOCATE** or **DEALLOCATE** statement with a *coarray allocate-object* is executed when the *current team*  
6 contains a *stopped image*, the *stat-variable* becomes defined with the value **STAT\_STOPPED\_IMAGE** from the  
7 intrinsic module **ISO\_FORTRAN\_ENV** (16.10.2). Otherwise, if an *allocate-object* is a *coarray*, the *current team*  
8 contains a *failed image*, and no other error condition occurs, the *stat-variable* becomes defined with value **STAT\_-**  
9 **FAILED\_IMAGE** from the intrinsic module **ISO\_FORTRAN\_ENV**. If any other error condition occurs during  
10 execution of the **ALLOCATE** or **DEALLOCATE** statement, the *stat-variable* becomes defined with a processor-  
11 dependent positive integer value different from **STAT\_STOPPED\_IMAGE** and **STAT\_FAILED\_IMAGE**.
- 12 5 If *stat-variable* became defined with the value **STAT\_FAILED\_IMAGE**, each *allocate-object* is successfully al-  
13 located or deallocated on all the *active images* of the *current team*. If any other error condition occurs, each  
14 *allocate-object* has a processor-dependent status:
- 15 • each *allocate-object* that was successfully allocated shall have an allocation status of allocated or a *pointer*  
16 *association* status of associated;
  - 17 • each *allocate-object* that was successfully deallocated shall have an allocation status of unallocated or a  
18 *pointer association* status of *disassociated*;
  - 19 • each *allocate-object* that was not successfully allocated or deallocated shall retain its previous allocation  
20 status or *pointer association* status.

**NOTE 9.25**

The status of objects that were not successfully allocated or deallocated can be individually checked with the intrinsic functions **ALLOCATED** or **ASSOCIATED**.

**9.7.5 ERRMSG= specifier**

- 21
- 22 1 The *errmsg-variable* shall not be allocated or deallocated within the **ALLOCATE** or **DEALLOCATE** statement  
23 in which it appears; nor shall it depend on the value, *bounds*, *deferred type parameters*, allocation status, or  
24 association status of any *allocate-object* in that statement. The *errmsg-variable* shall not depend on the value of  
25 the *stat-variable*.
- 26 2 If an error condition occurs during execution of an **ALLOCATE** or **DEALLOCATE** statement with an **ERRMSG=**  
27 specifier, the *errmsg-variable* is assigned an explanatory message, as if by *intrinsic assignment*. If no such condition  
28 occurs, the definition status and value of *errmsg-variable* are unchanged.

1

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# 10 Expressions and assignment

## 10.1 Expressions

### 10.1.1 Expression semantics

- 1 An expression represents either a data object reference or a computation, and its value is either a scalar or an array. Evaluation of an expression produces a value, which has a type, type parameters (if appropriate), and a shape (10.1.9). The *corank* of an expression that is not a variable is zero.

### 10.1.2 Form of an expression

#### 10.1.2.1 Overall expression syntax

- 1 An expression is formed from operands, operators, and parentheses. An operand is either a scalar or an array. An operation is either intrinsic (10.1.5) or defined (10.1.6). More complicated expressions can be formed using operands which are themselves expressions.
- 2 An expression is defined in terms of several categories: primary, level-1 expression, level-2 expression, level-3 expression, level-4 expression, and level-5 expression.
- 3 These categories are related to the different operator precedence levels and, in general, are defined in terms of other categories. The simplest form of each expression category is a *primary*.

#### 10.1.2.2 Primary

R1001 *primary* is *literal-constant*  
 or *designator*  
 or *array-constructor*  
 or *structure-constructor*  
 or *function-reference*  
 or *type-param-inquiry*  
 or *type-param-name*  
 or ( *expr* )

C1001 (R1001) The *type-param-name* shall be the name of a type parameter.

C1002 (R1001) The *designator* shall not be a whole *assumed-size array*.

C1003 (R1001) The *expr* shall not be a function reference that returns a procedure pointer.

#### NOTE 10.1

Examples of a *primary* are:

<u>Example</u>	<u>Syntactic class</u>
1.0	<i>constant</i>
'ABCDEFGHIJKLMNQRSTUWXYZ' (I:I)	<i>designator</i>
[ 1.0, 2.0 ]	<i>array-constructor</i>
PERSON ('Jones', 12)	<i>structure-constructor</i>
F (X, Y)	<i>function-reference</i>
X%KIND	<i>type-param-inquiry</i>
KIND	<i>type-param-name</i>
(S + T)	( <i>expr</i> )

### 1 10.1.2.3 Level-1 expressions

2 1 Defined unary operators have the highest operator precedence (Table 10.1). Level-1 expressions are primaries  
3 optionally operated on by defined unary operators:

4 R1002 *level-1-expr* is [ *defined-unary-op* ] *primary*

5 R1003 *defined-unary-op* is . *letter* [ *letter* ] ... .

6 C1004 (R1003) A *defined-unary-op* shall not contain more than 63 letters and shall not be the same as any  
7 *intrinsic-operator* or *logical-literal-constant*.

#### NOTE 10.2

Simple examples of a level-1 expression are:

<u>Example</u>	<u>Syntactic class</u>
A	<i>primary</i> (R1001)
. INVERSE. B	<i>level-1-expr</i> (R1002)

A more complicated example of a level-1 expression is:

. INVERSE. (A + B)

### 8 10.1.2.4 Level-2 expressions

9 1 Level-2 expressions are level-1 expressions optionally involving the numeric operators *power-op*, *mult-op*, and  
10 *add-op*.

11 R1004 *mult-operand* is *level-1-expr* [ *power-op mult-operand* ]

12 R1005 *add-operand* is [ *add-operand mult-op* ] *mult-operand*

13 R1006 *level-2-expr* is [ [ *level-2-expr* ] *add-op* ] *add-operand*

14 R1007 *power-op* is \*\*

15 R1008 *mult-op* is \*  
16 or /

17 R1009 *add-op* is +  
18 or -

#### NOTE 10.3

Simple examples of a level-2 expression are:

<u>Example</u>	<u>Syntactic class</u>	<u>Remarks</u>
A	<i>level-1-expr</i>	A is a <i>primary</i> . (R1002)
B ** C	<i>mult-operand</i>	B is a <i>level-1-expr</i> , ** is a <i>power-op</i> , and C is a <i>mult-operand</i> . (R1004)
D * E	<i>add-operand</i>	D is an <i>add-operand</i> , * is a <i>mult-op</i> , and E is a <i>mult-operand</i> . (R1005)
+1	<i>level-2-expr</i>	+ is an <i>add-op</i> and 1 is an <i>add-operand</i> . (R1006)
F - I	<i>level-2-expr</i>	F is a <i>level-2-expr</i> , - is an <i>add-op</i> , and I is an <i>add-operand</i> . (R1006)

**NOTE 10.3 (cont.)**

A more complicated example of a level-2 expression is:

- A + D \* E + B \*\* C

1 **10.1.2.5 Level-3 expressions**

2 1 Level-3 expressions are level-2 expressions optionally involving the character operator *concat-op*.

3 R1010 *level-3-expr* is [ *level-3-expr concat-op* ] *level-2-expr*

4 R1011 *concat-op* is //

**NOTE 10.4**

Simple examples of a level-3 expression are:

<u>Example</u>	<u>Syntactic class</u>
A	<i>level-2-expr</i> (R1006)
B // C	<i>level-3-expr</i> (R1010)

A more complicated example of a level-3 expression is:

X // Y // 'ABCD'

5 **10.1.2.6 Level-4 expressions**

6 1 Level-4 expressions are level-3 expressions optionally involving the relational operators *rel-op*.

7 R1012 *level-4-expr* is [ *level-3-expr rel-op* ] *level-3-expr*

8 R1013 *rel-op* is .EQ.  
 9 or .NE.  
 10 or .LT.  
 11 or .LE.  
 12 or .GT.  
 13 or .GE.  
 14 or ==  
 15 or /=  
 16 or <  
 17 or <=  
 18 or >  
 19 or >=

**NOTE 10.5**

Simple examples of a level-4 expression are:

<u>Example</u>	<u>Syntactic class</u>
A	<i>level-3-expr</i> (R1010)
B == C	<i>level-4-expr</i> (R1012)
D < E	<i>level-4-expr</i> (R1012)

A more complicated example of a level-4 expression is:

(A + B) /= C

1 **10.1.2.7 Level-5 expressions**

2 1 Level-5 expressions are level-4 expressions optionally involving the logical operators *not-op*, *and-op*, *or-op*, and  
3 *equiv-op*.

4	R1014	<i>and-operand</i>	is	[ <i>not-op</i> ] <i>level-4-expr</i>
5	R1015	<i>or-operand</i>	is	[ <i>or-operand and-op</i> ] <i>and-operand</i>
6	R1016	<i>equiv-operand</i>	is	[ <i>equiv-operand or-op</i> ] <i>or-operand</i>
7	R1017	<i>level-5-expr</i>	is	[ <i>level-5-expr equiv-op</i> ] <i>equiv-operand</i>
8	R1018	<i>not-op</i>	is	.NOT.
9	R1019	<i>and-op</i>	is	.AND.
10	R1020	<i>or-op</i>	is	.OR.
11	R1021	<i>equiv-op</i>	is	.EQV.
12			or	.NEQV.

**NOTE 10.6**

Simple examples of a level-5 expression are:

<u>Example</u>	<u>Syntactic class</u>
A	<i>level-4-expr</i> (R1012)
.NOT. B	<i>and-operand</i> (R1014)
C .AND. D	<i>or-operand</i> (R1015)
E .OR. F	<i>equiv-operand</i> (R1016)
G .EQV. H	<i>level-5-expr</i> (R1017)
S .NEQV. T	<i>level-5-expr</i> (R1017)

A more complicated example of a level-5 expression is:

A .AND. B .EQV. .NOT. C

13 **10.1.2.8 General form of an expression**

14 1 Expressions are level-5 expressions optionally involving defined binary operators. Defined binary operators have  
15 the lowest operator precedence (Table 10.1).

16 R1022 *expr* is [ *expr defined-binary-op* ] *level-5-expr*

17 R1023 *defined-binary-op* is . *letter* [ *letter* ] ... .

18 C1005 (R1023) A *defined-binary-op* shall not contain more than 63 letters and shall not be the same as any  
19 *intrinsic-operator* or *logical-literal-constant*.

**NOTE 10.7**

Simple examples of an expression are:

<u>Example</u>	<u>Syntactic class</u>
A	<i>level-5-expr</i> (R1017)
B.UNION.C	<i>expr</i> (R1022)

**NOTE 10.7 (cont.)**

More complicated examples of an expression are:

```
(B .INTERSECT. C) .UNION. (X - Y)
A + B == C * D
.INVERSE. (A + B)
A + B .AND. C * D
E // G == H (1:10)
```

### 1 10.1.3 Precedence of operators

- 2 1 There is a precedence among the intrinsic and extension operations corresponding to the form of expressions  
 3 specified in 10.1.2, which determines the order in which the operands are combined unless the order is changed  
 4 by the use of parentheses. This precedence order is summarized in Table 10.1.

Table 10.1: Categories of operations and relative precedence

Category of operation	Operators	Precedence
Extension	<i>defined-unary-op</i>	Highest
Numeric	**	.
Numeric	*, /	.
Numeric	unary +, -	.
Numeric	binary +, -	.
Character	//	.
Relational	.EQ., .NE., .LT., .LE., .GT., .GE., ==, /=, <, <=, >, >=	.
Logical	.NOT.	.
Logical	.AND.	.
Logical	.OR.	.
Logical	.EQV., .NEQV.	.
Extension	<i>defined-binary-op</i>	Lowest

- 5 2 The precedence of a *defined operation* is that of its operator.

**NOTE 10.8**

For example, in the expression

```
-A ** 2
```

the exponentiation operator (\*\*) has precedence over the negation operator (-); therefore, the operands of the exponentiation operator are combined to form an expression that is used as the operand of the negation operator. The interpretation of the above expression is the same as the interpretation of the expression

```
- (A ** 2)
```

- 6 3 The general form of an expression (10.1.2) also establishes a precedence among operators in the same syntactic  
 7 class. This precedence determines the order in which the operands are to be combined in determining the  
 8 interpretation of the expression unless the order is changed by the use of parentheses.

**NOTE 10.9**

In interpreting a *level-2-expr* containing two or more binary operators + or -, each operand (*add-operand*) is combined from left to right. Similarly, the same left-to-right interpretation for a *mult-operand* in *add-operand*, as well as for other kinds of expressions, is a consequence of the general form. However, for interpreting a *mult-operand* expression when two or more exponentiation operators \*\* combine *level-1-expr* operands, each *level-1-expr* is combined from right to left.

**NOTE 10.9 (cont.)**

For example, the expressions

```
2.1 + 3.4 + 4.9
2.1 * 3.4 * 4.9
2.1 / 3.4 / 4.9
2 ** 3 ** 4
'AB' // 'CD' // 'EF'
```

have the same interpretations as the expressions

```
(2.1 + 3.4) + 4.9
(2.1 * 3.4) * 4.9
(2.1 / 3.4) / 4.9
2 ** (3 ** 4)
('AB' // 'CD') // 'EF'
```

As a consequence of the general form (10.1.2), only the first *add-operand* of a *level-2-expr* can be preceded by the identity (+) or negation (-) operator. These formation rules do not permit expressions containing two consecutive numeric operators, such as  $A ** -B$  or  $A + -B$ . However, expressions such as  $A ** (-B)$  and  $A + (-B)$  are permitted. The rules do allow a binary operator or an intrinsic unary operator to be followed by a defined unary operator, such as:

```
A * .INVERSE. B
- .INVERSE. (B)
```

As another example, in the expression

```
A .OR. B .AND. C
```

the general form implies a higher precedence for the *.AND.* operator than for the *.OR.* operator; therefore, the interpretation of the above expression is the same as the interpretation of the expression

```
A .OR. (B .AND. C)
```

**NOTE 10.10**

An expression can contain more than one category of operator. The logical expression

```
L .OR. A + B >= C
```

where A, B, and C are of type real, and L is of type logical, contains a numeric operator, a relational operator, and a logical operator. This expression would be interpreted the same as the expression

```
L .OR. ((A + B) >= C)
```

**NOTE 10.11**

If

- the operator *\*\** is extended to type logical,
- the operator *.STARSTAR.* is defined to duplicate the function of *\*\** on type real,
- *.MINUS.* is defined to duplicate the unary operator *-*, and
- L1 and L2 are type logical and X and Y are type real,

then in precedence:  $L1 ** L2$  is higher than  $X * Y$ ;  $X * Y$  is higher than  $X .STARSTAR. Y$ ; and *.MINUS.* X is higher than  $-X$ .

## 10.1.4 Evaluation of operations

- 1 An intrinsic operation requires the values of its operands.
- 2 Execution of a function reference in the logical expression in an **IF statement** (11.1.8.4), the mask expression in a **WHERE statement** (10.2.3.1), or the *concurrent-limits* and *concurrent-steps* in a **FORALL statement** (10.2.4) is permitted to define variables in the subsidiary *action-stmt*, *where-assignment-stmt*, or *forall-assignment-stmt* respectively. Except in those cases:
- the evaluation of a function reference shall neither affect nor be affected by the evaluation of any other entity within the statement;
  - if a function reference causes definition or undefinition of an **actual argument** of the function, that argument or any associated entities shall not appear elsewhere in the same statement.

### NOTE 10.12

For example, the statements

```
A (I) = F (I)
Y = G (X) + X
```

are prohibited if the reference to F defines or undefines I or the reference to G defines or undefines X.

However, in the statements

```
IF (F (X)) A = X
WHERE (G (X)) B = X
```

the reference to F and/or the reference to G can define X.

- 3 The appearance of an array constructor requires the evaluation of each *scalar-int-expr* of the *ac-implied-do-control* in any *ac-implied-do* it may contain.
- 4 When an **elemental** binary operation is applied to a scalar and an array or to two arrays of the same shape, the operation is performed element-by-element on corresponding array elements of the array operands.

### NOTE 10.13

For example, the array expression

```
A + B
```

produces an array of the same shape as A and B. The individual array elements of the result have the values of the first element of A added to the first element of B, the second element of A added to the second element of B, etc.

- 5 When an **elemental** unary operator operates on an array operand, the operation is performed element-by-element, and the result is the same shape as the operand. If an **elemental operation** is intrinsically pure or is implemented by a pure **elemental** function (15.8), the element operations may be performed simultaneously or in any order.

## 10.1.5 Intrinsic operations

### 10.1.5.1 Intrinsic operation classification

- 1 An intrinsic operation is either a unary or binary operation. An intrinsic unary operation is an operation of the form *intrinsic-operator*  $x_2$  where  $x_2$  is of an intrinsic type (7.4) listed in Table 10.2 for the unary intrinsic operator.
- 2 An intrinsic binary operation is an operation of the form  $x_1$  *intrinsic-operator*  $x_2$  where  $x_1$  and  $x_2$  are **conformable** and of the intrinsic types (7.4) listed in Table 10.2 for the binary intrinsic operator.

- 1 3 A numeric intrinsic operation is an intrinsic operation for which the *intrinsic-operator* is a numeric operator (+, 2 -, \*, /, or \*\*). A numeric intrinsic operator is the operator in a numeric intrinsic operation.
- 3 4 The character intrinsic operation is the intrinsic operation for which the *intrinsic-operator* is (//) and both 4 operands are of type character with the same *kind type parameter*. The character intrinsic operator is the 5 operator in a character intrinsic operation.
- 6 5 A logical intrinsic operation is an intrinsic operation for which the *intrinsic-operator* is *.AND.*, *.OR.*, *.NOT.*, 7 *.EQV.*, or *.NEQV.* and both operands are of type logical. A logical intrinsic operator is the operator in a logical 8 intrinsic operation.
- 9 6 A relational intrinsic operator is an *intrinsic-operator* that is *.EQ.*, *.NE.*, *.GT.*, *.GE.*, *.LT.*, *.LE.*, *==*, */=*, *>*, 10 *>=*, *<*, or *<=*. A relational intrinsic operation is an intrinsic operation for which the *intrinsic-operator* is a 11 relational intrinsic operator. A numeric relational intrinsic operation is a relational intrinsic operation for which 12 both operands are of *numeric type*. A character relational intrinsic operation is a relational intrinsic operation for 13 which both operands are of type character. The *kind type parameters* of the operands of a character relational 14 intrinsic operation shall be the same.
- 15 7 The interpretations defined in subclause 10.1.5 apply to both scalars and arrays; the interpretation for arrays is 16 obtained by applying the interpretation for scalars element by element.

Table 10.2: Type of operands and results for intrinsic operators

Intrinsic operator <i>op</i>	Type of <i>x</i> <sub>1</sub>	Type of <i>x</i> <sub>2</sub>	Type of <i>[x</i> <sub>1</sub> <i>] op x</i> <sub>2</sub>
Unary +, -		I, R, Z	I, R, Z
Binary +, -, *, /, **	I	I, R, Z	I, R, Z
	R Z	I, R, Z I, R, Z	R, R, Z Z, Z, Z
//	C	C	C
<i>.EQ.</i> , <i>.NE.</i> , <i>==</i> , <i>/=</i>	I	I, R, Z	L, L, L
	R	I, R, Z	L, L, L
	Z C	I, R, Z C	L, L, L L
<i>.GT.</i> , <i>.GE.</i> , <i>.LT.</i> , <i>.LE.</i> <i>&gt;</i> , <i>&gt;=</i> , <i>&lt;</i> , <i>&lt;=</i>	I	I, R	L, L
	R	I, R	L, L
	C	C	L
<i>.NOT.</i>		L	L
<i>.AND.</i> , <i>.OR.</i> , <i>.EQV.</i> , <i>.NEQV.</i>	L	L	L

Note: The symbols I, R, Z, C, and L stand for the types integer, real, complex, character, and logical, respectively. Where more than one type for *x*<sub>2</sub> is given, the type of the result of the operation is given in the same relative position in the next column.

**NOTE 10.14**

For example, if X is of type real and J is of type integer, the expression X + J is of type real.

17 **10.1.5.2 Numeric intrinsic operations**18 **10.1.5.2.1 Interpretation of numeric intrinsic operations**

- 19 1 The two operands of numeric intrinsic binary operations may be of different *numeric types* or different *kind* 20 *type parameters*. Except for a value of type real or complex raised to an integer power, if the operands have 21 different types or *kind type parameters*, the effect is as if each operand that differs in type or *kind type parameter* 22 from those of the result is converted to the type and *kind type parameter* of the result before the operation is 23 performed. When a value of type real or complex is raised to an integer power, the integer operand need not be



- 1 converted.
- 2 2 A numeric operation is used to express a numeric computation. Evaluation of a numeric operation produces a  
3 numeric value. The permitted data types for operands of the numeric intrinsic operations are specified in 10.1.5.1.
- 4 3 The numeric operators and their interpretation in an expression are given in Table 10.3, where  $x_1$  denotes the  
5 operand to the left of the operator and  $x_2$  denotes the operand to the right of the operator.

Table 10.3: Interpretation of the numeric intrinsic operators

Operator	Representing	Use of operator	Interpretation
**	Exponentiation	$x_1 ** x_2$	Raise $x_1$ to the power $x_2$
/	Division	$x_1 / x_2$	Divide $x_1$ by $x_2$
*	Multiplication	$x_1 * x_2$	Multiply $x_1$ by $x_2$
-	Subtraction	$x_1 - x_2$	Subtract $x_2$ from $x_1$
-	Negation	$- x_2$	Negate $x_2$
+	Addition	$x_1 + x_2$	Add $x_1$ and $x_2$
+	Identity	$+ x_2$	Same as $x_2$

- 6 4 The interpretation of a division operation depends on the types of the operands (10.1.5.2.2).
- 7 5 If  $x_1$  and  $x_2$  are of type integer and  $x_2$  has a negative value, the interpretation of  $x_1 ** x_2$  is the same as the  
8 interpretation of  $1/(x_1 ** \text{ABS}(x_2))$ , which is subject to the rules of integer division (10.1.5.2.2).

**NOTE 10.15**

For example,  $2 ** (-3)$  has the value of  $1/(2 ** 3)$ , which is zero.

9 **10.1.5.2.2 Integer division**

- 10 1 One operand of type integer may be divided by another operand of type integer. Although the mathematical  
11 quotient of two integers is not necessarily an integer, Table 10.2 specifies that an expression involving the division  
12 operator with two operands of type integer is interpreted as an expression of type integer. The result of such an  
13 operation is the integer closest to the mathematical quotient and between zero and the mathematical quotient  
14 inclusively.

**NOTE 10.16**

For example, the expression  $(-8) / 3$  has the value  $(-2)$ .

15 **10.1.5.2.3 Complex exponentiation**

- 16 1 In the case of a complex value raised to a complex power, the value of the operation  $x_1 ** x_2$  is the principal  
17 value of  $x_1^{x_2}$ .

18 **10.1.5.2.4 Evaluation of numeric intrinsic operations**

- 19 1 The execution of any numeric operation whose result is not defined by the arithmetic used by the processor is  
20 prohibited. Raising a negative real value to a real power is prohibited.
- 21 2 Once the interpretation of a numeric intrinsic operation is established, the processor may evaluate any mathem-  
22 atically equivalent expression, provided that the integrity of parentheses is not violated.
- 23 3 Two expressions of a **numeric type** are mathematically equivalent if, for all possible values of their primaries, their  
24 mathematical values are equal. However, mathematically equivalent expressions of **numeric type** may produce  
25 different computational results.

**NOTE 10.17**

Any difference between the values of the expressions  $(1./3.)*3.$  and  $1.$  is a computational difference, not a mathematical difference. The difference between the values of the expressions  $5/2$  and  $5./2.$  is a mathematical difference, not a computational difference.

The mathematical definition of integer division is given in [10.1.5.2.2](#).

**NOTE 10.18**

The following are examples of expressions with allowable alternative forms that can be used by the processor in the evaluation of those expressions. A, B, and C represent arbitrary real or complex operands; I and J represent arbitrary integer operands; and X, Y, and Z represent arbitrary operands of [numeric type](#).

<u>Expression</u>	<u>Allowable alternative form</u>
$X + Y$	$Y + X$
$X * Y$	$Y * X$
$-X + Y$	$Y - X$
$X + Y + Z$	$X + (Y + Z)$
$X - Y + Z$	$X - (Y - Z)$
$X * A / Z$	$X * (A / Z)$
$X * Y - X * Z$	$X * (Y - Z)$
$A / B / C$	$A / (B * C)$
$A / 5.0$	$0.2 * A$

The following are examples of expressions with forbidden alternative forms that cannot be used by a processor in the evaluation of those expressions.

<u>Expression</u>	<u>Forbidden alternative form</u>
$I / 2$	$0.5 * I$
$X * I / J$	$X * (I / J)$
$I / J / A$	$I / (J * A)$
$(X + Y) + Z$	$X + (Y + Z)$
$(X * Y) - (X * Z)$	$X * (Y - Z)$
$X * (Y - Z)$	$X * Y - X * Z$

**NOTE 10.19**

In addition to the parentheses required to establish the desired interpretation, parentheses can be included to restrict the alternative forms that can be used by the processor in the actual evaluation of the expression. This is useful for controlling the magnitude and accuracy of intermediate values developed during the evaluation of an expression.

For example, in the expression

$$A + (B - C)$$

the parenthesized expression  $(B - C)$  is evaluated and then added to A.

The inclusion of parentheses could change the mathematical value of an expression. For example, the two expressions

$$A * I / J$$

$$A * (I / J)$$

could have different mathematical values if I and J are of type integer.

**NOTE 10.20**

Each operand in a numeric intrinsic operation has a type that can depend on the order of evaluation used by the processor.

For example, in the evaluation of the expression

$$Z + R + I$$

where Z, R, and I represent data objects of complex, real, and integer type, respectively, the type of the operand that is added to I could be either complex or real, depending on which pair of operands (Z and R, R and I, or Z and I) is added first.

1 **10.1.5.3 Character intrinsic operation**2 **10.1.5.3.1 Interpretation of the character intrinsic operation**

3 1 The character intrinsic operator // is used to concatenate two operands of type character with the same [kind](#)  
4 [type parameter](#). Evaluation of the character intrinsic operation produces a result of type character.

5 2 The interpretation of the character intrinsic operator // when used to form an expression is given in Table 10.4,  
6 where  $x_1$  denotes the operand to the left of the operator and  $x_2$  denotes the operand to the right of the operator.

Table 10.4: **Interpretation of the character intrinsic operator //**

Operator	Representing	Use of operator	Interpretation
//	Concatenation	$x_1 // x_2$	Concatenate $x_1$ with $x_2$

7 3 The result of the character intrinsic operation  $x_1 // x_2$  is a character string whose value is the value of  $x_1$   
8 concatenated on the right with the value of  $x_2$  and whose length is the sum of the lengths of  $x_1$  and  $x_2$ . Parentheses  
9 used to specify the order of evaluation have no effect on the value of a character expression.

**NOTE 10.21**

For example, the value of ('AB' // 'CDE') // 'F' is the string 'ABCDEF'. Also, the value of 'AB' // ('CDE' // 'F') is the string 'ABCDEF'.

10 **10.1.5.3.2 Evaluation of the character intrinsic operation**

11 1 A processor is only required to evaluate as much of the character intrinsic operation as is required by the context  
12 in which the expression appears.

**NOTE 10.22**

For example, the statements

```
CHARACTER (LEN = 2) C1, C2, C3, CF
C1 = C2 // CF (C3)
```

do not require the function CF to be evaluated, because only the value of C2 is needed to determine the value of C1 because C1 and C2 both have a length of 2.

13 **10.1.5.4 Logical intrinsic operations**14 **10.1.5.4.1 Interpretation of logical intrinsic operations**

15 1 A logical operation is used to express a logical computation. Evaluation of a logical operation produces a result  
16 of type logical. The permitted types for operands of the logical intrinsic operations are specified in [10.1.5.1](#).

1 2 The logical operators and their interpretation when used to form an expression are given in Table 10.5, where  $x_1$   
 2 denotes the operand to the left of the operator and  $x_2$  denotes the operand to the right of the operator.

Table 10.5: Interpretation of the logical intrinsic operators

Operator	Representing	Use of operator	Interpretation
.NOT.	Logical negation	.NOT. $x_2$	True if $x_2$ is false
.AND.	Logical conjunction	$x_1$ .AND. $x_2$	True if $x_1$ and $x_2$ are both true
.OR.	Logical inclusive disjunction	$x_1$ .OR. $x_2$	True if $x_1$ and/or $x_2$ is true
.EQV.	Logical equivalence	$x_1$ .EQV. $x_2$	True if both $x_1$ and $x_2$ are true or both are false
.NEQV.	Logical nonequivalence	$x_1$ .NEQV. $x_2$	True if either $x_1$ or $x_2$ is true, but not both

3 3 The values of the logical intrinsic operations are shown in Table 10.6.

Table 10.6: The values of operations involving logical intrinsic operators

$x_1$	$x_2$	.NOT. $x_2$	$x_1$ .AND. $x_2$	$x_1$ .OR. $x_2$	$x_1$ .EQV. $x_2$	$x_1$ .NEQV. $x_2$
true	true	false	true	true	true	false
true	false	true	false	true	false	true
false	true	false	false	true	false	true
false	false	true	false	false	true	false

4 **10.1.5.4.2 Evaluation of logical intrinsic operations**

5 1 Once the interpretation of a logical intrinsic operation is established, the processor may evaluate any other  
 6 expression that is logically equivalent, provided that the integrity of parentheses in any expression is not violated.

**NOTE 10.23**

For example, for the variables L1, L2, and L3 of type logical, the processor could choose to evaluate the expression

L1 .AND. L2 .AND. L3

as

L1 .AND. (L2 .AND. L3)

7 2 Two expressions of type logical are logically equivalent if their values are equal for all possible values of their  
 8 primaries.

9 **10.1.5.5 Relational intrinsic operations**

10 **10.1.5.5.1 Interpretation of relational intrinsic operations**

11 1 A relational intrinsic operation is used to compare values of two operands using the relational intrinsic operators  
 12 .LT., .LE., .GT., .GE., .EQ., .NE., <, <=, >, >=, ==, and /=. The permitted types for operands of the  
 13 relational intrinsic operators are specified in 10.1.5.1.

14 2 The operators <, <=, >, >=, ==, and /= always have the same interpretations as the operators .LT., .LE.,  
 15 .GT., .GE., .EQ., and .NE., respectively.

**NOTE 10.24**

As shown in Table 10.2, a relational intrinsic operator cannot be used to compare the value of an expression of a numeric type with one of type character or logical. Also, two operands of type logical cannot be

**NOTE 10.24 (cont.)**

compared, a complex operand can be compared with another numeric operand only when the operator is `.EQ.`, `.NE.`, `==`, or `/=`, and two character operands cannot be compared unless they have the same **kind type parameter** value.

- 1 3 Evaluation of a relational intrinsic operation produces a default logical result.
- 2 4 The interpretation of the relational intrinsic operators is given in Table 10.7, where  $x_1$  denotes the operand to  
3 the left of the operator and  $x_2$  denotes the operand to the right of the operator.

Table 10.7: **Interpretation of the relational intrinsic operators**

Operator	Representing	Use of operator	Interpretation
<code>.LT.</code>	Less than	$x_1$ <code>.LT.</code> $x_2$	$x_1$ less than $x_2$
<code>&lt;</code>	Less than	$x_1 < x_2$	$x_1$ less than $x_2$
<code>.LE.</code>	Less than or equal to	$x_1$ <code>.LE.</code> $x_2$	$x_1$ less than or equal to $x_2$
<code>&lt;=</code>	Less than or equal to	$x_1 <= x_2$	$x_1$ less than or equal to $x_2$
<code>.GT.</code>	Greater than	$x_1$ <code>.GT.</code> $x_2$	$x_1$ greater than $x_2$
<code>&gt;</code>	Greater than	$x_1 > x_2$	$x_1$ greater than $x_2$
<code>.GE.</code>	Greater than or equal to	$x_1$ <code>.GE.</code> $x_2$	$x_1$ greater than or equal to $x_2$
<code>&gt;=</code>	Greater than or equal to	$x_1 >= x_2$	$x_1$ greater than or equal to $x_2$
<code>.EQ.</code>	Equal to	$x_1$ <code>.EQ.</code> $x_2$	$x_1$ equal to $x_2$
<code>==</code>	Equal to	$x_1 == x_2$	$x_1$ equal to $x_2$
<code>.NE.</code>	Not equal to	$x_1$ <code>.NE.</code> $x_2$	$x_1$ not equal to $x_2$
<code>/=</code>	Not equal to	$x_1 /= x_2$	$x_1$ not equal to $x_2$

- 4 5 A numeric relational intrinsic operation is interpreted as having the logical value true if and only if the values of  
5 the operands satisfy the relation specified by the operator.
- 6 6 In the numeric relational operation  
7  $x_1$  *rel-op*  $x_2$   
8 if the types or **kind type parameters** of  $x_1$  and  $x_2$  differ, their values are converted to the type and **kind type**  
9 **parameter** of the expression  $x_1 + x_2$  before evaluation.
- 10 7 A character relational intrinsic operation is interpreted as having the logical value true if and only if the values  
11 of the operands satisfy the relation specified by the operator.
- 12 8 For a character relational intrinsic operation, the operands are compared one character at a time in order,  
13 beginning with the first character of each character operand. If the operands are of unequal length, the shorter  
14 operand is treated as if it were extended on the right with blanks to the length of the longer operand. If both  
15  $x_1$  and  $x_2$  are of zero length,  $x_1$  is equal to  $x_2$ ; if every character of  $x_1$  is the same as the character in the  
16 corresponding position in  $x_2$ ,  $x_1$  is equal to  $x_2$ . Otherwise, at the first position where the character operands  
17 differ, the character operand  $x_1$  is considered to be less than  $x_2$  if the character value of  $x_1$  at this position  
18 precedes the value of  $x_2$  in the **collating sequence** (3.31);  $x_1$  is greater than  $x_2$  if the character value of  $x_1$  at this  
19 position follows the value of  $x_2$  in the **collating sequence**.

**NOTE 10.25**

The **collating sequence** depends partially on the processor; however, the result of the use of the operators `.EQ.`, `.NE.`, `==`, and `/=` does not depend on the **collating sequence**.

For nondefault character kinds, the blank padding character is processor dependent.

**10.1.5.5.2 Evaluation of relational intrinsic operations**

- 20 1 Once the interpretation of a relational intrinsic operation is established, the processor may evaluate any other  
21 expression that is relationally equivalent, provided that the integrity of parentheses in any expression is not  
22

- 1 violated.
- 2 2 Two relational intrinsic operations are relationally equivalent if their logical values are equal for all possible values  
3 of their primaries.

**NOTE 10.26**

Whether an operand of a relational intrinsic operation could be an IEEE NaN affects whether expressions are equivalent. For example, if  $x$  or  $y$  could be a NaN, the expressions

.NOT. (x .LT. y) and x .GE. y

are not equivalent.

## 4 10.1.6 Defined operations

### 5 10.1.6.1 Definitions

- 6 1 A **defined operation** is either a unary operation or a binary operation. A unary **defined operation** is an operation  
7 that has the form *defined-unary-op*  $x_2$  or *intrinsic-operator*  $x_2$  and that is defined by a function and a **generic**  
8 **interface** (7.5.5, 15.4.3.4).
- 9 2 A function defines the unary operation  $op$   $x_2$  if
- 10 (1) the function is specified with a **FUNCTION** (15.6.2.2) or **ENTRY** (15.6.2.6) statement that specifies one  
11 **dummy argument**  $d_2$ ,
  - 12 (2) either
    - 13 (a) a **generic interface** (15.4.3.2) provides the function with a *generic-spec* of **OPERATOR** ( $op$ ),  
14 or
    - 15 (b) there is a generic **binding** (7.5.5) in the **declared type** of  $x_2$  with a *generic-spec* of **OPER-**  
16 **ATOR** ( $op$ ) and there is a corresponding **binding** to the function in the **dynamic type** of  $x_2$ ,
  - 17 (3) the type of  $d_2$  is compatible with the **dynamic type** of  $x_2$ ,
  - 18 (4) the type parameters, if any, of  $d_2$  match the corresponding type parameters of  $x_2$ , and
  - 19 (5) either
    - 20 (a) the **rank** of  $x_2$  matches that of  $d_2$  or
    - 21 (b) the function is **elemental** and there is no other function that defines the operation.
- 22 3 If  $d_2$  is an array, the shape of  $x_2$  shall match the shape of  $d_2$ .
- 23 4 A binary **defined operation** is an operation that has the form  $x_1$  *defined-binary-op*  $x_2$  or  $x_1$  *intrinsic-operator*  $x_2$   
24 and that is defined by a function and a **generic interface**.
- 25 5 A function defines the binary operation  $x_1$   $op$   $x_2$  if
- 26 (1) the function is specified with a **FUNCTION** (15.6.2.2) or **ENTRY** (15.6.2.6) statement that specifies  
27 two **dummy arguments**,  $d_1$  and  $d_2$ ,
  - 28 (2) either
    - 29 (a) a **generic interface** (15.4.3.2) provides the function with a *generic-spec* of **OPERATOR** ( $op$ ),  
30 or
    - 31 (b) there is a generic **binding** (7.5.5) in the **declared type** of  $x_1$  or  $x_2$  with a *generic-spec* of  
32 **OPERATOR** ( $op$ ) and there is a corresponding **binding** to the function in the **dynamic type**  
33 of  $x_1$  or  $x_2$ , respectively,
  - 34 (3) the types of  $d_1$  and  $d_2$  are compatible with the **dynamic types** of  $x_1$  and  $x_2$ , respectively,
  - 35 (4) the type parameters, if any, of  $d_1$  and  $d_2$  match the corresponding type parameters of  $x_1$  and  $x_2$ ,  
36 respectively, and
  - 37 (5) either

- 1 (a) the **ranks** of  $x_1$  and  $x_2$  match those of  $d_1$  and  $d_2$ , respectively, or  
 2 (b) the function is **elemental**,  $x_1$  and  $x_2$  are **conformable**, and there is no other function that defines  
 3 the operation.
- 4 6 If  $d_1$  or  $d_2$  is an array, the shapes of  $x_1$  and  $x_2$  shall match the shapes of  $d_1$  and  $d_2$ , respectively.

**NOTE 10.27**

An intrinsic operator can be used as the operator in a **defined operation**. In such a case, the generic properties of the operator are extended.

5 **10.1.6.2 Interpretation of a defined operation**

- 6 1 The interpretation of a **defined operation** is provided by the function that defines the operation.  
 7 2 The operators  $<$ ,  $<=$ ,  $>$ ,  $>=$ ,  $==$ , and  $/=$  always have the same interpretations as the operators **.LT.**, **.LE.**,  
 8 **.GT.**, **.GE.**, **.EQ.**, and **.NE.**, respectively.

9 **10.1.6.3 Evaluation of a defined operation**

- 10 1 Once the interpretation of a **defined operation** is established, the processor may evaluate any other expression  
 11 that is equivalent, provided that the integrity of parentheses is not violated.  
 12 2 Two expressions of derived type are equivalent if their values are equal for all possible values of their primaries.

13 **10.1.7 Evaluation of operands**

- 14 1 It is not necessary for a processor to evaluate all of the operands of an expression, or to evaluate entirely each  
 15 operand, if the value of the expression can be determined otherwise.

**NOTE 10.28**

This principle is most often applicable to logical expressions, zero-sized arrays, and zero-length strings, but it applies to all expressions.

For example, in evaluating the expression

$$X > Y \text{ .OR. } L(Z)$$

where  $X$ ,  $Y$ , and  $Z$  are real and  $L$  is a function of type logical, the function reference  $L(Z)$  need not be evaluated if  $X$  is greater than  $Y$ . Similarly, in the array expression

$$W(Z) + A$$

where  $A$  is of size zero and  $W$  is a function, the function reference  $W(Z)$  need not be evaluated.

- 16 2 If a statement contains a function reference in a part of an expression that need not be evaluated, all entities that  
 17 would have become defined in the execution of that reference become undefined at the completion of evaluation  
 18 of the expression containing the function reference.

**NOTE 10.29**

In the examples in Note 10.28, if  $L$  or  $W$  defines its argument, evaluation of the expressions under the specified conditions causes  $Z$  to become undefined, no matter whether or not  $L(Z)$  or  $W(Z)$  is evaluated.

- 19 3 If a statement contains a function reference in a part of an expression that need not be evaluated, no invocation  
 20 of that function in that part of the expression shall execute an **image control statement** other than **CRITICAL**  
 21 or **END CRITICAL**.

**NOTE 10.30**

This restriction is intended to avoid inadvertent deadlock caused by optimization.

**10.1.8 Integrity of parentheses**

1

- 2 1 The rules for evaluation specified in subclause 10.1.5 state certain conditions under which a processor may  
3 evaluate an expression that is different from the one specified by applying the rules given in 10.1.2 and rules for  
4 interpretation specified in subclause 10.1.5. However, any expression in parentheses shall be treated as a data  
5 entity.

**NOTE 10.31**

For example, in evaluating the expression  $A + (B - C)$  where A, B, and C are of **numeric types**, the difference of B and C shall be evaluated before the addition operation is performed; the processor shall not evaluate the mathematically equivalent expression  $(A + B) - C$ .

**10.1.9 Type, type parameters, and shape of an expression**

6

**10.1.9.1 General**

7

- 8 1 The type, type parameters, and shape of an expression depend on the operators and on the types, type parameters,  
9 and shapes of the primaries used in the expression, and are determined recursively from the syntactic form of the  
10 expression. The type of an expression is one of the intrinsic types (7.4) or a derived type (7.5).
- 11 2 If an expression is a **polymorphic primary** or **defined operation**, the type parameters and the **declared** and **dynamic**  
12 types of the expression are the same as those of the primary or **defined operation**. Otherwise the type parameters  
13 and **dynamic type** of the expression are the same as its **declared type** and type parameters; they are referred to  
14 simply as the type and type parameters of the expression.

15 R1024 *logical-expr* is *expr*

16 C1006 (R1024) *logical-expr* shall be of type logical.

17 R1025 *default-char-expr* is *expr*

18 C1007 (R1025) *default-char-expr* shall be default character.

19 R1026 *int-expr* is *expr*

20 C1008 (R1026) *int-expr* shall be of type integer.

21 R1027 *numeric-expr* is *expr*

22 C1009 (R1027) *numeric-expr* shall be of type integer, real, or complex.

**10.1.9.2 Type, type parameters, and shape of a primary**

23

- 24 1 The type, **type parameters**, and shape of a primary are determined according to whether the primary is a **literal**  
25 **constant**, **designator**, array constructor, **structure constructor**, function reference, **type parameter inquiry**, type  
26 parameter name, or parenthesized expression. If a primary is a **literal constant**, its type, **type parameters**, and  
27 shape are those of the **literal constant**. If it is a **structure constructor**, it is scalar and its type and **type parameters**  
28 are as described in 7.5.10. If it is an array constructor, its type, **type parameters**, and shape are as described in  
29 7.8. If it is a **designator** or function reference, its type, **type parameters**, and shape are those of the **designator**  
30 (8.2, 8.5) or the function reference (15.5.3), respectively. If the function reference is generic (15.4.3.2, 16.7)  
31 then its type, **type parameters**, and shape are those of the specific function referenced, which is determined by  
32 the declared types, **type parameters**, and **ranks** of its **actual arguments** as specified in 15.5.5.2. If it is a **type**  
33 **parameter inquiry** or type parameter name, it is a scalar integer with the kind of the type parameter.



- 1 2 If a primary is a parenthesized expression, its type, [type parameters](#), and shape are those of the expression.
- 2 3 The associated [target](#) object is referenced if a pointer appears as
- 3 • a primary in an intrinsic or [defined operation](#),
  - 4 • the *expr* of a parenthesized primary, or
  - 5 • the only primary on the right-hand side of an [intrinsic assignment statement](#).
- 6 4 The type, type parameters, and shape of the primary are those of the [target](#). If the pointer is not associated with
- 7 a [target](#), it may appear as a primary only as an [actual argument](#) in a reference to a procedure whose corresponding
- 8 [dummy argument](#) is declared to be a pointer, as the target in a [pointer assignment statement](#), or as explicitly
- 9 permitted elsewhere in this document.
- 10 5 A [disassociated array pointer](#) or an unallocated [allocatable](#) array has no shape but does have [rank](#). The type,
- 11 type parameters, and [rank](#) of the result of the intrinsic function [NULL](#) (16.9.144) depend on context.

### 12 10.1.9.3 Type, type parameters, and shape of the result of an operation

- 13 1 The type of the result of an intrinsic operation  $[x_1] \text{ op } x_2$  is specified by Table 10.2. The shape of the result of
- 14 an intrinsic operation is the shape of  $x_2$  if *op* is unary or if  $x_1$  is scalar, and is the shape of  $x_1$  otherwise.
- 15 2 The type, type parameters, and shape of the result of a defined operation  $[x_1] \text{ op } x_2$  are specified by the function
- 16 defining the operation (10.1.6).
- 17 3 An expression of an intrinsic type has a [kind type parameter](#). An expression of type character also has a character
- 18 length parameter.
- 19 4 The type parameters of the result of an intrinsic operation are as follows.
- 20 • For an expression  $x_1 // x_2$  where *//* is the character intrinsic operator and  $x_1$  and  $x_2$  are of type character,
  - 21 the character length parameter is the sum of the lengths of the operands and the [kind type parameter](#) is
  - 22 the [kind type parameter](#) of  $x_1$ , which shall be the same as the [kind type parameter](#) of  $x_2$ .
  - 23 • For an expression  $op x_2$  where *op* is an intrinsic unary operator and  $x_2$  is of type integer, real, complex, or
  - 24 logical, the [kind type parameter](#) of the expression is that of the operand.
  - 25 • For an expression  $x_1 \text{ op } x_2$  where *op* is a numeric intrinsic binary operator with one operand of type integer
  - 26 and the other of type real or complex, the [kind type parameter](#) of the expression is that of the real or
  - 27 complex operand.
  - 28 • For an expression  $x_1 \text{ op } x_2$  where *op* is a numeric intrinsic binary operator with both operands of the same
  - 29 type and [kind type parameters](#), or with one real and one complex with the same [kind type parameters](#), the
  - 30 [kind type parameter](#) of the expression is identical to that of each operand. In the case where both operands
  - 31 are integer with different [kind type parameters](#), the [kind type parameter](#) of the expression is that of the
  - 32 operand with the greater decimal exponent range if the decimal exponent ranges are different; if the decimal
  - 33 exponent ranges are the same, the [kind type parameter](#) of the expression is processor dependent, but it is
  - 34 the same as that of one of the operands. In the case where both operands are any of type real or complex
  - 35 with different [kind type parameters](#), the [kind type parameter](#) of the expression is that of the operand with
  - 36 the greater decimal precision if the decimal precisions are different; if the decimal precisions are the same,
  - 37 the [kind type parameter](#) of the expression is processor dependent, but it is the same as that of one of the
  - 38 operands.
  - 39 • For an expression  $x_1 \text{ op } x_2$  where *op* is a logical intrinsic binary operator with both operands of the same
  - 40 [kind type parameter](#), the [kind type parameter](#) of the expression is identical to that of each operand. In the
  - 41 case where both operands are of type logical with different [kind type parameters](#), the [kind type parameter](#)
  - 42 of the expression is processor dependent, but it is the same as that of one of the operands.
  - 43 • For an expression  $x_1 \text{ op } x_2$  where *op* is a relational intrinsic operator, the [kind type parameter](#) of the
  - 44 expression is default logical.

### 10.1.10 Conformability rules for elemental operations

- 1 An **elemental operation** is an intrinsic operation or a **defined operation** for which the function is **elemental** (15.8).
- 2 For all **elemental** binary operations, the two operands shall be conformable. In the case where one is a scalar and the other an array, the scalar is treated as if it were an array of the same shape as the array operand with every element, if any, of the array equal to the value of the scalar.

### 10.1.11 Specification expression

- 1 A **specification expression** is an expression with limitations that make it suitable for use in specifications such as length type parameters (C704) and array bounds (R817, R818). A *specification-expr* shall be a **constant expression** unless it is in an **interface body** (15.4.3.2), the specification part of a subprogram or **BLOCK construct**, a derived type definition, or the *declaration-type-spec* of a **FUNCTION statement** (15.6.2.2).

R1028 *specification-expr* is *scalar-int-expr*

C1010 (R1028) The *scalar-int-expr* shall be a restricted expression.

- 2 A restricted expression is an expression in which each operation is intrinsic or defined by a specification function and each primary is

- (1) a constant or subobject of a constant,
- (2) an **object designator** with a **base object** that is a **dummy argument** that has neither the **OPTIONAL** nor the **INTENT (OUT)** attribute,
- (3) an **object designator** with a **base object** that is in a **common block**,
- (4) an **object designator** with a **base object** that is made accessible by **use** or **host** association,
- (5) an array constructor where each element and each *scalar-int-expr* of each *ac-implied-do-control* is a restricted expression,
- (6) a **structure constructor** where each component is a restricted expression,
- (7) a specification inquiry where each **designator** or argument is
  - (a) a restricted expression or
  - (b) a variable that is not an optional dummy argument, and whose properties inquired about are not
    - (i) dependent on the upper bound of the last dimension of an **assumed-size array**,
    - (ii) deferred, or
    - (iii) defined by an expression that is not a restricted expression,
- (8) a specification inquiry that is a constant expression,
- (9) a reference to the intrinsic function **PRESENT**,
- (10) a reference to any other **standard intrinsic** function where each argument is a restricted expression,
- (11) a reference to a transformational function from the intrinsic module **IEEE\_ARITHMETIC**, **IEEE\_EXCEPTIONS**, or **ISO\_C\_BINDING**, where each argument is a restricted expression,
- (12) a reference to a specification function where each argument is a restricted expression,
- (13) a type parameter of the derived type being defined,
- (14) an *ac-do-variable* within an array constructor where each *scalar-int-expr* of the corresponding *ac-implied-do-control* is a restricted expression, or
- (15) a restricted expression enclosed in parentheses,

- 3 where each subscript, section subscript, substring starting point, substring ending point, and type parameter value is a restricted expression, and where any **final subroutine** that is invoked is pure.

- 4 A specification inquiry is a reference to

- (1) an intrinsic **inquiry function** other than **PRESENT**,

- 1 (2) a [type parameter inquiry](#) (9.4.5),  
 2 (3) an [inquiry function](#) from the intrinsic modules [IEEE\\_ARITHMETIC](#) and [IEEE\\_EXCEPTIONS](#)  
 3 (17.10),  
 4 (4) the function [C\\_SIZEOF](#) from the intrinsic module [ISO\\_C\\_BINDING](#) (18.2.3.7), or  
 5 (5) the [COMPILER\\_VERSION](#) or [COMPILER\\_OPTIONS](#) function from the intrinsic module [ISO\\_-](#)  
 6 [FORTRAN\\_ENV](#) (16.10.2.6, 16.10.2.7).
- 7 5 A function is a specification function if it is a pure function, is not a standard intrinsic function, is not an internal  
 8 function, is not a statement function, and does not have a [dummy procedure](#) argument.
- 9 6 Evaluation of a [specification expression](#) shall not directly or indirectly cause a procedure defined by the subpro-  
 10 gram in which it appears to be invoked.

**NOTE 10.32**

Specification functions are nonintrinsic functions that can be used in specification expressions to determine the attributes of data objects. The requirement that they be pure ensures that they cannot have side effects that could affect other objects being declared in the same *specification-part*. The requirement that they not be internal ensures that they cannot inquire, via [host association](#), about other objects being declared in the same *specification-part*. The prohibition against recursion avoids the creation of a new instance of a procedure while construction of one is in progress.

- 11 7 A variable in a [specification expression](#) shall have its type and type parameters, if any, specified by a previous  
 12 declaration in the same [scoping unit](#), by the implicit typing rules in effect for the [scoping unit](#), or by [host](#) or [use](#)  
 13 association. If a variable in a [specification expression](#) is typed by the implicit typing rules, its appearance in any  
 14 subsequent [type declaration statement](#) shall confirm the implied type and type parameters.
- 15 8 If a [specification expression](#) includes a specification inquiry that depends on a [type parameter](#), [array bound](#),  
 16 or [cobound](#) of an entity specified in the same *specification-part*, the [type parameter](#), [array bound](#), or [cobound](#)  
 17 shall be specified in a prior specification of the *specification-part*. The prior specification may be to the left of  
 18 the specification inquiry in the same statement, but shall not be within the same *entity-decl*. If a [specification](#)  
 19 [expression](#) includes a reference to the value of an element of an array specified in the same *specification-part*, the  
 20 array shall be completely specified in prior declarations.
- 21 9 A generic entity referenced in a [specification expression](#) in the *specification-part* of a [scoping unit](#) shall have no  
 22 specific procedures defined in the [scoping unit](#), or its [host scoping unit](#), subsequent to the [specification expression](#).

**NOTE 10.33**

The following are examples of [specification expressions](#):

```

LBOUND (B, 1) + 5 ! B is an assumed-shape dummy array
M + LEN (C)      ! M and C are dummy arguments
2 * PRECISION (A) ! A is a real variable made accessible
                  ! by a USE statement

```

**10.1.12 Constant expression**

- 23 1 A [constant expression](#) is an expression with limitations that make it suitable for use as a [kind type parameter](#),  
 24 initializer, or [named constant](#). It is an expression in which each operation is intrinsic, and each primary is  
 25
- 26 (1) a constant or subobject of a constant,  
 27 (2) an array constructor where each element and each *scalar-int-expr* of each *ac-implied-do-control* is a  
 28 [constant expression](#),  
 29 (3) a [structure constructor](#) where each *component-spec* corresponding to  
 30 (a) an [allocatable](#) component is a reference to the intrinsic function [NULL](#),

- 1 (b) a pointer component is an initialization target or a reference to the intrinsic function `NULL`,  
 2 and  
 3 (c) any other component is a `constant expression`,
- 4 (4) a specification inquiry where each `designator` or argument is  
 5 (a) a `constant expression` or  
 6 (b) a variable whose properties inquired about are not  
 7 (i) assumed,  
 8 (ii) deferred, or  
 9 (iii) defined by an expression that is not a `constant expression`,
- 10 (5) a reference to an `elemental` standard intrinsic function, where each argument is a `constant expression`,  
 11 (6) a reference to a `transformational` standard intrinsic function other than `COMMAND_ARGUMENT_`  
 12 `COUNT`, `NULL`, `NUM_IMAGES`, `THIS_IMAGE`, or `TRANSFER`, where each argument is a `con-`  
 13 `stant expression`,
- 14 (7) a reference to the intrinsic function `NULL` that does not have an argument with a type parameter  
 15 that is assumed or is defined by an expression that is not a `constant expression`,
- 16 (8) a reference to the intrinsic function `TRANSFER` where each argument is a `constant expression` and  
 17 each `ultimate pointer` component of the `SOURCE` argument is `disassociated`,
- 18 (9) a reference to a `transformational function` from the intrinsic module `IEEE_ARITHMETIC` or `IEEE_-`  
 19 `EXCEPTIONS`, where each argument is a `constant expression`,
- 20 (10) a previously declared `kind type parameter` of the derived type being defined,  
 21 (11) a `data-i-do-variable` within a `data-implied-do`,  
 22 (12) an `ac-do-variable` within an array constructor where each `scalar-int-expr` of the corresponding `ac-`  
 23 `implied-do-control` is a `constant expression`, or  
 24 (13) a `constant expression` enclosed in parentheses,

25 and where each subscript, section subscript, substring starting point, substring ending point, and type parameter  
 26 value is a `constant expression`.

27 R1029 `constant-expr` is `expr`

28 C1011 (R1029) `constant-expr` shall be a `constant expression`.

29 R1030 `default-char-constant-expr` is `default-char-expr`

30 C1012 (R1030) `default-char-constant-expr` shall be a `constant expression`.

31 R1031 `int-constant-expr` is `int-expr`

32 C1013 (R1031) `int-constant-expr` shall be a `constant expression`.

33 2 If a `constant expression` includes a specification inquiry that depends on a type parameter or an array bound of  
 34 an entity specified in the same `specification-part`, the type parameter or array bound shall be specified in a prior  
 35 specification of the `specification-part`. The prior specification may be to the left of the specification inquiry in the  
 36 same statement, but shall not be within the same `entity-decl` unless the specification inquiry appears within an  
 37 `initialization`.

38 3 A generic entity referenced in a `constant expression` in the `specification-part` of a `scoping unit` shall have no specific  
 39 procedures defined in that `scoping unit`, or its `host scoping unit`, subsequent to the `constant expression`.

#### NOTE 10.34

The following are examples of `constant expressions`:

3  
 -3 + 4

## NOTE 10.34 (cont.)

```
'AB'
'AB' // 'CD'
('AB' // 'CD') // 'EF'
SIZE (A)
DIGITS (X) + 4
4.0 * ATAN (1.0)
CEILING (number_of_decimal_digits / LOG10 (REAL (RADIX (0.0))))
```

where A is an [explicit-shape array](#) with constant bounds and X is default real.

## 10.2 Assignment

## 10.2.1 Assignment statement

## 10.2.1.1 General form

R1032 *assignment-stmt* is *variable* = *expr*

C1014 (R1032) The *variable* shall not be a whole [assumed-size array](#).

## NOTE 10.35

Examples of an assignment statement are:

```
A = 3.5 + X * Y
I = INT (A)
```

1 An *assignment-stmt* shall meet the requirements of either a [defined assignment](#) statement or an intrinsic assignment statement.

## 10.2.1.2 Intrinsic assignment statement

1 An intrinsic assignment statement is an assignment statement that is not a [defined assignment](#) statement (10.2.1.4). In an intrinsic assignment statement,

- (1) if the variable is [polymorphic](#) it shall be [allocatable](#) and not a [coarray](#),
- (2) if *expr* is an array then the variable shall also be an array,
- (3) the variable and *expr* shall be [conformable](#) unless the variable is an [allocatable](#) array that has the same [rank](#) as *expr* and is not a [coarray](#),
- (4) if the variable is [polymorphic](#) it shall be [type compatible](#) with *expr*; otherwise the [declared types](#) of the variable and *expr* shall conform as specified in Table 10.8,
- (5) if the variable is of type character and of [ISO 10646](#), [ASCII](#), or default character kind, *expr* shall be of [ISO 10646](#), [ASCII](#), or default character kind,
- (6) otherwise if the variable is of type character *expr* shall have the same [kind type parameter](#),
- (7) if the variable is of derived type each [kind type parameter](#) of the variable shall have the same value as the corresponding [kind type parameter](#) of *expr*, and
- (8) if the variable is of derived type each length type parameter of the variable shall have the same value as the corresponding type parameter of *expr* unless the variable is [allocatable](#), is not a [coarray](#), and its corresponding [type parameter](#) is [deferred](#).

Table 10.8: Intrinsic assignment type conformance

Type of the variable	Type of <i>expr</i>
integer	integer, real, complex

**Intrinsic assignment type conformance** (cont.)

Type of the variable	Type of <i>expr</i>
real	integer, real, complex
complex	integer, real, complex
character	character
logical	logical
derived type	same derived type as the variable

- 1 2 If the variable in an intrinsic assignment statement is a **coindexed object**,
- 2     • the variable shall not be **polymorphic**,
- 3     • the variable shall not have an **allocatable ultimate component**,
- 4     • the variable shall be **conformable** with *expr*, and
- 5     • each **deferred length type parameter** of the variable shall have the same value as the corresponding type
- 6     parameter of *expr*.

- 7 3 If the variable is a pointer, it shall be associated with a **definable target** such that the type, type parameters, and
- 8 shape of the **target** and *expr* conform.

### 9 10.2.1.3 Interpretation of intrinsic assignments

- 10 1 Execution of an intrinsic assignment causes, in effect, the evaluation of the expression *expr* and all expressions
- 11 within *variable* (10.1), the possible conversion of *expr* to the type and type parameters of the variable (Table
- 12 10.9), and the definition of the variable with the resulting value. The execution of the assignment shall have
- 13 the same effect as if the evaluation of *expr* and the evaluation of all expressions in *variable* occurred before any
- 14 portion of the variable is defined by the assignment. The evaluation of expressions within *variable* shall neither
- 15 affect nor be affected by the evaluation of *expr*.
- 16 2 If the variable is a pointer, the value of *expr* is assigned to the **target** of the variable.
- 17 3 If the variable is an unallocated **allocatable** array, *expr* shall have the same **rank**. If the variable is an allocated
- 18 **allocatable** variable, it is deallocated if *expr* is an array of different shape, any of the corresponding length type
- 19 parameter values of the variable and *expr* differ, or the variable is **polymorphic** and the **dynamic type** of the
- 20 variable and *expr* differ. If the variable is or becomes an unallocated **allocatable** variable, it is then allocated with
- 21     • the same **dynamic type** as *expr* if the variable is **polymorphic**,
- 22     • each **deferred type parameter** equal to the corresponding type parameter of *expr*,
- 23     • the same bounds as before if the variable is an array and *expr* is scalar, and
- 24     • the shape of *expr* with each lower bound equal to the corresponding element of **LBOUND** (*expr*) if *expr* is
- 25     an array.

**NOTE 10.36**

For example, given the declaration

```
CHARACTER(:), ALLOCATABLE :: NAME
```

then after the assignment statement

```
NAME = 'Dr. '//FIRST_NAME//' '//SURNAME
```

NAME will have the length **LEN** (FIRST\_NAME) + **LEN** (SURNAME) + 5, even if it had previously been unallocated, or allocated with a different length. However, the assignment statement

```
NAME(:) = 'Dr. '//FIRST_NAME//' '//SURNAME
```

**NOTE 10.36 (cont.)**

is only conforming if NAME is already allocated at the time of the assignment; the assigned value is truncated or blank padded to the previously allocated length of NAME.

- 1 4 Both *variable* and *expr* may contain references to any portion of the variable.

**NOTE 10.37**

For example, in the character intrinsic assignment statement:

```
STRING (2:5) = STRING (1:4)
```

the assignment of the first character of STRING to the second character does not affect the evaluation of STRING (1:4). If the value of STRING prior to the assignment was 'ABCDEF', the value following the assignment is 'AABCDF'.

- 2 5 If *expr* is a scalar and the variable is an array, the *expr* is treated as if it were an array of the same shape as the  
3 variable with every element of the array equal to the scalar value of *expr*.
- 4 6 If the variable is an array, the assignment is performed element-by-element on corresponding array elements of  
5 the variable and *expr*.

**NOTE 10.38**

For example, if A and B are arrays of the same shape, the array intrinsic assignment

```
A = B
```

assigns the corresponding elements of B to those of A; that is, the first element of B is assigned to the first element of A, the second element of B is assigned to the second element of A, etc.

If C is an *allocatable* array of *rank* 1, then

```
C = PACK (ARRAY, ARRAY>0)
```

will cause C to contain all the positive elements of ARRAY in array element order; if C is not allocated or is allocated with the wrong size, it will be re-allocated to be of the correct size to hold the result of PACK.

- 6 7 The processor may perform the element-by-element assignment in any order.

**NOTE 10.39**

For example, the following program segment results in the values of the elements of array X being reversed:

```
REAL X (10)
...
X (1:10) = X (10:1:-1)
```

- 7 8 For an intrinsic assignment statement where the variable is of numeric type, the *expr* can have a different *numeric*  
8 *type* or *kind type parameter*, in which case the value of *expr* is converted to the type and *kind type parameter*  
9 of the variable according to the rules of Table 10.9.

Table 10.9: **Numeric conversion and the assignment statement**

Type of the variable	Value assigned
integer	INT ( <i>expr</i> , KIND = KIND ( <i>variable</i> ))
real	REAL ( <i>expr</i> , KIND = KIND ( <i>variable</i> ))
complex	CMPLX ( <i>expr</i> , KIND = KIND ( <i>variable</i> ))

**Numeric conversion and the assignment statement** (cont.)

Type of the variable	Value Assigned
Note: <code>INT</code> , <code>REAL</code> , <code>CMPLX</code> , and <code>KIND</code> are the generic names of functions defined in 16.9.	

- 1 9 For an intrinsic assignment statement where the variable is of type logical, the *expr* can have a different **kind type parameter**, in which case the value of *expr* is converted to the **kind type parameter** of the variable.
- 2
- 3 10 For an intrinsic assignment statement where the variable is of type character, the *expr* can have a different character length parameter in which case the conversion of *expr* to the length of the variable is as follows.
- 4
- 5 (1) If the length of the variable is less than that of *expr*, the value of *expr* is truncated from the right until it is the same length as the variable.
- 6
- 7 (2) If the length of the variable is greater than that of *expr*, the value of *expr* is extended on the right with blanks until it is the same length as the variable.
- 8
- 9 11 For an intrinsic assignment statement where the variable is of type character, if *expr* has a different **kind type parameter**, each character *c* in *expr* is converted to the **kind type parameter** of the variable by `ACHAR (IACHAR(c), KIND (variable))`.
- 10
- 11

**NOTE 10.40**

For nondefault character kinds, the blank padding character is processor dependent. When assigning a character expression to a variable of a different kind, each character of the expression that is not representable in the kind of the variable is replaced by a processor-dependent character.

- 12 12 For an intrinsic assignment of the type `C_PTR` or `C_FUNPTR` from the intrinsic module `ISO_C_BINDING`, the variable becomes undefined if the variable and *expr* are not on the same **image**.
- 13

**NOTE 10.41**

An intrinsic assignment statement for a variable of type `C_PTR` or `C_FUNPTR` is not permitted to involve a **coindexed object**, see C915, which prevents inappropriate copying from one **image** to another. However, such copying can occur as an intrinsic assignment for a component in a derived-type assignment.

- 14 13 An intrinsic assignment where the variable is of derived type is performed as if each component of the variable were assigned from the corresponding component of *expr* using **pointer assignment** (10.2.2) for each pointer component, **defined assignment** for each nonpointer nonallocatable component of a type that has a **type-bound defined assignment** consistent with the component, intrinsic assignment for each other nonpointer nonallocatable component, and intrinsic assignment for each allocated **coarray** component. For unallocated **coarray** components, the corresponding component of the variable shall be unallocated. For a noncoarray **allocatable** component the following sequence of operations is applied.
- 15
- 16
- 17
- 18
- 19
- 20
- 21 (1) If the component of the variable is allocated, it is deallocated.
- 22 (2) If the component of the value of *expr* is allocated, the corresponding component of the variable is allocated with the same **dynamic type** and type parameters as the component of the value of *expr*. If it is an array, it is allocated with the same bounds. The value of the component of the value of *expr* is then assigned to the corresponding component of the variable using **defined assignment** if the **declared type** of the component has a **type-bound defined assignment** consistent with the component, and intrinsic assignment for the **dynamic type** of that component otherwise.
- 23
- 24
- 25
- 26
- 27
- 28 14 The processor may perform the component-by-component assignment in any order or by any means that has the same effect.
- 29



**NOTE 10.42**

For an example of a derived-type intrinsic assignment statement, if C and D are of the same derived type with a pointer component P and nonpointer components S, T, U, and V of type integer, logical, character, and another derived type, respectively, the intrinsic

$$C = D$$

pointer assigns D%P to C%P. It assigns D%S to C%S, D%T to C%T, and D%U to C%U using intrinsic assignment. It assigns D%V to C%V using [defined assignment](#) if objects of that type have a compatible [type-bound defined assignment](#), and intrinsic assignment otherwise.

**NOTE 10.43**

If an [allocatable](#) component of *expr* is unallocated, the corresponding component of the variable has an allocation status of unallocated after execution of the assignment.

1 **10.2.1.4 Defined assignment statement**

2 1 A [defined assignment](#) statement is an assignment statement that is defined by a subroutine and a [generic interface](#)  
3 (7.5.5, 15.4.3.4.3) that specifies [ASSIGNMENT](#) (=).

4 2 A subroutine defines the [defined assignment](#)  $x_1 = x_2$  if

5 (1) the subroutine is specified with a [SUBROUTINE](#) (15.6.2.3) or [ENTRY](#) (15.6.2.6) statement that specifies  
6 two [dummy arguments](#),  $d_1$  and  $d_2$ ,

7 (2) either

8 (a) a [generic interface](#) (15.4.3.2) provides the subroutine with a [generic-spec](#) of [ASSIGNMENT](#) (=),  
9 or

10 (b) there is a [generic binding](#) (7.5.5) in the [declared type](#) of  $x_1$  or  $x_2$  with a [generic-spec](#) of  
11 [ASSIGNMENT](#) (=) and there is a corresponding [binding](#) to the subroutine in the [dynamic](#)  
12 [type](#) of  $x_1$  or  $x_2$ , respectively,

13 (3) the types of  $d_1$  and  $d_2$  are compatible with the [dynamic types](#) of  $x_1$  and  $x_2$ , respectively,

14 (4) the type parameters, if any, of  $d_1$  and  $d_2$  match the corresponding type parameters of  $x_1$  and  $x_2$ ,  
15 respectively, and

16 (5) either

17 (a) the [ranks](#) of  $x_1$  and  $x_2$  match those of  $d_1$  and  $d_2$  or

18 (b) the subroutine is [elemental](#),  $x_2$  is scalar or has the same [rank](#) as  $x_1$ , and there is no other  
19 subroutine that defines the assignment.

20 3 If  $d_1$  or  $d_2$  is an array, the shapes of  $x_1$  and  $x_2$  shall match the shapes of  $d_1$  and  $d_2$ , respectively. If the subroutine  
21 is [elemental](#),  $x_2$  shall be [conformable](#) with  $x_1$ .

22 **10.2.1.5 Interpretation of defined assignment statements**

23 1 The interpretation of a [defined assignment](#) is provided by the subroutine that defines it.

24 2 If the [defined assignment](#) is an [elemental assignment](#) and the variable in the assignment is an array, the [defined](#)  
25 [assignment](#) is performed element-by-element, on corresponding elements of the variable and *expr*. If *expr* is a  
26 scalar, it is treated as if it were an array of the same shape as the variable with every element of the array equal  
27 to the scalar value of *expr*.

**NOTE 10.44**

The rules of [defined assignment](#) (15.4.3.4.3), [procedure references](#) (15.5), [subroutine references](#) (15.5.4), and [elemental](#) subroutine arguments (15.8.3) ensure that the [defined assignment](#) has the same effect as if the

## NOTE 10.44 (cont.)

evaluation of all operations in  $x_2$  and  $x_1$  occurs before any portion of  $x_1$  is defined. If an **elemental assignment** is defined by a pure **elemental** subroutine, the element assignments can be performed simultaneously or in any order.

## 10.2.2 Pointer assignment

## 10.2.2.1 General

- 1 **Pointer assignment** causes a pointer to become associated with a **target** or causes its **pointer association** status to become **disassociated** or undefined. Any previous association between the pointer and a **target** is broken.
- 2 **Pointer assignment** for a pointer component of a structure may also take place by execution of a derived-type **intrinsic assignment statement** (10.2.1.3).

## 10.2.2.2 Syntax of the pointer assignment statement

R1033 *pointer-assignment-stmt*      **is** *data-pointer-object* [ (*bounds-spec-list*) ] => *data-target*  
    **or** *data-pointer-object* (*bounds-remapping-list*) => *data-target*  
    **or** *proc-pointer-object* => *proc-target*

R1034 *data-pointer-object*            **is** *variable-name*  
    **or** *scalar-variable* % *data-pointer-component-name*

C1015 (R1033) If *data-target* is not **unlimited polymorphic**, *data-pointer-object* shall be **type compatible** (7.3.2.3) with it and the corresponding **kind type parameters** shall be equal.

C1016 (R1033) If *data-target* is **unlimited polymorphic**, *data-pointer-object* shall be **unlimited polymorphic**, or of a type with the **BIND attribute** or the **SEQUENCE attribute**.

C1017 (R1033) If *bounds-spec-list* is specified, the number of *bounds-specs* shall equal the **rank** of *data-pointer-object*.

C1018 (R1033) If *bounds-remapping-list* is specified, the number of *bounds-remappings* shall equal the **rank** of *data-pointer-object*.

C1019 (R1033) If *bounds-remapping-list* is not specified, the **ranks** of *data-pointer-object* and *data-target* shall be the same.

C1020 (R1033) A coarray *data-target* shall have the **VOLATILE attribute** if and only if the *data-pointer-object* has the **VOLATILE attribute**.

C1021 (R1034) A *variable-name* shall have the **POINTER attribute**.

C1022 (R1034) A *scalar-variable* shall be a *data-ref*.

C1023 (R1034) A *data-pointer-component-name* shall be the name of a component of *scalar-variable* that is a data pointer.

C1024 (R1034) A *data-pointer-object* shall not be a **coindexed object**.

R1035 *bounds-spec*                    **is** *lower-bound-expr* :

R1036 *bounds-remapping*            **is** *lower-bound-expr* : *upper-bound-expr*

R1037 *data-target*                   **is** *expr*

C1025 (R1037) The *expr* shall be a *designator* that designates a variable with either the **TARGET** or **POINTER** attribute and is not an array section with a vector subscript, or it shall be a reference to a function that

1 returns a data pointer.

2 C1026 (R1037) A *data-target* shall not be a *coindexed object*.

**NOTE 10.45**

A data pointer and its *target* are always on the same *image*. A *coarray* can be of a derived type with pointer or *allocatable subcomponents*. For example, if PTR is a pointer component, and Z%PTR on *image* P has been associated with a *target* by execution of an *ALLOCATE statement* or a *pointer assignment* on *image* P, Z[P]%PTR will be a reference to that *target*.

3 R1038 *proc-pointer-object* is *proc-pointer-name*  
4 or *proc-component-ref*

5 R1039 *proc-component-ref* is *scalar-variable* % *procedure-component-name*

6 C1027 (R1039) The *scalar-variable* shall be a *data-ref* that is not a *coindexed object*.

7 C1028 (R1039) The *procedure-component-name* shall be the name of a procedure pointer component of the  
8 *declared type* of *scalar-variable*.

9 R1040 *proc-target* is *expr*  
10 or *procedure-name*  
11 or *proc-component-ref*

12 C1029 (R1040) An *expr* shall be a reference to a function whose result is a procedure pointer.

13 C1030 (R1040) A *procedure-name* shall be the name of an *internal*, *module*, or *dummy* procedure, a *procedure*  
14 *pointer*, a *specific* intrinsic function listed in Table 16.2, or an *external procedure* that is accessed by *use* or *host*  
15 *association*, referenced in the *scoping unit* as a procedure, or that has the *EXTERNAL* attribute.

16 C1031 (R1040) The *proc-target* shall not be a nonintrinsic *elemental procedure*.

17 1 In a pointer assignment statement, *data-pointer-object* or *proc-pointer-object* denotes the pointer object and  
18 *data-target* or *proc-target* denotes the pointer target.

19 2 For pointer assignment performed by a derived-type *intrinsic assignment statement*, the pointer object is the  
20 pointer component of the variable and the pointer target is the corresponding component of *expr*.

21 **10.2.2.3 Data pointer assignment**

22 1 If the pointer object is not polymorphic (7.3.2.3) and the pointer target is *polymorphic* with *dynamic type* that  
23 differs from its *declared type*, the assignment target is the ancestor component of the pointer target that has the  
24 type of the pointer object. Otherwise, the assignment target is the pointer target.

25 2 If the pointer target is not a pointer, the pointer object becomes *pointer associated* with the assignment target;  
26 if the pointer target is a pointer with a target that is not on the same *image*, the *pointer association* status of the  
27 pointer object becomes undefined. Otherwise, the *pointer association* status of the pointer object becomes that  
28 of the pointer target; if the pointer target is associated with an object, the pointer object becomes associated  
29 with the assignment target. If the pointer target is *allocatable*, it shall be allocated.

**NOTE 10.46**

A *pointer assignment statement* is not permitted to involve a *coindexed* pointer or target, see C1024 and C1026. This prevents a pointer assignment statement from associating a pointer with a target on another *image*. If such an association would otherwise be implied, the association status of the pointer becomes undefined. For example, a derived-type intrinsic assignment where the variable and *expr* are on different images and the variable has an ultimate pointer component.

30 3 If the pointer object is *polymorphic*, it assumes the *dynamic type* of the pointer target. If the pointer object is

- 1 of a type with the [BIND attribute](#) or the [SEQUENCE attribute](#), the [dynamic type](#) of the pointer target shall be  
2 that type.
- 3 4 If the pointer target is a [disassociated](#) pointer, all nondeferred type parameters of the [declared type](#) of the pointer  
4 object that correspond to nondeferred type parameters of the pointer target shall have the same values as the  
5 corresponding type parameters of the pointer target.
- 6 5 Otherwise, all nondeferred type parameters of the [declared type](#) of the pointer object shall have the same values  
7 as the corresponding type parameters of the pointer target.
- 8 6 If the pointer object has nondeferred type parameters that correspond to [deferred type parameters](#) of the pointer  
9 target, the pointer target shall not be a pointer with undefined association status.
- 10 7 If the pointer object has the [CONTIGUOUS attribute](#), the pointer target shall be [contiguous](#).
- 11 8 If the target of a pointer is a [coarray](#), the pointer shall have the [VOLATILE attribute](#) if and only if the [coarray](#)  
12 has the [VOLATILE attribute](#).
- 13 9 If [bounds-remapping-list](#) appears, it specifies the upper and lower [bounds](#) of each dimension of the pointer,  
14 and thus the [extents](#); the pointer target shall be [simply contiguous](#) (9.5.4) or of [rank](#) one, and shall not be a  
15 [disassociated](#) or [undefined](#) pointer. The number of elements of the pointer target shall not be less than the  
16 number implied by the [bounds-remapping-list](#). The elements of the pointer object are associated with those of  
17 the pointer [target](#), in array element order; if the pointer target has more elements than specified for the pointer  
18 object, the remaining elements are not associated with the pointer object.
- 19 10 If no [bounds-remapping-list](#) appears, the extent of a dimension of the pointer object is the extent of the corres-  
20 ponding dimension of the pointer target. If [bounds-spec-list](#) appears, it specifies the lower bounds; otherwise,  
21 the lower bound of each dimension is the result of the intrinsic function [LBOUND](#) (16.9.109) applied to the  
22 corresponding dimension of the pointer target. The upper bound of each dimension is one less than the sum of  
23 the lower bound and the extent.

#### 24 10.2.2.4 Procedure pointer assignment

- 25 1 If the pointer target is not a pointer, the pointer object becomes [pointer associated](#) with the pointer target.  
26 Otherwise, the [pointer association](#) status of the pointer object becomes that of the pointer target; if the pointer  
27 target is associated with a procedure, the pointer object becomes associated with the same procedure.
- 28 2 The [host instance](#) (15.6.2.4) of an associated [procedure pointer](#) is the [host instance](#) of its target.
- 29 3 If the pointer object has an [explicit interface](#), its [characteristics](#) shall be the same as the pointer target except  
30 that the pointer target may be pure even if the pointer object is not pure and the pointer target may be an  
31 [elemental](#) intrinsic procedure even if the pointer object is not [elemental](#).
- 32 4 If the [characteristics](#) of the pointer object or the pointer target are such that an [explicit interface](#) is required,  
33 both the pointer object and the pointer target shall have an [explicit interface](#).
- 34 5 If the pointer object has an [implicit interface](#) and is explicitly typed or referenced as a function, the pointer target  
35 shall be a function. If the pointer object has an [implicit interface](#) and is referenced as a subroutine, the pointer  
36 target shall be a subroutine.
- 37 6 If the pointer object is a function with an [implicit interface](#), the pointer target shall be a function with the same  
38 type; corresponding type parameters shall have the same value.
- 39 7 If [procedure-name](#) is a specific procedure name that is also a generic name, only the specific procedure is associated  
40 with the pointer object.

#### 41 10.2.2.5 Examples

**NOTE 10.47**

The following are examples of pointer assignment statements. (See Note 15.14 for declarations of P and BESSEL.)

```

NEW_NODE % LEFT => CURRENT_NODE
SIMPLE_NAME => TARGET_STRUCTURE % SUBSTRUCT % COMPONENT
PTR => NULL ( )
ROW => MAT2D (N, :)
WINDOW => MAT2D (I-1:I+1, J-1:J+1)
POINTER_OBJECT => POINTER_FUNCTION (ARG_1, ARG_2)
EVERY_OTHER => VECTOR (1:N:2)
WINDOW2 (0:, 0:) => MAT2D (ML:MU, NL:NU)
! P is a procedure pointer and BESSEL is a procedure with a
! compatible interface.
P => BESSEL

! Likewise for a structure component.
STRUCT % COMPONENT => BESSEL

```

**NOTE 10.48**

It is possible to obtain different-rank views of parts of an object by specifying upper bounds in pointer assignment statements. This requires that the object be either rank one or contiguous. Consider the following example, in which a matrix is under consideration. The matrix is stored as a rank-one object in MYDATA because its diagonal is needed for some reason – the diagonal cannot be gotten as a single object from a rank-two representation. The matrix is represented as a rank-two view of MYDATA.

```

real, target :: MYDATA ( NR*NC )      ! An automatic array
real, pointer :: MATRIX ( :, : )      ! A rank-two view of MYDATA
real, pointer :: VIEW_DIAG ( : )
MATRIX( 1:NR, 1:NC ) => MYDATA        ! The MATRIX view of the data
VIEW_DIAG => MYDATA( 1::NR+1 )        ! The diagonal of MATRIX

```

Rows, columns, or blocks of the matrix can be accessed as sections of MATRIX.

Rank remapping can be applied to CONTIGUOUS arrays, for example:

```

REAL, CONTIGUOUS, POINTER :: A(:)
REAL, CONTIGUOUS, TARGET :: B(:, :) ! Dummy argument
A(1:SIZE(B)) => B                    ! Linear view of a rank-2 array

```

**10.2.3 Masked array assignment – WHERE****10.2.3.1 General form of the masked array assignment**

- 1 A masked array assignment is either a WHERE statement or a WHERE construct. It is used to mask the  
 2 evaluation of expressions and assignment of values in array assignment statements, according to the value of a  
 3 logical array expression.

6 R1041 *where-stmt* is WHERE ( *mask-expr* ) *where-assignment-stmt*

7 R1042 *where-construct* is *where-construct-stmt*  
 8 [ *where-body-construct* ] ...  
 9 [ *masked-elsewhere-stmt*  
 10 [ *where-body-construct* ] ... ] ...  
 11 [ *elsewhere-stmt*

- 1 [ *where-body-construct* ] ... ]  
 2 *end-where-stmt*
- 3 R1043 *where-construct-stmt* is [ *where-construct-name*: ] WHERE ( *mask-expr* )
- 4 R1044 *where-body-construct* is *where-assignment-stmt*  
 5 or *where-stmt*  
 6 or *where-construct*
- 7 R1045 *where-assignment-stmt* is *assignment-stmt*
- 8 R1046 *mask-expr* is *logical-expr*
- 9 R1047 *masked-elsewhere-stmt* is ELSEWHERE ( *mask-expr* ) [ *where-construct-name* ]
- 10 R1048 *elsewhere-stmt* is ELSEWHERE [ *where-construct-name* ]
- 11 R1049 *end-where-stmt* is END WHERE [ *where-construct-name* ]
- 12 C1032 (R1045) A *where-assignment-stmt* that is a defined assignment shall be elemental.
- 13 C1033 (R1042) If the *where-construct-stmt* is identified by a *where-construct-name*, the corresponding *end-where-stmt* shall specify the same *where-construct-name*. If the *where-construct-stmt* is not identified by a *where-construct-name*, the corresponding *end-where-stmt* shall not specify a *where-construct-name*. If an *elsewhere-stmt* or a *masked-elsewhere-stmt* is identified by a *where-construct-name*, the corresponding *where-construct-stmt* shall specify the same *where-construct-name*.
- 18 C1034 (R1044) A statement that is part of a *where-body-construct* shall not be a branch target statement.
- 19 2 If a *where-construct* contains a *where-stmt*, a *masked-elsewhere-stmt*, or another *where-construct* then each *mask-expr* within the *where-construct* shall have the same shape. In each *where-assignment-stmt*, the *mask-expr* and the variable being defined shall be arrays of the same shape.

**NOTE 10.49**

Examples of a masked array assignment are:

```
WHERE (TEMP > 100.0) TEMP = TEMP - REDUCE_TEMP
WHERE (PRESSURE <= 1.0)
  PRESSURE = PRESSURE + INC_PRESSURE
  TEMP = TEMP - 5.0
ELSEWHERE
  RAINING = .TRUE.
END WHERE
```

**10.2.3.2 Interpretation of masked array assignments**

- 23 1 When a WHERE statement or a *where-construct-stmt* is executed, a control mask is established. In addition,  
 24 when a WHERE construct statement is executed, a pending control mask is established. If the statement does  
 25 not appear as part of a *where-body-construct*, the *mask-expr* of the statement is evaluated, and the control mask is  
 26 established to be the value of *mask-expr*. The pending control mask is established to have the value *.NOT. mask-expr*  
 27 upon execution of a WHERE construct statement that does not appear as part of a *where-body-construct*.  
 28 The *mask-expr* is evaluated only once.
- 29 2 Each statement in a WHERE construct is executed in sequence.
- 30 3 Upon execution of a *masked-elsewhere-stmt*, the following actions take place in sequence.
- 31 (1) The control mask  $m_c$  is established to have the value of the pending control mask.  
 32 (2) The pending control mask is established to have the value  $m_c$  *.AND. (.NOT. mask-expr)*.

- 1 (3) The control mask  $m_c$  is established to have the value  $m_c$  *.AND. mask-expr*.
- 2 4 The *mask-expr* is evaluated at most once.
- 3 5 Upon execution of an ELSEWHERE statement, the control mask is established to have the value of the pending  
4 control mask. No new pending control mask value is established.
- 5 6 Upon execution of an ENDWHERE statement, the control mask and pending control mask are established to  
6 have the values they had prior to the execution of the corresponding WHERE construct statement. Following  
7 the execution of a WHERE statement that appears as a *where-body-construct*, the control mask is established to  
8 have the value it had prior to the execution of the WHERE statement.

**NOTE 10.50**

The establishment of control masks and the pending control mask is illustrated with the following example:

```

WHERE(cond1)      ! Statement 1
. . .
ELSEWHERE(cond2)  ! Statement 2
. . .
ELSEWHERE         ! Statement 3
. . .
END WHERE

```

Following execution of statement 1, the control mask has the value *cond1* and the pending control mask has the value *.NOT. cond1*. Following execution of statement 2, the control mask has the value *(.NOT. cond1) .AND. cond2* and the pending control mask has the value *(.NOT. cond1) .AND. (.NOT. cond2)*. Following execution of statement 3, the control mask has the value *(.NOT. cond1) .AND. (.NOT. cond2)*. The false condition values are propagated through the execution of the masked ELSEWHERE statement.

- 9 7 Upon execution of a WHERE construct statement that is part of a *where-body-construct*, the pending control  
10 mask is established to have the value  $m_c$  *.AND. (.NOT. mask-expr)*. The control mask is then established to  
11 have the value  $m_c$  *.AND. mask-expr*. The *mask-expr* is evaluated at most once.
- 12 8 Upon execution of a WHERE statement that is part of a *where-body-construct*, the control mask is established  
13 to have the value  $m_c$  *.AND. mask-expr*. The pending control mask is not altered.
- 14 9 If a nonelemental function reference occurs in the *expr* or *variable* of a *where-assignment-stmt* or in a *mask-expr*,  
15 the function is evaluated without any masked control; that is, all of its argument expressions are fully evaluated  
16 and the function is fully evaluated. If the result is an array and the reference is not within the argument list  
17 of a nonelemental function, elements corresponding to true values in the control mask are selected for use in  
18 evaluating the *expr*, *variable* or *mask-expr*.
- 19 10 If an elemental operation or function reference occurs in the *expr* or *variable* of a *where-assignment-stmt* or in a  
20 *mask-expr*, and is not within the argument list of a nonelemental function reference, the operation is performed  
21 or the function is evaluated only for the elements corresponding to true values of the control mask.
- 22 11 If an array constructor appears in a *where-assignment-stmt* or in a *mask-expr*, the array constructor is evaluated  
23 without any masked control and then the *where-assignment-stmt* is executed or the *mask-expr* is evaluated.
- 24 12 When a *where-assignment-stmt* is executed, the values of *expr* that correspond to true values of the control mask  
25 are assigned to the corresponding elements of the variable.
- 26 13 The value of the control mask is established by the execution of a WHERE statement, a WHERE construct  
27 statement, an ELSEWHERE statement, a masked ELSEWHERE statement, or an ENDWHERE statement.  
28 Subsequent changes to the value of entities in a *mask-expr* have no effect on the value of the control mask. The  
29 execution of a function reference in the mask expression of a WHERE statement is permitted to affect entities in

1 the assignment statement.

#### NOTE 10.51

Examples of function references in masked array assignments are:

```
WHERE (A > 0.0)
  A = LOG (A)           ! LOG is invoked only for positive elements.
  A = A / SUM (LOG (A)) ! LOG is invoked for all elements
                       ! because SUM is transformational
END WHERE
```

## 2 10.2.4 FORALL

### 3 10.2.4.1 Form of the FORALL Construct

4 1 The FORALL construct allows multiple assignments, masked array (WHERE) assignments, and nested FORALL constructs and  
5 statements to be controlled by a single *concurrent-control-list* and *scalar-mask-expr*.

6 R1050 *forall-construct* is *forall-construct-stmt*  
7 *[forall-body-construct ] ...*  
8 *end-forall-stmt*

9 R1051 *forall-construct-stmt* is *[forall-construct-name :]* FORALL *concurrent-header*

10 R1052 *forall-body-construct* is *forall-assignment-stmt*  
11 or *where-stmt*  
12 or *where-construct*  
13 or *forall-construct*  
14 or *forall-stmt*

15 R1053 *forall-assignment-stmt* is *assignment-stmt*  
16 or *pointer-assignment-stmt*

17 R1054 *end-forall-stmt* is END FORALL *[forall-construct-name ]*

18 C1035 (R1054) If the *forall-construct-stmt* has a *forall-construct-name*, the *end-forall-stmt* shall have the same *forall-construct-*  
19 *name*. If the *end-forall-stmt* has a *forall-construct-name*, the *forall-construct-stmt* shall have the same *forall-construct-*  
20 *name*.

21 C1036 (R1052) A statement in a *forall-body-construct* shall not define an *index-name* of the *forall-construct*.

22 C1037 (R1052) Any procedure referenced in a *forall-body-construct*, including one referenced by a defined operation, assignment,  
23 or *finalization*, shall be a *pure procedure*.

24 C1038 (R1052) A *forall-body-construct* shall not be a branch target.

25 2 The scope and attributes of an *index-name* in a *concurrent-header* in a FORALL construct or statement are described in 19.4.

### 26 10.2.4.2 Execution of the FORALL construct

#### 27 10.2.4.2.1 Execution stages

28 1 There are three stages in the execution of a FORALL construct:  
29 (1) determination of the values for *index-name* variables,  
30 (2) evaluation of the *scalar-mask-expr*, and  
31 (3) execution of the FORALL body constructs.

#### 32 10.2.4.2.2 Determination of the values for index variables

33 1 The values of the index variables are determined as they are for the *DO CONCURRENT statement* (11.1.7.4.2).



### 10.2.4.2.3 Evaluation of the mask expression

- 1 The mask expression is evaluated as it is for the [DO CONCURRENT statement](#) (11.1.7.4.2).

### 10.2.4.2.4 Execution of the FORALL body constructs

- 1 The *forall-body-constructs* are executed in the order in which they appear. Each construct is executed for all active combinations of the *index-name* values with the following interpretation:
- 2 Execution of a *forall-assignment-stmt* that is an *assignment-stmt* causes the evaluation of *expr* and all expressions within *variable* for all active combinations of *index-name* values. These evaluations may be done in any order. After all these evaluations have been performed, each *expr* value is assigned to the corresponding *variable*. The assignments may occur in any order.
  - 3 Execution of a *forall-assignment-stmt* that is a *pointer-assignment-stmt* causes the evaluation of all expressions within *data-target* and *data-pointer-object* or *proc-target* and *proc-pointer-object*, the determination of any pointers within *data-pointer-object* or *proc-pointer-object*, and the determination of the *target* for all active combinations of *index-name* values. These evaluations may be done in any order. After all these evaluations have been performed, each *data-pointer-object* or *proc-pointer-object* is associated with the corresponding *target*. These associations may occur in any order.
  - 4 In a *forall-assignment-stmt*, a *defined assignment* subroutine shall not reference any *variable* that becomes defined by the statement.

#### NOTE 10.52

If a variable defined in an assignment statement within a FORALL construct is referenced in a later statement in that construct, the later statement uses the value(s) computed in the preceding assignment statement, not the value(s) the variable had prior to execution of the FORALL.

- 5 Each statement in a *where-construct* (10.2.3) within a *forall-construct* is executed in sequence. When a *where-stmt*, *where-construct-stmt* or *masked-elsewhere-stmt* is executed, the statement's *mask-expr* is evaluated for all active combinations of *index-name* values as determined by the outer *forall-constructs*, masked by any control mask corresponding to outer *where-constructs*. Any *where-assignment-stmt* is executed for all active combinations of *index-name* values, masked by the control mask in effect for the *where-assignment-stmt*.
- 6 Execution of a *forall-stmt* or *forall-construct* causes the evaluation of the *concurrent-limit* and *concurrent-step* expressions in the *concurrent-control-list* for all active combinations of the *index-name* values of the outer FORALL construct. The set of combinations of *index-name* values for the inner FORALL is the union of the sets defined by these limits and steps for each active combination of the outer *index-name* values; it also includes the outer *index-name* values. The *scalar-mask-expr* is then evaluated for all combinations of the *index-name* values of the inner construct to produce a set of active combinations for the inner construct. If there is no *scalar-mask-expr*, it is as if it appeared with the value true. Each statement in the inner FORALL is then executed for each active combination of the *index-name* values.

### 10.2.4.3 The FORALL statement

- 1 The FORALL statement allows a single assignment statement or [pointer assignment statement](#) to be controlled by a set of index values and an optional mask expression.

R1055 *forall-stmt* is FORALL *concurrent-header forall-assignment-stmt*

- 2 A FORALL statement is equivalent to a FORALL construct containing a single *forall-body-construct* that is a *forall-assignment-stmt*.

- 3 The scope of an *index-name* in a *forall-stmt* is the statement itself (19.4).

### 10.2.4.4 Restrictions on FORALL constructs and statements

- 1 A many-to-one assignment is more than one assignment to the same object, or association of more than one *target* with the same pointer, whether the object is referenced directly or indirectly through a pointer. A many-to-one assignment shall not occur within a single statement in a FORALL construct or statement. It is possible to assign or pointer-assign to the same object in different [assignment](#) or [pointer assignment](#) statements in a FORALL construct.

**NOTE 10.53**

The appearance of each *index-name* in the identification of the left-hand side of an assignment statement is helpful in eliminating many-to-one assignments, but it is not sufficient to guarantee there will be none. For example, the following is allowed

```
FORALL (I = 1:10)
  A (INDEX (I)) = B(I)
END FORALL
```

if and only if INDEX(1:10) contains no repeated values.

- 1 2 Within the scope of a FORALL construct, a nested FORALL statement or FORALL construct shall not have the same *index-name*.
- 2 The *concurrent-header* expressions within a nested FORALL may depend on the values of outer *index-name* variables.

# 11 Execution control

## 11.1 Executable constructs containing blocks

### 11.1.1 Blocks

1 The following are executable constructs that contain blocks:

- 2 • ASSOCIATE construct;
- 3 • BLOCK construct;
- 4 • CHANGE TEAM construct;
- 5 • CRITICAL construct;
- 6 • DO construct;
- 7 • IF construct;
- 8 • SELECT CASE construct;
- 9 • SELECT RANK construct;
- 10 • SELECT TYPE construct.

14 R1101 *block* is [ *execution-part-construct* ] ...

15 2 Executable constructs may be used to control which blocks of a program are executed or how many times a block  
16 is executed. Blocks are always bounded by statements that are particular to the construct in which they are  
17 embedded.

#### NOTE 11.1

An example of a construct containing a block is:

```
IF (A > 0.0) THEN
  B = SQRT (A) ! These two statements
  C = LOG (A) ! form a block.
END IF
```

### 11.1.2 Rules governing blocks

#### 11.1.2.1 Control flow in blocks

1 Transfer of control to the interior of a block from outside the block is prohibited, except for the return from a  
2 procedure invoked within the block. Transfers within a block and transfers from the interior of a block to outside  
3 the block may occur.

4 2 Subroutine and function references (15.5.3, 15.5.4) may appear in a block.

#### 11.1.2.2 Execution of a block

1 Execution of a block begins with the execution of the first executable construct in the block.

2 Execution of the block is completed when

- 3 • execution of the last executable construct in the block completes without branching to a statement within  
4 the block,
- 5 • a branch (11.2) within the block that has a branch target outside the block occurs,

- a RETURN statement within the block is executed, or
- an EXIT statement or CYCLE statement statement that belongs to a construct that contains the block is executed.

**NOTE 11.2**

The action that takes place at the terminal boundary depends on the particular construct and on the block within that construct.

**11.1.3 ASSOCIATE construct****11.1.3.1 Purpose and form of the ASSOCIATE construct**

- 1 The ASSOCIATE construct associates named entities with expressions or variables during the execution of its block. These named construct entities (19.4) are associating entities (19.5.1.6). The names are associate names.

R1102 *associate-construct* is *associate-stmt*  
*block*  
*end-associate-stmt*

R1103 *associate-stmt* is [ *associate-construct-name* : ] ASSOCIATE ■  
 ■ (*association-list* )

R1104 *association* is *associate-name* => *selector*

R1105 *selector* is *expr*  
 or *variable*

C1101 (R1104) If *selector* is not a *variable* or is a *variable* that has a *vector subscript*, neither *associate-name* nor any subobject thereof shall appear in a variable definition context (19.6.7).

C1102 (R1104) An *associate-name* shall not be the same as another *associate-name* in the same *associate-stmt*.

C1103 (R1105) *variable* shall not be a *coindexed object*.

C1104 (R1105) *expr* shall not be a variable.

C1105 (R1105) *expr* shall not be a *designator* of a procedure pointer or a function reference that returns a procedure pointer.

R1106 *end-associate-stmt* is END ASSOCIATE [ *associate-construct-name* ]

C1106 (R1106) If the *associate-stmt* of an *associate-construct* specifies an *associate-construct-name*, the corresponding *end-associate-stmt* shall specify the same *associate-construct-name*. If the *associate-stmt* of an *associate-construct* does not specify an *associate-construct-name*, the corresponding *end-associate-stmt* shall not specify an *associate-construct-name*.

**11.1.3.2 Execution of the ASSOCIATE construct**

- 1 Execution of an ASSOCIATE construct causes evaluation of every expression within every *selector* that is a variable designator and evaluation of every other *selector*, followed by execution of its block. During execution of that block each *associate name* identifies an entity which is associated (19.5.1.6) with the corresponding selector. The *associating entity* assumes the *declared type* and type parameters of the selector. If and only if the selector is *polymorphic*, the *associating entity* is *polymorphic*.

- 2 The other attributes of the *associating entity* are described in 11.1.3.3.

- 3 It is permissible to branch to an *end-associate-stmt* only from within its ASSOCIATE construct.

### 11.1.3.3 Other attributes of associate names

- 1 Within an ASSOCIATE, CHANGE TEAM, or SELECT TYPE construct, each associating entity has the same rank as its associated selector. The lower bound of each dimension is the result of the intrinsic function LBOUND (16.9.109) applied to the corresponding dimension of selector. The upper bound of each dimension is one less than the sum of the lower bound and the extent. The associating entity does not have the ALLOCATABLE or POINTER attributes; it has the TARGET attribute if and only if the selector is a variable and has either the TARGET or POINTER attribute.
- 2 Within an ASSOCIATE, SELECT RANK, or SELECT TYPE construct, each associating entity has the same corank as its associated selector. If the selector is a coarray, the cobounds of each codimension of the associating entity are the same as those of the selector.
- 3 Within a CHANGE TEAM construct, the associating entity is a coarray. Its corank and cobounds are as specified in its codimension-decl.
- 4 Within an ASSOCIATE, CHANGE TEAM, SELECT RANK, or SELECT TYPE construct, the associating entity has the ASYNCHRONOUS or VOLATILE attribute if and only if the selector is a variable and has the attribute. If the associating entity is polymorphic, it assumes the dynamic type and type parameter values of the selector. If the selector has the OPTIONAL attribute, it shall be present. The associating entity is contiguous if and only if the selector is contiguous.
- 5 The associating entity itself is a variable, but if the selector is not a definable variable, the associating entity is not definable and shall not be defined or become undefined. If a selector is not permitted to appear in a variable definition context (19.6.7), neither the associate name nor any subobject thereof shall appear in a variable definition context.

### 11.1.3.4 Examples of the ASSOCIATE construct

#### NOTE 11.3

The following example illustrates an association with an expression.

```
ASSOCIATE ( Z => EXP (-(X**2+Y**2)) * COS (THETA) )
  PRINT *, A+Z, A-Z
END ASSOCIATE
```

The following example illustrates an association with a derived-type variable.

```
ASSOCIATE ( XC => AX%B(I,J)%C )
  XC%DV = XC%DV + PRODUCT (XC%EV(1:N))
END ASSOCIATE
```

The following example illustrates association with an array section.

```
ASSOCIATE ( ARRAY => AX%B(I,:)%C )
  ARRAY(N)%EV = ARRAY(N-1)%EV
END ASSOCIATE
```

The following example illustrates multiple associations.

```
ASSOCIATE ( W => RESULT(I,J)%W, ZX => AX%B(I,J)%D, ZY => AY%B(I,J)%D )
  W = ZX*X + ZY*Y
END ASSOCIATE
```

### 11.1.4 BLOCK construct

- 1 The BLOCK construct is an executable construct that may contain declarations.

- 1 R1107 *block-construct* is *block-stmt*  
 2 [ *block-specification-part* ]  
 3 *block*  
 4 *end-block-stmt*
- 5 R1108 *block-stmt* is [ *block-construct-name* : ] BLOCK
- 6 R1109 *block-specification-part* is [ *use-stmt* ]...  
 7 [ *import-stmt* ] ...  
 8 [ *implicit-part* ]  
 9 [ [ *declaration-construct* ] ...  
 10 *specification-construct* ]
- 11 R1110 *end-block-stmt* is END BLOCK [ *block-construct-name* ]
- 12 C1107 (R1107) A *block-specification-part* shall not contain a **COMMON**, **EQUIVALENCE**, **INTENT**, **NAMELIST**,  
 13 **OPTIONAL**, **statement function**, or **VALUE** statement.
- 14 C1108 (R1107) A **SAVE** statement in a BLOCK construct shall contain a *saved-entity-list* that does not specify a  
 15 *common-block-name*.
- 16 C1109 (R1107) If the *block-stmt* of a *block-construct* specifies a *block-construct-name*, the corresponding *end-*  
 17 *block-stmt* shall specify the same *block-construct-name*. If the *block-stmt* does not specify a *block-*  
 18 *construct-name*, the corresponding *end-block-stmt* shall not specify a *block-construct-name*.
- 19 2 Except for the **ASYNCHRONOUS** and **VOLATILE** statements, specifications in a BLOCK construct declare  
 20 **construct entities** whose scope is that of the BLOCK construct (19.4). The appearance of the name of an object  
 21 that is not a **construct entity** in an **ASYNCHRONOUS** or **VOLATILE** statement in a BLOCK construct specifies  
 22 that the object has the attribute within the construct even if it does not have the attribute outside the construct.
- 23 3 Execution of a BLOCK construct causes evaluation of the **specification expressions** within its specification part  
 24 in a processor-dependent order, followed by execution of its block.
- 25 4 It is permissible to branch to an *end-block-stmt* only from within its BLOCK construct.

**NOTE 11.4**

The following is an example of a BLOCK construct.

```

IF (swapxy) THEN
  BLOCK
    REAL (KIND (x)) tmp
    tmp = x
    x = y
    y = tmp
  END BLOCK
END IF

```

Actions on a variable local to a BLOCK construct do not affect any variable of the same name outside the construct. For example,

```

F = 254E-2
BLOCK
  REAL F
  F = 39.37
END BLOCK
! F is still equal to 254E-2.

```

A **SAVE** statement outside a BLOCK construct does not affect variables local to the BLOCK construct,

## NOTE 11.4 (cont.)

because a [SAVE statement](#) affects variables in its [scoping unit](#) rather than in its [inclusive scope](#). For example,

```

SUBROUTINE S
  ...
  SAVE
  ...
  BLOCK
    REAL X           ! Not saved.
    REAL,SAVE :: Y(100) ! SAVE attribute is allowed.
    Z = 3           ! Implicitly declared in S, thus saved.
    ...
  END BLOCK
  ...
END SUBROUTINE

```

## 11.1.5 CHANGE TEAM construct

## 11.1.5.1 Purpose and form of the CHANGE TEAM construct

1 The CHANGE TEAM construct changes the [current team](#). Named [construct entities](#) (19.4) can be associated (19.5.1.6) with [coarrays](#) in the containing [scoping unit](#), in the same way as for the [ASSOCIATE construct](#).

R1111 *change-team-construct* is *change-team-stmt*  
*block*  
*end-change-team-stmt*

R1112 *change-team-stmt* is [ *team-construct-name* : ] CHANGE TEAM ( *team-value* ■  
 ■ [ , *coarray-association-list* ] [ , *sync-stat-list* ] )

R1113 *coarray-association* is *codimension-decl* => *selector*

R1114 *end-change-team-stmt* is END TEAM [ ( [ *sync-stat-list* ] ) ] [ *team-construct-name* ]

R1115 *team-value* is *scalar-expr*

C1110 A branch (11.2) within a CHANGE TEAM construct shall not have a branch target that is outside the construct.

C1111 A RETURN statement shall not appear within a CHANGE TEAM construct.

C1112 If the *change-team-stmt* of a *change-team-construct* specifies a *team-construct-name*, the corresponding *end-change-team-stmt* shall specify the same *team-construct-name*. If the *change-team-stmt* of a *change-team-construct* does not specify a *team-construct-name*, the corresponding *end-change-team-stmt* shall not specify a *team-construct-name*.

C1113 In a *change-team-stmt*, a *coarray-name* in a *codimension-decl* shall not be the same as a *selector*, or another *coarray-name*, in that statement.

C1114 A *team-value* shall be of type `TEAM_TYPE` from the intrinsic module `ISO_FORTRAN_ENV`.

C1115 No *selector* shall appear more than once in a given *change-team-stmt*.

C1116 A *selector* in a *coarray-association* shall be a named coarray.

2 Each *coarray-name* in a *codimension-decl* in the CHANGE TEAM statement is an [associate name](#) which is associated with the corresponding selector. Each [associating entity](#) assumes the type and [type parameters](#) of

1 its selector; it is polymorphic if and only if the selector is polymorphic. The other attributes of the **associating**  
2 **entities** are described in 11.1.3.3.

### 3 11.1.5.2 Execution of a CHANGE TEAM construct

- 4 1 The *team-values* on the **active images** that execute the CHANGE TEAM statement shall be those of **team variables**  
5 defined by corresponding executions of the same **FORM TEAM statement** (11.6.9). When the CHANGE TEAM  
6 statement is executed, the **current team** shall be the team that was current when those **team variables** were defined.  
7 The current team for the statements of the CHANGE TEAM *block* is the team identified by the *team-value*. If  
8 *team-value* is a variable, the variable shall not be defined or become undefined during execution of the CHANGE  
9 TEAM construct. A CHANGE TEAM construct completes execution by executing its END TEAM statement,  
10 which restores the **current team** to the **original team that was current** for the CHANGE TEAM statement.
- 11 2 Execution of a CHANGE TEAM construct causes evaluation of the expressions within each *codimension-decl* in  
12 the CHANGE TEAM statement, followed by execution of its block. Each *selector* shall be an **established coarray**  
13 when the CHANGE TEAM statement begins execution.
- 14 3 It is permissible to branch to an *end-change-team-stmt* only from within its CHANGE TEAM construct.
- 15 4 An allocatable **coarray** that was allocated immediately before executing a CHANGE TEAM statement shall not  
16 be deallocated during execution of the construct. An allocatable **coarray** that was unallocated immediately before  
17 executing a CHANGE TEAM statement, and which is allocated immediately before executing the corresponding  
18 END TEAM statement, is deallocated by the execution of the END TEAM statement.
- 19 5 Successful execution of a CHANGE TEAM statement performs an implicit synchronization of all **images** of the  
20 new **team** that is identified by *team-value*. All **active images** of the new **team** shall execute the same CHANGE  
21 TEAM statement. On each **image** of the new **team**, execution of the segment following the CHANGE TEAM  
22 statement is delayed until all other **images** of that **team** have executed the same statement the same number of  
23 times in the original **team**.
- 24 6 If the new **team** contains a **failed image** and no other error condition occurs, there is an implicit synchronization  
25 of all **active images** of the new **team**. On each **active image** of the new **team**, execution of the segment following  
26 the CHANGE TEAM statement is delayed until all other **active images** of that **team** have executed the same  
27 statement the same number of times in the original **team**.
- 28 7 When a CHANGE TEAM construct completes execution, there is an implicit synchronization of all **active images**  
29 in the new **team**. On each **active image** of the new **team**, execution of the segment following the END TEAM  
30 statement is delayed until all other **active images** of this **team** have executed the same construct the same number  
31 of times in this team.

#### NOTE 11.5

Deallocation of an allocatable coarray that was not allocated at the beginning of a CHANGE TEAM construct, but is allocated at the end of execution of the construct, occurs even for allocatable coarrays with the SAVE attribute.

#### NOTE 11.6

Execution of a CHANGE TEAM statement includes a synchronization of the executing **image** with the other **images** that will be in the same **team** after execution of the CHANGE TEAM statement. Synchronization of these images occurs again when the corresponding END TEAM statement is executed.

If it is desired to synchronize all of the **images** in the **team that was current** when the CHANGE TEAM statement was executed, a **SYNC TEAM statement** that specifies the parent team can be executed immediately after the CHANGE TEAM statement. If similar semantics are desired following the END TEAM statement, a **SYNC ALL statement** could immediately follow the END TEAM statement.



### 11.1.6 CRITICAL construct

1 A CRITICAL construct limits execution of a block to one **image** at a time.

R1116 *critical-construct* is *critical-stmt*  
*block*  
*end-critical-stmt*

R1117 *critical-stmt* is [ *critical-construct-name* : ] CRITICAL [ ( [ *sync-stat-list* ] ) ]

R1118 *end-critical-stmt* is END CRITICAL [ *critical-construct-name* ]

C1117 (R1116) If the *critical-stmt* of a *critical-construct* specifies a *critical-construct-name*, the corresponding *end-critical-stmt* shall specify the same *critical-construct-name*. If the *critical-stmt* of a *critical-construct* does not specify a *critical-construct-name*, the corresponding *end-critical-stmt* shall not specify a *critical-construct-name*.

C1118 (R1116) The *block* of a *critical-construct* shall not contain a **RETURN statement** or an **image control statement**.

C1119 A branch (11.2) within a **CRITICAL construct** shall not have a branch target that is outside the construct.

2 Execution of the CRITICAL construct is completed when execution of its block is completed, or the executing image fails (5.3.6). A procedure invoked, directly or indirectly, from a CRITICAL construct shall not execute an **image control statement**.

3 The processor shall ensure that once an **image** has commenced executing *block*, no other **image** shall commence executing *block* until this **image** has completed execution of the construct. The **image** shall not execute an **image control statement** during the execution of *block*. The sequence of executed statements is therefore a segment (11.6.2). If **image** M completes execution of the construct without failing and **image** T is the next to execute the construct, the segment on **image** M precedes the segment on **image** T. Otherwise, if **image** M completes execution of the construct by failing, and **image** T is the next to execute the construct, the previous segment on **image** M precedes the segment on **image** T.

4 The effect of a **STAT=** or **ERRMSG=** specifier in a CRITICAL statement is specified in subclause 11.6.11.

5 It is permissible to branch to an *end-critical-stmt* only from within its CRITICAL construct.

#### NOTE 11.7

If more than one **image** executes the block of a CRITICAL construct without failing, its execution by one **image** always either precedes or succeeds its execution by another nonfailed **image**. Typically no other statement ordering is needed. Consider the following example:

```
CRITICAL
  GLOBAL_COUNTER[1] = GLOBAL_COUNTER[1] + 1
END CRITICAL
```

The definition of GLOBAL\_COUNTER [1] by a particular **image** will always precede the reference to the same variable by the next **image** to execute the block.

#### NOTE 11.8

The following example permits a large number of jobs to be shared among the **images**:

```
INTEGER :: NUM_JOBS[*], JOB

IF (THIS_IMAGE() == 1) READ(*,*) NUM_JOBS
SYNC ALL
```

## NOTE 11.8 (cont.)

```

DO
  CRITICAL
    JOB = NUM_JOBS[1]
    NUM_JOBS[1] = JOB - 1
  END CRITICAL
  IF (JOB > 0) THEN
    ! Work on JOB
  ELSE
    EXIT
  END IF
END DO
SYNC ALL

```

## 11.1.7 DO construct

## 11.1.7.1 Purpose and form of the DO construct

1 The DO construct specifies the repeated execution of a sequence of executable constructs. Such a repeated sequence is called a loop.

2 The number of iterations of a loop can be determined at the beginning of execution of the DO construct, or can be left indefinite (“DO forever” or DO WHILE). The execution order of the iterations can be left indeterminate (DO CONCURRENT); except in this case, the loop can be terminated immediately (11.1.7.4.5). An iteration of the loop can be curtailed by executing a [CYCLE statement](#) (11.1.7.4.4).

3 There are three phases in the execution of a DO construct: initiation of the loop, execution of each iteration of the loop, and termination of the loop.

4 The scope and attributes of an *index-name* in a *concurrent-header* (DO CONCURRENT) are described in 19.4.

## 11.1.7.2 Form of the DO construct

R1119 *do-construct* is *do-stmt*  
*block*  
*end-do*

R1120 *do-stmt* is *nonlabel-do-stmt*  
or *label-do-stmt*

R1121 *label-do-stmt* is [ *do-construct-name* : ] DO *label* [ *loop-control* ]

R1122 *nonlabel-do-stmt* is [ *do-construct-name* : ] DO [ *loop-control* ]

R1123 *loop-control* is [ , ] *do-variable* = *scalar-int-expr*, *scalar-int-expr* ■  
■ [ , *scalar-int-expr* ]  
or [ , ] WHILE ( *scalar-logical-expr* )  
or [ , ] CONCURRENT *concurrent-header* *concurrent-locality*

R1124 *do-variable* is *scalar-int-variable-name*

C1120 (R1124) The *do-variable* shall be a variable of type integer.

R1125 *concurrent-header* is ( [ *integer-type-spec* :: ] *concurrent-control-list* [ , *scalar-mask-expr* ] )

R1126 *concurrent-control* is *index-name* = *concurrent-limit* : *concurrent-limit* [ : *concurrent-step* ]

R1127 *concurrent-limit* is *scalar-int-expr*

- 1 R1128 *concurrent-step* is *scalar-int-expr*
- 2 R1129 *concurrent-locality* is [ *locality-spec* ]...
- 3 R1130 *locality-spec* is LOCAL ( *variable-name-list* )  
 4 or LOCAL\_INIT ( *variable-name-list* )  
 5 or SHARED ( *variable-name-list* )  
 6 or DEFAULT ( NONE )
- 7 C1121 (R1125) Any procedure referenced in the *scalar-mask-expr*, including one referenced by a [defined operation](#), shall be a [pure procedure](#) (15.7).
- 8
- 9 C1122 (R1126) The *index-name* shall be a named scalar variable of type integer.
- 10 C1123 (R1126) A *concurrent-limit* or *concurrent-step* in a *concurrent-control* shall not contain a reference to  
 11 any *index-name* in the *concurrent-control-list* in which it appears.
- 12 C1124 A *variable-name* in a *locality-spec* shall be the name of a variable in the innermost executable construct  
 13 or [scoping unit](#) that includes the DO CONCURRENT statement.
- 14 C1125 A *variable-name* in a *locality-spec* shall not be the same as an *index-name* in the *concurrent-header* of  
 15 the same DO CONCURRENT statement.
- 16 C1126 The name of a variable shall not appear in more than one *variable-name-list*, or more than once in a  
 17 *variable-name-list*, in a *locality-spec*.
- 18 C1127 A *variable-name* that appears in a LOCAL or LOCAL\_INIT *locality-spec* shall not have the [ALLOCAT-](#)  
 19 [ABLE](#), [INTENT \(IN\)](#), or [OPTIONAL](#) attribute, shall not be of [finalizable](#) type, shall not be a nonpointer  
 20 [polymorphic dummy argument](#), and shall not be a [coarray](#) or an [assumed-size array](#). A *variable-name* that  
 21 is not permitted to appear in a variable definition context shall not appear in a LOCAL or LOCAL\_INIT  
 22 *locality-spec*.
- 23 C1128 A variable that is referenced by the *scalar-mask-expr* of a *concurrent-header* or by any *concurrent-limit*  
 24 or *concurrent-step* in that *concurrent-header* shall not appear in a LOCAL *locality-spec* in the same DO  
 25 CONCURRENT statement.
- 26 C1129 If the *locality-spec* DEFAULT ( NONE ) appears in a DO CONCURRENT statement, a variable that  
 27 appears in the *block* of the construct and is not an *index-name* of that construct shall have its locality  
 28 explicitly specified.
- 29 R1131 *end-do* is *end-do-stmt*  
 30 or *continue-stmt*
- 31 R1132 *end-do-stmt* is END DO [ *do-construct-name* ]
- 32 C1130 (R1119) If the *do-stmt* of a *do-construct* specifies a *do-construct-name*, the corresponding *end-do* shall be  
 33 an *end-do-stmt* specifying the same *do-construct-name*. If the *do-stmt* of a *do-construct* does not specify  
 34 a *do-construct-name*, the corresponding *end-do* shall not specify a *do-construct-name*.
- 35 C1131 (R1119) If the *do-stmt* is a *nonlabel-do-stmt*, the corresponding *end-do* shall be an *end-do-stmt*.
- 36 C1132 (R1119) If the *do-stmt* is a *label-do-stmt*, the corresponding *end-do* shall be identified with the same *label*.
- 37 1 It is permissible to branch to an *end-do* only from within its DO construct.

### 11.1.7.3 Active and inactive DO constructs

- 39 1 A DO construct is either active or inactive. Initially inactive, a DO construct becomes active only when its [DO](#)  
 40 [statement](#) is executed.

1 2 Once active, the DO construct becomes inactive only when it terminates (11.1.7.4.5).

#### 2 11.1.7.4 Execution of a DO construct

##### 3 11.1.7.4.1 Loop initiation

4 1 When the DO statement is executed, the DO construct becomes active. If *loop-control* is

5  $[ , ] \textit{do-variable} = \textit{scalar-int-expr}_1 , \textit{scalar-int-expr}_2 [ , \textit{scalar-int-expr}_3 ]$

6 the following steps are performed in sequence.

- 7 (1) The initial parameter  $m_1$ , the terminal parameter  $m_2$ , and the incrementation parameter  $m_3$  are  
8 of type integer with the same kind type parameter as the *do-variable*. Their values are established  
9 by evaluating *scalar-int-expr*<sub>1</sub>, *scalar-int-expr*<sub>2</sub>, and *scalar-int-expr*<sub>3</sub>, respectively, including, if ne-  
10 cessary, conversion to the kind type parameter of the *do-variable* according to the rules for numeric  
11 conversion (Table 10.9). If *scalar-int-expr*<sub>3</sub> does not appear,  $m_3$  has the value 1. The value of  $m_3$   
12 shall not be zero.
- 13 (2) The DO variable becomes defined with the value of the initial parameter  $m_1$ .
- 14 (3) The iteration count is established and is the value of the expression  $(m_2 - m_1 + m_3)/m_3$ , unless that  
15 value is negative, in which case the iteration count is 0.

#### NOTE 11.9

The iteration count is zero whenever:

$$m_1 > m_2 \text{ and } m_3 > 0, \text{ or}$$

$$m_1 < m_2 \text{ and } m_3 < 0.$$

16 2 If *loop-control* is omitted, no iteration count is calculated. The effect is as if a large positive iteration count,  
17 impossible to decrement to zero, were established. If *loop-control* is  $[ , ]$  WHILE (*scalar-logical-expr*), the effect  
18 is as if *loop-control* were omitted and the following statement inserted as the first statement of the *block*:

19 IF (.NOT. (*scalar-logical-expr*)) EXIT

20 3 For a DO CONCURRENT construct, the values of the index variables for the iterations of the construct are  
21 determined by the rules in 11.1.7.4.2.

22 4 At the completion of the execution of the DO statement, the execution cycle begins.

##### 23 11.1.7.4.2 DO CONCURRENT loop control

24 1 The *concurrent-limit* and *concurrent-step* expressions in the *concurrent-control-list* are evaluated. These ex-  
25 pressions may be evaluated in any order. The set of values that a particular *index-name* variable assumes is  
26 determined as follows.

- 27 (1) The lower bound  $m_1$ , the upper bound  $m_2$ , and the step  $m_3$  are of type integer with the same kind  
28 type parameter as the *index-name*. Their values are established by evaluating the first *concurrent-*  
29 *limit*, the second *concurrent-limit*, and the *concurrent-step* expressions, respectively, including, if  
30 necessary, conversion to the kind type parameter of the *index-name* according to the rules for numeric  
31 conversion (Table 10.9). If *concurrent-step* does not appear,  $m_3$  has the value 1. The value  $m_3$  shall  
32 not be zero.
- 33 (2) Let the value of *max* be  $(m_2 - m_1 + m_3)/m_3$ . If  $\textit{max} \leq 0$  for some *index-name*, the execution of the  
34 construct is complete. Otherwise, the set of values for the *index-name* is  
35 
$$m_1 + (k - 1) \times m_3 \quad \text{where } k = 1, 2, \dots, \textit{max}.$$

36 2 The set of combinations of *index-name* values is the Cartesian product of the sets defined by each triplet specific-  
37 ation. An *index-name* becomes defined when this set is evaluated.

- 1 3 The *scalar-mask-expr*, if any, is evaluated for each combination of *index-name* values. If there is no *scalar-*  
 2 *mask-expr*, it is as if it appeared with the value true. The *index-name* variables may be primaries in the  
 3 *scalar-mask-expr*.
- 4 4 The set of active combinations of *index-name* values is the subset of all possible combinations for which the  
 5 *scalar-mask-expr* has the value true.

**NOTE 11.10**

The *index-name* variables can appear in the mask, for example

```
DO CONCURRENT (I=1:10, J=1:10, A(I) > 0.0 .AND. B(J) < 1.0)
. . .
```

6 **11.1.7.4.3 The execution cycle**

- 7 1 The execution cycle of a DO construct that is not a DO CONCURRENT construct consists of the following steps  
 8 performed in sequence repeatedly until termination.

- 9 (1) The iteration count, if any, is tested. If it is zero, the loop terminates and the DO construct becomes  
 10 inactive. If *loop-control* is [ , ] WHILE (*scalar-logical-expr*), the *scalar-logical-expr* is evaluated; if  
 11 the value of this expression is false, the loop terminates and the DO construct becomes inactive.
- 12 (2) The *block* of the loop is executed.
- 13 (3) The iteration count, if any, is decremented by one. The DO variable, if any, is incremented by the  
 14 value of the incrementation parameter  $m_3$ .

- 15 2 Except for the incrementation of the DO variable that occurs in step (3), the DO variable shall neither be redefined  
 16 nor become undefined while the DO construct is active.

- 17 3 The *block* of a DO CONCURRENT construct is executed for every active combination of the *index-name* values.  
 18 Each execution of the *block* is an iteration. The executions may occur in any order.

19 **11.1.7.4.4 CYCLE statement**

- 20 1 Execution of a loop iteration can be curtailed by executing a CYCLE statement that belongs to the construct.

21 R1133 *cycle-stmt* is CYCLE [ *do-construct-name* ]

- 22 C1133 If a *do-construct-name* appears on a CYCLE statement, the CYCLE statement shall be within that  
 23 *do-construct*; otherwise, it shall be within at least one *do-construct*.

- 24 C1134 A *cycle-stmt* shall not appear within a CHANGE TEAM, CRITICAL, or DO CONCURRENT construct  
 25 if it belongs to an outer construct.

- 26 2 A CYCLE statement belongs to a particular DO construct. If the CYCLE statement contains a DO construct  
 27 name, it belongs to that DO construct; otherwise, it belongs to the innermost DO construct in which it appears.

- 28 3 Execution of a CYCLE statement that belongs to a DO construct that is not a DO CONCURRENT construct  
 29 causes immediate progression to step (3) of the execution cycle of the DO construct to which it belongs.

- 30 4 Execution of a CYCLE statement that belongs to a DO CONCURRENT construct completes execution of that  
 31 iteration of the construct.

- 32 5 In a DO construct, a transfer of control to the *end-do* has the same effect as execution of a CYCLE statement  
 33 belonging to that construct.

34 **11.1.7.4.5 Loop termination**

- 35 1 For a DO construct that is not a DO CONCURRENT construct, the loop terminates, and the DO construct  
 36 becomes inactive, when any of the following occurs.

- 1       • The iteration count is determined to be zero or the *scalar-logical-expr* is false, when tested during step (1)
- 2       of the above execution cycle.
- 3       • An **EXIT statement** that belongs to the DO construct is executed.
- 4       • An **EXIT** or **CYCLE** statement that belongs to an outer construct and is within the DO construct is
- 5       executed.
- 6       • A branch occurs within the DO construct and the **branch target statement** is outside the construct.
- 7       • A **RETURN statement** within the DO construct is executed.

8   2 For a DO CONCURRENT construct, the loop terminates, and the DO construct becomes inactive when all of  
9   the iterations have completed execution.

10 3 When a DO construct becomes inactive, the DO variable, if any, of the DO construct retains its last defined  
11 value.

#### 12 11.1.7.5 Additional semantics for DO CONCURRENT constructs

13 C1135 A **RETURN statement** shall not appear within a DO CONCURRENT construct.

14 C1136 An **image control statement** shall not appear within a DO CONCURRENT construct.

15 C1137 A branch (11.2) within a DO CONCURRENT construct shall not have a branch target that is outside  
16 the construct.

17 C1138 A reference to an impure procedure shall not appear within a DO CONCURRENT construct.

18 C1139 A statement that might result in the deallocation of a **polymorphic** entity shall not appear within a DO  
19 CONCURRENT construct.

20 C1140 A reference to the procedure **IEEE\_GET\_FLAG**, **IEEE\_SET\_HALTING\_MODE**, or **IEEE\_GET\_-**  
21 **HALTING\_MODE** from the intrinsic module **IEEE\_EXCEPTIONS**, shall not appear within a DO  
22 CONCURRENT construct.

23 1 The locality of a variable that appears in a DO CONCURRENT construct is **LOCAL**, **LOCAL\_INIT**, **SHARED**,  
24 or unspecified.

25 2 A variable that has **LOCAL** or **LOCAL\_INIT** locality is a construct entity with the same type, type parameters,  
26 rank as the variable with the same name in the innermost executable construct or **scoping unit** that includes the  
27 DO CONCURRENT construct, and the outside variable is inaccessible by that name within the construct. The  
28 construct entity has the **ASYNCHRONOUS**, **CONTIGUOUS**, **POINTER**, **TARGET**, or **VOLATILE** attribute if  
29 and only if the outside variable has that attribute; it does not have the **BIND**, **INTENT**, **PROTECTED**, **SAVE**  
30 or **VALUE** attribute, even if the outside variable has that attribute. If it is not a pointer, it has the same bounds  
31 as the outside variable. At the beginning of execution of each iteration,

- 32       • if a variable with **LOCAL** locality is a pointer it has undefined pointer association status, and otherwise it
- 33       is undefined except for any subobjects that are default-initialized;
- 34       • a variable with **LOCAL\_INIT** locality has the pointer association status and definition status of the out-
- 35       side variable with that name; the outside variable shall not be an undefined pointer or a nonallocatable
- 36       nonpointer variable that is undefined.

37 If a variable with **LOCAL** or **LOCAL\_INIT** locality becomes an affector of a pending input/output operation,  
38 the operation shall have completed before the end of the iteration. If a variable with **LOCAL** or **LOCAL\_INIT**  
39 locality has the **TARGET attribute**, a pointer associated with it during an iteration becomes undefined when  
40 execution of that iteration completes.

41 3 If a variable has **SHARED** locality, appearances of the variable within the DO CONCURRENT construct refer  
42 to the variable in the innermost executable construct or **scoping unit** that includes the DO CONCURRENT  
43 construct. If it is defined or becomes undefined during any iteration, it shall not be referenced, defined, or  
44 become undefined during any other iteration. If it is allocated, deallocated, nullified, or pointer-assigned during  
45 an iteration it shall not have its allocation or **association** status, dynamic type, array bounds, shape, or a **deferred**

- 1 [type parameter](#) value inquired about in any other iteration. A noncontiguous array with [SHARED](#) locality shall  
2 not be supplied as an [actual argument](#) corresponding to a [contiguous INTENT \(INOUT\) dummy argument](#).
- 3 4 If a variable has unspecified locality,
- 4 • if it is referenced in an iteration it shall either be previously defined during that iteration, or shall not be  
5 defined or become undefined during any other iteration; if it is defined or becomes undefined by more than  
6 one iteration it becomes undefined when the loop terminates;
  - 7 • if it is noncontiguous and is supplied as an [actual argument](#) corresponding to a [contiguous INTENT \(IN-  
8 OUT\) dummy argument](#) in an iteration, it shall either be previously defined in that iteration or shall not  
9 be defined in any other iteration;
  - 10 • if it is a pointer and is used in an iteration other than as the pointer in pointer assignment, allocation,  
11 or nullification, it shall either be previously [pointer associated](#) during that iteration or shall not have its  
12 [pointer association](#) changed during any iteration;
  - 13 • if it is a pointer whose [pointer association](#) is changed in more than one iteration, it has an association status  
14 of undefined when the construct terminates;
  - 15 • if it is [allocatable](#) and is allocated in more than one iteration, it shall have an allocation status of unallocated  
16 at the end of every iteration;
  - 17 • if it is [allocatable](#) and is referenced, defined, deallocated, or has its allocation status, dynamic type, or a  
18 deferred type parameter value inquired about, in any iteration, it shall either be previously allocated in  
19 that iteration or shall not be allocated or deallocated in any other iteration.
- 20 5 A DO CONCURRENT construct shall not contain an input/output statement that has an ADVANCE= specifier.
- 21 6 If data are written to a file record or position in one iteration, that record or position in that file shall not be  
22 read from or written to in a different iteration. If records are written to a file connected for sequential access by  
23 more than one iteration, the ordering of records written by different iterations is processor dependent.

**NOTE 11.11**

The restrictions on referencing variables defined in an iteration of a DO CONCURRENT construct apply to any procedure invoked within the loop.

**NOTE 11.12**

The restrictions on the statements in a DO CONCURRENT construct are designed to ensure there are no data dependencies between iterations of the loop. This permits code optimizations that might otherwise be difficult or impossible because they would depend on properties of the program not visible to the compiler.

24 **11.1.7.6 Examples of DO constructs****NOTE 11.13**

The following program fragment computes a tensor product of two arrays:

```
DO I = 1, M
  DO J = 1, N
    C (I, J) = DOT_PRODUCT (A (I, J, :), B(:, I, J))
  END DO
END DO
```

**NOTE 11.14**

The following program fragment contains a DO construct that uses the [WHILE](#) form of [loop-control](#). The loop will continue to execute until an end-of-file or input/output error is encountered, at which point the [DO statement](#) terminates the loop. When a negative value of X is read, the program skips immediately to the next [READ statement](#), bypassing most of the [block](#) of the loop.

**NOTE 11.14 (cont.)**

```

READ (IUN, '(1X, G14.7)', IOSTAT = IOS) X
DO WHILE (IOS == 0)
  IF (X >= 0.) THEN
    CALL SUBA (X)
    CALL SUBB (X)
    . . .
    CALL SUBZ (X)
  ENDIF
  READ (IUN, '(1X, G14.7)', IOSTAT = IOS) X
END DO

```

**NOTE 11.15**

The following example behaves exactly the same as the one in Note 11.14. However, the [READ statement](#) has been moved to the interior of the loop, so that only one [READ statement](#) is needed. Also, a [CYCLE statement](#) has been used to avoid an extra level of [IF](#) nesting.

```

DO      ! A "DO WHILE + 1/2" loop
  READ (IUN, '(1X, G14.7)', IOSTAT = IOS) X
  IF (IOS /= 0) EXIT
  IF (X < 0.) CYCLE
  CALL SUBA (X)
  CALL SUBB (X)
  . . .
  CALL SUBZ (X)
END DO

```

**NOTE 11.16**

The following example illustrates a case in which the user knows that there are no repeated values in the index array IND. The DO CONCURRENT construct makes it easier for the processor to generate vector gather/scatter code, unroll the loop, or parallelize the code for this loop, potentially improving performance.

```

INTEGER :: A(N), IND(N)

DO CONCURRENT (I=1:M)
  A(IND(I)) = I
END DO

```

**NOTE 11.17**

The following code demonstrates the use of the [LOCAL](#) clause so that the X inside the [DO CONCURRENT construct](#) is a temporary variable, and will not affect the X outside the construct.

```

X = 1.0
DO CONCURRENT (I=1:10) LOCAL (X)
  IF (A (I) > 0) THEN
    X = SQRT (A (I))
    A (I) = A (I) - X**2
  END IF
  B (I) = B (I) - A (I)
END DO
PRINT *, X                                ! Always prints 1.0.

```



**NOTE 11.18**

Additional examples of DO constructs are in [C.6.3](#).

**11.1.8 IF construct and statement****11.1.8.1 Purpose and form of the IF construct**

- 1 The IF construct selects for execution at most one of its constituent blocks. The selection is based on a sequence of logical expressions.

```

R1134 if-construct           is if-then-stmt
                                block
                                [ else-if-stmt
                                block ] ...
                                [ else-stmt
                                block ]
                                end-if-stmt

R1135 if-then-stmt          is [ if-construct-name : ] IF ( scalar-logical-expr ) THEN

R1136 else-if-stmt         is ELSE IF ( scalar-logical-expr ) THEN [ if-construct-name ]

R1137 else-stmt           is ELSE [ if-construct-name ]

R1138 end-if-stmt         is END IF [ if-construct-name ]

```

C1141 (R1134) If the *if-then-stmt* of an *if-construct* specifies an *if-construct-name*, the corresponding *end-if-stmt* shall specify the same *if-construct-name*. If the *if-then-stmt* of an *if-construct* does not specify an *if-construct-name*, the corresponding *end-if-stmt* shall not specify an *if-construct-name*. If an *else-if-stmt* or *else-stmt* specifies an *if-construct-name*, the corresponding *if-then-stmt* shall specify the same *if-construct-name*.

**11.1.8.2 Execution of an IF construct**

- 1 At most one of the blocks in the IF construct is executed. If there is an ELSE statement in the construct, exactly one of the blocks in the construct is executed. The scalar logical expressions are evaluated in the order of their appearance in the construct until a true value is found or an ELSE statement or END IF statement is encountered. If a true value or an ELSE statement is found, the block immediately following is executed and this completes the execution of the construct. The scalar logical expressions in any remaining ELSE IF statements of the IF construct are not evaluated. If none of the evaluated expressions is true and there is no ELSE statement, the execution of the construct is completed without the execution of any block within the construct.
- 2 It is permissible to branch to an END IF statement only from within its IF construct. Execution of an END IF statement has no effect.

**11.1.8.3 Examples of IF constructs****NOTE 11.19**

```

IF (CVAR == 'RESET') THEN
  I = 0; J = 0; K = 0
END IF
PROOF_DONE: IF (PROP) THEN
  WRITE (3, '( 'QED' )')
  STOP
ELSE

```

## NOTE 11.19 (cont.)

```

PROP = NEXTPROP
END IF PROOF_DONE
IF (A > 0) THEN
  B = C/A
  IF (B > 0) THEN
    D = 1.0
  END IF
ELSE IF (C > 0) THEN
  B = A/C
  D = -1.0
ELSE
  B = ABS (MAX (A, C))
  D = 0
END IF

```

1 **11.1.8.4 IF statement**

2 1 The IF statement controls the execution of a single action statement based on a single logical expression.

3 R1139 *if-stmt* is IF ( *scalar-logical-expr* ) *action-stmt*

4 C1142 (R1139) The *action-stmt* in the *if-stmt* shall not be an *if-stmt*.

5 2 Execution of an IF statement causes evaluation of the scalar logical expression. If the value of the expression is  
6 true, the action statement is executed. If the value is false, the action statement is not executed.

7 3 The execution of a function reference in the scalar logical expression may affect entities in the action statement.

## NOTE 11.20

An example of an IF statement is:

```
IF (A > 0.0) A = LOG (A)
```

8 **11.1.9 SELECT CASE construct**9 **11.1.9.1 Purpose and form of the SELECT CASE construct**

10 1 The SELECT CASE construct selects for execution at most one of its constituent blocks. The selection is based  
11 on the value of an expression.

12 R1140 *case-construct* is *select-case-stmt*  
13 [ *case-stmt*  
14 *block* ] ...  
15 *end-select-stmt*

16 R1141 *select-case-stmt* is [ *case-construct-name* : ] SELECT CASE ( *case-expr* )

17 R1142 *case-stmt* is CASE *case-selector* [ *case-construct-name* ]

18 R1143 *end-select-stmt* is END SELECT [ *case-construct-name* ]

19 C1143 (R1140) If the *select-case-stmt* of a *case-construct* specifies a *case-construct-name*, the corresponding *end-*  
20 *select-stmt* shall specify the same *case-construct-name*. If the *select-case-stmt* of a *case-construct* does  
21 not specify a *case-construct-name*, the corresponding *end-select-stmt* shall not specify a *case-construct-*  
22 *name*. If a *case-stmt* specifies a *case-construct-name*, the corresponding *select-case-stmt* shall specify the  
23 same *case-construct-name*.

1 R1144 *case-expr* is *scalar-expr*

2 C1144 *case-expr* shall be of type character, integer, or logical.

3 R1145 *case-selector* is ( *case-value-range-list* )  
4 or DEFAULT

5 C1145 (R1140) No more than one of the selectors of one of the CASE statements shall be DEFAULT.

6 R1146 *case-value-range* is *case-value*  
7 or *case-value* :  
8 or : *case-value*  
9 or *case-value* : *case-value*

10 R1147 *case-value* is *scalar-constant-expr*

11 C1146 (R1140) For a given *case-construct*, each *case-value* shall be of the same type as *case-expr*. For character  
12 type, the *kind type parameters* shall be the same; character length differences are allowed.

13 C1147 (R1140) A *case-value-range* using a colon shall not be used if *case-expr* is of type logical.

14 C1148 (R1140) For a given *case-construct*, there shall be no possible value of the *case-expr* that matches more  
15 than one *case-value-range*.

#### 16 11.1.9.2 Execution of a SELECT CASE construct

17 1 The execution of the SELECT CASE statement causes the case expression to be evaluated. For a case value  
18 range list, a match occurs if the case expression value matches any of the case value ranges in the list. For a case  
19 expression with a value of *c*, a match is determined as follows.

- 20 (1) If the case value range contains a single value *v* without a colon, a match occurs for type logical if  
21 the expression *c* .EQV. *v* is true, and a match occurs for type integer or character if the expression  
22 *c* == *v* is true.
- 23 (2) If the case value range is of the form *low* : *high*, a match occurs if the expression *low* <= *c* .AND.  
24 *c* <= *high* is true.
- 25 (3) If the case value range is of the form *low* :, a match occurs if the expression *low* <= *c* is true.
- 26 (4) If the case value range is of the form : *high*, a match occurs if the expression *c* <= *high* is true.
- 27 (5) If no other selector matches and a DEFAULT selector appears, it matches the case index.
- 28 (6) If no other selector matches and the DEFAULT selector does not appear, there is no match.

29 2 The block following the CASE statement containing the matching selector, if any, is executed. This completes  
30 execution of the construct.

31 3 It is permissible to branch to an *end-select-stmt* only from within its SELECT CASE construct.

#### 32 11.1.9.3 Examples of SELECT CASE constructs

##### NOTE 11.21

An integer signum function:

```
INTEGER FUNCTION SIGNUM (N)
SELECT CASE (N)
CASE (:-1)
    SIGNUM = -1
CASE (0)
    SIGNUM = 0
CASE (1:)
```

**NOTE 11.21 (cont.)**

```
SIGNUM = 1
END SELECT
END
```

**NOTE 11.22**

A code fragment to check for balanced parentheses:

```
CHARACTER (80) :: LINE
...
LEVEL = 0
SCAN_LINE: DO I = 1, 80
  CHECK_PARENS: SELECT CASE (LINE (I:I))
    CASE ('(')
      LEVEL = LEVEL + 1
    CASE (')')
      LEVEL = LEVEL - 1
      IF (LEVEL < 0) THEN
        PRINT *, 'UNEXPECTED RIGHT PARENTHESIS'
        EXIT SCAN_LINE
      END IF
    CASE DEFAULT
      ! Ignore all other characters
  END SELECT CHECK_PARENS
END DO SCAN_LINE
IF (LEVEL > 0) THEN
  PRINT *, 'MISSING RIGHT PARENTHESIS'
END IF
```

**NOTE 11.23**

The following three fragments are equivalent:

```
IF (SILLY == 1) THEN
  CALL THIS
ELSE
  CALL THAT
END IF
SELECT CASE (SILLY == 1)
CASE (.TRUE.)
  CALL THIS
CASE (.FALSE.)
  CALL THAT
END SELECT
SELECT CASE (SILLY)
CASE DEFAULT
  CALL THAT
CASE (1)
  CALL THIS
END SELECT
```

**NOTE 11.24**

A code fragment showing several selections of one block:

```
SELECT CASE (N)
```

## NOTE 11.24 (cont.)

```

CASE (1, 3:5, 8) ! Selects 1, 3, 4, 5, 8
  CALL SUB
CASE DEFAULT
  CALL OTHER
END SELECT

```

## 11.1.10 SELECT RANK construct

## 11.1.10.1 Purpose and form of the SELECT RANK construct

1 The SELECT RANK construct selects for execution at most one of its constituent blocks. The selection is based on the rank of an [assumed-rank variable](#). A name is associated with the variable (19.4, 19.5.1.6), in the same way as for the [ASSOCIATE construct](#).

R1148 *select-rank-construct* is *select-rank-stmt*  
 [ *select-rank-case-stmt*  
*block* ]...  
*end-select-rank-stmt*

R1149 *select-rank-stmt* is [ *select-construct-name* : ] SELECT RANK ■  
 ■ ( [ *associate-name* => ] *selector* )

C1149 The *selector* in a *select-rank-stmt* shall be the name of an [assumed-rank](#) array.

R1150 *select-rank-case-stmt* is RANK ( *scalar-int-constant-expr* ) [ *select-construct-name* ]  
 or RANK ( \* ) [ *select-construct-name* ]  
 or RANK DEFAULT [ *select-construct-name* ]

C1150 A *scalar-int-constant-expr* in a *select-rank-case-stmt* shall be nonnegative and less than or equal to the maximum possible rank of *selector*.

C1151 For a given *select-rank-construct*, the same rank value shall not be specified in more than one *select-rank-case-stmt*.

C1152 For a given *select-rank-construct*, there shall be at most one RANK ( \* ) *select-rank-case-stmt* and at most one RANK DEFAULT *select-rank-case-stmt*.

C1153 If *select-construct-name* appears on a *select-rank-case-stmt* the corresponding *select-rank-stmt* shall specify the same *select-construct-name*.

C1154 A SELECT RANK construct shall not have a *select-rank-case-stmt* that is RANK ( \* ) if the selector has the [ALLOCATABLE](#) or [POINTER](#) attribute.

R1151 *end-select-rank-stmt* is END SELECT [ *select-construct-name* ]

C1155 If the *select-rank-stmt* of a *select-rank-construct* specifies a *select-construct-name*, the corresponding *end-select-rank-stmt* shall specify the same *select-construct-name*. If the *select-rank-stmt* of a *select-rank-construct* does not specify a *select-construct-name*, the corresponding *end-select-rank-stmt* shall not specify a *select-construct-name*.

2 The associate name of a SELECT RANK construct is the *associate-name* if specified; otherwise it is the name that constitutes the selector.

## 11.1.10.2 Execution of the SELECT RANK construct

1 A SELECT RANK construct selects at most one block to be executed. During execution of that block, the associate name identifies an entity which is associated (19.5.1.6) with the selector. A RANK ( \* ) statement

1 matches the selector if the selector is argument associated with an [assumed-size array](#). A RANK ( *scalar-int-constant-expr* ) statement matches the selector if the selector has that rank and is not argument associated with  
 2 an [assumed-size array](#). A RANK DEFAULT statement matches the selector if no other [select-rank-case-stmt](#)  
 3 of the construct matches the selector. If a [select-rank-case-stmt](#) matches the selector, the block following that  
 4 statement is executed; otherwise, control is transferred to the [end-select-rank-stmt](#).  
 5

6 2 It is permissible to branch to an [end-select-rank-stmt](#) only from within its SELECT RANK construct.

### 7 11.1.10.3 Attributes of a SELECT RANK associate name

8 1 Within the block following a RANK DEFAULT statement, the associating entity ([19.5.5](#)) is [assumed-rank](#) and has  
 9 exactly the same attributes as the selector. Within the block following a RANK ( \* ) statement, the associating  
 10 entity has rank 1 and is [assumed-size](#), as if it were declared with DIMENSION(1:\*). Within the block following  
 11 a RANK ( *scalar-int-constant-expr* ) statement, the associating entity has the specified rank; the lower bound of  
 12 each dimension is the result of the intrinsic function [LBOUND](#) ([16.9.109](#)) applied to the corresponding dimension  
 13 of the selector, and the upper bound of each dimension is the result of the intrinsic function [UBOUND](#) ([16.9.196](#))  
 14 applied to the corresponding dimension of the selector.

15 2 The associating entity has the [ALLOCATABLE](#), [POINTER](#), or [TARGET](#) attribute if the selector has that  
 16 attribute. The other attributes of the associating entity are described in [11.1.3.3](#).

### 17 11.1.10.4 Examples of the SELECT RANK construct

18 1 This example shows how to use a SELECT RANK construct to process scalars and rank-2 arrays; anything else  
 19 will be rejected as an error.

```
20      SUBROUTINE process(x)
21          REAL x(..)
22          !
23          SELECT RANK(x)
24          RANK (0)
25             x = 0
26          RANK (2)
27             IF (SIZE(x,2)>=2) x(:,2) = 2
28          RANK DEFAULT
29             Print *, 'I did not expect rank', RANK(x), 'shape', SHAPE(x)
30             ERROR STOP 'process bad arg'
31          END SELECT
```

32 2 The following example shows how to process [assumed-size arrays](#), including how to use [sequence association](#) if  
 33 you want to do multi-dimensional processing of an [assumed-size array](#).

```
34      SELECT RANK (y => x)
35      RANK (*)
36      IF (RANK(x)==2) THEN
37          ! Special code for the rank two case.
38          CALL sequence_assoc_2(y, LBOUND(x,1), UBOUND(x,1), LBOUND(x,2))
39      ELSE
40          ! We just do all the other ranks in array element order.
41          i = 1
42          DO
43             IF (y(i)==0) Exit
44             y(i) = -y(i)
45             i = i + 1
46          END DO
47      END IF
```

```

1      END SELECT
2      ...
3      CONTAINS
4      ...
5      SUBROUTINE sequence_assoc_2(a, lb1, ub1, lb2)
6          INTEGER, INTENT (IN) :: lb1, ub1, lb2
7          REAL a(lb1:ub1,lb2:*)
8          j = lb2
9      outer: DO
10         DO i=lb1,ub1
11             IF (a(i,j)==0) EXIT outer
12             a(i,j) = a(i,j)**2
13         END DO
14         j = j + 1
15         IF (ANY(a(:,j)==0)) EXIT
16         j = j + 1
17     END DO outer
18 END SUBROUTINE

```

### 19 11.1.11 SELECT TYPE construct

#### 20 11.1.11.1 Purpose and form of the SELECT TYPE construct

21 1 The SELECT TYPE construct selects for execution at most one of its constituent blocks. The selection is based  
 22 on the [dynamic type](#) of an expression. A [name](#) is associated with the expression or variable ([19.4](#), [19.5.1.6](#)), in  
 23 the same way as for the ASSOCIATE construct.

24 R1152 *select-type-construct*            **is** *select-type-stmt*  
 25    [ *type-guard-stmt*  
 26    *block* ] ...  
 27    *end-select-type-stmt*

28 R1153 *select-type-stmt*                **is** [ *select-construct-name* : ] SELECT TYPE ■  
 29    ■ ( [ *associate-name* => ] *selector* )

30 C1156 (R1153) If *selector* is not a named *variable*, *associate-name* => shall appear.

31 C1157 (R1153) If *selector* is not a *variable* or is a *variable* that has a [vector subscript](#), neither *associate-name*  
 32 nor any subobject thereof shall appear in a variable definition context ([19.6.7](#)).

33 C1158 (R1153) The *selector* in a *select-type-stmt* shall be [polymorphic](#).

34 R1154 *type-guard-stmt*                **is** TYPE IS ( *type-spec* ) [ *select-construct-name* ]  
 35    **or** CLASS IS ( *derived-type-spec* ) [ *select-construct-name* ]  
 36    **or** CLASS DEFAULT [ *select-construct-name* ]

37 C1159 (R1154) The *type-spec* or *derived-type-spec* shall specify that each length type parameter is assumed.

38 C1160 (R1154) The *type-spec* or *derived-type-spec* shall not specify a type with the [BIND attribute](#) or the  
 39 [SEQUENCE attribute](#).

40 C1161 (R1152) If *selector* is not [unlimited polymorphic](#), each TYPE IS or CLASS IS *type-guard-stmt* shall  
 41 specify an [extension](#) of the [declared type](#) of *selector*.

42 C1162 (R1152) For a given *select-type-construct*, the same type and [kind type parameter](#) values shall not be  
 43 specified in more than one TYPE IS *type-guard-stmt* and shall not be specified in more than one CLASS  
 44 IS *type-guard-stmt*.

1 C1163 (R1152) For a given *select-type-construct*, there shall be at most one CLASS DEFAULT *type-guard-stmt*.

2 R1155 *end-select-type-stmt* is END SELECT [ *select-construct-name* ]

3 C1164 (R1152) If the *select-type-stmt* of a *select-type-construct* specifies a *select-construct-name*, the correspond-  
 4 ing *end-select-type-stmt* shall specify the same *select-construct-name*. If the *select-type-stmt* of a *select-*  
 5 *type-construct* does not specify a *select-construct-name*, the corresponding *end-select-type-stmt* shall not  
 6 specify a *select-construct-name*. If a *type-guard-stmt* specifies a *select-construct-name*, the corresponding  
 7 *select-type-stmt* shall specify the same *select-construct-name*.

8 2 The *associate name* of a SELECT TYPE construct is the *associate-name* if specified; otherwise it is the *name*  
 9 that constitutes the *selector*.

#### 10 11.1.11.2 Execution of the SELECT TYPE construct

11 1 Execution of a SELECT TYPE construct causes evaluation of every expression within a selector that is a variable  
 12 designator, or evaluation of a selector that is not a variable designator.

13 2 A SELECT TYPE construct selects at most one block to be executed. During execution of that block, the  
 14 *associate name* identifies an entity which is associated (19.5.1.6) with the selector.

15 3 A TYPE IS type guard statement matches the selector if the *dynamic type* and *kind type parameter* values of  
 16 the selector are the same as those specified by the statement. A CLASS IS type guard statement matches the  
 17 selector if the *dynamic type* of the selector is an *extension* of the type specified by the statement and the *kind*  
 18 *type parameter* values specified by the statement are the same as the corresponding type parameter values of the  
 19 *dynamic type* of the selector.

20 4 The block to be executed is selected as follows.

21 (1) If a TYPE IS type guard statement matches the selector, the block following that statement is  
 22 executed.

23 (2) Otherwise, if exactly one CLASS IS type guard statement matches the selector, the block following  
 24 that statement is executed.

25 (3) Otherwise, if several CLASS IS type guard statements match the selector, one of these statements  
 26 will inevitably specify a type that is an *extension* of all the types specified in the others; the block  
 27 following that statement is executed.

28 (4) Otherwise, if there is a CLASS DEFAULT type guard statement, the block following that statement  
 29 is executed.

30 (5) Otherwise, no block is executed.

#### NOTE 11.25

This algorithm does not examine the type guard statements in source text order when it looks for a match; it selects the most particular type guard when there are several potential matches.

31 5 Within the block following a TYPE IS type guard statement, the *associating entity* (19.5.5) is not polymorphic  
 32 (7.3.2.3), has the type named in the type guard statement, and has the type parameter values of the selector.

33 6 Within the block following a CLASS IS type guard statement, the *associating entity* is *polymorphic* and has the  
 34 *declared type* named in the type guard statement. The type parameter values of the *associating entity* are the  
 35 corresponding type parameter values of the selector.

36 7 Within the block following a CLASS DEFAULT type guard statement, the *associating entity* is *polymorphic* and  
 37 has the same *declared type* as the selector. The type parameter values of the *associating entity* are those of the  
 38 *declared type* of the selector.



**NOTE 11.26**

If the *declared type* of the *selector* is T, specifying CLASS DEFAULT has the same effect as specifying CLASS IS (T).

- 1 8 The other attributes of the *associating entity* are described in 11.1.3.3.  
 2 9 It is permissible to branch to an *end-select-type-stmt* only from within its SELECT TYPE construct.

3 **11.1.11.3 Examples of the SELECT TYPE construct**

**NOTE 11.27**

```

TYPE POINT
  REAL :: X, Y
END TYPE POINT
TYPE, EXTENDS(POINT) :: POINT_3D
  REAL :: Z
END TYPE POINT_3D
TYPE, EXTENDS(POINT) :: COLOR_POINT
  INTEGER :: COLOR
END TYPE COLOR_POINT

TYPE(POINT), TARGET :: P
TYPE(POINT_3D), TARGET :: P3
TYPE(COLOR_POINT), TARGET :: C
CLASS(POINT), POINTER :: P_OR_C
P_OR_C => C
SELECT TYPE ( A => P_OR_C )
CLASS IS ( POINT )
  ! "CLASS ( POINT ) :: A" implied here
  PRINT *, A%X, A%Y ! This block gets executed
TYPE IS ( POINT_3D )
  ! "TYPE ( POINT_3D ) :: A" implied here
  PRINT *, A%X, A%Y, A%Z
END SELECT

```

**NOTE 11.28**

The following example illustrates the omission of *associate-name*. It uses the declarations from Note 11.27.

```

P_OR_C => P3
SELECT TYPE ( P_OR_C )
CLASS IS ( POINT )
  ! "CLASS ( POINT ) :: P_OR_C" implied here
  PRINT *, P_OR_C%X, P_OR_C%Y
TYPE IS ( POINT_3D )
  ! "TYPE ( POINT_3D ) :: P_OR_C" implied here
  PRINT *, P_OR_C%X, P_OR_C%Y, P_OR_C%Z ! This block gets executed
END SELECT

```

4 **11.1.12 EXIT statement**

- 5 1 The EXIT statement provides one way of terminating a loop, or completing execution of another construct.

6 R1156 *exit-stmt* is EXIT [ *construct-name* ]

1 C1165 If a *construct-name* appears on an EXIT statement, the EXIT statement shall be within that construct;  
2 otherwise, it shall be within at least one *do-construct*.

3 2 An EXIT statement belongs to a particular construct. If a construct name appears, the EXIT statement belongs  
4 to that construct; otherwise, it belongs to the innermost **DO construct** in which it appears.

5 C1166 An *exit-stmt* shall not appear within a **CHANGE TEAM**, **CRITICAL**, or **DO CONCURRENT** construct  
6 if it belongs to that construct or an outer construct.

7 3 When an EXIT statement that belongs to a **DO construct** is executed, it terminates the loop (11.1.7.4.5) and  
8 any active loops contained within the terminated loop. When an EXIT statement that belongs to a non-DO  
9 construct is executed, it terminates any active loops contained within that construct, and completes execution of  
10 that construct.

## 11 11.2 Branching

### 12 11.2.1 Branch concepts

13 1 Branching is used to alter the normal execution sequence. A branch causes a transfer of control from one statement  
14 to a labeled **branch target statement** in the same **inclusive scope**. Branching may be caused by a **GO TO state-**  
15 **ment**, a **computed GO TO statement**, a **CALL statement** that has an *alt-return-spec*, or an input/output statement that has  
16 an **END=**, **EOR=**, or **ERR=** specifier. Although procedure references and control constructs can cause transfer  
17 of control, they are not branches. A **branch target statement** is an *action-stmt*, *associate-stmt*, *end-associate-*  
18 *stmt*, *if-then-stmt*, *end-if-stmt*, *select-case-stmt*, *end-select-stmt*, *select-rank-stmt*, *end-select-rank-stmt*, *select-*  
19 *type-stmt*, *end-select-type-stmt*, *do-stmt*, *end-do-stmt*, *block-stmt*, *end-block-stmt*, *critical-stmt*, *end-critical-stmt*,  
20 *forall-construct-stmt*, *forall-stmt*, *where-construct-stmt*, *end-function-stmt*, *end-mp-subprogram-stmt*, *end-program-*  
21 *stmt*, or *end-subroutine-stmt*.

### 22 11.2.2 GO TO statement

23 R1157 *goto-stmt* is GO TO *label*

24 C1167 (R1157) The *label* shall be the statement label of a **branch target statement** that appears in the same  
25 **inclusive scope** as the *goto-stmt*.

26 1 Execution of a GO TO statement causes a branch to the **branch target statement** identified by the label.

### 27 11.2.3 Computed GO TO statement

28 R1158 *computed-goto-stmt* is GO TO ( *label-list* ) [ , ] *scalar-int-expr*

29 C1168 (R1158) Each *label* in *label-list* shall be the statement label of a **branch target statement** that appears in the same **inclusive**  
30 **scope** as the *computed-goto-stmt*.

31 1 Execution of a computed GO TO statement causes evaluation of the scalar integer expression. If this value is  $i$  such that  $1 \leq i \leq n$   
32 where  $n$  is the number of labels in *label-list*, a branch occurs to the **branch target statement** identified by the  $i^{th}$  label in the list of  
33 labels. If  $i$  is less than 1 or greater than  $n$ , the execution sequence continues as though a **CONTINUE statement** were executed.

## 34 11.3 CONTINUE statement

35 1 Execution of a CONTINUE statement has no effect.

36 R1159 *continue-stmt* is CONTINUE

## 11.4 STOP and ERROR STOP statements

R1160 *stop-stmt* is STOP [ *stop-code* ] [ , QUIET = *scalar-logical-expr* ]

R1161 *error-stop-stmt* is ERROR STOP [ *stop-code* ] [ , QUIET = *scalar-logical-expr* ]

R1162 *stop-code* is *scalar-default-char-expr*  
or *scalar-int-expr*

C1169 (R1162) The *scalar-int-expr* shall be of default kind.

1 Execution of a STOP statement initiates **normal termination** of execution. Execution of an ERROR STOP statement initiates **error termination** of execution.

2 When an **image** is terminated by a STOP or ERROR STOP statement, its stop code, if any, is made available in a processor-dependent manner. If the *stop-code* is an integer, it is recommended that the value be used as the process exit status, if the processor supports that concept. If the *stop-code* in a STOP statement is of type character or does not appear, or if an *end-program-stmt* is executed, it is recommended that the value zero be supplied as the process exit status, if the processor supports that concept. If the *stop-code* in an ERROR STOP statement is of type character or does not appear, it is recommended that a processor-dependent nonzero value be supplied as the process exit status, if the processor supports that concept.

3 If QUIET= is omitted or the *scalar-logical-expr* has the value false:

- if any exception (17) is signaling on that **image**, the processor shall issue a warning indicating which exceptions are signaling, and this warning shall be on the **unit** identified by the named constant **ERROR\_UNIT** from the intrinsic module **ISO\_FORTRAN\_ENV** (16.10.2.9);
- if a stop code is specified, it is recommended that it be made available by formatted output to the same **unit**.

4 If QUIET= appears and the *scalar-logical-expr* has the value true, no output of signaling exceptions or stop code shall be produced.

### NOTE 11.29

When **normal termination** occurs on more than one **image**, it is expected that a processor-dependent summary of any stop codes and signaling exceptions will be made available.

### NOTE 11.30

If the integer *stop-code* is used as the process exit status, the processor might be able to interpret only values within a limited range, or only a limited portion of the integer value (for example, only the least-significant 8 bits).

## 11.5 FAIL IMAGE statement

R1163 *fail-image-stmt* is FAIL IMAGE

1 Execution of a FAIL IMAGE statement causes the executing image to cease participating in program execution without initiating termination. No further statements are executed by that image.

### NOTE 11.31

The FAIL IMAGE statement allows a program to test a recovery algorithm without needing an actual failure.

On a processor that does not have the ability to detect that an image has failed, execution of a FAIL IMAGE statement might provide a simulated failure environment that provides debug information.

In a piece of code that executes about once a second, invoking this subroutine on an image

**NOTE 11.31 (cont.)**

```

SUBROUTINE FAIL
  REAL :: X
  CALL RANDOM_NUMBER (X)
  IF (X<0.001) FAIL IMAGE
END SUBROUTINE FAIL

```

will cause that image to have an independent 1/1000 chance of failure every second if the random number generators on different images are independent.

Note that FAIL IMAGE is not an [image control statement](#).

## 11.6 Image execution control

### 11.6.1 Image control statements

1 The execution sequence on each [image](#) is specified in [5.3.5](#).

2 Execution of an [image control statement](#) divides the execution sequence on an [image](#) into segments. Each of the following is an [image control statement](#):

- [SYNC ALL](#) statement;
- [SYNC IMAGES](#) statement;
- [SYNC MEMORY](#) statement;
- [SYNC TEAM](#) statement;
- [ALLOCATE](#) or [DEALLOCATE](#) statement that has a [coarray allocate-object](#);
- [CHANGE TEAM](#) or [END TEAM](#) statement ([11.1.5](#));
- [CRITICAL](#) or [END CRITICAL](#) statement ([11.1.6](#));
- [EVENT POST](#) or [EVENT WAIT](#) statement;
- [FORM TEAM](#) statement;
- [LOCK](#) or [UNLOCK](#) statement;
- any statement that completes execution of a block or procedure and which results in the implicit deallocation of a [coarray](#);
- a [CALL](#) statement that references the intrinsic subroutine [MOVE\\_ALLOC](#) with [coarray](#) arguments;
- [STOP](#) statement;
- [END](#) statement of a main program.

3 All image control statements except [CRITICAL](#), [END CRITICAL](#), [EVENT POST](#), [EVENT WAIT](#), [FORM TEAM](#), [LOCK](#), and [UNLOCK](#) include the effect of executing a [SYNC MEMORY](#) statement ([11.6.5](#)).

4 During an execution of a statement that invokes more than one procedure, at most one invocation shall cause execution of an [image control statement](#) other than [CRITICAL](#) or [END CRITICAL](#).

### 11.6.2 Segments

1 On each [image](#), the sequence of statements executed before the first execution of an [image control statement](#), between the execution of two [image control statements](#), or after the last execution of an [image control statement](#) is a segment. The segment executed immediately before the execution of an [image control statement](#) includes the evaluation of all expressions within the statement.

2 By execution of [image control statements](#) or user-defined ordering ([11.6.5](#)), the program can ensure that the execution of the  $i^{th}$  segment on [image](#) P,  $P_i$ , either precedes or succeeds the execution of the  $j^{th}$  segment on another [image](#) Q,  $Q_j$ . If the program does not ensure this, segments  $P_i$  and  $Q_j$  are unordered; depending on the

1 relative execution speeds of the **images**, some or all of the execution of the segment  $P_i$  may take place at the same  
2 time as some or all of the execution of the segment  $Q_j$ .

3 3 A **coarray** may be referenced or defined by execution of an **atomic subroutine** during the execution of a segment  
4 that is unordered relative to the execution of a segment in which the **coarray** is referenced or defined by execution  
5 of an **atomic subroutine**. An **event variable** may be referenced or defined during the execution of a segment that  
6 is unordered relative to the execution of another segment in which that **event variable** is defined. Otherwise,

- 7 • if a variable is defined or becomes undefined on an **image** in a segment, it shall not be referenced, defined,  
8 or become undefined in a segment on another **image** unless the segments are ordered,
- 9 • if the allocation of an **allocatable** subobject of a **coarray** or the **pointer association** of a pointer subobject  
10 of a **coarray** is changed on an **image** in a segment, that subobject shall not be referenced, defined, or have  
11 its allocation or **association** status, **dynamic type**, array **bounds**, **shape**, or a **deferred type parameter** value  
12 inquired about in a segment on another **image** unless the segments are ordered, and
- 13 • if a procedure invocation on **image** P is in execution in segments  $P_i, P_{i+1}, \dots, P_k$  and defines a noncoarray  
14 **dummy argument**, the **effective argument** shall not be referenced, defined, or become undefined on another  
15 **image** Q in a segment  $Q_j$  unless  $Q_j$  precedes  $P_i$  or succeeds  $P_k$ .

#### NOTE 11.32

The set of all segments on all **images** is partially ordered: the segment  $P_i$  precedes segment  $Q_j$  if and only if there is a sequence of segments starting with  $P_i$  and ending with  $Q_j$  such that each segment of the sequence precedes the next either because they are consecutive segments on the same **image** or because of the execution of **image control statements**.

#### NOTE 11.33

If the segments  $S_1, S_2, \dots, S_k$  on the distinct **images**  $P_1, P_2, \dots, P_k$  are all unordered with respect to each other, it is expected that the processor will ensure that each of these **images** is provided with an equitable share of resources for executing its segment.

#### NOTE 11.34

Because of the restrictions on references and definitions in unordered segments, the processor can apply code motion optimizations within a segment as if it were the only **image** in execution, provided calls to **atomic subroutines** are not involved.

#### NOTE 11.35

The model upon which the interpretation of a program is based is that there is a permanent memory location for each **coarray** and that all **images** on which it is established can access it.

In practice, apart from executions of **atomic subroutines**, the processor could make a copy of a nonvolatile **coarray** on an **image** (in cache or a register, for example) and, as an optimization, defer copying a changed value back to the permanent memory location while it is still being used. Since the variable is not volatile, it is safe to defer this transfer until the end of the segment and thereafter to reload from permanent memory any **coarray** that was not defined within the segment. It might not be safe to defer these actions beyond the end of the segment since another **image** might reference the variable then.

The value of the ATOM argument of an **atomic subroutine** might be accessed or modified by another concurrently executing **image**. Therefore, execution of an **atomic subroutine** that references the ATOM argument cannot rely on a local copy, but instead always gets its value from its permanent memory location. Execution of an **atomic subroutine** that defines the ATOM argument does not complete until the value of its ATOM argument has been sent to its permanent memory location.

**NOTE 11.36**

The incorrect sequencing of [image control statements](#) can suspend execution indefinitely. For example, one [image](#) might be executing a [SYNC ALL statement](#) while another is executing an [ALLOCATE statement](#) for a [coarray](#).

**11.6.3 SYNC ALL statement**

R1164 *sync-all-stmt* is SYNC ALL [ ( [ *sync-stat-list* ] ) ]

R1165 *sync-stat* is STAT = *stat-variable*  
or ERRMSG = *errmsg-variable*

C1170 No specifier shall appear more than once in a given *sync-stat-list*.

C1171 A *stat-variable* or *errmsg-variable* in a *sync-stat* shall not be a [coindexed object](#).

1 The [STAT=](#) and [ERRMSG=](#) specifiers for [image control statements](#) are described in [11.6.11](#).

2 Successful execution of a SYNC ALL statement performs a synchronization of all [images](#) in the [current team](#). Execution on an [image](#), M, of the segment following the SYNC ALL statement is delayed until each other [image](#) in the [current team](#) has executed a SYNC ALL statement as many times as has [image M](#) in this [team](#). The segments that executed before the SYNC ALL statement on an [image](#) precede the segments that execute after the SYNC ALL statement on another [image](#).

**NOTE 11.37**

The processor might have special hardware or employ an optimized algorithm to make the SYNC ALL statement execute efficiently.

Here is a simple example of its use. [Image 1](#) reads data and broadcasts it to other [images](#):

```
REAL :: P[*]
...
SYNC ALL
IF (THIS_IMAGE()==1) THEN
  READ (*,*) P
  DO I = 2, NUM_IMAGES()
    P[I] = P
  END DO
END IF
SYNC ALL
```

**11.6.4 SYNC IMAGES statement**

R1166 *sync-images-stmt* is SYNC IMAGES ( *image-set* [ , *sync-stat-list* ] )

R1167 *image-set* is *int-expr*  
or \*

C1172 An *image-set* that is an *int-expr* shall be scalar or of [rank](#) one.

C1173 The value of *image-set* shall not depend on the value of *stat-variable* or *errmsg-variable*.

1 If *image-set* is an array expression, the value of each element shall be positive and not greater than the number of [images](#) in the [current team](#), and there shall be no repeated values.

2 If *image-set* is a scalar expression, its value shall be positive and not greater than the number of [images](#) in the [current team](#).

- 1 3 An *image-set* that is an asterisk specifies all *images* in the *current team*.
- 2 4 Execution of a SYNC IMAGES statement performs a synchronization of the *image* with each of the other *images*  
 3 in the *image-set*. Executions of SYNC IMAGES statements on *images* M and T correspond if the number of  
 4 times *image* M has executed a SYNC IMAGES statement in the *current team* with T in its *image* set is the same  
 5 as the number of times *image* T has executed a SYNC IMAGES statement with M in its *image* set in this *team*.  
 6 The segments that executed before the SYNC IMAGES statement on either *image* precede the segments that  
 7 execute after the corresponding SYNC IMAGES statement on the other *image*.

**NOTE 11.38**

A SYNC IMAGES statement that specifies the single *image index* value `THIS_IMAGE ( )` in its *image* set is allowed. This simplifies writing programs for an arbitrary number of *images* by allowing correct execution in the limiting case of the number of *images* being equal to one.

**NOTE 11.39**

In a program that uses `SYNC ALL` as its only synchronization mechanism, every `SYNC ALL` statement could be replaced by a `SYNC IMAGES (*)` statement, but `SYNC ALL` might give better performance.

`SYNC IMAGES` statements are not required to specify the entire *image* set, or even the same *image* set, on all *images* participating in the synchronization. In the following example, *image* 1 will wait for each of the other *images* to execute the statement `SYNC IMAGES (1)`. The other *images* wait for *image* 1 to set up the data, but do not wait on any other *image*.

```
IF (THIS_IMAGE() == 1) then
  ! Set up coarray data needed by all other images.
  SYNC IMAGES(*)
ELSE
  SYNC IMAGES(1)
  ! Use the data set up by image 1.
END IF
```

When the following example runs on five or more *images*, each *image* synchronizes with both of its neighbors, in a circular fashion.

```
INTEGER :: up, down
...
IF (NUM_IMAGES () > 1) THEN
  up = THIS_IMAGE () + 1; IF (up>NUM_IMAGES ()) up = 1
  down = THIS_IMAGE () - 1; IF (down==0) down = NUM_IMAGES ()
  SYNC IMAGES ( (/ up, down /) )
END IF
```

This might appear to have the same effect as `SYNC ALL` but there is no ordering between the preceding and succeeding segments on non-adjacent *images*. For example, the segment preceding the `SYNC IMAGES` statement on *image* 3 will be ordered before those succeeding it on *images* 2 and 4, but not those on *images* 1 and 5.

**NOTE 11.40**

In the following example, each *image* synchronizes with its neighbor.

```
INTEGER :: ME, NE, STEP, NSTEPS
NE = NUM_IMAGES()
ME = THIS_IMAGE()
  ! Initial calculation
SYNC ALL
```

**NOTE 11.40 (cont.)**

```

DO STEP = 1, NSTEPS
  IF (ME > 1) SYNC IMAGES(ME-1)
    ! Perform calculation
  IF (ME < NE) SYNC IMAGES(ME+1)
END DO
SYNC ALL

```

The calculation starts on [image 1](#) since all the others will be waiting on SYNC IMAGES (ME-1). When this is done, [image 2](#) can start and [image 1](#) can perform its second calculation. This continues until they are all executing different steps at the same time. Eventually, [image 1](#) will finish and then the others will finish one by one.

1 **11.6.5 SYNC MEMORY statement**

2 1 Execution of a SYNC MEMORY statement ends one segment and begins another; those two segments can be  
3 ordered by a user-defined way with respect to segments on other [images](#).

4 R1168 *sync-memory-stmt* is SYNC MEMORY [ ( [ *sync-stat-list* ] ) ]

5 2 If, by execution of statements on [image P](#),

- 6 • a variable *X* on [image Q](#) is defined, referenced, becomes undefined, or has its allocation status, [pointer](#)  
7 [association](#) status, array bounds, [dynamic type](#), or type parameters changed or inquired about by execution  
8 of a statement,
- 9 • that statement precedes a successful execution of a SYNC MEMORY statement, and
- 10 • a variable *Y* on [image Q](#) is defined, referenced, becomes undefined, or has its allocation status, [pointer](#)  
11 [association](#) status, array bounds, [dynamic type](#), or type parameters changed or inquired about by execution  
12 of a statement that succeeds execution of that SYNC MEMORY statement,

13 then the action regarding *X* on [image Q](#) precedes the action regarding *Y* on [image Q](#).

14 3 User-defined ordering of segment  $P_i$  on [image P](#) to precede segment  $Q_j$  on [image Q](#) occurs when

- 15 • [image P](#) executes an [image control statement](#) that ends segment  $P_i$ , and then executes statements that  
16 initiate a cooperative synchronization between [images P](#) and [Q](#), and
- 17 • [image Q](#) executes statements that complete the cooperative synchronization between [images P](#) and [Q](#) and  
18 then executes an [image control statement](#) that begins segment  $Q_j$ .

19 4 Execution of the cooperative synchronization between [images P](#) and [Q](#) shall include a dependency that forces  
20 execution on [image P](#) of the statements that initiate the synchronization to precede the execution on [image Q](#) of  
21 the statements that complete the synchronization. The mechanisms available for creating such a dependency are  
22 processor dependent.

**NOTE 11.41**

SYNC MEMORY usually suppresses compiler optimizations that might reorder memory operations across the segment boundary defined by the SYNC MEMORY statement and ensures that all memory operations initiated in the preceding segments in its [image](#) complete before any memory operations in the subsequent segment in its [image](#) are initiated. It needs to do this unless it can establish that failure to do so could not alter processing on another [image](#).

**NOTE 11.42**

SYNC MEMORY can be used to implement specialized schemes for segment ordering. For example, the user might have access to an [external procedure](#) that performs synchronization between [images](#). That library procedure might not be aware of the mechanisms used by the processor to manage remote data references



**NOTE 11.42 (cont.)**

and definitions, and therefore not, by itself, be able to ensure the correct memory state before and after its reference. The SYNC MEMORY statement provides the needed memory ordering that enables the safe use of the external synchronization routine. For example:

```

INTEGER :: IAM
REAL    :: X[*]

IAM = THIS_IMAGE ()
IF (IAM == 1) X = 1.0
SYNC MEMORY
CALL EXTERNAL_SYNC ()
SYNC MEMORY
IF (IAM == 2) WRITE (*,*) X[1]

```

where executing the subroutine EXTERNAL\_SYNC has an *image* synchronization effect similar to executing a SYNC ALL statement.

**11.6.6 SYNC TEAM statement**

R1169 *sync-team-stmt* is SYNC TEAM ( *team-value* [ , *sync-stat-list* ] )

1 The *team-value* shall identify an *ancestor team*, the *current team*, or a *team* whose *parent* is the *current team*.  
 4 The executing *image* shall be a member of the specified *team*.

5 2 Successful execution of a SYNC TEAM statement performs a synchronization of the team identified by *team-value*.  
 6 Execution on an image, M, of the segment following the SYNC TEAM statement is delayed until each  
 7 other image of the specified team has executed a SYNC TEAM statement specifying the same team as many  
 8 times as has image M in this team. The segments that executed before the SYNC TEAM statement on an image  
 9 precede the segments that execute after the corresponding SYNC TEAM statement on another image.

**NOTE 11.43**

A SYNC TEAM statement synchronizes a particular team whereas a SYNC ALL statement synchronizes the current team.

**11.6.7 EVENT POST statement**

11 1 The EVENT POST statement posts an event.

12 R1170 *event-post-stmt* is EVENT POST ( *event-variable* [ , *sync-stat-list* ] )

13 R1171 *event-variable* is *scalar-variable*

14 C1174 (R1171) An *event-variable* shall be of type EVENT\_TYPE from the intrinsic module ISO\_FORTRAN\_  
 15 ENV (16.10.2).

16 2 The *event-variable* shall not depend on the value of *stat-variable* or *errmsg-variable*.

17 3 Successful execution of an EVENT POST statement atomically increments the count of the *event variable* by  
 18 one. If an error condition occurs during execution of an EVENT POST statement, the value of the count of the  
 19 *event variable* is processor dependent. The completion of an EVENT POST statement does not depend on the  
 20 execution of a corresponding EVENT WAIT statement.

**11.6.8 EVENT WAIT statement**

22 1 The EVENT WAIT statement waits until an event is posted.

1 R1172 *event-wait-stmt* is EVENT WAIT ( *event-variable* [ , *event-wait-spec-list* ] )

2 R1173 *event-wait-spec* is *until-spec*  
3 or *sync-stat*

4 R1174 *until-spec* is UNTIL\_COUNT = *scalar-int-expr*

5 C1175 (R1172) The *event-variable* in an *event-wait-stmt* shall not be coindexed.

6 C1176 No specifier shall appear more than once in a given *event-wait-spec-list*.

7 2 The *event-variable* shall not depend on the value of *stat-variable* or *errmsg-variable*.

8 3 Execution of an EVENT WAIT statement consists of the following sequence of actions:

9 1. if the UNTIL\_COUNT= specifier does not appear, the threshold value is set to one; otherwise, the threshold  
10 value is set to the maximum of the value of the *scalar-int-expr* and one;

11 2. the executing image waits until the count of the *event variable* is greater than or equal to the threshold  
12 value or an error condition occurs;

13 3. if no error condition occurs, the count of the *event variable* is atomically decremented by the threshold  
14 value.

15 4 If an error occurs during execution of an EVENT WAIT statement, the value of the count of its *event variable*  
16 is processor dependent.

17 5 An **EVENT POST statement** execution is initially unsatisfied. Successful execution of an EVENT WAIT state-  
18 ment with a threshold of *k* satisfies the first *k* unsatisfied **EVENT POST statement** executions for that *event*  
19 *variable*. This EVENT WAIT statement execution causes the segment following the EVENT WAIT statement  
20 execution to succeed the segments preceding those *k* **EVENT POST statement** executions.

### 21 11.6.9 FORM TEAM statement

22 1 The FORM TEAM statement creates new teams whose *parent team* is the *current team*.

23 R1175 *form-team-stmt* is FORM TEAM ( *team-number*, *team-variable* ■  
24 ■ [ , *form-team-spec-list* ] )

25 R1176 *team-number* is *scalar-int-expr*

26 R1177 *team-variable* is *scalar-variable*

27 C1177 A *team-variable* shall be of type **TEAM\_TYPE** from the intrinsic module **ISO\_FORTRAN\_ENV**.

28 R1178 *form-team-spec* is NEW\_INDEX = *scalar-int-expr*  
29 or *sync-stat*

30 C1178 No specifier shall appear more than once in a given *form-team-spec-list*.

31 2 Successful execution of a FORM TEAM statement creates a new team for each unique *team-number* value specified  
32 by the *active images* of the *current team*. The value of *team-number* shall be positive. Each executing image will  
33 belong to the team whose *team number* is equal to the value of *team-number* on that image, and the *team-variable*  
34 becomes defined with a value that identifies that team.

35 3 The value of the *scalar-int-expr* in a NEW\_INDEX= specifier specifies the *image index* that the executing image  
36 will have in its new *team*. It shall be positive, less than or equal to the number of *images* in the *team*, and  
37 different from the value specified by every other *image* that belongs to that *team*.

- 1 4 If the NEW\_INDEX= specifier does not appear, the *image index* of the executing image in the new *team* is  
 2 processor dependent. This *image index* will be positive, less than or equal to the number of *images* in the *team*,  
 3 and different from that of every other *image* in the *team*.
- 4 5 If the FORM TEAM statement is executed on one *image*, the same statement shall be executed on all *active*  
 5 *images* of the *current team*. When a FORM TEAM statement is executed, there is an implicit synchronization  
 6 of all *active images* in the *current team*. On those *images*, execution of the segment following the statement is  
 7 delayed until all other *active images* in the *current team* have executed the same statement the same number of  
 8 times in this team. If an error condition other than detection of a failed image occurs, the team variable becomes  
 9 undefined.
- 10 6 If execution of a FORM TEAM statement assigns the value STAT\_FAILED\_IMAGE to the *stat-variable*, the  
 11 effect is the same as for the successful execution of FORM TEAM except for the value assigned to *stat-variable*.

**NOTE 11.44**

Executing the statement

```
FORM TEAM ( 2 - MOD (THIS_IMAGE (), 2), ODD_EVEN )
```

will create two subteams of the *current team*, with *images* whose image index is odd being in the team with number 1, and those with an even image index being in the team with number 2.

**NOTE 11.45**

If the current team consists of  $P^2$  images, with corresponding coarrays on each image representing parts of a larger array spread over a  $P \times P$  square, the following code will establish teams for the rows with image indices equal to the column indices.

```
USE, INTRINSIC :: ISO_FORTRAN_ENV
TYPE(Team_Type) :: Row
REAL :: A [P, *]
INTEGER :: ME (2)
ME (:) = THIS_IMAGE (A)
FORM TEAM (ME(1), Row, NEW_INDEX=ME(2))
```

## 12 11.6.10 LOCK and UNLOCK statements

13 R1179 *lock-stmt* is LOCK ( *lock-variable* [ , *lock-stat-list* ] )

14 R1180 *lock-stat* is ACQUIRED\_LOCK = *scalar-logical-variable*  
 15 or *sync-stat*

16 C1179 No specifier shall appear more than once in a given *lock-stat-list*.

17 R1181 *unlock-stmt* is UNLOCK ( *lock-variable* [ , *sync-stat-list* ] )

18 R1182 *lock-variable* is *scalar-variable*

19 C1180 (R1182) A *lock-variable* shall be of type LOCK\_TYPE from the intrinsic module ISO\_FORTRAN\_ENV  
 20 (16.10.2.19).

21 1 The *lock-variable* shall not depend on the value of *stat-variable*, *errmsg-variable*, or the *scalar-logical-variable* in  
 22 the ACQUIRED\_LOCK= specifier. The *scalar-logical-variable* shall not depend on the value of the *lock-variable*,  
 23 *stat-variable*, or *errmsg-variable*.

24 2 A lock variable is unlocked if and only if the value of each component is the same as its default value. If it has any  
 25 other value, it is locked. A lock variable is locked by an *image* if it was locked by execution of a LOCK statement  
 26 on that *image*, has not been subsequently unlocked by execution of an UNLOCK statement on the same *image*,

- 1 and that `image` has not failed.
- 2 3 Successful execution of a LOCK statement without an ACQUIRED\_LOCK= specifier causes the lock variable  
3 to become locked by that `image`. If the lock variable is already locked by another `image`, that LOCK statement  
4 causes the lock variable to become locked after the other `image` causes the lock variable to become unlocked.
- 5 4 If the lock variable is unlocked, successful execution of a LOCK statement with an ACQUIRED\_LOCK= specifier  
6 causes the lock variable to become locked by that `image` and the scalar logical variable to become defined with the  
7 value true. If the lock variable is already locked by a different `image`, successful execution of a LOCK statement  
8 with an ACQUIRED\_LOCK= specifier leaves the lock variable unchanged and causes the scalar logical variable  
9 to become defined with the value false.
- 10 5 Successful execution of an UNLOCK statement causes the lock variable to become unlocked. Failure of an `image`  
11 causes all lock variables that are locked by that `image` to become unlocked.
- 12 6 During execution of the program, the value of a lock variable changes through a sequence of locked and unlocked  
13 states due to the execution of LOCK and UNLOCK statements, and by `failure of an image` while it is locked by  
14 that `image`. If a lock variable becomes unlocked by execution of an UNLOCK statement on `image M` and next  
15 becomes locked by execution of a LOCK statement on `image T`, the segments preceding the UNLOCK statement  
16 on `image M` precede the segments following the LOCK statement on `image T`. Execution of a LOCK statement  
17 that does not cause the lock variable to become locked does not affect segment ordering.
- 18 7 An error condition occurs if the lock variable in a LOCK statement is already locked by the executing `image`.  
19 An error condition occurs if the lock variable in an UNLOCK statement is not already locked by the executing  
20 `image`. If an error condition occurs during execution of a LOCK or UNLOCK statement, the value of the lock  
21 variable is not changed and the value of the ACQUIRED\_LOCK variable, if any, is not changed.

**NOTE 11.46**

A lock variable is effectively defined atomically by a LOCK or UNLOCK statement. If LOCK statements on two `images` both attempt to acquire a lock, one will succeed and the other will either fail if an ACQUIRED\_LOCK= specifier appears, or will wait until the lock is later released if an ACQUIRED\_LOCK= specifier does not appear.

**NOTE 11.47**

An `image` might wait for a LOCK statement to successfully complete for a long period of time if other `images` frequently lock and unlock the same lock variable. This situation might result from executing LOCK statements with ACQUIRED\_LOCK= specifiers inside a spin loop.

**NOTE 11.48**

The following example illustrates the use of LOCK and UNLOCK statements to manage a work queue:

```
USE, INTRINSIC :: ISO_FORTRAN_ENV

TYPE(LOCK_TYPE) :: queue_lock[*] ! Lock on each image to manage its work queue
INTEGER :: work_queue_size[*]
TYPE(Task) :: work_queue(100)[*] ! List of tasks to perform

TYPE(Task) :: job ! Current task working on
INTEGER :: me

me = THIS_IMAGE()
DO
  ! Process the next item in your work queue

  LOCK (queue_lock) ! New segment A starts
```

## NOTE 11.48 (cont.)

```

! This segment A is ordered with respect to
! segment B executed by image me-1 below because of lock exclusion
IF (work_queue_size>0) THEN
  ! Fetch the next job from the queue
  job = work_queue(work_queue_size)
  work_queue_size = work_queue_size-1
END IF
UNLOCK (queue_lock) ! Segment ends
... ! Actually process the task

! Add a new task on neighbors queue:
LOCK(queue_lock[me+1]) ! Starts segment B
! This segment B is ordered with respect to
! segment A executed by image me+1 above because of lock exclusion
IF (work_queue_size[me+1]<SIZE (work_queue)) THEN
  work_queue_size[me+1] = work_queue_size[me+1]+1
  work_queue(work_queue_size[me+1])[me+1] = job
END IF
UNLOCK (queue_lock[me+1]) ! Ends segment B

END DO

```

1 **11.6.11 STAT= and ERRMSG= specifiers in image control statements**

- 2 1 In an *image control statement*, the *stat-variable* in a *sync-stat* shall not depend on the value of an *errmsg-variable*  
3 in a *sync-stat*, *event-variable*, *lock-variable*, *team-variable*, or the *scalar-logical-variable* in the ACQUIRED\_  
4 LOCK= specifier. The *errmsg-variable* in a *sync-stat* shall not depend on the value of a *stat-variable* in a  
5 *sync-stat*, *event-variable*, *lock-variable*, *team-variable*, or the *scalar-logical-variable* in the ACQUIRED\_LOCK=  
6 specifier.
- 7 2 If a STAT= specifier appears in a *sync-stat* in an *image control statement*, the *stat-variable* is assigned the value  
8 zero if execution of the statement is successful.
- 9 3 If the STAT= specifier appears in a *sync-stat* in an *EVENT WAIT* or *SYNC MEMORY* statement and an error  
10 condition occurs, *stat-variable* is assigned a processor-dependent positive value that is different from the value  
11 of *STAT\_FAILED\_IMAGE* (16.10.2.26) and *STAT\_STOPPED\_IMAGE* (16.10.2.29) from the intrinsic module  
12 *ISO\_FORTRAN\_ENV* (16.10.2).
- 13 4 The *images* involved in execution of an *END TEAM*, *FORM TEAM*, or *SYNC ALL* statement are those in the  
14 *current team*. The *images* involved in execution of a *CHANGE TEAM* or *SYNC TEAM* statement are those of  
15 the specified *team*. The *images* involved in execution of a *SYNC IMAGES* statement are the *images* specified  
16 and the executing *image*. The *images* involved in execution of an *EVENT POST* statement are the *image* on  
17 which the *event variable* is located and the executing *image*.
- 18 5 If the STAT= specifier appears in a *sync-stat* in a *CHANGE TEAM*, *END TEAM*, *EVENT POST*, *FORM*  
19 *TEAM*, *SYNC ALL*, *SYNC IMAGES*, or *SYNC TEAM* statement,
- 20 • if one of the *images* involved has stopped, *stat-variable* is assigned the value *STAT\_STOPPED\_IMAGE*  
21 (16.10.2.29) from the intrinsic module *ISO\_FORTRAN\_ENV*;
  - 22 • otherwise, if one of the *images* involved has failed and no other error condition occurs, the intended action  
23 is performed on the *active images* involved and *stat-variable* is assigned the value *STAT\_FAILED\_IMAGE*  
24 (16.10.2.26) from the intrinsic module *ISO\_FORTRAN\_ENV*;
  - 25 • otherwise, if any other error condition occurs, *stat-variable* is assigned a processor-dependent positive value  
26 that is different from the values of *STAT\_STOPPED\_IMAGE* and *STAT\_FAILED\_IMAGE*.

- 1 6 If the STAT= specifier appears in a *sync-stat* in a SYNC ALL, SYNC IMAGES, or SYNC TEAM statement  
2 and the error condition STAT\_STOPPED\_IMAGE occurs, the effect is the same as that of executing the SYNC  
3 MEMORY statement, except for defining the *stat-variable*.
- 4 7 If the STAT= specifier appears in a *sync-stat* in a LOCK statement,
- 5 • if the image on which the lock variable is located has failed, the *stat-variable* becomes defined with the  
6 value STAT\_FAILED\_IMAGE;
  - 7 • otherwise, if the lock variable is locked by the executing image, the *stat-variable* becomes defined with the  
8 value of STAT\_LOCKED (16.10.2.27) from the intrinsic module ISO\_FORTRAN\_ENV;
  - 9 • otherwise, if the lock variable is unlocked because of the failure of the image that locked it, *stat-variable*  
10 becomes defined with the value STAT\_UNLOCKED\_FAILED\_IMAGE (16.10.2.31) from the intrinsic  
11 module ISO\_FORTRAN\_ENV.
- 12 8 If the STAT= specifier appears in a *sync-stat* in an UNLOCK statement,
- 13 • if the image on which the lock variable is located has failed, the *stat-variable* becomes defined with the  
14 value STAT\_FAILED\_IMAGE;
  - 15 • otherwise, if the lock variable has the value unlocked, the *stat-variable* becomes defined with the value of  
16 STAT\_UNLOCKED (16.10.2.30) from the intrinsic module ISO\_FORTRAN\_ENV;
  - 17 • otherwise, if the lock variable is locked by a different image, the *stat-variable* becomes defined with the  
18 value STAT\_LOCKED\_OTHER\_IMAGE (16.10.2.28) from the intrinsic module ISO\_FORTRAN\_ENV.
- 19 9 If the STAT= specifier appears in a *sync-stat* in a LOCK or UNLOCK statement and any other error condi-  
20 tion occurs during execution of that statement, the *stat-variable* becomes defined with a processor-dependent  
21 positive integer value that is different from STAT\_LOCKED, STAT\_LOCKED\_OTHER\_IMAGE, STAT\_UN-  
22 LOCKED, and STAT\_UNLOCKED\_FAILED\_IMAGE.
- 23 10 If an image completes execution of a CRITICAL statement that has a *sync-stat* that is a STAT= specifier and the  
24 previous image to have entered the construct failed while executing it, the *stat-variable* becomes defined with the  
25 value STAT\_FAILED\_IMAGE and execution of the construct continues normally. If any other error condition  
26 occurs during execution of a CRITICAL statement that has a STAT= specifier, the *stat-variable* becomes defined  
27 with a processor-dependent value other than STAT\_FAILED\_IMAGE.
- 28 11 If an error condition occurs during execution of an image control statement that does not contain the STAT=  
29 specifier in a *sync-stat*, error termination is initiated.
- 30 12 If an ERRMSG= specifier appears in an image control statement and an error condition occurs, *errmsg-variable*  
31 is assigned an explanatory message, as if by intrinsic assignment. If no such condition occurs, the definition status  
32 and value of *errmsg-variable* are unchanged.
- 33 13 The set of error conditions that can occur in an image control statement is processor dependent.

**NOTE 11.49**

A processor might detect communication failure between images and treat it as an error condition. A processor might also treat an invalid set of images in a SYNC IMAGES statement as an error condition.

## 12 Input/output statements

### 12.1 Input/output concepts

1 **Input statements** provide the means of transferring data from external media to internal storage or from an **internal file** to internal storage. This process is called reading. **Output statements** provide the means of transferring data from internal storage to external media or from internal storage to an **internal file**. This process is called writing. Some **input/output statements** specify that editing of the data is to be performed.

2 In addition to the statements that transfer data, there are auxiliary input/output statements to manipulate the external medium, or to describe or inquire about the properties of the connection to the external medium.

3 The input/output statements are the **BACKSPACE**, **CLOSE**, **ENDFILE**, **FLUSH**, **INQUIRE**, **OPEN**, **PRINT**, **READ**, **REWIND**, **WAIT**, and **WRITE** statements.

4 A file is composed of either a sequence of **file storage units** (12.3.5) or a sequence of records, which provide an extra level of organization to the file. A file composed of records is called a **record file**. A file composed of **file storage units** is called a **stream file**. A processor may allow a file to be viewed both as a **record file** and as a **stream file**; in this case the relationship between the **file storage units** when viewed as a **stream file** and the records when viewed as a **record file** is processor dependent.

5 A file is either an **external file** (12.3) or an **internal file** (12.4).

### 12.2 Records

#### 12.2.1 Definition of a record

1 A **record** is a sequence of values or a sequence of characters. For example, a line on a terminal is usually considered to be a record. However, a record does not necessarily correspond to a physical entity. There are three kinds of records:

- (1) formatted;
- (2) unformatted;
- (3) endfile.

#### NOTE 12.1

What is called a “record” in Fortran is commonly called a “logical record”. There is no concept in Fortran of a “physical record.”

#### 12.2.2 Formatted record

1 A formatted record consists of a sequence of characters that are representable in the processor; however, a processor may prohibit some control characters (6.1.1) from appearing in a formatted record. The length of a formatted record is measured in characters and depends primarily on the number of characters put into the record when it is written; however, it may depend on the processor and the external medium. The length may be zero. Formatted records shall be read or written only by **formatted input/output statements**.

#### 12.2.3 Unformatted record

1 An unformatted record consists of a sequence of values in a processor-dependent form and may contain data of any type or may contain no data. The length of an unformatted record is measured in **file storage units**

1 (12.3.5) and depends on the output list (12.6.3) used when it is written, as well as on the processor and the  
2 external medium. The length may be zero. Unformatted records may be read or written only by [unformatted](#)  
3 [input/output statements](#).

#### 4 12.2.4 Endfile record

5 1 An endfile record is written explicitly by the [ENDFILE statement](#); the file shall be [connected](#) for sequential  
6 access. An endfile record is written implicitly to a file [connected](#) for sequential access when the most recent [data](#)  
7 [transfer statement](#) referring to the file is an [output statement](#), no intervening [file positioning statement](#) referring to  
8 the file has been executed, and

- 9 • a [REWIND](#) or [BACKSPACE](#) statement references the [unit](#) to which the file is [connected](#), or
- 10 • the [unit](#) is closed, either explicitly by a [CLOSE statement](#), implicitly by [normal termination](#), or implicitly  
11 by another [OPEN statement](#) for the same [unit](#).

12 2 An endfile record may occur only as the last record of a file. An endfile record does not have a length property.

##### NOTE 12.2

An endfile record does not necessarily have any physical embodiment. The processor can use a record count or any other means to register the position of the file at the time an [ENDFILE statement](#) is executed, so that it can take appropriate action when that position is reached again during a read operation. The endfile record, however it is implemented, is considered to exist for the [BACKSPACE statement](#) (12.8.2).

### 13 12.3 External files

#### 14 12.3.1 External file concepts

- 15 1 An [external file](#) is any file that exists in a medium external to the program.
- 16 2 At any given time, there is a processor-dependent set of allowed access methods, a processor-dependent set of  
17 allowed forms, a processor-dependent set of allowed actions, and a processor-dependent set of allowed record  
18 lengths for a file.

##### NOTE 12.3

For example, the processor-dependent set of allowed actions for a printer would likely include the write action, but not the read action.

- 19 3 A file may have a name; a file that has a name is called a named file. The name of a named file is represented by  
20 a character string value. The set of allowable names for a file is processor dependent. Whether a named file on  
21 one [image](#) is the same as a file with the same name on another [image](#) is processor dependent.

##### NOTE 12.4

If different files are needed on each [image](#), using a different file name on each [image](#) will improve portability of the code. One technique is to incorporate the [image index](#) as part of the name.

- 22 4 An [external file](#) that is [connected](#) to a [unit](#) has a position property (12.3.4).

##### NOTE 12.5

For more explanatory information on [external files](#), see C.7.1.

#### 23 12.3.2 File existence

- 24 1 At any given time, there is a processor-dependent set of [external files](#) that exist for a program. A file may be  
25 known to the processor, yet not exist for a program at a particular time.



1 2 To create a file means to cause a file to exist that did not exist previously. To delete a file means to terminate  
2 the existence of the file.

3 3 All input/output statements may refer to files that exist. A **CLOSE**, **ENDFILE**, **FLUSH**, **INQUIRE**, **OPEN**,  
4 **PRINT**, **REWIND**, or **WRITE** statement is permitted to refer to a file that does not exist. No other input/output  
5 statement shall refer to a file that does not exist. Execution of a **WRITE**, **PRINT**, or **ENDFILE** statement  
6 referring to a **preconnected** file that does not exist creates the file. This file is a different file from one **preconnected**  
7 on any other **image**.

### 8 **12.3.3 File access**

#### 9 **12.3.3.1 File access methods**

10 1 There are three methods of accessing the data of an **external file**: sequential, direct, and stream. Some files may  
11 have more than one allowed access method; other files may be restricted to one access method.

#### **NOTE 12.6**

For example, a processor might provide only sequential access to a file on magnetic tape. Thus, the set of allowed access methods depends on the file and the processor.

12 2 The method of accessing a file is determined when the file is **connected** to a **unit** (12.5.4) or when the file is  
13 created if the file is **preconnected** (12.5.5).

#### 14 **12.3.3.2 Sequential access**

15 1 Sequential access is a method of accessing the records of an external **record file** in order.

16 2 While **connected** for sequential access, an **external file** has the following properties.

- 17 • The order of the records is the order in which they were written if the direct access method is not a member  
18 of the set of allowed access methods for the file. If the direct access method is also a member of the set of  
19 allowed access methods for the file, the order of the records is the same as that specified for direct access.  
20 In this case, the first record accessible by sequential access is the record whose record number is 1 for direct  
21 access. The second record accessible by sequential access is the record whose record number is 2 for direct  
22 access, etc. A record that has not been written since the file was created shall not be read.
- 23 • The records of the file are either all formatted or all unformatted, except that the last record of the file may  
24 be an endfile record. Unless the previous reference to the file was an **output statement**, the last record, if  
25 any, of the file shall be an endfile record.
- 26 • The records of the file shall be read or written only by sequential access **data transfer statements**.

#### 27 **12.3.3.3 Direct access**

28 1 Direct access is a method of accessing the records of an external **record file** in arbitrary order.

29 2 While **connected** for direct access, an **external file** has the following properties.

- 30 • Each record of the file is uniquely identified by a positive integer called the record number. The record  
31 number of a record is specified when the record is written. Once established, the record number of a record  
32 can never be changed. The order of the records is the order of their record numbers.
- 33 • The records of the file are either all formatted or all unformatted. If the sequential access method is also a  
34 member of the set of allowed access methods for the file, its endfile record, if any, is not considered to be  
35 part of the file while it is **connected** for direct access. If the sequential access method is not a member of  
36 the set of allowed access methods for the file, the file shall not contain an endfile record.
- 37 • The records of the file shall be read or written only by direct access **data transfer statements**.
- 38 • All records of the file have the same length.

- Records need not be read or written in the order of their record numbers. Any record may be written into the file while it is **connected** to a **unit**. For example, it is permissible to write record 3, even though records 1 and 2 have not been written. Any record may be read from the file while it is **connected** to a **unit**, provided that the record has been written since the file was created, and if a **READ statement** for this connection is permitted.
- The records of the file shall not be read or written using list-directed formatting (13.10), namelist formatting (13.11), or a nonadvancing **data transfer statement** (12.3.4.2).

**NOTE 12.7**

A record cannot be deleted; however, a record can be rewritten.

**12.3.3.4 Stream access**

- 1 Stream access is a method of accessing the **file storage units** (12.3.5) of an external **stream file**.
- 2 The properties of an **external file connected** for stream access depend on whether the connection is for unformatted or formatted access. While connected for stream access, the **file storage units** of the file shall be read or written only by stream access **data transfer statements**.
- 3 While **connected** for unformatted stream access, an **external file** has the following properties.
  - Each **file storage unit** in the file is uniquely identified by a positive integer called the position. The first **file storage unit** in the file is at position 1. The position of each subsequent **file storage unit** is one greater than that of its preceding **file storage unit**.
  - If it is possible to position the file, the **file storage units** need not be read or written in order of their position. For example, it might be permissible to write the **file storage unit** at position 3, even though the **file storage units** at positions 1 and 2 have not been written. Any **file storage unit** may be read from the file while it is **connected** to a **unit**, provided that the **file storage unit** has been written since the file was created, and if a **READ statement** for this connection is permitted.
- 4 While **connected** for formatted stream access, an **external file** has the following properties.
  - Some **file storage units** of the file may contain record markers; this imposes a record structure on the file in addition to its stream structure. There might or might not be a record marker at the end of the file. If there is no record marker at the end of the file, the final record is incomplete.
  - No maximum length (12.5.6.15) is applicable to these records.
  - Writing an empty record with no record marker has no effect.
  - Each **file storage unit** in the file is uniquely identified by a positive integer called the position. The first **file storage unit** in the file is at position 1. The relationship between positions of successive **file storage units** is processor dependent; not all positive integers need correspond to valid positions.
  - If it is possible to position the file, the file position can be set to a position that was previously identified by the **POS= specifier** in an **INQUIRE statement**.
  - A processor may prohibit some control characters (6.1.1) from appearing in a formatted **stream file**.

**NOTE 12.8**

Because the record structure is determined from the record markers that are stored in the file itself, an incomplete record at the end of the file is necessarily not empty.

**NOTE 12.9**

There might be some character positions in the file that do not correspond to characters written; this is because on some processors a record marker could be written to the file as a carriage-return/line-feed or other sequence. The means of determining the position in a file **connected** for stream access is via the **POS= specifier** in an **INQUIRE statement** (12.10.2.22).

## 1 12.3.4 File position

### 2 12.3.4.1 General

3 1 Execution of certain input/output statements affects the position of an [external file](#). Certain circumstances can  
4 cause the position of a file to become indeterminate.

5 2 The initial point of a file is the position just before the first record or [file storage unit](#). The terminal point is the  
6 position just after the last record or [file storage unit](#). If there are no records or [file storage units](#) in the file, the  
7 initial point and the terminal point are the same position.

8 3 If a [record file](#) is positioned within a record, that record is the current record; otherwise, there is no current  
9 record.

10 4 Let  $n$  be the number of records in the file. If  $1 < i \leq n$  and a file is positioned within the  $i$ th record or between  
11 the  $(i - 1)$ th record and the  $i$ th record, the  $(i - 1)$ th record is the preceding record. If  $n \geq 1$  and the file is  
12 positioned at its terminal point, the preceding record is the  $n$ th and last record. If  $n = 0$  or if a file is positioned  
13 at its initial point or within the first record, there is no preceding record.

14 5 If  $1 \leq i < n$  and a file is positioned within the  $i$ th record or between the  $i$ th and  $(i + 1)$ th record, the  $(i + 1)$ th  
15 record is the next record. If  $n \geq 1$  and the file is positioned at its initial point, the first record is the next record.  
16 If  $n = 0$  or if a file is positioned at its terminal point or within the  $n$ th (last) record, there is no next record.

17 6 For a file [connected](#) for stream access, the file position is either between two [file storage units](#), at the initial point  
18 of the file, at the terminal point of the file, or undefined.

### 19 12.3.4.2 Advancing and nonadvancing input/output

20 1 An advancing input/output statement always positions a [record file](#) after the last record read or written, unless  
21 there is an error condition.

22 2 A nonadvancing input/output statement may position a [record file](#) at a character position within the current  
23 record, or a subsequent record ([13.8.2](#)). Using nonadvancing input/output, it is possible to read or write a record  
24 of the file by a sequence of [data transfer statements](#), each accessing a portion of the record. If a nonadvancing  
25 [output statement](#) leaves a file positioned within a current record and no further [output statement](#) is executed for  
26 the file before it is closed or a [BACKSPACE](#), [ENDFILE](#), or [REWIND](#) statement is executed for it, the effect is  
27 as if the [output statement](#) were the corresponding advancing [output statement](#).

### 28 12.3.4.3 File position prior to data transfer

29 1 The positioning of the file prior to data transfer depends on the method of access: sequential, direct, or stream.

30 2 For sequential access on input, if there is a current record, the file position is not changed. Otherwise, the file is  
31 positioned at the beginning of the next record and this record becomes the current record. Input shall not occur  
32 if there is no next record or if there is a current record and the last [data transfer statement](#) accessing the file  
33 performed output.

34 3 If the file contains an endfile record, the file shall not be positioned after the endfile record prior to data transfer.  
35 However, a [REWIND](#) or [BACKSPACE](#) statement may be used to reposition the file.

36 4 For sequential access on output, if there is a current record, the file position is not changed and the current record  
37 becomes the last record of the file. Otherwise, a new record is created as the next record of the file; this new  
38 record becomes the last and current record of the file and the file is positioned at the beginning of this record.

39 5 For direct access, the file is positioned at the beginning of the record specified by the [REC= specifier](#). This record  
40 becomes the current record.

41 6 For stream access, the file is positioned immediately before the [file storage unit](#) specified by the [POS= specifier](#);  
42 if there is no [POS= specifier](#), the file position is not changed.

1 7 File positioning for child [data transfer statements](#) is described in [12.6.4.8](#).

#### 2 **12.3.4.4 File position after data transfer**

3 1 If an error condition ([12.11](#)) occurred, the position of the file is indeterminate. If no error condition occurred,  
4 but an end-of-file condition ([12.11](#)) occurred as a result of reading an endfile record, the file is positioned after  
5 the endfile record.

6 2 For unformatted stream input/output, if no error condition occurred, the file position is not changed. For  
7 unformatted stream output, if the file position exceeds the previous terminal point of the file, the terminal point  
8 is set to the file position.

##### **NOTE 12.10**

An unformatted stream [output statement](#) with a [POS= specifier](#) and an empty output list can have the effect of extending the terminal point of a file without actually writing any data.

9 3 For formatted stream input, if an end-of-file condition occurred, the file position is not changed.

10 4 For nonadvancing input, if no error condition or end-of-file condition occurred, but an end-of-record condition  
11 ([12.11](#)) occurred, the file is positioned after the record just read. If no error condition, end-of-file condition, or  
12 end-of-record condition occurred in a nonadvancing [input statement](#), the file position is not changed. If no error  
13 condition occurred in a nonadvancing [output statement](#), the file position is not changed.

14 5 In all other cases, the file is positioned after the record just read or written and that record becomes the preceding  
15 record.

16 6 For a formatted stream [output statement](#), if no error condition occurred, the terminal point of the file is set to  
17 the next position after the highest-numbered position to which a datum was transferred by the statement.

##### **NOTE 12.11**

The highest-numbered position might not be the current one if the output involved a T, TL, TR, or X edit descriptor ([13.8.1](#)) and the statement is a nonadvancing [output statement](#).

#### 18 **12.3.5 File storage units**

19 1 A [file storage unit](#) is the basic unit of storage in a [stream file](#) or an unformatted [record file](#). It is the unit of file  
20 position for stream access, the unit of record length for unformatted files, and the unit of file size for all [external](#)  
21 [files](#).

22 2 Every value in a [stream file](#) or an unformatted [record file](#) shall occupy an integer number of [file storage units](#); if  
23 the [stream](#) or [record](#) file is unformatted, this number shall be the same for all scalar values of the same type and  
24 type parameters. The number of [file storage units](#) required for an item of a given type and type parameters may  
25 be determined using the [IOLENGTH= specifier](#) of the [INQUIRE statement](#) ([12.10.3](#)).

26 3 For a file [connected](#) for unformatted stream access, the processor shall not have alignment restrictions that prevent  
27 a value of any type from being stored at any positive integer file position.

28 4 The number of bits in a [file storage unit](#) is given by the constant `FILE_STORAGE_SIZE` ([16.10.2.11](#)) defined  
29 in the intrinsic module `ISO_FORTRAN_ENV`. It is recommended that the [file storage unit](#) be an 8-bit octet  
30 where this choice is practical.

##### **NOTE 12.12**

The requirement that every data value occupy an integer number of [file storage units](#) implies that data items inherently smaller than a [file storage unit](#) will require padding. This suggests that the [file storage unit](#) be small to avoid wasted space. Ideally, the file storage unit would be chosen such that padding is never required. A [file storage unit](#) of one bit would always meet this goal, but would likely be impractical because of the alignment requirements.

**NOTE 12.12 (cont.)**

The prohibition on alignment restrictions prohibits the processor from requiring data alignments larger than the [file storage unit](#).

The 8-bit octet is recommended as a good compromise that is small enough to accommodate the requirements of many applications, yet not so small that the data alignment requirements are likely to cause significant performance problems.

**12.4 Internal files**

1 [Internal files](#) provide a means of transferring and converting data from internal storage to internal storage.

2 An [internal file](#) is a [record file](#) with the following properties.

- 3 • The file is a variable of default, [ASCII](#), or [ISO 10646](#) character kind that is not an [array section](#) with a [vector subscript](#).
- 4 • A record of an [internal file](#) is a scalar character variable.
- 5 • If the file is a scalar character variable, it consists of a single record whose length is the same as the length
- 6 of the scalar character variable. If the file is a character array, it is treated as a sequence of character array
- 7 elements. Each array element, if any, is a record of the file. The ordering of the records of the file is the
- 8 same as the ordering of the array elements in the array ([9.5.3.2](#)) or the [array section](#) ([9.5.3.3](#)). Every record
- 9 of the file has the same length, which is the length of an array element in the array.
- 10 • A record of the [internal file](#) becomes defined by writing the record. If the number of characters written in
- 11 a record is less than the length of the record, the remaining portion of the record is filled with blanks. The
- 12 number of characters to be written shall not exceed the length of the record.
- 13 • A record may be read only if the record is defined.
- 14 • A record of an [internal file](#) may become defined (or undefined) by means other than an [output statement](#).
- 15 For example, the character variable may become defined by a character [assignment statement](#).
- 16 • An [internal file](#) is always positioned at the beginning of the first record prior to data transfer, except for
- 17 child [data transfer statements](#) ([12.6.4.8](#)). This record becomes the current record.
- 18 • The initial value of a connection mode ([12.5.2](#)) is the value that would be implied by an initial [OPEN](#)
- 19 [statement](#) without the corresponding keyword.
- 20 • Reading and writing records shall be accomplished only by sequential access formatted [data transfer state-](#)
- 21 [ments](#).
- 22 • An [internal file](#) shall not be specified as the [unit](#) in a [CLOSE](#), [INQUIRE](#), or [OPEN](#) statement.

**12.5 File connection****12.5.1 Referring to a file**

1 A [unit](#), specified by an [io-unit](#), provides a means for referring to a file.

2 R1201 *io-unit* is *file-unit-number*  
 3 or \*  
 4 or *internal-file-variable*

5 R1202 *file-unit-number* is *scalar-int-expr*

6 R1203 *internal-file-variable* is *char-variable*

7 C1201 (R1203) The *char-variable* shall not be an [array section](#) with a [vector subscript](#).

8 C1202 (R1203) The *char-variable* shall be default character, [ASCII character](#), or [ISO 10646 character](#).

- 1 2 A **unit** is either an **external unit** or an **internal unit**. An **external unit** is used to refer to an **external file** and  
 2 is specified by an asterisk or a *file-unit-number*. The value of *file-unit-number* shall be nonnegative, equal to  
 3 one of the **named constants** `INPUT_UNIT`, `OUTPUT_UNIT`, or `ERROR_UNIT` of the intrinsic module `ISO_`  
 4 `FORTRAN_ENV` (16.10.2), the **unit** argument of an active **defined input/output** procedure (12.6.4.8), or a  
 5 `NEWUNIT` value (12.5.6.12). An **internal unit** is used to refer to an **internal file** and is specified by an *internal-*  
 6 *file-variable* or a *file-unit-number* whose value is equal to the **unit** argument of an active **defined input/output**  
 7 procedure. The value of a *file-unit-number* shall identify a valid **unit**.
- 8 3 On an **image**, the **external unit** identified by a particular value of a *scalar-int-expr* is the same **external unit** in  
 9 all **program units**.

**NOTE 12.13**

In the example:

```
SUBROUTINE A
  READ (6) X
  . . .
SUBROUTINE B
  N = 6
  REWIND N
```

the value 6 used in both **program units** identifies the same **external unit**.

- 10 4 In a **READ statement**, an *io-unit* that is an asterisk identifies an **external unit** that is **preconnected** for sequential  
 11 formatted input on **image 1** in the **initial team** only (12.6.4.3). This **unit** is also identified by the value of  
 12 the **named constant** `INPUT_UNIT` of the intrinsic module `ISO_FORTRAN_ENV` (16.10.2.13). In a **WRITE**  
 13 **statement**, an *io-unit* that is an asterisk identifies an **external unit** that is **preconnected** for sequential formatted  
 14 output. This **unit** is also identified by the value of the **named constant** `OUTPUT_UNIT` of the intrinsic module  
 15 `ISO_FORTRAN_ENV` (16.10.2.22).
- 16 5 This document identifies a processor-dependent **external unit** for the purpose of error reporting. This **unit** shall  
 17 be **preconnected** for sequential formatted output. The processor may define this to be the same as the output  
 18 **unit** identified by an asterisk. This **unit** is also identified by a unit number defined by the **named constant**  
 19 `ERROR_UNIT` of the intrinsic module `ISO_FORTRAN_ENV`.

**NOTE 12.14**

Even though `OUTPUT_UNIT` is connected to a separate file on each **image**, it is expected that the processor could merge the sequences of records from these files into a single sequence of records that is sent to the physical device associated with this **unit**, such as the user's terminal. If `ERROR_UNIT` is associated with the same physical device, the sequences of records from files connected to `ERROR_UNIT` on each of the **images** could be merged into the same sequence generated from the `OUTPUT_UNIT` files. Otherwise, it is expected that the sequence of records in the files connected to `ERROR_UNIT` on each **image** could be merged into a single sequence of records that is sent to the physical device associated with `ERROR_UNIT`.

## 20 12.5.2 Connection modes

- 21 1 A connection for formatted input/output has several changeable modes: these are the blank interpretation mode  
 22 (13.8.6), delimiter mode (13.10.4, 13.11.4.2), sign mode (13.8.4), decimal edit mode (13.8.8), input/output round-  
 23 ing mode (13.7.2.3.8), pad mode (12.6.4.5.3), and scale factor (13.8.5). A connection for unformatted input/output  
 24 has no changeable modes.
- 25 2 Values for the modes of a connection are established when the connection is initiated. If the connection is initiated  
 26 by an **OPEN statement**, the values are as specified, either explicitly or implicitly, by the **OPEN statement**. If the  
 27 connection is initiated other than by an **OPEN statement** (that is, if the file is an **internal file** or **preconnected file**)  
 28 the values established are those that would be implied by an initial **OPEN statement** without the corresponding  
 29 keywords.

- 1 3 The scale factor cannot be explicitly specified in an **OPEN statement**; it is implicitly 0.
- 2 4 The modes of a connection to an **external file** may be changed by a subsequent **OPEN statement** that modifies  
3 the connection.
- 4 5 The modes of a connection may be temporarily changed by a corresponding keyword specifier in a **data transfer  
5 statement** or by an edit descriptor. Keyword specifiers take effect at the beginning of execution of the **data  
6 transfer statement**. Edit descriptors take effect when they are encountered in format processing. When a **data  
7 transfer statement** terminates, the values for the modes are reset to the values in effect immediately before the  
8 **data transfer statement** was executed.

### 9 12.5.3 Unit existence

- 10 1 At any given time, there is a processor-dependent set of **external units** that exist for an **image**.
- 11 2 All input/output statements are permitted to refer to **units** that exist. The **CLOSE**, **INQUIRE**, and **WAIT**  
12 statements are also permitted to refer to **units** that do not exist. No other input/output statement shall refer to  
13 a **unit** that does not exist.

### 14 12.5.4 Connection of a file to a unit

- 15 1 An **external unit** has a property of being **connected** or not connected. If **connected**, it refers to an **external file**. An  
16 **external unit** may become **connected** by **preconnection** or by the execution of an **OPEN statement**. The property  
17 of connection is symmetric; the unit is **connected** to a file if and only if the file is **connected** to the **unit**.
- 18 2 Every input/output statement except an **OPEN**, **CLOSE**, **INQUIRE**, or **WAIT** statement shall refer to a **unit**  
19 that is **connected** to a file and thereby make use of or affect that file.
- 20 3 A file may be **connected** and not exist (12.3.2).

#### NOTE 12.15

An example is a **preconnected external file** that has not yet been written.

- 21 4 A **unit** shall not be **connected** to more than one file at the same time. However, means are provided to change  
22 the status of an **external unit** and to connect a **unit** to a different file. It is processor dependent whether a file  
23 can be **connected** to more than one **unit** at the same time.
- 24 5 This document defines means of portable interoperation with C. C streams are described in 7.21.2 of ISO/IEC  
25 9899:2011. Whether a **unit** can be **connected** to a file that is also **connected** to a C stream is processor dependent.  
26 If a **unit** is **connected** to a file that is also **connected** to a C stream, the results of performing input/output  
27 operations on such a file are processor dependent. It is processor dependent whether the files **connected** to  
28 the **units** **INPUT\_UNIT**, **OUTPUT\_UNIT**, and **ERROR\_UNIT** correspond to the predefined C text streams  
29 standard input, standard output, and standard error. If a main program or procedure defined by means of Fortran  
30 and a main program or procedure defined by means other than Fortran perform input/output operations on the  
31 same **external file**, the results are processor dependent. A main program or procedure defined by means of Fortran  
32 and a main program or procedure defined by means other than Fortran can perform input/output operations on  
33 different **external files** without interference.
- 34 6 If input/output operations are performed on more than one unit while they are connected to the same **external  
35 file**, the results are processor dependent.
- 36 7 After an **external unit** has been disconnected by the execution of a **CLOSE statement**, it may be **connected** again  
37 within the same program to the same file or to a different file. After an **external file** has been disconnected by  
38 the execution of a **CLOSE statement**, it may be **connected** again within the same program to the same **unit** or  
39 to a different **unit**.

**NOTE 12.16**

The only means of referencing a file that has been disconnected is by the appearance of its name in an **OPEN** or **INQUIRE** statement. There might be no means of reconnecting an unnamed file once it is disconnected.

1 8 An **internal unit** is always **connected** to the **internal file** designated by the variable that identifies the **unit**.

**NOTE 12.17**

For more explanatory information on file connection properties, see **C.7.4**.

## 2 12.5.5 Preconnection

3 1 Preconnection means that the **unit** is **connected** to a file at the beginning of execution of the program and therefore  
4 it may be specified in input/output statements without the prior execution of an **OPEN statement**.

## 5 12.5.6 OPEN statement

### 6 12.5.6.1 General

7 1 An **OPEN statement** initiates or modifies the connection between an **external file** and a specified **unit**. The **OPEN**  
8 statement may be used to connect an existing file to a **unit**, create a file that is **preconnected**, create a file and  
9 connect it to a **unit**, or change certain modes of a connection between a file and a **unit**.

10 2 An **external unit** may be **connected** by an **OPEN statement** in the main program or any subprogram.

11 3 If the file to be **connected** to the **unit** does not exist but is the same as the file to which the **unit** is **preconnected**,  
12 the modes specified by an **OPEN statement** become a part of the connection.

13 4 If the file to be **connected** to the **unit** is not the same as the file to which the **unit** is **connected**, the effect is as  
14 if a **CLOSE statement** without a **STATUS= specifier** had been executed for the **unit** immediately prior to the  
15 execution of an **OPEN statement**.

16 5 If a **unit** is **connected** to a file that exists, execution of an **OPEN statement** for that **unit** is permitted. If the  
17 **FILE= specifier** is not included in such an **OPEN statement**, the file to be **connected** to the **unit** is the same as  
18 the file to which the **unit** is already **connected**.

19 6 If the file to be **connected** to the **unit** is the same as the file to which the **unit** is **connected**, a new connection is not  
20 established and values for any changeable modes (**12.5.2**) specified come into effect for the established connection;  
21 the current file position is unaffected. Before any effect on changeable modes, a wait operation is performed for  
22 any pending asynchronous data transfer operations for the specified **unit**. If the **POSITION= specifier** appears  
23 in such an **OPEN statement**, the value specified shall not disagree with the current position of the file. If the  
24 **STATUS= specifier** is included in such an **OPEN statement**, it shall be specified with the value **OLD**. Other than  
25 **ERR=**, **IOSTAT=**, and **IOMSG=**, and the changeable modes, the values of all other specifiers in such an **OPEN**  
26 statement shall not differ from those in effect for the established connection.

27 7 A **STATUS= specifier** with a value of **OLD** is always allowed when the file to be **connected** to the **unit** is the same  
28 as the file to which the **unit** is **connected**. In this case, if the status of the file was **SCRATCH** before execution of  
29 the **OPEN statement**, the file will still be deleted when the **unit** is closed, and the file is still considered to have  
30 a status of **SCRATCH**.

### 31 12.5.6.2 Syntax of the OPEN statement

32 R1204 *open-stmt* is OPEN ( *connect-spec-list* )

33 R1205 *connect-spec* is [ UNIT = ] *file-unit-number*  
34 or ACCESS = *scalar-default-char-expr*  
35 or ACTION = *scalar-default-char-expr*



- 1 or ASYNCHRONOUS = *scalar-default-char-expr*  
 2 or BLANK = *scalar-default-char-expr*  
 3 or DECIMAL = *scalar-default-char-expr*  
 4 or DELIM = *scalar-default-char-expr*  
 5 or ENCODING = *scalar-default-char-expr*  
 6 or ERR = *label*  
 7 or FILE = *file-name-expr*  
 8 or FORM = *scalar-default-char-expr*  
 9 or IOMSG = *iomsg-variable*  
 10 or IOSTAT = *scalar-int-variable*  
 11 or NEWUNIT = *scalar-int-variable*  
 12 or PAD = *scalar-default-char-expr*  
 13 or POSITION = *scalar-default-char-expr*  
 14 or RECL = *scalar-int-expr*  
 15 or ROUND = *scalar-default-char-expr*  
 16 or SIGN = *scalar-default-char-expr*  
 17 or STATUS = *scalar-default-char-expr*
- 18 R1206 *file-name-expr* is *scalar-default-char-expr*
- 19 R1207 *iomsg-variable* is *scalar-default-char-variable*
- 20 C1203 No specifier shall appear more than once in a given *connect-spec-list*.
- 21 C1204 (R1204) If the **NEWUNIT= specifier** does not appear, a *file-unit-number* shall be specified; if the optional  
 22 characters UNIT= are omitted, the *file-unit-number* shall be the first item in the *connect-spec-list*.
- 23 C1205 (R1204) If a **NEWUNIT= specifier** appears, a *file-unit-number* shall not appear.
- 24 C1206 (R1204) The *label* used in the **ERR= specifier** shall be the statement label of a **branch target statement**  
 25 that appears in the same **inclusive scope** as the OPEN statement.
- 26 1 A specifier that requires a *scalar-default-char-expr* may have a limited list of character values. These values are  
 27 listed for each such specifier. Any trailing blanks are ignored. The value specified is without regard to case. Some  
 28 specifiers have a default value if the specifier is omitted.
- 29 2 The **IOSTAT=**, **ERR=**, and **IOMSG=** specifiers are described in 12.11.

**NOTE 12.18**

An example of an OPEN statement is:

```
OPEN (10, FILE = 'employee.names', ACTION = 'READ', PAD = 'YES')
```

**NOTE 12.19**

For more explanatory information on the OPEN statement, see C.7.3.

### 30 12.5.6.3 ACCESS= specifier in the OPEN statement

- 31 1 The *scalar-default-char-expr* shall evaluate to SEQUENTIAL, DIRECT, or STREAM. The ACCESS= specifier  
 32 specifies the access method for the connection of the file as being sequential, direct, or stream. If this specifier is  
 33 omitted, the default value is SEQUENTIAL. For an existing file, the specified access method shall be included in  
 34 the set of allowed access methods for the file. For a new file, the processor creates the file with a set of allowed  
 35 access methods that includes the specified method.

#### 1 12.5.6.4 ACTION= specifier in the OPEN statement

2 1 The *scalar-default-char-expr* shall evaluate to READ, WRITE, or READWRITE. READ specifies that the  
3 WRITE, PRINT, and ENDFILE statements shall not refer to this connection. WRITE specifies that READ  
4 statements shall not refer to this connection. READWRITE permits any input/output statements to refer to this  
5 connection. If this specifier is omitted, the default value is processor dependent. If READWRITE is included in  
6 the set of allowable actions for a file, both READ and WRITE also shall be included in the set of allowed actions  
7 for that file. For an existing file, the specified action shall be included in the set of allowed actions for the file.  
8 For a new file, the processor creates the file with a set of allowed actions that includes the specified action.

#### 9 12.5.6.5 ASYNCHRONOUS= specifier in the OPEN statement

10 1 The *scalar-default-char-expr* shall evaluate to YES or NO. If YES is specified, asynchronous input/output on  
11 the *unit* is allowed. If NO is specified, asynchronous input/output on the *unit* is not allowed. If this specifier is  
12 omitted, the default value is NO.

#### 13 12.5.6.6 BLANK= specifier in the OPEN statement

14 1 The *scalar-default-char-expr* shall evaluate to NULL or ZERO. The BLANK= specifier is permitted only for a  
15 connection for formatted input/output. It specifies the blank interpretation mode (13.8.6, 12.6.2.6) for input for  
16 this connection. This mode has no effect on output. It is a changeable mode (12.5.2). If this specifier is omitted  
17 in an OPEN statement that initiates a connection, the default value is NULL.

#### 18 12.5.6.7 DECIMAL= specifier in the OPEN statement

19 1 The *scalar-default-char-expr* shall evaluate to COMMA or POINT. The DECIMAL= specifier is permitted only  
20 for a connection for formatted input/output. It specifies the decimal edit mode (13.6, 13.8.8, 12.6.2.7) for this  
21 connection. It is a changeable mode (12.5.2). If this specifier is omitted in an OPEN statement that initiates a  
22 connection, the default value is POINT.

#### 23 12.5.6.8 DELIM= specifier in the OPEN statement

24 1 The *scalar-default-char-expr* shall evaluate to APOSTROPHE, QUOTE, or NONE. The DELIM= specifier is  
25 permitted only for a connection for formatted input/output. It specifies the delimiter mode (12.6.2.8) for list-  
26 directed (13.10.4) and namelist (13.11.4.2) output for the connection. This mode has no effect on input. It is  
27 a changeable mode (12.5.2). If this specifier is omitted in an OPEN statement that initiates a connection, the  
28 default value is NONE.

#### 29 12.5.6.9 ENCODING= specifier in the OPEN statement

30 1 The *scalar-default-char-expr* shall evaluate to UTF-8 or DEFAULT. The ENCODING= specifier is permitted  
31 only for a connection for formatted input/output. The value UTF-8 specifies that the encoding form of the file  
32 is UTF-8 as specified in ISO/IEC 10646. Such a file is called a Unicode file, and all characters therein are of ISO  
33 10646 character kind. The value UTF-8 shall not be specified if the processor does not support the ISO 10646  
34 character kind. The value DEFAULT specifies that the encoding form of the file is processor dependent. If this  
35 specifier is omitted in an OPEN statement that initiates a connection, the default value is DEFAULT.

#### 36 12.5.6.10 FILE= specifier in the OPEN statement

37 1 The value of the FILE= specifier is the name of the file to be connected to the specified *unit*. Any trailing blanks  
38 are ignored. The *file-name-expr* shall be a name that is allowed by the processor. The interpretation of case is  
39 processor dependent.

40 2 This specifier shall appear if the STATUS= specifier has the value NEW or REPLACE. This specifier shall not  
41 appear if the STATUS= specifier has the value SCRATCH. If the STATUS= specifier has the value OLD, this  
42 specifier shall appear unless the *unit* is connected and the file connected to the *unit* exists. If this specifier

1 is omitted and the `unit` is not connected to a file, the `STATUS= specifier` shall be specified with a value of  
2 SCRATCH; in this case, the connection is made to a processor-dependent file.

### 3 12.5.6.11 FORM= specifier in the OPEN statement

4 1 The *scalar-default-char-expr* shall evaluate to FORMATTED or UNFORMATTED. The FORM= specifier de-  
5 termines whether the file is being connected for formatted or unformatted input/output. If this specifier is  
6 omitted, the default value is UNFORMATTED if the file is being connected for direct access or stream access,  
7 and the default value is FORMATTED if the file is being connected for sequential access. For an existing file,  
8 the specified form shall be included in the set of allowed forms for the file. For a new file, the processor creates  
9 the file with a set of allowed forms that includes the specified form.

### 10 12.5.6.12 NEWUNIT= specifier in the OPEN statement

11 1 If this specifier appears in an OPEN statement, either the `FILE= specifier` shall appear, or the `STATUS= specifier`  
12 shall appear with a value of SCRATCH.

13 2 The variable is defined with a processor determined NEWUNIT value if no error occurs during the execution of  
14 the OPEN statement. If an error occurs, the processor shall not change the value of the variable.

15 3 A NEWUNIT value is a negative number, and shall not be equal to  $-1$ , any of the `named constants` ER-  
16 ROR\_UNIT, INPUT\_UNIT, or OUTPUT\_UNIT from the intrinsic module ISO\_FORTRAN\_ENV (16.10.2),  
17 any value used by the processor for the `unit` argument to a `defined input/output` procedure, nor any previous  
18 NEWUNIT value that identifies a file that is `connected`. The `unit` identified by a NEWUNIT value shall not be  
19 `preconnected`.

### 20 12.5.6.13 PAD= specifier in the OPEN statement

21 1 The *scalar-default-char-expr* shall evaluate to YES or NO. The PAD= specifier is permitted only for a connection  
22 for formatted input/output. It specifies the pad mode (12.6.4.5.3, 12.6.2.10) for input for this connection. This  
23 mode has no effect on output. It is a changeable mode (12.5.2). If this specifier is omitted in an OPEN statement  
24 that initiates a connection, the default value is YES.

### 25 12.5.6.14 POSITION= specifier in the OPEN statement

26 1 The *scalar-default-char-expr* shall evaluate to ASIS, REWIND, or APPEND. The connection shall be for sequen-  
27 tial or stream access. A new file is positioned at its initial point. REWIND positions an existing file at its initial  
28 point. APPEND positions an existing file such that the endfile record is the next record, if it has one. If an  
29 existing file does not have an endfile record, APPEND positions the file at its terminal point. ASIS leaves the  
30 position unchanged if the file exists and already is `connected`. If the file exists but is not connected, the position  
31 resulting from ASIS is processor dependent. If this specifier is omitted, the default value is ASIS.

### 32 12.5.6.15 RECL= specifier in the OPEN statement

33 1 The value of the RECL= specifier shall be positive. It specifies the length of each record in a file being connected  
34 for direct access, or specifies the maximum length of a record in a file being connected for sequential access. This  
35 specifier shall not appear when a file is being connected for stream access. This specifier shall appear when a  
36 file is being connected for direct access. If this specifier is omitted when a file is being connected for sequential  
37 access, the default value is processor dependent. If the file is being connected for formatted input/output, the  
38 length is the number of characters for all records that contain only characters of default kind. When a record  
39 contains any nondefault characters, the effect of the RECL= specifier is processor dependent. If the file is being  
40 connected for unformatted input/output, the length is measured in `file storage units`. For an existing file, the  
41 value of the RECL= specifier shall be included in the set of allowed record lengths for the file. For a new file, the  
42 processor creates the file with a set of allowed record lengths that includes the specified value.

### 1 12.5.6.16 ROUND= specifier in the OPEN statement

- 2 1 The *scalar-default-char-expr* shall evaluate to one of UP, DOWN, ZERO, NEAREST, COMPATIBLE, or PRO-  
3 CESSOR\_DEFINED. The ROUND= specifier is permitted only for a connection for formatted input/output.  
4 It specifies the input/output rounding mode (13.7.2.3.8, 12.6.2.13) for this connection. It is a changeable mode  
5 (12.5.2). If this specifier is omitted in an OPEN statement that initiates a connection, the input/output rounding  
6 mode is processor dependent; it shall be one of the above modes.

#### NOTE 12.20

A processor is free to select any input/output rounding mode for the default mode. The mode might correspond to UP, DOWN, ZERO, NEAREST, or COMPATIBLE; or it might be a completely different input/output rounding mode.

### 7 12.5.6.17 SIGN= specifier in the OPEN statement

- 8 1 The *scalar-default-char-expr* shall evaluate to one of PLUS, SUPPRESS, or PROCESSOR\_DEFINED. The  
9 SIGN= specifier is permitted only for a connection for formatted input/output. It specifies the sign mode  
10 (13.8.4, 12.6.2.14) for this connection. It is a changeable mode (12.5.2). If this specifier is omitted in an OPEN  
11 statement that initiates a connection, the default value is PROCESSOR\_DEFINED.

### 12 12.5.6.18 STATUS= specifier in the OPEN statement

- 13 1 The *scalar-default-char-expr* shall evaluate to OLD, NEW, SCRATCH, REPLACE, or UNKNOWN. If OLD is  
14 specified, the file shall exist. If NEW is specified, the file shall not exist.
- 15 2 Successful execution of an OPEN statement with NEW specified creates the file and changes the status to OLD.  
16 If REPLACE is specified and the file does not already exist, the file is created and the status is changed to OLD.  
17 If REPLACE is specified and the file does exist, the file is deleted, a new file is created with the same name, and  
18 the status is changed to OLD. If SCRATCH is specified, the file is created and connected to the specified **unit**  
19 for use by the program but is deleted at the execution of a **CLOSE statement** referring to the same **unit** or at  
20 the **normal termination** of the program.
- 21 3 If UNKNOWN is specified, the status is processor dependent. If this specifier is omitted, the default value is  
22 UNKNOWN.

#### NOTE 12.21

SCRATCH cannot be specified if the **FILE= specifier** appears (12.5.6.10).

## 23 12.5.7 CLOSE statement

### 24 12.5.7.1 General

- 25 1 The CLOSE statement is used to terminate the connection of a specified **unit** to an **external file**.
- 26 2 Execution of a CLOSE statement for a **unit** may occur in any **program unit** of a program and need not occur in  
27 the same **program unit** as the execution of an **OPEN statement** referring to that **unit**.
- 28 3 Execution of a CLOSE statement performs a wait operation for any pending asynchronous data transfer operations  
29 for the specified **unit**.
- 30 4 Execution of a CLOSE statement specifying a **unit** that does not exist, exists but is connected to a file that does  
31 not exist, or has no file **connected** to it, is permitted and affects no file or **unit**.
- 32 5 After a **unit** has been disconnected by execution of a CLOSE statement, it may be **connected** again within the  
33 same program, either to the same file or to a different file. After a named file has been disconnected by execution  
34 of a CLOSE statement, it may be **connected** again within the same program, either to the same **unit** or to a  
35 different **unit**, provided that the file still exists.

- 1 6 During the completion step (5.3.7) of normal termination, all units that are connected are closed. Each unit is  
 2 closed with status KEEP unless the file status prior to termination of execution was SCRATCH, in which case  
 3 the unit is closed with status DELETE.

**NOTE 12.22**

The effect is as though a CLOSE statement without a STATUS= specifier were executed on each connected unit.

4 **12.5.7.2 Syntax**

5 R1208 *close-stmt* is CLOSE ( *close-spec-list* )

6 R1209 *close-spec* is [ UNIT = ] *file-unit-number*  
 7 or IOSTAT = *scalar-int-variable*  
 8 or IOMSG = *iomsg-variable*  
 9 or ERR = *label*  
 10 or STATUS = *scalar-default-char-expr*

11 C1207 No specifier shall appear more than once in a given *close-spec-list*.

12 C1208 A *file-unit-number* shall be specified in a *close-spec-list*; if the optional characters UNIT= are omitted,  
 13 the *file-unit-number* shall be the first item in the *close-spec-list*.

14 C1209 (R1209) The *label* used in the ERR= specifier shall be the statement label of a branch target statement  
 15 that appears in the same inclusive scope as the CLOSE statement.

16 1 The *scalar-default-char-expr* has a limited list of character values. Any trailing blanks are ignored. The value  
 17 specified is without regard to case.

18 2 The IOSTAT=, ERR=, and IOMSG= specifiers are described in 12.11.

**NOTE 12.23**

An example of a CLOSE statement is:

```
CLOSE (10, STATUS = 'KEEP')
```

19 **12.5.7.3 STATUS= specifier in the CLOSE statement**

- 20 1 The *scalar-default-char-expr* shall evaluate to KEEP or DELETE. The STATUS= specifier determines the dis-  
 21 position of the file that is connected to the specified unit. KEEP shall not be specified for a file whose status prior  
 22 to execution of a CLOSE statement is SCRATCH. If KEEP is specified for a file that exists, the file continues  
 23 to exist after the execution of a CLOSE statement. If KEEP is specified for a file that does not exist, the file  
 24 will not exist after the execution of a CLOSE statement. If DELETE is specified, the file will not exist after the  
 25 execution of a CLOSE statement. If this specifier is omitted, the default value is KEEP, unless the file status  
 26 prior to execution of the CLOSE statement is SCRATCH, in which case the default value is DELETE.

27 **12.6 Data transfer statements**28 **12.6.1 Form of input and output statements**

- 29 1 The READ statement is the data transfer input statement. The WRITE statement and the PRINT statement  
 30 are the data transfer output statements.

31 R1210 *read-stmt* is READ ( *io-control-spec-list* ) [ *input-item-list* ]  
 32 or READ *format* [ , *input-item-list* ]

33 R1211 *write-stmt* is WRITE ( *io-control-spec-list* ) [ *output-item-list* ]

1 R1212 *print-stmt* is PRINT *format* [ , *output-item-list* ]

#### NOTE 12.24

Examples of data transfer statements are:

```

READ (6, *) SIZE
READ 10, A, B
WRITE (6, 10) A, S, J
PRINT 10, A, S, J
10 FORMAT (2E16.3, I5)

```

## 2 12.6.2 Control information list

### 3 12.6.2.1 Syntax

4 1 A control information list is an *io-control-spec-list*. It governs data transfer.

5 R1213 *io-control-spec* is [ UNIT = ] *io-unit*  
6 or [ FMT = ] *format*  
7 or [ NML = ] *namelist-group-name*  
8 or ADVANCE = *scalar-default-char-expr*  
9 or ASYNCHRONOUS = *scalar-default-char-constant-expr*  
10 or BLANK = *scalar-default-char-expr*  
11 or DECIMAL = *scalar-default-char-expr*  
12 or DELIM = *scalar-default-char-expr*  
13 or END = *label*  
14 or EOR = *label*  
15 or ERR = *label*  
16 or ID = *id-variable*  
17 or IOMSG = *iomsg-variable*  
18 or IOSTAT = *scalar-int-variable*  
19 or PAD = *scalar-default-char-expr*  
20 or POS = *scalar-int-expr*  
21 or REC = *scalar-int-expr*  
22 or ROUND = *scalar-default-char-expr*  
23 or SIGN = *scalar-default-char-expr*  
24 or SIZE = *scalar-int-variable*

25 R1214 *id-variable* is *scalar-int-variable*

26 C1210 No specifier shall appear more than once in a given *io-control-spec-list*.

27 C1211 An *io-unit* shall be specified in an *io-control-spec-list*; if the optional characters UNIT= are omitted, the  
28 *io-unit* shall be the first item in the *io-control-spec-list*.

29 C1212 (R1213) A DELIM= or SIGN= specifier shall not appear in a *read-stmt*.

30 C1213 (R1213) A BLANK=, PAD=, END=, EOR=, or SIZE= specifier shall not appear in a *write-stmt*.

31 C1214 (R1213) The *label* in the ERR=, EOR=, or END= specifier shall be the statement label of a *branch*  
32 *target statement* that appears in the same *inclusive scope* as the data transfer statement.

33 C1215 (R1213) A *namelist-group-name* shall be the name of a namelist group.

34 C1216 (R1213) A *namelist-group-name* shall not appear if a REC= specifier, *format*, *input-item-list*, or an  
35 *output-item-list* appears in the data transfer statement.

36 C1217 (R1213) If *format* appears without a preceding FMT=, it shall be the second item in the *io-control-spec-*

- 1 *list* and the first item shall be *io-unit*.
- 2 C1218 (R1213) If *namelist-group-name* appears without a preceding NML=, it shall be the second item in the  
3 *io-control-spec-list* and the first item shall be *io-unit*.
- 4 C1219 (R1213) If *io-unit* is not a *file-unit-number*, the *io-control-spec-list* shall not contain a REC= specifier or  
5 a POS= specifier.
- 6 C1220 (R1213) If the REC= specifier appears, an END= specifier shall not appear, and the *format*, if any, shall  
7 not be an asterisk.
- 8 C1221 (R1213) An ADVANCE= specifier may appear only in a formatted sequential or stream data transfer  
9 statement with explicit format specification (13.2) whose *io-control-spec-list* does not contain an *internal-*  
10 *file-variable* as the *io-unit*.
- 11 C1222 (R1213) If an EOR= specifier appears, an ADVANCE= specifier also shall appear.
- 12 C1223 (R1213) The *scalar-default-char-constant-expr* in an ASYNCHRONOUS= specifier shall have the value  
13 YES or NO.
- 14 C1224 (R1213) An ASYNCHRONOUS= specifier with a value YES shall not appear unless *io-unit* is a *file-*  
15 *unit-number*.
- 16 C1225 (R1213) If an ID= specifier appears, an ASYNCHRONOUS= specifier with the value YES shall also  
17 appear.
- 18 C1226 (R1213) If a POS= specifier appears, the *io-control-spec-list* shall not contain a REC= specifier.
- 19 C1227 (R1213) If a DECIMAL=, BLANK=, PAD=, SIGN=, or ROUND= specifier appears, a *format* or  
20 *namelist-group-name* shall also appear.
- 21 C1228 (R1213) If a DELIM= specifier appears, either *format* shall be an asterisk or *namelist-group-name* shall  
22 appear.
- 23 C1229 (R1214) The *scalar-int-variable* shall have a decimal exponent range no smaller than that of default  
24 integer.
- 25 2 If an EOR= specifier appears, an ADVANCE= specifier with the value NO shall also appear.
- 26 3 If the data transfer statement contains a *format* or *namelist-group-name*, the statement is a formatted in-  
27 put/output statement; otherwise, it is an unformatted input/output statement.
- 28 4 The ADVANCE=, ASYNCHRONOUS=, DECIMAL=, BLANK=, DELIM=, PAD=, SIGN=, and ROUND=  
29 specifiers have a limited list of character values. Any trailing blanks are ignored. The values specified are without  
30 regard to case.
- 31 5 The IOSTAT=, ERR=, EOR=, END=, and IOMSG= specifiers are described in 12.11.

**NOTE 12.25**

An example of a READ statement is:

```
READ (IOSTAT = IOS, UNIT = 6, FMT = '(10F8.2)') A, B
```

**12.6.2.2 Format specification in a data transfer statement**

- 32 1 The *format* specifier supplies a format specification or specifies list-directed formatting for a formatted in-  
33 put/output statement.  
34

1 R1215 *format* is *default-char-expr*  
 2 or *label*  
 3 or \*

4 C1230 (R1215) The *label* shall be the label of a **FORMAT statement** that appears in the same **inclusive scope**  
 5 as the statement containing the FMT= specifier.

6 2 The *default-char-expr* shall evaluate to a valid format specification (13.2.1 and 13.2.2).

7 3 If *default-char-expr* is an array, it is treated as if all of the elements of the array were specified in array element  
 8 order and were concatenated.

9 4 If *format* is \*, the statement is a list-directed input/output statement.

#### NOTE 12.26

An example in which the format is a character expression is:

```
READ (6, FMT = "(" // CHAR_FMT // ")" ) X, Y, Z
```

where CHAR\_FMT is a default character variable.

#### 10 12.6.2.3 NML= specifier in a data transfer statement

11 1 The NML= specifier supplies the *namelist-group-name* (8.9). This name identifies a particular collection of data  
 12 objects on which transfer is to be performed.

13 2 If a *namelist-group-name* appears, the statement is a namelist input/output statement.

#### 14 12.6.2.4 ADVANCE= specifier in a data transfer statement

15 1 The *scalar-default-char-expr* shall evaluate to YES or NO. The ADVANCE= specifier determines whether advan-  
 16 cing input/output occurs for a nonchild data transfer statement. If YES is specified for a nonchild data transfer  
 17 statement, advancing input/output occurs. If NO is specified, nonadvancing input/output occurs (12.3.4.2). If  
 18 this specifier is omitted from a nonchild data transfer statement that allows the specifier, the default value is  
 19 YES. A formatted child data transfer statement is a nonadvancing input/output statement, and any ADVANCE=  
 20 specifier is ignored.

#### 21 12.6.2.5 ASYNCHRONOUS= specifier in a data transfer statement

22 1 The ASYNCHRONOUS= specifier determines whether this data transfer statement is synchronous or asynchron-  
 23 ous. If YES is specified, the statement and the input/output operation are asynchronous. If NO is specified or if  
 24 the specifier is omitted, the statement and the input/output operation are synchronous.

25 2 Asynchronous input/output is permitted only for **external files** opened with an ASYNCHRONOUS= specifier  
 26 with the value YES in the **OPEN statement**.

#### NOTE 12.27

Both synchronous and asynchronous input/output are allowed for files opened with an ASYNCHRONOUS= specifier of YES. For other files, only synchronous input/output is allowed; this includes files opened with an ASYNCHRONOUS= specifier of NO, files opened without an ASYNCHRONOUS= specifier, **preconnected** files accessed without an **OPEN statement**, and **internal files**.

The ASYNCHRONOUS= specifier value in a data transfer statement is a **constant expression** because it effects compiler optimizations and, therefore, needs to be known at compile time.

27 3 The processor may perform an asynchronous data transfer operation asynchronously, but it is not required to do  
 28 so. For each **external file**, records and **file storage units** read or written by asynchronous data transfer statements



1 are read, written, and processed in the same order as they would have been if the data transfer statements were  
2 synchronous. The documentation of the Fortran processor should describe when input/output will be performed  
3 asynchronously.

4 4 If a variable is used in an asynchronous data transfer statement as

- 5 • an item in an input/output list,
- 6 • a group object in a namelist, or
- 7 • a SIZE= specifier,

8 the [base object](#) of the *data-ref* is implicitly given the [ASYNCHRONOUS](#) attribute in the [scoping unit](#) of the data  
9 transfer statement. This attribute may be confirmed by explicit declaration.

10 5 When an asynchronous input/output statement is executed, the set of [storage units](#) specified by the item list or  
11 NML= specifier, plus the [storage units](#) specified by the SIZE= specifier, is defined to be the pending input/output  
12 storage sequence for the data transfer operation.

#### NOTE 12.28

A pending input/output storage sequence is not necessarily a [contiguous](#) set of [storage units](#).

13 6 A pending input/output storage sequence affector is a variable of which any part is associated with a [storage unit](#)  
14 in a pending input/output storage sequence.

#### 15 12.6.2.6 BLANK= specifier in a data transfer statement

16 1 The *scalar-default-char-expr* shall evaluate to NULL or ZERO. The BLANK= specifier temporarily changes  
17 (12.5.2) the blank interpretation mode (13.8.6, 12.5.6.6) for the connection. If the specifier is omitted, the mode  
18 is not changed.

#### 19 12.6.2.7 DECIMAL= specifier in a data transfer statement

20 1 The *scalar-default-char-expr* shall evaluate to COMMA or POINT. The DECIMAL= specifier temporarily changes  
21 (12.5.2) the decimal edit mode (13.6, 13.8.8, 12.5.6.7) for the connection. If the specifier is omitted, the mode is  
22 not changed.

#### 23 12.6.2.8 DELIM= specifier in a data transfer statement

24 1 The *scalar-default-char-expr* shall evaluate to APOSTROPHE, QUOTE, or NONE. The DELIM= specifier tem-  
25 porarily changes (12.5.2) the delimiter mode (13.10.4, 13.11.4.2, 12.5.6.8) for the connection. If the specifier is  
26 omitted, the mode is not changed.

#### 27 12.6.2.9 ID= specifier in a data transfer statement

28 1 Successful execution of an asynchronous data transfer statement containing an ID= specifier causes the variable  
29 specified in the ID= specifier to become defined with a processor determined value. If this value is zero, the  
30 data transfer operation has been completed. A nonzero value is referred to as the identifier of the data transfer  
31 operation. This identifier is different from the identifier of any other pending data transfer operation for this [unit](#).  
32 It can be used in a subsequent [WAIT](#) or [INQUIRE](#) statement to identify the particular data transfer operation.

33 2 If an error occurs during the execution of a data transfer statement containing an ID= specifier, the variable  
34 specified in the ID= specifier becomes undefined.

35 3 A child data transfer statement shall not specify the ID= specifier.

#### 1 12.6.2.10 PAD= specifier in a data transfer statement

2 1 The *scalar-default-char-expr* shall evaluate to YES or NO. The PAD= specifier temporarily changes (12.5.2) the  
3 pad mode (12.6.4.5.3, 12.5.6.13) for the connection. If the specifier is omitted, the mode is not changed.

#### 4 12.6.2.11 POS= specifier in a data transfer statement

5 1 The POS= specifier specifies the file position in [file storage units](#). This specifier may appear in a data transfer  
6 statement only if the statement specifies a [unit connected](#) for stream access. A child data transfer statement shall  
7 not specify this specifier.

8 2 A processor may prohibit the use of POS= with particular files that do not have the properties necessary to  
9 support random positioning. A processor may also prohibit positioning a particular file to any position prior to  
10 its current file position if the file does not have the properties necessary to support such positioning.

#### NOTE 12.29

A [unit](#) that is [connected](#) to a device or data stream might not be positionable.

11 3 If the file is [connected](#) for formatted stream access, the file position specified by POS= shall be equal to either 1  
12 (the beginning of the file) or a value previously returned by a [POS= specifier](#) in an [INQUIRE statement](#) for the  
13 file.

#### 14 12.6.2.12 REC= specifier in a data transfer statement

15 1 The REC= specifier specifies the number of the record that is to be read or written. This specifier may appear  
16 only in a data transfer statement that specifies a [unit connected](#) for direct access; it shall not appear in a child  
17 data transfer statement. If the *io-control-spec-list* contains a REC= specifier, the statement is a direct access  
18 data transfer statement. A child data transfer statement is a direct access data transfer statement if the parent  
19 is a direct access data transfer statement. Any other data transfer statement is a sequential access data transfer  
20 statement or a stream access data transfer statement, depending on whether the file connection is for sequential  
21 access or stream access.

#### 22 12.6.2.13 ROUND= specifier in a data transfer statement

23 1 The *scalar-default-char-expr* shall evaluate to UP, DOWN, ZERO, NEAREST, COMPATIBLE or PROCESSOR\_  
24 DEFINED. The ROUND= specifier temporarily changes (12.5.2) the input/output rounding mode (13.7.2.3.8,  
25 12.5.6.16) for the connection. If the specifier is omitted, the mode is not changed.

#### 26 12.6.2.14 SIGN= specifier in a data transfer statement

27 1 The *scalar-default-char-expr* shall evaluate to PLUS, SUPPRESS, or PROCESSOR\_DEFINED. The SIGN=  
28 specifier temporarily changes (12.5.2) the sign mode (13.8.4, 12.5.6.17) for the connection. If the specifier is  
29 omitted, the mode is not changed.

#### 30 12.6.2.15 SIZE= specifier in a data transfer statement

31 1 The SIZE= specifier in an [input statement](#) causes the variable specified to become [defined](#) with the count of  
32 the characters transferred from the file by data edit descriptors during the input operation. Blanks inserted as  
33 padding are not counted.

34 2 For a synchronous [input statement](#), this definition occurs when execution of the statement completes. For an  
35 asynchronous [input statement](#), this definition occurs when the corresponding wait operation is performed.

### 36 12.6.3 Data transfer input/output list

37 1 An input/output list specifies the entities whose values are transferred by a data transfer statement.

- 1 R1216 *input-item* is *variable*  
 2 or *io-implied-do*
- 3 R1217 *output-item* is *expr*  
 4 or *io-implied-do*
- 5 R1218 *io-implied-do* is ( *io-implied-do-object-list* , *io-implied-do-control* )
- 6 R1219 *io-implied-do-object* is *input-item*  
 7 or *output-item*
- 8 R1220 *io-implied-do-control* is *do-variable* = *scalar-int-expr* , ■  
 9 ■ *scalar-int-expr* [ , *scalar-int-expr* ]
- 10 C1231 (R1216) A variable that is an *input-item* shall not be a whole *assumed-size* array.
- 11 C1232 (R1219) In an *input-item-list*, an *io-implied-do-object* shall be an *input-item*. In an *output-item-list*, an  
 12 *io-implied-do-object* shall be an *output-item*.
- 13 C1233 (R1217) An expression that is an *output-item* shall not have a value that is a procedure pointer.
- 14 2 An *input-item* shall not appear as, nor be associated with, the *do-variable* of any *io-implied-do* that contains the  
 15 *input-item*.

**NOTE 12.30**

A constant, an expression involving operators or function references that does not have a pointer result, or an expression enclosed in parentheses shall not appear as an input list item.

- 16 3 If an input item is a pointer, it shall be associated with a *definable target* and data are transferred from the file to  
 17 the associated *target*. If an output item is a pointer, it shall be associated with a *target* and data are transferred  
 18 from the *target* to the file.

**NOTE 12.31**

Data transfers always involve the movement of values between a file and internal storage. A pointer as such cannot be read or written. Therefore, a pointer shall not appear as an item in an input/output list unless it is associated with a *target* that can receive a value (input) or can deliver a value (output).

- 19 4 If an input item or an output item is *allocatable*, it shall be allocated.
- 20 5 A list item shall not be *polymorphic* unless it is processed by a *defined input/output* procedure (12.6.4.8).
- 21 6 The *do-variable* of an *io-implied-do* that is in another *io-implied-do* shall not appear as, nor be associated with,  
 22 the *do-variable* of the containing *io-implied-do*.
- 23 7 The following rules describing whether to expand an input/output list item are re-applied to each expanded list  
 24 item until none of the rules apply.
- 25 • If an array appears as an input/output list item, it is treated as if the elements, if any, were specified in  
 26 array element order (9.5.3.2). However, no element of that array may affect the value of any expression in  
 27 the *input-item*, nor may any element appear more than once in a given *input-item*.

**NOTE 12.32**

For example:

```

INTEGER A (100), J (100)
...
READ *, A (A) ! Not allowed
READ *, A (LBOUND (A, 1) : UBOUND (A, 1)) ! Allowed
  
```

**NOTE 12.32 (cont.)**

READ *, A (J)	! Allowed if no two elements ! of J have the same value
A(1) = 1; A(10) = 10	
READ *, A (A (1) : A (10))	! Not allowed

- 1 • If a list item of derived type in an unformatted input/output statement is not processed by a [defined input/output](#) procedure (12.6.4.8), and if any subobject of that list item would be processed by a [defined input/output](#) procedure, the list item is treated as if all of the components of the object were specified in the list in [component order](#) (7.5.4.7); those components shall be accessible in the [scoping unit](#) containing the data transfer statement and shall not be pointers or [allocatable](#).
- 2
- 3
- 4
- 5
- 6 • An effective item of derived type in an unformatted input/output statement is treated as a single value in a processor-dependent form unless the list item or a subobject thereof is processed by a [defined input/output](#) procedure (12.6.4.8).
- 7
- 8

**NOTE 12.33**

The appearance of a derived-type object as an input/output list item in an unformatted input/output statement is not equivalent to the list of its components.

Unformatted input/output involving derived-type list items forms the single exception to the rule that the appearance of an aggregate list item (such as an array) is equivalent to the appearance of its expanded list of component parts. This exception permits the processor greater latitude in improving efficiency or in matching the processor-dependent sequence of values for a derived-type object to similar sequences for aggregate objects used by means other than Fortran. However, formatted input/output of all list items and unformatted input/output of list items other than those of derived types adhere to the above rule.

- 9 • If a list item of derived type in a formatted input/output statement is not processed by a [defined input/output](#) procedure, that list item is treated as if all of the components of the list item were specified in the list in [component order](#); those components shall be accessible in the [scoping unit](#) containing the input/output statement and shall not be pointers or [allocatable](#).
- 10
- 11
- 12
- 13 • If a derived-type list item is not processed by a [defined input/output](#) procedure and is not treated as a list of its individual components, all the subcomponents of that list item shall be accessible in the [scoping unit](#) containing the data transfer statement and shall not be pointers or [allocatable](#).
- 14
- 15
- 16 • For an *io-implied-do*, the loop initialization and execution are the same as for a [DO construct](#) (11.1.7.4).

**NOTE 12.34**

An example of an output list with an implied DO is:

```
WRITE (LP, FMT = '(10F8.2)') (LOG (A (I)), I = 1, N + 9, K), G
```

- 17 8 The scalar objects resulting when a data transfer statement's list items are expanded according to the rules in this subclause for handling array and derived-type list items are called [effective items](#). Zero-sized arrays and *io-implied-dos* with an iteration count of zero do not contribute to the list of [effective items](#). A scalar character item of zero length is an [effective item](#).
- 18
- 19
- 20

**NOTE 12.35**

In a formatted input/output statement, edit descriptors are associated with [effective items](#), which are always scalar. The rules in 12.6.3 determine the set of [effective items](#) corresponding to each actual list item in the statement. These rules might have to be applied repetitively until all of the [effective items](#) are scalar items.

- 21 9 An input/output list shall not contain an effective item of nondefault character kind if the data transfer statement specifies an [internal file](#) of default character kind. An input/output list shall not contain an effective item that is nondefault character except for [ISO 10646](#) or [ASCII character](#) if the data transfer statement specifies an [internal file](#) of [ISO 10646 character](#) kind. An input/output list shall not contain an effective item of type character of any kind other than [ASCII](#) if the data transfer statement specifies an [ASCII character internal file](#).
- 22
- 23
- 24
- 25

## 12.6.4 Execution of a data transfer input/output statement

### 12.6.4.1 Data transfer sequence of operations

1 Execution of a WRITE or PRINT statement for a unit connected to a file that does not exist creates the file unless an error condition occurs.

2 The effect of executing a synchronous data transfer statement shall be as if the following operations were performed in the order specified.

- (1) Determine the direction of data transfer (12.6.4.2).
- (2) Identify the *unit* (12.6.4.3).
- (3) Perform a wait operation for all pending input/output operations for the *unit*. If an error, end-of-file, or end-of-record condition occurs during any of the wait operations, steps 4 through 8 are skipped.
- (4) Establish the format if one is specified.
- (5) If the statement is not a child data transfer statement (12.6.4.8),
  - (a) position the file prior to data transfer (12.3.4.3), and
  - (b) for formatted data transfer, set the left tab limit (13.8.1.2).
- (6) Transfer data between the file and the entities specified by the input/output list (if any) or namelist, possibly mediated by *defined input/output* procedures (12.6.4.8).
- (7) Determine whether an error, end-of-file, or end-of-record condition has occurred.
- (8) Position the file after data transfer (12.3.4.4) unless the statement is a child data transfer statement (12.6.4.8).
- (9) Cause any variable specified in a SIZE= specifier to become defined.
- (10) If an error, end-of-file, or end-of-record condition occurred, processing continues as specified in 12.11; otherwise, any variable specified in an IOSTAT= specifier is assigned the value zero.

3 The effect of executing an asynchronous data transfer statement shall be as if the following operations were performed in the order specified.

- (1) Determine the direction of data transfer (12.6.4.2).
- (2) Identify the *unit* (12.6.4.3).
- (3) Optionally, perform wait operations for one or more pending input/output operations for the *unit*. If an error, end-of-file, or end-of-record condition occurs during any of the wait operations, steps 4 through 9 are skipped.
- (4) Establish the format if one is specified.
- (5) Position the file prior to data transfer (12.3.4.3) and, for formatted data transfer, set the left tab limit (13.8.1.2).
- (6) Establish the set of *storage units* identified by the input/output list. For an *input statement*, this might require some or all of the data in the file to be read if an input variable is used as a *scalar-int-expr* in an *io-implied-do-control* in the input/output list, as a *subscript*, *substring-range*, *stride*, or is otherwise referenced.
- (7) Initiate an asynchronous data transfer between the file and the entities specified by the input/output list (if any) or namelist. The asynchronous data transfer may complete (and an error, end-of-file, or end-of-record condition may occur) during the execution of this data transfer statement or during a later wait operation.
- (8) Determine whether an error, end-of-file, or end-of-record condition has occurred. The conditions may occur during the execution of this data transfer statement or during the corresponding wait operation, but not both.
- (9) Position the file as if the data transfer had finished (12.3.4.4).
- (10) Cause any variable specified in a SIZE= specifier to become undefined.
- (11) If an error, end-of-file, or end-of-record condition occurred, processing continues as specified in 12.11; otherwise, any variable specified in an IOSTAT= specifier is assigned the value zero.

- 1 4 For an asynchronous data transfer statement, the data transfers may occur during execution of the statement,  
2 during execution of the corresponding wait operation, or anywhere between. The data transfer operation is  
3 considered to be pending until a corresponding wait operation is performed.
- 4 5 For asynchronous output, a pending input/output storage sequence affector (12.6.2.5) shall not be redefined,  
5 become undefined, or have its [pointer association](#) status changed.
- 6 6 For asynchronous input, a pending input/output storage sequence affector shall not be referenced, become defined,  
7 become undefined, become associated with a [dummy argument](#) that has the [VALUE attribute](#), or have its [pointer  
8 association](#) status changed.
- 9 7 Error, end-of-file, and end-of-record conditions in an asynchronous data transfer operation may occur during  
10 execution of either the data transfer statement or the corresponding wait operation. If an ID= specifier does not  
11 appear in the initiating data transfer statement, the conditions may occur during the execution of any subsequent  
12 data transfer or wait operation for the same [unit](#). When a condition occurs for a previously executed asynchronous  
13 data transfer statement, a wait operation is performed for all pending data transfer operations on that [unit](#). When  
14 a condition occurs during a subsequent statement, any actions specified by IOSTAT=, IOMSG=, ERR=, END=,  
15 and EOR= specifiers for that statement are taken.
- 16 8 If execution of the program is terminated during execution of an [output statement](#), the contents of the file become  
17 undefined.

**NOTE 12.36**

Because end-of-file and error conditions for asynchronous data transfer statements without an ID= specifier can be reported by the processor during the execution of a subsequent data transfer statement, it might be impossible for the user to determine which data transfer statement caused the condition. Reliably detecting which [input statement](#) caused an end-of-file condition requires that all asynchronous [input statements](#) for the [unit](#) include an ID= specifier.

**12.6.4.2 Direction of data transfer**

- 18  
19 1 Execution of a READ statement causes values to be transferred from a file to the entities specified by the input  
20 list, if any, or specified within the file itself for namelist input. Execution of a WRITE or PRINT statement  
21 causes values to be transferred to a file from the entities specified by the output list and format specification, if  
22 any, or by the *namelist-group-name* for namelist output.

**12.6.4.3 Identifying a unit**

- 23  
24 1 A data transfer statement that contains an input/output control list includes a UNIT= specifier that identifies  
25 an [external](#) or [internal](#) unit. A READ statement that does not contain an input/output control list specifies a  
26 particular processor-dependent [unit](#), which is the same as the [unit](#) identified by \* in a [READ statement](#) that  
27 contains an input/output control list (12.5.1) and is the same as the [unit](#) identified by the value of the [named  
28 constant](#) INPUT\_UNIT of the intrinsic module ISO\_FORTRAN\_ENV (16.10.2.13). The [PRINT statement](#)  
29 specifies some other processor-dependent [unit](#), which is the same as the [unit](#) identified by \* in a [WRITE statement](#)  
30 and is the same as the [unit](#) identified by the value of the [named constant](#) OUTPUT\_UNIT of the intrinsic module  
31 ISO\_FORTRAN\_ENV (16.10.2.22). Thus, each data transfer statement identifies an [external](#) or [internal](#) unit.
- 32 2 The [unit](#) identified by an unformatted data transfer statement shall be an [external unit](#).
- 33 3 The [unit](#) identified by a data transfer statement shall be [connected](#) to a file when execution of the statement  
34 begins.

**NOTE 12.37**

The [unit](#) could be [preconnected](#).

#### 1 12.6.4.4 Establishing a format

- 2 1 If the input/output control list contains \* as a format, list-directed formatting is established. If *namelist-group-*  
3 *name* appears, namelist formatting is established. If no *format* or *namelist-group-name* is specified, unformatted  
4 data transfer is established. Otherwise, the format specified by *format* is established.
- 5 2 For output to an *internal file*, a format specification that is in the file or is associated with the file shall not be  
6 specified.
- 7 3 An input list item, or an entity associated with it, shall not contain any portion of an established format spe-  
8 cification.

#### 9 12.6.4.5 Data transfer

##### 10 12.6.4.5.1 General

- 11 1 Data are transferred between the file and the entities specified by the input/output list or namelist. The list items  
12 are processed in the order of the input/output list for all data transfer statements except namelist data transfer  
13 statements. The list items for a namelist *input statement* are processed in the order of the entities specified  
14 within the input records. The list items for a namelist *output statement* are processed in the order in which the  
15 variables are specified in the *namelist-group-object-list*. *Effective items* are derived from the input/output list  
16 items as described in 12.6.3.
- 17 2 All values needed to determine which entities are specified by an input/output list item are determined at the  
18 beginning of the processing of that item.
- 19 3 All values are transmitted to or from the entities specified by a list item prior to the processing of any succeeding  
20 list item for all data transfer statements.

#### NOTE 12.38

In the example

```
READ (N) N, X (N)
```

the old value of N identifies the *unit*, but the new value of N is the subscript of X.

- 21 4 All values following the *name=* part of the namelist entity (13.11) within the input records are transmitted to  
22 the matching entity specified in the *namelist-group-object-list* prior to processing any succeeding entity within  
23 the input record for namelist *input statements*. If an entity is specified more than once within the input record  
24 during a namelist *input statement*, the last occurrence of the entity specifies the value or values to be used for  
25 that entity.
- 26 5 If the input/output item is a pointer, data are transferred between the file and the associated *target*.
- 27 6 If an *internal file* has been specified, an input/output list item shall not be in the file or associated with the file.
- 28 7 During the execution of an *output statement* that specifies an *internal file*, no part of that *internal file* shall be  
29 referenced, defined, or become undefined as the result of evaluating any output list item.
- 30 8 During the execution of an *input statement* that specifies an *internal file*, no part of that *internal file* shall be  
31 defined or become undefined as the result of transferring a value to any input list item.
- 32 9 A DO variable becomes defined and its iteration count established at the beginning of processing of the *io-implied-*  
33 *do-object-list* an *io-implied-do*.
- 34 10 On output, every entity whose value is to be transferred shall be defined.

#### 12.6.4.5.2 Unformatted data transfer

- 1 If the file is not **connected** for unformatted input/output, unformatted data transfer is prohibited.
- 2 During unformatted data transfer, data are transferred without editing between the file and the entities specified by the input/output list. If the file is **connected** for sequential or direct access, exactly one record is read or written.
- 3 A value in the file is stored in a **contiguous** sequence of **file storage units**, beginning with the **file storage unit** immediately following the current file position.
- 4 After each value is transferred, the current file position is moved to a point immediately after the last **file storage unit** of the value.
- 5 On input from a file **connected** for sequential or direct access, the number of **file storage units** required by the input list shall be less than or equal to the number of **file storage units** in the record.
- 6 On input, if the **file storage units** transferred do not contain a value with the same type and type parameters as the input list entity, then the resulting value of the entity is processor dependent except in the following cases.
  - A complex entity may correspond to two real values with the same **kind type parameter** as the complex entity.
  - A default character list entity of length *n* may correspond to *n* default characters stored in the file, regardless of the length parameters of the entities that were written to these **storage units** of the file. If the file is **connected** for stream input, the characters may have been written by formatted stream output.
- 7 On output to a file **connected** for unformatted direct access, the output list shall not specify more values than can fit into the record. If the file is **connected** for direct access and the values specified by the output list do not fill the record, the remainder of the record is undefined.
- 8 If the file is **connected** for unformatted sequential access, the record is created with a length sufficient to hold the values from the output list. This length shall be one of the set of allowed record lengths for the file and shall not exceed the value specified in the **RECL= specifier**, if any, of the **OPEN statement** that established the connection.

#### 12.6.4.5.3 Formatted data transfer

- 1 If the file is not **connected** for formatted input/output, formatted data transfer is prohibited.
- 2 During formatted data transfer, data are transferred with editing between the file and the entities specified by the input/output list or by the *namelist-group-name*. Format control is initiated and editing is performed as described in Clause 13.
- 3 The current record and possibly additional records are read or written.
- 4 During advancing input when the pad mode has the value NO, the input list and format specification shall not require more characters from the record than the record contains.
- 5 During advancing input when the pad mode has the value YES, blank characters are supplied by the processor if the input list and format specification require more characters from the record than the record contains.
- 6 During nonadvancing input when the pad mode has the value NO, an end-of-record condition (12.11) occurs if the input list and format specification require more characters from the record than the record contains, and the record is complete (12.3.3.4). If the record is incomplete, an end-of-file condition occurs instead of an end-of-record condition.
- 7 During nonadvancing input when the pad mode has the value YES, blank characters are supplied by the processor if an **effective item** and its corresponding data edit descriptors require more characters from the record than the record contains. If the record is incomplete, an end-of-file condition occurs; otherwise, an end-of-record condition



1 occurs.

2 8 If the file is [connected](#) for direct access, the record number is increased by one as each succeeding record is read  
3 or written.

4 9 On output, if the file is [connected](#) for direct access or is an internal file and the characters specified by the output  
5 list and format do not fill a record, blank characters are added to fill the record.

6 10 On output, the output list and format specification shall not specify more characters for a record than have been  
7 specified by a [RECL= specifier](#) in the [OPEN statement](#) or the record length of an [internal file](#).

#### 8 **12.6.4.6 List-directed formatting**

9 1 If list-directed formatting has been established, editing is performed as described in [13.10](#).

#### 10 **12.6.4.7 Namelist formatting**

11 1 If namelist formatting has been established, editing is performed as described in [13.11](#).

12 2 Every [allocatable namelist-group-object](#) in the namelist group shall be allocated and every [namelist-group-object](#)  
13 that is a pointer shall be associated with a [target](#). If a [namelist-group-object](#) is [polymorphic](#) or has an [ultimate](#)  
14 [component](#) that is [allocatable](#) or a pointer, that object shall be processed by a [defined input/output](#) procedure  
15 ([12.6.4.8](#)).

#### 16 **12.6.4.8 Defined input/output**

##### 17 **12.6.4.8.1 General**

18 1 [Defined input/output](#) allows a program to override the default handling of derived-type objects and values in  
19 data transfer statements described in [12.6.3](#).

20 2 A [defined input/output](#) procedure is a procedure accessible by a [defined-io-generic-spec](#) ([15.4.3.2](#)). A particular  
21 [defined input/output](#) procedure is selected as described in [12.6.4.8.4](#).

##### 22 **12.6.4.8.2 Defined input/output procedures**

23 1 For a particular derived type and a particular set of [kind type parameter](#) values, there are four possible sets of  
24 [characteristics](#) for [defined input/output](#) procedures; one each for formatted input, formatted output, unformatted  
25 input, and unformatted output. The program need not supply all four procedures. The procedures are specified  
26 to be used for derived-type input/output by [interface blocks](#) ([15.4.3.2](#)) or by generic [bindings](#) ([7.5.5](#)), with a  
27 [defined-io-generic-spec](#) (R1509). The [defined-io-generic-specs](#) for these procedures are [READ \(FORMATTED\)](#),  
28 [READ \(UNFORMATTED\)](#), [WRITE \(FORMATTED\)](#), and [WRITE \(UNFORMATTED\)](#), for formatted input,  
29 unformatted input, formatted output, and unformatted output respectively.

30 2 In the four [interfaces](#), which specify the [characteristics](#) of [defined input/output](#) procedures, the following syntax  
31 term is used:

```
32 R1221 dtv-type-spec                is TYPE( derived-type-spec )
33                                     or CLASS( derived-type-spec )
```

34 C1234 (R1221) If [derived-type-spec](#) specifies an [extensible type](#), the [CLASS](#) keyword shall be used; otherwise,  
35 the TYPE keyword shall be used.

36 C1235 (R1221) All length type parameters of [derived-type-spec](#) shall be assumed.

37 3 If the [defined-io-generic-spec](#) is [READ \(FORMATTED\)](#), the [characteristics](#) shall be the same as those specified  
38 by the following [interface](#):

```
39 4          SUBROUTINE my_read_routine_formatted                &
```

```

1          (dtv,          &
2          unit,         &
3          iotype, v_list, &
4          iostat, iomsg)
5          ! the derived-type variable
6          dtv-type-spec, INTENT(INOUT) :: dtv
7          INTEGER, INTENT(IN) :: unit ! unit number
8          ! the edit descriptor string
9          CHARACTER (LEN=*), INTENT(IN) :: iotype
10         INTEGER, INTENT(IN) :: v_list(:)
11         INTEGER, INTENT(OUT) :: iostat
12         CHARACTER (LEN=*), INTENT(INOUT) :: iomsg
13     END

```

14 5 If the *defined-io-generic-spec* is **READ (UNFORMATTED)**, the *characteristics* shall be the same as those specified  
15 by the following *interface*:

```

16 6          SUBROUTINE my_read_routine_unformatted          &
17          (dtv,          &
18          unit,         &
19          iostat, iomsg)
20          ! the derived-type variable
21          dtv-type-spec, INTENT(INOUT) :: dtv
22          INTEGER, INTENT(IN) :: unit
23          INTEGER, INTENT(OUT) :: iostat
24          CHARACTER (LEN=*), INTENT(INOUT) :: iomsg
25     END

```

26 7 If the *defined-io-generic-spec* is **WRITE (FORMATTED)**, the *characteristics* shall be the same as those specified  
27 by the following *interface*:

```

28 8          SUBROUTINE my_write_routine_formatted          &
29          (dtv,          &
30          unit,         &
31          iotype, v_list, &
32          iostat, iomsg)
33          ! the derived-type value/variable
34          dtv-type-spec, INTENT(IN) :: dtv
35          INTEGER, INTENT(IN) :: unit
36          ! the edit descriptor string
37          CHARACTER (LEN=*), INTENT(IN) :: iotype
38          INTEGER, INTENT(IN) :: v_list(:)
39          INTEGER, INTENT(OUT) :: iostat
40          CHARACTER (LEN=*), INTENT(INOUT) :: iomsg
41     END

```

42 9 If the *defined-io-generic-spec* is **WRITE (UNFORMATTED)**, the *characteristics* shall be the same as those  
43 specified by the following *interface*:

```

44 10         SUBROUTINE my_write_routine_unformatted          &
45         (dtv,          &
46         unit,         &
47         iostat, iomsg)
48         ! the derived-type value/variable
49         dtv-type-spec, INTENT(IN) :: dtv

```

```

1      INTEGER, INTENT(IN) :: unit
2      INTEGER, INTENT(OUT) :: iostat
3      CHARACTER (LEN=*), INTENT(INOUT) :: iomsg
4      END

```

5 11 The actual specific procedure names (the `my_..._routine_...` procedure names above) are not significant. In  
6 the discussion here and elsewhere, the [dummy arguments](#) in these [interfaces](#) are referred to by the names given  
7 above; the names are, however, arbitrary.

### 8 12.6.4.8.3 Executing defined input/output data transfers

9 1 If a [defined input/output](#) procedure is selected for an effective item as specified in 12.6.4.8.4, the processor shall  
10 call the selected [defined input/output](#) procedure for that item. The [defined input/output](#) procedure controls the  
11 actual data transfer operations for the derived-type list item.

12 2 A data transfer statement that includes a derived-type list item and that causes a [defined input/output](#) procedure  
13 to be invoked is called a parent data transfer statement. A data transfer statement that is executed while a parent  
14 data transfer statement is being processed and that specifies the [unit](#) passed into a [defined input/output](#) procedure  
15 is called a child data transfer statement.

#### NOTE 12.39

A [defined input/output](#) procedure will usually contain child data transfer statements that read values from or write values to the current record or at the current file position. The effect of executing the [defined input/output](#) procedure is similar to that of substituting the list items from any child data transfer statements into the parent data transfer statement's list items, along with similar substitutions in the format specification.

#### NOTE 12.40

A particular execution of a READ, WRITE or PRINT statement can be both a parent and a child data transfer statement. A [defined input/output](#) procedure can indirectly call itself or another [defined input/output](#) procedure by executing a child data transfer statement containing a list item of derived type, where a matching [interface](#) is accessible for that derived type. If a [defined input/output](#) procedure calls itself indirectly in this manner, it cannot be declared NON\_RECURSIVE.

16 3 A child data transfer statement is processed differently from a nonchild data transfer statement in the following  
17 ways.

- 18 • Executing a child data transfer statement does not position the file prior to data transfer.
- 19 • An unformatted child data transfer statement does not position the file after data transfer is complete.
- 20 • Any [ADVANCE= specifier](#) in a child input/output statement is ignored.

21 4 When a [defined input/output](#) procedure is invoked, the processor shall pass a [unit](#) argument that has a value as  
22 follows.

- 23 • If the parent data transfer statement uses a [file-unit-number](#), the value of the [unit](#) argument shall be that  
24 of the [file-unit-number](#).
- 25 • If the parent data transfer statement is a [WRITE statement](#) with an asterisk [unit](#) or a [PRINT statement](#),  
26 the [unit](#) argument shall have the same value as the [named constant OUTPUT\\_UNIT](#) of the intrinsic  
27 module [ISO\\_FORTRAN\\_ENV](#) (16.10.2).
- 28 • If the parent data transfer statement is a [READ statement](#) with an asterisk [unit](#) or a [READ statement](#)  
29 without an [io-control-spec-list](#), the [unit](#) argument shall have the same value as the [INPUT\\_UNIT](#) named  
30 constant of the intrinsic module [ISO\\_FORTRAN\\_ENV](#) (16.10.2).
- 31 • Otherwise the parent data transfer statement accesses an [internal file](#), in which case the [unit](#) argument  
32 shall have a processor-dependent negative value.

**NOTE 12.41**

The `unit` argument passed to a [defined input/output](#) procedure will be negative when the parent data transfer statement specified an [internal unit](#), or specified an [external unit](#) that is a `NEWUNIT` value. When an [internal unit](#) is used with the `INQUIRE` statement, an error condition will occur, and any variable specified in an `IOSTAT= specifier` will be assigned the value `IOSTAT_INQUIRE_INTERNAL_UNIT` from the intrinsic module `ISO_FORTRAN_ENV` (16.10.2).

- 1 5 For formatted data transfer, the processor shall pass an `iotype` argument that has the value
- 2     • “LISTDIRECTED” if the parent data transfer statement specified list directed formatting,
- 3     • “NAMELIST” if the parent data transfer statement specified namelist formatting, or
- 4     • “DT” concatenated with the *char-literal-constant*, if any, of the DT edit descriptor in the format specification
- 5     of the parent data transfer statement.
- 6 6 If the parent data transfer statement is an [input statement](#), the `dtv dummy argument` is argument associated
- 7 with the [effective item](#) that caused the [defined input](#) procedure to be invoked, as if the [effective item](#) were an
- 8 [actual argument](#) in this [procedure reference](#) (5.4.5).
- 9 7 If the parent data transfer statement is an [output statement](#), the processor shall provide the value of the [effective](#)
- 10 [item](#) in the `dtv dummy argument`.
- 11 8 If the *v-list* of the edit descriptor appears in the parent data transfer statement, the processor shall provide the
- 12 values from it in the `v_list dummy argument`, with the same number of elements in the same order as *v-list*.
- 13 If there is no *v-list* in the edit descriptor or if the data transfer statement specifies list-directed or namelist
- 14 formatting, the processor shall provide `v_list` as a zero-sized array.

**NOTE 12.42**

The user’s procedure might choose to interpret an element of the `v_list` argument as a field width, but this is not required. If it does, it would be appropriate to fill an output field with “\*”s if the width is too small.

- 15 9 The `iostat` argument is used to report whether an error, end-of-record, or end-of-file condition (12.11) occurs.
- 16 If an error condition occurs, the [defined input/output](#) procedure shall assign a positive value to the `iostat`
- 17 argument. Otherwise, if an end-of-file condition occurs, the [defined input](#) procedure shall assign the value of the
- 18 [named constant IOSTAT\\_END](#) (16.10.2.16) to the `iostat` argument. Otherwise, if an end-of-record condition
- 19 occurs, the [defined input](#) procedure shall assign the value of the [named constant IOSTAT\\_EOR](#) (16.10.2.17) to
- 20 `iostat`. Otherwise, the [defined input/output](#) procedure shall assign the value zero to the `iostat` argument.
- 21 10 If the [defined input/output](#) procedure returns a nonzero value for the `iostat` argument, the procedure shall also
- 22 return an explanatory message in the `iomsg` argument. Otherwise, the procedure shall not change the value of
- 23 the `iomsg` argument.

**NOTE 12.43**

The values of the `iostat` and `iomsg` arguments set in a [defined input/output](#) procedure need not be passed to all of the parent data transfer statements.

- 24 11 If the `iostat` argument of the [defined input/output](#) procedure has a nonzero value when that procedure returns,
- 25 and the processor therefore [terminates execution](#) of the program as described in 12.11, the processor shall make
- 26 the value of the `iomsg` argument available in a processor-dependent manner.
- 27 12 While a parent `READ` statement is active, an input/output statement shall not read from any [external unit](#) other
- 28 than the one specified by the `unit dummy argument` and shall not perform output to any [external unit](#).
- 29 13 While a parent `WRITE` or `PRINT` statement is active, an input/output statement shall not perform output to
- 30 any [external unit](#) other than the one specified by the `unit dummy argument` and shall not read from any [external](#)
- 31 [unit](#).

- 1 14 While a parent data transfer statement is active, a data transfer statement that specifies an **internal file** is  
2 permitted.
- 3 15 **OPEN**, **CLOSE**, **BACKSPACE**, **ENDFILE**, and **REWIND** statements shall not be executed while a parent data  
4 transfer statement is active.
- 5 16 A **defined input/output** procedure may use a format specification with a **DT edit descriptor** for handling a  
6 component of the derived type that is itself of a derived type. A child data transfer statement that is a list  
7 directed or namelist input/output statement may contain a list item of derived type.
- 8 17 Because a child data transfer statement does not position the file prior to data transfer, the child data transfer  
9 statement starts transferring data from where the file was positioned by the parent data transfer statement's  
10 most recently processed **effective item** or edit descriptor. This is not necessarily at the beginning of a record.
- 11 18 The edit descriptors **T** and **TL** used on **unit** by a child data transfer statement shall not cause the file to be  
12 positioned before the file position at the time the **defined input/output** procedure was invoked.

**NOTE 12.44**

A **defined input/output** procedure could use **INQUIRE** to determine the settings of **BLANK=**, **PAD=**, **ROUND=**, **DECIMAL=**, and **DELIM=** for an **external unit**. The **INQUIRE statement** provides values as specified in 12.10.

- 13 19 Neither a parent nor child data transfer statement shall be asynchronous.
- 14 20 A **defined input/output** procedure, and any procedures invoked therefrom, shall not define, nor cause to become  
15 undefined, any **storage unit** referenced by any input/output list item, the corresponding format, or any specifier  
16 in any active parent data transfer statement, except through the **dtv** argument.

**NOTE 12.45**

A data transfer statement with an **ID=**, **POS=**, or **REC=** specifier cannot be a child data transfer statement in a standard-conforming program.

**NOTE 12.46**

A simple example of derived type formatted output follows. The derived type variable **chairman** has two components. The type and an associated write formatted procedure are defined in a module so as to be accessible from wherever they might be needed. It would also be possible to check that **iotype** indeed has the value 'DT' and to set **iostat** and **iomsg** accordingly.

```

MODULE p

  TYPE :: person
    CHARACTER (LEN=20) :: name
    INTEGER :: age
  CONTAINS
    PROCEDURE,PRIVATE :: pwf
    GENERIC          :: WRITE(FORMATTED) => pwf
  END TYPE person

CONTAINS

  SUBROUTINE pwf (dtv,unit,iotype,vlist,iostat,iomsg)
! argument definitions
    CLASS(person), INTENT(IN) :: dtv
    INTEGER, INTENT(IN) :: unit
    CHARACTER (LEN=*), INTENT(IN) :: iotype

```

## NOTE 12.46 (cont.)

```

    INTEGER, INTENT(IN) :: vlist(:)
    INTEGER, INTENT(OUT) :: iostat
    CHARACTER (LEN=*), INTENT(INOUT) :: iomsg
! local variable
    CHARACTER (LEN=9) :: pfmt

! vlist(1) and (2) are to be used as the field widths of the two
! components of the derived type variable. First set up the format to
! be used for output.
    WRITE(pfmt,'(A,I2,A,I2,A)' ) '(A', vlist(1), ',I', vlist(2), ' )'

! now the basic output statement
    WRITE(unit, FMT=pfmt, IOSTAT=iostat) dtv%name, dtv%age

END SUBROUTINE pwf

END MODULE p

PROGRAM committee
    USE p
    INTEGER id, members
    TYPE (person) :: chairman
    ...
    WRITE(6, FMT="(I2, DT (15,6), I5)" ) id, chairman, members
! this writes a record with four fields, with lengths 2, 15, 6, 5
! respectively

END PROGRAM

```

## NOTE 12.47

In the following example, the variables of the derived type `node` form a linked list, with a single value at each node. The subroutine `pwf` is used to write the values in the list, one per line.

```

MODULE p

    TYPE node
        INTEGER :: value = 0
        TYPE (NODE), POINTER :: next_node => NULL ( )
    CONTAINS
        PROCEDURE,PRIVATE :: pwf
        GENERIC          :: WRITE(FORMATTED) => pwf
    END TYPE node

CONTAINS

    SUBROUTINE pwf (dtv,unit,iotype,vlist,iostat,iomsg)
! Write the chain of values, each on a separate line in I9 format.
        CLASS(node), INTENT(IN) :: dtv
        INTEGER, INTENT(IN) :: unit
        CHARACTER (LEN=*), INTENT(IN) :: iotype
        INTEGER, INTENT(IN) :: vlist(:)
        INTEGER, INTENT(OUT) :: iostat
        CHARACTER (LEN=*), INTENT(INOUT) :: iomsg

```

## NOTE 12.47 (cont.)

```

WRITE(unit,'(i9 /)', IOSTAT = iostat) dtv%value
IF(iostat/=0) RETURN
IF(ASSOCIATED(dtv%next_node)) WRITE(unit,'(dt)', IOSTAT=iostat) dtv%next_node
END SUBROUTINE pwf

END MODULE p

```

1 **12.6.4.8.4 Resolving defined input/output procedure references**

- 2 1 A suitable [generic interface](#) for [defined input/output](#) of an [effective item](#) is one that has a [defined-io-generic-spec](#)  
3 that is appropriate to the direction (read or write) and form (formatted or unformatted) of the data transfer  
4 as specified in [12.6.4.8.2](#), and has a [specific interface](#) whose `dtv` argument is compatible with the [effective item](#)  
5 according to the rules for argument association in [15.5.2.4](#).
- 6 2 When an [effective item](#) ([12.6.3](#)) that is of derived type is encountered during a data transfer, [defined input/output](#)  
7 occurs if both of the following conditions are true.
- 8 (1) The circumstances of the input/output are such that [defined input/output](#) is permitted; that is,  
9 either
- 10 (a) the transfer was initiated by a list-directed, namelist, or unformatted input/output statement,  
11 or
- 12 (b) a format specification is supplied for the data transfer statement, and the edit descriptor  
13 corresponding to the [effective item](#) is a DT edit descriptor.
- 14 (2) A suitable [defined input/output](#) procedure is available; that is, either
- 15 (a) the [declared type](#) of the [effective item](#) has a suitable generic [type-bound procedure](#), or  
16 (b) a suitable [generic interface](#) is accessible.
- 17 3 If (2a) is true, the procedure referenced is determined as for explicit [type-bound procedure](#) references ([15.5](#)); that  
18 is, the [binding](#) with the appropriate [specific interface](#) is located in the [declared type](#) of the [effective item](#), and the  
19 corresponding [binding](#) in the [dynamic type](#) of the [effective item](#) is selected.
- 20 4 If (2a) is false and (2b) is true, the reference is to the procedure identified by the appropriate [specific interface](#)  
21 in the [interface block](#).

22 **12.6.5 Termination of data transfer statements**

- 23 1 Termination of a data transfer statement occurs when
- 24 • format processing encounters a colon or data edit descriptor and there are no remaining elements in the  
25 [input-item-list](#) or [output-item-list](#),
  - 26 • unformatted or list-directed data transfer exhausts the [input-item-list](#) or [output-item-list](#),
  - 27 • namelist output exhausts the [namelist-group-object-list](#),
  - 28 • an error condition occurs,
  - 29 • an end-of-file condition occurs,
  - 30 • a slash (/) is encountered as a value separator ([13.10](#), [13.11](#)) in the record being read during list-directed  
31 or namelist input, or
  - 32 • an end-of-record condition occurs during execution of a nonadvancing input statement ([12.11](#)).

## 12.7 Waiting on pending data transfer

### 12.7.1 Wait operation

- 1 Execution of an asynchronous [data transfer statement](#) in which neither an error, end-of-record, nor end-of-file condition occurs initiates a pending data transfer operation. There may be multiple pending data transfer operations for the same or multiple [units](#) simultaneously. A pending data transfer operation remains pending until a corresponding wait operation is performed. A wait operation may be performed by a [BACKSPACE](#), [CLOSE](#), [ENDFILE](#), [FLUSH](#), [INQUIRE](#), [PRINT](#), [READ](#), [REWIND](#), [WAIT](#), or [WRITE](#) statement.
- 2 A wait operation completes the processing of a pending data transfer operation. Each wait operation completes only a single data transfer operation, although a single statement may perform multiple wait operations.
- 3 If the actual data transfer is not yet complete, the wait operation first waits for its completion. If the data transfer operation is an input operation that completed without error, the [storage units](#) of the [input/output storage sequence](#) then become defined with the values as described in [12.6.2.15](#) and [12.6.4.5](#).
- 4 If any error, end-of-file, or end-of-record conditions occur, the applicable actions specified by the [IOSTAT=](#), [IOMSG=](#), [ERR=](#), [END=](#), and [EOR=](#) specifiers of the statement that performs the wait operation are taken.
- 5 If an error or end-of-file condition occurs during a wait operation for a [unit](#), the processor performs a wait operation for all pending data transfer operations for that [unit](#).

#### NOTE 12.48

Error, end-of-file, and end-of-record conditions can be raised either during the [data transfer statement](#) that initiates asynchronous input/output, a subsequent asynchronous [data transfer statement](#) for the same [unit](#), or during the wait operation. If such conditions are raised during a [data transfer statement](#), they trigger actions according to the [IOSTAT=](#), [ERR=](#), [END=](#), and [EOR=](#) specifiers of that statement; if they are raised during the wait operation, the actions are in accordance with the specifiers of the statement that performs the wait operation.

- 6 After completion of the wait operation, the data transfer operation and its input/output storage sequence are no longer considered to be pending.

### 12.7.2 WAIT statement

- 1 A WAIT statement performs a wait operation for specified pending asynchronous data transfer operations.

R1222 *wait-stmt*                                    **is** WAIT (*wait-spec-list*)

R1223 *wait-spec*                                    **is** [ UNIT = ] *file-unit-number*  
**or** END = *label*  
**or** EOR = *label*  
**or** ERR = *label*  
**or** ID = *scalar-int-expr*  
**or** IOMSG = *iomsg-variable*  
**or** IOSTAT = *scalar-int-variable*

C1236 No specifier shall appear more than once in a given *wait-spec-list*.

C1237 A *file-unit-number* shall be specified in a *wait-spec-list*; if the optional characters UNIT= are omitted, the *file-unit-number* shall be the first item in the *wait-spec-list*.

C1238 (R1223) The *label* in the [ERR=](#), [EOR=](#), or [END=](#) specifier shall be the statement label of a [branch target statement](#) that appears in the same [inclusive scope](#) as the WAIT statement.

- 2 The [IOSTAT=](#), [ERR=](#), [EOR=](#), [END=](#), and [IOMSG=](#) specifiers are described in [12.11](#).



- 1 3 The value of the expression specified in the ID= specifier shall be zero or the identifier of a pending data transfer  
 2 operation for the specified **unit**. If the ID= specifier appears, a wait operation for the specified data transfer  
 3 operation, if any, is performed. If the ID= specifier is omitted, wait operations for all pending data transfers for  
 4 the specified **unit** are performed.
- 5 4 Execution of a WAIT statement specifying a **unit** that does not exist, has no file connected to it, or is not open  
 6 for asynchronous input/output is permitted, provided that the WAIT statement has no ID= specifier; such a  
 7 WAIT statement does not cause an error or end-of-file condition to occur.

**NOTE 12.49**

An **EOR= specifier** has no effect if the pending data transfer operation is not a nonadvancing read. An **END= specifier** has no effect if the pending data transfer operation is not a READ.

## 8 12.8 File positioning statements

### 9 12.8.1 Syntax

10 R1224 *backspace-stmt* is BACKSPACE *file-unit-number*  
 11 or BACKSPACE ( *position-spec-list* )

12 R1225 *endfile-stmt* is ENDFILE *file-unit-number*  
 13 or ENDFILE ( *position-spec-list* )

14 R1226 *rewind-stmt* is REWIND *file-unit-number*  
 15 or REWIND ( *position-spec-list* )

- 16 1 A **unit** that is connected for direct access shall not be referred to by a BACKSPACE, ENDFILE, or REWIND  
 17 statement. A **unit** that is connected for unformatted stream access shall not be referred to by a BACKSPACE  
 18 statement. A **unit** that is connected with an ACTION= specifier having the value READ shall not be referred  
 19 to by an ENDFILE statement.

20 R1227 *position-spec* is [ UNIT = ] *file-unit-number*  
 21 or IOMSG = *iomsg-variable*  
 22 or IOSTAT = *scalar-int-variable*  
 23 or ERR = *label*

24 C1239 No specifier shall appear more than once in a given *position-spec-list*.

25 C1240 A *file-unit-number* shall be specified in a *position-spec-list*; if the optional characters UNIT= are omitted,  
 26 the *file-unit-number* shall be the first item in the *position-spec-list*.

27 C1241 (R1227) The *label* in the **ERR= specifier** shall be the statement label of a **branch target statement** that  
 28 appears in the same **inclusive scope** as the file positioning statement.

29 2 The **IOSTAT=**, **ERR=**, and **IOMSG=** specifiers are described in 12.11.

30 3 Execution of a file positioning statement performs a wait operation for all pending asynchronous data transfer  
 31 operations for the specified **unit**.

### 32 12.8.2 BACKSPACE statement

- 33 1 Execution of a BACKSPACE statement causes the file connected to the specified **unit** to be positioned before  
 34 the current record if there is a current record, or before the preceding record if there is no current record. If the  
 35 file is at its initial point, the position of the file is not changed.

**NOTE 12.50**

If the preceding record is an endfile record, the file is positioned before the endfile record.

- 1 2 If a BACKSPACE statement causes the implicit writing of an endfile record, the file is positioned before the  
2 record that precedes the endfile record.
- 3 3 Backspacing a file that is connected but does not exist is prohibited.
- 4 4 Backspacing over records written using list-directed or namelist formatting is prohibited.

**NOTE 12.51**

An example of a BACKSPACE statement is:

```
BACKSPACE (10, IOSTAT = N)
```

**5 12.8.3 ENDFILE statement**

- 6 1 Execution of an ENDFILE statement for a file connected for sequential access writes an endfile record as the next  
7 record of the file. The file is then positioned after the endfile record, which becomes the last record of the file.  
8 If the file also may be connected for direct access, only those records before the endfile record are considered to  
9 have been written. Thus, only those records may be read during subsequent direct access connections to the file.
- 10 2 After execution of an ENDFILE statement for a file connected for sequential access, a **BACKSPACE** or **REWIND**  
11 statement shall be used to reposition the file prior to execution of any **data transfer** input/output statement or  
12 ENDFILE statement.
- 13 3 Execution of an ENDFILE statement for a file connected for stream access causes the terminal point of the file  
14 to become equal to the current file position. Only **file storage units** before the current position are considered  
15 to have been written; thus only those **file storage units** may be subsequently read. Subsequent stream **output**  
16 **statements** may be used to write further data to the file.
- 17 4 Execution of an ENDFILE statement for a file that is connected but does not exist creates the file; if the file is  
18 connected for sequential access, it is created prior to writing the endfile record.

**NOTE 12.52**

An example of an ENDFILE statement is:

```
ENDFILE K
```

**19 12.8.4 REWIND statement**

- 20 1 Execution of a REWIND statement causes the specified file to be positioned at its initial point.

**NOTE 12.53**

If the file is already positioned at its initial point, execution of this statement has no effect on the position of the file.

- 21 2 Execution of a REWIND statement for a file that is connected but does not exist is permitted and has no effect  
22 on any file.

**NOTE 12.54**

An example of a REWIND statement is:

```
REWIND 10
```

## 12.9 FLUSH statement

R1228 *flush-stmt* is FLUSH *file-unit-number*  
or FLUSH ( *flush-spec-list* )

R1229 *flush-spec* is [UNIT =] *file-unit-number*  
or IOSTAT = *scalar-int-variable*  
or IOMSG = *iomsg-variable*  
or ERR = *label*

C1242 No specifier shall appear more than once in a given *flush-spec-list*.

C1243 A *file-unit-number* shall be specified in a *flush-spec-list*; if the optional characters UNIT= are omitted from the unit specifier, the *file-unit-number* shall be the first item in the *flush-spec-list*.

C1244 (R1229) The *label* in the ERR= specifier shall be the statement label of a branch target statement that appears in the same inclusive scope as the FLUSH statement.

1 The IOSTAT=, IOMSG= and ERR= specifiers are described in 12.11.

2 Execution of a FLUSH statement causes data written to an external file to be available to other processes, or causes data placed in an external file by means other than Fortran to be available to a READ statement. These actions are processor dependent.

3 Execution of a FLUSH statement for a file that is connected but does not exist is permitted and has no effect on any file. A FLUSH statement has no effect on file position.

4 Execution of a FLUSH statement performs a wait operation for all pending asynchronous data transfer operations for the specified unit.

### NOTE 12.55

Because this document does not specify the mechanism of file storage, the exact meaning of the flush operation is not precisely defined. It is expected that the flush operation will make all data written to a file available to other processes or devices, or make data recently added to a file by other processes or devices available to the program via a subsequent read operation. This is commonly called “flushing input/output buffers”.

### NOTE 12.56

An example of a FLUSH statement is:

```
FLUSH (10, IOSTAT = N)
```

## 12.10 File inquiry statement

### 12.10.1 Forms of the INQUIRE statement

1 The INQUIRE statement can be used to inquire about properties of a particular named file, of the connection to a particular unit, or the number of file storage units required for an output list. There are three forms of the INQUIRE statement: inquire by file, which uses the FILE= specifier, inquire by unit, which uses the UNIT= specifier, and inquire by output list, which uses only the IOLENGTH= specifier. Assignments to specifier variables are performed as if by intrinsic assignment.

2 For inquiry by unit, the unit specified need not exist or be connected to a file. If it is connected to a file, the inquiry is being made about the connection and about the file connected.

3 For inquiry by file, the file specified need not exist or be connected to a unit. If it is connected to a unit, the inquiry is being made about the connection as well as about the file.

1 4 An INQUIRE statement may be executed before, while, or after a file is connected to a **unit**. All values assigned  
2 by an INQUIRE statement are those that are current at the time the statement is executed.

3 R1230 *inquire-stmt* is INQUIRE ( *inquire-spec-list* )  
4 or INQUIRE ( IOLENGTH = *scalar-int-variable* ) ■  
5 ■ *output-item-list*

#### NOTE 12.57

Examples of INQUIRE statements are:

```
INQUIRE (IOLENGTH = IOL) A (1:N)
INQUIRE (UNIT = JOAN, OPENED = LOG_01, NAMED = LOG_02, &
FORM = CHAR_VAR, IOSTAT = IOS)
```

## 6 12.10.2 Inquiry specifiers

### 7 12.10.2.1 Syntax

8 1 Unless constrained, the following inquiry specifiers may be used in either of the inquire by file or inquire by unit  
9 forms of the INQUIRE statement.

10 R1231 *inquire-spec* is [ UNIT = ] *file-unit-number*  
11 or FILE = *file-name-expr*  
12 or ACCESS = *scalar-default-char-variable*  
13 or ACTION = *scalar-default-char-variable*  
14 or ASYNCHRONOUS = *scalar-default-char-variable*  
15 or BLANK = *scalar-default-char-variable*  
16 or DECIMAL = *scalar-default-char-variable*  
17 or DELIM = *scalar-default-char-variable*  
18 or DIRECT = *scalar-default-char-variable*  
19 or ENCODING = *scalar-default-char-variable*  
20 or ERR = *label*  
21 or EXIST = *scalar-logical-variable*  
22 or FORM = *scalar-default-char-variable*  
23 or FORMATTED = *scalar-default-char-variable*  
24 or ID = *scalar-int-expr*  
25 or IOMSG = *iomsg-variable*  
26 or IOSTAT = *scalar-int-variable*  
27 or NAME = *scalar-default-char-variable*  
28 or NAMED = *scalar-logical-variable*  
29 or NEXTREC = *scalar-int-variable*  
30 or NUMBER = *scalar-int-variable*  
31 or OPENED = *scalar-logical-variable*  
32 or PAD = *scalar-default-char-variable*  
33 or PENDING = *scalar-logical-variable*  
34 or POS = *scalar-int-variable*  
35 or POSITION = *scalar-default-char-variable*  
36 or READ = *scalar-default-char-variable*  
37 or READWRITE = *scalar-default-char-variable*  
38 or RECL = *scalar-int-variable*  
39 or ROUND = *scalar-default-char-variable*  
40 or SEQUENTIAL = *scalar-default-char-variable*  
41 or SIGN = *scalar-default-char-variable*  
42 or SIZE = *scalar-int-variable*  
43 or STREAM = *scalar-default-char-variable*  
44 or UNFORMATTED = *scalar-default-char-variable*



1 or if the connection is not for formatted input/output, the *scalar-default-char-variable* is assigned the value  
2 UNDEFINED.

### 3 **12.10.2.8 DELIM= specifier in the INQUIRE statement**

4 1 The *scalar-default-char-variable* in the DELIM= specifier is assigned the value APOSTROPHE, QUOTE, or  
5 NONE, corresponding to the delimiter mode in effect for a connection for formatted input/output. If there is no  
6 connection or if the connection is not for formatted input/output, the *scalar-default-char-variable* is assigned the  
7 value UNDEFINED.

### 8 **12.10.2.9 DIRECT= specifier in the INQUIRE statement**

9 1 The *scalar-default-char-variable* in the DIRECT= specifier is assigned the value YES if DIRECT is included in  
10 the set of allowed access methods for the file, NO if DIRECT is not included in the set of allowed access methods  
11 for the file, and UNKNOWN if the processor is unable to determine whether DIRECT is included in the set of  
12 allowed access methods for the file or if the *unit* identified by *file-unit-number* is not connected to a file.

### 13 **12.10.2.10 ENCODING= specifier in the INQUIRE statement**

14 1 The *scalar-default-char-variable* in the ENCODING= specifier is assigned the value UTF-8 if the connection is  
15 for formatted input/output with an encoding form of UTF-8, and is assigned the value UNDEFINED if the  
16 connection is for unformatted input/output. If there is no connection, it is assigned the value UTF-8 if the  
17 processor is able to determine that the encoding form of the file is UTF-8; if the processor is unable to determine  
18 the encoding form of the file or if the *unit* identified by *file-unit-number* is not connected to a file, the variable is  
19 assigned the value UNKNOWN.

#### **NOTE 12.58**

The value assigned could be something other than UTF-8, UNDEFINED, or UNKNOWN if the processor supports other specific encoding forms (e.g. UTF-16BE).
--

### 20 **12.10.2.11 EXIST= specifier in the INQUIRE statement**

21 1 Execution of an INQUIRE by file statement causes the *scalar-logical-variable* in the EXIST= specifier to be  
22 assigned the value true if there exists a file with the specified name; otherwise, false is assigned. Execution of an  
23 INQUIRE by unit statement causes true to be assigned if the specified *unit* exists; otherwise, false is assigned.

### 24 **12.10.2.12 FORM= specifier in the INQUIRE statement**

25 1 The *scalar-default-char-variable* in the FORM= specifier is assigned the value FORMATTED if the connection  
26 is for formatted input/output, and is assigned the value UNFORMATTED if the connection is for unformatted  
27 input/output. If there is no connection, it is assigned the value UNDEFINED.

### 28 **12.10.2.13 FORMATTED= specifier in the INQUIRE statement**

29 1 The *scalar-default-char-variable* in the FORMATTED= specifier is assigned the value YES if FORMATTED is  
30 included in the set of allowed forms for the file, NO if FORMATTED is not included in the set of allowed forms  
31 for the file, and UNKNOWN if the processor is unable to determine whether FORMATTED is included in the  
32 set of allowed forms for the file or if the *unit* identified by *file-unit-number* is not connected to a file.

### 33 **12.10.2.14 ID= specifier in the INQUIRE statement**

34 1 The value of the expression specified in the ID= specifier shall be the identifier of a pending data transfer operation  
35 for the specified *unit*. This specifier interacts with the PENDING= specifier (12.10.2.21).

1     **12.10.2.15 NAME= specifier in the INQUIRE statement**

- 2     1 The *scalar-default-char-variable* in the NAME= specifier is assigned the value of the name of the file if the file  
3     has a name; otherwise, it becomes undefined. The value assigned shall be suitable for use as the value of the  
4     *file-name-expr* in the FILE= specifier in an OPEN statement.

**NOTE 12.59**

If this specifier appears in an INQUIRE by file statement, its value is not necessarily the same as the name given in the FILE= specifier.

The processor could assign a file name qualified by a user identification, device, directory, or other relevant information.

- 5     2 The case of the characters assigned to *scalar-default-char-variable* is processor dependent.

6     **12.10.2.16 NAMED= specifier in the INQUIRE statement**

- 7     1 The *scalar-logical-variable* in the NAMED= specifier is assigned the value true if the file has a name; otherwise,  
8     it is assigned the value false.

9     **12.10.2.17 NEXTREC= specifier in the INQUIRE statement**

- 10    1 The *scalar-int-variable* in the NEXTREC= specifier is assigned the value  $n + 1$ , where  $n$  is the record number of  
11    the last record read from or written to the connection for direct access. If there is a connection but no records have  
12    been read or written since the connection, the *scalar-int-variable* is assigned the value 1. If there is no connection,  
13    the connection is not for direct access, or the position is indeterminate because of a previous error condition, the  
14    *scalar-int-variable* becomes undefined. If there are pending data transfer operations for the specified **unit**, the  
15    value assigned is computed as if all the pending data transfers had already completed.

16    **12.10.2.18 NUMBER= specifier in the INQUIRE statement**

- 17    1 Execution of an INQUIRE by file statement causes the *scalar-int-variable* in the NUMBER= specifier to be  
18    assigned the value of the **external unit** number of the unit that is connected to the file. If more than one unit  
19    on an **image** is connected to the file, which of the connected external unit numbers is assigned to the *scalar-int-*  
20    *variable* is processor dependent. If there is no **unit** connected to the file, the value  $-1$  is assigned. Execution of  
21    an INQUIRE by **unit** statement causes the *scalar-int-variable* to be assigned the value of *file-unit-number*.

22    **12.10.2.19 OPENED= specifier in the INQUIRE statement**

- 23    1 Execution of an INQUIRE by file statement causes the *scalar-logical-variable* in the OPENED= specifier to be  
24    assigned the value true if the file specified is connected to a **unit**; otherwise, false is assigned. Execution of an  
25    INQUIRE by unit statement causes the *scalar-logical-variable* to be assigned the value true if the specified **unit**  
26    is connected to a file; otherwise, false is assigned.

27    **12.10.2.20 PAD= specifier in the INQUIRE statement**

- 28    1 The *scalar-default-char-variable* in the PAD= specifier is assigned the value YES or NO, corresponding to the  
29    pad mode in effect for a connection for formatted input/output. If there is no connection or if the connection is  
30    not for formatted input/output, the *scalar-default-char-variable* is assigned the value UNDEFINED.

31    **12.10.2.21 PENDING= specifier in the INQUIRE statement**

- 32    1 The PENDING= specifier is used to determine whether previously pending asynchronous data transfers are  
33    complete. A data transfer operation is previously pending if it is pending at the beginning of execution of the  
34    INQUIRE statement.

- 1 2 If an ID= specifier appears and the specified data transfer operation is complete, then the variable specified in  
2 the PENDING= specifier is assigned the value false and the INQUIRE statement performs the wait operation  
3 for the specified data transfer.
- 4 3 If the ID= specifier is omitted and all previously pending data transfer operations for the specified **unit** are  
5 complete, then the variable specified in the PENDING= specifier is assigned the value false and the INQUIRE  
6 statement performs wait operations for all previously pending data transfers for the specified **unit**.
- 7 4 In all other cases, the variable specified in the PENDING= specifier is assigned the value true, no wait operations  
8 are performed, and the previously pending data transfers remain pending after the execution of the INQUIRE  
9 statement.

**NOTE 12.60**

The processor has considerable flexibility in defining when it considers a transfer to be complete. Any of the following approaches could be used:

- The INQUIRE statement could consider an asynchronous data transfer to be incomplete until after the corresponding wait operation. In this case PENDING= would always return true unless there were no previously pending data transfers for the **unit**.
- The INQUIRE statement could wait for all specified data transfers to complete and then always return false for PENDING=.
- The INQUIRE statement could actually test the state of the specified data transfer operations.

**12.10.2.22 POS= specifier in the INQUIRE statement**

- 10  
11 1 The *scalar-int-variable* in the POS= specifier is assigned the number of the **file storage unit** immediately following  
12 the current position of a file connected for stream access. If the file is positioned at its terminal position, the  
13 variable is assigned a value one greater than the number of the highest-numbered **file storage unit** in the file.  
14 If there are pending data transfer operations for the specified unit, the value assigned is computed as if all the  
15 pending data transfers had already completed. If there is no connection, the file is not connected for stream  
16 access, or if the position of the file is indeterminate because of previous error conditions, the variable becomes  
17 undefined.

**12.10.2.23 POSITION= specifier in the INQUIRE statement**

- 18  
19 1 The *scalar-default-char-variable* in the POSITION= specifier is assigned the value REWIND if the connection  
20 was opened for positioning at its initial point, APPEND if the connection was opened for positioning before its  
21 endfile record or at its terminal point, and ASIS if the connection was opened without changing its position.  
22 If there is no connection or if the file is connected for direct access, the *scalar-default-char-variable* is assigned  
23 the value UNDEFINED. If the file has been repositioned since the connection, the *scalar-default-char-variable*  
24 is assigned a processor-dependent value, which shall not be REWIND unless the file is positioned at its initial  
25 point and shall not be APPEND unless the file is positioned so that its endfile record is the next record or at its  
26 terminal point if it has no endfile record.

**12.10.2.24 READ= specifier in the INQUIRE statement**

- 27  
28 1 The *scalar-default-char-variable* in the READ= specifier is assigned the value YES if READ is included in the  
29 set of allowed actions for the file, NO if READ is not included in the set of allowed actions for the file, and  
30 UNKNOWN if the processor is unable to determine whether READ is included in the set of allowed actions for  
31 the file or if the **unit** identified by *file-unit-number* is not connected to a file.

**12.10.2.25 READWRITE= specifier in the INQUIRE statement**

- 32  
33 1 The *scalar-default-char-variable* in the READWRITE= specifier is assigned the value YES if READWRITE is  
34 included in the set of allowed actions for the file, NO if READWRITE is not included in the set of allowed actions



1 for the file, and UNKNOWN if the processor is unable to determine whether READWRITE is included in the  
2 set of allowed actions for the file or if the **unit** identified by *file-unit-number* is not connected to a file.

### 3 **12.10.2.26 RECL= specifier in the INQUIRE statement**

4 1 The *scalar-int-variable* in the RECL= specifier is assigned the value of the record length of a connection for direct  
5 access, or the value of the maximum record length of a connection for sequential access. If the connection is for  
6 formatted input/output, the length is the number of characters for all records that contain only characters of  
7 default kind. If the connection is for unformatted input/output, the length is measured in **file storage units**. If  
8 there is no connection, the *scalar-int-variable* is assigned the value  $-1$ , and if the connection is for stream access,  
9 the *scalar-int-variable* is assigned the value  $-2$ .

### 10 **12.10.2.27 ROUND= specifier in the INQUIRE statement**

11 1 The *scalar-default-char-variable* in the ROUND= specifier is assigned the value UP, DOWN, ZERO, NEAREST,  
12 COMPATIBLE, or PROCESSOR\_DEFINED, corresponding to the input/output rounding mode in effect for  
13 a connection for formatted input/output. If there is no connection or if the connection is not for formatted  
14 input/output, the *scalar-default-char-variable* is assigned the value UNDEFINED. The processor shall return the  
15 value PROCESSOR\_DEFINED only if the behavior of the input/output rounding mode is different from that  
16 of the UP, DOWN, ZERO, NEAREST, and COMPATIBLE modes.

### 17 **12.10.2.28 SEQUENTIAL= specifier in the INQUIRE statement**

18 1 The *scalar-default-char-variable* in the SEQUENTIAL= specifier is assigned the value YES if SEQUENTIAL is  
19 included in the set of allowed access methods for the file, NO if SEQUENTIAL is not included in the set of allowed  
20 access methods for the file, and UNKNOWN if the processor is unable to determine whether SEQUENTIAL is  
21 included in the set of allowed access methods for the file or if the **unit** identified by *file-unit-number* is not  
22 connected to a file.

### 23 **12.10.2.29 SIGN= specifier in the INQUIRE statement**

24 1 The *scalar-default-char-variable* in the SIGN= specifier is assigned the value PLUS, SUPPRESS, or PRO-  
25 CESSOR\_DEFINED, corresponding to the sign mode in effect for a connection for formatted input/output.  
26 If there is no connection, or if the connection is not for formatted input/output, the *scalar-default-char-variable*  
27 is assigned the value UNDEFINED.

### 28 **12.10.2.30 SIZE= specifier in the INQUIRE statement**

29 1 The *scalar-int-variable* in the SIZE= specifier is assigned the size of the file in **file storage units**. If the file size  
30 cannot be determined or if the **unit** identified by *file-unit-number* is not connected to a file, the variable is assigned  
31 the value  $-1$ .  
32 2 For a file that may be connected for stream access, the file size is the number of the highest-numbered **file storage**  
33 **unit** in the file.  
34 3 For a file that may be connected for sequential or direct access, the file size may be different from the number of  
35 **storage units** implied by the data in the records; the exact relationship is processor dependent.  
36 4 If there are pending data transfer operations for the specified unit, the value assigned is computed as if all the  
37 pending data transfers had already completed.

### 38 **12.10.2.31 STREAM= specifier in the INQUIRE statement**

39 1 The *scalar-default-char-variable* in the STREAM= specifier is assigned the value YES if STREAM is included in  
40 the set of allowed access methods for the file, NO if STREAM is not included in the set of allowed access methods  
41 for the file, and UNKNOWN if the processor is unable to determine whether STREAM is included in the set of  
42 allowed access methods for the file or if the **unit** identified by *file-unit-number* is not connected to a file.

### 1 12.10.2.32 UNFORMATTED= specifier in the INQUIRE statement

2 1 The *scalar-default-char-variable* in the UNFORMATTED= specifier is assigned the value YES if UNFORMAT-  
3 TED is included in the set of allowed forms for the file, NO if UNFORMATTED is not included in the set of  
4 allowed forms for the file, and UNKNOWN if the processor is unable to determine whether UNFORMATTED is  
5 included in the set of allowed forms for the file or if the *unit* identified by *file-unit-number* is not connected to a  
6 file.

### 7 12.10.2.33 WRITE= specifier in the INQUIRE statement

8 1 The *scalar-default-char-variable* in the WRITE= specifier is assigned the value YES if WRITE is included in the  
9 set of allowed actions for the file, NO if WRITE is not included in the set of allowed actions for the file, and  
10 UNKNOWN if the processor is unable to determine whether WRITE is included in the set of allowed actions for  
11 the file or if the *unit* identified by *file-unit-number* is not connected to a file.

## 12 12.10.3 Inquire by output list

13 1 The *scalar-int-variable* in the IOLENGTH= specifier is assigned the processor-dependent number of *file storage*  
14 *units* that would be required to store the data of the output list in an unformatted file. The value shall be suitable  
15 as a *RECL= specifier* in an *OPEN statement* that connects a file for unformatted direct access if data will be  
16 read from or written to the file using data transfer statements with an input/output list that specifies transfer of  
17 a sequence of objects having the same types, type parameters, and extents, in the same order as the output list  
18 in the INQUIRE statement.

19 2 The output list in an INQUIRE statement shall not contain any derived-type list items that require a *defined*  
20 *input/output* procedure as described in subclause 12.6.3. If a derived-type list item appears in the output list, the  
21 value returned for the IOLENGTH= specifier assumes that no *defined input/output* procedure will be invoked.

## 22 12.11 Error, end-of-record, and end-of-file conditions

### 23 12.11.1 Occurrence of input/output conditions

24 1 The set of input/output error conditions is processor dependent. Except as otherwise specified, when an error  
25 condition occurs or is detected is processor dependent.

26 2 An end-of-record condition occurs when a nonadvancing *input statement* attempts to transfer data from a position  
27 beyond the end of the current record, unless the file is a *stream file* and the current record is at the end of the  
28 file (an end-of-file condition occurs instead).

29 3 An end-of-file condition occurs when

- 30 • an endfile record is encountered during the reading of a file connected for sequential access,
- 31 • an attempt is made to read a record beyond the end of an *internal file*, or
- 32 • an attempt is made to read beyond the end of a *stream file*.

33 4 An end-of-file condition may occur at the beginning of execution of an *input statement*. An end-of-file condition  
34 also may occur during execution of a formatted *input statement* when more than one record is required by the  
35 interaction of the input list and the format. An end-of-file condition also may occur during execution of a stream  
36 *input statement*.

### 37 12.11.2 Error conditions and the ERR= specifier

38 1 If an error condition occurs during execution of an input/output statement, the position of the file becomes  
39 indeterminate.

1 2 If an error condition occurs during execution of an input/output statement that contains neither an ERR= nor  
2 IOSTAT= specifier, **error termination** is initiated. If an error condition occurs during execution of an input/output  
3 statement that contains either an ERR= specifier or an IOSTAT= specifier then:

- 4 (1) processing of the input/output list, if any, terminates;
- 5 (2) if the statement is a **data transfer statement** or the error occurs during a wait operation, all *do-*  
6 *variables* in the statement that initiated the transfer become undefined;
- 7 (3) if an **IOSTAT= specifier** appears, the *scalar-int-variable* in the **IOSTAT= specifier** becomes defined  
8 as specified in 12.11.5;
- 9 (4) if an **IOMSG= specifier** appears, the *iomsg-variable* becomes defined as specified in 12.11.6;
- 10 (5) if the statement is a **READ statement** and it contains a **SIZE= specifier**, the *scalar-int-variable* in  
11 the **SIZE= specifier** becomes defined as specified in 12.6.2.15;
- 12 (6) if the statement is a **READ statement** or the error condition occurs in a wait operation for a transfer  
13 initiated by a **READ statement**, all input items or namelist group objects in the statement that  
14 initiated the transfer become undefined;
- 15 (7) if an ERR= specifier appears, a branch to the statement labeled by the *label* in the ERR= specifier  
16 occurs.

### 17 12.11.3 End-of-file condition and the END= specifier

18 1 If an end-of-file condition occurs during execution of an input/output statement that contains neither an END=  
19 specifier nor an IOSTAT= specifier, **error termination** is initiated. If an end-of-file condition occurs during  
20 execution of an input/output statement that contains either an END= specifier or an IOSTAT= specifier, and  
21 an error condition does not occur then:

- 22 (1) processing of the input list, if any, terminates;
- 23 (2) if the statement is a **data transfer statement** or the end-of-file condition occurs during a wait operation,  
24 all *do-variables* in the statement that initiated the transfer become undefined;
- 25 (3) if the statement is an **input statement** or the end-of-file condition occurs during a wait operation  
26 for a transfer initiated by an **input statement**, all input list items or namelist group objects in the  
27 statement that initiated the transfer become undefined;
- 28 (4) if the file specified in the **input statement** is an external record file, it is positioned after the endfile  
29 record;
- 30 (5) if an **IOSTAT= specifier** appears, the *scalar-int-variable* in the **IOSTAT= specifier** becomes defined  
31 as specified in 12.11.5;
- 32 (6) if an **IOMSG= specifier** appears, the *iomsg-variable* becomes defined as specified in 12.11.6;
- 33 (7) if an END= specifier appears, a branch to the statement labeled by the *label* in the END= specifier  
34 occurs.

### 35 12.11.4 End-of-record condition and the EOR= specifier

36 1 If an end-of-record condition occurs during execution of an input/output statement that contains neither an  
37 EOR= specifier nor an IOSTAT= specifier, **error termination** is initiated. If an end-of-record condition occurs  
38 during execution of an input/output statement that contains either an EOR= specifier or an IOSTAT= specifier,  
39 and an error condition does not occur then:

- 40 (1) if the pad mode has the value
  - 41 (a) YES, the record is padded with blanks to satisfy the **effective item** (12.6.4.5.3) and correspond-  
42 ing data edit descriptors that require more characters than the record contains,
  - 43 (b) NO, the input list item becomes undefined;
- 44 (2) processing of the input list, if any, terminates;
- 45 (3) if the statement is a **data transfer statement** or the end-of-record condition occurs during a wait  
46 operation, all *do-variables* in the statement that initiated the transfer become undefined;
- 47 (4) the file specified in the input statement is positioned after the current record;

- 1 (5) if an `IOSTAT= specifier` appears, the *scalar-int-variable* in the `IOSTAT= specifier` becomes defined  
 2 as specified in 12.11.5;  
 3 (6) if an `IOMSG= specifier` appears, the *iomsg-variable* becomes defined as specified in 12.11.6;  
 4 (7) if a `SIZE= specifier` appears, the *scalar-int-variable* in the `SIZE= specifier` becomes defined as spe-  
 5 cified in (12.6.2.15);  
 6 (8) if an `EOR= specifier` appears, a branch to the statement labeled by the *label* in the `EOR= specifier`  
 7 occurs.

### 8 12.11.5 IOSTAT= specifier

9 1 Execution of an input/output statement containing the `IOSTAT= specifier` causes the *scalar-int-variable* in the  
 10 `IOSTAT= specifier` to become defined with

- 11 • a zero value if neither an error condition, an end-of-file condition, nor an end-of-record condition occurs,
- 12 • the processor-dependent positive integer value of the constant `IOSTAT_INQUIRE_INTERNAL_UNIT`  
 13 from the intrinsic module `ISO_FORTRAN_ENV` (16.10.2) if a `unit` number in an `INQUIRE statement`  
 14 identifies an `internal file`,
- 15 • a processor-dependent positive integer value different from `IOSTAT_INQUIRE_INTERNAL_UNIT` if any  
 16 other error condition occurs,
- 17 • the processor-dependent negative integer value of the constant `IOSTAT_END` (16.10.2.16) from the intrinsic  
 18 module `ISO_FORTRAN_ENV` if an end-of-file condition occurs and no error condition occurs,
- 19 • the processor-dependent negative integer value of the constant `IOSTAT_EOR` (16.10.2.17) from the intrinsic  
 20 module `ISO_FORTRAN_ENV` if an end-of-record condition occurs and no error condition or end-of-file  
 21 condition occurs, or
- 22 • a processor-dependent negative integer value different from `IOSTAT_EOR` and `IOSTAT_END`, if the IO-  
 23 `STAT= specifier` appears in a `FLUSH statement` and the processor does not support the flush operation for  
 24 the specified unit.

#### NOTE 12.61

An end-of-file condition can occur only for sequential or stream input and an end-of-record condition can occur only for nonadvancing input.

For example,

```

READ (FMT = "(E8.3)", UNIT = 3, IOSTAT = IOSS) X
IF (IOSS < 0) THEN
    ! Perform end-of-file processing on the file connected to unit 3.
    CALL END_PROCESSING
ELSE IF (IOSS > 0) THEN
    ! Perform error processing
    CALL ERROR_PROCESSING
END IF

```

### 25 12.11.6 IOMSG= specifier

- 26 1 If an error, end-of-file, or end-of-record condition occurs during execution of an input/output statement, *iomsg-*  
 27 *variable* is assigned an explanatory message as if by `intrinsic assignment`. If no such condition occurs, the definition  
 28 status and value of *iomsg-variable* are unchanged.

## 29 12.12 Restrictions on input/output statements

- 30 1 If a `unit`, or a file connected to a `unit`, does not have all of the properties required for the execution of certain  
 31 input/output statements, those statements shall not refer to the `unit`.

- 1 2 An input/output statement that is executed while another input/output statement is being executed is a recursive  
2 input/output statement. A recursive input/output statement shall not identify an **external unit** that is identified  
3 by another input/output statement being executed except that a child **data transfer statement** may identify its  
4 parent **data transfer statement external unit**.
- 5 3 An input/output statement shall not cause the value of any established format specification to be modified.
- 6 4 A recursive input/output statement shall not modify the value of any **internal unit** except that a recursive **WRITE**  
7 **statement** may modify the **internal unit** identified by that recursive **WRITE statement**.
- 8 5 The value of a specifier in an input/output statement shall not depend on the definition or evaluation of any other  
9 specifier in the *io-control-spec-list* or *inquire-spec-list* in that statement. The value of an *internal-file-variable* or  
10 of a **FMT=**, **ID=**, **IOMSG=**, **IOSTAT=**, or **SIZE=** specifier shall not depend on the value of any *input-item* or  
11 *io-implied-do do-variable* in the same statement.
- 12 6 The value of any subscript or substring bound of a variable that appears in a specifier in an input/output  
13 statement shall not depend on any *input-item*, *io-implied-do do-variable*, or on the definition or evaluation of any  
14 other specifier in the *io-control-spec-list* or *inquire-spec-list* in that statement.
- 15 7 In a **data transfer statement**, the variable specified in an **IOSTAT=**, **IOMSG=**, or **SIZE=** specifier, if any, shall  
16 not be associated with any entity in the data transfer input/output list (12.6.3) or *namelist-group-object-list*, nor  
17 with a *do-variable* of an *io-implied-do* in the data transfer input/output list.
- 18 8 In a **data transfer statement**, if a variable specified in an **IOSTAT=**, **IOMSG=**, or **SIZE=** specifier is an array  
19 element reference, its subscript values shall not be affected by the data transfer, the *io-implied-do* processing, or  
20 the definition or evaluation of any other specifier in the *io-control-spec-list*.
- 21 9 A variable that may become defined or undefined as a result of its use in a specifier in an **INQUIRE statement**,  
22 or any associated entity, shall not appear in another specifier in the same **INQUIRE statement**.

**NOTE 12.62**

Restrictions on the evaluation of expressions (10.1.4) prohibit certain side effects.

1

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## 13 Input/output editing

### 13.1 Format specifications

- 1 A format used in conjunction with a [data transfer statement](#) provides information that directs the editing between the internal representation of data and the characters of a sequence of formatted records.
- 2 A *format* ([12.6.2.2](#)) in a [data transfer statement](#) may refer to a [FORMAT statement](#) or to a character expression that contains a format specification. A format specification provides explicit editing information. The *format* alternatively may be an asterisk (\*), which indicates list-directed formatting ([13.10](#)). Namelist formatting ([13.11](#)) may be indicated by specifying a *namelist-group-name* instead of a *format*.

### 13.2 Explicit format specification methods

#### 13.2.1 FORMAT statement

R1301 *format-stmt* is [FORMAT \*format-specification\*](#)

R1302 *format-specification* is ( [ [format-items](#) ] )  
or ( [ [format-items](#), ] [unlimited-format-item](#) )

C1301 (R1301) The *format-stmt* shall be labeled.

- 1 Blank characters may precede the initial left parenthesis of the format specification. Additional blank characters may appear at any point within the format specification, with no effect on the interpretation of the format specification, except within a character string edit descriptor ([13.9](#)).

#### NOTE 13.1

Examples of FORMAT statements are:

```
5      FORMAT (1PE12.4, I10)
9      FORMAT (I12, /, ' Dates: ', 2 (2I3, I5))
```

#### 13.2.2 Character format specification

- 1 A character expression used as a *format* in a formatted input/output statement shall evaluate to a character string whose leading part is a valid format specification.

#### NOTE 13.2

The format specification begins with a left parenthesis and ends with a right parenthesis.

- 2 All character positions up to and including the final right parenthesis of the format specification shall be defined at the time the [data transfer statement](#) is executed, and shall not become redefined or undefined during the execution of the statement. Character positions, if any, following the right parenthesis that ends the format specification need not be defined and may contain any character data with no effect on the interpretation of the format specification.
- 3 If the *format* is a character array, it is treated as if all of the elements of the array were specified in array element order and were concatenated. However, if a *format* is a character array element, the format specification shall be entirely within that array element.

**NOTE 13.3**

If a character constant is used as a *format* in *data transfer statement*, care shall be taken that the value of the character constant is a valid format specification. In particular, if a format specification delimited by apostrophes contains a character constant edit descriptor delimited with apostrophes, two apostrophes shall be written to delimit the edit descriptor and four apostrophes shall be written for each apostrophe that occurs within the edit descriptor. For example, the text:

```
2 ISN'T 3
```

can be written by various combinations of *output statements* and format specifications:

```
WRITE (6, 100) 2, 3
100 FORMAT (1X, I1, 1X, 'ISN'T', 1X, I1)
WRITE (6, '(1X, I1, 1X, ''ISN''''T'', 1X, I1)') 2, 3
WRITE (6, '(A)') ' 2 ISN'T 3'
```

Doubling of internal apostrophes usually can be avoided by using quotation marks to delimit the format specification and doubling of internal quotation marks usually can be avoided by using apostrophes as delimiters.

## 13.3 Form of a format item list

### 13.3.1 Syntax

R1303 *format-items* is *format-item* [ [ , ] *format-item* ] ...

R1304 *format-item* is [ *r* ] *data-edit-desc*  
 or *control-edit-desc*  
 or *char-string-edit-desc*  
 or [ *r* ] ( *format-items* )

R1305 *unlimited-format-item* is \* ( *format-items* )

R1306 *r* is *int-literal-constant*

C1302 (R1303) The optional comma shall not be omitted except

- between a P edit descriptor and an immediately following F, E, EN, ES, D, or G edit descriptor (13.8.5), possibly preceded by a repeat specification,
- before a slash edit descriptor when the optional repeat specification does not appear (13.8.2),
- after a slash edit descriptor, or
- before or after a colon edit descriptor (13.8.3)

C1303 (R1305) An *unlimited-format-item* shall contain at least one data edit descriptor.

C1304 (R1306) *r* shall be positive.

C1305 (R1306) A kind parameter shall not be specified for *r*.

1 The integer literal constant *r* is called a repeat specification.

### 13.3.2 Edit descriptors

1 An edit descriptor is a data edit descriptor (*data-edit-desc*), control edit descriptor (*control-edit-desc*), or character string edit descriptor (*char-string-edit-desc*).



- 1 R1307 *data-edit-desc* is I *w* [ . *m* ]  
 2 or B *w* [ . *m* ]  
 3 or O *w* [ . *m* ]  
 4 or Z *w* [ . *m* ]  
 5 or F *w* . *d*  
 6 or E *w* . *d* [ E *e* ]  
 7 or EN *w* . *d* [ E *e* ]  
 8 or ES *w* . *d* [ E *e* ]  
 9 or EX *w* . *d* [ E *e* ]  
 10 or G *w* [ . *d* [ E *e* ] ]  
 11 or L *w*  
 12 or A [ *w* ]  
 13 or D *w* . *d*  
 14 or DT [ *char-literal-constant* ] [ ( *v-list* ) ]
- 15 R1308 *w* is *int-literal-constant*
- 16 R1309 *m* is *int-literal-constant*
- 17 R1310 *d* is *int-literal-constant*
- 18 R1311 *e* is *int-literal-constant*
- 19 R1312 *v* is *signed-int-literal-constant*
- 20 C1306 (R1308) *w* shall be zero or positive for the I, B, O, Z, D, E, EN, ES, EX, F, and G edit descriptors. *w*  
 21 shall be positive for all other edit descriptors.
- 22 C1307 (R1307) For the G edit descriptor, *d* shall be specified if *w* is not zero.
- 23 C1308 (R1307) For the G edit descriptor, *e* shall not be specified if *w* is zero.
- 24 C1309 (R1307) A kind parameter shall not be specified for the *char-literal-constant* in the DT edit descriptor,  
 25 or for *w*, *m*, *d*, *e*, and *v*.
- 26 2 An I, B, O, Z, F, E, EN, ES, EX, G, L, A, D, or DT edit descriptor indicates the manner of editing.
- 27 R1313 *control-edit-desc* is *position-edit-desc*  
 28 or [ *r* ] /  
 29 or :  
 30 or *sign-edit-desc*  
 31 or *k* P  
 32 or *blank-interp-edit-desc*  
 33 or *round-edit-desc*  
 34 or *decimal-edit-desc*
- 35 R1314 *k* is *signed-int-literal-constant*
- 36 C1310 (R1314) A kind parameter shall not be specified for *k*.
- 37 3 In *k* P, *k* is called the scale factor.
- 38 R1315 *position-edit-desc* is T *n*  
 39 or TL *n*  
 40 or TR *n*  
 41 or *n* X
- 42 R1316 *n* is *int-literal-constant*
- 43 C1311 (R1316) *n* shall be positive.

1 C1312 (R1316) A kind parameter shall not be specified for *n*.

2 R1317 *sign-edit-desc* is SS  
3 or SP  
4 or S

5 R1318 *blank-interp-edit-desc* is BN  
6 or BZ

7 R1319 *round-edit-desc* is RU  
8 or RD  
9 or RZ  
10 or RN  
11 or RC  
12 or RP

13 R1320 *decimal-edit-desc* is DC  
14 or DP

15 4 A T, TL, TR, X, slash, colon, SS, SP, S, P, BN, BZ, RU, RD, RZ, RN, RC, RP, DC, or DP edit descriptor  
16 indicates the manner of editing.

17 R1321 *char-string-edit-desc* is *char-literal-constant*

18 C1313 (R1321) A kind parameter shall not be specified for the *char-literal-constant*.

19 5 Each *rep-char* in a character string edit descriptor shall be capable of representation by the processor.

20 6 A character string edit descriptor provides constant data to be output, and is not valid for input.

21 7 The edit descriptors are without regard to case except within a character string edit descriptor.

### 22 13.3.3 Fields

23 1 A field is a part of a record that is read on input or written on output when format control encounters a data  
24 edit descriptor or a character string edit descriptor. The field width is the size in characters of the field.

## 25 13.4 Interaction between input/output list and format

26 1 The start of formatted data transfer using a format specification initiates format control (12.6.4.5.3). Each action  
27 of format control depends on information jointly provided by the next edit descriptor in the format specification  
28 and the next *effective item* in the input/output list, if one exists.

29 2 If an input/output list specifies at least one *effective item*, at least one data edit descriptor shall exist in the  
30 format specification.

#### NOTE 13.4

An empty format specification of the form ( ) can be used only if the input/output list has no *effective item* (12.6.4.5). A zero length character item is an *effective item*, but a zero sized array and an implied DO list with an iteration count of zero is not.

31 3 A format specification is interpreted from left to right. The exceptions are format items preceded by a repeat  
32 specification *r*, and format reversion (described below).

33 4 A format item preceded by a repeat specification is processed as a list of *r* items, each identical to the format  
34 item but without the repeat specification and separated by commas.

**NOTE 13.5**

An omitted repeat specification is treated in the same way as a repeat specification whose value is one.

- 1 5 To each data edit descriptor interpreted in a format specification, there corresponds one **effective item** specified by  
 2 the input/output list (12.6.3), except that an input/output list item of type complex requires the interpretation  
 3 of two F, E, EN, ES, D, or G edit descriptors. For each control edit descriptor or character edit descriptor,  
 4 there is no corresponding item specified by the input/output list, and format control communicates information  
 5 directly with the record.
- 6 6 Whenever format control encounters a data edit descriptor in a format specification, it determines whether  
 7 there is a corresponding **effective item** specified by the input/output list. If there is such an item, it transmits  
 8 appropriately edited information between the item and the record, and then format control proceeds. If there is  
 9 no such item, format control terminates.
- 10 7 If format control encounters a colon edit descriptor in a format specification and another effective item is not  
 11 specified, format control terminates.
- 12 8 If format control encounters the rightmost parenthesis of an unlimited format item, control reverts to the leftmost  
 13 parenthesis of that unlimited format item. This reversion of format control has no effect on the changeable modes  
 14 (12.5.2).
- 15 9 If format control encounters the rightmost parenthesis of a complete format specification and another effective  
 16 item is not specified, format control terminates. However, if another effective item is specified, format control  
 17 then reverts to the beginning of the format item terminated by the last preceding right parenthesis that is not  
 18 part of a DT edit descriptor. If there is no such preceding right parenthesis, format control reverts to the first  
 19 left parenthesis of the format specification. If any reversion occurs, the reused portion of the format specification  
 20 shall contain at least one data edit descriptor. If format control reverts to a parenthesis that is preceded by a  
 21 repeat specification, the repeat specification is reused. Reversion of format control, of itself, has no effect on  
 22 the **changeable modes**. The file is positioned in a manner identical to the way it is positioned when a slash edit  
 23 descriptor is processed (13.8.2).

**NOTE 13.6**

Example: The format specification:

```
10 FORMAT (1X, 2(F10.3, I5))
```

with an output list of

```
WRITE (10,10) 10.1, 3, 4.7, 1, 12.4, 5, 5.2, 6
```

produces the same output as the format specification:

```
10 FORMAT (1X, F10.3, I5, F10.3, I5/F10.3, I5, F10.3, I5)
```

**NOTE 13.7**

The effect of an **unlimited-format-item** is as if its enclosed list were preceded by a very large repeat count. There is no file positioning implied by **unlimited-format-item** reversion. This can be used to write what is commonly called a comma separated value record.

For example,

```
WRITE( 10, '( "IARRAY =", *( I0, :, ",") )') IARRAY
```

produces a single record with a header and a comma separated list of integer values.

## 1 13.5 Positioning by format control

- 2 1 After each data edit descriptor or character string edit descriptor is processed, the file is positioned after the last  
3 character read or written in the current record.
- 4 2 After each T, TL, TR, or X edit descriptor is processed, the file is positioned as described in 13.8.1.1. After each  
5 slash edit descriptor is processed, the file is positioned as described in 13.8.2.
- 6 3 During formatted stream output, processing of an A edit descriptor can cause file positioning to occur (13.7.4).
- 7 4 If format control reverts as described in 13.4, the file is positioned in a manner identical to the way it is positioned  
8 when a slash edit descriptor is processed (13.8.2).
- 9 5 During a read operation, any unprocessed characters of the current record are skipped whenever the next record  
10 is read.

## 11 13.6 Decimal symbol

- 12 1 The [decimal symbol](#) is the character that separates the whole and fractional parts in the decimal representation  
13 of a real number in an internal or [external](#) file. When the decimal edit mode is POINT, the [decimal symbol](#) is a  
14 decimal point. When the decimal edit mode is COMMA, the [decimal symbol](#) is a comma.
- 15 2 If the decimal edit mode is COMMA during list-directed input/output, the character used as a value separator  
16 is a semicolon in place of a comma.

## 17 13.7 Data edit descriptors

### 18 13.7.1 Purpose of data edit descriptors

- 19 1 Data edit descriptors cause the conversion of data to or from its internal representation; during formatted stream  
20 output, the A data edit descriptor may also cause file positioning. On input, the specified variable becomes  
21 defined unless an error condition, an end-of-file condition, or an end-of-record condition occurs. On output, the  
22 specified expression is evaluated.
- 23 2 During input from a Unicode file,
  - 24 • characters in the record that correspond to an [ASCII character](#) variable shall have a position in the [ISO](#)  
25 [10646 character collating sequence](#) of 127 or less, and
  - 26 • characters in the record that correspond to a default character variable shall be representable as default  
27 characters.
- 28 3 During input from a non-Unicode file,
  - 29 • characters in the record that correspond to a character variable shall have the kind of the character variable,  
30 and
  - 31 • characters in the record that correspond to a numeric or logical variable shall be default characters.
- 32 4 During output to a Unicode file, all characters transmitted to the record are of [ISO 10646 character](#) kind. If a  
33 character input/output list item or character string edit descriptor contains a character that is not representable  
34 as an [ISO 10646 character](#), the result is processor dependent.
- 35 5 During output to a non-Unicode file, characters transmitted to the record as a result of processing a character  
36 string edit descriptor or as a result of evaluating a numeric, logical, or default character data entity, are of default  
37 kind.

## 13.7.2 Numeric editing

### 13.7.2.1 General rules

1 The I, B, O, Z, F, E, EN, ES, EX, D, and G edit descriptors may be used to specify the input/output of integer, real, and complex data. The following general rules apply.

- (1) On input, leading blanks are not significant. When the input field is not an IEEE exceptional specification or hexadecimal-significand number (13.7.2.3.2), the interpretation of blanks, other than leading blanks, is determined by the blank interpretation mode (13.8.6). Plus signs may be omitted. A field containing only blanks is considered to be zero.
- (2) On input, with F, E, EN, ES, EX, D, and G editing, a **decimal symbol** appearing in the input field overrides the portion of an edit descriptor that specifies the **decimal symbol** location. The input field may have more digits than the processor uses to approximate the value of the datum.
- (3) On output with I, F, E, EN, ES, EX, D, and G editing, the representation of a positive or zero internal value in the field may be prefixed with a plus sign, as controlled by the S, SP, and SS edit descriptors or the processor. The representation of a negative internal value in the field shall be prefixed with a minus sign.
- (4) On output, the representation is right justified in the field. If the number of characters produced by the editing is smaller than the field width, leading blanks are inserted in the field.
- (5) On output, if an exponent exceeds its specified or implied width using the E, EN, ES, EX, D, or G edit descriptor, or the number of characters produced exceeds the field width, the processor shall fill the entire field of width *w* with asterisks. However, the processor shall not produce asterisks if the field width is not exceeded when optional characters are omitted.

#### NOTE 13.8

When the sign mode is PLUS, a plus sign is not optional.

- (6) On output, with I, B, O, Z, D, E, EN, ES, EX, F, and G editing, the specified value of the field width *w* may be zero. In such cases, the processor selects the smallest positive actual field width that does not result in a field filled with asterisks. The specified value of *w* shall not be zero on input.
- (7) On output of a real zero value, the digits in the exponent field shall all be zero.

### 13.7.2.2 Integer editing

1 The *Iw* and *Iw.m* edit descriptors indicate that the field to be edited occupies *w* positions, except when *w* is zero. When *w* is zero, the processor selects the field width. On input, *w* shall not be zero. The specified input/output list item shall be of type integer. The G, B, O, and Z edit descriptor also may be used to edit integer data (13.7.5.2.1, 13.7.2.4).

2 On input, *m* has no effect.

3 In the standard form of the input field for the I edit descriptor, the character string is a *signed-digit-string* (R710), except for the interpretation of blanks. If the input field does not have the standard form and is not acceptable to the processor, an error condition occurs.

4 The output field for the *Iw* edit descriptor consists of zero or more leading blanks followed by a minus sign if the internal value is negative, or an optional plus sign otherwise, followed by the magnitude of the internal value as a *digit-string* without leading zeros.

#### NOTE 13.9

A *digit-string* always consists of at least one digit.

5 The output field for the *Iw.m* edit descriptor is the same as for the *Iw* edit descriptor, except that the *digit-string* consists of at least *m* digits. If necessary, sufficient leading zeros are included to achieve the minimum of *m* digits. The value of *m* shall not exceed the value of *w*, except when *w* is zero. If *m* is zero and the internal value is zero,

1 the output field consists of only blank characters, regardless of the sign control in effect. When  $m$  and  $w$  are both  
2 zero, and the internal value is zero, one blank character is produced.

### 3 13.7.2.3 Real and complex editing

#### 4 13.7.2.3.1 General

5 1 The F, E, EN, ES, and D edit descriptors specify the editing of real and complex data. An input/output list  
6 item corresponding to an F, E, EN, ES, or D edit descriptor shall be real or complex. The G, B, O, and Z edit  
7 descriptors also may be used to edit real and complex data (13.7.5.2.2, 13.7.2.4).

#### 8 13.7.2.3.2 F editing

9 1 The F $w.d$  edit descriptor indicates that the field occupies  $w$  positions, except when  $w$  is zero in which case the  
10 processor selects the field width. The fractional part of the field consists of  $d$  digits. On input,  $w$  shall not be  
11 zero.

12 2 A lower-case letter is equivalent to the corresponding upper-case letter in an IEEE exceptional specification or  
13 the exponent in a numeric input field.

14 3 The standard form of the input field is an IEEE exceptional specification, a hexadecimal-significand number, or  
15 consists of a mantissa optionally followed by an exponent. The form of the mantissa is an optional sign, followed  
16 by a string of one or more digits optionally containing a decimal symbol, including any blanks interpreted as  
17 zeros. The  $d$  has no effect on input if the input field contains a decimal symbol. If the decimal symbol is omitted,  
18 the rightmost  $d$  digits of the string, with leading zeros assumed if necessary, are interpreted as the fractional part  
19 of the value represented. The string of digits may contain more digits than a processor uses to approximate the  
20 value. The form of the exponent is one of the following:

- 21 • a *sign* followed by a *digit-string*;
- 22 • the letter E followed by zero or more blanks, followed by a *signed-digit-string*;
- 23 • the letter D followed by zero or more blanks, followed by a *signed-digit-string*.

24 4 An exponent containing a D is processed identically to an exponent containing an E.

#### NOTE 13.10

If the input field does not contain an exponent, the effect is as if the basic form were followed by an exponent with a value of  $-k$ , where  $k$  is the established scale factor (13.8.5).

25 5 An input field that is an IEEE exceptional specification consists of optional blanks, followed by either

- 26 • an optional sign, followed by the string 'INF' or the string 'INFINITY', or
- 27 • an optional sign, followed by the string 'NAN', optionally followed by zero or more alphanumeric characters  
28 enclosed in parentheses,

29 optionally followed by blanks.

30 6 The value specified by 'INF' or 'INFINITY' is an IEEE infinity; this form shall not be used if the processor does  
31 not support IEEE infinities for the input variable. The value specified by 'NAN' is an IEEE NaN; this form shall  
32 not be used if the processor does not support IEEE NaNs for the input variable. The NaN value is a quiet NaN if  
33 the only nonblank characters in the field are 'NAN' or 'NAN()'; otherwise, the NaN value is processor dependent.  
34 The interpretation of a sign in a NaN input field is processor dependent.

35 7 An input field that is a hexadecimal-significand number consists of an optional sign, followed by the hexadecimal  
36 indicator which is the digit 0 immediately followed by the letter X, followed by a hexadecimal significand followed  
37 by a hexadecimal exponent. A hexadecimal significand is a string of one or more hexadecimal characters optionally  
38 containing a decimal symbol. The decimal symbol indicates the position of the hexadecimal point; if no decimal  
39 symbol appears, the hexadecimal point implicitly follows the last hexadecimal symbol. A hexadecimal exponent

- 1 is the letter P followed by a (decimal) *signed-digit-string*. Embedded blanks are not permitted in a hexadecimal-  
 2 significand number. The value is equal to the significand multiplied by two raised to the power of the exponent,  
 3 negated if the optional sign is minus.
- 4 8 If the input field does not have one of the standard forms, and is not acceptable to the processor, an error  
 5 condition occurs.
- 6 9 For an internal value that is an IEEE infinity, the output field consists of blanks, if necessary, followed by a minus  
 7 sign for negative infinity or an optional plus sign otherwise, followed by the letters 'Inf' or 'Infinity', right justified  
 8 within the field. The minimum field width required for output of the form 'Inf' is 3 if no sign is produced, and  
 9 4 otherwise. The minimum field width required for output of the form 'Infinity' is 8 if no sign is produced, and  
 10 9 otherwise. If  $w$  is greater than or equal to the minimum required for the form 'Infinity', the form 'Infinity' is  
 11 output. If  $w$  is zero or  $w$  is less than the minimum required for the form 'Infinity' and greater than or equal to  
 12 the minimum required for the form 'Inf', the form 'Inf' is output. Otherwise ( $w$  is greater than zero but less than  
 13 the minimum required for any form), the field is filled with asterisks.
- 14 10 For an internal value that is an IEEE NaN, the output field consists of blanks, if necessary, followed by the  
 15 letters 'NaN' and optionally followed by one to  $w-5$  alphanumeric processor-dependent characters enclosed in  
 16 parentheses, right justified within the field. If  $w$  is greater than zero and less than 3, the field is filled with  
 17 asterisks. If  $w$  is zero, the output field is 'NaN'.

**NOTE 13.11**

The processor-dependent characters following 'NaN' might convey additional information about that particular NaN.

- 18 11 For an internal value that is neither an IEEE infinity nor a NaN, the output field consists of blanks, if necessary,  
 19 followed by a minus sign if the internal value is negative, or an optional plus sign otherwise, followed by a string  
 20 of digits that contains a *decimal symbol* and represents the magnitude of the internal value, as modified by the  
 21 established scale factor and rounded (13.7.2.3.8) to  $d$  fractional digits. Leading zeros are not permitted except  
 22 for an optional zero immediately to the left of the *decimal symbol* if the magnitude of the value in the output  
 23 field is less than one. The optional zero shall appear if there would otherwise be no digits in the output field.

**13.7.2.3.3 E and D editing**

- 24 1 The  $Ew.d$ ,  $Dw.d$ , and  $Ew.d Ee$  edit descriptors indicate that the external field occupies  $w$  positions, except when  
 25  $w$  is zero in which case the processor selects the field width. The fractional part of the field contains  $d$  digits,  
 26 unless a scale factor greater than one is in effect. If  $e$  is positive the exponent part contains  $e$  digits, otherwise it  
 27 contains the minimum number of digits required to represent the exponent value. The  $e$  has no effect on input.  
 28
- 29 2 The form and interpretation of the input field is the same as for  $Fw.d$  editing (13.7.2.3.2).
- 30 3 For an internal value that is an IEEE infinity or NaN, the form of the output field is the same as for  $Fw.d$ .
- 31 4 For an internal value that is neither an IEEE infinity nor a NaN, the form of the output field for a scale factor  
 32 of zero is  
 33 
$$[ \pm ] [0].x_1x_2 \dots x_d exp$$
  
 34 where:
- 35 •  $\pm$  signifies a plus sign or a minus sign;
  - 36 •  $.$  signifies a *decimal symbol* (13.6);
  - 37 •  $x_1x_2 \dots x_d$  are the  $d$  most significant digits of the internal value after rounding (13.7.2.3.8);
  - 38 •  $exp$  is a decimal exponent having one of the forms specified in table 13.1.

Table 13.1: **E and D exponent forms**

Edit Descriptor	Absolute Value of Exponent	Form of Exponent <sup>1</sup>
<i>Ew.d</i>	$ exp  \leq 99$	$E\pm z_1 z_2$ or $\pm 0z_1 z_2$
	$99 <  exp  \leq 999$	$\pm z_1 z_2 z_3$
<i>Ew.d Ee</i> with $e > 0$	$ exp  \leq 10^e - 1$	$E\pm z_1 z_2 \dots z_e$
<i>Ew.d E0</i>	any	$E\pm z_1 z_2 \dots z_s$
<i>Dw.d</i>	$ exp  \leq 99$	$D\pm z_1 z_2$ or $E\pm z_1 z_2$ or $\pm 0z_1 z_2$
	$99 <  exp  \leq 999$	$\pm z_1 z_2 z_3$

(1) where each  $z$  is a digit, and  $s$  is the minimum number of digits required to represent the exponent.

- 1 5 The sign in the exponent is produced. A plus sign is produced if the exponent value is zero.
- 2 6 The scale factor  $k$  controls the decimal normalization (13.3.2, 13.8.5). If  $-d < k \leq 0$ , the output field contains  
3 exactly  $|k|$  leading zeros and  $d - |k|$  significant digits after the decimal symbol. If  $0 < k < d + 2$ , the output field  
4 contains exactly  $k$  significant digits to the left of the decimal symbol and  $d - k + 1$  significant digits to the right  
5 of the decimal symbol. Other values of  $k$  are not permitted.

#### 6 13.7.2.3.4 EN editing

- 7 1 The EN edit descriptor produces an output field in the form of a real number in engineering notation such that  
8 the decimal exponent is divisible by three and the absolute value of the significand (R715) is greater than or  
9 equal to 1 and less than 1000, except when the output value is zero. The scale factor has no effect on output.
- 10 2 The forms of the edit descriptor are *ENw.d* and *ENw.d Ee* indicating that the external field occupies  $w$  positions,  
11 except when  $w$  is zero in which case the processor selects the field width. The fractional part of the field contains  
12  $d$  digits. If  $e$  is positive the exponent part contains  $e$  digits, otherwise it contains the minimum number of digits  
13 required to represent the exponent value.
- 14 3 The form and interpretation of the input field is the same as for *Fw.d* editing (13.7.2.3.2).
- 15 4 For an internal value that is an IEEE infinity or NaN, the form of the output field is the same as for *Fw.d*.
- 16 5 For an internal value that is neither an IEEE infinity nor a NaN, the form of the output field is  
17  $[\pm] yyy . x_1 x_2 \dots x_d exp$   
18 where:
- 19 •  $\pm$  signifies a plus sign or a minus sign;
  - 20 •  $yyy$  are the 1 to 3 decimal digits representative of the most significant digits of the internal value after  
21 rounding (13.7.2.3.8);
  - 22 •  $yyy$  is an integer such that  $1 \leq yyy < 1000$  or, if the output value is zero,  $yyy = 0$ ;
  - 23 •  $.$  signifies a decimal symbol (13.6);
  - 24 •  $x_1 x_2 \dots x_d$  are the  $d$  next most significant digits of the internal value after rounding;
  - 25 •  $exp$  is a decimal exponent, divisible by three, having one of the forms specified in table 13.2.



Table 13.2: EN exponent forms

Edit Descriptor	Absolute Value of Exponent	Form of Exponent <sup>1</sup>
EN <i>w.d</i>	$ exp  \leq 99$	E± <i>z</i> <sub>1</sub> <i>z</i> <sub>2</sub> or ±0 <i>z</i> <sub>1</sub> <i>z</i> <sub>2</sub>
	$99 <  exp  \leq 999$	± <i>z</i> <sub>1</sub> <i>z</i> <sub>2</sub> <i>z</i> <sub>3</sub>
EN <i>w.d</i> E <i>e</i> with <i>e</i> > 0	$ exp  \leq 10^e - 1$	E± <i>z</i> <sub>1</sub> <i>z</i> <sub>2</sub> ... <i>z</i> <sub><i>e</i></sub>
EN <i>w.d</i> E0	any	E± <i>z</i> <sub>1</sub> <i>z</i> <sub>2</sub> ... <i>z</i> <sub><i>s</i></sub>

(1) where each *z* is a digit, and *s* is the minimum number of digits required to represent the exponent.

1 6 The sign in the exponent is produced. A plus sign is produced if the exponent value is zero.

**NOTE 13.12**

Examples:

Internal Value	Output field Using SS, EN12.3
6.421	6.421E+00
-.5	-500.000E-03
.00217	2.170E-03
4721.3	4.721E+03

2 **13.7.2.3.5 ES editing**

3 1 The ES edit descriptor produces an output field in the form of a real number in scientific notation such that the  
4 absolute value of the significand (R715) is greater than or equal to 1 and less than 10, except when the output  
5 value is zero. The scale factor has no effect on output.

6 2 The forms of the edit descriptor are ES*w.d* and ES*w.d* E*e* indicating that the external field occupies *w* positions,  
7 except when *w* is zero in which case the processor selects the field width. The fractional part of the field contains  
8 *d* digits. If *e* is positive the exponent part contains *e* digits, otherwise it contains the minimum number of digits  
9 required to represent the exponent value.

10 3 The form and interpretation of the input field is the same as for F*w.d* editing (13.7.2.3.2).

11 4 For an internal value that is an IEEE infinity or NaN, the form of the output field is the same as for F*w.d*.

12 5 For an internal value that is neither an IEEE infinity nor a NaN, the form of the output field is

$$[\pm] y . x_1 x_2 \dots x_d exp$$

13 where:

- 14 • ± signifies a plus sign or a minus sign;
- 15 • *y* is a decimal digit representative of the most significant digit of the internal value after rounding (13.7.2.3.8);
- 16 • . signifies a decimal symbol (13.6);
- 17 • *x*<sub>1</sub>*x*<sub>2</sub>...*x*<sub>*d*</sub> are the *d* next most significant digits of the internal value after rounding;
- 18 • *exp* is a decimal exponent having one of the forms specified in table 13.3.

Table 13.3: ES exponent forms

Edit Descriptor	Absolute Value of Exponent	Form of Exponent <sup>1</sup>
ES <i>w.d</i>	$ exp  \leq 99$	E± <i>z</i> <sub>1</sub> <i>z</i> <sub>2</sub> or ±0 <i>z</i> <sub>1</sub> <i>z</i> <sub>2</sub>
	$99 <  exp  \leq 999$	± <i>z</i> <sub>1</sub> <i>z</i> <sub>2</sub> <i>z</i> <sub>3</sub>
ES <i>w.d</i> E <i>e</i> with <i>e</i> > 0	$ exp  \leq 10^e - 1$	E± <i>z</i> <sub>1</sub> <i>z</i> <sub>2</sub> ... <i>z</i> <sub><i>e</i></sub>
ES <i>w.d</i> E0	any	E± <i>z</i> <sub>1</sub> <i>z</i> <sub>2</sub> ... <i>z</i> <sub><i>s</i></sub>

(1) where each *z* is a digit, and *s* is the minimum number of digits required to represent the exponent.

1 6 The sign in the exponent is produced. A plus sign is produced if the exponent value is zero.

**NOTE 13.13**

Examples:

Internal Value	Output field Using SS, ES12.3
6.421	6.421E+00
-.5	-5.000E-01
.00217	2.170E-03
4721.3	4.721E+03

2 **13.7.2.3.6 EX editing**

3 1 The EX edit descriptor produces an output field in the form of a hexadecimal-significand number.

4 2 The EX $w.d$  and EX $w.dEe$  edit descriptors indicate that the external field occupies  $w$  positions, except when  $w$   
5 is zero in which case the processor selects the field width. The fractional part of the field contains  $d$  hexadecimal  
6 digits, except when  $d$  is zero in which case the processor selects the number of hexadecimal digits to be the  
7 minimum required so that the output field is equal to the internal value;  $d$  shall not be zero if the radix of the  
8 internal value is not a power of two. The hexadecimal point, represented by a decimal symbol, appears after  
9 the first hexadecimal digit. For the form EX $w.d$ , and for EX $w.dE0$ , the exponent part contains the minimum  
10 number of digits needed to represent the exponent; otherwise the exponent contains  $e$  digits. The  $e$  has no effect  
11 on input. The scale factor has no effect on output.

12 3 The form and interpretation of the input field is the same as for F $w.d$  editing (13.7.2.3.2).

13 4 For an internal value that is an IEEE infinity or NaN, the form of the output field is the same as for F $w.d$ .

14 5 For an internal value that is neither an IEEE infinity nor a NaN, the form of the output field is

15  $[ \pm ] 0X x_0 . x_1x_2 \dots exp$

16 where:

- 17 •  $\pm$  signifies a plus sign or a minus sign;
- 18 •  $.$  signifies a [decimal symbol \(13.6\)](#);
- 19 •  $x_0x_1x_2 \dots$  are the most significant hexadecimal digits of the internal value, after rounding if  $d$  is not zero  
20 ([13.7.2.3.8](#));
- 21 •  $exp$  is a binary exponent expressed as a decimal integer; for EX $w.d$  and EX $w.dE0$ , the form is  $P \pm z_1 \dots z_n$ ,  
22 where  $n$  is the minimum number of digits needed to represent  $exp$ , and for EX $w.dEe$  with  $e$  greater than  
23 zero the form is  $P \pm z_1 \dots z_e$ . The choice of binary exponent is processor dependent. If the most significant  
24 binary digits of the internal value are  $b_0b_1b_2 \dots$ , the binary exponent might make the value of  $x_0$  be that of  
25  $b_0, b_0b_1, b_0b_1b_2$ , or  $b_0b_1b_2b_3$ .

26 6 The sign in the exponent is produced. A plus sign is produced if the exponent value is zero.

**NOTE 13.14**

Examples:

Internal value	Edit descriptor	Possible output with SS in effect
1.375	EX0.1	0X1.6P+0
-15.625	EX14.4E3	-0X1.F400P+003
1048580.0	EX0.0	0X1.00004P+20
2.375	EX0.1	0X2.6P+0

27 **13.7.2.3.7 Complex editing**

28 1 A complex datum consists of a pair of separate real data. The editing of a scalar datum of complex type is  
29 specified by two edit descriptors each of which specifies the editing of real data. The first edit descriptor specifies

1 the editing for the real part; the second specifies it for the imaginary part. The two edit descriptors may be  
2 different. Control and character string edit descriptors may be processed between the edit descriptor for the real  
3 part and the edit descriptor for the imaginary part.

#### 4 13.7.2.3.8 Input/output rounding mode

5 1 The input/output rounding mode can be specified by an [OPEN statement \(12.5.2\)](#), a [data transfer statement](#)  
6 [\(12.6.2.13\)](#), or an edit descriptor [\(13.8.7\)](#).

7 2 In what follows, the term “decimal value” means the exact decimal number as given by the character string, while  
8 the term “internal value” means the number actually stored in the processor. For example, in dealing with the  
9 decimal constant 0.1, the decimal value is the mathematical quantity 1/10, which has no exact representation  
10 in binary form. Formatted output of real data involves conversion from an internal value to a decimal value;  
11 formatted input involves conversion from a decimal value to an internal value.

12 3 When the input/output rounding mode is UP, the value resulting from conversion shall be the smallest represent-  
13 able value that is greater than or equal to the original value. When the input/output rounding mode is DOWN,  
14 the value resulting from conversion shall be the largest representable value that is less than or equal to the original  
15 value. When the input/output rounding mode is ZERO, the value resulting from conversion shall be the value  
16 closest to the original value and no greater in magnitude than the original value. When the input/output rounding  
17 mode is NEAREST, the value resulting from conversion shall be the closer of the two nearest representable values  
18 if one is closer than the other. If the two nearest representable values are equidistant from the original value, it is  
19 processor dependent which one of them is chosen. When the input/output rounding mode is COMPATIBLE, the  
20 value resulting from conversion shall be the closer of the two nearest representable values or the value away from  
21 zero if halfway between them. When the input/output rounding mode is PROCESSOR\_DEFINED, rounding  
22 during conversion shall be a processor-dependent default mode, which may correspond to one of the other modes.

23 4 On processors that support IEEE rounding on conversions [\(17.4\)](#), NEAREST shall correspond to round to nearest,  
24 as specified in ISO/IEC/IEEE 60559:2011.

#### NOTE 13.15

<p>On processors that support IEEE rounding on conversions, the input/output rounding modes COMPATIBLE and NEAREST will produce the same results except when the datum is halfway between the two representable values. In that case, NEAREST will pick the even value, but COMPATIBLE will pick the value away from zero. The input/output rounding modes UP, DOWN, and ZERO have the same effect as those specified in ISO/IEC/IEEE 60559:2011 for round toward <math>+\infty</math>, round toward <math>-\infty</math>, and round toward 0, respectively.</p>
--

#### 25 13.7.2.4 B, O, and Z editing

26 1 The *Bw*, *Bw.m*, *Ow*, *Ow.m*, *Zw*, and *Zw.m* edit descriptors indicate that the field to be edited occupies *w*  
27 positions, except when *w* is zero. When *w* is zero, the processor selects the field width. On input, *w* shall not be  
28 zero. The corresponding input/output list item shall be of type integer, real, or complex.

29 2 On input, *m* has no effect.

30 3 In the standard form of the input field for the B, O, and Z edit descriptors the character string consists of binary,  
31 octal, or hexadecimal digits (as in [R765](#), [R766](#), [R767](#)) in the respective input field. The lower-case hexadecimal  
32 digits a through f in a hexadecimal input field are equivalent to the corresponding upper-case hexadecimal digits.  
33 If the input field does not have the standard form, and is not acceptable to the processor, an error condition  
34 occurs.

35 4 The value is INT (X) if the input list item is of type integer and REAL (X) if the input list item is of type real  
36 or complex, where X is a *boz-literal-constant* that specifies the same bit sequence as the digits of the input field.

37 5 The output field for the *Bw*, *Ow*, and *Zw* descriptors consists of zero or more leading blanks followed by the  
38 internal value in a form identical to the digits of a binary, octal, or hexadecimal constant, respectively, that

1 specifies the same bit sequence but without leading zero bits.

**NOTE 13.16**

A binary, octal, or hexadecimal constant always consists of at least one digit or hexadecimal digit.

2 R1322 *hex-digit-string* is *hex-digit* [ *hex-digit* ] ...

3 6 The output field for the *Bw.m*, *Ow.m*, and *Zw.m* edit descriptor is the same as for the *Bw*, *Ow*, and *Zw* edit  
4 descriptor, except that the *digit-string* or *hex-digit-string* consists of at least *m* digits. If necessary, sufficient  
5 leading zeros are included to achieve the minimum of *m* digits. The value of *m* shall not exceed the value of *w*,  
6 except when *w* is zero. If *m* is zero and the internal value consists of all zero bits, the output field consists of  
7 only blank characters. When *m* and *w* are both zero, and the internal value consists of all zero bits, one blank  
8 character is produced.

9 **13.7.3 Logical editing**

- 10 1 The *Lw* edit descriptor indicates that the field occupies *w* positions. The specified input/output list item shall  
11 be of type logical. The *G* edit descriptor also may be used to edit logical data (13.7.5.3).
- 12 2 The standard form of the input field consists of optional blanks, optionally followed by a period, followed by a *T*  
13 for true or *F* for false. The *T* or *F* may be followed by additional characters in the field, which are ignored. If the  
14 input field does not have the standard form, and is not acceptable to the processor, an error condition occurs.
- 15 3 A lower-case letter is equivalent to the corresponding upper-case letter in a logical input field.

**NOTE 13.17**

The logical constants *.TRUE.* and *.FALSE.* are acceptable input forms.

- 16 4 The output field consists of *w*−1 blanks followed by a *T* or *F*, depending on whether the internal value is true or  
17 false, respectively.

18 **13.7.4 Character editing**

- 19 1 The *A[w]* edit descriptor is used with an input/output list item of type character. The *G* edit descriptor also may  
20 be used to edit character data (13.7.5.4). The kind type parameter of all characters transferred and converted  
21 under control of one *A* or *G* edit descriptor is implied by the kind of the corresponding list item.
- 22 2 If a field width *w* is specified with the *A* edit descriptor, the field consists of *w* characters. If a field width *w* is  
23 not specified with the *A* edit descriptor, the number of characters in the field is the length of the corresponding  
24 list item, regardless of the value of the kind type parameter.
- 25 3 Let *len* be the length of the input/output list item. If the specified field width *w* for an *A* edit descriptor  
26 corresponding to an input item is greater than or equal to *len*, the rightmost *len* characters will be taken from the  
27 input field. If the specified field width *w* is less than *len*, the *w* characters will appear left justified with *len*−*w*  
28 trailing blanks in the internal value.
- 29 4 If the specified field width *w* for an *A* edit descriptor corresponding to an output item is greater than *len*, the  
30 output field will consist of *w*−*len* blanks followed by the *len* characters from the internal value. If the specified  
31 field width *w* is less than or equal to *len*, the output field will consist of the leftmost *w* characters from the internal  
32 value.

**NOTE 13.18**

For nondefault character kinds, the blank padding character is processor dependent.

- 33 5 If the file is connected for stream access, the output may be split across more than one record if it contains  
34 newline characters. A newline character is a nonblank character returned by the intrinsic function *NEW\_LINE*.

1 Beginning with the first character of the output field, each character that is not a newline is written to the current  
 2 record in successive positions; each newline character causes file positioning at that point as if by slash editing  
 3 (the current record is terminated at that point, a new empty record is created following the current record, this  
 4 new record becomes the last and current record of the file, and the file is positioned at the beginning of this new  
 5 record).

**NOTE 13.19**

If the intrinsic function `NEW_LINE` returns a blank character for a particular character kind, then the processor does not support using a character of that kind to cause record termination in a formatted stream file.

6 **13.7.5 Generalized editing**

7 **13.7.5.1 Overview**

8 1 The `Gw`, `Gw.d` and `Gw.d Ee` edit descriptors are used with an input/output list item of any intrinsic type. When  
 9 `w` is nonzero, these edit descriptors indicate that the external field occupies `w` positions. For real or complex  
 10 data the fractional part consists of a maximum of `d` digits and the exponent part consists of `e` digits. When these  
 11 edit descriptors are used to specify the input/output of integer, logical, or character data, `d` and `e` have no effect.  
 12 When `w` is zero the processor selects the field width. On input, `w` shall not be zero.

13 **13.7.5.2 Generalized numeric editing**

14 1 When used to specify the input/output of integer, real, and complex data, the `Gw`, `Gw.d` and `Gw.d Ee` edit  
 15 descriptors follow the general rules for numeric editing (13.7.2).

**NOTE 13.20**

The `Gw.d Ee` edit descriptor follows any additional rules for the `Ew.d Ee` edit descriptor.

16 **13.7.5.2.1 Generalized integer editing**

17 1 When used to specify the input/output of integer data, the `Gw.d` and `Gw.d Ee` edit descriptors follow the rules  
 18 for the `Iw` edit descriptor (13.7.2.2), except that `w` shall not be zero. When used to specify the output of integer  
 19 data, the `G0` and `G0.d` edit descriptors follow the rules for the `I0` edit descriptor.

20 **13.7.5.2.2 Generalized real and complex editing**

21 1 The form and interpretation of the input field is the same as for `Fw.d` editing (13.7.2.3.2).  
 22 2 If `d` is zero, `kPEw.0` or `kPEw.0Ee` editing is used for `Gw.0` editing or `Gw.0Ee` editing respectively.  
 23 3 When used to specify the output of real or complex data that is not an IEEE infinity or NaN, the `G0` and `G0.d`  
 24 edit descriptors follow the rules for the `Gw.dEe` edit descriptor, except that any leading or trailing blanks are  
 25 removed. Reasonable processor-dependent values of `w`, `d` (if not specified), and `e` are used with each output value.  
 26 4 For an internal value that is an IEEE infinity or NaN, the form of the output field for the `Gw.d` and `Gw.d Ee`  
 27 edit descriptors is the same as for `Fw.d`, and the form of the output field for the `G0` and `G0.d` edit descriptors is  
 28 the same as for `F0.0`.  
 29 5 Otherwise, the method of representation in the output field depends on the magnitude of the internal value  
 30 being edited. If the internal value is zero, let `s` be one. If the internal value is a number other than zero, let `N`  
 31 be the decimal value that is the result of converting the internal value to `d` significant digits according to the  
 32 input/output rounding mode and let `s` be the integer such that  $10^{s-1} \leq |N| < 10^s$ . If  $s < 0$  or  $s > d$ , `kPEw.d` or  
 33 `kPEw.dEe` editing is used for `Gw.d` editing or `Gw.dEe` editing respectively, where `k` is the scale factor (13.8.5).  
 34 If  $0 \leq s \leq d$ , the scale factor has no effect and `F(w-n).(d-s),n('b')` editing is used where `b` is a blank and `n` is  
 35 4 for `Gw.d` editing, `e+2` for `Gw.dEe` editing if `e > 0`, and 4 for `Gw.dE0` editing.

- 1 6 The value of  $w-n$  shall be positive.

**NOTE 13.21**

The scale factor has no effect on output unless the magnitude of the datum to be edited is outside the range that permits effective use of F editing.

2 **13.7.5.3 Generalized logical editing**

- 3 1 When used to specify the input/output of logical data, the  $Gw.d$  and  $Gw.d Ee$  edit descriptors follow the rules for  
4 the  $Lw$  edit descriptor (13.7.3). When used to specify the output of logical data, the  $G0$  and  $G0.d$  edit descriptors  
5 follow the rules for the  $L1$  edit descriptor.

6 **13.7.5.4 Generalized character editing**

- 7 1 When used to specify the input/output of character data, the  $Gw.d$  and  $Gw.d Ee$  edit descriptors follow the rules  
8 for the  $Aw$  edit descriptor (13.7.4). When used to specify the output of character data, the  $G0$  and  $G0.d$  edit  
9 descriptors follow the rules for the  $A$  edit descriptor with no field width.

10 **13.7.6 User-defined derived-type editing**

- 11 1 The DT edit descriptor specifies that a user-provided procedure shall be used instead of the processor's default  
12 input/output formatting for processing a list item of derived type.
- 13 2 The DT edit descriptor may include a character literal constant. The character value "DT" concatenated with the  
14 character literal constant is passed to the [defined input/output](#) procedure as the `iotype` argument (12.6.4.8). The  
15  $v$  values of the edit descriptor are passed to the [defined input/output](#) procedure as the `v_list` array argument.

**NOTE 13.22**

For the edit descriptor DT'Link List'(10, 4, 2), `iotype` is "DTLink List" and `v_list` is [10, 4, 2].

- 16 3 If a derived-type variable or value corresponds to a DT edit descriptor, there shall be an accessible [interface](#) to  
17 a corresponding [defined input/output](#) procedure for that derived type (12.6.4.8). A DT edit descriptor shall not  
18 correspond to a list item that is not of a derived type.

19 **13.8 Control edit descriptors**

20 **13.8.1 Position edit descriptors**

21 **13.8.1.1 Position editing**

- 22 1 The position edit descriptors T, TL, TR, and X, specify the position at which the next character will be transmit-  
23 ted to or from the record. If any character skipped by a position edit descriptor is of type nondefault character,  
24 and the `unit` is a default character [internal file](#) or an [external](#) non-Unicode file, the result of that position editing  
25 is processor dependent.
- 26 2 On input, if the position specified by a position edit descriptor is before the current position, portions of a record  
27 can be processed more than once, possibly with different editing.
- 28 3 On input, a position beyond the last character of the record may be specified if no characters are transmitted  
29 from such positions.
- 30 4 On output, a position edit descriptor does not by itself cause characters to be transmitted and therefore does not  
31 by itself affect the length of the record. If characters are transmitted to positions at or after the position specified  
32 by a position edit descriptor, positions skipped and not previously filled are filled with blanks. The result is as if  
33 the entire record were initially filled with blanks.

1 5 On output, a character in the record can be replaced. A position edit descriptor never directly causes a character  
2 already placed in the record to be replaced, but it might result in positioning such that subsequent editing causes  
3 a replacement.

#### 4 13.8.1.2 T, TL, and TR editing

5 1 The left tab limit affects file positioning by the T and TL edit descriptors. Immediately prior to nonchild data  
6 transfer (12.6.4.8.3), the left tab limit becomes defined as the character position of the current record or the  
7 current position of the stream file. If, during data transfer, the file is positioned to another record, the left tab  
8 limit becomes defined as character position one of that record.

9 2 The T $n$  edit descriptor indicates that the transmission of the next character to or from a record is to occur at  
10 the  $n$ th character position of the record, relative to the left tab limit. This position can be in either direction  
11 from the current position.

12 3 The TL $n$  edit descriptor indicates that the transmission of the next character to or from the record is to occur at  
13 the character position  $n$  characters backward from the current position. However, if  $n$  is greater than the difference  
14 between the current position and the left tab limit, the TL $n$  edit descriptor indicates that the transmission of  
15 the next character to or from the record is to occur at the left tab limit.

16 4 The TR $n$  edit descriptor indicates that the transmission of the next character to or from the record is to occur  
17 at the character position  $n$  characters forward from the current position.

#### 18 13.8.1.3 X editing

19 1 The  $n$ X edit descriptor indicates that the transmission of the next character to or from a record is to occur at  
20 the character position  $n$  characters forward from the current position.

##### NOTE 13.23

An  $n$ X edit descriptor has the same effect as a TR $n$  edit descriptor.

#### 21 13.8.2 Slash editing

22 1 The slash edit descriptor indicates the end of data transfer to or from the current record.

23 2 On input from a file connected for sequential or stream access, the remaining portion of the current record is  
24 skipped and the file is positioned at the beginning of the next record. This record becomes the current record.  
25 On output to a file connected for sequential or stream access, a new empty record is created following the current  
26 record; this new record then becomes the last and current record of the file and the file is positioned at the  
27 beginning of this new record.

28 3 For a file connected for direct access, the record number is increased by one and the file is positioned at the  
29 beginning of the record that has that record number, if there is such a record, and this record becomes the  
30 current record.

##### NOTE 13.24

A record that contains no characters can be written on output. If the file is an [internal file](#) or a file connected  
for direct access, the record is filled with blank characters.

An entire record can be skipped on input.

31 4 The repeat specification is optional in the slash edit descriptor. If it is not specified, the default value is one.

#### 32 13.8.3 Colon editing

33 1 The colon edit descriptor terminates format control if there are no more [effective items](#) in the input/output list  
34 (12.6.3). The colon edit descriptor has no effect if there are more [effective items](#) in the input/output list.

### 13.8.4 SS, SP, and S editing

- 1 The SS, SP, and S edit descriptors temporarily change (12.5.2) the sign mode (12.5.6.17, 12.6.2.14) for the connection. The edit descriptors SS, SP, and S set the sign mode corresponding to the SIGN= specifier values SUPPRESS, PLUS, and PROCESSOR\_DEFINED, respectively.
- 2 The sign mode controls optional plus characters in numeric output fields. When the sign mode is PLUS, the processor shall produce a plus sign in any position that normally contains an optional plus sign. When the sign mode is SUPPRESS, the processor shall not produce a plus sign in such positions. When the sign mode is PROCESSOR\_DEFINED, the processor has the option of producing a plus sign or not in such positions, subject to 13.7.2(5).
- 3 The SS, SP, and S edit descriptors affect only I, F, E, EN, ES, D, and G editing during the execution of an output statement. The SS, SP, and S edit descriptors have no effect during the execution of an input statement.

### 13.8.5 P editing

- 1 The  $k$ P edit descriptor temporarily changes (12.5.2) the scale factor for the connection to  $k$ . The scale factor affects the editing done by the F, E, EN, ES, EX, D, and G edit descriptors for numeric quantities.
- 2 The scale factor  $k$  affects the appropriate editing in the following manner.
  - On input, with F, E, EN, ES, EX, D, and G editing (provided that no exponent exists in the field), the effect is that the externally represented number equals the internally represented number multiplied by  $10^k$ ; the scale factor is applied to the external decimal value and then this is converted using the input/output rounding mode.
  - On input, with F, E, EN, ES, EX, D, and G editing, the scale factor has no effect if there is an exponent in the field.
  - On output, with F output editing, the effect is that the externally represented number equals the internally represented number multiplied by  $10^k$ ; the internal value is converted using the input/output rounding mode and then the scale factor is applied to the converted decimal value.
  - On output, with E and D editing, the effect is that the significand (R715) part of the quantity to be produced is multiplied by  $10^k$  and the exponent is reduced by  $k$ .
  - On output, with G editing, the effect is suspended unless the magnitude of the datum to be edited is outside the range that permits the use of F editing. If the use of E editing is required, the scale factor has the same effect as with E output editing.
  - On output, with EN, ES, and EX editing, the scale factor has no effect.

### 13.8.6 BN and BZ editing

- 1 The BN and BZ edit descriptors temporarily change (12.5.2) the blank interpretation mode (12.5.6.6, 12.6.2.6) for the connection. The edit descriptors BN and BZ set the blank interpretation mode corresponding to the BLANK= specifier values NULL and ZERO, respectively.
- 2 The blank interpretation mode controls the interpretation of nonleading blanks in numeric input fields. Such blank characters are interpreted as zeros when the blank interpretation mode has the value ZERO; they are ignored when the blank interpretation mode has the value NULL. The effect of ignoring blanks is to treat the input field as if blanks had been removed, the remaining portion of the field right justified, and the blanks replaced as leading blanks. However, a field containing only blanks has the value zero.
- 3 The blank interpretation mode affects only numeric editing (13.7.2) and generalized numeric editing (13.7.5.2) on input. It has no effect on output.

### 13.8.7 RU, RD, RZ, RN, RC, and RP editing

- 1 The round edit descriptors temporarily change (12.5.2) the connection's input/output rounding mode (12.5.6.16, 12.6.2.13, 13.7.2.3.8). The round edit descriptors RU, RD, RZ, RN, RC, and RP set the input/output rounding



1 mode corresponding to the `ROUND= specifier` values UP, DOWN, ZERO, NEAREST, COMPATIBLE, and  
2 `PROCESSOR_DEFINED`, respectively. The input/output rounding mode affects the conversion of real and  
3 complex values in formatted input/output. It affects only D, E, EN, ES, F, and G editing.

### 4 **13.8.8 DC and DP editing**

5 1 The decimal edit descriptors temporarily change (12.5.2) the decimal edit mode (12.5.6.7, 12.6.2.7, 13.6) for  
6 the connection. The edit descriptors DC and DP set the decimal edit mode corresponding to the `DECIMAL=`  
7 `specifier` values `COMMA` and `POINT`, respectively.

8 2 The decimal edit mode controls the representation of the `decimal symbol` (13.6) during conversion of real and  
9 complex values in formatted input/output. The decimal edit mode affects only D, E, EN, ES, F, and G editing.

## 10 **13.9 Character string edit descriptors**

11 1 A character string edit descriptor shall not be used on input.

12 2 The character string edit descriptor causes characters to be written from the enclosed characters of the edit  
13 descriptor itself, including blanks. For a character string edit descriptor, the width of the field is the number of  
14 characters between the delimiting characters. Within the field, two consecutive delimiting characters are counted  
15 as a single character.

### NOTE 13.25

A delimiter for a character string edit descriptor is either an apostrophe or quote.

## 16 **13.10 List-directed formatting**

### 17 **13.10.1 Purpose of list-directed formatting**

18 1 List-directed input/output allows data editing according to the type of the list item instead of by a format  
19 specification. It also allows data to be free-field, that is, separated by commas (or semicolons) or blanks.

### 20 **13.10.2 Values and value separators**

21 1 The characters in one or more list-directed records constitute a sequence of values and value separators. The end  
22 of a record has the same effect as a blank character, unless it is within a character constant. Any sequence of two  
23 or more consecutive blanks is treated as a single blank, unless it is within a character constant.

24 2 Each value is either a null value,  $c$ ,  $r^*c$ , or  $r^*$ , where  $c$  is a literal constant, optionally signed if integer or real,  
25 or an undelimited character constant and  $r$  is an unsigned, nonzero, integer literal constant. Neither  $c$  nor  $r$   
26 shall have kind type parameters specified. The constant  $c$  is interpreted as though it had the same kind type  
27 parameter as the corresponding list item. The  $r^*c$  form is equivalent to  $r$  successive appearances of the constant  
28  $c$ , and the  $r^*$  form is equivalent to  $r$  successive appearances of the null value. Neither of these forms may contain  
29 embedded blanks, except where permitted within the constant  $c$ .

30 3 A value separator is

- 31 • a comma optionally preceded by one or more contiguous blanks and optionally followed by one or more  
32 contiguous blanks, unless the decimal edit mode is `COMMA`, in which case a semicolon is used in place of  
33 the comma,
- 34 • a slash optionally preceded by one or more contiguous blanks and optionally followed by one or more  
35 contiguous blanks, or
- 36 • one or more contiguous blanks between two nonblank values or following the last nonblank value, where a  
37 nonblank value is a constant, an  $r^*c$  form, or an  $r^*$  form.

**NOTE 13.26**

Although a slash encountered in an input record is referred to as a separator, it actually causes termination of list-directed and namelist [input statements](#); it does not actually separate two values.

**NOTE 13.27**

If no list items are specified in a list-directed input/output statement, one input record is skipped or one empty output record is written.

**13.10.3 List-directed input****13.10.3.1 List-directed input forms**

- 1 Input forms acceptable to edit descriptors for a given type are acceptable for list-directed formatting, except as noted below. If the form of the input value is not acceptable to the processor for the type of the next [effective item](#) in the list, an error condition occurs. Blanks are never used as zeros, and embedded blanks are not permitted in constants, except within character constants and complex constants as specified below.
- 2 For the  $r^*c$  form of an input value, the constant  $c$  is interpreted as an un delimited character constant if the first list item corresponding to this value is default, [ASCII](#), or [ISO 10646 character](#), there is a nonblank character immediately after  $r^*$ , and that character is not an apostrophe or a quotation mark; otherwise,  $c$  is interpreted as a literal constant.

**NOTE 13.28**

The end of a record has the effect of a blank, except when it appears within a character constant.

- 3 When the next [effective item](#) is of type integer, the value in the input record is interpreted as if an  $Iw$  edit descriptor with a suitable value of  $w$  were used.
- 4 When the next [effective item](#) is of type real, the input form is that of a numeric input field. A numeric input field is a field suitable for F editing ([13.7.2.3.2](#)) that is assumed to have no fractional digits unless a [decimal symbol](#) appears within the field.
- 5 When the next [effective item](#) is of type complex, the input form consists of a left parenthesis followed by an ordered pair of numeric input fields separated by a comma (if the decimal edit mode is POINT) or semicolon (if the decimal edit mode is COMMA), and followed by a right parenthesis. The first numeric input field is the real part of the complex constant and the second is the imaginary part. Each of the numeric input fields may be preceded or followed by any number of blanks and ends of records. The end of a record may occur after the real part or before the imaginary part.
- 6 When the next [effective item](#) is of type logical, the input form shall not include value separators among the optional characters permitted for L editing.
- 7 When the next [effective item](#) is of type character, the input form consists of a possibly delimited sequence of zero or more *rep-chars* whose kind type parameter is implied by the kind of the [effective item](#). Character sequences may be continued from the end of one record to the beginning of the next record, but the end of record shall not occur between a doubled apostrophe in an apostrophe-delimited character sequence, nor between a doubled quote in a quote-delimited character sequence. The end of the record does not cause a blank or any other character to become part of the character sequence. The character sequence may be continued on as many records as needed. The characters blank, comma, semicolon, and slash may appear in default, [ASCII](#), or [ISO 10646 character](#) sequences.
- 8 If the next [effective item](#) is default, [ASCII](#), or [ISO 10646 character](#) and
  - the character sequence does not contain value separators,
  - the character sequence does not cross a record boundary,
  - the first nonblank character is not a quotation mark or an apostrophe,

- 1           • the leading characters are not *digits* followed by an asterisk, and  
 2           • the character sequence contains at least one character,

3 the delimiting apostrophes or quotation marks are not required. If the delimiters are omitted, the character  
 4 sequence is terminated by the first blank, comma (if the decimal edit mode is POINT), semicolon (if the decimal  
 5 edit mode is COMMA), slash, or end of record; in this case apostrophes and quotation marks within the datum  
 6 are not to be doubled.

- 7 9 Let *len* be the current length of the next **effective item**, and let *w* be the length of the character sequence. If *len*  
 8 is less than or equal to *w*, the leftmost *len* characters of the sequence are transmitted to the next **effective item**.  
 9 If *len* is greater than *w*, the sequence is transmitted to the leftmost *w* characters of the next **effective item** and  
 10 the remaining *len*−*w* characters of the next **effective item** are filled with blanks.

**NOTE 13.29**

An **allocatable**, **deferred-length** character **effective item** does not have its allocation status or allocated length changed as a result of list-directed input.

11 **13.10.3.2 Null values**

- 12 1 A null value is specified by

- 13           • the *r*\* form,  
 14           • no characters between consecutive value separators, or  
 15           • no characters before the first value separator in the first record read by each execution of a list-directed  
 16 **input statement**.

**NOTE 13.30**

The end of a record following any other value separator, with or without separating blanks, does not specify a null value in list-directed input.

- 17 2 A null value has no effect on the definition status of the next **effective item**. A null value shall not be used for  
 18 either the real or imaginary part of a complex constant, but a single null value may represent an entire complex  
 19 constant.

- 20 3 A slash encountered as a value separator during execution of a list-directed **input statement** causes termination  
 21 of execution of that **input statement** after the transference of the previous value. Any characters remaining in the  
 22 current record are ignored. If there are additional items in the input list, the effect is as if null values had been  
 23 supplied for them. Any *do-variable* in the input list becomes defined as if enough null values had been supplied  
 24 for any remaining input list items.

**NOTE 13.31**

All blanks encountered during list-directed input are considered to be part of some value separator except for

- blanks embedded in a character sequence,
- embedded blanks surrounding the real or imaginary part of a complex constant, and
- leading blanks in the first record read by each execution of a list-directed **input statement**, unless immediately followed by a slash or comma.

**NOTE 13.32**

List-directed input example:

```
INTEGER I; REAL X (8); CHARACTER (11) P;
```

**NOTE 13.32 (cont.)**

```

COMPLEX Z; LOGICAL G
...
READ *, I, X, P, Z, G
...

```

The input data records are:

```

12345,12345,,2*1.5,4*
ISN'T_BOB'S,(123,0),.TEXAS$

```

The results are:

Variable	Value
I	12345
X (1)	12345.0
X (2)	unchanged
X (3)	1.5
X (4)	1.5
X (5) – X (8)	unchanged
P	ISN'T_BOB'S
Z	(123.0,0.0)
G	true

#### 1 13.10.4 List-directed output

- 2 1 The form of the values produced is the same as that required for input, except as noted otherwise. With the  
3 exception of adjacent undelimited character sequences, the values are separated by one or more blanks or by a  
4 comma, or a semicolon if the decimal edit mode is COMMA, optionally preceded by one or more blanks and  
5 optionally followed by one or more blanks. Two undelimited character sequences are considered adjacent when  
6 both were written using list-directed input/output, no intervening [data transfer](#) or [file positioning](#) operations on  
7 that [unit](#) occurred, and both were written either by a single [data transfer statement](#), or during the execution of  
8 a parent [data transfer statement](#) along with its child [data transfer statements](#). The form of the values produced  
9 by [defined output](#) (12.6.4.8) is determined by the [defined output](#) procedure; this form need not be compatible  
10 with list-directed input.
- 11 2 The processor may begin new records as necessary, but the end of record shall not occur within a constant except  
12 as specified for complex constants and character sequences. The processor shall not insert blanks within character  
13 sequences or within constants, except as specified for complex constants.
- 14 3 Logical output values are T for the value true and F for the value false.
- 15 4 Integer output constants are produced with the effect of an *Iw* edit descriptor.
- 16 5 Real constants are produced with the effect of either an F edit descriptor or an E edit descriptor, depending on  
17 the magnitude  $x$  of the value and a range  $10^{d_1} \leq x < 10^{d_2}$ , where  $d_1$  and  $d_2$  are processor-dependent integers. If  
18 the magnitude  $x$  is within this range or is zero, the constant is produced using *0PFw.d*; otherwise, *1PEw.d Ee* is  
19 used.
- 20 6 For numeric output, reasonable processor-dependent values of *w*, *d*, and *e* are used for each of the numeric  
21 constants output.
- 22 7 Complex constants are enclosed in parentheses with a separator between the real and imaginary parts, each  
23 produced as defined above for real constants. The separator is a comma if the decimal edit mode is POINT; it  
24 is a semicolon if the decimal edit mode is COMMA. The end of a record may occur between the separator and  
25 the imaginary part only if the entire constant is as long as, or longer than, an entire record. The only embedded  
26 blanks permitted within a complex constant are between the separator and the end of a record and one blank at

- 1 the beginning of the next record.
- 2 8 Character sequences produced when the delimiter mode has a value of NONE
- 3 • are not delimited by apostrophes or quotation marks,
  - 4 • are not separated from each other by value separators,
  - 5 • have each internal apostrophe or quotation mark represented externally by one apostrophe or quotation
  - 6 mark, and
  - 7 • have a blank character inserted by the processor at the beginning of any record that begins with the
  - 8 continuation of a character sequence from the preceding record.
- 9 9 Character sequences produced when the delimiter mode has a value of QUOTE are delimited by quotes, are
- 10 preceded and followed by a value separator, and have each internal quote represented on the external medium by
- 11 two contiguous quotes.
- 12 10 Character sequences produced when the delimiter mode has a value of APOSTROPHE are delimited by apo-
- 13 strophes, are preceded and followed by a value separator, and have each internal apostrophe represented on the
- 14 external medium by two contiguous apostrophes.
- 15 11 If two or more successive values in an output record have identical values, the processor has the option of producing
- 16 a repeated constant of the form  $r*c$  instead of the sequence of identical values.
- 17 12 Slashes, as value separators, and null values are not produced as output by list-directed formatting.
- 18 13 Except for new records created by explicit formatting within a [defined output](#) procedure or by continuation of
- 19 delimited character sequences, each output record begins with a blank character.

**NOTE 13.33**

The length of the output records is not specified and is processor dependent.

## 20 13.11 Namelist formatting

### 21 13.11.1 Purpose of namelist formatting

- 22 1 Namelist input/output allows data editing with name-value subsequences. This facilitates documentation of input
- 23 and output files and more flexibility on input.

### 24 13.11.2 Name-value subsequences

- 25 1 The characters in one or more namelist records constitute a sequence of name-value subsequences, each of which
- 26 consists of an [object designator](#) followed by an equals and followed by one or more values and value separators.
- 27 The equals may optionally be preceded or followed by one or more contiguous blanks. The end of a record has the
- 28 same effect as a blank character, unless it is within a character constant. Any sequence of two or more consecutive
- 29 blanks is treated as a single blank, unless it is within a character constant.
- 30 2 Each object designator shall begin with a name from the [namelist-group-object-list](#) (8.9) and shall follow the
- 31 syntax of [designator](#) (R901). It shall not contain a [vector subscript](#) or an [image-selector](#) and shall not designate a
- 32 zero-sized array, a zero-sized array section, or a zero-length character string. Each subscript, stride, and substring
- 33 range expression shall be an optionally signed integer literal constant with no [kind type parameter](#) specified. If
- 34 a section subscript list appears, the number of section subscripts shall be equal to the rank of the object. If
- 35 the namelist group object is of derived type, the designator in the input record may be either the name of the
- 36 variable or the designator of one of its components, indicated by qualifying the variable name with the appropriate
- 37 component name. Successive qualifications may be applied as appropriate to the shape and type of the variable
- 38 represented. Each designator may be preceded and followed by one or more optional blanks but shall not contain
- 39 embedded blanks.

1 3 A value separator for namelist formatting is the same as for list-directed formatting (13.10.2), or one or more  
2 contiguous blanks between a nonblank value and the following *object designator* or namelist comment (13.11.3.6).

### 3 13.11.3 Namelist input

#### 4 13.11.3.1 Overall syntax

5 1 Input for a namelist *input statement* consists of

- 6 (1) optional blanks and namelist comments,
- 7 (2) the character & followed immediately by the *namelist-group-name* as specified in the *NAMELIST*  
8 *statement*,
- 9 (3) one or more blanks,
- 10 (4) a sequence of zero or more name-value subsequences separated by value separators, and
- 11 (5) a slash to terminate the namelist input.

#### NOTE 13.34

A slash encountered in a namelist input record causes the *input statement* to terminate. A slash cannot be used to separate two values in a namelist *input statement*.

12 2 The order of the name-value subsequences in the input records need not match the order of the *namelist-group-*  
13 *object-list*. The input records need not specify all objects in the *namelist-group-object-list*. They may specify a  
14 part of an object more than once.

15 3 A group name or object name is without regard to case.

#### 16 13.11.3.2 Namelist input processing

17 1 The name-value subsequences are evaluated serially, in left-to-right order. A namelist group object *designator*  
18 may appear in more than one name-value subsequence. The definition status of an object that is not a subobject  
19 of a *designator* in any name-value subsequence remains unchanged.

20 2 When the *designator* in the input record represents an array variable or a variable of derived type, the effect is  
21 as if the variable represented were expanded into a sequence of scalar list items, in the same way that formatted  
22 input/output list items are expanded (12.6.3). The number of values following the equals shall not exceed the  
23 number of list items in the expanded sequence, but may be less; in the latter case, the effect is as if sufficient  
24 null values had been appended to match any remaining list items in the expanded sequence. Except as noted  
25 elsewhere in this subclause, if an input value is not acceptable to the processor for the type of the list item in the  
26 corresponding position in the expanded sequence, an error condition occurs.

#### NOTE 13.35

For example, if the designator in the input record designates an integer array of size 100, at most 100 values, each of which is either a digit string or a null value, can follow the equals; these values would then be assigned to the elements of the array in array element order.

27 3 A slash encountered as a value separator during the execution of a namelist *input statement* causes termination  
28 of execution of that *input statement* after transference of the previous value. If there are additional items in the  
29 namelist group object being transferred, the effect is as if null values had been supplied for them.

30 4 Successive namelist records are read by namelist input until a slash is encountered; the remainder of the record  
31 is ignored.

32 5 A namelist comment may appear after any value separator except a slash (which terminates namelist input). A  
33 namelist comment is also permitted to start in the first nonblank position of an input record except within a  
34 character literal constant.

### 13.11.3.3 Namelist input values

- 1
- 2 1 Each value is either a null value (13.11.3.4),  $c$ ,  $r*c$ , or  $r*$ , where  $c$  is a literal constant, optionally signed if integer  
3 or real, and  $r$  is an unsigned, nonzero, integer literal constant. A kind type parameter shall not be specified for  $c$   
4 or  $r$ . The constant  $c$  is interpreted as though it had the same kind type parameter as the corresponding **effective**  
5 **item**. The  $r*c$  form is equivalent to  $r$  successive appearances of the constant  $c$ , and the  $r*$  form is equivalent to  $r$   
6 successive null values. Neither of these forms may contain embedded blanks, except where permitted within the  
7 constant  $c$ .
- 8 2 The datum  $c$  (13.11) is any input value acceptable to format specifications for a given type, except for a restriction  
9 on the form of input values corresponding to list items of types logical, integer, and character as specified in this  
10 subclause. The form of a real or complex value is dependent on the decimal edit mode in effect (13.6). The form  
11 of an input value shall be acceptable for the type of the namelist group object list item. The number and forms  
12 of the input values that may follow the equals in a name-value subsequence depend on the shape and type of  
13 the object represented by the name in the input record. When the name in the input record is that of a scalar  
14 variable of an intrinsic type, the equals shall not be followed by more than one value. Blanks are never used  
15 as zeros, and embedded blanks are not permitted in constants except within character constants and complex  
16 constants as specified in this subclause.
- 17 3 When the next **effective item** is of type real, the input form of the input value is that of a numeric input field. A  
18 numeric input field is a field suitable for F editing (13.7.2.3.2) that is assumed to have no fractional digits unless  
19 a decimal symbol appears within the field.
- 20 4 When the next **effective item** is of type complex, the input form of the input value consists of a left parenthesis  
21 followed by an ordered pair of numeric input fields separated by a comma (if the decimal edit mode is POINT) or  
22 a semicolon (if the decimal edit mode is COMMA), and followed by a right parenthesis. The first numeric input  
23 field is the real part of the complex constant and the second field is the imaginary part. Each of the numeric  
24 input fields may be preceded or followed by any number of blanks and ends of records. The end of a record may  
25 occur between the real part and the comma or semicolon, or between the comma or semicolon and the imaginary  
26 part.
- 27 5 When the next **effective item** is of type logical, the input form of the input value shall not include equals or value  
28 separators among the optional characters permitted for L editing (13.7.3).
- 29 6 When the next **effective item** is of type integer, the value in the input record is interpreted as if an  $Iw$  edit  
30 descriptor with a suitable value of  $w$  were used.
- 31 7 When the next **effective item** is of type character, the input form consists of a sequence of zero or more *rep-chars*  
32 whose kind type parameter is implied by the kind of the corresponding list item, delimited by apostrophes or  
33 quotes. Such a sequence may be continued from the end of one record to the beginning of the next record, but the  
34 end of record shall not occur between a doubled apostrophe in an apostrophe-delimited sequence, nor between a  
35 doubled quote in a quote-delimited sequence. The end of the record does not cause a blank or any other character  
36 to become part of the sequence. The sequence may be continued on as many records as needed. The characters  
37 blank, comma, semicolon, and slash may appear in such character sequences.

#### NOTE 13.36

The delimiters in the input form for a namelist input item of type character avoid the ambiguity that could arise between undelimited character sequences and object names. The value of the **DELIM= specifier**, if any, in the **OPEN statement** for an **external file** is ignored during namelist input (12.5.6.8).

- 38 8 Let  $len$  be the length of the next **effective item**, and let  $w$  be the length of the character sequence. If  $len$  is less  
39 than or equal to  $w$ , the leftmost  $len$  characters of the sequence are transmitted to the next **effective item**. If  $len$   
40 is greater than  $w$ , the constant is transmitted to the leftmost  $w$  characters of the next **effective item** and the  
41 remaining  $len-w$  characters of the next **effective item** are filled with blanks. The effect is as though the sequence  
42 were assigned to the next **effective item** in an **intrinsic assignment statement** (10.2.1.3).

1 **13.11.3.4 Null values**

2 1 A null value is specified by

- 3     • the  $r^*$  form,  
 4     • blanks between two consecutive nonblank value separators following an equals,  
 5     • a value separator that is the first nonblank character following an equals, or  
 6     • two consecutive nonblank value separators.

7 2 A null value has no effect on the definition status of the corresponding input list item. If the namelist group  
 8 object list item is defined, it retains its previous value; if it is undefined, it remains undefined. A null value shall  
 9 not be used as either the real or imaginary part of a complex constant, but a single null value may represent an  
 10 entire complex constant.

**NOTE 13.37**

The end of a record following a value separator, with or without intervening blanks, does not specify a null value in namelist input.

11 **13.11.3.5 Blanks**

12 1 All blanks in a namelist input record are considered to be part of some value separator except for

- 13     • blanks embedded in a character constant,  
 14     • embedded blanks surrounding the real or imaginary part of a complex constant,  
 15     • leading blanks following the equals unless followed immediately by a slash or comma, or a semicolon if the  
 16     decimal edit mode is COMMA, and  
 17     • blanks between a name and the following equals.

18 **13.11.3.6 Namelist comments**

19 1 Except within a character literal constant, a “!” character after a value separator or in the first nonblank position  
 20 of a namelist input record initiates a comment. The comment extends to the end of the record and may contain  
 21 any graphic character in the processor-dependent character set. The comment is ignored. A slash within the  
 22 namelist comment does not terminate execution of the namelist [input statement](#). Namelist comments are not  
 23 allowed in stream input because comments depend on record structure.

**NOTE 13.38**

Namelist input example:

```
INTEGER I; REAL X (8); CHARACTER (11) P; COMPLEX Z; LOGICAL G
NAMELIST / TODAY / G, I, P, Z, X
READ (*, NML = TODAY)
```

The input data records are:

```
&TODAY I = 12345, X(1) = 12345, X(3:4) = 2*1.5, I=6, ! This is a comment.
P = ' ' ISN'T_BOB'S', Z = (123,0)/
```

The results stored are:



**NOTE 13.38 (cont.)**

Variable	Value
I	6
X (1)	12345.0
X (2)	unchanged
X (3)	1.5
X (4)	1.5
X (5) – X (8)	unchanged
P	ISN'T_BOB'S
Z	(123.0,0.0)
G	unchanged

1 **13.11.4 Namelist output**2 **13.11.4.1 Form of namelist output**

3 1 The form of the output produced by intrinsic namelist output shall be suitable for input, except for character  
4 output. The names in the output are in upper case. With the exception of adjacent undelimited character  
5 values, the values are separated by one or more blanks or by a comma, or a semicolon if the decimal edit mode is  
6 COMMA, optionally preceded by one or more blanks and optionally followed by one or more blanks. The form  
7 of the output produced by [defined output \(12.6.4.8\)](#) is determined by the [defined output](#) procedure; this form  
8 need not be compatible with namelist input.

9 2 Namelist output shall not include namelist comments.

10 3 The processor may begin new records as necessary. However, except for complex constants and character values,  
11 the end of a record shall not occur within a constant, character value, or name, and blanks shall not appear  
12 within a constant, character value, or name.

**NOTE 13.39**

The length of the output records is not specified exactly and is processor dependent.

13 **13.11.4.2 Namelist output editing**

14 1 Values in namelist output records are edited as for list-directed output ([13.10.4](#)).

**NOTE 13.40**

Namelist output records produced with a [DELIM= specifier](#) with a value of NONE and which contain a character sequence might not be acceptable as namelist input records.

15 **13.11.4.3 Namelist output records**

16 1 If two or more successive values for the same namelist group item in an output record produced have identical  
17 values, the processor has the option of producing a repeated constant of the form  $r*c$  instead of the sequence of  
18 identical values.

19 2 The name of each namelist group object list item is placed in the output record followed by an equals and a list  
20 of values of the namelist group object list item.

21 3 An ampersand character followed immediately by a *namelist-group-name* is placed at the start of the first output  
22 record to indicate which particular group of data objects is being output. A slash is placed in the output record  
23 to indicate the end of the namelist formatting.

24 4 A null value is not produced by namelist formatting.

25 5 Except for new records created by explicit formatting within a [defined output](#) procedure or by continuation of

1 delimited character sequences, each output record begins with a blank character.

## 14 Program units

### 14.1 Main program

- 1 A Fortran main program is a [program unit](#) that does not contain a [SUBROUTINE](#), [FUNCTION](#), [MODULE](#), [SUBMODULE](#), or [BLOCK DATA](#) statement as its first statement.

R1401 *main-program* is [ *program-stmt* ]  
   [ *specification-part* ]  
   [ *execution-part* ]  
   [ *internal-subprogram-part* ]  
   *end-program-stmt*

R1402 *program-stmt* is PROGRAM *program-name*

R1403 *end-program-stmt* is END [ PROGRAM [ *program-name* ] ]

- C1401 (R1401) The *program-name* may be included in the *end-program-stmt* only if the optional *program-stmt* is used and, if included, shall be identical to the *program-name* specified in the *program-stmt*.

#### NOTE 14.1

The program name is global to the program (19.2). For explanatory information about uses for the program name, see subclause C.9.1.

#### NOTE 14.2

An example of a main program is:

```
PROGRAM ANALYZE
  REAL A, B, C (10,10)    ! Specification part
  CALL FIND              ! Execution part
CONTAINS
  SUBROUTINE FIND        ! Internal subprogram
  ...
  END SUBROUTINE FIND
END PROGRAM ANALYZE
```

- 2 The main program may be defined by means other than Fortran; in that case, the program shall not contain a *main-program program unit*.
- 3 A reference to a Fortran *main-program* shall not appear in any [program unit](#) in the program, including itself.

## 14.2 Modules

### 14.2.1 Module syntax and semantics

- 1 A [module](#) contains declarations, specifications, and definitions. Public identifiers of module entities are accessible to other [program units](#) by [use association](#) as specified in 14.2.2. A module that is provided as an inherent part of the processor is an intrinsic module. A nonintrinsic module is defined by a module [program unit](#) or a means other than Fortran.
- 2 Procedures and types defined in an intrinsic module are not themselves intrinsic.

1	R1404	<i>module</i>	<b>is</b>	<i>module-stmt</i>
2				[ <i>specification-part</i> ]
3				[ <i>module-subprogram-part</i> ]
4				<i>end-module-stmt</i>
5	R1405	<i>module-stmt</i>	<b>is</b>	MODULE <i>module-name</i>
6	R1406	<i>end-module-stmt</i>	<b>is</b>	END [ MODULE [ <i>module-name</i> ] ]
7	R1407	<i>module-subprogram-part</i>	<b>is</b>	<i>contains-stmt</i>
8				[ <i>module-subprogram</i> ] ...
9	R1408	<i>module-subprogram</i>	<b>is</b>	<i>function-subprogram</i>
10			<b>or</b>	<i>subroutine-subprogram</i>
11			<b>or</b>	<i>separate-module-subprogram</i>

12 C1402 (R1404) If the *module-name* is specified in the *end-module-stmt*, it shall be identical to the *module-name*  
 13 specified in the *module-stmt*.

14 C1403 (R1404) A module *specification-part* shall not contain a *stmt-function-stmt*, an *entry-stmt*, or a *format-stmt*.

15 3 If a procedure declared in the *scoping unit* of a module has an *implicit interface*, it shall be given the **EXTERNAL**  
 16 **attribute** in that *scoping unit*; if it is a function, its type and type parameters shall be explicitly declared in a  
 17 *type declaration statement* in that *scoping unit*.

18 4 If an intrinsic procedure is declared in the *scoping unit* of a module, it shall explicitly be given the **INTRINSIC**  
 19 **attribute** in that *scoping unit* or be used as an intrinsic procedure in that *scoping unit*.

#### NOTE 14.3

The module name is global to the program (19.2).

#### NOTE 14.4

Although *statement function definitions*, **ENTRY** statements, and **FORMAT** statements cannot appear in the specification part of a module, they can appear in the specification part of a module subprogram in the module.

#### NOTE 14.5

For a discussion of the impact of modules on dependent compilation, see subclause C.9.2.

#### NOTE 14.6

For examples of the use of modules, see subclause C.9.3.

## 20 14.2.2 The USE statement and use association

21 1 The USE statement specifies *use association*. A USE statement is a *reference* to the module it specifies. At the  
 22 time a USE statement is processed, the public portions of the specified module shall be available. A module shall  
 23 not reference itself, either directly or indirectly.

24 2 The USE statement provides the means by which a *scoping unit* accesses named data objects, derived types,  
 25 procedures, *abstract interfaces*, *generic identifiers*, and namelist groups in a module. The entities in the *scoping*  
 26 *unit* are use associated with the entities in the module. The accessed entities have the attributes specified in the  
 27 module, except that a local entity may have a different *accessibility attribute*, it may have the **ASYNCHRONOUS**  
 28 **attribute** even if the associated module entity does not, and if it is not a *coarray* it may have the **VOLATILE**  
 29 **attribute** even if the associated module entity does not. The entities made accessible are identified by the names  
 30 or *generic identifiers* used to identify them in the module. By default, the local entities are identified by the  
 31 same identifiers in the *scoping unit* containing the USE statement, but it is possible to specify that different local

1 identifiers are used. A use-associated entity is considered to have been previously defined.

#### NOTE 14.7

The accessibility of module entities can be controlled by accessibility attributes (7.5.2.2, 8.5.2), and the ONLY option of the USE statement. Definability of module entities can be controlled by the PROTECTED attribute (8.5.15).

2 R1409 *use-stmt* is USE [ [ , *module-nature* ] :: ] *module-name* [ , *rename-list* ]  
 3 or USE [ [ , *module-nature* ] :: ] *module-name* , ■  
 4 ■ ONLY : [ *only-list* ]

5 R1410 *module-nature* is INTRINSIC  
 6 or NON\_INTRINSIC

7 R1411 *rename* is *local-name* => *use-name*  
 8 or OPERATOR (*local-defined-operator*) => ■  
 9 ■ OPERATOR (*use-defined-operator*)

10 R1412 *only* is *generic-spec*  
 11 or *only-use-name*  
 12 or *rename*

13 R1413 *only-use-name* is *use-name*

14 C1404 (R1409) If *module-nature* is INTRINSIC, *module-name* shall be the name of an intrinsic module.

15 C1405 (R1409) If *module-nature* is NON\_INTRINSIC, *module-name* shall be the name of a nonintrinsic module.

16 C1406 (R1409) A *scoping unit* shall not directly reference an intrinsic module and a nonintrinsic module of the  
 17 same name.

18 C1407 (R1411) OPERATOR (*use-defined-operator*) shall not identify a *type-bound generic interface*.

19 C1408 (R1412) The *generic-spec* shall not identify a *type-bound generic interface*.

#### NOTE 14.8

Constraints C1407 and C1408 do not prevent accessing a *generic-spec* that is declared by an *interface block*, even if a *type-bound generic interface* has the same *generic-spec*.

20 C1409 Each *generic-spec*, *use-name*, and *use-defined-operator* in a USE statement shall be a public identifier of  
 21 the module.

22 C1410 An *only-use-name* shall be a nongeneric name.

23 R1414 *local-defined-operator* is *defined-unary-op*  
 24 or *defined-binary-op*

25 R1415 *use-defined-operator* is *defined-unary-op*  
 26 or *defined-binary-op*

27 3 A *use-stmt* without a *module-nature* provides access either to an intrinsic or to a nonintrinsic module. If the  
 28 *module-name* is the name of both an intrinsic and a nonintrinsic module, the nonintrinsic module is accessed.

29 4 The USE statement without the ONLY option provides access to all public entities in the specified module.

30 5 A USE statement with the ONLY option provides access only to those entities that appear as *generic-specs*,  
 31 *use-names*, or *use-defined-operators* in the *only-list*.

32 6 More than one USE statement for a given module may appear in a specification part. If one of the USE statements

- 1 is without an **ONLY** option, all public entities in the module are accessible. If all the USE statements have **ONLY**  
 2 options, only those entities in one or more of the *only-lists* are accessible.
- 3 7 An accessible entity in the referenced module has one or more local identifiers. These identifiers are
- 4 • the identifier of the entity in the referenced module if that identifier appears as an *only-use-name* or as the  
 5 *defined-operator* of a *generic-spec* in any *only* for that module,
  - 6 • each of the *local-names* or *local-defined-operators* that the entity is given in any *rename* for that module,  
 7 and
  - 8 • the identifier of the entity in the referenced module if that identifier does not appear as a *use-name* or  
 9 *use-defined-operator* in any *rename* for that module.
- 10 8 An ultimate entity is a module entity that is not accessed by use association. Two or more accessible ultimate  
 11 entities, other than **generic interfaces** or defined operators, may have the same local identifier only if the identifier  
 12 is not used. **Generic interfaces** and defined operators are handled as described in 15.4.3.4. Except for these cases,  
 13 the local identifier of any entity given accessibility by a USE statement shall differ from the local identifiers of all  
 14 other entities accessible to the **scoping unit**.

**NOTE 14.9**

There is no prohibition against a *use-name* or *use-defined-operator* appearing multiple times in one USE statement or in multiple USE statements involving the same module. As a result, it is possible for one use-associated entity to be accessible by more than one local identifier.

- 15 9 The local identifier of an entity made accessible by a USE statement shall not appear in any other nonexecutable  
 16 statement that would cause any **attribute** (8.5) of the entity to be specified in the **scoping unit** that contains the  
 17 USE statement, except that it may appear in a **PUBLIC** or **PRIVATE** statement in the **scoping unit** of a module  
 18 and it may be given the **ASYNCHRONOUS** or **VOLATILE** attribute.
- 19 10 An entity in a **scoping unit** that is accessed by use association through more than one use path, has the **ASYN-**  
 20 **CHRONOUS** or **VOLATILE** attribute in any of those use paths, and is not given that attribute in that **scoping**  
 21 **unit**, shall have that attribute in all use paths.

**NOTE 14.10**

The constraints in subclauses 8.10.1, 8.10.2, and 8.9 prohibit the *local-name* from appearing as a *common-block-object* in a **COMMON** statement, an *equivalence-object* in an **EQUIVALENCE** statement, or a *namelist-group-name* in a **NAMELIST** statement, respectively. There is no prohibition against the *local-name* appearing as a *common-block-name* or a *namelist-group-object*.

**NOTE 14.11**

For a discussion of the impact of the **ONLY** option and renaming on dependent compilation, see subclause C.9.2.1.

**NOTE 14.12**

Examples:

```
USE STATS_LIB
```

provides access to all public entities in the module STATS\_LIB.

```
USE MATH_LIB; USE STATS_LIB, SPROD => PROD
```

provides access to all public identifiers in both MATH\_LIB and STATS\_LIB. If MATH\_LIB contains an entity named PROD, it can be accessed by that name, while the entity PROD of STATS\_LIB can be accessed by the name SPROD.

**NOTE 14.12 (cont.)**

```
USE STATS_LIB, ONLY: YPROD; USE STATS_LIB, ONLY : PROD
```

provides access to YPROD and PROD in STAT\_LIB.

```
USE STATS_LIB, ONLY : YPROD; USE STATS_LIB
```

provides access to all public identifiers in STAT\_LIB.

1 **14.2.3 Submodules**

2 1 A **submodule** is a **program unit** that extends a module or another submodule. The **program unit** that it extends  
3 is its **host**, and is specified by the *parent-identifier* in the *submodule-stmt*.

4 2 A module or submodule is an ancestor program unit of all of its descendants, which are its submodules and their  
5 descendants. The submodule identifier is the ordered pair whose first element is the ancestor module name and  
6 whose second element is the submodule name; the submodule name by itself is not a local or global identifier.

**NOTE 14.13**

A module and its submodules stand in a tree-like relationship one to another, with the module at the root.  
Therefore, a submodule has exactly one ancestor module and can have one or more ancestor submodules.

7 3 A submodule may provide implementations for separate module procedures (15.6.2.5), each of which is declared  
8 (15.4.3.2) within that submodule or one of its ancestors, and declarations and definitions of other entities that  
9 are accessible by **host association** in its **descendants**.

10 R1416 *submodule*                            **is** *submodule-stmt*  
11    [ *specification-part* ]  
12    [ *module-subprogram-part* ]  
13    *end-submodule-stmt*

14 R1417 *submodule-stmt*                   **is** SUBMODULE ( *parent-identifier* ) *submodule-name*

15 R1418 *parent-identifier*               **is** *ancestor-module-name* [ : *parent-submodule-name* ]

16 R1419 *end-submodule-stmt*           **is** END [ SUBMODULE [ *submodule-name* ] ]

17 C1411 (R1416) A submodule *specification-part* shall not contain a *format-stmt*, *entry-stmt*, or *stmt-function-stmt*.

18 C1412 (R1418) The *ancestor-module-name* shall be the name of a nonintrinsic module that declares a separate  
19 module procedure; the *parent-submodule-name* shall be the name of a **descendant** of that module.

20 C1413 (R1416) If a *submodule-name* appears in the *end-submodule-stmt*, it shall be identical to the one in the  
21 *submodule-stmt*.

22 **14.3 Block data program units**

23 1 A block data program unit is used to provide initial values for data objects in named **common blocks**.

24 R1420 *block-data*                           **is** *block-data-stmt*  
25    [ *specification-part* ]  
26    *end-block-data-stmt*

27 R1421 *block-data-stmt*                   **is** BLOCK DATA [ *block-data-name* ]

28 R1422 *end-block-data-stmt*           **is** END [ BLOCK DATA [ *block-data-name* ] ]

29 C1414 (R1420) The *block-data-name* shall be included in the *end-block-data-stmt* only if it was provided in the *block-data-stmt*  
30 and, if included, shall be identical to the *block-data-name* in the *block-data-stmt*.

- 1 C1415 (R1420) A *block-data specification-part* shall contain only derived-type definitions and ASYNCHRONOUS, BIND, COM-  
2 MON, DATA, DIMENSION, EQUIVALENCE, IMPLICIT, INTRINSIC, PARAMETER, POINTER, SAVE, TARGET,  
3 USE, VOLATILE, and type declaration statements.
- 4 C1416 (R1420) A type declaration statement in a *block-data specification-part* shall not contain ALLOCATABLE, EXTERNAL,  
5 or BIND attribute specifiers.
- 6 2 If an object in a named common block is initially defined, all storage units in the common block storage sequence shall be specified  
7 even if they are not all initially defined. More than one named common block may have objects initially defined in a single block  
8 data program unit.
- 9 3 Only an object in a named common block may be initially defined in a block data program unit.
- 10 4 The same named common block shall not be specified in more than one block data program unit in a program.
- 11 5 There shall not be more than one unnamed block data program unit in a program.



## 15 Procedures

### 15.1 Concepts

- 1 The concept of a procedure was introduced in 5.2.3. This clause contains a complete description of procedures. The actions specified by a procedure are performed when the procedure is invoked by execution of a reference to it.
- 2 The sequence of actions encapsulated by a procedure has access to entities in the procedure reference by way of [argument association](#) (15.5.2). A name that appears as a *dummy-arg-name* in the `SUBROUTINE`, `FUNCTION`, or `ENTRY` statement in the declaration of a procedure (R1536) is a [dummy argument](#). [Dummy arguments](#) are also specified for intrinsic procedures and procedures in intrinsic modules in Clauses 16, 17, and 18.

### 15.2 Procedure classifications

#### 15.2.1 Procedure classification by reference

- 1 The definition of a procedure specifies it to be a function or a subroutine. A reference to a function either appears explicitly as a primary within an expression, or is implied by a [defined operation](#) (10.1.6) within an expression. A reference to a subroutine is a [CALL statement](#), a [defined assignment](#) statement (10.2.1.4), the appearance of an object processed by [defined input/output](#) (12.6.4.8) in an input/output list, or [finalization](#) (7.5.6).
- 2 A procedure is classified as [elemental](#) if it is a procedure that may be referenced [elementally](#) (15.8).

#### 15.2.2 Procedure classification by means of definition

##### 15.2.2.1 Intrinsic procedures

- 1 A procedure that is provided as an inherent part of the processor is an intrinsic procedure.

##### 15.2.2.2 External, internal, and module procedures

- 1 An [external procedure](#) is a procedure that is defined by an external subprogram or by a means other than Fortran.
- 2 An [internal procedure](#) is a procedure that is defined by an internal subprogram. Internal subprograms may appear in the main program, in an external subprogram, or in a module subprogram. Internal subprograms shall not appear in other internal subprograms. Internal subprograms are the same as external subprograms except that the name of the [internal procedure](#) is not a global identifier, an internal subprogram shall not contain an [ENTRY statement](#), and the internal subprogram has access to host entities by [host association](#).
- 3 A [module procedure](#) is a procedure that is defined by a module subprogram.
- 4 A subprogram defines a procedure for the `SUBROUTINE` or `FUNCTION` statement. If the subprogram has one or more [ENTRY statements](#), it also defines a procedure for each of them.

##### 15.2.2.3 Dummy procedures

- 1 A [dummy argument](#) that is specified to be a procedure or appears as the procedure designator in a [procedure reference](#) is a [dummy procedure](#). A [dummy procedure](#) with the [POINTER attribute](#) is a [dummy procedure pointer](#).

#### 1 15.2.2.4 Procedure pointers

2 1 A [procedure pointer](#) is a procedure that has the [EXTERNAL](#) and [POINTER](#) attributes; it may be [pointer](#)  
3 [associated](#) with an [external procedure](#), an [internal procedure](#), an intrinsic procedure, a [module procedure](#), or a  
4 [dummy procedure](#) that is not a [procedure pointer](#).

#### 5 15.2.2.5 Statement functions

6 1 A function that is defined by a single statement is a statement function ([15.6.4](#)).

### 7 15.3 Characteristics

#### 8 15.3.1 Characteristics of procedures

9 1 The [characteristics](#) of a procedure are the classification of the procedure as a function or subroutine, whether it  
10 is pure, whether it is [elemental](#), whether it has the [BIND attribute](#), the [characteristics](#) of its [dummy arguments](#),  
11 and the [characteristics](#) of its [function result](#) if it is a function.

#### 12 15.3.2 Characteristics of dummy arguments

##### 13 15.3.2.1 General

14 1 Each [dummy argument](#) has the [characteristic](#) that it is a [dummy data object](#), a [dummy procedure](#), or an asterisk  
15 (alternate return indicator).

##### 16 15.3.2.2 Characteristics of dummy data objects

17 1 The [characteristics](#) of a [dummy data object](#) are its [declared type](#), its [type parameters](#), its shape (unless it is  
18 [assumed-rank](#)), its [corank](#), its [codimensions](#), its intent ([8.5.10](#), [8.6.9](#)), whether it is optional ([8.5.12](#), [8.6.10](#)),  
19 whether it is [allocatable](#) ([8.5.3](#)), whether it has the [ASYNCHRONOUS](#) ([8.5.4](#)), [CONTIGUOUS](#) ([8.5.7](#)), [VALUE](#)  
20 ([8.5.18](#)), or [VOLATILE](#) ([8.5.19](#)) attributes, whether it is [polymorphic](#), and whether it is a [pointer](#) ([8.5.14](#), [8.6.12](#))  
21 or a target ([8.5.17](#), [8.6.15](#)). If a [type parameter](#) of an object or a bound of an array is not a [constant expression](#),  
22 the exact dependence on the entities in the expression is a characteristic. If a rank, shape, size, type, or type  
23 parameter is assumed or [deferred](#), it is a characteristic.

##### 24 15.3.2.3 Characteristics of dummy procedures

25 1 The [characteristics](#) of a [dummy procedure](#) are the explicitness of its [interface](#) ([15.4.2](#)), its [characteristics](#) as a  
26 procedure if the [interface](#) is [explicit](#), whether it is a pointer, and whether it is optional ([8.5.12](#), [8.6.10](#)).

##### 27 15.3.2.4 Characteristics of asterisk dummy arguments

28 1 A [dummy argument](#) that is an asterisk has no other characteristic.

#### 29 15.3.3 Characteristics of function results

30 1 The [characteristics](#) of a function result are its [declared type](#), [type parameters](#), [rank](#), whether it is [polymorphic](#),  
31 whether it is [allocatable](#), whether it is a [pointer](#), whether it has the [CONTIGUOUS attribute](#), and whether it is a  
32 [procedure pointer](#). If a function result is an array that is not [allocatable](#) or a pointer, its shape is a characteristic.  
33 If a type parameter of a function result or a bound of a function result array is not a [constant expression](#), the  
34 exact dependence on the entities in the expression is a characteristic. If type parameters of a function result are  
35 [deferred](#), which parameters are [deferred](#) is a characteristic. Whether the length of a character function result is assumed  
36 is a characteristic.

## 15.4 Procedure interface

### 15.4.1 Interface and abstract interface

- 1 The **interface** of a procedure determines the forms of reference through which it may be invoked. The procedure's interface consists of its name, **binding label**, generic identifiers, characteristics, and the names of its dummy arguments. The **characteristics** and **binding label** of a procedure are fixed, but the remainder of the interface may differ in differing contexts, except that for a separate module procedure body (15.6.2.5), the **dummy argument** names and whether it is recursive shall be the same as in its corresponding module procedure interface body (15.4.3.2).
- 2 An **abstract interface** is a set of procedure **characteristics** with the **dummy argument** names.

### 15.4.2 Implicit and explicit interfaces

#### 15.4.2.1 Interfaces and scopes

- 1 The interface of a procedure is either explicit or implicit. It is explicit if it is
  - an **internal procedure**, **module procedure**, or **intrinsic** procedure,
  - a subroutine, or a function with a separate result name, within the **scoping unit** that defines it, or
  - a procedure declared by a **procedure declaration statement** that specifies an explicit interface, or by an interface body.
 Otherwise, the interface of the identifier is implicit. The interface of a statement function is always implicit.

#### NOTE 15.1

For example, the subroutine LLS of C.9.3.4 has an **explicit interface**.

#### 15.4.2.2 Explicit interface

- 1 Within the scope of a procedure identifier, the procedure shall have an **explicit interface** if it is not a statement function and
  - (1) a reference to the procedure appears with an **argument keyword** (15.5.2),
  - (2) the procedure is used in a context that requires it to be pure (15.7),
  - (3) the procedure has a **dummy argument** that
    - (a) has the **ALLOCATABLE**, **ASYNCHRONOUS**, **OPTIONAL**, **POINTER**, **TARGET**, **VALUE**, or **VOLATILE** attribute,
    - (b) is an **assumed-shape array**,
    - (c) is **assumed-rank**,
    - (d) is a **coarray**,
    - (e) is of a parameterized derived type, or
    - (f) is **polymorphic**,
  - (4) the procedure has a result that
    - (a) is an array,
    - (b) is a **pointer** or is **allocatable**, or
    - (c) has a nonassumed type parameter value that is not a **constant expression**,
  - (5) the procedure is **elemental**, or
  - (6) the procedure has the **BIND attribute**.

## 15.4.3 Specification of the procedure interface

### 15.4.3.1 General

- 1 The interface for an [internal](#), [external](#), [module](#), or [dummy procedure](#) is specified by a [FUNCTION](#), [SUBROUTINE](#), or [ENTRY](#) statement and by specification statements for the [dummy arguments](#) and the result of a function. These statements may appear in the procedure definition, in an [interface body](#), or both, except that the [ENTRY statement](#) shall not appear in an [interface body](#).

#### NOTE 15.2

An [interface body](#) cannot be used to describe the interface of an [internal procedure](#), a [module procedure](#) that is not a separate [module procedure](#), or an intrinsic procedure because the interfaces of such procedures are already [explicit](#). However, the name of a procedure can appear in a [PROCEDURE statement](#) in an [interface block](#) (15.4.3.2).

### 15.4.3.2 Interface block

- |    |       |   |    |   |
|----|-------|---|----|---|
| 8  | R1501 | <i>interface-block</i>  | is | <i>interface-stmt</i><br>[ <i>interface-specification</i> ] ...<br><i>end-interface-stmt</i>  |
| 11 | R1502 | <i>interface-specification</i>  | is | <i>interface-body</i><br>or <i>procedure-stmt</i>   |
| 13 | R1503 | <i>interface-stmt</i>   | is | INTERFACE [ <i>generic-spec</i> ]<br>or ABSTRACT INTERFACE  |
| 15 | R1504 | <i>end-interface-stmt</i>   | is | END INTERFACE [ <i>generic-spec</i> ]   |
| 16 | R1505 | <i>interface-body</i>   | is | <i>function-stmt</i><br>[ <i>specification-part</i> ]<br><i>end-function-stmt</i><br>or <i>subroutine-stmt</i><br>[ <i>specification-part</i> ]<br><i>end-subroutine-stmt</i> |
| 22 | R1506 | <i>procedure-stmt</i>   | is | [ MODULE ] PROCEDURE [ :: ] <i>specific-procedure-list</i>  |
| 23 | R1507 | <i>specific-procedure</i>   | is | <i>procedure-name</i>   |
| 24 | R1508 | <i>generic-spec</i>   | is | <i>generic-name</i><br>or OPERATOR ( <i>defined-operator</i> )<br>or ASSIGNMENT ( = )<br>or <i>defined-io-generic-spec</i>  |
| 28 | R1509 | <i>defined-io-generic-spec</i>  | is | READ (FORMATTED)<br>or READ (UNFORMATTED)<br>or WRITE (FORMATTED)<br>or WRITE (UNFORMATTED)   |
| 32 | C1501 | (R1501) An <i>interface-block</i> in a subprogram shall not contain an <i>interface-body</i> for a procedure defined by that subprogram.  |    |   |
| 34 | C1502 | (R1501) If the <i>end-interface-stmt</i> includes a <i>generic-spec</i> , the <i>interface-stmt</i> shall specify the same <i>generic-spec</i> , except that if one <i>generic-spec</i> has a <i>defined-operator</i> that is <i>.LT.</i> , <i>.LE.</i> , <i>.GT.</i> , <i>.GE.</i> , <i>.EQ.</i> , or <i>.NE.</i> , the other <i>generic-spec</i> may have a <i>defined-operator</i> that is the corresponding operator <i>&lt;</i> , <i>&lt;=</i> , <i>&gt;</i> , <i>&gt;=</i> , <i>==</i> , or <i>/=</i> . |    |   |

- 1 C1503 (R1503) If the *interface-stmt* is ABSTRACT INTERFACE, then the *function-name* in the *function-stmt*  
2 or the *subroutine-name* in the *subroutine-stmt* shall not be the same as a keyword that specifies an  
3 intrinsic type.
- 4 C1504 (R1502) A *procedure-stmt* is allowed only in an interface block that has a *generic-spec*.
- 5 C1505 (R1505) An *interface-body* of a pure procedure shall specify the intents of all *dummy arguments* except  
6 alternate return indicators, *dummy procedures*, and arguments with the POINTER or VALUE attribute.
- 7 C1506 (R1505) An *interface-body* shall not contain a *data-stmt*, *format-stmt*, *entry-stmt*, or *stmt-function-stmt*.
- 8 C1507 (R1506) If MODULE appears in a *procedure-stmt*, each *procedure-name* in that statement shall denote a  
9 *module procedure*.
- 10 C1508 (R1507) A *procedure-name* shall denote a nonintrinsic procedure that has an *explicit interface*.
- 11 C1509 (R1501) An *interface-specification* in a generic interface block shall not specify a procedure that was  
12 specified previously in any accessible interface with the same *generic identifier*.
- 13 1 An external or module subprogram specifies a *specific interface* for each procedure defined in that subprogram.
- 14 2 An interface block introduced by ABSTRACT INTERFACE is an *abstract interface block*. An interface body  
15 in an *abstract interface block* specifies an *abstract interface*. An interface block with a generic specification is  
16 a *generic interface block*. An interface block with neither ABSTRACT nor a generic specification is a *specific*  
17 *interface block*.
- 18 3 The name of the entity declared by an interface body is the *function-name* in the *function-stmt* or the *subroutine-*  
19 *name* in the *subroutine-stmt* that begins the interface body.
- 20 4 A module procedure interface body is an interface body whose initial statement contains the keyword MODULE.  
21 It specifies the interface for a separate *module procedure* (15.6.2.5). A separate *module procedure* is accessible  
22 by *use association* if and only if its interface body is declared in the specification part of a module and is public.  
23 If a corresponding (15.6.2.5) separate *module procedure* is not defined, the interface may be used to specify an  
24 explicit *specific interface* but the procedure shall not be used in any other way.
- 25 5 An interface body in a generic or specific interface block specifies the EXTERNAL attribute and an explicit  
26 specific interface for an *external procedure*, *dummy procedure*, or *procedure pointer*. If the name of the declared  
27 procedure is that of a *dummy argument* in the subprogram containing the interface body, the procedure is a  
28 *dummy procedure*. If the procedure has the POINTER attribute, it is a *procedure pointer*. If it is not a *dummy*  
29 *procedure* or *procedure pointer*, it is an *external procedure*.
- 30 6 An interface body specifies all of the *characteristics* of the explicit specific interface or *abstract interface*. The  
31 specification part of an interface body may specify attributes or define values for data entities that do not  
32 determine *characteristics* of the procedure. Such specifications have no effect.
- 33 7 If an explicit specific interface for an *external procedure* is specified by an interface body or a *procedure declaration*  
34 *statement* (15.4.3.6), the *characteristics* shall be consistent with those specified in the procedure definition, except  
35 that the interface may specify a procedure that is not pure even if the procedure is defined to be pure. An interface  
36 for a procedure defined by an ENTRY statement may be specified by using the entry name as the procedure name in the interface body.  
37 If an *external procedure* does not exist in the program, an interface body for it may be used to specify an explicit  
38 specific interface but the procedure shall not be used in any other way. A procedure shall not have more than  
39 one explicit specific interface in a given *scoping unit*, except that if the interface is accessed by *use association*,  
40 there may be more than one local name for the procedure. If a procedure is accessed by *use association*, each  
41 access shall be to the same procedure declaration or definition.

**NOTE 15.3**

The *dummy argument* names in an interface body can be different from the corresponding *dummy argument* names in the procedure definition because the name of a *dummy argument* is not a characteristic.

**NOTE 15.4**

An example of a specific interface block is:

```

INTERFACE
  SUBROUTINE EXT1 (X, Y, Z)
    REAL, DIMENSION (100, 100) :: X, Y, Z
  END SUBROUTINE EXT1
  SUBROUTINE EXT2 (X, Z)
    REAL X
    COMPLEX (KIND = 4) Z (2000)
  END SUBROUTINE EXT2
  FUNCTION EXT3 (P, Q)
    LOGICAL EXT3
    INTEGER P (1000)
    LOGICAL Q (1000)
  END FUNCTION EXT3
END INTERFACE

```

This interface block specifies [explicit interfaces](#) for the three [external procedures](#) EXT1, EXT2, and EXT3. Invocations of these procedures can use [argument keywords](#) (15.5.2); for example:

```
PRINT *, EXT3 (Q = P_MASK (N+1 : N+1000), P = ACTUAL_P)
```

### 1 15.4.3.3 GENERIC statement

2 1 A GENERIC statement specifies a generic identifier for one or more specific procedures, in the same way as a  
3 generic interface block that does not contain interface bodies.

4 R1510 *generic-stmt* is GENERIC [ , *access-spec* ] :: *generic-spec* => *specific-procedure-list*

5 C1510 (R1510) A *specific-procedure* in a GENERIC statement shall not specify a procedure that was specified  
6 previously in any accessible interface with the same generic identifier.

7 2 If *access-spec* appears, it specifies the accessibility (8.5.2) of *generic-spec*.

### 8 15.4.3.4 Generic interfaces

#### 9 15.4.3.4.1 Generic identifiers

10 1 A [generic interface block](#) specifies a [generic interface](#) for each of the procedures in the interface block. The  
11 PROCEDURE statement lists [procedure pointers](#), [external procedures](#), [dummy procedures](#), or [module procedures](#)  
12 that have this [generic interface](#). A [GENERIC statement](#) specifies a generic interface for each of the procedures  
13 named in its *specific-procedure-list*. A [generic interface](#) is always [explicit](#).

14 2 The *generic-spec* in an *interface-stmt* is a [generic identifier](#) for all the procedures in the interface block. The  
15 *generic-spec* in a [GENERIC statement](#) is a [generic identifier](#) for all of the procedures named in its *specific-*  
16 *procedure-list*. The rules specifying how any two procedures with the same [generic identifier](#) shall differ are given  
17 in 15.4.3.4.5. They ensure that any generic invocation applies to at most one specific procedure. If a specific  
18 procedure in a [generic interface](#) has a function [dummy argument](#), that argument shall have its type and type  
19 parameters explicitly declared in the [specific interface](#).

20 3 A generic name is a [generic identifier](#) that refers to all of the procedure names in the [generic interface](#). A generic  
21 name may be the same as any one of the procedure names in the [generic interface](#), or the same as any accessible  
22 generic name.

23 4 A generic name may be the same as a derived-type name, in which case all of the procedures in the [generic](#)  
24 [interface](#) shall be functions.

- 1 5 An *interface-stmt* having a *defined-io-generic-spec* is an interface for a *defined input/output* procedure (12.6.4.8).

**NOTE 15.5**

An example of a generic procedure interface is:

```
INTERFACE SWITCH
  SUBROUTINE INT_SWITCH (X, Y)
    INTEGER, INTENT (INOUT) :: X, Y
  END SUBROUTINE INT_SWITCH
  SUBROUTINE REAL_SWITCH (X, Y)
    REAL, INTENT (INOUT) :: X, Y
  END SUBROUTINE REAL_SWITCH
  SUBROUTINE COMPLEX_SWITCH (X, Y)
    COMPLEX, INTENT (INOUT) :: X, Y
  END SUBROUTINE COMPLEX_SWITCH
END INTERFACE SWITCH
```

Any of these three subroutines (INT\_SWITCH, REAL\_SWITCH, COMPLEX\_SWITCH) can be referenced with the generic name SWITCH, as well as by its *specific name*. For example, a reference to INT\_SWITCH could take the form:

```
CALL SWITCH (MAX_VAL, LOC_VAL) ! MAX_VAL and LOC_VAL are of type INTEGER
```

**NOTE 15.6**

A *type-bound-generic-stmt* within a derived-type definition (7.5.5) specifies a *generic identifier* for a set of *type-bound procedures*.

#### 2 15.4.3.4.2 Defined operations

- 3 1 If **OPERATOR** is specified in a generic specification, all of the procedures specified in the *generic interface* shall  
 4 be functions that may be referenced as *defined operations* (10.1.6, 15.5). In the case of functions of two arguments,  
 5 infix binary operator notation is implied. In the case of functions of one argument, prefix operator notation is  
 6 implied. **OPERATOR** shall not be specified for functions with no arguments or for functions with more than two  
 7 arguments. The *dummy arguments* shall be nonoptional *dummy data objects* and shall have the **INTENT (IN)**  
 8 or **VALUE** attribute. The function result shall not have assumed character length. If the operator is an *intrinsic-operator*  
 9 (R608), the number of *dummy arguments* shall be consistent with the intrinsic uses of that operator, and the  
 10 types, kind type parameters, or *ranks* of the *dummy arguments* shall differ from those required for the intrinsic  
 11 operation (10.1.5).
- 12 2 A *defined operation* is treated as a reference to the function. For a unary *defined operation*, the operand  
 13 corresponds to the function's *dummy argument*; for a binary operation, the left-hand operand corresponds to the  
 14 first *dummy argument* of the function and the right-hand operand corresponds to the second *dummy argument*.  
 15 All restrictions and constraints that apply to *actual arguments* in a *reference* to the function also apply to the  
 16 corresponding operands in the expression as if they were used as *actual arguments*.
- 17 3 A given defined operator may, as with generic names, apply to more than one function, in which case it is generic  
 18 in exact analogy to generic procedure names. For intrinsic operator symbols, the generic properties include the  
 19 intrinsic operations they represent. Because both forms of each relational operator have the same interpretation  
 20 (10.1.6.2), extending one form (such as <=) has the effect of defining both forms (<= and **.LE.**).

**NOTE 15.7**

An example of the use of the **OPERATOR** generic specification is:

```
INTERFACE OPERATOR ( * )
  FUNCTION BOOLEAN_AND (B1, B2)
```

**NOTE 15.7 (cont.)**

```

    LOGICAL, INTENT (IN) :: B1 (:), B2 (SIZE (B1))
    LOGICAL :: BOOLEAN_AND (SIZE (B1))
    END FUNCTION BOOLEAN_AND
END INTERFACE OPERATOR ( * )

```

This allows, for example

```
SENSOR (1:N) * ACTION (1:N)
```

as an alternative to the function reference

```

BOOLEAN_AND (SENSOR (1:N), ACTION (1:N))    ! SENSOR and ACTION are
                                             ! of type LOGICAL

```

1 **15.4.3.4.3 Defined assignments**

- 2 1 If **ASSIGNMENT** ( = ) is specified in a generic specification, all the procedures in the **generic interface** shall  
3 be subroutines that may be referenced as **defined assignments** (10.2.1.4, 10.2.1.5). **Defined assignment** may, as  
4 with generic names, apply to more than one subroutine, in which case it is generic in exact analogy to generic  
5 procedure names.
- 6 2 Each of these subroutines shall have exactly two **dummy arguments**. The **dummy arguments** shall be nonoptional  
7 **dummy data objects**. The first argument shall have **INTENT (OUT)** or **INTENT (INOUT)** and the second  
8 argument shall have the **INTENT (IN)** or **VALUE** attribute. Either the second argument shall be an array whose  
9 **rank** differs from that of the first argument, the **declared types** and **kind type parameters** of the arguments shall  
10 not conform as specified in Table 10.8, or the first argument shall be of derived type. A **defined assignment** is  
11 treated as a reference to the subroutine, with the left-hand side as the first argument and the right-hand side  
12 enclosed in parentheses as the second argument. All restrictions and constraints that apply to **actual arguments**  
13 in a reference to the subroutine also apply to the left-hand-side and to the right-hand-side enclosed in parentheses  
14 as if they were used as **actual arguments**. The **ASSIGNMENT** generic specification specifies that assignment is  
15 extended or redefined.

**NOTE 15.8**

An example of the use of the **ASSIGNMENT** generic specification is:

```

INTERFACE ASSIGNMENT ( = )

    SUBROUTINE LOGICAL_TO_NUMERIC (N, B)
        INTEGER, INTENT (OUT) :: N
        LOGICAL, INTENT (IN)  :: B
    END SUBROUTINE LOGICAL_TO_NUMERIC
    SUBROUTINE CHAR_TO_STRING (S, C)
        USE STRING_MODULE      ! Contains definition of type STRING
        TYPE (STRING), INTENT (OUT) :: S ! A variable-length string
        CHARACTER (*), INTENT (IN) :: C
    END SUBROUTINE CHAR_TO_STRING
END INTERFACE ASSIGNMENT ( = )

```

Example assignments are:

```

KOUNT = SENSOR (J)    ! CALL LOGICAL_TO_NUMERIC (KOUNT, (SENSOR (J)))
NOTE  = '89AB'       ! CALL CHAR_TO_STRING (NOTE, ('89AB'))

```



**NOTE 15.9**

A procedure which has a generic identifier of **ASSIGNMENT** ( = ) and whose second **dummy argument** has the **ALLOCATABLE** or **POINTER** attribute cannot be directly invoked by **defined assignment**. This is because the **actual argument** associated with that **dummy argument** is the right-hand side of the assignment enclosed in parentheses, which makes the **actual argument** an expression that does not have the **ALLOCATABLE**, **POINTER**, or **TARGET** attribute.

1 **15.4.3.4.4 Defined input/output procedure interfaces**

2 1 All of the procedures specified in an interface block for a **defined input/output** procedure shall be subroutines  
3 that have interfaces as described in 12.6.4.8.2.

4 **15.4.3.4.5 Restrictions on generic declarations**

5 1 This subclause contains the rules that shall be satisfied by every pair of specific procedures that have the same  
6 **generic identifier** within the scope of the identifier. If a generic procedure is accessed from a module, the rules  
7 apply to all the specific versions even if some of them are inaccessible by their **specific names**.

**NOTE 15.10**

In most **scoping units**, the possible sources of procedures with a particular **generic identifier** are the accessible generic identifiers specified by **generic interface blocks** or **GENERIC statements** and the generic **bindings** other than names for the accessible objects in that **scoping unit**. In a type definition, they are the generic **bindings**, including those from a **parent type**.

8 2 A **dummy argument** is type, kind, and **rank** compatible, or TKR compatible, with another **dummy argument** if  
9 the first is **type compatible** with the second, the kind type parameters of the first have the same values as the  
10 corresponding kind type parameters of the second, and both have the same **rank** or either is **assumed-rank**.

11 3 Two **dummy arguments** are distinguishable if

- 12 • one is a procedure and the other is a data object,
- 13 • they are both data objects or known to be functions, and neither is TKR compatible with the other,
- 14 • one has the **ALLOCATABLE attribute** and the other has the **POINTER attribute** and not the **INTENT**  
15 **(IN) attribute**, or
- 16 • one is a function with nonzero **rank** and the other is not known to be a function.

17 C1511 Within the scope of a generic operator, if two procedures with that identifier have the same number of  
18 arguments, one shall have a **dummy argument** that corresponds by position in the argument list to a  
19 **dummy argument** of the other that is distinguishable from it.

20 C1512 Within the scope of the generic **ASSIGNMENT** (=) identifier, if two procedures have that identifier, one  
21 shall have a **dummy argument** that corresponds by position in the argument list to a **dummy argument**  
22 of the other that is distinguishable from it.

23 C1513 Within the scope of a *defined-io-generic-spec*, if two procedures have that **generic identifier**, their **dtv**  
24 arguments (12.6.4.8.2) shall be distinguishable.

25 C1514 Within the scope of a generic name, each pair of procedures identified by that name shall both be  
26 subroutines or both be functions, and

- 27 (1) there is a non-passed-object **dummy data object** in one or the other of them such that
  - 28 (a) the number of **dummy data objects** in one that are nonoptional, are not passed-object, and  
29 with which that **dummy data object** is TKR compatible, possibly including that **dummy**  
30 **data object** itself,  
31 exceeds

- 1 (b) the number of non-passed-object **dummy data objects**, both optional and nonoptional, in  
 2 the other that are not distinguishable from that **dummy data object**,
- 3 (2) the number of nonoptional **dummy procedures** in one of them exceeds the number of **dummy**  
 4 **procedures** in the other,
- 5 (3) both have **passed-object dummy arguments** and the **passed-object dummy arguments** are distin-  
 6 guishable, or
- 7 (4) at least one of them shall have both
- 8 (a) a nonoptional non-passed-object **dummy argument** at an effective position such that either  
 9 the other procedure has no **dummy argument** at that effective position or the **dummy argu-**  
 10 **ment** at that position is distinguishable from it, and
- 11 (b) a nonoptional non-passed-object **dummy argument** whose name is such that either the other  
 12 procedure has no **dummy argument** with that name or the **dummy argument** with that name  
 13 is distinguishable from it,
- 14 and the **dummy argument** that disambiguates by position shall either be the same as or occur  
 15 earlier in the argument list than the one that disambiguates by name.
- 16 4 The effective position of a **dummy argument** is its position in the argument list after any **passed-object dummy**  
 17 **argument** has been removed.
- 18 5 Within the scope of a generic name that is the same as the generic name of an intrinsic procedure, the intrinsic  
 19 procedure is not accessible by its generic name if the procedures in the interface and the intrinsic procedure are  
 20 not all functions or not all subroutines. If a generic invocation is consistent with both a specific procedure from  
 21 an interface and an accessible intrinsic procedure, it is the specific procedure from the interface that is referenced.

**NOTE 15.11**

An extensive explanation of the application of these rules is in [C.10.6](#).

**15.4.3.5 EXTERNAL statement**

- 22 1 An EXTERNAL statement specifies the **EXTERNAL attribute** (8.5.9) for a list of names.

23 R1511 *external-stmt* is EXTERNAL [ :: ] *external-name-list*

- 24 2 The appearance of the name of a **block data program unit** in an EXTERNAL statement confirms that the **block**  
 25 **data program unit** is a part of the program.

**NOTE 15.12**

For explanatory information on potential portability problems with **external procedures**, see subclause [C.10.1](#).

**NOTE 15.13**

An example of an EXTERNAL statement is:

EXTERNAL FOCUS

**15.4.3.6 Procedure declaration statement**

- 27 1 A procedure declaration statement declares **procedure pointers**, **dummy procedures**, and **external procedures**. It  
 28 specifies the **EXTERNAL attribute** (8.5.9) for all entities in the *proc-decl-list*.

29 R1512 *procedure-declaration-stmt* is PROCEDURE ( [ *proc-interface* ] ) ■  
 30 ■ [ [ , *proc-attr-spec* ] ... :: ] *proc-decl-list*

31 R1513 *proc-interface* is *interface-name*  
 32 or *declaration-type-spec*

- 1 R1514 *proc-attr-spec* is *access-spec*  
 2 or *proc-language-binding-spec*  
 3 or INTENT ( *intent-spec* )  
 4 or OPTIONAL  
 5 or POINTER  
 6 or PROTECTED  
 7 or SAVE
- 8 R1515 *proc-decl* is *procedure-entity-name* [ => *proc-pointer-init* ]
- 9 R1516 *interface-name* is *name*
- 10 R1517 *proc-pointer-init* is *null-init*  
 11 or *initial-proc-target*
- 12 R1518 *initial-proc-target* is *procedure-name*  
 13
- 14 C1515 (R1516) The *name* shall be the name of an **abstract interface** or of a procedure that has an **explicit**  
 15 **interface**. If *name* is declared by a *procedure-declaration-stmt* it shall be previously declared. If *name*  
 16 denotes an intrinsic procedure it shall be one that is listed in Table 16.2.
- 17 C1516 (R1516) The *name* shall not be the same as a keyword that specifies an intrinsic type.
- 18 C1517 (R1512) If a *proc-interface* describes an **elemental procedure**, each *procedure-entity-name* shall specify an  
 19 **external procedure**.
- 20 C1518 (R1515) If => appears in *proc-decl*, the procedure entity shall have the **POINTER attribute**.
- 21 C1519 (R1518) The *procedure-name* shall be the name of a **nonelemental external** or **module procedure**, or a  
 22 specific intrinsic function listed in Table 16.2.
- 23 C1520 (R1512) If *proc-language-binding-spec* with NAME= is specified, then *proc-decl-list* shall contain exactly  
 24 one *proc-decl*, which shall neither have the **POINTER attribute** nor be a **dummy procedure**.
- 25 C1521 (R1512) If *proc-language-binding-spec* is specified, the *proc-interface* shall appear, it shall be an *interface-*  
 26 *name*, and *interface-name* shall be declared with a *proc-language-binding-spec*.
- 27 2 If *proc-interface* appears and consists of *interface-name*, it specifies an explicit **specific interface** (15.4.3.2) for the  
 28 declared procedure entities. The **abstract interface** (15.4) is that specified by the interface named by *interface-*  
 29 *name*. The interface specified by *interface-name* shall not depend on any characteristic of a procedure identified  
 30 by a *procedure-entity-name* in the *proc-decl-list* of the same procedure declaration statement.
- 31 3 If *proc-interface* appears and consists of *declaration-type-spec*, it specifies that the declared procedure entities are  
 32 functions having **implicit interfaces** and the specified result type. If a type is specified for an external function,  
 33 its function definition (15.6.2.2) shall specify the same result type and type parameters.
- 34 4 If *proc-interface* does not appear, the procedure declaration statement does not specify whether the declared  
 35 procedure entities are subroutines or functions.
- 36 5 If a *proc-attr-spec* other than a *proc-language-binding-spec* appears, it specifies that the declared procedure entities  
 37 have that attribute. These attributes are described in 8.5. If a *proc-language-binding-spec* with NAME= appears,  
 38 it specifies a **binding label** or its absence, as described in 18.10.2. A *proc-language-binding-spec* without NAME=  
 39 is allowed, but is redundant with the *proc-interface* required by C1521.
- 40 6 If => appears in a *proc-decl* in a *procedure-declaration-stmt* it specifies the initial association status of the  
 41 corresponding procedure entity, and implies the **SAVE attribute**, which may be confirmed by explicit specification.  
 42 If => *null-init* appears, the procedure entity is initially **disassociated**. If => *initial-proc-target* appears, the  
 43 procedure entity is initially associated with the **target**.

- 1 7 If *procedure-entity-name* has an **explicit interface**, its **characteristics** shall be the same as *initial-proc-target* except  
 2 that *initial-proc-target* may be pure even if *procedure-entity-name* is not pure and *initial-proc-target* may be an **elemental**  
 3 **intrinsic procedure**.
- 4 8 If the **characteristics** of *procedure-entity-name* or *initial-proc-target* are such that an **explicit interface** is required,  
 5 both *procedure-entity-name* and *initial-proc-target* shall have an **explicit interface**.
- 6 9 If *procedure-entity-name* has an **implicit interface** and is explicitly typed or referenced as a function, *initial-proc-*  
 7 *target* shall be a function. If *procedure-entity-name* has an **implicit interface** and is referenced as a subroutine,  
 8 *initial-proc-target* shall be a subroutine.
- 9 10 If *initial-proc-target* and *procedure-entity-name* are functions, their results shall have the same **characteristics**.

**NOTE 15.14**

The following code illustrates procedure declaration statements. Note 10.47 illustrates the use of the P and BESSEL defined by this code.

```

ABSTRACT INTERFACE
  FUNCTION REAL_FUNC (X)
    REAL, INTENT (IN) :: X
    REAL :: REAL_FUNC
  END FUNCTION REAL_FUNC
END INTERFACE

INTERFACE
  SUBROUTINE SUB (X)
    REAL, INTENT (IN) :: X
  END SUBROUTINE SUB
END INTERFACE

!-- Some external or dummy procedures with explicit interface.
PROCEDURE (REAL_FUNC) :: BESSEL, GFUN
PROCEDURE (SUB) :: PRINT_REAL
!-- Some procedure pointers with explicit interface,
!-- one initialized to NULL().
PROCEDURE (REAL_FUNC), POINTER :: P, R => NULL ()
PROCEDURE (REAL_FUNC), POINTER :: PTR_TO_GFUN
!-- A derived type with a procedure pointer component ...
TYPE STRUCT_TYPE
  PROCEDURE (REAL_FUNC), POINTER, NOPASS :: COMPONENT
END TYPE STRUCT_TYPE
!-- ... and a variable of that type.
TYPE(STRUCT_TYPE) :: STRUCT
!-- An external or dummy function with implicit interface
PROCEDURE (REAL) :: PSI

```

10 **15.4.3.7 INTRINSIC statement**

11 1 An INTRINSIC statement specifies the **INTRINSIC attribute** (8.5.11) for a list of names.

12 R1519 *intrinsic-stmt* is INTRINSIC [ :: ] *intrinsic-procedure-name-list*

13 C1522 (R1519) Each *intrinsic-procedure-name* shall be the name of an intrinsic procedure.

### 1 15.4.3.8 Implicit interface specification

- 2 1 If the interface of a function is **implicit**, the type and type parameters of the function result are specified by an  
 3 implicit or explicit type specification of the function name. The type, type parameters, and shape of the **dummy**  
 4 **arguments** of a procedure invoked from where the interface of the procedure is **implicit** shall be such that each  
 5 **actual argument** is consistent with the **characteristics** of the corresponding **dummy argument**.

## 6 15.5 Procedure reference

### 7 15.5.1 Syntax of a procedure reference

- 8 1 The form of a procedure reference is dependent on the **interface** of the procedure or **procedure pointer**, but is  
 9 independent of the means by which the procedure is defined. The forms of procedure references are as follows.

10 R1520 *function-reference* is *procedure-designator* ( [ *actual-arg-spec-list* ] )

11 C1523 (R1520) The *procedure-designator* shall designate a function.

12 C1524 (R1520) The *actual-arg-spec-list* shall not contain an *alt-return-spec*.

13 R1521 *call-stmt* is CALL *procedure-designator* [ ( [ *actual-arg-spec-list* ] ) ]

14 C1525 (R1521) The *procedure-designator* shall designate a subroutine.

15 R1522 *procedure-designator* is *procedure-name*  
 16 or *proc-component-ref*  
 17 or *data-ref* % *binding-name*

18 C1526 (R1522) A *procedure-name* shall be a generic name or the name of a procedure.

19 C1527 (R1522) A *binding-name* shall be a **binding name** (7.5.5) of the **declared type** of *data-ref*.

20 C1528 (R1522) A *data-ref* shall not be a **polymorphic** subobject of a **coindexed object**.

21 C1529 (R1522) If *data-ref* is an array, the referenced **type-bound procedure** shall have the **PASS attribute**.

- 22 2 The *data-ref* in a *procedure-designator* shall not be an unallocated **allocatable** variable or a pointer that is not  
 23 associated.

- 24 3 Resolving references to **type-bound procedures** is described in 15.5.6.

- 25 4 A function may also be referenced as a **defined operation** (10.1.6). A subroutine may also be referenced as a  
 26 **defined assignment** (10.2.1.4, 10.2.1.5), by **defined input/output** (12.6.4.8), or by **finalization** (7.5.6).

#### NOTE 15.15

When resolving **type-bound procedure** references, constraints on the use of **coindexed objects** ensure that the **coindexed object** (on the remote **image**) has the same **dynamic type** as the corresponding object on the local **image**. Thus a processor can resolve the **type-bound procedure** using the **coarray** variable on its own **image** and pass the **coindexed object** as the **actual argument**.

27 R1523 *actual-arg-spec* is [ *keyword* = ] *actual-arg*

28 R1524 *actual-arg* is *expr*  
 29 or *variable*  
 30 or *procedure-name*  
 31 or *proc-component-ref*  
 32 or *alt-return-spec*

- 1 R1525 *alt-return-spec* is \* *label*
- 2 C1530 (R1523) The *keyword* = shall not appear if the *interface* of the procedure is *implicit*.
- 3 C1531 (R1523) The *keyword* = shall not be omitted from an *actual-arg-spec* unless it has been omitted from  
4 each preceding *actual-arg-spec* in the argument list.
- 5 C1532 (R1523) Each *keyword* shall be the name of a *dummy argument* in the *explicit interface* of the procedure.
- 6 C1533 (R1524) A nonintrinsic *elemental procedure* shall not be used as an *actual argument*.
- 7 C1534 (R1524) A *procedure-name* shall be the name of an *external*, *internal*, *module*, or *dummy* procedure, a  
8 specific intrinsic function listed in Table 16.2, or a *procedure pointer*.
- 9 C1535 (R1524) *expr* shall not be a variable.
- 10 C1536 (R1525) The *label* shall be the *statement label* of a *branch target statement* that appears in the same *inclusive scope* as the  
11 *call-stmt*.
- 12 C1537 An actual argument that is a *coindexed object* shall not have a pointer ultimate component.

**NOTE 15.16**

Examples of procedure reference using *procedure pointers*:

```
P => BESSEL
WRITE (*, *) P(2.5)      !-- BESSEL(2.5)

S => PRINT_REAL
CALL S(3.14)
```

**NOTE 15.17**

An *internal procedure* cannot be invoked using a *procedure pointer* from either Fortran or C after the *host instance* completes execution, because the pointer is then undefined. While the *host instance* is active, however, if an *internal procedure* was passed as an *actual argument* or is the *target* of a *procedure pointer*, it could be invoked from outside of the host subprogram.

Assume there is a procedure with the following *interface* that calculates  $\int_a^b f(x) dx$ .

```
INTERFACE
  FUNCTION INTEGRATE(F, A, B) RESULT(INTEGRAL) BIND(C)
    USE ISO_C_BINDING
    INTERFACE
      FUNCTION F(X) BIND(C) ! Integrand
        USE ISO_C_BINDING
        REAL(C_FLOAT), VALUE :: X
        REAL(C_FLOAT) :: F
      END FUNCTION
    END INTERFACE
    REAL(C_FLOAT), VALUE :: A, B ! Bounds
    REAL(C_FLOAT) :: INTEGRAL
  END FUNCTION INTEGRATE
END INTERFACE
```

This procedure can be called from Fortran or C, and could be written in either Fortran or C. The argument F representing the mathematical function  $f(x)$  can be written as an *internal procedure*; this *internal procedure* will have access to any *host instance local variables* necessary to actually calculate  $f(x)$ . For example:

**NOTE 15.17 (cont.)**

```

REAL FUNCTION MY_INTEGRATION(N, A, B) RESULT(INTEGRAL)
  ! Integrate f(x)=x^n over [a,b]
  USE ISO_C_BINDING
  INTEGER, INTENT(IN) :: N
  REAL, INTENT(IN) :: A, B

  INTEGRAL = INTEGRATE(MY_F, REAL (A, C_FLOAT), REAL (B, C_FLOAT))
  ! This will call the internal function MY_F to calculate f(x).
  ! The above interface of INTEGRATE needs to be explicit and available.

CONTAINS

  REAL(C_FLOAT) FUNCTION MY_F(X) BIND(C) ! Integrand
    REAL(C_FLOAT), VALUE :: X
    MY_F = X**N ! N is taken from the host instance of MY_INTEGRATION.
  END FUNCTION

END FUNCTION MY_INTEGRATION

```

The function INTEGRATE cannot retain a function pointer to MY\_F and use it after INTEGRATE has finished execution, because the [host instance](#) of MY\_F might no longer exist, making the pointer undefined. If such a pointer is retained, then it can only be used to invoke MY\_F during the execution of the instance of MY\_INTEGRATION that called INTEGRATE.

1 **15.5.2 Actual arguments, dummy arguments, and argument association**

2 **15.5.2.1 Argument correspondence**

- 3 1 In either a subroutine reference or a function reference, the [actual argument](#) list identifies the correspondence  
 4 between the [actual arguments](#) and the [dummy arguments](#) of the procedure. This correspondence may be estab-  
 5 lished either by keyword or by position. If an [argument keyword](#) appears, the [actual argument](#) corresponds to  
 6 the [dummy argument](#) whose name is the same as the [argument keyword](#) (using the [dummy argument](#) names from  
 7 the [interface](#) accessible by the [procedure reference](#)). In the absence of an [argument keyword](#), an [actual argument](#)  
 8 corresponds to the [dummy argument](#) occupying the corresponding position in the reduced [dummy argument](#) list;  
 9 that is, the first [actual argument](#) corresponds to the first [dummy argument](#) in the reduced list, the second [actual](#)  
 10 [argument](#) corresponds to the second [dummy argument](#) in the reduced list, etc. The reduced [dummy argument](#)  
 11 list is either the full [dummy argument](#) list or, if there is a [passed-object dummy argument](#) (7.5.4.5), the [dummy argument](#)  
 12 list with the [passed-object dummy argument](#) omitted. Exactly one [actual argument](#) shall correspond  
 13 to each nonoptional [dummy argument](#). At most one [actual argument](#) shall correspond to each optional [dummy](#)  
 14 [argument](#). Each [actual argument](#) shall correspond to a [dummy argument](#).

**NOTE 15.18**

For example, the procedure defined by

```

SUBROUTINE SOLVE (FUNCT, SOLUTION, METHOD, STRATEGY, PRINT)
  INTERFACE
    FUNCTION FUNCT (X)
      REAL FUNCT, X
    END FUNCTION FUNCT
  END INTERFACE
  REAL SOLUTION
  INTEGER, OPTIONAL :: METHOD, STRATEGY, PRINT
  ...

```

**NOTE 15.18 (cont.)**

can be invoked with

```
CALL SOLVE (FUN, SOL, PRINT = 6)
```

provided its [interface](#) is [explicit](#), and if the [interface](#) is specified by an [interface body](#), the name of the last argument is PRINT.

**15.5.2.2 The passed-object dummy argument and argument correspondence**

- 1 In a reference to a [type-bound procedure](#), or a [procedure pointer](#) component, that has a [passed-object dummy argument](#) (7.5.4.5), the [data-ref](#) of the [function-reference](#) or [call-stmt](#) corresponds, as an [actual argument](#), with the [passed-object dummy argument](#).

**15.5.2.3 Argument association**

- 1 Except in references to intrinsic [inquiry functions](#), a pointer [actual argument](#) that corresponds to a nonoptional nonpointer [dummy argument](#) shall be [pointer associated](#) with a [target](#).
- 2 If a nonpointer [dummy argument](#) without the [VALUE attribute](#) corresponds to a pointer [actual argument](#) that is [pointer associated](#) with a [target](#),
- if the [dummy argument](#) is [polymorphic](#), it becomes argument associated with that [target](#);
  - if the [dummy argument](#) is nonpolymorphic, it becomes argument associated with the [declared type](#) part of that [target](#).
- 3 If a present nonpointer [dummy argument](#) without the [VALUE attribute](#) corresponds to a nonpointer [actual argument](#),
- if the [dummy argument](#) is [polymorphic](#), it becomes argument associated with that [actual argument](#);
  - if the [dummy argument](#) is nonpolymorphic, it becomes argument associated with the [declared type](#) part of that [actual argument](#).
- 4 A present [dummy argument](#) with the [VALUE attribute](#) becomes argument associated with a [definable](#) anonymous data object whose initial value is the value of the [actual argument](#).
- 5 A present pointer [dummy argument](#) that corresponds to a pointer [actual argument](#) becomes argument associated with that [actual argument](#). A present pointer [dummy argument](#) that does not correspond to a pointer [actual argument](#) is not argument associated.
- 6 The entity that is argument associated with a [dummy argument](#) is called its [effective argument](#).
- 7 The [ultimate argument](#) is the [effective argument](#) if the [effective argument](#) is not a [dummy argument](#) or a subobject of a [dummy argument](#). If the [effective argument](#) is a dummy argument, the [ultimate argument](#) is the [ultimate argument](#) of that dummy argument. If the [effective argument](#) is a subobject of a dummy argument, the [ultimate argument](#) is the corresponding subobject of the [ultimate argument](#) of that dummy argument.

**NOTE 15.19**

For the sequence of subroutine calls

```
INTEGER :: X(100)
CALL SUBA (X)
...
SUBROUTINE SUBA(A)
INTEGER :: A(:)
CALL SUBB (A(1:5), A(5:1:-1))
...
SUBROUTINE SUBB(B, C)
```



**NOTE 15.19 (cont.)**

```
INTEGER :: B(:), C(:)
```

the [ultimate argument](#) of B is X(1:5). The [ultimate argument](#) of C is X(5:1:-1) and this is not the same object as the [ultimate argument](#) of B.

**NOTE 15.20**

Fortran [argument association](#) is usually similar to call by reference and call by value-result. If the [VALUE attribute](#) is specified, the effect is as if the [actual argument](#) were assigned to a temporary variable, and that variable were then argument associated with the [dummy argument](#). Subsequent changes to the value or definition status of the dummy argument do not affect the [actual argument](#). The actual mechanism by which this happens is determined by the processor.

1 **15.5.2.4 Ordinary dummy variables**

- 2 1 The requirements in this subclause apply to [actual arguments](#) that correspond to nonallocatable nonpointer  
3 dummy data objects.
- 4 2 The [dummy argument](#) shall be [type compatible](#) with the [actual argument](#). If the [actual argument](#) is a [polymorphic  
5 coindexed object](#), the dummy argument shall not be [polymorphic](#). If the [actual argument](#) is a [polymorphic  
6 assumed-size array](#), the [dummy argument](#) shall be [polymorphic](#). If the [actual argument](#) is of a derived type that  
7 has [type parameters](#), [type-bound procedures](#), or [final subroutines](#), the [dummy argument](#) shall not be [assumed-  
8 type](#).
- 9 3 The [kind type parameter](#) values of the [actual argument](#) shall agree with the corresponding ones of the [dummy  
10 argument](#). The [length type parameter](#) values of a present [actual argument](#) shall agree with the corresponding  
11 ones of the [dummy argument](#) that are not assumed, except for the case of the character length parameter of  
12 an [actual argument](#) of type character with default kind or C character kind ([18.2.2](#)) associated with a [dummy  
13 argument](#) that is not [assumed-shape](#) or [assumed-rank](#).
- 14 4 If a present scalar [dummy argument](#) is of type character with default kind or C character kind, the length *len* of  
15 the [dummy argument](#) shall be less than or equal to the length of the [actual argument](#). The [dummy argument](#)  
16 becomes associated with the leftmost *len* characters of the [actual argument](#). If a present array [dummy argument](#)  
17 is of type character with default kind or C character kind and is not [assumed-shape](#) or [assumed-rank](#), it becomes  
18 associated with the leftmost characters of the [actual argument](#) element sequence ([15.5.2.11](#)).
- 19 5 The values of [assumed type parameters](#) of a [dummy argument](#) are assumed from the corresponding type para-  
20 meters of its [effective argument](#).
- 21 6 If the [actual argument](#) is a [coindexed object](#) with an [allocatable ultimate component](#), the dummy argument shall  
22 have the [INTENT \(IN\)](#) or the [VALUE](#) attribute.

**NOTE 15.21**

If the [actual argument](#) is a [coindexed object](#), a processor that uses distributed memory might create a copy on the executing [image](#) of the [actual argument](#), including copies of any allocated [allocatable subobjects](#), and associate the [dummy argument](#) with that copy. If necessary, on return from the procedure, the value of the copy would be copied back to the [actual argument](#).

- 23 7 Except in references to intrinsic [inquiry functions](#), if the dummy argument is nonoptional and the [actual argument](#)  
24 is [allocatable](#), the corresponding [actual argument](#) shall be allocated.
- 25 8 If the dummy argument does not have the [TARGET attribute](#), any pointers associated with the [effective argument](#)  
26 do not become associated with the corresponding dummy argument on invocation of the procedure. If such a  
27 dummy argument is used as an [actual argument](#) that corresponds to a dummy argument with the [TARGET  
28 attribute](#), whether any pointers associated with the original [effective argument](#) become associated with the dummy  
29 argument with the [TARGET attribute](#) is processor dependent.

- 1 9 If the dummy argument has the **TARGET attribute**, does not have the **VALUE attribute**, and either the **effective**  
2 **argument** is **simply contiguous** or the **dummy argument** is scalar, **assumed-rank**, or **assumed-shape**, and does not  
3 have the **CONTIGUOUS attribute**, and the **effective argument** has the **TARGET attribute** but is not a **coindexed**  
4 **object** or an **array section** with a **vector subscript** then
- 5 • any pointers associated with the **effective argument** become associated with the corresponding dummy  
6 argument on invocation of the procedure, and
  - 7 • when execution of the procedure completes, any pointers that do not become undefined (19.5.2.5) and are  
8 associated with the dummy argument remain associated with the **effective argument**.
- 9 10 If the dummy argument has the **TARGET attribute** and is an **explicit-shape array**, an **assumed-shape array** with  
10 the **CONTIGUOUS attribute**, an **assumed-rank object** with the **CONTIGUOUS attribute**, or an **assumed-size**  
11 **array**, and the **effective argument** has the **TARGET attribute** but is not **simply contiguous** and is not an **array**  
12 **section** with a **vector subscript** then
- 13 • on invocation of the procedure, whether any pointers associated with the **effective argument** become asso-  
14 ciated with the corresponding dummy argument is processor dependent, and
  - 15 • when execution of the procedure completes, the **pointer association** status of any pointer that is **pointer**  
16 **associated** with the **dummy argument** is processor dependent.
- 17 11 If the dummy argument has the **TARGET attribute** and the **effective argument** does not have the **TARGET**  
18 **attribute** or is an **array section** with a **vector subscript**, any pointers associated with the dummy argument  
19 become undefined when execution of the procedure completes.
- 20 12 If the **dummy argument** has the **TARGET attribute** and the **VALUE attribute**, any pointers associated with the  
21 **dummy argument** become undefined when execution of the procedure completes.
- 22 13 If the **actual argument** is a **coindexed** scalar, the corresponding **dummy argument** shall be scalar. If the **actual**  
23 **argument** is a noncoindexed scalar, the corresponding **dummy argument** shall be scalar unless the **actual argument**  
24 is default character, of type character with the C character kind (18.2.2), or is an element or substring of an element  
25 of an array that is not an **assumed-shape**, **pointer**, or **polymorphic** array. If the procedure is **nonelemental** and is  
26 referenced by a generic name or as a defined operator or **defined assignment**, the **ranks** of the **actual arguments**  
27 and corresponding **dummy arguments** shall agree.
- 28 14 If a dummy argument is an **assumed-shape array**, the **rank** of the **actual argument** shall be the same as the **rank**  
29 of the dummy argument, and the **actual argument** shall not be an **assumed-size array**.
- 30 15 An **actual argument** of any **rank** may correspond to an **assumed-rank** dummy argument. The **rank** and **extents** of  
31 the dummy argument are the **rank** and **extents** of the corresponding actual argument. The lower bound of each  
32 dimension of the dummy argument is equal to one. The upper bound is equal to the **extent**, except for the last  
33 dimension when the actual argument is **assumed-size**.
- 34 16 Except when a procedure reference is **elemental** (15.8), each element of an array **actual argument** or of a sequence  
35 in a sequence association (15.5.2.11) is associated with the element of the dummy array that has the same position  
36 in array element order (9.5.3.2).

**NOTE 15.22**

For default character sequence associations, the interpretation of element is provided in 15.5.2.11.

- 37 17 A scalar dummy argument of a **nonelemental** procedure shall correspond only to a scalar actual argument.
- 38 18 If a dummy argument has **INTENT (OUT)** or **INTENT (INOUT)**, the **actual argument** shall be **definable**. If a  
39 dummy argument has **INTENT (OUT)**, the **effective argument** becomes undefined at the time the association is  
40 established, except for **direct components** of an object of derived type for which **default initialization** has been  
41 specified.
- 42 19 If the procedure is **nonelemental**, the **dummy argument** does not have the **VALUE attribute**, and the **actual**

1 [argument](#) is an [array section](#) having a [vector subscript](#), the dummy argument is not [definable](#) and shall not have  
2 the [ASYNCHRONOUS](#), [INTENT \(OUT\)](#), [INTENT \(INOUT\)](#), or [VOLATILE](#) attributes.

3 20 If the [dummy argument](#) has a [coarray ultimate component](#), the corresponding [actual argument](#) shall have the  
4 [VOLATILE](#) attribute if and only if the [dummy argument](#) has the [VOLATILE](#) attribute.

#### NOTE 15.23

Argument intent specifications serve several purposes. See Note 8.18.

#### NOTE 15.24

For more explanatory information on [targets](#) as dummy arguments, see subclause C.10.4.

5 C1538 An [actual argument](#) that is a [coindexed object](#) with the [ASYNCHRONOUS](#) or [VOLATILE](#) attribute shall  
6 not correspond to a [dummy argument](#) that has either the [ASYNCHRONOUS](#) or [VOLATILE](#) attribute,  
7 unless the [dummy argument](#) has the [VALUE](#) attribute.

8 C1539 (R1524) If an [actual argument](#) is a nonpointer array that has the [ASYNCHRONOUS](#) or [VOLATILE](#)  
9 attribute but is not [simply contiguous](#) (9.5.4), and the corresponding [dummy argument](#) has either the  
10 [ASYNCHRONOUS](#) or [VOLATILE](#) attribute, but does not have the [VALUE](#) attribute, that [dummy](#)  
11 [argument](#) shall be [assumed-shape](#) or [assumed-rank](#) and shall not have the [CONTIGUOUS](#) attribute.

12 C1540 (R1524) If an [actual argument](#) is an array pointer that has the [ASYNCHRONOUS](#) or [VOLATILE](#)  
13 attribute but does not have the [CONTIGUOUS](#) attribute, and the corresponding [dummy argument](#)  
14 has either the [ASYNCHRONOUS](#) or [VOLATILE](#) attribute, but does not have the [VALUE](#) attribute,  
15 that [dummy argument](#) shall be an array pointer, an [assumed-shape](#) array without the [CONTIGUOUS](#)  
16 [attribute](#), or an [assumed-rank](#) entity without the [CONTIGUOUS](#) attribute.

#### NOTE 15.25

The constraints on an [actual argument](#) with the [ASYNCHRONOUS](#) or [VOLATILE](#) attribute that corresponds to a dummy argument with either the [ASYNCHRONOUS](#) or [VOLATILE](#) attribute are designed to avoid forcing a processor to use the so-called copy-in/copy-out argument passing mechanism. Making a copy of an [actual argument](#) whose value is likely to change due to an asynchronous input/output operation completing or in some unpredictable manner will cause the new value to be lost when a called procedure returns and the copy-out overwrites the [actual argument](#).

### 17 15.5.2.5 Allocatable and pointer dummy variables

18 1 The requirements in this subclause apply to an [actual argument](#) with the [ALLOCATABLE](#) or [POINTER](#) attribute  
19 that corresponds to a [dummy argument](#) with the same attribute.

20 2 The [actual argument](#) shall be [polymorphic](#) if and only if the associated [dummy argument](#) is [polymorphic](#), and  
21 either both the [actual](#) and [dummy](#) arguments shall be [unlimited polymorphic](#), or the [declared type](#) of the [actual](#)  
22 [argument](#) shall be the same as the [declared type](#) of the [dummy argument](#).

#### NOTE 15.26

The [dynamic type](#) of a [polymorphic allocatable](#) or pointer dummy argument can change as a result of execution of an [ALLOCATE](#) statement or [pointer assignment](#) in the subprogram. Because of this the corresponding [actual argument](#) needs to be [polymorphic](#) and have a [declared type](#) that is the same as the [declared type](#) of the dummy argument or an [extension](#) of that type. However, type compatibility requires that the [declared type](#) of the dummy argument be the same as, or an [extension](#) of, the type of the [actual argument](#). Therefore, the dummy and [actual](#) arguments need to have the same [declared type](#).

[Dynamic type](#) information is not maintained for a nonpolymorphic [allocatable](#) or pointer dummy argument. However, allocating or pointer-assigning such a dummy argument would require maintenance of this information if the corresponding [actual argument](#) is [polymorphic](#). Therefore, the corresponding [actual argument](#) needs to be nonpolymorphic.

1 3 The **rank** of the **actual argument** shall be the same as that of the dummy argument, unless the dummy argument  
2 is **assumed-rank**. The type parameter values of the **actual argument** shall agree with the corresponding ones of  
3 the dummy argument that are not **assumed** or **deferred**. The values of **assumed type parameters** of the **dummy**  
4 **argument** are assumed from the corresponding type parameters of its **effective argument**.

5 4 The **actual argument** shall have **deferred** the same type parameters as the dummy argument.

#### 6 **15.5.2.6 Allocatable dummy variables**

7 1 The requirements in this subclause apply to **actual arguments** that correspond to **allocatable** dummy data objects.

8 2 The **actual argument** shall be **allocatable**. It is permissible for the **actual argument** to have an allocation status  
9 of unallocated.

10 3 The **corank** of the **actual argument** shall be the same as that of the **dummy argument**.

11 4 If the **actual argument** is a **coindexed object**, the dummy argument shall have the **INTENT (IN)** attribute.

12 5 If the **dummy argument** does not have the **TARGET attribute**, any pointers associated with the **actual argument**  
13 do not become associated with the corresponding **dummy argument** on invocation of the procedure. If such a  
14 **dummy argument** is used as an **actual argument** that is associated with a **dummy argument** with the **TARGET**  
15 **attribute**, whether any pointers associated with the original **actual argument** become associated with the **dummy**  
16 **argument** with the **TARGET attribute** is processor dependent.

17 6 If the dummy argument has the **TARGET attribute**, does not have the **INTENT (OUT)** or **VALUE** attribute,  
18 and the corresponding **actual argument** has the **TARGET attribute** then

- 19 • any pointers associated with the **actual argument** become associated with the corresponding dummy argu-  
20 ment on invocation of the procedure, and
- 21 • when execution of the procedure completes, any pointers that do not become undefined (19.5.2.5) and are  
22 associated with the dummy argument remain associated with the **actual argument**.

23 7 If a dummy argument has **INTENT (OUT)** or **INTENT (INOUT)**, the **actual argument** shall be **definable**. If a  
24 dummy argument has **INTENT (OUT)** and its associated **actual argument** is allocated, the **actual argument** is  
25 deallocated on procedure invocation (9.7.3.2).

#### 26 **15.5.2.7 Pointer dummy variables**

27 1 The requirements in this subclause apply to **actual arguments** that correspond to dummy data pointers.

28 C1541 The **actual argument** corresponding to a dummy pointer with the **CONTIGUOUS attribute** shall be  
29 **simply contiguous** (9.5.4).

30 C1542 The **actual argument** corresponding to a dummy pointer shall not be a **coindexed object**.

31 2 If the dummy argument does not have **INTENT (IN)**, the **actual argument** shall be a pointer. Otherwise, the  
32 **actual argument** shall be a pointer or a valid **target** for the dummy pointer in a **pointer assignment statement**. If  
33 the **actual argument** is not a pointer, the dummy pointer becomes **pointer associated** with the **actual argument**.

34 3 If the dummy argument has **INTENT (OUT)**, the **pointer association** status of the **actual argument** becomes  
35 undefined on invocation of the procedure.

#### **NOTE 15.27**

For more explanatory information on pointers as dummy arguments, see subclause C.10.4.

#### 36 **15.5.2.8 Coarray dummy variables**

37 1 If the **dummy argument** is a **coarray**, the corresponding **actual argument** shall be a **coarray** and shall have the  
38 **VOLATILE attribute** if and only if the **dummy argument** has the **VOLATILE attribute**.

- 1 2 If the dummy argument is an array `coarray` that has the `CONTIGUOUS` attribute or is not of `assumed shape`,  
2 the corresponding `actual argument` shall be `simply contiguous` or an element of a `simply contiguous` array.

**NOTE 15.28**

Consider the invocation of a procedure on a particular `image`. Each dummy `coarray` is associated with its `ultimate argument` on the `image`. In addition, during this execution of the procedure, this `image` can access the `coarray` corresponding to the `ultimate argument` on any other `image`. For example, consider

```
INTERFACE
  SUBROUTINE SUB(X)
    REAL :: X[*]
  END SUBROUTINE SUB
END INTERFACE
...
REAL :: A(1000)[*]
...
CALL SUB(A(10))
```

During execution of this invocation of SUB, the executing `image` has access through the syntax X[P] to A(10) on `image` P.

**NOTE 15.29**

Each invocation of a procedure with a nonallocatable `coarray dummy argument` establishes a dummy `coarray` for the `image` with its own bounds and `cobounds`. During this execution of the procedure, this `image` can use its own bounds and `cobounds` to access the `coarray` corresponding to the `ultimate argument` on any other `image`. For example, consider

```
INTERFACE
  SUBROUTINE SUB(X,N)
    INTEGER :: N
    REAL :: X(N,N)[N,*]
  END SUBROUTINE SUB
END INTERFACE
...
REAL :: A(1000)[*]
...
CALL SUB(A,10)
```

During execution of this invocation of SUB, the executing `image` has access through the syntax X(1,2)[3,4] to A(11) on the `image` with `image index` 33.

**NOTE 15.30**

The requirements on an `actual argument` that corresponds to a dummy `coarray` that is not of `assumed-shape` or has the `CONTIGUOUS` attribute are designed to avoid forcing a processor to use the so-called copy-in/copy-out argument passing mechanism.

### 3 15.5.2.9 Actual arguments associated with dummy procedure entities

- 4 1 If the `interface` of a `dummy procedure` is `explicit`, its `characteristics` as a procedure (15.3.1) shall be the same as  
5 those of its `effective argument`, except that a pure `effective argument` may be associated with a dummy argument  
6 that is not pure and an `elemental` intrinsic `actual` procedure may be associated with a `dummy procedure` (which cannot be  
7 `elemental`).
- 8 2 If the `interface` of a `dummy procedure` is `implicit` and either the dummy argument is explicitly typed or referenced

- 1 as a function, it shall not be referenced as a subroutine and any corresponding **actual argument** shall be a function,  
2 function **procedure pointer**, or **dummy procedure**. If both the **actual argument** and **dummy argument** are known  
3 to be functions, they shall have the same type and type parameters. If only the **dummy argument** is known to  
4 be a function, the function that would be invoked by a reference to the **dummy argument** shall have the same  
5 type and type parameters, except that an external function with assumed character length may be associated with a dummy  
6 argument with explicit character length.
- 7 3 If the **interface** of a **dummy procedure** is **implicit** and a reference to it appears as a subroutine reference, any  
8 corresponding **actual argument** shall be a subroutine, subroutine **procedure pointer**, or **dummy procedure**.
- 9 4 If a dummy argument is a **dummy procedure** without the **POINTER attribute**, its **effective argument** shall be an  
10 **external**, **internal**, **module**, or **dummy procedure**, or a specific intrinsic procedure listed in Table 16.2. If the **specific name**  
11 is also a generic name, only the specific procedure is associated with the dummy argument.
- 12 5 If a dummy argument is a **procedure pointer**, the corresponding **actual argument** shall be a **procedure pointer**, a  
13 reference to a function that returns a **procedure pointer**, a reference to the intrinsic function **NULL**, or a valid  
14 **target** for the dummy pointer in a **pointer assignment statement**. If the **actual argument** is not a pointer, the  
15 dummy argument shall have **INTENT (IN)** and becomes **pointer associated** with the **actual argument**.
- 16 6 When the **actual argument** is a procedure, the **host instance** of the **dummy argument** is the **host instance** of the  
17 **actual argument** (15.6.2.4).
- 18 7 If an **external procedure** or a **dummy procedure** is used as an **actual argument**, its **interface** shall be **explicit** or it  
19 shall be explicitly declared to have the **EXTERNAL attribute**.

#### 20 15.5.2.10 Actual arguments and alternate return indicators

- 21 1 If a dummy argument is an asterisk (15.6.2.3), the corresponding **actual argument** shall be an alternate return specifier (R1525).

#### 22 15.5.2.11 Sequence association

- 23 1 An **actual argument** represents an element sequence if it is an array expression, an array element **designator**, a  
24 default character scalar, or a scalar of type character with the C character kind (18.2.2). If the **actual argument** is  
25 an array expression, the element sequence consists of the elements in array element order. If the **actual argument**  
26 is an array element **designator**, the element sequence consists of that array element and each element that follows  
27 it in array element order.
- 28 2 If the **actual argument** is default character or of type character with the C character kind, and is an array  
29 expression, array element, or array element substring **designator**, the element sequence consists of the **storage**  
30 **units** beginning with the first **storage unit** of the **actual argument** and continuing to the end of the array. The  
31 **storage units** of an array element substring **designator** are viewed as array elements consisting of consecutive  
32 groups of **storage units** having the character length of the dummy array.
- 33 3 If the **actual argument** is default character or of type character with the C character kind, and is a scalar that is  
34 not an array element or array element substring **designator**, the element sequence consists of the **storage units** of  
35 the **actual argument**.

#### NOTE 15.31

Some of the elements in the element sequence might consist of **storage units** from different elements of the original array.

- 36 4 An **actual argument** that represents an element sequence and corresponds to a dummy argument that is an array  
37 is sequence associated with the dummy argument if the dummy argument is an **explicit-shape** or **assumed-size**  
38 array. The **rank** and shape of the **actual argument** need not agree with the **rank** and shape of the dummy  
39 argument, but the number of elements in the dummy argument shall not exceed the number of elements in the  
40 element sequence of the **actual argument**. If the dummy argument is **assumed-size**, the number of elements in the  
41 dummy argument is exactly the number of elements in the element sequence.

### 15.5.2.12 Argument presence and restrictions on arguments not present

1 A **dummy argument** or an entity that is **host associated** with a **dummy argument** is not present if the **dummy argument**

- does not correspond to an **actual argument**,
- corresponds to an **actual argument** that is not present, or
- does not have the **ALLOCATABLE** or **POINTER** attribute, and corresponds to an **actual argument** that
  - has the **ALLOCATABLE** attribute and is not allocated, or
  - has the **POINTER** attribute and is **disassociated**;

otherwise, it is present.

2 A nonoptional **dummy argument** shall be present. If an optional nonpointer **dummy argument** corresponds to a present pointer **actual argument**, the **pointer association** status of the **actual argument** shall not be undefined.

3 An optional dummy argument that is not present is subject to the following restrictions.

- (1) If it is a data object, it shall not be referenced or be defined. If it is of a type that has default initialization, the initialization has no effect.
- (2) It shall not be used as the *data-target* or *proc-target* of a **pointer assignment**.
- (3) If it is a procedure or **procedure pointer**, it shall not be invoked.
- (4) It shall not be supplied as an **actual argument** corresponding to a nonoptional dummy argument other than as the argument of the intrinsic function **PRESENT** or as an argument of a function reference that is a **constant expression**.
- (5) A **designator** with it as the **base object** and with one or more subobject selectors shall not be supplied as an **actual argument**.
- (6) If it is an array, it shall not be supplied as an **actual argument** to an **elemental procedure** unless an array of the same **rank** is supplied as an **actual argument** corresponding to a nonoptional dummy argument of that **elemental procedure**.
- (7) If it is a pointer, it shall not be allocated, deallocated, nullified, pointer-assigned, or supplied as an **actual argument** corresponding to an optional nonpointer dummy argument.
- (8) If it is **allocatable**, it shall not be allocated, deallocated, or supplied as an **actual argument** corresponding to an optional nonallocatable **dummy argument**.
- (9) If it has length type parameters, they shall not be the subject of an inquiry.
- (10) It shall not be used as the *selector* in an **ASSOCIATE**, **SELECT RANK**, or **SELECT TYPE** construct.
- (11) It shall not be supplied as the *data-ref* in a *procedure-designator*.
- (12) It shall not be supplied as the *scalar-variable* in a *proc-component-ref*.

4 Except as noted in the list above, it may be supplied as an **actual argument** corresponding to an optional dummy argument, which is then also considered not to be present.

### 15.5.2.13 Restrictions on entities associated with dummy arguments

1 While an entity is associated with a dummy argument, the following restrictions hold.

- (1) Action that affects the allocation status of the entity or a subobject thereof shall be taken through the dummy argument.
- (2) If the allocation status of the entity or a subobject thereof is affected through the dummy argument, then at any time during the invocation and execution of the procedure, either before or after the allocation or deallocation, it shall be referenced only through the dummy argument.
- (3) Action that affects the value of the entity or any subobject of it shall be taken only through the dummy argument unless
  - (a) the dummy argument has the **POINTER** attribute,

- 1 (b) the dummy argument is a scalar, [assumed-shape](#), or [assumed-rank](#) object, and has the [TARGET attribute](#) but not the [INTENT \(IN\)](#) or [CONTIGUOUS](#) attributes, and the [actual argument](#) is a target other than a [coindexed object](#) or an [array section](#) with a [vector subscript](#),
- 2
- 3
- 4 (c) the dummy argument is an [assumed-rank](#) object with the [TARGET attribute](#) and not the
- 5 [INTENT \(IN\) attribute](#), and the [actual argument](#) is a scalar target, or
- 6
- 7 (d) the dummy argument is a [coarray](#) and the action is a coindexed definition of the corresponding
- 8 [ultimate argument coarray](#) by a different [image](#).
- 9
- 10 (4) If the value of the entity or any subobject of it is affected through the dummy argument, then at
- 11 any time during the invocation and execution of the procedure, either before or after the definition,
- 12 it may be referenced only through that dummy argument unless
- 13
- 14 (a) the dummy argument has the [POINTER attribute](#),
- 15
- 16 (b) the dummy argument is a scalar, [assumed-shape](#), or [assumed-rank](#) object, and has the [TARGET attribute](#) but not the [INTENT \(IN\)](#) or [CONTIGUOUS](#) attributes, and the [actual argument](#) is a target other than a [coindexed object](#) or an [array section](#) with a [vector subscript](#),
- 17
- 18 (c) the dummy argument is an [assumed-rank](#) object with the [TARGET attribute](#) and not the
- [INTENT \(IN\) attribute](#), and the [actual argument](#) is a scalar target, or
- (d) the dummy argument is a [coarray](#) and the reference is a coindexed reference of its corresponding
- [ultimate argument coarray](#) by a different [image](#).

**NOTE 15.32**

In

```

SUBROUTINE OUTER
  REAL, POINTER :: A (:)
  ...
  ALLOCATE (A (1:N))
  ...
  CALL INNER (A)
  ...
CONTAINS
  SUBROUTINE INNER (B)
    REAL :: B (:)
    ...
  END SUBROUTINE INNER
  SUBROUTINE SET (C, D)
    REAL, INTENT (OUT) :: C
    REAL, INTENT (IN) :: D
    C = D
  END SUBROUTINE SET
END SUBROUTINE OUTER

```

an [assignment statement](#) such as

```
A (1) = 1.0
```

would not be permitted during the execution of INNER because this would be changing A without using B, but statements such as

```
B (1) = 1.0
```

or

```
CALL SET (B (1), 1.0)
```



**NOTE 15.32 (cont.)**

would be allowed. Similarly,

```
DEALLOCATE (A)
```

would not be allowed because this affects the allocation of B without using B. In this case,

```
DEALLOCATE (B)
```

also would not be permitted. If B were declared with the **POINTER attribute**, either of the statements

```
DEALLOCATE (A)
```

and

```
DEALLOCATE (B)
```

would be permitted, but not both.

**NOTE 15.33**

If there is a partial or complete overlap between the **effective arguments** of two different dummy arguments of the same procedure and the dummy arguments have neither the **POINTER** nor **TARGET** attribute, the overlapped portions cannot be defined, redefined, or become undefined during the execution of the procedure. For example, in

```
CALL SUB (A (1:5), A (3:9))
```

the array section A (3:5) cannot be defined, redefined, or become undefined through the first dummy argument because it is part of the argument associated with the second dummy argument and cannot be defined, redefined, or become undefined through the second dummy argument because it is part of the argument associated with the first dummy argument. The array section A (1:2) remains **definable** through the first dummy argument and A (6:9) remains **definable** through the second dummy argument.

This restriction applies equally to pointer targets. In

```
REAL, DIMENSION (10), TARGET :: A
REAL, DIMENSION (:), POINTER :: B, C
B => A (1:5)
C => A (3:9)
CALL SUB (B, C) ! The dummy arguments of SUB are neither pointers nor targets.
```

the array section B (3:5) cannot be defined because it is part of the argument associated with the second dummy argument. The array section C (1:3) cannot be defined because it is part of the argument associated with the first dummy argument. The array section A (1:2), which is associated with B (1:2), remains **definable** through the first dummy argument and A (6:9), which is associated with C (4:7), remains **definable** through the second dummy argument.

**NOTE 15.34**

In

```
MODULE DATA
  REAL :: W, X, Y, Z
END MODULE DATA

PROGRAM MAIN
```

**NOTE 15.34 (cont.)**

```

USE DATA
...
CALL INIT (X)
...
END PROGRAM MAIN
SUBROUTINE INIT (V)
  USE DATA
  ...
  READ (*, *) V
  ...
END SUBROUTINE INIT

```

variable X cannot be directly referenced at any time during the execution of INIT because it is being defined through the dummy argument V. X can be (indirectly) referenced through V. W, Y, and Z can be directly referenced. X can, of course, be directly referenced once execution of INIT is complete.

**NOTE 15.35**

The restrictions on entities associated with dummy arguments are intended to facilitate a variety of optimizations in the translation of the subprogram, including implementations of [argument association](#) in which the value of an [actual argument](#) that is neither a pointer nor a target is maintained in a register or in local storage.

**NOTE 15.36**

The exception to the aliasing restrictions for dummy [coarrays](#) enables cross-image access while the procedure is executing. Because nonatomic accesses from different [images](#) typically need to be separated by an [image control statement](#), code optimization within segments is not unduly inhibited.

### 1 15.5.3 Function reference

- 2 1 A function is invoked during expression evaluation by a [function-reference](#) or by a defined operation (10.1.6).  
 3 When it is invoked, all [actual argument](#) expressions are evaluated, then the arguments are associated, and then  
 4 the function is executed. When execution of the function is complete, the value of the function result is available  
 5 for use in the expression that caused the function to be invoked. The [characteristics](#) of the function result (15.3.3)  
 6 are determined by the [interface](#) of the function. If a reference to an [elemental function](#) (15.8) is an [elemental](#)  
 7 [reference](#), all array arguments shall have the same shape.

### 8 15.5.4 Subroutine reference

- 9 1 A subroutine is invoked by execution of a [CALL statement](#), execution of a [defined assignment statement](#) (10.2.1.4),  
 10 [defined input/output](#) (12.6.4.8.3), or [finalization](#)(7.5.6). When a subroutine is invoked, all [actual argument](#)  
 11 expressions are evaluated, then the arguments are associated, and then the subroutine is executed. When the  
 12 actions specified by the subroutine are completed, the execution of the [CALL statement](#), the execution of the  
 13 [defined assignment statement](#), the processing of an input or output list item, or [finalization](#) of an object is  
 14 also completed. If a [CALL statement](#) includes one or more alternate return specifiers among its arguments, a branch to one  
 15 of the statements indicated might occur, depending on the action specified by the subroutine. If a reference to an [elemental](#)  
 16 subroutine (15.8) is an [elemental reference](#), at least one [actual argument](#) shall correspond to an [INTENT \(OUT\)](#)  
 17 or [INTENT \(INOUT\)](#) dummy argument, all such [actual arguments](#) shall be arrays, and all [actual arguments](#)  
 18 shall be [conformable](#).

## 15.5.5 Resolving named procedure references

### 15.5.5.1 Establishment of procedure names

1 The rules for interpreting a procedure reference depend on whether the procedure name in the reference is established by the available declarations and specifications to be generic in the **scoping unit** containing the reference, is established to be only specific in the **scoping unit** containing the reference, or is not established.

2 A procedure name is established to be generic in a **scoping unit**

- (1) if that **scoping unit** contains an **interface block** with that name;
- (2) if that **scoping unit** contains a **GENERIC statement** with a *generic-spec* that is that name;
- (3) if that **scoping unit** contains an **INTRINSIC attribute** specification for that name and it is the generic name of an intrinsic procedure;
- (4) if that **scoping unit** contains a **USE statement** that makes that procedure name accessible and the corresponding name in the module is established to be generic; or
- (5) if that **scoping unit** contains no declarations of that name, that **scoping unit** has a **host scoping unit**, and that name is established to be generic in the **host scoping unit**.

3 A procedure name is established to be only specific in a **scoping unit** if it is established to be specific and not established to be generic. It is established to be specific

- (1) if that **scoping unit** contains a **module subprogram**, **internal subprogram**, or **statement function statement** that defines a procedure with that name;
- (2) if that **scoping unit** is of a subprogram that defines a procedure with that name;
- (3) if that **scoping unit** contains an **INTRINSIC attribute** specification for that name and it is the name of a specific intrinsic procedure;
- (4) if that **scoping unit** contains an explicit **EXTERNAL attribute** specification for that name;
- (5) if that **scoping unit** contains a **USE statement** that makes that procedure name accessible and the corresponding name in the module is established to be specific; or
- (6) if that **scoping unit** contains no declarations of that name, that **scoping unit** has a **host scoping unit**, and that name is established to be specific in the **host scoping unit**.

4 A procedure name is not established in a **scoping unit** if it is neither established to be generic nor established to be specific.

### 15.5.5.2 Resolving procedure references to names established to be generic

1 If the reference is consistent with a **nonelemental reference** to one of the **specific interfaces** of a **generic interface** that has that name and either is defined in the **scoping unit** in which the reference appears or is made accessible by a **USE statement** in the **scoping unit**, the reference is to the specific procedure in the **interface block** that provides that **interface**. The rules in 15.4.3.4.5 ensure that there can be at most one such specific procedure.

2 Otherwise, if the reference is consistent with an **elemental reference** to one of the **specific interfaces** of a **generic interface** that has that name and either is defined in the **scoping unit** in which the reference appears or is made accessible by a **USE statement** in the **scoping unit**, the reference is to the specific **elemental procedure** in the **interface block** that provides that **interface**. The rules in 15.4.3.4.5 ensure that there can be at most one such specific **elemental procedure**.

3 Otherwise, if the **scoping unit** contains either an **INTRINSIC attribute** specification for that name or a **USE statement** that makes that name accessible from a module in which the corresponding name is specified to have the **INTRINSIC attribute**, and if the reference is consistent with the **interface** of that **intrinsic procedure**, the reference is to that **intrinsic procedure**.

4 Otherwise, if the **scoping unit** has a **host scoping unit**, the name is established to be generic in that **host scoping unit**, and there is agreement between the **scoping unit** and the **host scoping unit** as to whether the name is a

- 1 function name or a subroutine name, the name is resolved by applying the rules in this subclause to the [host](#)  
 2 [scoping unit](#) as if the reference appeared there.
- 3 5 Otherwise, if the name is that of an intrinsic procedure and the reference is consistent with that intrinsic procedure,  
 4 the reference is to that intrinsic procedure.

**NOTE 15.37**

Because of the renaming facility of the [USE statement](#), the name in the reference can be different from the usual name of the intrinsic procedure.

**NOTE 15.38**

These rules allow particular specific procedures with the same [generic identifier](#) to be used for particular array [ranks](#) and a general [elemental](#) version to be used for other [ranks](#). For example, given an [interface block](#) such as:

```
INTERFACE RANF
  ELEMENTAL FUNCTION SCALAR_RANF(X)
    REAL, INTENT(IN) :: X
  END FUNCTION SCALAR_RANF
  FUNCTION VECTOR_RANDOM(X)
    REAL X(:)
    REAL VECTOR_RANDOM(SIZE(X))
  END FUNCTION VECTOR_RANDOM
END INTERFACE RANF
```

and a declaration such as:

```
REAL A(10,10), AA(10,10)
```

then the statement

```
A = RANF(AA)
```

is an [elemental reference](#) to SCALAR\_RANF. The statement

```
A(6:10,2) = RANF(AA(6:10,2))
```

is a [nonelemental reference](#) to VECTOR\_RANDOM.

### 5 15.5.5.3 Resolving procedure references to names established to be only specific

- 6 1 If the name has the [EXTERNAL attribute](#),
- 7     • if it is a [procedure pointer](#), the reference is to its target;
- 8     • if it is a [dummy procedure](#) that is not a [procedure pointer](#), the reference is to the [effective argument](#)  
 9     corresponding to that name;
- 10    • otherwise, the reference is to the [external procedure](#) with that name.
- 11 2 If the name is that of an accessible [external procedure](#), [internal procedure](#), [module procedure](#), [intrinsic procedure](#),  
 12 or statement function, the reference is to that procedure.

**NOTE 15.39**

Because of the renaming facility of the [USE statement](#), the name in the reference can be different from the original name of the procedure.

#### 1 15.5.5.4 Resolving procedure references to names not established

- 2 1 If the name is the name of a dummy argument of the [scoping unit](#), the dummy argument is a [dummy procedure](#)  
3 and the reference is to that [dummy procedure](#). That is, the procedure invoked by executing that reference is the  
4 [effective argument](#) corresponding to that [dummy procedure](#).
- 5 2 Otherwise, if the name is the name of an intrinsic procedure, and if there is agreement between the reference and  
6 the status of the intrinsic procedure as being a function or subroutine, the reference is to that intrinsic procedure.
- 7 3 Otherwise, the reference is to an [external procedure](#) with that name.

#### 8 15.5.6 Resolving type-bound procedure references

- 9 1 If the *binding-name* in a *procedure-designator* (R1522) is that of a specific [type-bound procedure](#), the procedure  
10 referenced is the one bound to that name in the [dynamic type](#) of the *data-ref*.
- 11 2 If the *binding-name* in a *procedure-designator* is that of a generic type bound procedure, the generic [binding](#) with  
12 that name in the [declared type](#) of the *data-ref* is used to select a specific [binding](#) using the following criteria.
  - 13 • If the reference is consistent with one of the specific [bindings](#) of that generic [binding](#), that specific [binding](#)  
14 is selected.
  - 15 • Otherwise, the reference shall be consistent with an [elemental reference](#) to one of the specific [bindings](#) of  
16 that generic [binding](#); that specific [binding](#) is selected.
- 17 3 The reference is to the procedure bound to the same name as the selected specific [binding](#) in the [dynamic type](#)  
18 of the *data-ref*.

## 19 15.6 Procedure definition

### 20 15.6.1 Intrinsic procedure definition

- 21 1 Intrinsic procedures are defined as an inherent part of the processor. A standard-conforming processor shall  
22 include the intrinsic procedures described in Clause 16, but may include others. However, a standard-conforming  
23 program shall not make use of intrinsic procedures other than those described in Clause 16.

### 24 15.6.2 Procedures defined by subprograms

#### 25 15.6.2.1 General

- 26 1 A procedure is defined by the initial [SUBROUTINE](#) or [FUNCTION](#) statement of a subprogram, and each [ENTRY](#)  
27 [statement](#) defines an additional procedure (15.6.2.6).
- 28 2 A subprogram is specified to be [elemental](#) (15.8), [pure](#) (15.7), recursive, or a separate module subprogram  
29 (15.6.2.5) by a *prefix* in its initial [SUBROUTINE](#) or [FUNCTION](#) statement.

30 R1526 *prefix* is *prefix-spec* [ *prefix-spec* ] ...

31 R1527 *prefix-spec* is *declaration-type-spec*  
32 or ELEMENTAL  
33 or IMPURE  
34 or MODULE  
35 or NON\_RECURSIVE  
36 or PURE  
37 or RECURSIVE

38 C1543 (R1526) A *prefix* shall contain at most one of each *prefix-spec*.

- 1 C1544 (R1526) A *prefix* shall not specify both PURE and IMPURE.
- 2 C1545 (R1526) A *prefix* shall not specify both NON\_RECURSIVE and RECURSIVE.
- 3 C1546 An *elemental procedure* shall not have the *BIND* attribute.
- 4 C1547 (R1526) *MODULE* shall appear only in the *function-stmt* or *subroutine-stmt* of a module subprogram or  
5 of a nonabstract *interface body* that is declared in the *scoping unit* of a module or submodule.
- 6 C1548 (R1526) If *MODULE* appears in the *prefix* of a module subprogram, it shall have been declared to be a  
7 separate *module procedure* in the containing program unit or an ancestor of that program unit.
- 8 C1549 (R1526) If *MODULE* appears in the *prefix* of a module subprogram, the subprogram shall specify the  
9 same *characteristics* and dummy argument names as its corresponding module procedure *interface body*.
- 10 C1550 (R1526) If *MODULE* appears in the *prefix* of a module subprogram and a *binding label* is specified, it  
11 shall be the same as the *binding label* specified in the corresponding module procedure *interface body*.
- 12 C1551 (R1526) If *MODULE* appears in the *prefix* of a module subprogram, NON\_RECURSIVE shall appear  
13 if and only if NON\_RECURSIVE appears in the *prefix* in the corresponding module procedure *interface*  
14 *body*.
- 15 3 The *NON\_RECURSIVE prefix-spec* shall not appear if any procedure defined by the subprogram directly or  
16 indirectly invokes itself or any other procedure defined by the subprogram.
- 17 4 If the *prefix-spec PURE* appears, or the *prefix-spec ELEMENTAL* appears and *IMPURE* does not appear, the  
18 subprogram is a pure subprogram and shall meet the additional constraints of 15.7.
- 19 5 If the *prefix-spec ELEMENTAL* appears, the subprogram is an *elemental subprogram* and shall meet the additional  
20 constraints of 15.8.1.
- 21 R1528 *proc-language-binding-spec* is *language-binding-spec*
- 22 6 A *proc-language-binding-spec* specifies that the procedure defined or declared by the statement is interoperable  
23 (18.3.7).
- 24 C1552 A *proc-language-binding-spec* with a *NAME= specifier* shall not be specified in the *function-stmt* or  
25 *subroutine-stmt* of an *internal procedure*, or of an *interface body* for an *abstract interface* or a *dummy*  
26 *procedure*.
- 27 C1553 If *proc-language-binding-spec* is specified for a function, the function result shall be an *interoperable scalar*  
28 *variable*.
- 29 C1554 If *proc-language-binding-spec* is specified for a procedure, each of its dummy arguments shall be an *inter-*  
30 *operable* procedure (18.3.7) or a variable that is *interoperable* (18.3.5, 18.3.6), *assumed-shape*, *assumed-*  
31 *rank*, *assumed-type*, of type CHARACTER with assumed length, or that has the *ALLOCATABLE* or  
32 *POINTER* attribute.
- 33 C1555 A variable that is a dummy argument of a procedure that has a *proc-language-binding-spec* shall be  
34 *assumed-type* or of *interoperable* type.
- 35 C1556 If *proc-language-binding-spec* is specified for a procedure, it shall not have a *default-initialized dummy*  
36 *argument* with the *ALLOCATABLE* or *POINTER* attribute.
- 37 C1557 If *proc-language-binding-spec* is specified for a procedure, it shall not have a dummy argument that is a  
38 *coarray*.
- 39 C1558 If *proc-language-binding-spec* is specified for a procedure, it shall not have an array dummy argument  
40 with the *VALUE* attribute.

1 **15.6.2.2 Function subprogram**

2 1 A function subprogram is a subprogram that has a FUNCTION statement as its first statement.

3 R1529 *function-subprogram* is *function-stmt*  
 4 [ *specification-part* ]  
 5 [ *execution-part* ]  
 6 [ *internal-subprogram-part* ]  
 7 *end-function-stmt*

8 R1530 *function-stmt* is [ *prefix* ] FUNCTION *function-name* ■  
 9 ■ ( [ *dummy-arg-name-list* ] ) [ *suffix* ]

10 C1559 (R1530) If RESULT appears, *result-name* shall not be the same as *function-name* and shall not be the same  
 11 as the *entry-name* in any ENTRY statement in the subprogram.

12 C1560 (R1530) If RESULT appears, the *function-name* shall not appear in any specification statement in the  
 13 *scoping unit* of the function subprogram.

14 R1531 *dummy-arg-name* is *name*

15 C1561 (R1531) A *dummy-arg-name* shall be the name of a dummy argument.

16 R1532 *suffix* is *proc-language-binding-spec* [ RESULT ( *result-name* ) ]  
 17 or RESULT ( *result-name* ) [ *proc-language-binding-spec* ]

18 R1533 *end-function-stmt* is END [ FUNCTION [ *function-name* ] ]

19 C1562 (R1529) An internal function subprogram shall not contain an *internal-subprogram-part*.

20 C1563 (R1533) If a *function-name* appears in the *end-function-stmt*, it shall be identical to the *function-name*  
 21 specified in the *function-stmt*.

22 2 The name of the function is *function-name*.

23 3 The type and type parameters (if any) of the result of the function defined by a function subprogram may be  
 24 specified by a type specification in the FUNCTION statement or by the name of the *function result* appearing  
 25 in a *type declaration statement* in the specification part of the function subprogram. They shall not be specified  
 26 both ways. If they are not specified either way, they are determined by the implicit typing rules in effect within  
 27 the function subprogram. If the *function result* is an array, *allocatable*, or a pointer, this shall be specified by  
 28 specifications of the name of the *function result* within the function body. The specifications of the *function*  
 29 *result* attributes, the specification of dummy argument attributes, and the information in the procedure heading  
 30 collectively define the *characteristics* of the function (15.3.1).

31 4 If RESULT appears, the name of the *function result* of the function is *result-name* and all occurrences of the  
 32 function name in *execution-part* statements in its scope refer to the function itself. If RESULT does not appear,  
 33 the name of the *function result* is *function-name* and all occurrences of the function name in *execution-part*  
 34 statements in its scope are references to the *function result*. On completion of execution of the function, the value  
 35 returned is that of its *function result*. If the *function result* is a data pointer, the shape of the value returned by  
 36 the function is determined by the shape of the *function result* when the execution of the function is completed. If  
 37 the *function result* is not a pointer, its value shall be defined by the function. If the *function result* is a pointer,  
 38 on return the *pointer association* status of the *function result* shall not be undefined.

**NOTE 15.40**

The *function result* is similar to any other entity (variable or procedure pointer) local to a function subprogram. Its existence begins when execution of the function is initiated and ends when execution of the function is terminated. However, because the final value of this entity is used subsequently in the evaluation of the expression that invoked the function, an implementation might defer releasing the storage occupied

**NOTE 15.40 (cont.)**

by that entity until after its value has been used in expression evaluation.

**NOTE 15.41**

The following is an example of the declaration of an *interface body* with the *BIND attribute*, and a reference to the procedure declared.

```
USE, INTRINSIC :: ISO_C_BINDING

INTERFACE
  FUNCTION JOE (I, J, R) BIND(C,NAME="FrEd")
    USE, INTRINSIC :: ISO_C_BINDING
    INTEGER(C_INT) :: JOE
    INTEGER(C_INT), VALUE :: I, J
    REAL(C_FLOAT), VALUE :: R
  END FUNCTION JOE
END INTERFACE

INT = JOE(1_C_INT, 3_C_INT, 4.0_C_FLOAT)
END PROGRAM
```

The invocation of the function JOE results in a reference to a function with a *binding label* "FrEd". FrEd could be a C function described by the C prototype

```
int FrEd(int n, int m, float x);
```

1 **15.6.2.3 Subroutine subprogram**

2 1 A subroutine subprogram is a subprogram that has a SUBROUTINE statement as its first statement.

3 R1534 *subroutine-subprogram* is *subroutine-stmt*  
 4 [ *specification-part* ]  
 5 [ *execution-part* ]  
 6 [ *internal-subprogram-part* ]  
 7 *end-subroutine-stmt*

8 R1535 *subroutine-stmt* is [ *prefix* ] SUBROUTINE *subroutine-name* ■  
 9 ■ [ ( [ *dummy-arg-list* ] ) [ *proc-language-binding-spec* ] ]

10 C1564 (R1535) The *prefix* of a *subroutine-stmt* shall not contain a *declaration-type-spec*.

11 R1536 *dummy-arg* is *dummy-arg-name*  
 12 or \*

13 R1537 *end-subroutine-stmt* is END [ SUBROUTINE [ *subroutine-name* ] ]

14 C1565 (R1534) An internal subroutine subprogram shall not contain an *internal-subprogram-part*.

15 C1566 (R1537) If a *subroutine-name* appears in the *end-subroutine-stmt*, it shall be identical to the *subroutine-*  
 16 *name* specified in the *subroutine-stmt*.

17 2 The name of the subroutine is *subroutine-name*.



#### 1 15.6.2.4 Instances of a subprogram

2 1 When a procedure defined by a subprogram is invoked, an instance of that subprogram is created. Each instance  
3 has an independent sequence of execution and an independent set of dummy arguments, [unsaved local variables](#),  
4 and [unsaved local procedure pointers](#). Saved local entities are shared by all instances of the subprogram.

5 2 When a statement function is invoked, an instance of that statement function is created.

6 3 When execution of an instance completes it ceases to exist.

7 4 The caller of an instance of a procedure is the instance of the main program, subprogram, or statement function  
8 that invoked it. The call sequence of an instance of a procedure is its caller, followed by the call sequence of its  
9 caller. The call sequence of the main program is empty. The [host instance](#) of an instance of a statement function  
10 or an [internal procedure](#) that is invoked by its name is the first element of the call sequence that is an instance  
11 of the host of the statement function or internal subprogram. The [host instance](#) of an [internal procedure](#) that is  
12 invoked via a [dummy procedure](#) or [procedure pointer](#) is the [host instance](#) of the [associating entity](#) from when the  
13 [argument association](#) or [pointer association](#) was established (19.5.5). The [host instance](#) of a [module procedure](#) is  
14 the module or submodule in which it is defined. A main program or external subprogram has no [host instance](#).

#### 15 15.6.2.5 Separate module procedures

16 1 A separate module procedure is a [module procedure](#) defined by a [separate-module-subprogram](#), by a [function-](#)  
17 [subprogram](#) whose initial statement contains the keyword MODULE, or by a [subroutine-subprogram](#) whose initial  
18 statement contains the keyword MODULE.

19 R1538 *separate-module-subprogram* **is** *mp-subprogram-stmt*  
20 [ *specification-part* ]  
21 [ *execution-part* ]  
22 [ *internal-subprogram-part* ]  
23 *end-mp-subprogram-stmt*

24 R1539 *mp-subprogram-stmt* **is** MODULE PROCEDURE *procedure-name*

25 R1540 *end-mp-subprogram-stmt* **is** END [PROCEDURE [*procedure-name*]]

26 C1567 (R1538) The *procedure-name* shall have been declared to be a separate [module procedure](#) in the containing  
27 [program unit](#) or an ancestor of that [program unit](#).

28 C1568 (R1540) If a *procedure-name* appears in the *end-mp-subprogram-stmt*, it shall be identical to the *procedure-*  
29 *name* in the *mp-subprogram-stmt*.

30 2 A separate [module procedure](#) shall not be defined more than once.

31 3 The [interface](#) of a procedure defined by a [separate-module-subprogram](#) is explicitly declared by the *mp-subprogram-*  
32 *stmt* to be the same as its module procedure [interface body](#). It is recursive if and only if it is declared to be  
33 recursive by the [interface body](#), and if it is a function its result name is determined by the [FUNCTION statement](#)  
34 in the [interface body](#).

#### NOTE 15.42

A separate [module procedure](#) can be accessed by [use association](#) only if its [interface body](#) is declared in the  
specification part of a module and is public.

#### 35 15.6.2.6 ENTRY statement

36 1 An ENTRY statement permits a procedure reference to begin with a particular executable statement within the function or subroutine  
37 subprogram in which the ENTRY statement appears.

38 R1541 *entry-stmt* **is** ENTRY *entry-name* [ ( [ *dummy-arg-list* ] ) [ *suffix* ] ]

- 1 C1569 (R1541) If **RESULT** appears, the *entry-name* shall not appear in any specification or **type declaration statement** in the  
2 **scoping unit** of the function program.
- 3 C1570 (R1541) An *entry-stmt* shall appear only in an *external-subprogram* or a *module-subprogram* that does not define a separate  
4 **module procedure**. An *entry-stmt* shall not appear within an *executable-construct*.
- 5 C1571 (R1541) **RESULT** shall appear only if the *entry-stmt* is in a function subprogram.
- 6 C1572 (R1541) A *dummy-arg* shall not be an alternate return indicator if the **ENTRY** statement is in a function subprogram.
- 7 C1573 (R1541) If **RESULT** appears, *result-name* shall not be the same as the *function-name* in the **FUNCTION statement** and  
8 shall not be the same as the *entry-name* in any **ENTRY** statement in the subprogram.
- 9 2 Optionally, a subprogram may have one or more **ENTRY** statements.
- 10 3 If the **ENTRY** statement is in a function subprogram, an additional function is defined by that subprogram. The name of the  
11 function is *entry-name* and the name of its **result** is *result-name* or is *entry-name* if no *result-name* is provided. The **dummy**  
12 **arguments** of the function are those specified in the **ENTRY** statement. If the **characteristics** of the result of the function named in  
13 the **ENTRY** statement are the same as the **characteristics** of the result of the function named in the **FUNCTION statement**, their  
14 **result** names identify the same entity, although their names need not be the same. Otherwise, they are **storage associated** and shall  
15 all be nonpointer, nonallocatable scalar variables that are default integer, default real, double precision real, default complex, or  
16 default logical.
- 17 4 If the **ENTRY** statement is in a subroutine subprogram, an additional subroutine is defined by that subprogram. The name of the  
18 subroutine is *entry-name*. The dummy arguments of the subroutine are those specified in the **ENTRY** statement.
- 19 5 The order, number, types, kind type parameters, and names of the dummy arguments in an **ENTRY** statement may differ from the  
20 order, number, types, kind type parameters, and names of the dummy arguments in the **FUNCTION** or **SUBROUTINE** statement  
21 in the containing subprogram.
- 22 6 Because an **ENTRY** statement defines an additional function or an additional subroutine, it is referenced in the same manner as any  
23 other function or subroutine (15.5).
- 24 7 In a subprogram, a dummy argument specified in an **ENTRY** statement shall not appear in an executable statement preceding that  
25 **ENTRY** statement, unless it also appears in a **FUNCTION**, **SUBROUTINE**, or **ENTRY** statement that precedes the executable  
26 statement. A function result specified by a *result-name* in an **ENTRY** statement shall not appear in any executable statement that  
27 precedes the first **RESULT** clause with that name.
- 28 8 In a subprogram, a **dummy argument** specified in an **ENTRY** statement shall not appear in the expression of a **statement function**  
29 that precedes the first *dummy-arg* with that name in the subprogram. A function result specified by a *result-name* in an **ENTRY**  
30 statement shall not appear in the expression of a **statement function** that precedes the first **RESULT** clause with that name.
- 31 9 If a **dummy argument** appears in an executable statement, the execution of the executable statement is permitted during the  
32 execution of a reference to the function or subroutine only if the **dummy argument** appears in the **dummy argument** list of the  
33 referenced procedure.
- 34 10 If a **dummy argument** is used in a **specification expression** to specify an array bound or character length of an object, the appearance  
35 of the object in a statement that is executed during a procedure reference is permitted only if the **dummy argument** appears in the  
36 **dummy argument** list of the referenced procedure and it is present (15.5.2.12).
- 37 11 The **NON\_RECURSIVE** and **RECURSIVE** keywords are not used in an **ENTRY** statement. Instead, the presence or absence of  
38 **NON\_RECURSIVE** in the initial **SUBROUTINE** or **FUNCTION** statement controls whether the procedure defined by an **ENTRY**  
39 statement is permitted to reference itself or another procedure defined by the subprogram.
- 40 12 The keywords **PURE** and **IMPURE** are not used in an **ENTRY** statement. Instead, the procedure defined by an **ENTRY** statement  
41 is pure if and only if the subprogram is a pure subprogram.
- 42 13 The keyword **ELEMENTAL** is not used in an **ENTRY** statement. Instead, the procedure defined by an **ENTRY** statement is **elemental**  
43 if and only if **ELEMENTAL** is specified in the **SUBROUTINE** or **FUNCTION** statement.

1 **15.6.2.7 RETURN statement**2 R1542 *return-stmt* is RETURN [ *scalar-int-expr* ]3 C1574 (R1542) The *return-stmt* shall be in the **inclusive scope** of a function or subroutine subprogram.4 C1575 (R1542) The *scalar-int-expr* is allowed only in the **inclusive scope** of a subroutine subprogram.5 1 Execution of the RETURN statement completes execution of the instance of the subprogram in which it appears.  
6 If the expression appears and has a value *n* between 1 and the number of asterisks in the dummy argument list, the **CALL statement**  
7 that invoked the subroutine branches (11.2) to the **branch target statement** identified by the *n<sup>th</sup>* alternate return specifier in the  
8 **actual argument** list of the referenced procedure. If the expression is omitted or has a value outside the required range, there is no  
9 transfer of control to an alternate return.10 2 Execution of an *end-function-stmt*, *end-mp-subprogram-stmt*, or *end-subroutine-stmt* is equivalent to execution  
11 of a RETURN statement with no expression.12 **15.6.2.8 CONTAINS statement**13 R1543 *contains-stmt* is CONTAINS14 1 The CONTAINS statement separates the body of a main program, module, submodule, or subprogram from any  
15 internal or module subprograms it might contain, or it introduces the type-bound procedure part of a derived-type  
16 definition (7.5.5). The CONTAINS statement is not executable.17 **15.6.3 Definition and invocation of procedures by means other than Fortran**18 1 A procedure may be defined by means other than Fortran. The **interface** of a procedure defined by means other  
19 than Fortran may be specified by an **interface body** or **procedure declaration statement**. A reference to such a  
20 procedure is made as though it were defined by an external subprogram.21 2 A procedure defined by means other than Fortran that is invoked by a Fortran procedure and does not cause  
22 **termination of execution** shall return to its caller.**NOTE 15.43**

Examples of code that might cause a transfer of control that bypasses the normal return mechanism of a Fortran procedure are setjmp and longjmp in C and exception handling in other languages. No such behavior is permitted by this document.

23 3 If the **interface** of a procedure has a *proc-language-binding-spec*, the procedure is **interoperable** (18.10).

24 4 Interoperation with C functions is described in 18.10.

**NOTE 15.44**

For explanatory information on definition of procedures by means other than Fortran, see subclause C.10.2.

25 **15.6.4 Statement function**

26 1 A statement function is a function defined by a single statement.

27 R1544 *stmt-function-stmt* is *function-name* ( [ *dummy-arg-name-list* ] ) = *scalar-expr*28 C1576 (R1544) Each *primary* in *scalar-expr* shall be a constant (literal or named), a reference to a variable, a reference to a  
29 function, or an expression in parentheses. Each operation shall be intrinsic. If *scalar-expr* contains a reference to a  
30 function, the reference shall not require an **explicit interface**, the function shall not require an **explicit interface** unless it is  
31 an intrinsic function, the function shall not be a **transformational** intrinsic, and the result shall be scalar. If an argument to  
32 a function is an array, it shall be an array name. If a reference to a statement function appears in *scalar-expr*, its definition  
33 shall have been provided earlier in the **scoping unit** and shall not be the name of the statement function being defined.

- 1 C1577 (R1544) *Named constants* in *scalar-expr* shall have been declared earlier in the *scoping unit* or made accessible by use  
2 or *host* association. If array elements appear in *scalar-expr*, the array shall have been declared as an array earlier in the  
3 *scoping unit* or made accessible by use or *host* association.
- 4 C1578 (R1544) If a *dummy-arg-name*, variable, function reference, or dummy function reference is typed by the implicit typing  
5 rules, its appearance in any subsequent *type declaration statement* shall confirm this implied type and the values of any  
6 implied type parameters.
- 7 C1579 (R1544) The *function-name* and each *dummy-arg-name* shall be specified, explicitly or implicitly, to be scalar.
- 8 C1580 (R1544) A given *dummy-arg-name* shall not appear more than once in a given *dummy-arg-name-list*.
- 9 C1581 A statement function shall not be of a parameterized derived type.
- 10 2 The definition of a statement function with the same name as an accessible entity from the *host* shall be preceded by the declaration  
11 of its type in a *type declaration statement*.
- 12 3 The dummy arguments have a scope of the statement function statement. Each dummy argument has the same type and type  
13 parameters as the entity of the same name in the *scoping unit* containing the statement function statement.
- 14 4 A statement function shall not be supplied as an actual argument.
- 15 5 Execution of a statement function consists of evaluating the expression using the values of the *actual arguments* for the values of the  
16 corresponding dummy arguments and, if necessary, converting the result to the *declared type* and type parameters of the function.
- 17 6 A function reference in the scalar expression shall not cause a dummy argument of the statement function to become redefined or  
18 undefined.

## 19 15.7 Pure procedures

- 20 1 A *pure procedure* is
- 21 • a pure intrinsic procedure (16.1),
  - 22 • a *module procedure* in an intrinsic module, if it is specified to be pure,
  - 23 • defined by a pure subprogram,
  - 24 • a *dummy procedure* that has been specified to be **PURE**, or
  - 25 • a statement function that references only pure functions and does not contain the *designator* of a variable with the **VOLATILE**  
26 *attribute*.
- 27 2 A pure subprogram is a subprogram that has the *prefix-spec* **PURE** or that has the *prefix-spec* **ELEMENTAL**  
28 and does not have the *prefix-spec* **IMPURE**. The following additional constraints apply to pure subprograms.
- 29 C1582 The *specification-part* of a pure function subprogram shall specify that all its nonpointer dummy data  
30 objects have the **INTENT (IN)** or the **VALUE** attribute.
- 31 C1583 The function result of a pure function shall not be such that *finalization* of a reference to the function  
32 would reference an impure procedure.
- 33 C1584 The function result of a pure function shall not be both **polymorphic** and allocatable, or have a **poly-**  
34 **morphic** allocatable *ultimate component*.
- 35 C1585 The *specification-part* of a pure subroutine subprogram shall specify the intents of all its nonpointer  
36 dummy data objects that do not have the **VALUE** attribute.
- 37 C1586 An **INTENT (OUT)** *dummy argument* of a pure procedure shall not be such that *finalization* of the  
38 actual argument would reference an impure procedure.
- 39 C1587 An **INTENT (OUT)** *dummy argument* of a pure procedure shall not be **polymorphic** or have a **poly-**  
40 **morphic** allocatable *ultimate component*.

1 C1588 A **local variable** of a pure subprogram, or of a **BLOCK construct** within a pure subprogram, shall not  
2 have the **SAVE** or **VOLATILE** attribute.

**NOTE 15.45**

Variable initialization in a *type-declaration-stmt* or a *data-stmt* implies the **SAVE attribute**; therefore, such initialization is also disallowed.

3 C1589 The *specification-part* of a pure subprogram shall specify that all its **dummy procedures** are pure.

4 C1590 If a procedure that is neither an intrinsic procedure nor a statement function is used in a context that requires  
5 it to be pure, then its **interface** shall be **explicit** in the scope of that use. The **interface** shall specify that  
6 the procedure is pure.

7 C1591 All internal subprograms in a pure subprogram shall be pure.

8 C1592 A *designator* of a variable with the **VOLATILE** attribute shall not appear in a pure subprogram.

9 C1593 In a pure subprogram any **designator** with a **base object** that is in common or accessed by **host** or **use**  
10 association, is a **pointer dummy argument** of a pure function, is a **dummy argument** with the **INTENT**  
11 **(IN) attribute**, is a **coindexed object**, or an object that is **storage associated** with any such variable, shall  
12 not be used

- 13 (1) in a variable definition context (19.6.7),
- 14 (2) in a pointer association context (19.6.8),
- 15 (3) as the *data-target* in a *pointer-assignment-stmt*,
- 16 (4) as the *expr* corresponding to a component in a *structure-constructor* if the component has the  
17 **POINTER attribute** or has a pointer component at any level of component selection,
- 18 (5) as the *expr* of an **intrinsic assignment statement** in which the variable is of a derived type if the  
19 derived type has a pointer component at any level of component selection,
- 20 (6) as the *source-expr* in a **SOURCE= specifier** if the designator is of a derived type that has a pointer  
21 component at any level of component selection,
- 22 (7) as an **actual argument** corresponding to a dummy argument with the **POINTER attribute**, or
- 23 (8) as the **actual argument** to the function **C\_LOC** from the intrinsic module **ISO\_C\_BINDING**.

**NOTE 15.46**

Item 5 requires that processors be able to determine if entities with the **PRIVATE attribute** or with private components have a pointer component.

24 C1594 Any procedure referenced in a pure subprogram, including one referenced via a **defined operation**, **defined**  
25 **assignment**, **defined input/output**, or **finalization**, shall be pure.

26 C1595 A statement that might result in the deallocation of a **polymorphic** entity is not permitted in a pure  
27 procedure.

**NOTE 15.47**

This includes **intrinsic assignment** to a variable that has a **potential subobject component** that is **polymorphic** and **allocatable**.

28 C1596 A pure subprogram shall not contain a *print-stmt*, *open-stmt*, *close-stmt*, *backspace-stmt*, *endfile-stmt*,  
29 *rewind-stmt*, *flush-stmt*, *wait-stmt*, or *inquire-stmt*.

30 C1597 A pure subprogram shall not contain a *read-stmt* or *write-stmt* whose *io-unit* is a *file-unit-number* or \*.

31 C1598 A pure subprogram shall not contain an **image control statement** (11.6.1).

**NOTE 15.48**

The above constraints are designed to guarantee that a **pure procedure** is free from side effects (modifications of data visible outside the procedure), which means that it is safe to reference it in constructs such as **DO CONCURRENT** and **FORALL**, where there is no explicit order of evaluation.

The constraints on pure subprograms appear to be complicated, but it is not necessary for a programmer to be intimately familiar with them. From the programmer's point of view, these constraints can be summarized as follows: a pure subprogram cannot contain any operation that could conceivably result in an assignment or **pointer assignment** to a common variable, a variable accessed by use or **host** association, or an **INTENT (IN)** dummy argument; nor can a pure subprogram contain any operation that could conceivably perform any **external file** input/output or execute an image control statement (including a **STOP statement**). Note the use of the word conceivably; it is not sufficient for a pure subprogram merely to be side-effect free in practice. For example, a function that contains an assignment to a global variable but in a block that is not executed in any invocation of the function is nevertheless not a pure function. The exclusion of functions of this nature is required if strict compile-time checking is to be used.

It is expected that most library procedures will conform to the constraints required of **pure procedures**, and so can be declared pure and referenced in **DO CONCURRENT** constructs, **FORALL** statements and constructs, and within user-defined **pure procedures**.

**NOTE 15.49**

Pure subroutines are included to allow subroutine calls from **pure procedures** in a safe way, and to allow *forall-assignment-stmts* to be **defined assignments**. The constraints for pure subroutines are based on the same principles as for pure functions, except that side effects to **INTENT (OUT)**, **INTENT (INOUT)**, and pointer dummy arguments are permitted.

**15.8 Elemental procedures****15.8.1 Elemental procedure declaration and interface**

1 An **elemental procedure** is an elemental intrinsic procedure or a procedure that is defined by an **elemental subprogram**. An **elemental procedure** has only **scalar dummy arguments**, but may have array **actual arguments**.

2 An **elemental subprogram** has the *prefix-spec* **ELEMENTAL**. An **elemental subprogram** is a pure subprogram unless it has the *prefix-spec* **IMPURE**. The following additional constraints apply to **elemental subprograms**.

C1599 All **dummy arguments** of an **elemental procedure** shall be scalar noncoarray dummy data objects and shall not have the **POINTER** or **ALLOCATABLE** attribute.

C15100 The **result** of an **elemental** function shall be scalar, and shall not have the **POINTER** or **ALLOCATABLE** attribute.

C15101 The *specification-part* of an **elemental subprogram** shall specify the intents of all of its **dummy arguments** that do not have the **VALUE** attribute.

C15102 In the *specification-expr* that specifies a **type parameter** value of the result of an elemental function, an **object designator** with a **dummy argument** of the function as the **base object** shall appear only as the subject of a specification inquiry (10.1.11), and that specification inquiry shall not depend on a property that is deferred.

3 In a reference to an elemental procedure, if any argument is an array, each **actual argument** that corresponds to an **INTENT (OUT)** or **INTENT (INOUT)** **dummy argument** shall be an array. All **actual arguments** shall be **conformable**. An array **actual argument** is considered to be associated with the scalar **dummy arguments** of the procedure throughout the entire execution of the elemental reference; thus, the restrictions on actions specified in 15.5.2.13 apply to the entirety of the **actual array argument**.

## 1 15.8.2 Elemental function actual arguments and results

- 2 1 If a generic name or a [specific name](#) is used to reference an [elemental](#) function, the shape of the result is the  
3 same as the shape of the [actual argument](#) with the greatest [rank](#). If there are no [actual arguments](#) or the [actual](#)  
4 [arguments](#) are all scalar, the result is scalar. In the array case, the values of the elements, if any, of the result are  
5 the same as would have been obtained if the scalar function had been applied separately, in array element order,  
6 to corresponding elements of each array [actual argument](#).

### NOTE 15.50

An example of an [elemental reference](#) to the intrinsic function [MAX](#):

if X and Y are arrays of shape (M, N),

[MAX](#) (X, 0.0, Y)

is an array expression of shape (M, N) whose elements have values

[MAX](#) (X(I, J), 0.0, Y(I, J)), I = 1, 2, ..., M, J = 1, 2, ..., N

## 7 15.8.3 Elemental subroutine actual arguments

- 8 1 In a reference to an elemental subroutine, if the [actual arguments](#) corresponding to [INTENT \(OUT\)](#) and [INTENT](#)  
9 [\(INOUT\) dummy arguments](#) are arrays, the values of the elements, if any, of the results are the same as would  
10 be obtained if the subroutine had been applied separately, in array element order, to corresponding elements of  
11 each array [actual argument](#).

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## 16 Intrinsic procedures and modules

### 16.1 Classes of intrinsic procedures

1 Intrinsic procedures are divided into eight classes: [inquiry functions](#), [elemental functions](#), [transformational functions](#), [elemental subroutines](#), [pure subroutines](#), [atomic subroutines](#), [collective subroutines](#), and (impure) subroutines.

2 An intrinsic [inquiry function](#) is one whose result depends on the properties of one or more of its arguments instead of their values; in fact, these argument values may be undefined. Unless the description of an intrinsic [inquiry function](#) states otherwise, these arguments are permitted to be unallocated [allocatable](#) variables or pointers that are undefined or [disassociated](#). An [elemental](#) intrinsic function is one that is specified for scalar arguments, but may be applied to array arguments as described in [15.8](#). All other intrinsic functions are [transformational functions](#); they almost all have one or more array arguments or an array result. All standard intrinsic functions are pure.

3 An [atomic subroutine](#) is an intrinsic subroutine that performs an atomic action. The semantics of atomic actions are described in [16.5](#).

4 A [collective subroutine](#) is an intrinsic subroutine that performs a cooperative calculation on a [team](#) of [images](#) without requiring synchronization. The semantics of [collective subroutines](#) are described in [16.6](#).

5 The subroutine [MOVE\\_ALLOC](#) with noncoarray argument FROM, and the [elemental](#) subroutine [MVBITS](#), are [pure](#). No other standard intrinsic subroutine is pure.

6 [Generic names](#) of standard intrinsic procedures are listed in [16.7](#). In most cases, generic functions accept arguments of more than one type and the type of the result is the same as the type of the arguments. [Specific names](#) of standard intrinsic functions with corresponding generic names are listed in [16.8](#).

7 If an intrinsic procedure is used as an [actual argument](#) to a procedure, its [specific name](#) shall be used and it may be referenced in the called procedure only with scalar arguments. If an intrinsic procedure does not have a [specific name](#), it shall not be used as an [actual argument](#) ([15.5.2.9](#)).

8 [Elemental](#) intrinsic procedures behave as described in [15.8](#).

### 16.2 Arguments to intrinsic procedures

#### 16.2.1 General rules

1 All intrinsic procedures may be invoked with either positional arguments or [argument keywords](#) ([15.5](#)). The descriptions in [16.7](#) through [16.9](#) give the [argument keyword](#) names and positional sequence for standard intrinsic procedures.

2 Many of the intrinsic procedures have optional arguments. These arguments are identified by the notation “optional” in the argument descriptions. In addition, the names of the optional arguments are enclosed in square brackets in description headings and in lists of procedures. The valid forms of reference for procedures with optional arguments are described in [15.5.2](#).

#### NOTE 16.1

The text `CMPLX (X [, Y, KIND])` indicates that Y and KIND are both optional arguments. Valid reference forms include `CMPLX(x)`, `CMPLX(x, y)`, `CMPLX(x, KIND=kind)`, `CMPLX(x, y, kind)`, and `CMPLX(KIND=kind, X=x, Y=y)`.

**NOTE 16.2**

Some intrinsic procedures impose additional requirements on their optional arguments. For example, [SELECTED\\_REAL\\_KIND](#) requires that at least one of its optional arguments be present, and [RANDOM\\_SEED](#) requires that at most one of its optional arguments be present.

- 1 3 The dummy arguments of the specific intrinsic procedures in 16.8 have **INTENT (IN)**. The dummy arguments of the intrinsic  
2 procedures in 16.9 have **INTENT (IN)** if the intent is not stated explicitly.
- 3 4 The **actual argument** corresponding to an intrinsic function dummy argument named KIND shall be a scalar  
4 integer **constant expression** and its value shall specify a representation method for the function result that exists  
5 on the processor.
- 6 5 Intrinsic subroutines that assign values to arguments of type character do so in accordance with the rules of  
7 intrinsic assignment (10.2.1.3).
- 8 6 In a reference to the intrinsic subroutine **MVBITS**, the **actual arguments** corresponding to the TO and FROM  
9 dummy arguments may be the same variable and may be associated scalar variables or associated array variables  
10 all of whose corresponding elements are associated. Apart from this, the **actual arguments** in a reference to an  
11 **intrinsic** subroutine shall be such that the execution of the intrinsic subroutine would satisfy the restrictions of  
12 15.5.2.13.
- 13 7 An argument to an intrinsic procedure other than **ASSOCIATED**, **NULL**, or **PRESENT** shall be a data object.

**16.2.2 The shape of array arguments**

- 15 1 Unless otherwise specified, the intrinsic **inquiry functions** accept array arguments for which the shape need not  
16 be defined. The shape of array arguments to **transformational** and **elemental** intrinsic functions shall be defined.

**16.2.3 Mask arguments**

- 18 1 Some array intrinsic functions have an optional MASK argument of type logical that is used by the function to  
19 select the elements of one or more arguments to be operated on by the function. Any element not selected by the  
20 mask need not be defined at the time the function is invoked.
- 21 2 The MASK affects only the value of the function, and does not affect the evaluation, prior to invoking the  
22 function, of arguments that are array expressions.

**16.2.4 DIM arguments and reduction functions**

- 24 1 Some array intrinsic functions are “reduction” functions; that is, they reduce the rank of an array by collapsing  
25 one dimension (or all dimensions, usually producing a scalar result). These functions have a DIM argument that  
26 can specify the dimension to be reduced.
- 27 2 The process of reducing a dimension usually combines the selected elements with a simple operation such as  
28 addition or an intrinsic function such as **MAX**, but more sophisticated reductions are also provided, e.g. by  
29 **COUNT** and **MAXLOC**.

**16.3 Bit model****16.3.1 General**

- 32 1 The bit manipulation procedures are described in terms of a model for the representation and behavior of bits  
33 on a processor.
- 34 2 For the purposes of these procedures, a bit is defined to be a binary digit  $w$  located at position  $k$  of a nonnegative  
35 integer scalar object based on a model nonnegative integer defined by

$$j = \sum_{k=0}^{z-1} w_k \times 2^k$$

1 and for which  $w_k$  may have the value 0 or 1. This defines a sequence of bits  $w_{z-1} \dots w_0$ , with  $w_{z-1}$  the leftmost  
 2 bit and  $w_0$  the rightmost bit. The positions of bits in the sequence are numbered from right to left, with the  
 3 position of the rightmost bit being zero. The length of a sequence of bits is  $z$ . An example of a model number  
 4 compatible with the examples used in 16.4 would have  $z = 32$ , thereby defining a 32-bit integer.

5 3 The interpretation of a negative integer as a sequence of bits is processor dependent.

6 4 The [inquiry function BIT\\_SIZE](#) provides the value of the parameter  $z$  of the model.

7 5 Effectively, this model defines an integer object to consist of  $z$  bits in sequence numbered from right to left from  
 8 0 to  $z - 1$ . This model is valid only in the context of the use of such an object as the argument or result of an  
 9 intrinsic procedure that interprets that object as a sequence of bits. In all other contexts, the model defined for  
 10 an integer in 16.4 applies. In particular, whereas the models are identical for  $r = 2$  and  $w_{z-1} = 0$ , they do not  
 11 correspond for  $r \neq 2$  or  $w_{z-1} = 1$  and the interpretation of bits in such objects is processor dependent.

### 12 16.3.2 Bit sequence comparisons

13 1 When bit sequences of unequal length are compared, the shorter sequence is considered to be extended to the  
 14 length of the longer sequence by padding with zero bits on the left.

15 2 Bit sequences are compared from left to right, one bit at a time, until unequal bits are found or all bits have been  
 16 compared and found to be equal. If unequal bits are found, the sequence with zero in the unequal position is  
 17 considered to be less than the sequence with one in the unequal position. Otherwise the sequences are considered  
 18 to be equal.

### 19 16.3.3 Bit sequences as arguments to INT and REAL

20 1 When a *boz-literal-constant* is the argument A of the intrinsic function [INT](#) or [REAL](#),

- 21 • if the length of the sequence of bits specified by A is less than the size in bits of a scalar variable of the  
 22 same type and kind type parameter as the result, the *boz-literal-constant* is treated as if it were extended  
 23 to a length equal to the size in bits of the result by padding on the left with zero bits, and
- 24 • if the length of the sequence of bits specified by A is greater than the size in bits of a scalar variable of the  
 25 same type and kind type parameter as the result, the *boz-literal-constant* is treated as if it were truncated  
 26 from the left to a length equal to the size in bits of the result.

27 C1601 If a *boz-literal-constant* is truncated as an argument to the intrinsic function [REAL](#), the discarded bits  
 28 shall all be zero.

#### NOTE 16.3

The result values of the intrinsic functions [CMPLX](#) and [DBLE](#) are defined by references to the intrinsic function [REAL](#) with the same arguments. Therefore, the padding and truncation of *boz-literal-constant* arguments to those intrinsic functions is the same as for the intrinsic function [REAL](#).

## 29 16.4 Numeric models

30 1 The numeric manipulation and [inquiry functions](#) are described in terms of a model for the representation and  
 31 behavior of numbers on a processor. The model has parameters that are determined so as to make the model  
 32 best fit the machine on which the program is executed.

33 2 The model set for integer  $i$  is defined by

$$i = s \times \sum_{k=0}^{q-1} w_k \times r^k$$

1 where  $r$  is an integer exceeding one,  $q$  is a positive integer, each  $w_k$  is a nonnegative integer less than  $r$ , and  $s$  is  
 2 +1 or  $-1$ . The integer parameters  $r$  and  $q$  determine the set of model integers.

3 3 The model set for real  $x$  is defined by

$$x = \begin{cases} 0 \text{ or} \\ s \times b^e \times \sum_{k=1}^p f_k \times b^{-k} \end{cases}$$

4 where  $b$  and  $p$  are integers exceeding one; each  $f_k$  is a nonnegative integer less than  $b$ , with  $f_1$  nonzero;  $s$  is  
 5 +1 or  $-1$ ; and  $e$  is an integer that lies between some integer maximum  $e_{\max}$  and some integer minimum  $e_{\min}$   
 6 inclusively. For  $x = 0$ , its exponent  $e$  and digits  $f_k$  are defined to be zero. The integer parameters  $b$ ,  $p$ ,  $e_{\min}$ , and  
 7  $e_{\max}$  determine the set of model floating-point numbers.

8 4 The parameters of the integer and real models are available for each representation method of the integer and  
 9 real types. The parameters characterize the set of available numbers in the definition of the model. Intrinsic  
 10 functions provide the values of some parameters and other values related to the models.

11 5 There is also an extended model set for each kind of real  $x$ ; this extended model is the same as the ordinary  
 12 model except that there are no limits on the range of the exponent  $e$ .

#### NOTE 16.4

Some of the function descriptions use the models

$$i = s \times \sum_{k=0}^{30} w_k \times 2^k$$

and

$$x = 0 \text{ or } s \times 2^e \times \left( \frac{1}{2} + \sum_{k=2}^{24} f_k \times 2^{-k} \right), \quad -126 \leq e \leq 127$$

## 13 16.5 Atomic subroutines

14 1 An [atomic subroutine](#) is an intrinsic subroutine that performs an action on its ATOM argument, and if it has an  
 15 OLD argument, determination of the value to be assigned to that argument, atomically. Definition or evaluation  
 16 of any argument other than ATOM is not performed atomically.

17 2 For any two executions in unordered segments of atomic subroutines whose ATOM argument is the same object,  
 18 the effect is as if one of the executions is performed completely before the other execution begins. Which execution  
 19 is performed first is processor dependent. The sequence of [atomic](#) actions within ordered segments is specified in  
 20 [5.3.5](#). If successive atomic subroutine invocations on image  $P$  redefine a variable atomically in segments  $P_i$  and  
 21  $P_j$ , atomic references to that variable from image  $Q$  in a segment  $Q_k$  that is unordered relative to  $P_i$  and  $P_j$  may  
 22 observe the changes in the value of that variable in any order.

23 3 Atomic operations shall make asynchronous progress. If a variable  $X$  on image  $P$  is defined by an [atomic](#)  
 24 [subroutine](#) on image  $Q$ , image  $R$  repeatedly references  $X$  [ $P$ ] by an atomic subroutine in an unordered segment,  
 25 and no other image defines  $X$  [ $P$ ] in an unordered segment, image  $R$  shall eventually receive the value assigned

- 1 by image *Q*, even if none of the images *P*, *Q*, or *R* execute an **image control statement** until after the definition  
2 of *X* [*P*] by image *Q* and the reception of that value by image *R*.
- 3 4 If the STAT argument is present in an invocation of an atomic subroutine and no error condition occurs, it is  
4 assigned the value zero.
- 5 5 If the STAT argument is present in an invocation of an atomic subroutine and an error condition occurs, any  
6 other argument that is not **INTENT (IN)** becomes undefined. If the ATOM argument is on a failed image, an  
7 error condition occurs and the value **STAT\_FAILED\_IMAGE** from the intrinsic module **ISO\_FORTRAN\_-**  
8 **ENV** is assigned to the STAT argument. If any other error condition occurs, the STAT argument is assigned a  
9 processor-dependent positive value that is different from the value of **STAT\_FAILED\_IMAGE**.
- 10 6 If the STAT argument is not present in an invocation of an atomic subroutine and an error condition occurs,  
11 **error termination** is initiated.

#### NOTE 16.5

The properties of atomic subroutines are intended to support custom synchronization mechanisms. The program will need to handle all possible orderings of sequences of atomic subroutine executions that can arise as a consequence of the above rules; note that the orderings can appear to be different on different images even in the same program execution.

## 12 16.6 Collective subroutines

- 13 1 Successful execution of a **collective subroutine** performs a calculation on all the **images** of the **current team** and  
14 assigns a computed value on one or all of them. If it is invoked by one **image**, it shall be invoked by the same  
15 statement on all **active images** of its **current team** in segments that are not ordered with respect to each other;  
16 corresponding references participate in the same collective computation.
- 17 2 Before execution of the first **CHANGE TEAM** statement on an **image**, in between executions of **CHANGE**  
18 **TEAM** and/or **END TEAM** statements, and after the last execution of an **END TEAM** statement, the sequence  
19 of invocations of **collective subroutines** shall be the same on all **active images** of a **team**. A **collective subroutine**  
20 shall not be referenced when an **image control statement** is not permitted to be executed (for example, in a  
21 procedure invoked from a **CRITICAL** construct).
- 22 C1602 A reference to a **collective subroutine** shall not appear in a context where an **image control statement** is  
23 not permitted to appear.
- 24 3 If the A argument in a reference to a **collective subroutine** is a **coarray**, the corresponding **ultimate arguments** on  
25 all **active images** of the **current team** shall be corresponding **coarrays** as described in 5.4.7.
- 26 4 If the STAT argument is present in a reference to a **collective subroutine** on one image, it shall be present on all  
27 the corresponding references.
- 28 5 If the STAT argument is present in a reference to a **collective subroutine** and its execution is successful, it is  
29 assigned the value zero.
- 30 6 If the STAT argument is present in a reference to a **collective subroutine** and an error condition occurs, it is  
31 assigned a positive value and the A argument becomes undefined.
- 32 7 If the STAT argument is present and the **current team** contains **stopped images**, an error condition occurs and  
33 the value **STAT\_STOPPED\_IMAGE** from the intrinsic module **ISO\_FORTRAN\_ENV** is assigned to the STAT  
34 argument. Otherwise, if the **current team** contained **failed images**, an error condition occurs and the value  
35 **STAT\_FAILED\_IMAGE** from the intrinsic module **ISO\_FORTRAN\_ENV** is assigned to the STAT argument.  
36 If any other error condition occurs, the STAT argument is assigned a processor-dependent positive value that is  
37 different from the value of **STAT\_STOPPED\_IMAGE** or **STAT\_FAILED\_IMAGE**.
- 38 8 If the STAT argument is not present in a reference to a **collective subroutine** and an error condition occurs, error

- 1 termination is initiated.
- 2 9 If the ERRMSG argument is present in a reference to a [collective subroutine](#) and an error condition occurs, it is
- 3 assigned an explanatory message, truncated or padded with blanks if required. If no error condition occurs, the
- 4 definition status and value of ERRMSG are unchanged.

**NOTE 16.6**

The argument A becomes undefined if an error condition occurs during execution of a [collective subroutine](#) because it is intended to allow the processor to use A for intermediate values during calculation.

**NOTE 16.7**

Although the calculations performed by a collective subroutine have some internal synchronizations, a reference to a collective subroutine is not an [image control statement](#).

## 5 16.7 Standard generic intrinsic procedures

- 6 1 For all of the standard intrinsic procedures, the arguments shown are the names that shall be used for [argument](#)
- 7 [keywords](#) if the keyword form is used for [actual arguments](#).

**NOTE 16.8**

For example, a reference to CMPLX can be written in the form CMPLX (A, B, M) or in the form CM-PLX (Y = B, KIND = M, X = A).

**NOTE 16.9**

Many of the [argument keywords](#) have names that are indicative of their usage. For example:

KIND	Describes the kind type parameter of the result
STRING, STRING_A	An arbitrary character string
BACK	Controls the direction of string scan (forward or backward)
MASK	A mask to be applied to the arguments
DIM	A selected dimension of an array argument

- 8 2 In the Class column of Table 16.1,
- 9 A indicates that the procedure is an [atomic subroutine](#),
- 10 C indicates that the procedure is a [collective subroutine](#),
- 11 E indicates that the procedure is an [elemental](#) function,
- 12 ES indicates that the procedure is an [elemental](#) subroutine,
- 13 I indicates that the procedure is an [inquiry function](#),
- 14 PS indicates that the procedure is a [pure](#) subroutine when the FROM argument is not a [coarray](#),
- 15 S indicates that the procedure is an impure subroutine, and
- 16 T indicates that the procedure is in a [transformational function](#).

Table 16.1: **Standard generic intrinsic procedure summary**

Procedure	Arguments	Class	Description
ABS	(A)	E	Absolute value.
ACHAR	(I [, KIND])	E	Character from ASCII code value.
ACOS	(X)	E	Arccosine (inverse cosine) function.
ACOSH	(X)	E	Inverse hyperbolic cosine function.
ADJUSTL	(STRING)	E	Left-justified string value.
ADJUSTR	(STRING)	E	Right-justified string value.

Table 16.1: Standard generic intrinsic procedure summary (cont.)

Procedure	Arguments	Class	Description
AIMAG	(Z)	E	Imaginary part of a complex number.
AINT	(A [, KIND])	E	Truncation toward 0 to a whole number.
ALL	(MASK) or (MASK, DIM)	T	Array reduced by <b>.AND.</b> operator.
ALLOCATED	(ARRAY) or (SCALAR)	I	Allocation status of allocatable variable.
ANINT	(A [, KIND])	E	Nearest whole number.
ANY	(MASK) or (MASK, DIM)	T	Array reduced by <b>.OR.</b> operator.
ASIN	(X)	E	Arcsine (inverse sine) function.
ASINH	(X)	E	Inverse hyperbolic sine function.
ASSOCIATED	(POINTER [, TARGET])	I	<b>Pointer association</b> status inquiry.
ATAN	(X) or (Y, X)	E	Arctangent (inverse tangent) function.
ATAN2	(Y, X)	E	Arctangent (inverse tangent) function.
ATANH	(X)	E	Inverse hyperbolic tangent function.
ATOMIC_ADD	(ATOM, VALUE [, STAT])	A	Atomic addition.
ATOMIC_AND	(ATOM, VALUE [, STAT])	A	Atomic bitwise AND.
ATOMIC_CAS	(ATOM, OLD, COMPARE, VALUE [, STAT])	A	Atomic compare and swap.
ATOMIC_DEFINE	(ATOM, VALUE [, STAT])	A	Define a variable atomically.
ATOMIC_FETCH_ ADD	(ATOM, VALUE, OLD [, STAT])	A	Atomic fetch and add.
ATOMIC_FETCH_ AND	(ATOM, VALUE, OLD [, STAT])	A	Atomic fetch and bitwise AND.
ATOMIC_FETCH_ OR	(ATOM, VALUE, OLD [, STAT])	A	Atomic fetch and bitwise OR.
ATOMIC_FETCH_ XOR	(ATOM, VALUE, OLD [, STAT])	A	Atomic fetch and bitwise exclusive OR.
ATOMIC_OR	(ATOM, VALUE [, STAT])	A	Atomic bitwise OR.
ATOMIC_REF	(VALUE, ATOM [, STAT])	A	Reference a variable atomically.
ATOMIC_XOR	(ATOM, VALUE [, STAT])	A	Atomic bitwise exclusive OR.
BESSEL_J0	(X)	E	Bessel function of the 1 <sup>st</sup> kind, order 0.
BESSEL_J1	(X)	E	Bessel function of the 1 <sup>st</sup> kind, order 1.
BESSEL_JN	(N, X)	E	Bessel function of the 1 <sup>st</sup> kind, order N.
BESSEL_JN	(N1, N2, X)	T	Bessel functions of the 1 <sup>st</sup> kind.
BESSEL_Y0	(X)	E	Bessel function of the 2 <sup>nd</sup> kind, order 0.
BESSEL_Y1	(X)	E	Bessel function of the 2 <sup>nd</sup> kind, order 1.
BESSEL_YN	(N, X)	E	Bessel function of the 2 <sup>nd</sup> kind, order N.
BESSEL_YN	(N1, N2, X)	T	Bessel functions of the 2 <sup>nd</sup> kind.
BGE	(I, J)	E	Bitwise greater than or equal to.
BGT	(I, J)	E	Bitwise greater than.
BIT_SIZE	(I)	I	Number of bits in integer model <a href="#">16.3</a> .
BLE	(I, J)	E	Bitwise less than or equal to.
BLT	(I, J)	E	Bitwise less than.
BTEST	(I, POS)	E	Test single bit in an integer.
CEILING	(A [, KIND])	E	Least integer greater than or equal to A.
CHAR	(I [, KIND])	E	Character from code value.
CMPLX	(X [, KIND]) or (X [, Y, KIND])	E	Conversion to complex type.
CO_BROADCAST	(A, SOURCE_IMAGE [, STAT, ERRMSG])	C	Broadcast value to images.
CO_MAX	(A [, RESULT_IMAGE, STAT, ERRMSG])	C	Compute maximum value across images.
CO_MIN	(A [, RESULT_IMAGE, STAT, ERRMSG])	C	Compute minimum value across images.

Table 16.1: Standard generic intrinsic procedure summary (cont.)

Procedure	Arguments	Class	Description
CO_REDUCE	(A, OPERATION [, RESULT_IMAGE, STAT, ERRMSG])	C	Generalized reduction across images.
CO_SUM	(A [, RESULT_IMAGE, STAT, ERRMSG])	C	Compute sum across images.
COMMAND_ARGUMENT_COUNT	( )	T	Number of command arguments.
CONJG	(Z)	E	Conjugate of a complex number.
COS	(X)	E	Cosine function.
COSH	(X)	E	Hyperbolic cosine function.
COSHAPE	(COARRAY [, KIND])	I	Sizes of codimensions of a coarray.
COUNT	(MASK [, DIM, KIND])	T	Logical array reduced by counting true values.
CPU_TIME	(TIME)	S	Processor time used.
CSHIFT	(ARRAY, SHIFT [, DIM])	T	Circular shift of an array.
DATE_AND_TIME	([DATE, TIME, ZONE, VALUES])	S	Date and time.
DBLE	(A)	E	Conversion to double precision real.
DIGITS	(X)	I	Significant digits in numeric model.
DIM	(X, Y)	E	Maximum of X – Y and zero.
DOT_PRODUCT	(VECTOR_A, VECTOR_B)	T	Dot product of two vectors.
DPROD	(X, Y)	E	Double precision real product.
DSHIFTL	(I, J, SHIFT)	E	Combined left shift.
DSHIFTR	(I, J, SHIFT)	E	Combined right shift.
EOSHIFT	(ARRAY, SHIFT [, BOUNDARY, DIM])	T	End-off shift of the elements of an array.
EPSILON	(X)	I	Model number that is small compared to 1.
ERF	(X)	E	Error function.
ERFC	(X)	E	Complementary error function.
ERFC_SCALED	(X)	E	Scaled complementary error function.
EVENT_QUERY	(EVENT, COUNT [, STAT])	S	Query event count.
EXECUTE_COMMAND_LINE	(COMMAND [, WAIT, EXITSTAT, CMDSTAT, CMDMSG])	S	Execute a command line.
EXP	(X)	E	Exponential function.
EXPONENT	(X)	E	Exponent of floating-point number.
EXTENDS_TYPE_OF	(A, MOLD)	I	Dynamic type extension inquiry.
FAILED_IMAGES	([TEAM, KIND])	T	Indices of failed images.
FINDLOC	(ARRAY, VALUE, DIM [, MASK, KIND, BACK]) or (ARRAY, VALUE [, MASK, KIND, BACK])	T	Location(s) of a specified value.
FLOOR	(A [, KIND])	E	Greatest integer less than or equal to A.
FRACTION	(X)	E	Fractional part of number.
GAMMA	(X)	E	Gamma function.
GET_COMMAND	([COMMAND, LENGTH, STATUS, ERRMSG])	S	Get program invocation command.
GET_COMMAND_ARGUMENT	(NUMBER [, VALUE, LENGTH, STATUS, ERRMSG])	S	Get program invocation argument.



Table 16.1: Standard generic intrinsic procedure summary (cont.)

Procedure	Arguments	Class	Description
GET_ENVIRON- MENT_VARIABLE	(NAME [, VALUE, LENGTH, STATUS, TRIM_NAME, ERRMSG])	S	Get environment variable.
GET_TEAM	([LEVEL])	T	Team.
HUGE	(X)	I	Largest model number.
HYPOT	(X, Y)	E	Euclidean distance function.
IACHAR	(C [, KIND])	E	ASCII code value for character.
IALL	(ARRAY, DIM [, MASK]) or (ARRAY [, MASK])	T	Array reduced by IAND function.
IAND	(I, J)	E	Bitwise AND.
IANY	(ARRAY, DIM [, MASK]) or (ARRAY [, MASK])	T	Array reduced by IOR function.
IBCLR	(I, POS)	E	I with bit POS replaced by zero.
IBITS	(I, POS, LEN)	E	Specified sequence of bits.
IBSET	(I, POS)	E	I with bit POS replaced by one.
ICHAR	(C [, KIND])	E	Code value for character.
IEOR	(I, J)	E	Bitwise exclusive OR.
IMAGE_INDEX	(COARRAY, SUB) or (COARRAY, SUB, TEAM) or (COARRAY, SUB, TEAM_NUMBER)	I	<a href="#">Image index</a> from <a href="#">cosubscripts</a> .
IMAGE_STATUS	(IMAGE [, TEAM])	T	Image execution state.
INDEX	(STRING, SUBSTRING [, BACK, KIND])	E	Character string search.
INT	(A [, KIND])	E	Conversion to integer type.
IOR	(I, J)	E	Bitwise inclusive OR.
IPARITY	(ARRAY, DIM [, MASK]) or (ARRAY [, MASK])	T	Array reduced by IEOB function.
ISHFT	(I, SHIFT)	E	Logical shift.
ISHFTC	(I, SHIFT [, SIZE])	E	Circular shift of the rightmost bits.
IS_CONTIGUOUS	(ARRAY)	I	Array contiguity test (8.5.7).
IS_IOSTAT_END	(I)	E	IOSTAT value test for end of file.
IS_IOSTAT_EOR	(I)	E	IOSTAT value test for end of record.
KIND	(X)	I	Value of the kind type parameter of X.
LBOUND	(ARRAY [, DIM, KIND])	I	Lower bound(s).
LCOBOUND	(COARRAY [, DIM, KIND])	I	Lower <a href="#">cobound</a> (s) of a <a href="#">coarray</a> .
LEADZ	(I)	E	Number of leading zero bits.
LEN	(STRING [, KIND])	I	Length of a character entity.
LEN_TRIM	(STRING [, KIND])	E	Length without trailing blanks.
LGE	(STRING_A, STRING_B)	E	ASCII greater than or equal.
LGT	(STRING_A, STRING_B)	E	ASCII greater than.
LLE	(STRING_A, STRING_B)	E	ASCII less than or equal.
LLT	(STRING_A, STRING_B)	E	ASCII less than.
LOG	(X)	E	Natural logarithm.
LOG_GAMMA	(X)	E	Logarithm of the absolute value of the gamma function.
LOG10	(X)	E	Common logarithm.
LOGICAL	(L [, KIND])	E	Conversion between kinds of logical.
MASKL	(I [, KIND])	E	Left justified mask.
MASKR	(I [, KIND])	E	Right justified mask.
MATMUL	(MATRIX_A, MATRIX_B)	T	Matrix multiplication.
MAX	(A1, A2 [, A3, ...])	E	Maximum value.
MAXEXPONENT	(X)	I	Maximum exponent of a real model.

Table 16.1: Standard generic intrinsic procedure summary (cont.)

Procedure	Arguments	Class	Description
MAXLOC	(ARRAY, DIM [, MASK, KIND, BACK]) or (ARRAY [, MASK, KIND, BACK])	T	Location(s) of maximum value.
MAXVAL	(ARRAY, DIM [, MASK]) or (ARRAY [, MASK])	T	Maximum value(s) of array.
MERGE	(TSOURCE, FSOURCE, MASK)	E	Expression value selection.
MERGE_BITS	(I, J, MASK)	E	Merge of bits under mask.
MIN	(A1, A2 [, A3, ...])	E	Minimum value.
MINEXPONENT	(X)	I	Minimum exponent of a real model.
MINLOC	(ARRAY, DIM [, MASK, KIND, BACK]) or (ARRAY [, MASK, KIND, BACK])	T	Location(s) of minimum value.
MINVAL	(ARRAY, DIM [, MASK]) or (ARRAY [, MASK])	T	Minimum value(s) of array.
MOD	(A, P)	E	Remainder function.
MODULO	(A, P)	E	Modulo function.
MOVE_ALLOC	(FROM, TO [, STAT, ERRMSG])	PS	Move an allocation.
MVBITS	(FROM, FROMPOS, LEN, TO, TOPOS)	ES	Copy a sequence of bits.
NEAREST	(X, S)	E	Adjacent machine number.
NEW_LINE	(A)	I	Newline character.
NINT	(A [, KIND])	E	Nearest integer.
NORM2	(X) or (X, DIM)	T	$L_2$ norm of an array.
NOT	(I)	E	Bitwise complement.
NULL	([MOLD])	T	<a href="#">Disassociated</a> pointer or unallocated allocatable entity.
NUM_IMAGES	( ) or (TEAM) or (TEAM_NUMBER)	T	Number of <a href="#">images</a> .
OUT_OF_RANGE	(X, MOLD [, ROUND])	E	Whether a value cannot be converted safely.
PACK	(ARRAY, MASK [, VECTOR])	T	Array packed into a vector.
PARITY	(MASK) or (MASK, DIM)	T	Array reduced by <a href="#">.NEQV.</a> operator.
POPCNT	(I)	E	Number of one bits.
POPPAR	(I)	E	Parity expressed as 0 or 1.
PRECISION	(X)	I	Decimal precision of a real model.
PRESENT	(A)	I	Presence of optional argument.
PRODUCT	(ARRAY, DIM [, MASK]) or (ARRAY [, MASK])	T	Array reduced by multiplication.
RADIX	(X)	I	Base of a numeric model.
RANDOM_INIT	(REPEATABLE, IMAGE_DISTINCT)	S	Initialise the pseudorandom number generator.
RANDOM_NUMBER	(HARVEST)	S	Generate pseudorandom number(s).
RANDOM_SEED	([SIZE, PUT, GET])	S	Restart or query the pseudorandom number generator.
RANGE	(X)	I	Decimal exponent range of a numeric model ( <a href="#">16.4</a> ).
RANK	(A)	I	Rank of a data object.
REAL	(A [, KIND])	E	Conversion to real type.

Table 16.1: Standard generic intrinsic procedure summary (cont.)

Procedure	Arguments	Class	Description
REDUCE	(ARRAY, OPERATION, DIM [, MASK, IDENTITY, ORDERED]) or (ARRAY, OPERATION [, MASK, IDENTITY, ORDERED])	T	General reduction of array
REPEAT	(STRING, NCOPIES)	T	Repetitive string concatenation.
RESHAPE	(SOURCE, SHAPE [, PAD, ORDER])	T	Arbitrary shape array construction.
RRSPACING	(X)	E	Reciprocal of relative spacing of model numbers.
SAME_TYPE_AS	(A, B)	I	<a href="#">Dynamic type</a> equality test.
SCALE	(X, I)	E	Real number scaled by radix power.
SCAN	(STRING, SET [, BACK, KIND])	E	Character set membership search.
SELECTED_CHAR_- KIND	(NAME)	T	Character kind selection.
SELECTED_INT_- KIND	(R)	T	Integer kind selection.
SELECTED_REAL_- KIND	([P, R, RADIX])	T	Real kind selection.
SET_EXPONENT	(X, I)	E	Real value with specified exponent.
SHAPE	(SOURCE [, KIND])	I	Shape of an array or a scalar.
SHIFTA	(I, SHIFT)	E	Right shift with fill.
SHIFTL	(I, SHIFT)	E	Left shift.
SHIFTR	(I, SHIFT)	E	Right shift.
SIGN	(A, B)	E	Magnitude of A with the sign of B.
SIN	(X)	E	Sine function.
SINH	(X)	E	Hyperbolic sine function.
SIZE	(ARRAY [, DIM, KIND])	I	Size of an array or one extent.
SPACING	(X)	E	Spacing of model numbers.
SPREAD	(SOURCE, DIM, NCOPIES)	T	Value replicated in a new dimension.
SQRT	(X)	E	Square root.
STOPPED_IMAGES	([TEAM, KIND])	T	Indices of stopped images.
STORAGE_SIZE	(A [, KIND])	I	Storage size in bits.
SUM	(ARRAY, DIM [, MASK]) or (ARRAY [, MASK])	T	Array reduced by addition.
SYSTEM_CLOCK	([COUNT, COUNT_RATE, COUNT_MAX])	S	Query system clock.
TAN	(X)	E	Tangent function.
TANH	(X)	E	Hyperbolic tangent function.
TEAM_NUMBER	([TEAM])	T	Team number.
THIS_IMAGE	([TEAM])	T	<a href="#">Index</a> of the invoking <a href="#">image</a> .
THIS_IMAGE	(COARRAY [, TEAM]) or (COARRAY, DIM [, TEAM])	T	<a href="#">Cosubscript(s)</a> for this <a href="#">image</a> .
TINY	(X)	I	Smallest positive model number.
TRAILZ	(I)	E	Number of trailing zero bits.
TRANSFER	(SOURCE, MOLD [, SIZE])	T	Transfer physical representation.
TRANSPOSE	(MATRIX)	T	Transpose of an array of <a href="#">rank</a> two.
TRIM	(STRING)	T	String without trailing blanks.
UBOUND	(ARRAY [, DIM, KIND])	I	Upper bound(s).
UCBOUND	(COARRAY [, DIM, KIND])	I	Upper <a href="#">cobound(s)</a> of a <a href="#">coarray</a> .
UNPACK	(VECTOR, MASK, FIELD)	T	Vector unpacked into an array.

Table 16.1: Standard generic intrinsic procedure summary (cont.)

Procedure	Arguments	Class	Description
VERIFY	(STRING, SET [, BACK, KIND])	E	Character set non-membership search.

- 1 3 The effect of calling EXECUTE\_COMMAND\_LINE on any image other than image 1 in the initial team is  
2 processor dependent.
- 3 4 If RANDOM\_INIT or RANDOM\_SEED is called in a segment A, and RANDOM\_INIT, RANDOM\_SEED, or  
4 RANDOM\_NUMBER is called in segment B, then segments A and B shall be ordered. It is processor dependent  
5 whether each image uses a separate random number generator, or if some or all images use common random  
6 number generators.
- 7 5 The use of all other standard intrinsic procedures in unordered segments is subject only to their argument use  
8 following the rules in 11.6.2.

## 9 16.8 Specific names for standard intrinsic functions

- 10 1 Except for AMAX0, AMIN0, MAX1, and MIN1, the result type of the specific function is the same that the result type of the  
11 corresponding generic function reference would be if it were invoked with the same arguments as the specific function.
- 12 2 A function listed in Table 16.3 is not permitted to be used as an actual argument (15.5.1, C1534), as a target in a procedure  
13 pointer assignment statement (10.2.2.2, C1030), as an initial target in a procedure declaration statement (15.4.3.6, C1519), or to  
14 specify an interface (15.4.3.6, C1515).

Table 16.2: Unrestricted specific intrinsic functions

Specific name	Generic name	Argument type and kind
ABS	ABS	default real
ACOS	ACOS	default real
AIMAG	AIMAG	default complex
AINT	AINT	default real
ALOG	LOG	default real
ALOG10	LOG10	default real
AMOD	MOD	default real
ANINT	ANINT	default real
ASIN	ASIN	default real
ATAN	ATAN (X)	default real
ATAN2	ATAN2	default real
CABS	ABS	default complex
CCOS	COS	default complex
CEXP	EXP	default complex
CLOG	LOG	default complex
CONJG	CONJG	default complex
COS	COS	default real
COSH	COSH	default real
CSIN	SIN	default complex
CSQRT	SQRT	default complex
DABS	ABS	double precision real
DACOS	ACOS	double precision real
DASIN	ASIN	double precision real
DATAN	ATAN	double precision real
DATAN2	ATAN2	double precision real
DCOS	COS	double precision real
DCOSH	COSH	double precision real
DDIM	DIM	double precision real
DEXP	EXP	double precision real
DIM	DIM	default real
DINT	AINT	double precision real
DLOG	LOG	double precision real
DLOG10	LOG10	double precision real
DMOD	MOD	double precision real
DNINT	ANINT	double precision real
DPROD	DPROD	default real

Unrestricted specific intrinsic functions (cont.)

Specific name	Generic name	Argument type and kind
DSIGN	SIGN	double precision real
DSIN	SIN	double precision real
DSINH	SINH	double precision real
DSQRT	SQRT	double precision real
DTAN	TAN	double precision real
DTANH	TANH	double precision real
EXP	EXP	default real
IABS	ABS	default integer
IDIM	DIM	default integer
IDNINT	NINT	double precision real
INDEX	INDEX	default character
ISIGN	SIGN	default integer
LEN	LEN	default character
MOD	MOD	default integer
NINT	NINT	default real
SIGN	SIGN	default real
SIN	SIN	default real
SINH	SINH	default real
SQRT	SQRT	default real
TAN	TAN	default real
TANH	TANH	default real

Table 16.3: Restricted specific intrinsic functions

Specific name	Generic name	Argument type and kind
AMAX0 (...)	REAL (MAX (...))	default integer
AMAX1	MAX	default real
AMIN0 (...)	REAL (MIN (...))	default integer
AMIN1	MIN	default real
CHAR	CHAR	default integer
DMAX1	MAX	double precision real
DMIN1	MIN	double precision real
FLOAT	REAL	default integer
ICHAR	ICHAR	default character
IDINT	INT	double precision real
IFIX	INT	default real
INT	INT	default real
LGE	LGE	default character
LGT	LGT	default character
LLE	LLE	default character
LLT	LLT	default character
MAX0	MAX	default integer
MAX1 (...)	INT (MAX (...))	default real
MIN0	MIN	default integer
MIN1 (...)	INT (MIN (...))	default real
REAL	REAL	default integer
SNGL	REAL	double precision real

## 16.9 Specifications of the standard intrinsic procedures

### 16.9.1 General

- 1 Detailed specifications of the standard generic intrinsic procedures are provided in [16.9](#) in alphabetical order.
- 2 The types and type parameters of standard intrinsic procedure arguments and function results are determined by these specifications. The “Argument(s)” paragraphs specify requirements on the [actual arguments](#) of the procedures. The result [characteristics](#) are sometimes specified in terms of the [characteristics](#) of dummy arguments. A program shall not invoke an intrinsic procedure under circumstances where a value to be assigned to a subroutine argument or returned as a function result is not representable by objects of the specified type and type parameters.

1 3 If an IEEE infinity is assigned or returned by an intrinsic procedure, the intrinsic module `IEEE_ARITHMETIC`  
2 is accessible, and the actual arguments were finite numbers, the flag `IEEE_OVERFLOW` or `IEEE_DIVIDE_-`  
3 `BY_ZERO` shall signal. If an IEEE NaN is assigned or returned, the actual arguments were finite numbers, the  
4 intrinsic module `IEEE_ARITHMETIC` is accessible, and the exception `IEEE_INVALID` is supported, the flag  
5 `IEEE_INVALID` shall signal. If no IEEE infinity or NaN is assigned or returned, these flags shall have the same  
6 status as when the intrinsic procedure was invoked.

### 7 16.9.2 ABS (A)

8 1 **Description.** Absolute value.

9 2 **Class.** `Elemental` function.

10 3 **Argument.** A shall be of type integer, real, or complex.

11 4 **Result Characteristics.** The same as A except that if A is complex, the result is real.

12 5 **Result Value.** If A is of type integer or real, the value of the result is  $|A|$ ; if A is complex with value  $(x, y)$ ,  
13 the result is equal to a processor-dependent approximation to  $\sqrt{x^2 + y^2}$  computed without undue overflow or  
14 underflow.

15 6 **Example.** `ABS ((3.0, 4.0))` has the value 5.0 (approximately).

### 16 16.9.3 ACHAR (I [, KIND])

17 1 **Description.** Character from ASCII code value.

18 2 **Class.** `Elemental` function.

19 3 **Arguments.**

20 I shall be of type integer.

21 KIND (optional) shall be a scalar integer `constant expression`.

22 4 **Result Characteristics.** Character of length one. If KIND is present, the kind type parameter is that specified  
23 by the value of KIND; otherwise, the kind type parameter is that of default character.

24 5 **Result Value.** If I has a value in the range  $0 \leq I \leq 127$ , the result is the character in position I of the `ASCII`  
25 `collating sequence`, provided the processor is capable of representing that character in the character kind of the  
26 result; otherwise, the result is processor dependent. `ACHAR (IACHAR (C))` shall have the value C for any  
27 character C capable of representation as a default character.

28 6 **Example.** `ACHAR (88)` has the value 'X'.

### 29 16.9.4 ACOS (X)

30 1 **Description.** Arccosine (inverse cosine) function.

31 2 **Class.** `Elemental` function.

32 3 **Argument.** X shall be of type real with a value that satisfies the inequality  $|X| \leq 1$ , or of type complex.

33 4 **Result Characteristics.** Same as X.

34 5 **Result Value.** The result has a value equal to a processor-dependent approximation to  $\arccos(X)$ . If it is real  
35 it is expressed in radians and lies in the range  $0 \leq \text{ACOS}(X) \leq \pi$ . If it is complex the real part is expressed in  
36 radians and lies in the range  $0 \leq \text{REAL}(\text{ACOS}(X)) \leq \pi$ .

37 6 **Example.** `ACOS (0.54030231)` has the value 1.0 (approximately).

### 16.9.5 ACOSH (X)

1 **Description.** Inverse hyperbolic cosine function.

2 **Class.** [Elemental](#) function.

3 **Argument.** X shall be of type real or complex.

4 **Result Characteristics.** Same as X.

5 **Result Value.** The result has a value equal to a processor-dependent approximation to the inverse hyperbolic cosine function of X. If the result is complex the imaginary part is expressed in radians and lies in the range  $0 \leq \text{AIMAG}(\text{ACOSH}(X)) \leq \pi$

6 **Example.** ACOSH (1.5430806) has the value 1.0 (approximately).

### 16.9.6 ADJUSTL (STRING)

1 **Description.** Left-justified string value.

2 **Class.** [Elemental](#) function.

3 **Argument.** STRING shall be of type character.

4 **Result Characteristics.** Character of the same length and kind type parameter as STRING.

5 **Result Value.** The value of the result is the same as STRING except that any leading blanks have been deleted and the same number of trailing blanks have been inserted.

6 **Example.** ADJUSTL (' WORD') has the value 'WORD '.

### 16.9.7 ADJUSTR (STRING)

1 **Description.** Right-justified string value.

2 **Class.** [Elemental](#) function.

3 **Argument.** STRING shall be of type character.

4 **Result Characteristics.** Character of the same length and kind type parameter as STRING.

5 **Result Value.** The value of the result is the same as STRING except that any trailing blanks have been deleted and the same number of leading blanks have been inserted.

6 **Example.** ADJUSTR ('WORD ') has the value ' WORD'.

### 16.9.8 AIMAG (Z)

1 **Description.** Imaginary part of a complex number.

2 **Class.** [Elemental](#) function.

3 **Argument.** Z shall be of type complex.

4 **Result Characteristics.** Real with the same kind type parameter as Z.

5 **Result Value.** If Z has the value  $(x, y)$ , the result has the value  $y$ .

6 **Example.** AIMAG ((2.0, 3.0)) has the value 3.0.

### 16.9.9 AINT (A [, KIND])

1 **Description.** Truncation toward 0 to a whole number.

2 **Class.** [Elemental](#) function.

3 **Arguments.**

4 A shall be of type real.

5 KIND (optional) shall be a scalar integer [constant expression](#).

6 **Result Characteristics.** The result is of type real. If KIND is present, the kind type parameter is that specified by the value of KIND; otherwise, the kind type parameter is that of A.

7 **Result Value.** If  $|A| < 1$ , AINT (A) has the value 0; if  $|A| \geq 1$ , AINT (A) has a value equal to the integer whose magnitude is the largest integer that does not exceed the magnitude of A and whose sign is the same as the sign of A.

8 **Examples.** AINT (2.783) has the value 2.0. AINT (-2.783) has the value -2.0.

### 16.9.10 ALL (MASK) or ALL (MASK, DIM)

1 **Description.** Array reduced by [.AND.](#) operator.

2 **Class.** [Transformational](#) function.

3 **Arguments.**

4 MASK shall be a logical array.

5 DIM shall be an integer scalar with value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the [rank](#) of MASK.

6 **Result Characteristics.** The result is of type logical with the same kind type parameter as MASK. It is scalar if DIM does not appear or  $n = 1$ ; otherwise, the result has [rank](#)  $n - 1$  and shape  $[d_1, d_2, \dots, d_{\text{DIM}-1}, d_{\text{DIM}+1}, \dots, d_n]$  where  $[d_1, d_2, \dots, d_n]$  is the shape of MASK.

7 **Result Value.**

8 *Case (i):* The result of ALL (MASK) has the value true if all elements of MASK are true or if MASK has size zero, and the result has value false if any element of MASK is false.

9 *Case (ii):* If MASK has [rank](#) one, ALL (MASK, DIM) is equal to ALL (MASK). Otherwise, the value of element  $(s_1, s_2, \dots, s_{\text{DIM}-1}, s_{\text{DIM}+1}, \dots, s_n)$  of ALL (MASK, DIM) is equal to ALL (MASK  $(s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n)$ ).

10 **Examples.**

11 *Case (i):* The value of ALL ( [.TRUE., .FALSE., .TRUE.] ) is false.

12 *Case (ii):* If B is the array  $\begin{bmatrix} 1 & 3 & 5 \\ 2 & 4 & 6 \end{bmatrix}$  and C is the array  $\begin{bmatrix} 0 & 3 & 5 \\ 7 & 4 & 8 \end{bmatrix}$  then ALL (B /= C, DIM = 1) is [true, false, false] and ALL (B /= C, DIM = 2) is [false, false].

### 16.9.11 ALLOCATED (ARRAY) or ALLOCATED (SCALAR)

1 **Description.** Allocation status of allocatable variable.

2 **Class.** [Inquiry](#) function.

3 **Arguments.**

4 ARRAY shall be an [allocatable](#) array.

5 SCALAR shall be an [allocatable](#) scalar.



1 4 **Result Characteristics.** Default logical scalar.

2 5 **Result Value.** The result has the value true if the argument (ARRAY or SCALAR) is allocated and has the  
3 value false if the argument is unallocated.

#### 4 16.9.12 ANINT (A [, KIND])

5 1 **Description.** Nearest whole number.

6 2 **Class.** [Elemental](#) function.

7 3 **Arguments.**

8 A shall be of type real.

9 KIND (optional) shall be a scalar integer [constant expression](#).

10 4 **Result Characteristics.** The result is of type real. If KIND is present, the kind type parameter is that specified  
11 by the value of KIND; otherwise, the kind type parameter is that of A.

12 5 **Result Value.** The result is the integer nearest A, or if there are two integers equally near A, the result is  
13 whichever such integer has the greater magnitude.

14 6 **Examples.** ANINT (2.783) has the value 3.0. ANINT (−2.783) has the value −3.0.

#### 15 16.9.13 ANY (MASK) or ANY (MASK, DIM)

16 1 **Description.** Array reduced by [.OR.](#) operator.

17 2 **Class.** [Transformational](#) function.

18 3 **Arguments.**

19 MASK shall be a logical array.

20 DIM shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the [rank](#) of MASK.

21 4 **Result Characteristics.** The result is of type logical with the same kind type parameter as MASK. It is scalar  
22 if DIM does not appear or  $n = 1$ ; otherwise, the result has [rank](#)  $n - 1$  and shape  $[d_1, d_2, \dots, d_{\text{DIM}-1}, d_{\text{DIM}+1},$   
23  $\dots, d_n]$  where  $[d_1, d_2, \dots, d_n]$  is the shape of MASK.

24 5 **Result Value.**

25 *Case (i):* The result of ANY (MASK) has the value true if any element of MASK is true and has the value  
26 false if no elements are true or if MASK has size zero.

27 *Case (ii):* If MASK has [rank](#) one, ANY (MASK, DIM) is equal to ANY (MASK). Otherwise, the value of  
28 element  $(s_1, s_2, \dots, s_{\text{DIM}-1}, s_{\text{DIM}+1}, \dots, s_n)$  of ANY (MASK, DIM) is equal to ANY (MASK  $(s_1,$   
29  $s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n)$ ).

30 6 **Examples.**

31 *Case (i):* The value of ANY ([.TRUE.](#), [.FALSE.](#), [.TRUE.](#)) is true.

32 *Case (ii):* If B is the array  $\begin{bmatrix} 1 & 3 & 5 \\ 2 & 4 & 6 \end{bmatrix}$  and C is the array  $\begin{bmatrix} 0 & 3 & 5 \\ 7 & 4 & 8 \end{bmatrix}$  then ANY (B /= C, DIM = 1) is  
33 [true, false, true] and ANY (B /= C, DIM = 2) is [true, true].

#### 34 16.9.14 ASIN (X)

35 1 **Description.** Arcsine (inverse sine) function.

36 2 **Class.** [Elemental](#) function.

37 3 **Argument.** X shall be of type real with a value that satisfies the inequality  $|X| \leq 1$ , or of type complex.

1 4 **Result Characteristics.** Same as X.

2 5 **Result Value.** The result has a value equal to a processor-dependent approximation to  $\arcsin(X)$ . If it is real it  
3 is expressed in radians and lies in the range  $-\pi/2 \leq \text{ASIN}(X) \leq \pi/2$ . If it is complex the real part is expressed  
4 in radians and lies in the range  $-\pi/2 \leq \text{REAL}(\text{ASIN}(X)) \leq \pi/2$ .

5 6 **Example.**  $\text{ASIN}(0.84147098)$  has the value 1.0 (approximately).

## 6 16.9.15 ASINH (X)

7 1 **Description.** Inverse hyperbolic sine function.

8 2 **Class.** [Elemental function](#).

9 3 **Argument.** X shall be of type real or complex.

10 4 **Result Characteristics.** Same as X.

11 5 **Result Value.** The result has a value equal to a processor-dependent approximation to the inverse hyperbolic  
12 sine function of X. If the result is complex the imaginary part is expressed in radians and lies in the range  
13  $-\pi/2 \leq \text{AIMAG}(\text{ASINH}(X)) \leq \pi/2$ .

14 6 **Example.**  $\text{ASINH}(1.1752012)$  has the value 1.0 (approximately).

## 15 16.9.16 ASSOCIATED (POINTER [, TARGET])

16 1 **Description.** [Pointer association](#) status inquiry.

17 2 **Class.** [Inquiry function](#).

18 3 **Arguments.**

19 POINTER shall be a pointer. It may be of any type or may be a procedure pointer. Its [pointer association](#)  
20 status shall not be undefined.

21 TARGET (optional) shall be allowable as the *data-target* or *proc-target* in a [pointer assignment statement \(10.2.2\)](#)  
22 in which POINTER is *data-pointer-object* or *proc-pointer-object*. If TARGET is a pointer then its  
23 [pointer association](#) status shall not be undefined.

24 4 **Result Characteristics.** Default logical scalar.

25 5 **Result Value.**

26 *Case (i):* If TARGET is absent, the result is true if and only if POINTER is associated with a target.

27 *Case (ii):* If TARGET is present and is a procedure, the result is true if and only if POINTER is associated  
28 with TARGET and, if TARGET is an [internal procedure](#), they have the same [host instance](#).

29 *Case (iii):* If TARGET is present and is a procedure pointer, the result is true if and only if POINTER and  
30 TARGET are associated with the same procedure and, if the procedure is an [internal procedure](#),  
31 they have the same [host instance](#).

32 *Case (iv):* If TARGET is present and is a scalar target, the result is true if and only if TARGET is not a zero-  
33 sized [storage sequence](#) and POINTER is associated with a target that occupies the same [storage](#)  
34 [units](#) as TARGET.

35 *Case (v):* If TARGET is present and is an array target, the result is true if and only if POINTER is associated  
36 with a target that has the same shape as TARGET, is neither of size zero nor an array whose elements  
37 are zero-sized [storage sequences](#), and occupies the same [storage units](#) as TARGET in array element  
38 order.

39 *Case (vi):* If TARGET is present and is a scalar pointer, the result is true if and only if POINTER and  
40 TARGET are associated, the targets are not zero-sized [storage sequences](#), and they occupy the  
41 same [storage units](#).

1 *Case (vii):* If TARGET is present and is an **array pointer**, the result is true if and only if POINTER and  
 2 TARGET are both associated, have the same shape, are neither of size zero nor arrays whose  
 3 elements are zero-sized **storage sequences**, and occupy the same **storage units** in array element  
 4 order.

#### NOTE 16.10

The references to TARGET in the above cases are referring to properties that might be possessed by the actual argument, so the case of TARGET being a disassociated pointer will be covered by case *(iii)*, *(vi)*, or *(vii)*.

5 6 **Examples.** ASSOCIATED (CURRENT, HEAD) is true if CURRENT is associated with the target HEAD.  
 6 After the execution of  
 7     A\_PART => A (:N)  
 8 ASSOCIATED (A\_PART, A) is true if N is equal to UBOUND (A, DIM = 1). After the execution of  
 9     NULLIFY (CUR); NULLIFY (TOP)  
 10 ASSOCIATED (CUR, TOP) is false.

### 11 16.9.17 ATAN (X) or ATAN (Y, X)

12 1 **Description.** Arctangent (inverse tangent) function.

13 2 **Class.** **Elemental** function.

14 3 **Arguments.**

15 Y shall be of type real.

16 X If Y appears, X shall be of type real with the same kind type parameter as Y. If Y has the value  
 17 zero, X shall not have the value zero. If Y does not appear, X shall be of type real or complex.

18 4 **Result Characteristics.** Same as X.

19 5 **Result Value.** If Y appears, the result is the same as the result of **ATAN2** (Y,X). If Y does not appear, the  
 20 result has a value equal to a processor-dependent approximation to arctan(X) whose real part is expressed in  
 21 radians and lies in the range  $-\pi/2 \leq \text{ATAN}(X) \leq \pi/2$ .

22 6 **Example.** ATAN (1.5574077) has the value 1.0 (approximately).

### 23 16.9.18 ATAN2 (Y, X)

24 1 **Description.** Arctangent (inverse tangent) function.

25 2 **Class.** **Elemental** function.

26 3 **Arguments.**

27 Y shall be of type real.

28 X shall be of the same type and kind type parameter as Y. If Y has the value zero, X shall not have  
 29 the value zero.

30 4 **Result Characteristics.** Same as X.

31 5 **Result Value.** The result has a value equal to a processor-dependent approximation to the principal value of  
 32 the argument of the complex number (X, Y), expressed in radians. It lies in the range  $-\pi \leq \text{ATAN2}(Y,X)$   
 33  $\leq \pi$  and is equal to a processor-dependent approximation to a value of arctan(Y/X) if  $X \neq 0$ . If  $Y > 0$ , the  
 34 result is positive. If  $Y = 0$  and  $X > 0$ , the result is Y. If  $Y = 0$  and  $X < 0$ , then the result is approximately  
 35  $\pi$  if Y is positive real zero or the processor does not distinguish between positive and negative real zero, and  
 36 approximately  $-\pi$  if Y is negative real zero. If  $Y < 0$ , the result is negative. If  $X = 0$ , the absolute value of the  
 37 result is approximately  $\pi/2$ .

1 6 **Examples.** ATAN2 (1.5574077, 1.0) has the value 1.0 (approximately). If Y has the value  $\begin{bmatrix} 1 & 1 \\ -1 & -1 \end{bmatrix}$  and X  
 2 has the value  $\begin{bmatrix} -1 & 1 \\ -1 & 1 \end{bmatrix}$ , the value of ATAN2 (Y, X) is approximately  $\begin{bmatrix} 3\pi/4 & \pi/4 \\ -3\pi/4 & -\pi/4 \end{bmatrix}$ .

### 3 16.9.19 ATANH (X)

4 1 **Description.** Inverse hyperbolic tangent function.

5 2 **Class.** [Elemental](#) function.

6 3 **Argument.** X shall be of type real or complex.

7 4 **Result Characteristics.** Same as X.

8 5 **Result Value.** The result has a value equal to a processor-dependent approximation to the inverse hyperbolic  
 9 tangent function of X. If the result is complex the imaginary part is expressed in radians and lies in the range  
 10  $-\pi/2 \leq \text{AIMAG}(\text{ATANH}(X)) \leq \pi/2$ .

11 6 **Example.** ATANH (0.76159416) has the value 1.0 (approximately).

### 12 16.9.20 ATOMIC\_ADD (ATOM, VALUE [, STAT])

13 1 **Description.** Atomic addition.

14 2 **Class.** [Atomic subroutine](#).

15 3 **Arguments.**

16 ATOM shall be a scalar [coarray](#) or [coindexed object](#). It shall be of type integer with kind [ATOMIC\\_INT\\_-](#)  
 17 [KIND](#) from the intrinsic module [ISO\\_FORTRAN\\_ENV](#). It is an [INTENT \(INOUT\)](#) argument. If  
 18 an error condition occurs, ATOM becomes undefined; otherwise, it becomes defined with the value  
 19 of ATOM + VALUE.

20 VALUE shall be an integer scalar. It is an [INTENT \(IN\)](#) argument. The values of VALUE and ATOM +  
 21 VALUE shall be representable in kind [ATOMIC\\_INT\\_KIND](#).

22 STAT (optional) shall be a noncoindexed integer scalar with a decimal range of at least four. It is an [INTENT](#)  
 23 [\(OUT\)](#) argument. It is assigned a value as specified in [16.5](#). If an error condition occurs and STAT  
 24 is not present, [error termination](#) is initiated.

25 4 **Example.** CALL ATOMIC\_ADD (I [3], 42) will cause I on [image](#) 3 to become defined with the value 46 if the  
 26 value of I [3] is 4 when the atomic operation is executed.

### 27 16.9.21 ATOMIC\_AND (ATOM, VALUE [, STAT])

28 1 **Description.** Atomic bitwise AND.

29 2 **Class.** [Atomic subroutine](#).

30 3 **Arguments.**

31 ATOM shall be a scalar [coarray](#) or [coindexed object](#). It shall be of type integer with kind [ATOMIC\\_INT\\_-](#)  
 32 [KIND](#) from the intrinsic module [ISO\\_FORTRAN\\_ENV](#). It is an [INTENT \(INOUT\)](#) argument. If  
 33 an error condition occurs, ATOM becomes undefined; otherwise, it becomes defined with the value  
 34 of [IAND](#) (ATOM, [INT](#) (VALUE, [ATOMIC\\_INT\\_KIND](#))).

35 VALUE shall be an integer scalar. It is an [INTENT \(IN\)](#) argument. The value of VALUE shall be repres-  
 36 entable in kind [ATOMIC\\_INT\\_KIND](#).

37 STAT (optional) shall be a noncoindexed integer scalar with a decimal range of at least four. It is an [INTENT](#)  
 38 [\(OUT\)](#) argument. It is assigned a value as specified in [16.5](#). If an error condition occurs and STAT  
 39 is not present, [error termination](#) is initiated.

1 4 **Example.** CALL ATOMIC\_AND (I [3], 6) will cause I on [image 3](#) to become defined with the value 4 if the  
2 value of I [3] is 5 when the atomic operation is executed.

### 3 16.9.22 ATOMIC\_CAS (ATOM, OLD, COMPARE, NEW [, STAT])

4 1 **Description.** Atomic compare and swap.

5 2 **Class.** [Atomic subroutine](#).

6 3 **Arguments.**

7 ATOM shall be a scalar [coarray](#) or [coindexed object](#). It shall be of type integer with kind [ATOMIC\\_INT\\_](#)  
8 [KIND](#) from the intrinsic module [ISO\\_FORTRAN\\_ENV](#), or of type logical with kind [ATOMIC\\_](#)  
9 [LOGICAL\\_KIND](#) from the intrinsic module [ISO\\_FORTRAN\\_ENV](#). It is an [INTENT \(INOUT\)](#)  
10 argument. If an error condition occurs, ATOM becomes undefined; otherwise, if ATOM is of type  
11 integer and equal to COMPARE, or of type logical and equivalent to COMPARE, it becomes defined  
12 with the value of NEW.

13 OLD shall be scalar and of the same type and kind as ATOM. It is an [INTENT \(OUT\)](#) argument. If  
14 an error condition occurs, it becomes undefined; otherwise, it becomes defined with the value that  
15 ATOM had at the start of the atomic operation.

16 COMPARE shall be scalar and of the same type and kind as ATOM. It is an [INTENT \(IN\)](#) argument.

17 NEW shall be scalar and of the same type and kind as ATOM. It is an [INTENT \(IN\)](#) argument.

18 STAT (optional) shall be a noncoindexed integer scalar with a decimal range of at least four. It is an [INTENT](#)  
19 [\(OUT\)](#) argument. It is assigned a value as specified in [16.5](#). If an error condition occurs and STAT  
20 is not present, [error termination](#) is initiated.

21 4 **Example.** If the value of I on [image 3](#) is equal to 13 at the beginning of the atomic operation performed by  
22 CALL ATOMIC\_CAS (I [3], OLD, 0, 1), the value of I on [image 3](#) will be unchanged, and OLD will become  
23 defined with the value 13. If the value of I on [image 3](#) is equal to 0 at the beginning of the atomic operation  
24 performed by CALL ATOMIC\_CAS (I [3], OLD, 0, 1), I on [image 3](#) will become defined with the value 1, and  
25 OLD will become defined with the value 0.

### 26 16.9.23 ATOMIC\_DEFINE (ATOM, VALUE [, STAT])

27 1 **Description.** Define a variable atomically.

28 2 **Class.** [Atomic subroutine](#).

29 3 **Arguments.**

30 ATOM shall be a scalar [coarray](#) or [coindexed object](#). It shall be of type integer with kind [ATOMIC\\_INT\\_](#)  
31 [KIND](#) from the intrinsic module [ISO\\_FORTRAN\\_ENV](#), or of type logical with kind [ATOMIC\\_](#)  
32 [LOGICAL\\_KIND](#) from the intrinsic module [ISO\\_FORTRAN\\_ENV](#). It is an [INTENT \(OUT\)](#)  
33 argument. On successful execution, it becomes defined with the value of VALUE. If an error  
34 condition occurs, it becomes undefined.

35 VALUE shall be scalar and of the same type as ATOM. It is an [INTENT \(IN\)](#) argument.

36 STAT (optional) shall be a noncoindexed integer scalar with a decimal range of at least four. It is an [INTENT](#)  
37 [\(OUT\)](#) argument. It is assigned a value as specified in [16.5](#). If an error condition occurs and STAT  
38 is not present, [error termination](#) is initiated.

39 4 **Example.** CALL ATOMIC\_DEFINE (I [3], 4) causes I on [image 3](#) to become defined with the value 4.

### 40 16.9.24 ATOMIC\_FETCH\_ADD (ATOM, VALUE, OLD [, STAT])

41 1 **Description.** Atomic fetch and add.

42 2 **Class.** [Atomic subroutine](#).

1 3 **Arguments.**

2 ATOM shall be a scalar [coarray](#) or [coindexed object](#). It shall be of type integer with kind [ATOMIC\\_INT\\_KIND](#) from the intrinsic module [ISO\\_FORTRAN\\_ENV](#). It is an [INTENT \(INOUT\)](#) argument. If  
3 an error condition occurs, ATOM becomes undefined; otherwise, it becomes defined with the value  
4 of ATOM + VALUE.

5  
6 VALUE shall be an integer scalar. It is an [INTENT \(IN\)](#) argument. The values of VALUE and ATOM +  
7 VALUE shall be representable in kind [ATOMIC\\_INT\\_KIND](#).

8 OLD shall be scalar and of the same type and kind as ATOM. It is an [INTENT \(OUT\)](#) argument. If  
9 an error condition occurs, it becomes undefined; otherwise, it becomes defined with the value that  
10 ATOM had at the start of the atomic operation.

11 STAT (optional) shall be a noncoindexed integer scalar with a decimal range of at least four. It is an [INTENT](#)  
12 [\(OUT\)](#) argument. It is assigned a value as specified in [16.5](#). If an error condition occurs and STAT  
13 is not present, [error termination](#) is initiated.

14 4 **Example.** CALL ATOMIC\_FETCH\_ADD (I [3], 7, J) will cause I on [image 3](#) to become defined with the value  
15 12, and J to become defined with the value 5, if the value of I [3] is 5 when the atomic operation is executed.

16 **16.9.25 ATOMIC\_FETCH\_AND (ATOM, VALUE, OLD [, STAT])**

17 1 **Description.** Atomic fetch and bitwise AND.

18 2 **Class.** [Atomic subroutine](#).

19 3 **Arguments.**

20 ATOM shall be a scalar [coarray](#) or [coindexed object](#). It shall be of type integer with kind [ATOMIC\\_INT\\_KIND](#)  
21 from the intrinsic module [ISO\\_FORTRAN\\_ENV](#). It is an [INTENT \(INOUT\)](#) argument. If  
22 an error condition occurs, ATOM becomes undefined; otherwise, it becomes defined with the value  
23 of [IAND](#) (ATOM, [INT](#) (VALUE, [ATOMIC\\_INT\\_KIND](#))).

24 VALUE shall be an integer scalar. It is an [INTENT \(IN\)](#) argument. The value of VALUE shall be repre-  
25 sentable in kind [ATOMIC\\_INT\\_KIND](#).

26 OLD shall be scalar and of the same type and kind as ATOM. It is an [INTENT \(OUT\)](#) argument. If  
27 an error condition occurs, it becomes undefined; otherwise, it becomes defined with the value that  
28 ATOM had at the start of the atomic operation.

29 STAT (optional) shall be a noncoindexed integer scalar with a decimal range of at least four. It is an [INTENT](#)  
30 [\(OUT\)](#) argument. It is assigned a value as specified in [16.5](#). If an error condition occurs and STAT  
31 is not present, [error termination](#) is initiated.

32 4 **Example.** CALL ATOMIC\_FETCH\_AND (I [3], 6, J) will cause I on [image 3](#) to become defined with the value  
33 4, and J to become defined with the value 5, if the value of I [3] is 5 when the atomic operation is executed.

34 **16.9.26 ATOMIC\_FETCH\_OR (ATOM, VALUE, OLD [, STAT])**

35 1 **Description.** Atomic fetch and bitwise OR.

36 2 **Class.** [Atomic subroutine](#).

37 3 **Arguments.**

38 ATOM shall be a scalar [coarray](#) or [coindexed object](#). It shall be of type integer with kind [ATOMIC\\_INT\\_KIND](#)  
39 from the intrinsic module [ISO\\_FORTRAN\\_ENV](#). It is an [INTENT \(INOUT\)](#) argument. If  
40 an error condition occurs, ATOM becomes undefined; otherwise, it becomes defined with the value  
41 of [IOR](#) (ATOM, [INT](#) (VALUE, [ATOMIC\\_INT\\_KIND](#))).

42 VALUE shall be an integer scalar. It is an [INTENT \(IN\)](#) argument. The value of VALUE shall be repre-  
43 sentable in kind [ATOMIC\\_INT\\_KIND](#).

44 OLD shall be scalar and of the same type and kind as ATOM. It is an [INTENT \(OUT\)](#) argument. If

1 an error condition occurs, it becomes undefined; otherwise, it becomes defined with the value that  
2 ATOM had at the start of the atomic operation.

3 STAT (optional) shall be a noncoindexed integer scalar with a decimal range of at least four. It is an **INTENT**  
4 **(OUT)** argument. It is assigned a value as specified in 16.5. If an error condition occurs and STAT  
5 is not present, **error termination** is initiated.

6 4 **Example.** CALL ATOMIC\_FETCH\_OR (I [3], 1, J) will cause I on **image 3** to become defined with the value  
7 3, and J to become defined with the value 2, if the value of I [3] is 2 when the atomic operation is executed.

## 8 16.9.27 ATOMIC\_FETCH\_XOR (ATOM, VALUE, OLD [, STAT])

9 1 **Description.** Atomic fetch and bitwise exclusive OR.

10 2 **Class.** **Atomic subroutine.**

11 3 **Arguments.**

12 ATOM shall be a scalar **coarray** or **coindexed object**. It shall be of type integer with kind **ATOMIC\_INT\_-**  
13 **KIND** from the intrinsic module **ISO\_FORTRAN\_ENV**. It is an **INTENT (INOUT)** argument. If  
14 an error condition occurs, ATOM becomes undefined; otherwise, it becomes defined with the value  
15 of **IEOR** (ATOM, **INT** (VALUE, **ATOMIC\_INT\_KIND**)).

16 VALUE shall be an integer scalar. It is an **INTENT (IN)** argument. The value of VALUE shall be repres-  
17 entable in kind **ATOMIC\_INT\_KIND**.

18 OLD shall be scalar and of the same type and kind as ATOM. It is an **INTENT (OUT)** argument. If  
19 an error condition occurs, it becomes undefined; otherwise, it becomes defined with the value that  
20 ATOM had at the start of the atomic operation.

21 STAT (optional) shall be a noncoindexed integer scalar with a decimal range of at least four. It is an **INTENT**  
22 **(OUT)** argument. It is assigned a value as specified in 16.5. If an error condition occurs and STAT  
23 is not present, **error termination** is initiated.

24 4 **Example.** CALL ATOMIC\_FETCH\_XOR (I [3], 1, J) will cause I on **image 3** to become defined with the value  
25 2, and J to become defined with the value 3, if the value of I [3] is 3 when the atomic operation is executed.

## 26 16.9.28 ATOMIC\_OR (ATOM, VALUE [, STAT])

27 1 **Description.** Atomic bitwise OR.

28 2 **Class.** **Atomic subroutine.**

29 3 **Arguments.**

30 ATOM shall be a scalar **coarray** or **coindexed object**. It shall be of type integer with kind **ATOMIC\_INT\_-**  
31 **KIND** from the intrinsic module **ISO\_FORTRAN\_ENV**. It is an **INTENT (INOUT)** argument. If  
32 an error condition occurs, ATOM becomes undefined; otherwise, it becomes defined with the value  
33 of **IOR** (ATOM, **INT** (VALUE, **ATOMIC\_INT\_KIND**)).

34 VALUE shall be an integer scalar. It is an **INTENT (IN)** argument. The value of VALUE shall be repres-  
35 entable in kind **ATOMIC\_INT\_KIND**.

36 STAT (optional) shall be a noncoindexed integer scalar with a decimal range of at least four. It is an **INTENT**  
37 **(OUT)** argument. It is assigned a value as specified in 16.5. If an error condition occurs and STAT  
38 is not present, **error termination** is initiated.

39 4 **Example.** CALL ATOMIC\_OR (I [3], 1) will cause I on **image 3** to become defined with the value 3 if the value  
40 of I [3] is 2 when the atomic operation is executed.

### 16.9.29 ATOMIC\_REF (VALUE, ATOM [, STAT])

1 **Description.** Reference a variable atomically.

2 **Class.** [Atomic subroutine](#).

3 **Arguments.**

4 VALUE shall be scalar and of the same type as ATOM. It is an [INTENT \(OUT\)](#) argument. On successful  
5 execution, it becomes defined with the value of ATOM. If an error condition occurs, it becomes  
6 undefined.

7  
8 ATOM shall be a scalar [coarray](#) or [coindexed object](#). It shall be of type integer with kind [ATOMIC\\_INT\\_-](#)  
9 [KIND](#) from the intrinsic module [ISO\\_FORTRAN\\_ENV](#), or of type logical with kind [ATOMIC\\_LO-](#)  
10 [GICAL\\_KIND](#) from the intrinsic module [ISO\\_FORTRAN\\_ENV](#). It is an [INTENT \(IN\)](#) argument.

11 STAT (optional) shall be a noncoindexed integer scalar with a decimal range of at least four. It is an [INTENT](#)  
12 [\(OUT\)](#) argument. It is assigned a value as specified in [16.5](#). If an error condition occurs and STAT  
13 is not present, [error termination](#) is initiated.

14 **Example.** CALL ATOMIC\_REF (VAL, I [3]) causes VAL to become defined with the value of I on [image 3](#).

### 16.9.30 ATOMIC\_XOR (ATOM, VALUE [, STAT])

15 **Description.** Atomic bitwise exclusive OR.

16 **Class.** [Atomic subroutine](#).

17 **Arguments.**

18 ATOM shall be a scalar [coarray](#) or [coindexed object](#). It shall be of type integer with kind [ATOMIC\\_INT\\_-](#)  
19 [KIND](#) from the intrinsic module [ISO\\_FORTRAN\\_ENV](#). It is an [INTENT \(INOUT\)](#) argument. If  
20 an error condition occurs, ATOM becomes undefined; otherwise, it becomes defined with the value  
21 of [IEOR](#) (ATOM, [INT](#) (VALUE, [ATOMIC\\_INT\\_KIND](#))).

22  
23 VALUE shall be an integer scalar. It is an [INTENT \(IN\)](#) argument. The value of VALUE shall be repres-  
24 entable in kind [ATOMIC\\_INT\\_KIND](#).

25 STAT (optional) shall be a noncoindexed integer scalar with a decimal range of at least four. It is an [INTENT](#)  
26 [\(OUT\)](#) argument. It is assigned a value as specified in [16.5](#). If an error condition occurs and STAT  
27 is not present, [error termination](#) is initiated.

28 **Example.** CALL ATOMIC\_XOR (I [3], 1) will cause I on [image 3](#) to become defined with the value 2 if the  
29 value of I [3] is 3 when the atomic operation is executed.

### 16.9.31 BESSEL\_J0 (X)

30 **Description.** Bessel function of the 1<sup>st</sup> kind, order 0.

31 **Class.** [Elemental function](#).

32 **Argument.** X shall be of type real.

33 **Result Characteristics.** Same as X.

34 **Result Value.** The result has a value equal to a processor-dependent approximation to the Bessel function of  
35 the first kind and order zero of X.

36 **Example.** BESSEL\_J0 (1.0) has the value 0.765 (approximately).



### 16.9.32 BESSEL\_J1 (X)

1 **Description.** Bessel function of the 1<sup>st</sup> kind, order 1.

2 **Class.** [Elemental](#) function.

3 **Argument.** X shall be of type real.

4 **Result Characteristics.** Same as X.

5 **Result Value.** The result has a value equal to a processor-dependent approximation to the Bessel function of the first kind and order one of X.

6 **Example.** BESSEL\_J1 (1.0) has the value 0.440 (approximately).

### 16.9.33 BESSEL\_JN (N, X) or BESSEL\_JN (N1, N2, X)

1 **Description.** Bessel functions of the 1<sup>st</sup> kind.

2 **Class.**

3 *Case (i):* BESSEL\_JN (N,X) is an [elemental](#) function.

4 *Case (ii):* BESSEL\_JN (N1,N2,X) is a [transformational function](#).

5 **Arguments.**

6 N shall be of type integer and nonnegative.

7 N1 shall be an integer scalar with a nonnegative value.

8 N2 shall be an integer scalar with a nonnegative value.

9 X shall be of type real; if the function is transformational, X shall be scalar.

10 **Result Characteristics.** Same type and kind as X. The result of BESSEL\_JN (N1, N2, X) is a rank-one array with extent [MAX](#) (N2–N1+1, 0).

11 **Result Value.**

12 *Case (i):* The result value of BESSEL\_JN (N, X) is a processor-dependent approximation to the Bessel function of the first kind and order N of X.

13 *Case (ii):* Element *i* of the result value of BESSEL\_JN (N1, N2, X) is a processor-dependent approximation to the Bessel function of the first kind and order N1+i – 1 of X.

14 **Example.** BESSEL\_JN (2, 1.0) has the value 0.115 (approximately).

### 16.9.34 BESSEL\_Y0 (X)

1 **Description.** Bessel function of the 2<sup>nd</sup> kind, order 0.

2 **Class.** [Elemental](#) function.

3 **Argument.** X shall be of type real. Its value shall be greater than zero.

4 **Result Characteristics.** Same as X.

5 **Result Value.** The result has a value equal to a processor-dependent approximation to the Bessel function of the second kind and order zero of X.

6 **Example.** BESSEL\_Y0 (1.0) has the value 0.088 (approximately).

### 16.9.35 BESSEL\_Y1 (X)

1 **Description.** Bessel function of the 2<sup>nd</sup> kind, order 1.

2 **Class.** [Elemental](#) function.

3 **Argument.** X shall be of type real. Its value shall be greater than zero.

4 **Result Characteristics.** Same as X.

5 **Result Value.** The result has a value equal to a processor-dependent approximation to the Bessel function of the second kind and order one of X.

6 **Example.** BESSEL\_Y1 (1.0) has the value  $-0.781$  (approximately).

### 16.9.36 BESSEL\_YN (N, X) or BESSEL\_YN (N1, N2, X)

1 **Description.** Bessel functions of the 2<sup>nd</sup> kind.

2 **Class.**

3 *Case (i):* BESSEL\_YN (N, X) is an [elemental](#) function.

4 *Case (ii):* BESSEL\_YN (N1, N2, X) is a [transformational function](#).

5 **Arguments.**

6 N shall be of type integer and nonnegative.

7 N1 shall be an integer scalar with a nonnegative value.

8 N2 shall be an integer scalar with a nonnegative value.

9 X shall be of type real; if the function is transformational, X shall be scalar. Its value shall be greater than zero.

10 **Result Characteristics.** Same type and kind as X. The result of BESSEL\_YN (N1, N2, X) is a rank-one array with extent [MAX](#) (N2–N1+1, 0).

11 **Result Value.**

12 *Case (i):* The result value of BESSEL\_YN (N, X) is a processor-dependent approximation to the Bessel function of the second kind and order N of X.

13 *Case (ii):* Element *i* of the result value of BESSEL\_YN (N1, N2, X) is a processor-dependent approximation to the Bessel function of the second kind and order  $N1+i-1$  of X.

14 **Example.** BESSEL\_YN (2, 1.0) has the value  $-1.651$  (approximately).

### 16.9.37 BGE (I, J)

1 **Description.** Bitwise greater than or equal to.

2 **Class.** [Elemental](#) function.

3 **Arguments.**

4 I shall be of type integer or a [boz-literal-constant](#).

5 J shall be of type integer or a [boz-literal-constant](#).

6 **Result Characteristics.** Default logical.

7 **Result Value.** The result is true if the sequence of bits represented by I is greater than or equal to the sequence of bits represented by J, according to the method of bit sequence comparison in [16.3.2](#); otherwise the result is false.

1 6 The interpretation of a *boz-literal-constant* as a sequence of bits is described in 7.7. The interpretation of an  
2 integer value as a sequence of bits is described in 16.3.

3 7 **Example.** If BIT\_SIZE (J) has the value 8, BGE (Z'FF', J) has the value true for any value of J. BGE (0, -1)  
4 has the value false.

### 5 **16.9.38 BGT (I, J)**

6 1 **Description.** Bitwise greater than.

7 2 **Class.** [Elemental](#) function.

8 3 **Arguments.**

9 I shall be of type integer or a *boz-literal-constant*.

10 J shall be of type integer or a *boz-literal-constant*.

11 4 **Result Characteristics.** Default logical.

12 5 **Result Value.** The result is true if the sequence of bits represented by I is greater than the sequence of bits  
13 represented by J, according to the method of bit sequence comparison in 16.3.2; otherwise the result is false.

14 6 The interpretation of a *boz-literal-constant* as a sequence of bits is described in 7.7. The interpretation of an  
15 integer value as a sequence of bits is described in 16.3.

16 7 **Example.** BGT (Z'FF', Z'FC') has the value true. BGT (0, -1) has the value false.

### 17 **16.9.39 BIT\_SIZE (I)**

18 1 **Description.** Number of bits in integer model 16.3.

19 2 **Class.** [Inquiry](#) function.

20 3 **Argument.** I shall be of type integer. It may be a scalar or an array.

21 4 **Result Characteristics.** Scalar integer with the same kind type parameter as I.

22 5 **Result Value.** The result has the value of the number of bits *z* of the model integer defined for bit manipulation  
23 contexts in 16.3.

24 6 **Example.** BIT\_SIZE (1) has the value 32 if *z* of the model is 32.

### 25 **16.9.40 BLE (I, J)**

26 1 **Description.** Bitwise less than or equal to.

27 2 **Class.** [Elemental](#) function.

28 3 **Arguments.**

29 I shall be of type integer or a *boz-literal-constant*.

30 J shall be of type integer or a *boz-literal-constant*.

31 4 **Result Characteristics.** Default logical.

32 5 **Result Value.** The result is true if the sequence of bits represented by I is less than or equal to the sequence of  
33 bits represented by J, according to the method of bit sequence comparison in 16.3.2; otherwise the result is false.

34 6 The interpretation of a *boz-literal-constant* as a sequence of bits is described in 7.7. The interpretation of an  
35 integer value as a sequence of bits is described in 16.3.

1 7 **Example.** BLE (0, J) has the value true for any value of J. BLE (−1, 0) has the value false.

## 2 16.9.41 BLT (I, J)

3 1 **Description.** Bitwise less than.

4 2 **Class.** [Elemental](#) function.

5 3 **Arguments.**

6 I shall be of type integer or a *boz-literal-constant*.

7 J shall be of type integer or a *boz-literal-constant*.

8 4 **Result Characteristics.** Default logical.

9 5 **Result Value.** The result is true if the sequence of bits represented by I is less than the sequence of bits  
10 represented by J, according to the method of bit sequence comparison in [16.3.2](#); otherwise the result is false.

11 6 The interpretation of a *boz-literal-constant* as a sequence of bits is described in [7.7](#). The interpretation of an  
12 integer value as a sequence of bits is described in [16.3](#).

13 7 **Example.** BLT (0, −1) has the value true. BLT (Z'FF', Z'FC') has the value false.

## 14 16.9.42 BTEST (I, POS)

15 1 **Description.** Test single bit in an integer.

16 2 **Class.** [Elemental](#) function.

17 3 **Arguments.**

18 I shall be of type integer.

19 POS shall be of type integer. It shall be nonnegative and be less than BIT\_SIZE (I).

20 4 **Result Characteristics.** Default logical.

21 5 **Result Value.** The result has the value true if bit POS of I has the value 1 and has the value false if bit POS  
22 of I has the value 0. The model for the interpretation of an integer value as a sequence of bits is in [16.3](#).

23 6 **Examples.** BTEST (8, 3) has the value true. If A has the value  $\begin{bmatrix} 1 & 2 \\ 3 & 4 \end{bmatrix}$ , the value of BTEST (A, 2) is  
24  $\begin{bmatrix} \text{false} & \text{false} \\ \text{false} & \text{true} \end{bmatrix}$  and the value of BTEST (2, A) is  $\begin{bmatrix} \text{true} & \text{false} \\ \text{false} & \text{false} \end{bmatrix}$ .

## 25 16.9.43 CEILING (A [, KIND])

26 1 **Description.** Least integer greater than or equal to A.

27 2 **Class.** [Elemental](#) function.

28 3 **Arguments.**

29 A shall be of type real.

30 KIND (optional) shall be a scalar integer [constant expression](#).

31 4 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of  
32 KIND; otherwise, the kind type parameter is that of default integer type.

33 5 **Result Value.** The result has a value equal to the least integer greater than or equal to A.

1 6 **Examples.** CEILING (3.7) has the value 4. CEILING (−3.7) has the value −3.

## 2 16.9.44 CHAR (I [, KIND])

3 1 **Description.** Character from code value.

4 2 **Class.** [Elemental](#) function.

5 3 **Arguments.**

6 I shall be of type integer with a value in the range  $0 \leq I \leq n - 1$ , where  $n$  is the number of characters  
7 in the [collating sequence](#) associated with the specified kind type parameter.

8 KIND (optional) shall be a scalar integer [constant expression](#).

9 4 **Result Characteristics.** Character of length one. If KIND is present, the kind type parameter is that specified  
10 by the value of KIND; otherwise, the kind type parameter is that of default character.

11 5 **Result Value.** The result is the character in position I of the [collating sequence](#) associated with the spe-  
12 cified kind type parameter. ICHAR (CHAR (I, KIND (C))) shall have the value I for  $0 \leq I \leq n - 1$  and  
13 CHAR (ICHAR (C), KIND (C)) shall have the value C for any character C capable of representation in the  
14 processor.

15 6 **Example.** CHAR (88) has the value 'X' on a processor using the [ASCII collating sequence](#) for default characters.

## 16 16.9.45 CMPLX (X [, KIND]) or CMPLX (X [, Y, KIND])

17 1 **Description.** Conversion to complex type.

18 2 **Class.** [Elemental](#) function.

19 3 **Arguments for CMPLX(X [, KIND]).**

20 X shall be of type complex.

21 KIND (optional) shall be a scalar integer [constant expression](#).

22 4 **Arguments for CMPLX(X [, Y, KIND]).**

23 X shall be of type integer or real, or a *boz-literal-constant*.

24 Y (optional) shall be of type integer or real, or a *boz-literal-constant*.

25 KIND (optional) shall be a scalar integer [constant expression](#).

26 5 **Result Characteristics.** The result is of type complex. If KIND is present, the kind type parameter is that  
27 specified by the value of KIND; otherwise, the kind type parameter is that of default real kind.

28 6 **Result Value.** If Y is absent and X is not complex, it is as if Y were present with the value zero. If KIND is  
29 absent, it is as if KIND were present with the value [KIND](#) (0.0). If X is complex, the result is the same as that  
30 of CMPLX ([REAL](#) (X), [AIMAG](#) (X), KIND). The result of CMPLX (X, Y, KIND) has the complex value whose  
31 real part is [REAL](#) (X, KIND) and whose imaginary part is [REAL](#) (Y, KIND).

32 7 **Example.** CMPLX (−3) has the value (−3.0, 0.0).

## 33 16.9.46 CO\_BROADCAST(A, SOURCE\_IMAGE [, STAT, ERRMSG])

34 1 **Description.** Broadcast value to images.

35 2 **Class.** [Collective subroutine](#).

36 3 **Arguments.**

37 A shall have the same shape, dynamic type, and type parameter values, in corresponding references.  
38 It shall not be a [coindexed object](#). It is an [INTENT \(INOUT\)](#) argument. If no error condition

1 occurs, A becomes defined, as if by intrinsic assignment, on all images in the **current team** with the  
2 value of A on image SOURCE\_IMAGE.

3 SOURCE\_IMAGE shall be an integer scalar. It is an **INTENT (IN)** argument. Its value shall be that of an  
4 image index of an image in the current team. The value shall be the same in all corresponding  
5 references.

6 STAT (optional) shall be a noncoindexed integer scalar with a decimal exponent range of at least four. It is an  
7 **INTENT (OUT)** argument.

8 ERRMSG (optional) shall be a noncoindexed default character scalar. It is an **INTENT (INOUT)** argument.

9 4 The semantics of STAT and ERRMSG are described in 16.6.

10 5 **Example.** If A is the array [1, 5, 3] on image one, after execution of

11 CALL CO\_BROADCAST (A, 1)

12 the value of A on all images of the current team is [1, 5, 3].

### 13 16.9.47 CO\_MAX(A [, RESULT\_IMAGE, STAT, ERRMSG])

14 1 **Description.** Compute maximum value across images.

15 2 **Class.** Collective subroutine.

16 3 **Arguments.**

17 A shall be of type integer, real, or character. It shall have the same shape, type, and type parameter  
18 values, in corresponding references. It shall not be a **coindexed object**. It is an **INTENT (INOUT)**  
19 argument. If it is scalar, the computed value is equal to the maximum value of A in all corresponding  
20 references. If it is an array each element of the computed value is equal to the maximum value of  
21 all corresponding elements of A in all corresponding references.

22 The computed value is assigned to A if no error condition occurs, and either RESULT\_IMAGE is  
23 absent, or the executing image is the one identified by RESULT\_IMAGE. Otherwise, A becomes  
24 undefined.

25 RESULT\_IMAGE (optional) shall be an integer scalar. It is an **INTENT (IN)** argument. Its presence, and value  
26 if present, shall be the same in all corresponding references. If it is present, its value shall be that  
27 of an image index in the current team.

28 STAT (optional) shall be a noncoindexed integer scalar with a decimal exponent range of at least four. It is an  
29 **INTENT (OUT)** argument.

30 ERRMSG (optional) shall be a noncoindexed default character scalar. It is an **INTENT (INOUT)** argument.

31 4 The semantics of STAT and ERRMSG are described in 16.6.

32 5 **Example.** If the number of images in the current team is two and A is the array [1, 5, 3] on one image and [4,  
33 1, 6] on the other image, the value of A after executing the statement CALL CO\_MAX (A) is [4, 5, 6] on both  
34 images.

### 35 16.9.48 CO\_MIN(A [, RESULT\_IMAGE, STAT, ERRMSG])

36 1 **Description.** Compute minimum value across images.

37 2 **Class.** Collective subroutine.

38 3 **Arguments.**

39 A shall be of type integer, real, or character. It shall have the same shape, type, and type parameter  
40 values, in corresponding references. It shall not be a **coindexed object**. It is an **INTENT (INOUT)**  
41 argument. If it is scalar, the computed value is equal to the minimum value of A in all corresponding

1 references. If it is an array each element of the computed value is equal to the minimum value of  
2 all corresponding elements of A in all corresponding references.

3 The computed value is assigned to A if no error condition occurs, and either RESULT\_IMAGE is  
4 absent, or the executing image is the one identified by RESULT\_IMAGE. Otherwise, A becomes  
5 undefined.

6 RESULT\_IMAGE (optional) shall be an integer scalar. It is an **INTENT (IN)** argument. Its presence, and value  
7 if present, shall be the same in all corresponding references. If it is present, its value shall be that  
8 of an image index in the current team.

9 STAT (optional) shall be a noncoindexed integer scalar with a decimal exponent range of at least four. It is an  
10 **INTENT (OUT)** argument.

11 ERRMSG (optional) shall be a noncoindexed default character scalar. It is an **INTENT (INOUT)** argument.

12 4 The semantics of STAT and ERRMSG are described in 16.6.

13 5 **Example.** If the number of images in the current team is two and A is the array [1, 5, 3] on one image and [4,  
14 1, 6] on the other image, the value of A after executing the statement CALL CO\_MIN (A) is [1, 1, 3] on both  
15 images.

### 16 16.9.49 CO\_REDUCE(A, OPERATION [, RESULT\_IMAGE, STAT, ERRMSG])

17 1 **Description.** Generalized reduction across images.

18 2 **Class.** Collective subroutine.

19 3 **Arguments.**

20 A shall not be polymorphic. It shall have the same shape, type, and type parameter values, in  
21 corresponding references. It shall not be a **coindexed object**. It is an **INTENT (INOUT)** argument.  
22 If A is scalar, the computed value is the result of the reduction operation of applying OPERATION  
23 to the values of A in all corresponding references. If A is an array, each element of the computed  
24 value is equal to the result of the reduction operation of applying OPERATION to corresponding  
25 elements of A in all corresponding references.

26 The computed value is assigned to A if no error condition occurs, and either RESULT\_IMAGE is  
27 absent, or the executing image is the one identified by RESULT\_IMAGE. Otherwise, A becomes  
28 undefined.

29 OPERATION shall be a pure function with exactly two arguments; the result and each argument shall be  
30 a scalar, nonallocatable, nonpointer, nonpolymorphic data object with the same type and type  
31 parameters as A. The arguments shall not be optional. If one argument has the **ASYNCHRONOUS**,  
32 **TARGET**, or **VALUE** attribute, the other shall have that attribute. OPERATION shall implement  
33 a mathematically associative operation. OPERATION shall be the same function on all images in  
34 corresponding references.

35 The computed value of a reduction operation over a set of values is the result of an iterative process.  
36 Each iteration involves the evaluation of OPERATION ( $x, y$ ) for  $x$  and  $y$  in the set, the removal of  
37  $x$  and  $y$  from the set, and the addition of the value of OPERATION ( $x, y$ ) to the set. The process  
38 terminates when the set has only one element; this is the computed value.

39 RESULT\_IMAGE (optional) shall be an integer scalar. It is an **INTENT (IN)** argument. Its presence, and value  
40 if present, shall be the same in all corresponding references. If it is present, its value shall be that  
41 of an image index in the current team.

42 STAT (optional) shall be a noncoindexed integer scalar with a decimal exponent range of at least four. It is an  
43 **INTENT (OUT)** argument.

44 ERRMSG (optional) shall be a noncoindexed default character scalar. It is an **INTENT (INOUT)** argument.

45 4 The semantics of STAT and ERRMSG are described in 16.6.

46 5 **Example.** The subroutine below demonstrates how to use CO\_REDUCE to create a collective counterpart to

```

1 the intrinsic function ALL:
2     SUBROUTINE co_all(boolea)
3         LOGICAL, INTENT(INOUT) :: boolea
4         CALL CO_REDUCE(boolea,both)
5     CONTAINS
6         PURE FUNCTION both(lhs,rhs) RESULT(lhs_and_rhs)
7             LOGICAL, INTENT(IN) :: lhs,rhs
8             LOGICAL :: lhs_and_rhs
9             lhs_and_rhs = lhs .AND. rhs
10        END FUNCTION both
11    END SUBROUTINE co_all

```

**NOTE 16.11**

If the OPERATION function is not mathematically commutative, the result of calling CO\_REDUCE can depend on the order of evaluations.

**12 16.9.50 CO\_SUM(A [, RESULT\_IMAGE, STAT, ERRMSG])**

13 1 **Description.** Compute sum across images.

14 2 **Class.** [Collective subroutine.](#)

15 3 **Arguments.**

16 A shall be of numeric type. It shall have the same shape, type, and type parameter values, in corresponding references. It shall not be a [coindexed object](#). It is an [INTENT \(INOUT\)](#) argument. If it is scalar, the computed value is equal to a processor-dependent approximation to the sum of the values of A in corresponding references. If it is an array, each element of the computed value is equal to a processor-dependent approximation to the sum of all corresponding elements of A in corresponding references.

22 The computed value is assigned to A if no error condition occurs, and either RESULT\_IMAGE is absent, or the executing image is the one identified by RESULT\_IMAGE. Otherwise, A becomes undefined.

25 RESULT\_IMAGE (optional) shall be an integer scalar. It is an [INTENT \(IN\)](#) argument. Its presence, and value if present, shall be the same in all corresponding references. If it is present, its value shall be that of an image index in the current team.

28 STAT (optional) shall be a noncoindexed integer scalar with a decimal exponent range of at least four. It is an [INTENT \(OUT\)](#) argument.

30 ERRMSG (optional) shall be a noncoindexed default character scalar. It is an [INTENT \(INOUT\)](#) argument.

31 4 The semantics of STAT and ERRMSG are described in [16.6](#).

32 5 **Example.** If the number of images in the current team is two and A is the array [1, 5, 3] on one image and [4, 1, 6] on the other image, the value of A after executing the statement CALL CO\_SUM(A) is [5, 6, 9] on both images.

**35 16.9.51 COMMAND\_ARGUMENT\_COUNT ( )**

36 1 **Description.** Number of command arguments.

37 2 **Class.** [Transformational function.](#)

38 3 **Argument.** None.



1 4 **Result Characteristics.** Default integer scalar.

2 5 **Result Value.** The result value is equal to the number of command arguments available. If there are no  
3 command arguments available or if the processor does not support command arguments, then the result has the  
4 value zero. If the processor has a concept of a command name, the command name does not count as one of the  
5 command arguments.

6 6 **Example.** See 16.9.83.

## 7 16.9.52 CONJG (Z)

8 1 **Description.** Conjugate of a complex number.

9 2 **Class.** [Elemental](#) function.

10 3 **Argument.** Z shall be of type complex.

11 4 **Result Characteristics.** Same as Z.

12 5 **Result Value.** If Z has the value  $(x, y)$ , the result has the value  $(x, -y)$ .

13 6 **Example.** CONJG ((2.0, 3.0)) has the value (2.0, -3.0).

## 14 16.9.53 COS (X)

15 1 **Description.** Cosine function.

16 2 **Class.** [Elemental](#) function.

17 3 **Argument.** X shall be of type real or complex.

18 4 **Result Characteristics.** Same as X.

19 5 **Result Value.** The result has a value equal to a processor-dependent approximation to  $\cos(X)$ . If X is of type  
20 real, it is regarded as a value in radians. If X is of type complex, its real part is regarded as a value in radians.

21 6 **Example.** COS (1.0) has the value 0.54030231 (approximately).

## 22 16.9.54 COSH (X)

23 1 **Description.** Hyperbolic cosine function.

24 2 **Class.** [Elemental](#) function.

25 3 **Argument.** X shall be of type real or complex.

26 4 **Result Characteristics.** Same as X.

27 5 **Result Value.** The result has a value equal to a processor-dependent approximation to  $\cosh(X)$ . If X is of type  
28 complex its imaginary part is regarded as a value in radians.

29 6 **Example.** COSH (1.0) has the value 1.5430806 (approximately).

## 30 16.9.55 COSHAPE (COARRAY [, KIND])

31 1 **Description.** Sizes of [codimensions](#) of a [coarray](#).

32 2 **Class.** [Inquiry function](#).

33 3 **Arguments.**

34 COARRAY shall be a [coarray](#) of any type. It shall not be an unallocated [allocatable coarray](#).

1 KIND (optional) shall be a scalar integer [constant expression](#).

2 4 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value  
3 of KIND; otherwise the kind type parameter is that of default integer type. The result is an array of [rank](#) one  
4 whose size is equal to the [corank](#) of COARRAY.

5 5 **Result Value.** The result has a value whose  $i^{th}$  element is equal to the size of the  $i^{th}$  [codimension](#) of COARRAY,  
6 as given by [UCOBOUND](#) (COARRAY,  $i$ ) – [LCOBOUND](#) (COARRAY,  $i$ ) +1.

7 6 **Example.**

8 The following code allocates the coarray D with the same size in each codimension as that of the coarray C, with  
9 the lower cobound 1.

```
10 REAL, ALLOCATABLE :: C[:,:], D[:,:]
11 INTEGER, ALLOCATABLE :: COSHAPE_C(:)
12 ...
13 COSHAPE_C = COSHAPE(C)
14 ALLOCATE ( D[COSHAPE_C(1),*] )
```

## 15 16.9.56 COUNT (MASK [, DIM, KIND])

16 1 **Description.** Logical array reduced by counting true values.

17 2 **Class.** [Transformational function](#).

18 3 **Arguments.**

19 MASK shall be a logical array.

20 DIM (optional) shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the [rank](#) of MASK.  
21 The corresponding [actual argument](#) shall not be an optional dummy argument, a disassociated  
22 pointer, or an unallocated allocatable.

23 KIND (optional) shall be a scalar integer [constant expression](#).

24 4 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of  
25 KIND; otherwise the kind type parameter is that of default integer type. The result is scalar if DIM is absent or  
26  $n = 1$ ; otherwise, the result has [rank](#)  $n - 1$  and shape  $[d_1, d_2, \dots, d_{\text{DIM}-1}, d_{\text{DIM}+1}, \dots, d_n]$  where  $[d_1, d_2, \dots, d_n]$   
27 is the shape of MASK.

28 5 **Result Value.**

29 *Case (i):* If DIM is absent or MASK has rank one, the result has a value equal to the number of true elements  
30 of MASK or has the value zero if MASK has size zero.

31 *Case (ii):* If DIM is present and MASK has rank  $n > 1$ , the value of element  $(s_1, s_2, \dots, s_{\text{DIM}-1}, s_{\text{DIM}+1}, \dots,$   
32  $s_n)$  of the result is equal to the number of true elements of MASK  $(s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1},$   
33  $\dots, s_n)$ .

34 6 **Examples.**

35 *Case (i):* The value of COUNT (.TRUE., .FALSE., .TRUE.) is 2.

36 *Case (ii):* If B is the array  $\begin{bmatrix} 1 & 3 & 5 \\ 2 & 4 & 6 \end{bmatrix}$  and C is the array  $\begin{bmatrix} 0 & 3 & 5 \\ 7 & 4 & 8 \end{bmatrix}$ , COUNT (B /= C, DIM = 1) is  
37 [2, 0, 1] and COUNT (B /= C, DIM = 2) is [1, 2].

## 16.9.57 CPU\_TIME (TIME)

1 **Description.** Processor time used.

2 **Class.** Subroutine.

3 **Argument.** TIME shall be a real scalar. It is an **INTENT (OUT)** argument. If the processor cannot provide a meaningful value for the time, it is assigned a processor-dependent negative value; otherwise, it is assigned a processor-dependent approximation to the processor time in seconds. Whether the value assigned is an approximation to the amount of time used by the invoking **image**, or the amount of time used by the whole program, is processor dependent.

4 **Example.**

```

10     REAL T1, T2
11     ...
12     CALL CPU_TIME(T1)
13     ... ! Code to be timed.
14     CALL CPU_TIME(T2)
15     WRITE (*,*) 'Time taken by code was ', T2-T1, ' seconds'
```

16 writes the processor time taken by a piece of code.

### NOTE 16.12

A processor for which a single result is inadequate (for example, a parallel processor) might choose to provide an additional version for which time is an array.

The exact definition of time is left imprecise because of the variability in what different processors are able to provide. The primary purpose is to compare different algorithms on the same processor or discover which parts of a calculation are the most expensive.

The start time is left imprecise because the purpose is to time sections of code, as in the example.

Most computer systems have multiple concepts of time. One common concept is that of time expended by the processor for a given program. This might or might not include system overhead, and has no obvious connection to elapsed “wall clock” time.

## 16.9.58 CSHIFT (ARRAY, SHIFT [, DIM])

1 **Description.** Circular shift of an array.

2 **Class.** Transformational function.

3 **Arguments.**

4 **ARRAY** may be of any type. It shall be an array.

5 **SHIFT** shall be of type integer and shall be scalar if **ARRAY** has **rank** one; otherwise, it shall be scalar or of **rank**  $n - 1$  and of shape  $[d_1, d_2, \dots, d_{DIM-1}, d_{DIM+1}, \dots, d_n]$  where  $[d_1, d_2, \dots, d_n]$  is the shape of **ARRAY**.

6 **DIM** (optional) shall be an integer scalar with a value in the range  $1 \leq DIM \leq n$ , where  $n$  is the **rank** of **ARRAY**. If **DIM** is absent, it is as if it were present with the value 1.

7 **Result Characteristics.** The result is of the type and type parameters of **ARRAY**, and has the shape of **ARRAY**.

8 **Result Value.**

9 *Case (i):* If **ARRAY** has **rank** one, element  $i$  of the result is  $ARRAY(1 + \text{MODULO}(i + \text{SHIFT} - 1, \text{SIZE}(\text{ARRAY})))$ .

1 *Case (ii):* If ARRAY has **rank** greater than one, section  $(s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n)$  of the result  
 2 has a value equal to CSHIFT (ARRAY  $(s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n)$ , *sh*, 1), where *sh* is  
 3 SHIFT or SHIFT  $(s_1, s_2, \dots, s_{\text{DIM}-1}, s_{\text{DIM}+1}, \dots, s_n)$ .

#### 4 6 Examples.

5 *Case (i):* If V is the array [1, 2, 3, 4, 5, 6], the effect of shifting V circularly to the left by two positions is  
 6 achieved by CSHIFT (V, SHIFT = 2) which has the value [3, 4, 5, 6, 1, 2]; CSHIFT (V, SHIFT =  
 7 -2) achieves a circular shift to the right by two positions and has the value [5, 6, 1, 2, 3, 4].

8 *Case (ii):* The rows of an array of **rank** two may all be shifted by the same amount or by different amounts.

9 If M is the array  $\begin{bmatrix} 1 & 2 & 3 \\ 4 & 5 & 6 \\ 7 & 8 & 9 \end{bmatrix}$ , the value of

10 CSHIFT (M, SHIFT = -1, DIM = 2) is  $\begin{bmatrix} 3 & 1 & 2 \\ 6 & 4 & 5 \\ 9 & 7 & 8 \end{bmatrix}$ , and the value of

11 CSHIFT (M, SHIFT = [-1, 1, 0], DIM = 2) is  $\begin{bmatrix} 3 & 1 & 2 \\ 5 & 6 & 4 \\ 7 & 8 & 9 \end{bmatrix}$ .

### 12 16.9.59 DATE\_AND\_TIME ([DATE, TIME, ZONE, VALUES])

13 1 **Description.** Date and time.

14 2 **Class.** Subroutine.

15 3 **Arguments.**

16 DATE (optional) shall be a default character scalar. It is an **INTENT (OUT)** argument. It is assigned a value  
 17 of the form *YYYYMMDD*, where *YYYY* is the year in the Gregorian calendar, *MM* is the month  
 18 within the year, and *DD* is the day within the month. The characters of this value shall all be  
 19 decimal digits. If there is no date available, DATE is assigned all blanks.

20 TIME (optional) shall be a default character scalar. It is an **INTENT (OUT)** argument. It is assigned a value  
 21 of the form *hhmmss.sss*, where *hh* is the hour of the day, *mm* is the minutes of the hour, and *ss.sss*  
 22 is the seconds and milliseconds of the minute. Except for the decimal point, the characters of this  
 23 value shall all be decimal digits. If there is no clock available, TIME is assigned all blanks.

24 ZONE (optional) shall be a default character scalar. It is an **INTENT (OUT)** argument. It is assigned a value of  
 25 the form *+hhmm* or *-hhmm*, where *hh* and *mm* are the time difference with respect to Coordinated  
 26 Universal Time (UTC) in hours and minutes, respectively. The characters of this value following  
 27 the sign character shall all be decimal digits. If this information is not available, ZONE is assigned  
 28 all blanks.

29 VALUES (optional) shall be a rank-one array of type integer with a decimal exponent range of at least four. It  
 30 is an **INTENT (OUT)** argument. Its size shall be at least 8. The values assigned to VALUES are  
 31 as follows:

32 VALUES (1) the year, including the century (for example, 2008), or -HUGE (VALUES) if there is no date  
 33 available;

34 VALUES (2) the month of the year, or -HUGE (VALUES) if there is no date available;

35 VALUES (3) the day of the month, or -HUGE (VALUES) if there is no date available;

36 VALUES (4) the time difference with respect to Coordinated Universal Time (UTC) in minutes, or -HUGE (VAL-  
 37 UES) if this information is not available;

38 VALUES (5) the hour of the day, in the range of 0 to 23, or -HUGE (VALUES) if there is no clock;

39 VALUES (6) the minutes of the hour, in the range 0 to 59, or -HUGE (VALUES) if there is no clock;

40 VALUES (7) the seconds of the minute, in the range 0 to 60, or -HUGE (VALUES) if there is no clock;

41 VALUES (8) the milliseconds of the second, in the range 0 to 999, or -HUGE (VALUES) if there is no clock.

1 4 The date, clock, and time zone information might be available on some [images](#) and not others. If the date, clock,  
2 or time zone information is available on more than one [image](#), it is processor dependent whether or not those  
3 [images](#) share the same information.

4 5 **Example.** If run in Geneva, Switzerland on April 12, 2008 at 15:27:35.5 with a system configured for the  
5 local time zone, this example would have assigned the value 20080412 to BIG\_BEN (1), the value 152735.500 to  
6 BIG\_BEN (2), the value +0100 to BIG\_BEN (3), and the value [2008, 4, 12, 60, 15, 27, 35, 500] to DATE\_TIME.

```
7         INTEGER DATE_TIME (8)
8         CHARACTER (LEN = 10) BIG_BEN (3)
9         CALL DATE_AND_TIME (BIG_BEN (1), BIG_BEN (2), BIG_BEN (3), DATE_TIME)
```

#### NOTE 16.13

These forms are compatible with the representations defined in ISO 8601:2004. UTC is established by the International Bureau of Weights and Measures (BIPM, i.e. Bureau International des Poids et Mesures) and the International Earth Rotation Service (IERS).

### 10 16.9.60 DBLE (A)

11 1 **Description.** Conversion to double precision real.

12 2 **Class.** [Elemental](#) function.

13 3 **Argument.** A shall be of type integer, real, complex, or a [boz-literal-constant](#).

14 4 **Result Characteristics.** Double precision real.

15 5 **Result Value.** The result has the value REAL (A, KIND (0.0D0)).

16 6 **Example.** DBLE (−3) has the value −3.0D0.

### 17 16.9.61 DIGITS (X)

18 1 **Description.** Significant digits in numeric model.

19 2 **Class.** [Inquiry function](#).

20 3 **Argument.** X shall be of type integer or real. It may be a scalar or an array.

21 4 **Result Characteristics.** Default integer scalar.

22 5 **Result Value.** The result has the value  $q$  if X is of type integer and  $p$  if X is of type real, where  $q$  and  $p$  are as  
23 defined in [16.4](#) for the model representing numbers of the same type and kind type parameter as X.

24 6 **Example.** DIGITS (X) has the value 24 for real X whose model is as in Note [16.4](#).

### 25 16.9.62 DIM (X, Y)

26 1 **Description.** Maximum of X – Y and zero.

27 2 **Class.** [Elemental](#) function.

28 3 **Arguments.**

29 X shall be of type integer or real.

30 Y shall be of the same type and kind type parameter as X.

31 4 **Result Characteristics.** Same as X.

1 5 **Result Value.** The value of the result is the maximum of  $X - Y$  and zero.

2 6 **Example.** DIM (-3.0, 2.0) has the value 0.0.

### 3 16.9.63 DOT\_PRODUCT (VECTOR\_A, VECTOR\_B)

4 1 **Description.** Dot product of two vectors.

5 2 **Class.** Transformational function.

6 3 **Arguments.**

7 VECTOR\_A shall be of **numeric type** (integer, real, or complex) or of logical type. It shall be a rank-one array.

8 VECTOR\_B shall be of **numeric type** if VECTOR\_A is of **numeric type** or of type logical if VECTOR\_A is of  
9 type logical. It shall be a rank-one array. It shall be of the same size as VECTOR\_A.

10 4 **Result Characteristics.** If the arguments are of **numeric type**, the type and **kind type parameter** of the result  
11 are those of the expression VECTOR\_A \* VECTOR\_B determined by the types and kinds of the arguments  
12 according to 10.1.9.3. If the arguments are of type logical, the result is of type logical with the kind type parameter  
13 of the expression VECTOR\_A .AND. VECTOR\_B according to 10.1.9.3. The result is scalar.

14 5 **Result Value.**

15 *Case (i):* If VECTOR\_A is of type integer or real, the result has the value SUM (VECTOR\_A\*VECTOR\_-  
16 B). If the vectors have size zero, the result has the value zero.

17 *Case (ii):* If VECTOR\_A is of type complex, the result has the value SUM (CONJG (VECTOR\_A)\*VECTOR\_-  
18 B). If the vectors have size zero, the result has the value zero.

19 *Case (iii):* If VECTOR\_A is of type logical, the result has the value ANY (VECTOR\_A .AND. VECTOR\_B).  
20 If the vectors have size zero, the result has the value false.

21 6 **Example.** DOT\_PRODUCT ([1, 2, 3], [2, 3, 4]) has the value 20.

### 22 16.9.64 DPROD (X, Y)

23 1 **Description.** Double precision real product.

24 2 **Class.** Elemental function.

25 3 **Arguments.**

26 X shall be default real.

27 Y shall be default real.

28 4 **Result Characteristics.** Double precision real.

29 5 **Result Value.** The result has a value equal to a processor-dependent approximation to the product of X and  
30 Y. DPROD (X, Y) should have the same value as DBLE (X) \* DBLE (Y).

31 6 **Example.** DPROD (-3.0, 2.0) has the value -6.0D0.

### 32 16.9.65 DSHIFTL (I, J, SHIFT)

33 1 **Description.** Combined left shift.

34 2 **Class.** Elemental function.

35 3 **Arguments.**

36 I shall be of type integer or a *boz-literal-constant*.

37 J shall be of type integer or a *boz-literal-constant*. If both I and J are of type integer, they shall have  
38 the same kind type parameter. I and J shall not both be *boz-literal-constants*.

1 SHIFT shall be of type integer. It shall be nonnegative and less than or equal to BIT\_SIZE (I) if I is of  
2 type integer; otherwise, it shall be less than or equal to BIT\_SIZE (J).

3 4 **Result Characteristics.** Same as I if I is of type integer; otherwise, same as J.

4 5 **Result Value.** If either I or J is a *boz-literal-constant*, it is first converted as if by the intrinsic function INT to  
5 type integer with the kind type parameter of the other. The rightmost SHIFT bits of the result value are the same  
6 as the leftmost bits of J, and the remaining bits of the result value are the same as the rightmost bits of I. This  
7 is equal to IOR (SHIFTL (I, SHIFT), SHIFTR (J, BIT\_SIZE (J)–SHIFT)). The model for the interpretation of  
8 an integer value as a sequence of bits is in 16.3.

9 6 **Examples.** DSHIFTL (1, 2\*\*30, 2) has the value 5 if default integer has 32 bits. DSHIFTL (I, I, SHIFT) has  
10 the same result value as ISHFTC (I, SHIFT).

### 11 16.9.66 DSHIFTR (I, J, SHIFT)

12 1 **Description.** Combined right shift.

13 2 **Class.** Elemental function.

14 3 **Arguments.**

15 I shall be of type integer or a *boz-literal-constant*.

16 J shall be of type integer or a *boz-literal-constant*. If both I and J are of type integer, they shall have  
17 the same kind type parameter. I and J shall not both be *boz-literal-constants*.

18 SHIFT shall be of type integer. It shall be nonnegative and less than or equal to BIT\_SIZE (I) if I is of  
19 type integer; otherwise, it shall be less than or equal to BIT\_SIZE (J).

20 4 **Result Characteristics.** Same as I if I is of type integer; otherwise, same as J.

21 5 **Result Value.** If either I or J is a *boz-literal-constant*, it is first converted as if by the intrinsic function INT to  
22 type integer with the kind type parameter of the other. The leftmost SHIFT bits of the result value are the same  
23 as the rightmost bits of I, and the remaining bits of the result value are the same as the leftmost bits of J. This  
24 is equal to IOR (SHIFTL (I, BIT\_SIZE (I)–SHIFT), SHIFTR (J, SHIFT)). The model for the interpretation of  
25 an integer value as a sequence of bits is in 16.3.

26 6 **Examples.** DSHIFTR (1, 16, 3) has the value  $2^{29} + 2$  if default integer has 32 bits. DSHIFTR (I, I, SHIFT) has  
27 the same result value as ISHFTC (I,–SHIFT).

### 28 16.9.67 EOSHIFT (ARRAY, SHIFT [, BOUNDARY, DIM])

29 1 **Description.** End-off shift of the elements of an array.

30 2 **Class.** Transformational function.

31 3 **Arguments.**

32 ARRAY shall be an array be of any type.

33 SHIFT shall be of type integer and shall be scalar if ARRAY has rank one; otherwise, it shall be scalar or  
34 of rank  $n - 1$  and of shape  $[d_1, d_2, \dots, d_{\text{DIM}-1}, d_{\text{DIM}+1}, \dots, d_n]$  where  $[d_1, d_2, \dots, d_n]$  is the shape  
35 of ARRAY.

36 BOUNDARY (optional) shall be of the same type and type parameters as ARRAY and shall be scalar if ARRAY  
37 has rank one; otherwise, it shall be either scalar or of rank  $n - 1$  and of shape  $[d_1, d_2, \dots, d_{\text{DIM}-1},$   
38  $d_{\text{DIM}+1}, \dots, d_n]$ . BOUNDARY is permitted to be absent only for the types in Table 16.4, and in  
39 this case it is as if it were present with the scalar value shown, converted if necessary to the kind  
40 type parameter value of ARRAY.

Table 16.4: Default BOUNDARY values for EOSHIFT

Type of ARRAY	Value of BOUNDARY
Integer	0
Real	0.0
Complex	(0.0, 0.0)
Logical	.FALSE.
Character ( <i>len</i> )	<i>len</i> blanks

1 DIM (optional) shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the **rank** of ARRAY.  
 2 If DIM is absent, it is as if it were present with the value 1.

3 4 **Result Characteristics.** The result has the type, type parameters, and shape of ARRAY.

4 5 **Result Value.** Element  $(s_1, s_2, \dots, s_n)$  of the result has the value ARRAY  $(s_1, s_2, \dots, s_{\text{DIM}-1}, s_{\text{DIM}} + sh,$   
 5  $s_{\text{DIM}+1}, \dots, s_n)$  where  $sh$  is SHIFT or SHIFT  $(s_1, s_2, \dots, s_{\text{DIM}-1}, s_{\text{DIM}+1}, \dots, s_n)$  provided the inequality  
 6 LBOUND (ARRAY, DIM)  $\leq s_{\text{DIM}} + sh \leq$  UBOUND (ARRAY, DIM) holds and is otherwise BOUNDARY or  
 7 BOUNDARY  $(s_1, s_2, \dots, s_{\text{DIM}-1}, s_{\text{DIM}+1}, \dots, s_n)$ .

8 6 **Examples.**

9 *Case (i):* If V is the array [1, 2, 3, 4, 5, 6], the effect of shifting V end-off to the left by 3 positions is achieved  
 10 by EOSHIFT (V, SHIFT = 3), which has the value [4, 5, 6, 0, 0, 0]; EOSHIFT (V, SHIFT = -2,  
 11 BOUNDARY = 99) achieves an end-off shift to the right by 2 positions with the boundary value of  
 12 99 and has the value [99, 99, 1, 2, 3, 4].

13 *Case (ii):* The rows of an array of **rank** two may all be shifted by the same amount or by different amounts  
 14 and the boundary elements can be the same or different. If M is the array  $\begin{bmatrix} A & B & C \\ D & E & F \\ G & H & I \end{bmatrix}$ , then the  
 15 value of EOSHIFT (M, SHIFT = -1, BOUNDARY = '\*', DIM = 2) is  $\begin{bmatrix} * & A & B \\ * & D & E \\ * & G & H \end{bmatrix}$ , and the value  
 16 of EOSHIFT (M, SHIFT = [-1, 1, 0], BOUNDARY = ['\*', '/', '?'], DIM = 2) is  $\begin{bmatrix} * & A & B \\ E & F & / \\ G & H & I \end{bmatrix}$ .

### 17 16.9.68 EPSILON (X)

18 1 **Description.** Model number that is small compared to 1.

19 2 **Class.** [Inquiry function](#).

20 3 **Argument.** X shall be of type real. It may be a scalar or an array.

21 4 **Result Characteristics.** Scalar of the same type and kind type parameter as X.

22 5 **Result Value.** The result has the value  $b^{1-p}$  where  $b$  and  $p$  are as defined in 16.4 for the model representing  
 23 numbers of the same type and kind type parameter as X.

24 6 **Example.** EPSILON (X) has the value  $2^{-23}$  for real X whose model is as in Note 16.4.

### 25 16.9.69 ERF (X)

26 1 **Description.** Error function.

27 2 **Class.** [Elemental function](#).

28 3 **Argument.** X shall be of type real.



1 4 **Result Characteristics.** Same as X.

2 5 **Result Value.** The result has a value equal to a processor-dependent approximation to the error function of X,  
3  $\frac{2}{\sqrt{\pi}} \int_0^X \exp(-t^2) dt$ .

4 6 **Example.** ERF (1.0) has the value 0.843 (approximately).

## 5 16.9.70 ERFC (X)

6 1 **Description.** Complementary error function.

7 2 **Class.** [Elemental](#) function.

8 3 **Argument.** X shall be of type real.

9 4 **Result Characteristics.** Same as X.

10 5 **Result Value.** The result has a value equal to a processor-dependent approximation to the complementary error  
11 function of X,  $1 - \text{ERF}(X)$ ; this is equivalent to  $\frac{2}{\sqrt{\pi}} \int_X^\infty \exp(-t^2) dt$ .

12 6 **Example.** ERFC (1.0) has the value 0.157 (approximately).

## 13 16.9.71 ERFC\_SCALED (X)

14 1 **Description.** Scaled complementary error function.

15 2 **Class.** [Elemental](#) function.

16 3 **Argument.** X shall be of type real.

17 4 **Result Characteristics.** Same as X.

18 5 **Result Value.** The result has a value equal to a processor-dependent approximation to the exponentially-scaled  
19 complementary error function of X,  $\exp(X^2) \frac{2}{\sqrt{\pi}} \int_X^\infty \exp(-t^2) dt$ .

20 6 **Example.** ERFC\_SCALED (20.0) has the value 0.02817434874 (approximately).

### NOTE 16.14

The complementary error function is asymptotic to  $\exp(-X^2)/(X\sqrt{\pi})$ . As such it underflows for  $X > \approx 9$  when using ISO/IEC/IEEE 60559:2011 single precision arithmetic. The exponentially-scaled complementary error function is asymptotic to  $1/(X\sqrt{\pi})$ . As such it does not underflow until  $X > \text{HUGE}(X)/\sqrt{\pi}$ .

## 21 16.9.72 EVENT\_QUERY (EVENT, COUNT [, STAT])

22 1 **Description.** Query event count.

23 2 **Class.** Subroutine.

24 3 **Arguments.**

25 EVENT shall be an [event variable](#) (16.10.2.10). It shall not be [coindexed](#). It is an [INTENT \(IN\)](#) argument.  
26 The EVENT argument is accessed atomically with respect to the execution of [EVENT POST](#)  
27 [statements](#) in unordered segments, in exact analogy to [atomic subroutines](#).

28 COUNT shall be an integer scalar with a decimal range no smaller than that of default integer. It is an  
29 [INTENT \(OUT\)](#) argument. If no error condition occurs, COUNT is assigned the value of the count  
30 of EVENT; otherwise, it is assigned the value  $-1$ .

31 STAT (optional) shall be a noncoindexed integer scalar with a decimal exponent range of at least four. It is an  
32 [INTENT \(OUT\)](#) argument. If the STAT argument is present, it is assigned a processor-dependent

1 positive value if an error condition occurs; otherwise it is assigned the value zero. If the STAT  
2 argument is not present and an error condition occurs, [error termination](#) is initiated.

3 4 **Example.** If EVENT is an [event variable](#) for which there have been no successful posts or waits in preceding  
4 segments, and for which there are no posts or waits in an unordered segment, after execution of

5 CALL EVENT\_QUERY (EVENT, COUNT)

6 the integer variable COUNT will have the value zero. If there have been ten successful posts to EVENT and two  
7 successful waits without an [UNTIL\\_COUNT= specifier](#) in preceding segments, and for which there are no posts  
8 or waits in an unordered segment, after execution of

9 CALL EVENT\_QUERY (EVENT, COUNT)

10 the variable COUNT will have the value eight.

#### NOTE 16.15

Execution of EVENT_QUERY does not imply any synchronization.
--

### 11 16.9.73 EXECUTE\_COMMAND\_LINE (COMMAND [, WAIT, EXITSTAT, 12 CMDSTAT, CMDMSG ])

12 1 **Description.** Execute a command line.

13 2 **Class.** Subroutine.

14 3 **Arguments.**

15 COMMAND shall be a default character scalar. It is an [INTENT \(IN\)](#) argument. Its value is the command line  
16 to be executed. The interpretation is processor dependent.

17 WAIT (optional) shall be a logical scalar. It is an [INTENT \(IN\)](#) argument. If WAIT is present with the value  
18 false, and the processor supports asynchronous execution of the command, the command is executed  
19 asynchronously; otherwise it is executed synchronously.

20 EXITSTAT (optional) shall be a scalar of type integer with a decimal exponent range of at least nine. It is an  
21 [INTENT \(INOUT\)](#) argument. If the command is executed synchronously, it is assigned the value  
22 of the processor-dependent exit status. Otherwise, the value of EXITSTAT is unchanged.

23 CMDSTAT (optional) shall be a scalar of type integer with a decimal exponent range of at least four. It is an  
24 [INTENT \(OUT\)](#) argument. It is assigned the value  $-1$  if the processor does not support command  
25 line execution, a processor-dependent positive value if an error condition occurs, or the value  $-2$   
26 if no error condition occurs but WAIT is present with the value false and the processor does not  
27 support asynchronous execution. Otherwise it is assigned the value 0.

28 CMDMSG (optional) shall be a default character scalar. It is an [INTENT \(INOUT\)](#) argument. If an error condi-  
29 tion occurs, it is assigned a processor-dependent explanatory message. Otherwise, it is unchanged.

30 4 If the processor supports command line execution, it shall support synchronous and may support asynchronous  
31 execution of the command line.

32 5 When the command is executed synchronously, EXECUTE\_COMMAND\_LINE returns after the command line  
33 has completed execution. Otherwise, EXECUTE\_COMMAND\_LINE returns without waiting.

34 6 If a condition occurs that would assign a nonzero value to CMDSTAT but the CMDSTAT variable is not present,  
35 [error termination](#) is initiated.

**16.9.74 EXP (X)**

1 **Description.** Exponential function.

2 **Class.** [Elemental](#) function.

3 **Argument.** X shall be of type real or complex.

4 **Result Characteristics.** Same as X.

5 **Result Value.** The result has a value equal to a processor-dependent approximation to  $e^X$ . If X is of type complex, its imaginary part is regarded as a value in radians.

6 **Example.** EXP (1.0) has the value 2.7182818 (approximately).

**16.9.75 EXPONENT (X)**

1 **Description.** Exponent of floating-point number.

2 **Class.** [Elemental](#) function.

3 **Argument.** X shall be of type real.

4 **Result Characteristics.** Default integer.

5 **Result Value.** The result has a value equal to the exponent  $e$  of the representation for the value of X in the extended real model for the kind of X ([16.4](#)), provided X is nonzero and  $e$  is within the range for default integers. If X has the value zero, the result has the value zero. If X is an IEEE infinity or NaN, the result has the value [HUGE](#) (0).

6 **Examples.** EXPONENT (1.0) has the value 1 and EXPONENT (4.1) has the value 3 for reals whose model is as in Note [16.4](#).

**16.9.76 EXTENDS\_TYPE\_OF (A, MOLD)**

1 **Description.** [Dynamic type](#) extension inquiry.

2 **Class.** [Inquiry function](#).

3 **Arguments.**

A shall be an object of [extensible declared](#) type or [unlimited polymorphic](#). If it is a [polymorphic](#) pointer, it shall not have an undefined association status.

MOLD shall be an object of [extensible declared](#) type or [unlimited polymorphic](#). If it is a [polymorphic](#) pointer, it shall not have an undefined association status.

4 **Result Characteristics.** Default logical scalar.

5 **Result Value.** If MOLD is [unlimited polymorphic](#) and is either a [disassociated](#) pointer or unallocated [allocatable](#) variable, the result is true; otherwise if A is [unlimited polymorphic](#) and is either a [disassociated](#) pointer or unallocated [allocatable](#) variable, the result is false; otherwise if the [dynamic type](#) of A or MOLD is [extensible](#), the result is true if and only if the [dynamic type](#) of A is an [extension type](#) of the [dynamic type](#) of MOLD; otherwise the result is processor dependent.

**NOTE 16.16**

The [dynamic type](#) of a [disassociated](#) pointer or unallocated [allocatable](#) variable is its [declared type](#).

6 **Example.** Given the declarations and assignments

TYPE T1

```

1         REAL C
2         END TYPE
3         TYPE, EXTENDS(T1) :: T2
4         END TYPE
5         CLASS(T1), POINTER :: P, Q
6         ALLOCATE (P)
7         ALLOCATE (T2 :: Q)

```

8 the result of `EXTENDS_TYPE_OF (P, Q)` will be false, and the result of `EXTENDS_TYPE_OF (Q, P)` will  
9 be true.

## 10 16.9.77 FAILED\_IMAGES([TEAM, KIND])

11 1 **Description.** Indices of failed images.

12 2 **Class.** [Transformational function](#).

13 3 **Arguments.**

14 TEAM (optional) shall be a scalar of type [TEAM\\_TYPE](#) from the intrinsic module [ISO\\_FORTRAN\\_ENV](#). Its  
15 value shall be that of the [current](#) or an [ancestor](#) team. If TEAM is absent, the team specified is the  
16 [current team](#).

17 KIND (optional) shall be a scalar integer [constant expression](#).

18 4 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value  
19 of KIND; otherwise, the kind type parameter is that of default integer type. The result is an array of rank one  
20 whose size is equal to the number of images in the specified team that are known by the invoking image to have  
21 failed.

22 5 **Result Value.** The elements of the result are the values of the image indices of the known failed images in the  
23 specified team, in numerically increasing order. If the executing image has previously executed an [image control](#)  
24 [statement](#) whose STAT= specifier assigned the value [STAT\\_FAILED\\_IMAGE](#) from the intrinsic module [ISO\\_-](#)  
25 [FORTRAN\\_ENV](#), or referenced a [collective subroutine](#) whose STAT argument was set to [STAT\\_FAILED\\_-](#)  
26 [IMAGE](#), at least one image in the set of images participating in that [image control statement](#) or [collective](#)  
27 [subroutine](#) reference shall be known to have failed.

28 6 **Examples.** If image 3 is the only image in the [current team](#) that is known by the invoking image to have failed,  
29 `FAILED_IMAGES()` will have the value [3]. If there are no images in the current team that are known by the  
30 invoking image to have failed, the value of `FAILED_IMAGES()` will be a zero-sized array.

## 31 16.9.78 FINDLOC (ARRAY, VALUE, DIM [, MASK, KIND, BACK]) or 32 FINDLOC (ARRAY, VALUE [, MASK, KIND, BACK])

32 1 **Description.** Location(s) of a specified value.

33 2 **Class.** [Transformational function](#).

34 3 **Arguments.**

35 ARRAY shall be an array of intrinsic type.

36 VALUE shall be scalar and in type conformance with ARRAY, as specified in Table 10.2 for the operator  
37 == or the operator [.EQV.](#)

38 DIM shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the [rank](#) of ARRAY.

39 MASK (optional) shall be of type logical and shall be [conformable](#) with ARRAY.

40 KIND (optional) shall be a scalar integer [constant expression](#).

41 BACK (optional) shall be a logical scalar.

1 4 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of  
 2 KIND; otherwise the kind type parameter is that of default integer type. If DIM does not appear, the result is  
 3 an array of **rank** one and of size equal to the **rank** of ARRAY; otherwise, the result is of **rank**  $n - 1$  and shape  
 4  $[d_1, d_2, \dots, d_{\text{DIM}-1}, d_{\text{DIM}+1}, \dots, d_n]$ , where  $[d_1, d_2, \dots, d_n]$  is the shape of ARRAY.

5 5 **Result Value.**

6 *Case (i):* The result of FINDLOC (ARRAY, VALUE) is a rank-one array whose element values are the values  
 7 of the subscripts of an element of ARRAY whose value matches VALUE. If there is such a value,  
 8 the  $i^{\text{th}}$  element value is in the range 1 to  $e_i$ , where  $e_i$  is the extent of the  $i^{\text{th}}$  dimension of ARRAY.  
 9 If no elements match VALUE or ARRAY has size zero, all elements of the result are zero.

10 *Case (ii):* The result of FINDLOC (ARRAY, VALUE, MASK = MASK) is a rank-one array whose element  
 11 values are the values of the subscripts of an element of ARRAY, corresponding to a true element  
 12 of MASK, whose value matches VALUE. If there is such a value, the  $i^{\text{th}}$  element value is in the  
 13 range 1 to  $e_i$ , where  $e_i$  is the extent of the  $i^{\text{th}}$  dimension of ARRAY. If no elements match VALUE,  
 14 ARRAY has size zero, or every element of MASK has the value false, all elements of the result are  
 15 zero.

16 *Case (iii):* If ARRAY has **rank** one, the result of  
 17 FINDLOC (ARRAY, VALUE, DIM=DIM [, MASK = MASK]) is a scalar whose value is equal to  
 18 that of the first element of FINDLOC (ARRAY, VALUE [, MASK = MASK]). Otherwise, the value  
 19 of element  $(s_1, s_2, \dots, s_{\text{DIM}-1}, s_{\text{DIM}+1}, \dots, s_n)$  of the result is equal to FINDLOC (ARRAY  $(s_1,$   
 20  $s_2, \dots, s_{\text{DIM}-1}, \text{;}, s_{\text{DIM}+1}, \dots, s_n)$ , VALUE, DIM=1 [, MASK = MASK  $(s_1, s_2, \dots, s_{\text{DIM}-1}, \text{;},$   
 21  $s_{\text{DIM}+1}, \dots, s_n)$ )).

22 6 If both ARRAY and VALUE are of type logical, the comparison is performed with the **.EQV.** operator; otherwise,  
 23 the comparison is performed with the **==** operator. If the value of the comparison is true, that element of ARRAY  
 24 matches VALUE.

25 7 If DIM is not present, more than one element matches VALUE, and BACK is absent or present with the value  
 26 false, the value returned indicates the first such element, taken in array element order. If DIM is not present and  
 27 BACK is present with the value true, the value returned indicates the last such element, taken in array element  
 28 order.

29 8 **Examples.**

30 *Case (i):* The value of FINDLOC ([2, 6, 4, 6], VALUE=6) is [2], and the value of FINDLOC ([2, 6, 4, 6],  
 31 VALUE=6, BACK = **.TRUE.**) is [4].

32 *Case (ii):* If A has the value  $\begin{bmatrix} 0 & -5 & 7 & 7 \\ 3 & 4 & -1 & 2 \\ 1 & 5 & 6 & 7 \end{bmatrix}$ , and M has the value  $\begin{bmatrix} T & T & F & T \\ T & T & F & T \\ T & T & F & T \end{bmatrix}$ , FINDLOC (A, 7,  
 33 MASK = M) has the value [1, 4] and FINDLOC (A, 7, MASK = M, BACK = **.TRUE.**) has the  
 34 value [3, 4]. This is independent of the declared lower bounds for A.

35 *Case (iii):* The value of FINDLOC ([2, 6, 4], VALUE=6, DIM=1) is 2. If B has the value  
 36  $\begin{bmatrix} 1 & 2 & -9 \\ 2 & 2 & 6 \end{bmatrix}$ , FINDLOC (B, VALUE=2, DIM=1) has the value [2, 1, 0] and FINDLOC (B,  
 37 VALUE=2, DIM=2) has the value [2, 1]. This is independent of the declared lower bounds for B.

### 38 16.9.79 FLOOR (A [, KIND])

39 1 **Description.** Greatest integer less than or equal to A.

40 2 **Class.** **Elemental** function.

41 3 **Arguments.**

42 A shall be of type real.

43 KIND (optional) shall be a scalar integer **constant expression**.

1 4 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of  
2 KIND; otherwise, the kind type parameter is that of default integer type.

3 5 **Result Value.** The result has a value equal to the greatest integer less than or equal to A.

4 6 **Examples.** FLOOR (3.7) has the value 3. FLOOR (-3.7) has the value -4.

## 5 16.9.80 FRACTION (X)

6 1 **Description.** Fractional part of number.

7 2 **Class.** [Elemental](#) function.

8 3 **Argument.** X shall be of type real.

9 4 **Result Characteristics.** Same as X.

10 5 **Result Value.** The result has the value  $X \times b^{-e}$ , where  $b$  and  $e$  are as defined in 16.4 for the representation of  
11 X in the extended real model for the kind of X. If X has the value zero, the result is zero. If X is an [IEEE NaN](#),  
12 the result is that [NaN](#). If X is an IEEE infinity, the result is an [IEEE NaN](#).

13 6 **Example.** FRACTION (3.0) has the value 0.75 for reals whose model is as in Note 16.4.

## 14 16.9.81 GAMMA (X)

15 1 **Description.** Gamma function.

16 2 **Class.** [Elemental](#) function.

17 3 **Argument.** X shall be of type real. Its value shall not be a negative integer or zero.

18 4 **Result Characteristics.** Same as X.

5 **Result Value.** The result has a value equal to a processor-dependent approximation to the gamma function of  
X,

$$\Gamma(X) = \begin{cases} \int_0^{\infty} t^{X-1} \exp(-t) dt & X > 0 \\ \int_0^{\infty} t^{X-1} \left( \exp(-t) - \sum_{k=0}^n \frac{(-t)^k}{k!} \right) dt & -n - 1 < X < -n, n \text{ an integer } \geq 0 \end{cases}$$

19 6 **Example.** GAMMA (1.0) has the value 1.000 (approximately).

## 20 16.9.82 GET\_COMMAND ([COMMAND, LENGTH, STATUS, ERRMSG])

21 1 **Description.** Get program invocation command.

22 2 **Class.** Subroutine.

23 3 **Arguments.**

24 COMMAND (optional) shall be a default character scalar. It is an [INTENT \(OUT\)](#) argument. It is assigned  
25 the entire command by which the program was invoked. If the command cannot be determined,  
26 COMMAND is assigned all blanks.

27 LENGTH (optional) shall be a scalar of type integer with a decimal exponent range of at least four. It is an  
28 [INTENT \(OUT\)](#) argument. It is assigned the significant length of the command by which the  
29 program was invoked. The significant length may include trailing blanks if the processor allows  
30 commands with significant trailing blanks. This length does not consider any possible truncation or  
31 padding in assigning the command to the COMMAND argument; in fact the COMMAND argument  
32 need not even be present. If the command length cannot be determined, a length of 0 is assigned.

1 STATUS (optional) shall be a scalar of type integer with a decimal exponent range of at least four. It is an  
 2 **INTENT (OUT)** argument. It is assigned the value  $-1$  if the COMMAND argument is present and  
 3 has a length less than the significant length of the command. It is assigned a processor-dependent  
 4 positive value if the command retrieval fails. Otherwise it is assigned the value 0.

5 ERRMSG (optional) shall be a default character scalar. It is an **INTENT (INOUT)** argument. It is assigned a  
 6 processor-dependent explanatory message if the command retrieval fails. Otherwise, it is unchanged.

7 4 **Example.** If the program below is invoked with the command “example” on a processor that supports command  
 8 retrieval, it will display “Hello example”.

```

9         PROGRAM hello
10        CHARACTER(:), ALLOCATABLE :: cmd
11        INTEGER :: cmdlen
12        CALL GET_COMMAND(LENGTH=cmdlen)
13        IF (cmdlen>0) THEN
14            ALLOCATE(CHARACTER(cmdlen) :: cmd)
15            CALL GET_COMMAND(cmd)
16            PRINT *, 'Hello ', cmd
17        END IF
18    END PROGRAM
  
```

### 19 16.9.83 GET\_COMMAND\_ARGUMENT (NUMBER [, VALUE, LENGTH, 20 STATUS, ERRMSG])

21 1 **Description.** Get program invocation argument.

22 2 **Class.** Subroutine.

23 3 **Arguments.**

24 NUMBER shall be an integer scalar. It is an **INTENT (IN)** argument that specifies the number of the command  
 25 argument that the other arguments give information about.

26 Command argument 0 always exists, and is the command name by which the program was invoked  
 27 if the processor has such a concept; otherwise, the value of command argument 0 is processor  
 28 dependent. The remaining command arguments are numbered consecutively from 1 to the argument  
 29 count in an order determined by the processor.

30 VALUE (optional) shall be a default character scalar. It is an **INTENT (OUT)** argument. If the command  
 31 argument specified by NUMBER exists, its value is assigned to VALUE; otherwise, VALUE is  
 32 assigned all blanks.

33 LENGTH (optional) shall be a scalar of type integer with a decimal exponent range of at least four. It is an  
 34 **INTENT (OUT)** argument. If the command argument specified by NUMBER exists, its significant  
 35 length is assigned to LENGTH; otherwise, LENGTH is assigned the value zero. It is processor  
 36 dependent whether the significant length includes trailing blanks. This length does not consider any  
 37 possible truncation or padding in assigning the command argument value to the VALUE argument;  
 38 in fact the VALUE argument need not even be present.

39 STATUS (optional) shall be a scalar of type integer with a decimal exponent range of at least four. It is an  
 40 **INTENT (OUT)** argument. If NUMBER is less than zero or greater than the argument count that  
 41 would be returned by the intrinsic function **COMMAND\_ARGUMENT\_COUNT**, or command  
 42 retrieval fails, STATUS is assigned a processor-dependent positive value. Otherwise, if VALUE is  
 43 present and has a length less than the significant length of the specified command argument, it is  
 44 assigned the value  $-1$ . Otherwise it is assigned the value 0.

45 ERRMSG (optional) shall be a default character scalar. It is an **INTENT (INOUT)** argument. It is assigned  
 a processor-dependent explanatory message if the optional argument STATUS is, or would be if

1 present, assigned a positive value. Otherwise, it is unchanged.

2 4 **Example.** On a processor that supports command arguments, the following program displays the arguments of  
3 the command by which it was invoked.

```
4 PROGRAM show_arguments
5     INTEGER :: i
6     CHARACTER :: command*32, arg*128
7     CALL get_command_argument(0, command)
8     WRITE (*,*) "Command name is: ", command
9     DO i = 1, command_argument_count()
10        CALL get_command_argument(i, arg)
11        WRITE (*,*) "Argument ", i, " is ", arg
12    END DO
13 END PROGRAM show_arguments
```

#### 14 16.9.84 GET\_ENVIRONMENT\_VARIABLE (NAME [, VALUE, LENGTH, STATUS, TRIM\_NAME, ERRMSG])

15 1 **Description.** Get environment variable.

16 2 **Class.** Subroutine.

17 3 **Arguments.**

18 NAME shall be a default character scalar. It is an **INTENT (IN)** argument. The interpretation of case is  
19 processor dependent.

20 VALUE (optional) shall be a default character scalar. It is an **INTENT (OUT)** argument. It is assigned the value  
21 of the environment variable specified by NAME. VALUE is assigned all blanks if the environment  
22 variable does not exist or does not have a value or if the processor does not support environment  
23 variables.

24 LENGTH (optional) shall be a scalar of type integer with a decimal exponent range of at least four. It is an  
25 **INTENT (OUT)** argument. If the specified environment variable exists and has a value, LENGTH  
26 is assigned the value of its length. Otherwise LENGTH is assigned the value zero.

27 STATUS (optional) shall be a scalar of type integer with a decimal exponent range of at least four. It is an  
28 **INTENT (OUT)** argument. If the environment variable exists and either has no value, its value is  
29 successfully assigned to VALUE, or the VALUE argument is not present, STATUS is assigned the  
30 value zero. STATUS is assigned the value -1 if the VALUE argument is present and has a length  
31 less than the significant length of the environment variable. It is assigned the value 1 if the specified  
32 environment variable does not exist, or 2 if the processor does not support environment variables.  
33 Processor-dependent values greater than 2 may be assigned for other error conditions.

34 TRIM\_NAME (optional) shall be a logical scalar. It is an **INTENT (IN)** argument. If TRIM\_NAME is present  
35 with the value false then trailing blanks in NAME are considered significant if the processor sup-  
36 ports trailing blanks in environment variable names. Otherwise trailing blanks in NAME are not  
37 considered part of the environment variable's name.

38 ERRMSG (optional) shall be a default character scalar. It is an **INTENT (INOUT)** argument. It is assigned  
39 a processor-dependent explanatory message if the optional argument STATUS is, or would be if  
40 present, assigned a positive value. Otherwise, it is unchanged.

41 4 It is processor dependent whether an environment variable that exists on an **image** also exists on another **image**,  
42 and if it does exist on both **images**, whether the values are the same or different.

43 5 **Example.** If the value of the environment variable DATAFILE is datafile.dat, executing the statement sequence  
44 below will assign the value 'datafile.dat' to FILENAME.



```

1      CHARACTER(:),ALLOCATABLE :: FILENAME
2      INTEGER :: NAMELEN
3      CALL GET_ENVIRONMENT_VARIABLE ("DATAFILE", LENGTH=NAMELEN)
4      IF (LENGTH>0) THEN
5          ALLOCATE(CHARACTER(LENGTH) :: FILENAME)
6          CALL GET_ENVIRONMENT_VARIABLE("DATAFILE", FILENAME)
7      END IF

```

## 8 16.9.85 GET\_TEAM ([LEVEL])

9 1 **Description.** Team.

10 2 **Class.** Transformational function.

11 3 **Argument.** LEVEL (optional) shall be a scalar integer whose value is equal to one of the named constants  
12 [INITIAL\\_TEAM](#), [PARENT\\_TEAM](#), or [CURRENT\\_TEAM](#) from the intrinsic module [ISO\\_FORTRAN\\_ENV](#).

13 4 **Result Characteristics.** Scalar of type [TEAM\\_TYPE](#) from the intrinsic module [ISO\\_FORTRAN\\_ENV](#).

14 5 **Result Value.** The result is a [TEAM\\_TYPE](#) value that identifies the [current team](#) if LEVEL is not present,  
15 present with the value [CURRENT\\_TEAM](#), or if the [current team](#) is the [initial team](#). Otherwise, the result  
16 identifies the [parent team](#) if LEVEL is present with the value [PARENT\\_TEAM](#), and identifies the [initial team](#)  
17 if LEVEL is present with the value [INITIAL\\_TEAM](#).

18 6 **Examples.**

```

19      PROGRAM EXAMPLE1
20          USE,INTRINSIC :: ISO_FORTRAN_ENV, ONLY: TEAM_TYPE
21          TYPE(Team_Type) :: WORLD_TEAM, TEAM2
22
23          ! Define a team variable representing the initial team
24          WORLD_TEAM = GET_TEAM()
25      END PROGRAM
26
27      SUBROUTINE EXAMPLE2 (A)
28          USE,INTRINSIC :: ISO_FORTRAN_ENV, ONLY: TEAM_TYPE
29          REAL A[*]
30          TYPE(Team_Type) :: NEW_TEAM, PARENT_TEAM
31
32          ... ! Form NEW_TEAM
33
34          PARENT_TEAM = GET_TEAM ()
35
36          CHANGE TEAM (NEW_TEAM)
37
38          ! Reference image 1 in parent's team
39          A [1,TEAM=PARENT_TEAM] = 4.2
40
41          ! Reference image 1 in current team
42          A [1] = 9.0
43      END TEAM

```

1           END SUBROUTINE EXAMPLE2

## 2    **16.9.86    HUGE (X)**

3    1 **Description.** Largest model number.

4    2 **Class.** [Inquiry function](#).

5    3 **Argument.** X shall be of type integer or real. It may be a scalar or an array.

6    4 **Result Characteristics.** Scalar of the same type and kind type parameter as X.

7    5 **Result Value.** The result has the value  $r^q - 1$  if X is of type integer and  $(1 - b^{-p})b^{e_{\max}}$  if X is of type real,  
8    where  $r$ ,  $q$ ,  $b$ ,  $p$ , and  $e_{\max}$  are as defined in [16.4](#) for the model representing numbers of the same type and kind  
9    type parameter as X.

10   6 **Example.** HUGE (X) has the value  $(1 - 2^{-24}) \times 2^{127}$  for real X whose model is as in Note [16.4](#).

## 11   **16.9.87    HYPOT (X, Y)**

12   1 **Description.** Euclidean distance function.

13   2 **Class.** [Elemental](#) function.

14   3 **Arguments.**

15    X           shall be of type real.

16    Y           shall be of type real with the same kind type parameter as X.

17   4 **Result Characteristics.** Same as X.

18   5 **Result Value.** The result has a value equal to a processor-dependent approximation to the Euclidean distance,  
19     $\sqrt{X^2 + Y^2}$ , without undue overflow or underflow.

20   6 **Example.** HYPOT (3.0, 4.0) has the value 5.0 (approximately).

## 21   **16.9.88    IACHAR (C [, KIND])**

22   1 **Description.** ASCII code value for character.

23   2 **Class.** [Elemental](#) function.

24   3 **Arguments.**

25    C           shall be of type character and of length one.

26    KIND (optional) shall be a scalar integer [constant expression](#).

27   4 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of  
28    KIND; otherwise, the kind type parameter is that of default integer type.

29   5 **Result Value.** If C is in the [collating sequence](#) defined by the codes specified in ISO/IEC 646:1991 (International  
30    Reference Version), the result is the position of C in that sequence; it is nonnegative and less than or equal to  
31    127. The value of the result is processor dependent if C is not in the [ASCII collating sequence](#). The results  
32    are consistent with the LGE, LGT, LLE, and LLT comparison functions. For example, if LLE (C, D) is true,  
33    IACHAR (C) <= IACHAR (D) is true where C and D are any two characters representable by the processor.

34   6 **Example.** IACHAR ('X') has the value 88.

## 16.9.89 IALL (ARRAY, DIM [, MASK]) or IALL (ARRAY [, MASK])

1 **Description.** Array reduced by IAND function.

2 **Class.** Transformational function.

3 **Arguments.**

4 ARRAY shall be an array of type integer.

5 DIM shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the [rank](#) of ARRAY.

6 MASK (optional) shall be of type logical and shall be [conformable](#) with ARRAY.

7 **Result Characteristics.** The result is of the same type and kind type parameter as ARRAY. It is scalar if DIM does not appear or if ARRAY has [rank](#) one; otherwise, the result is an array of [rank](#)  $n - 1$  and shape  $[d_1, d_2, \dots, d_{\text{DIM}-1}, d_{\text{DIM}+1}, \dots, d_n]$  where  $[d_1, d_2, \dots, d_n]$  is the shape of ARRAY.

8 **Result Value.**

9 *Case (i):* If ARRAY has size zero the result value is equal to NOT (INT (0, KIND (ARRAY))). Otherwise, the result of IALL (ARRAY) has a value equal to the bitwise AND of all the elements of ARRAY.

10 *Case (ii):* The result of IALL (ARRAY, MASK=MASK) has a value equal to IALL (PACK (ARRAY, MASK)).

11 *Case (iii):* The result of IALL (ARRAY, DIM=DIM [, MASK=MASK]) has a value equal to that of IALL (ARRAY [, MASK=MASK]) if ARRAY has [rank](#) one. Otherwise, the value of element  $(s_1, s_2, \dots, s_{\text{DIM}-1}, s_{\text{DIM}+1}, \dots, s_n)$  of the result is equal to IALL (ARRAY  $(s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n)$  [, MASK = MASK  $(s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n)$ ]).

12 **Examples.** IALL ([14, 13, 11]) has the value 8. IALL ([14, 13, 11], MASK=[.true., .false., .true]) has the value 10.

## 16.9.90 IAND (I, J)

1 **Description.** Bitwise AND.

2 **Class.** Elemental function.

3 **Arguments.**

4 I shall be of type integer or a [boz-literal-constant](#).

5 J shall be of type integer or a [boz-literal-constant](#). If both I and J are of type integer, they shall have the same kind type parameter. I and J shall not both be [boz-literal-constants](#).

6 **Result Characteristics.** Same as I if I is of type integer; otherwise, same as J.

7 **Result Value.** If either I or J is a [boz-literal-constant](#), it is first converted as if by the intrinsic function INT to type integer with the kind type parameter of the other. The result has the value obtained by combining I and J bit-by-bit according to the following table:

I	J	IAND (I, J)
1	1	1
1	0	0
0	1	0
0	0	0

8 The model for the interpretation of an integer value as a sequence of bits is in [16.3](#).

9 **Example.** IAND (1, 3) has the value 1.

### 16.9.91 IANY (ARRAY, DIM [, MASK]) or IANY (ARRAY [, MASK])

1 **Description.** Reduce array with bitwise OR operation.

2 **Class.** Transformational function.

3 **Arguments.**

4 ARRAY shall be of type integer. It shall be an array.

5 DIM shall be an integer scalar with a value in the range  $1 \leq DIM \leq n$ , where  $n$  is the rank of ARRAY.

6 MASK (optional) shall be of type logical and shall be conformable with ARRAY.

7 **Result Characteristics.** The result is of the same type and kind type parameter as ARRAY. It is scalar if DIM does not appear or if ARRAY has rank one; otherwise, the result is an array of rank  $n - 1$  and shape  $[d_1, d_2, \dots, d_{DIM-1}, d_{DIM+1}, \dots, d_n]$  where  $[d_1, d_2, \dots, d_n]$  is the shape of ARRAY.

8 **Result Value.**

9 *Case (i):* The result of IANY (ARRAY) is the bitwise OR of all the elements of ARRAY. If ARRAY has size zero the result value is equal to zero.

10 *Case (ii):* The result of IANY (ARRAY, MASK=MASK) has a value equal to IANY (PACK (ARRAY, MASK)).

11 *Case (iii):* The result of IANY (ARRAY, DIM=DIM [, MASK=MASK]) has a value equal to that of IANY (ARRAY [, MASK=MASK]) if ARRAY has rank one. Otherwise, the value of element  $(s_1, s_2, \dots, s_{DIM-1}, s_{DIM+1}, \dots, s_n)$  of the result is equal to IANY (ARRAY  $(s_1, s_2, \dots, s_{DIM-1}, :, s_{DIM+1}, \dots, s_n)$  [, MASK = MASK  $(s_1, s_2, \dots, s_{DIM-1}, :, s_{DIM+1}, \dots, s_n)$ ]).

12 **Examples.** IANY ([14, 13, 8]) has the value 15. IANY ([14, 13, 8], MASK=[.true., .false., .true]) has the value 14.

### 16.9.92 IBCLR (I, POS)

1 **Description.** I with bit POS replaced by zero.

2 **Class.** Elemental function.

3 **Arguments.**

4 I shall be of type integer.

5 POS shall be of type integer. It shall be nonnegative and less than BIT\_SIZE (I).

6 **Result Characteristics.** Same as I.

7 **Result Value.** The result has the value of the sequence of bits of I, except that bit POS is zero. The model for the interpretation of an integer value as a sequence of bits is in 16.3.

8 **Examples.** IBCLR (14, 1) has the value 12. If V has the value [1, 2, 3, 4], the value of IBCLR (POS = V, I = 31) is [29, 27, 23, 15].

### 16.9.93 IBITS (I, POS, LEN)

1 **Description.** Specified sequence of bits.

2 **Class.** Elemental function.

3 **Arguments.**

4 I shall be of type integer.

5 POS shall be of type integer. It shall be nonnegative and POS + LEN shall be less than or equal to BIT\_SIZE (I).

1        LEN            shall be of type integer and nonnegative.

2    4 **Result Characteristics.** Same as I.

3    5 **Result Value.** The result has the value of the sequence of LEN bits in I beginning at bit POS, right-adjusted  
4        and with all other bits zero. The model for the interpretation of an integer value as a sequence of bits is in 16.3.

5    6 **Example.** IBITS (14, 1, 3) has the value 7.

## 6    6 **16.9.94    IBSET (I, POS)**

7    1 **Description.** I with bit POS replaced by one.

8    2 **Class.** [Elemental](#) function.

9    3 **Arguments.**

10    I                shall be of type integer.

11    POS             shall be of type integer. It shall be nonnegative and less than BIT\_SIZE (I).

12    4 **Result Characteristics.** Same as I.

13    5 **Result Value.** The result has the value of the sequence of bits of I, except that bit POS is one. The model for  
14        the interpretation of an integer value as a sequence of bits is in 16.3.

15    6 **Examples.** IBSET (12, 1) has the value 14. If V has the value [1, 2, 3, 4], the value of IBSET (POS = V, I = 0)  
16        is [2, 4, 8, 16].

## 17    6 **16.9.95    ICHAR (C [, KIND])**

18    1 **Description.** Code value for character.

19    2 **Class.** [Elemental](#) function.

20    3 **Arguments.**

21    C                shall be of type character and of length one. Its value shall be that of a character capable of  
22        representation in the processor.

23    KIND (optional) shall be a scalar integer [constant expression](#).

24    4 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of  
25        KIND; otherwise, the kind type parameter is that of default integer type.

26    5 **Result Value.** The result is the position of C in the processor [collating sequence](#) associated with the kind type  
27        parameter of C; it is nonnegative and less than *n*, where *n* is the number of characters in the [collating sequence](#).  
28        The [kind type parameter](#) of the result shall specify an integer kind that is capable of representing *n*. For any char-  
29        acters C and D capable of representation in the processor, C <= D is true if and only if ICHAR (C) <= ICHAR (D)  
30        is true and C == D is true if and only if ICHAR (C) == ICHAR (D) is true.

31    6 **Example.** ICHAR ('X') has the value 88 on a processor using the [ASCII collating sequence](#) for default characters.

## 32    6 **16.9.96    IEOR (I, J)**

33    1 **Description.** Bitwise exclusive OR.

34    2 **Class.** [Elemental](#) function.

35    3 **Arguments.**

36    I                shall be of type integer or a [boz-literal-constant](#).

1 J shall be of type integer or a *boz-literal-constant*. If both I and J are of type integer, they shall have  
2 the same kind type parameter. I and J shall not both be *boz-literal-constants*.

3 4 **Result Characteristics.** Same as I if I is of type integer; otherwise, same as J.

4 5 **Result Value.** If either I or J is a *boz-literal-constant*, it is first converted as if by the intrinsic function `INT` to  
5 type integer with the kind type parameter of the other. The result has the value obtained by combining I and J  
6 bit-by-bit according to the following table:

I	J	IEOR (I, J)
1	1	0
1	0	1
0	1	1
0	0	0

7 6 The model for the interpretation of an integer value as a sequence of bits is in [16.3](#).

8 7 **Example.** `IEOR (1, 3)` has the value 2.

### 9 16.9.97 **IMAGE\_INDEX (COARRAY, SUB) or (COARRAY, SUB, TEAM) or (COARRAY, SUB, TEAM\_NUMBER)**

10 1 **Description.** Image index from `cosubscripts`.

11 2 **Class.** Inquiry function.

12 3 **Arguments.**

13 COARRAY shall be a `coarray` of any type. If `TEAM_NUMBER` appears and the `current team` is not the `initial`  
14 `team`, it shall be established in an `ancestor` of the `current team`. Otherwise, if `TEAM` appears, it  
15 shall be established in that `team`. Otherwise, it shall be established in the `current team`.

16 SUB shall be a rank-one integer array of size equal to the `corank` of COARRAY.

17 TEAM shall be a scalar of type `TEAM_TYPE` from the intrinsic module `ISO_FORTRAN_ENV`, with a  
18 value that identifies the `current` or an `ancestor team`.

19 TEAM\_NUMBER shall be an integer scalar. It shall identify the `initial team` or a `team` whose `parent` is the  
20 same as that of the `current team`.

21 4 **Result Characteristics.** Default integer scalar.

22 5 **Result Value.** If the value of SUB is a valid sequence of `cosubscripts` for COARRAY in the `team` specified by  
23 TEAM or TEAM\_NUMBER, or the `current team` if neither TEAM nor TEAM\_NUMBER appears, the result  
24 is the `index` of the corresponding `image` in that `team`. Otherwise, the result is zero.

25 6 **Examples.** If A and B are declared as A [0:\*] and B (10, 20) [10, 0:9, 0:\*] respectively, `IMAGE_INDEX (A, [0])`  
26 has the value 1 and `IMAGE_INDEX (B, [3, 1, 2])` has the value 213 (on any `image`).

### 27 16.9.98 **IMAGE\_STATUS (IMAGE [, TEAM])**

28 1 **Description.** Image execution state.

29 2 **Class.** Elemental function.

30 3 **Arguments.**

31 IMAGE shall be of type integer. Its value shall be positive and less than or equal to the number of `images`  
32 in the specified `team`.

1 TEAM (optional) shall be a scalar of type `TEAM_TYPE` from the intrinsic module `ISO_FORTRAN_ENV`. Its  
 2 value shall represent the `current` or an `ancestor` team. If TEAM is absent, the team specified is the  
 3 `current team`.

4 **4 Result Characteristics.** Default integer.

5 **5 Result Value.** The result value is `STAT_FAILED_IMAGE` from the intrinsic module `ISO_FORTRAN_ENV`  
 6 if the specified image has failed, `STAT_STOPPED_IMAGE` from the intrinsic module `ISO_FORTRAN_ENV`  
 7 if that image has initiated `normal termination`, and zero otherwise.

8 **6 Example.** If image 3 of the current team has failed, `IMAGE_STATUS (3)` has the value `STAT_FAILED_-`  
 9 `IMAGE`.

## 10 **16.9.99 INDEX (STRING, SUBSTRING [, BACK, KIND])**

11 **1 Description.** Character string search.

12 **2 Class.** Elemental function.

13 **3 Arguments.**

14 `STRING` shall be of type character.

15 `SUBSTRING` shall be of type character with the same `kind type parameter` as `STRING`.

16 `BACK` (optional) shall be of type logical.

17 `KIND` (optional) shall be a scalar integer `constant expression`.

18 **4 Result Characteristics.** Integer. If `KIND` is present, the `kind type parameter` is that specified by the value of  
 19 `KIND`; otherwise the `kind type parameter` is that of default integer type.

20 **5 Result Value.**

21 *Case (i):* If `STRING % LEN < SUBSTRING % LEN`, the result has the value zero.

22 *Case (ii):* Otherwise, if there is an integer `I` in the range  $1 \leq I \leq \text{STRING \% LEN} - \text{SUBSTRING \% LEN} + 1$ , such that `STRING(I : I + SUBSTRING % LEN - 1)` is equal to `SUBSTRING`, the result has  
 23 the value of the smallest such `I` if `BACK` is absent or present with the value `false`, and the greatest  
 24 such `I` if `BACK` is present with the value `true`.  
 25

26 *Case (iii):* Otherwise, the result has the value zero.

27 **6 Examples.** `INDEX ('FORTRAN', 'R')` has the value 3.

28 `INDEX ('FORTRAN', 'R', BACK = .TRUE.)` has the value 5.

## 29 **16.9.100 INT (A [, KIND])**

30 **1 Description.** Conversion to integer type.

31 **2 Class.** Elemental function.

32 **3 Arguments.**

33 `A` shall be of type integer, real, or complex, or a *boz-literal-constant*.

34 `KIND` (optional) shall be a scalar integer `constant expression`.

35 **4 Result Characteristics.** Integer. If `KIND` is present, the `kind type parameter` is that specified by the value of  
 36 `KIND`; otherwise, the `kind type parameter` is that of default integer type.

37 **5 Result Value.**

38 *Case (i):* If `A` is of type integer, `INT (A) = A`.

- 1 *Case (ii):* If A is of type real, there are two cases: if  $|A| < 1$ , INT (A) has the value 0; if  $|A| \geq 1$ , INT (A)  
 2 is the integer whose magnitude is the largest integer that does not exceed the magnitude of A and  
 3 whose sign is the same as the sign of A.
- 4 *Case (iii):* If A is of type complex, INT (A) = INT (REAL (A, KIND (A))).
- 5 *Case (iv):* If A is a *boz-literal-constant*, the value of the result is the value whose bit sequence according to the  
 6 model in 16.3 is the same as that of A as modified by padding or truncation according to 16.3.3.  
 7 The interpretation of a bit sequence whose most significant bit is 1 is processor dependent.

8 6 **Example.** INT (-3.7) has the value -3.

## 9 16.9.101 IOR (I, J)

10 1 **Description.** Bitwise inclusive OR.

11 2 **Class.** Elemental function.

12 3 **Arguments.**

13 I shall be of type integer or a *boz-literal-constant*.

14 J shall be of type integer or a *boz-literal-constant*. If both I and J are of type integer, they shall have  
 15 the same kind type parameter. I and J shall not both be *boz-literal-constants*.

16 4 **Result Characteristics.** Same as I if I is of type integer; otherwise, same as J.

17 5 **Result Value.** If either I or J is a *boz-literal-constant*, it is first converted as if by the intrinsic function INT to  
 18 type integer with the kind type parameter of the other. The result has the value obtained by combining I and J  
 19 bit-by-bit according to the following table:

I	J	IOR (I, J)
1	1	1
1	0	1
0	1	1
0	0	0

20 6 The model for the interpretation of an integer value as a sequence of bits is in 16.3.

21 7 **Example.** IOR (5, 3) has the value 7.

## 22 16.9.102 IPARITY (ARRAY, DIM [, MASK]) or IPARITY (ARRAY [, MASK])

23 1 **Description.** Array reduced by IEXOR function.

24 2 **Class.** Transformational function.

25 3 **Arguments.**

26 ARRAY shall be of type integer. It shall be an array.

27 DIM shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the rank of ARRAY.

28 MASK (optional) shall be of type logical and shall be conformable with ARRAY.

29 4 **Result Characteristics.** The result is of the same type and kind type parameter as ARRAY. It is scalar if  
 30 DIM does not appear; otherwise, the result has rank  $n - 1$  and shape  $[d_1, d_2, \dots, d_{\text{DIM}-1}, d_{\text{DIM}+1}, \dots, d_n]$  where  
 31  $[d_1, d_2, \dots, d_n]$  is the shape of ARRAY.

32 5 **Result Value.**

33 *Case (i):* The result of IPARITY (ARRAY) has a value equal to the bitwise exclusive OR of all the elements  
 34 of ARRAY. If ARRAY has size zero the result has the value zero.



1 *Case (ii):* The result of IPARITY (ARRAY, MASK=MASK) has a value equal to that of IPARITY (PACK  
2 (ARRAY, MASK)).

3 *Case (iii):* The result of IPARITY (ARRAY, DIM=DIM [, MASK=MASK]) has a value equal to that of  
4 IPARITY (ARRAY [, MASK=MASK]) if ARRAY has **rank** one. Otherwise, the value of element  
5 ( $s_1, s_2, \dots, s_{DIM-1}, s_{DIM+1}, \dots, s_n$ ) of the result is equal to IPARITY (ARRAY ( $s_1, s_2, \dots,$   
6  $s_{DIM-1}, :, s_{DIM+1}, \dots, s_n$ ) [, MASK = MASK ( $s_1, s_2, \dots, s_{DIM-1}, :, s_{DIM+1}, \dots, s_n$ )]).

7 6 **Examples.** IPARITY ([14, 13, 8]) has the value 11. IPARITY ([14, 13, 8], MASK=[.true., .false., .true]) has the  
8 value 6.

### 9 16.9.103 ISHFT (I, SHIFT)

10 1 **Description.** Logical shift.

11 2 **Class.** [Elemental](#) function.

12 3 **Arguments.**

13 I shall be of type integer.

14 SHIFT shall be of type integer. The absolute value of SHIFT shall be less than or equal to BIT\_SIZE (I).

15 4 **Result Characteristics.** Same as I.

16 5 **Result Value.** The result has the value obtained by shifting the bits of I by SHIFT positions. If SHIFT is  
17 positive, the shift is to the left; if SHIFT is negative, the shift is to the right; if SHIFT is zero, no shift is  
18 performed. Bits shifted out from the left or from the right, as appropriate, are lost. Zeros are shifted in from the  
19 opposite end. The model for the interpretation of an integer value as a sequence of bits is in [16.3](#).

20 6 **Example.** ISHFT (3, 1) has the value 6.

### 21 16.9.104 ISHFTC (I, SHIFT [, SIZE])

22 1 **Description.** Circular shift of the rightmost bits.

23 2 **Class.** [Elemental](#) function.

24 3 **Arguments.**

25 I shall be of type integer.

26 SHIFT shall be of type integer. The absolute value of SHIFT shall be less than or equal to SIZE.

27 SIZE (optional) shall be of type integer. The value of SIZE shall be positive and shall not exceed BIT\_SIZE (I).  
28 If SIZE is absent, it is as if it were present with the value of BIT\_SIZE (I).

29 4 **Result Characteristics.** Same as I.

30 5 **Result Value.** The result has the value obtained by shifting the SIZE rightmost bits of I circularly by SHIFT  
31 positions. If SHIFT is positive, the shift is to the left; if SHIFT is negative, the shift is to the right; and if SHIFT  
32 is zero, no shift is performed. No bits are lost. The unshifted bits are unaltered. The model for the interpretation  
33 of an integer value as a sequence of bits is in [16.3](#).

34 6 **Example.** ISHFTC (3, 2, 3) has the value 5.

### 35 16.9.105 IS\_CONTIGUOUS (ARRAY)

36 1 **Description.** Array contiguity test ([8.5.7](#)).

37 2 **Class.** [Inquiry](#) function.

38 3 **Argument.** ARRAY may be of any type. It shall be [assumed-rank](#) or an array. If it is a pointer it shall be

1 associated.

2 4 **Result Characteristics.** Default logical scalar.

3 5 **Result Value.** The result has the value true if ARRAY has rank zero or is [contiguous](#), and false otherwise.

4 6 **Example.** After the [pointer assignment](#) AP => TARGET (1:10:2), IS\_CONTIGUOUS (AP) has the value  
5 false.

### 6 16.9.106 IS\_IOSTAT\_END (I)

7 1 **Description.** IOSTAT value test for end of file.

8 2 **Class.** [Elemental](#) function.

9 3 **Argument.** I shall be of type integer.

10 4 **Result Characteristics.** Default logical.

11 5 **Result Value.** The result has the value true if and only if I is a value for the *scalar-int-variable* in an [IOSTAT=](#)  
12 [specifier](#) (12.11.5) that would indicate an end-of-file condition.

### 13 16.9.107 IS\_IOSTAT\_EOR (I)

14 1 **Description.** IOSTAT value test for end of record.

15 2 **Class.** [Elemental](#) function.

16 3 **Argument.** I shall be of type integer.

17 4 **Result Characteristics.** Default logical.

18 5 **Result Value.** The result has the value true if and only if I is a value for the *scalar-int-variable* in an [IOSTAT=](#)  
19 [specifier](#) (12.11.5) that would indicate an end-of-record condition.

### 20 16.9.108 KIND (X)

21 1 **Description.** Value of the kind type parameter of X.

22 2 **Class.** [Inquiry function](#).

23 3 **Argument.** X may be of any intrinsic type. It may be a scalar or an array.

24 4 **Result Characteristics.** Default integer scalar.

25 5 **Result Value.** The result has a value equal to the kind type parameter value of X.

26 6 **Example.** KIND (0.0) has the kind type parameter value of default real.

### 27 16.9.109 LBOUND (ARRAY [, DIM, KIND])

28 1 **Description.** Lower bound(s).

29 2 **Class.** [Inquiry function](#).

30 3 **Arguments.**

31 ARRAY shall be [assumed-rank](#) or an array. It shall not be an unallocated [allocatable](#) variable or a pointer  
32 that is not associated.

33 DIM (optional) shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the [rank](#) of ARRAY.  
34 The corresponding [actual argument](#) shall not be an optional dummy argument, a disassociated

1 pointer, or an unallocated allocatable.

2 KIND (optional) shall be a scalar integer [constant expression](#).

3 4 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of  
4 KIND; otherwise the kind type parameter is that of default integer type. The result is scalar if DIM is present;  
5 otherwise, the result is an array of [rank](#) one and size  $n$ , where  $n$  is the [rank](#) of ARRAY.

6 5 **Result Value.**

7 *Case (i):* If DIM is present, ARRAY is a [whole array](#), and either ARRAY is an [assumed-size array](#) of [rank](#)  
8 DIM or dimension DIM of ARRAY has nonzero extent, the result has a value equal to the lower  
9 bound for subscript DIM of ARRAY. Otherwise, if DIM is present, the result value is 1.

10 *Case (ii):* LBOUND (ARRAY) has a value whose  $i^{th}$  element is equal to LBOUND (ARRAY,  $i$ ), for  $i = 1, 2,$   
11  $\dots, n$ , where  $n$  is the [rank](#) of ARRAY. LBOUND (ARRAY, KIND=KIND) has a value whose  $i^{th}$   
12 element is equal to LBOUND (ARRAY,  $i$ , KIND), for  $i = 1, 2, \dots, n$ , where  $n$  is the [rank](#) of  
13 ARRAY.

#### NOTE 16.17

If ARRAY is [assumed-rank](#) and has rank zero, DIM cannot be present since it cannot satisfy the requirement  $1 \leq \text{DIM} \leq 0$ .

14 6 **Examples.** If A is declared by the statement

15 7 REAL A (2:3, 7:10)

16 8 then LBOUND (A) is [2, 7] and LBOUND (A, DIM=2) is 7.

### 17 16.9.110 LCOBOUND (COARRAY [, DIM, KIND])

18 1 **Description.** Lower [cobound](#)(s) of a [coarray](#).

19 2 **Class.** [Inquiry function](#).

20 3 **Arguments.**

21 COARRAY shall be a [coarray](#) and may be of any type. It may be a scalar or an array. If it is [allocatable](#) it  
22 shall be allocated.

23 DIM (optional) shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the [corank](#)  
24 of COARRAY. The corresponding [actual argument](#) shall not be an optional dummy argument, a  
25 disassociated pointer, or an unallocated allocatable.

26 KIND (optional) shall be a scalar integer [constant expression](#).

27 4 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of  
28 KIND; otherwise, the kind type parameter is that of default integer type. The result is scalar if DIM is present;  
29 otherwise, the result is an array of [rank](#) one and size  $n$ , where  $n$  is the [corank](#) of COARRAY.

30 5 **Result Value.**

31 *Case (i):* If DIM is present, the result has a value equal to the lower [cobound](#) for [codimension](#) DIM of  
32 COARRAY.

33 *Case (ii):* If DIM is absent, the result has a value whose  $i^{th}$  element is equal to the lower [cobound](#) for [codi-](#)  
34 [mension](#)  $i$  of COARRAY, for  $i = 1, 2, \dots, n$ , where  $n$  is the [corank](#) of COARRAY.

35 6 **Examples.** If A is allocated by the statement ALLOCATE (A [2:3, 7:\*) then LCOBOUND (A) is [2, 7] and  
36 LCOBOUND (A, DIM=2) is 7.

### 16.9.111 LEADZ (I)

1 **Description.** Number of leading zero bits.

2 **Class.** [Elemental](#) function.

3 **Argument.** I shall be of type integer.

4 **Result Characteristics.** Default integer.

5 **Result Value.** If all of the bits of I are zero, the result has the value BIT\_SIZE (I). Otherwise, the result has  
6 the value BIT\_SIZE (I) – 1 – k, where k is the position of the leftmost 1 bit in I. The model for the interpretation  
7 of an integer value as a sequence of bits is in [16.3](#).

8 **Examples.** LEADZ (1) has the value 31 if BIT\_SIZE (1) has the value 32.

### 16.9.112 LEN (STRING [, KIND])

1 **Description.** Length of a character entity.

2 **Class.** [Inquiry](#) function.

3 **Arguments.**

4 STRING shall be of type character. If it is an unallocated [allocatable](#) variable or a pointer that is not  
5 associated, its [length type parameter](#) shall not be [deferred](#).

6 KIND (optional) shall be a scalar integer [constant expression](#).

7 **Result Characteristics.** Integer scalar. If KIND is present, the kind type parameter is that specified by the  
8 value of KIND; otherwise the kind type parameter is that of default integer type.

9 **Result Value.** The result has a value equal to the number of characters in STRING if it is scalar or in an  
10 element of STRING if it is an array.

11 **Example.** If C is declared by the statement

```
12 CHARACTER (11) C (100)
```

13 LEN (C) has the value 11.

### 16.9.113 LEN\_TRIM (STRING [, KIND])

1 **Description.** Length without trailing blanks.

2 **Class.** [Elemental](#) function.

3 **Arguments.**

4 STRING shall be of type character.

5 KIND (optional) shall be a scalar integer [constant expression](#).

6 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of  
7 KIND; otherwise the kind type parameter is that of default integer type.

8 **Result Value.** The result has a value equal to the number of characters remaining after any trailing blanks in  
9 STRING are removed. If the argument contains no nonblank characters, the result is zero.

10 **Examples.** LEN\_TRIM (' A B ') has the value 4 and LEN\_TRIM (' ') has the value 0.

### 16.9.114 LGE (STRING\_A, STRING\_B)

1 **Description.** ASCII greater than or equal.

2 **Class.** [Elemental](#) function.

3 **Arguments.**

4 STRING\_A shall be default character or [ASCII character](#).

5 STRING\_B shall be of type character with the same kind type parameter as STRING\_A.

6 **Result Characteristics.** Default logical.

7 **Result Value.** If the strings are of unequal length, the comparison is made as if the shorter string were extended on the right with blanks to the length of the longer string. If either string contains a character not in the [ASCII character](#) set, the result is processor dependent. The result is true if the strings are equal or if STRING\_A follows STRING\_B in the [ASCII collating sequence](#); otherwise, the result is false.

#### NOTE 16.18

The result is true if both STRING\_A and STRING\_B are of zero length.

8 **Example.** LGE ('ONE', 'TWO') has the value false.

### 16.9.115 LGT (STRING\_A, STRING\_B)

9 **Description.** ASCII greater than.

10 **Class.** [Elemental](#) function.

11 **Arguments.**

12 STRING\_A shall be default character or [ASCII character](#).

13 STRING\_B shall be of type character with the same kind type parameter as STRING\_A.

14 **Result Characteristics.** Default logical.

15 **Result Value.** If the strings are of unequal length, the comparison is made as if the shorter string were extended on the right with blanks to the length of the longer string. If either string contains a character not in the [ASCII character](#) set, the result is processor dependent. The result is true if STRING\_A follows STRING\_B in the [ASCII collating sequence](#); otherwise, the result is false.

#### NOTE 16.19

The result is false if both STRING\_A and STRING\_B are of zero length.

16 **Example.** LGT ('ONE', 'TWO') has the value false.

### 16.9.116 LLE (STRING\_A, STRING\_B)

17 **Description.** ASCII less than or equal.

18 **Class.** [Elemental](#) function.

19 **Arguments.**

20 STRING\_A shall be default character or [ASCII character](#).

21 STRING\_B shall be of type character with the same kind type parameter as STRING\_A.

22 **Result Characteristics.** Default logical.

1 5 **Result Value.** If the strings are of unequal length, the comparison is made as if the shorter string were extended  
 2 on the right with blanks to the length of the longer string. If either string contains a character not in the [ASCII](#)  
 3 [character](#) set, the result is processor dependent. The result is true if the strings are equal or if `STRING_A`  
 4 precedes `STRING_B` in the [ASCII collating sequence](#); otherwise, the result is false.

**NOTE 16.20**

The result is true if both `STRING_A` and `STRING_B` are of zero length.

5 6 **Example.** `LLE ('ONE', 'TWO')` has the value true.

**16.9.117 LLT (STRING\_A, STRING\_B)**

7 1 **Description.** ASCII less than.

8 2 **Class.** [Elemental](#) function.

9 3 **Arguments.**

10 `STRING_A` shall be default character or [ASCII character](#).

11 `STRING_B` shall be of type character with the same kind type parameter as `STRING_A`.

12 4 **Result Characteristics.** Default logical.

13 5 **Result Value.** If the strings are of unequal length, the comparison is made as if the shorter string were extended  
 14 on the right with blanks to the length of the longer string. If either string contains a character not in the [ASCII](#)  
 15 [character](#) set, the result is processor dependent. The result is true if `STRING_A` precedes `STRING_B` in the  
 16 [ASCII collating sequence](#); otherwise, the result is false.

**NOTE 16.21**

The result is false if both `STRING_A` and `STRING_B` are of zero length.

17 6 **Example.** `LLT ('ONE', 'TWO')` has the value true.

**16.9.118 LOG (X)**

19 1 **Description.** Natural logarithm.

20 2 **Class.** [Elemental](#) function.

21 3 **Argument.** `X` shall be of type real or complex. If `X` is real, its value shall be greater than zero. If `X` is complex,  
 22 its value shall not be zero.

23 4 **Result Characteristics.** Same as `X`.

24 5 **Result Value.** The result has a value equal to a processor-dependent approximation to  $\log_e X$ . A result of type  
 25 complex is the principal value with imaginary part  $\omega$  in the range  $-\pi \leq \omega \leq \pi$ . If the real part of `X` is less  
 26 than zero and the imaginary part of `X` is zero, then the imaginary part of the result is approximately  $\pi$  if the  
 27 imaginary part of `X` is positive real zero or the processor does not distinguish between positive and negative real  
 28 zero, and approximately  $-\pi$  if the imaginary part of `X` is negative real zero.

29 6 **Example.** `LOG (10.0)` has the value 2.3025851 (approximately).

**16.9.119 LOG\_GAMMA (X)**

31 1 **Description.** Logarithm of the absolute value of the gamma function.

32 2 **Class.** [Elemental](#) function.

1 3 **Argument.** X shall be of type real. Its value shall not be a negative integer or zero.

2 4 **Result Characteristics.** Same as X.

3 5 **Result Value.** The result has a value equal to a processor-dependent approximation to the natural logarithm  
4 of the absolute value of the gamma function of X.

5 6 **Example.** LOG\_GAMMA (3.0) has the value 0.693 (approximately).

### 6 16.9.120 LOG10 (X)

7 1 **Description.** Common logarithm.

8 2 **Class.** [Elemental](#) function.

9 3 **Argument.** X shall be of type real. The value of X shall be greater than zero.

10 4 **Result Characteristics.** Same as X.

11 5 **Result Value.** The result has a value equal to a processor-dependent approximation to  $\log_{10}X$ .

12 6 **Example.** LOG10 (10.0) has the value 1.0 (approximately).

### 13 16.9.121 LOGICAL (L [, KIND])

14 1 **Description.** Conversion between kinds of logical.

15 2 **Class.** [Elemental](#) function.

16 3 **Arguments.**

17 L shall be of type logical.

18 KIND (optional) shall be a scalar integer [constant expression](#).

19 4 **Result Characteristics.** Logical. If KIND is present, the kind type parameter is that specified by the value of  
20 KIND; otherwise, the kind type parameter is that of default logical.

21 5 **Result Value.** The value is that of L.

22 6 **Example.** LOGICAL (L [.OR.](#) [.NOT.](#) L) has the value true and is default logical, regardless of the kind type  
23 parameter of the logical variable L.

### 24 16.9.122 MASKL (I [, KIND])

25 1 **Description.** Left justified mask.

26 2 **Class.** [Elemental](#) function.

27 3 **Arguments.**

28 I shall be of type integer. It shall be nonnegative and less than or equal to the number of bits  $z$  of  
29 the model integer defined for bit manipulation contexts in [16.3](#) for the kind of the result.

30 KIND (optional) shall be a scalar integer [constant expression](#).

31 4 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of  
32 KIND; otherwise, the kind type parameter is that of default integer type.

33 5 **Result Value.** The result value has its leftmost I bits set to 1 and the remaining bits set to 0. The model for  
34 the interpretation of an integer value as a sequence of bits is in [16.3](#).

35 6 **Example.** MASKL (3) has the value SHIFTL (7, BIT\_SIZE (0) – 3).

### 16.9.123 MASKR (I [, KIND])

1 **Description.** Right justified mask.

2 **Class.** [Elemental](#) function.

3 **Arguments.**

4 I shall be of type integer. It shall be nonnegative and less than or equal to the number of bits  $z$  of  
5 the model integer defined for bit manipulation contexts in [16.3](#) for the kind of the result.

6 KIND (optional) shall be a scalar integer [constant expression](#).

7 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of  
8 KIND; otherwise, the kind type parameter is that of default integer type.

9 **Result Value.** The result value has its rightmost I bits set to 1 and the remaining bits set to 0. The model for  
10 the interpretation of an integer value as a sequence of bits is in [16.3](#).

11 **Example.** MASKR (3) has the value 7.

### 16.9.124 MATMUL (MATRIX\_A, MATRIX\_B)

12 **Description.** Matrix multiplication.

13 **Class.** [Transformational](#) function.

14 **Arguments.**

15 MATRIX\_A shall be a rank-one or rank-two array of [numeric type](#) or logical type.

16 MATRIX\_B shall be of [numeric type](#) if MATRIX\_A is of [numeric type](#) and of logical type if MATRIX\_A is of  
17 logical type. It shall be an array of [rank](#) one or two. MATRIX\_A and MATRIX\_B shall not both  
18 have [rank](#) one. The size of the first (or only) dimension of MATRIX\_B shall equal the size of the  
19 last (or only) dimension of MATRIX\_A.

20 **Result Characteristics.** If the arguments are of [numeric type](#), the type and kind type parameter of the result  
21 are determined by the types of the arguments as specified in [10.1.9.3](#) for the \* operator. If the arguments are of  
22 type logical, the result is of type logical with the kind type parameter of the arguments as specified in [10.1.9.3](#)  
23 for the [.AND.](#) operator. The shape of the result depends on the shapes of the arguments as follows:

24 *Case (i):* If MATRIX\_A has shape  $[n, m]$  and MATRIX\_B has shape  $[m, k]$ , the result has shape  $[n, k]$ .

25 *Case (ii):* If MATRIX\_A has shape  $[m]$  and MATRIX\_B has shape  $[m, k]$ , the result has shape  $[k]$ .

26 *Case (iii):* If MATRIX\_A has shape  $[n, m]$  and MATRIX\_B has shape  $[m]$ , the result has shape  $[n]$ .

27 **Result Value.**

28 *Case (i):* Element  $(i, j)$  of the result has the value SUM (MATRIX\_A  $(i, :)$  \* MATRIX\_B  $(:, j)$ ) if the  
29 arguments are of [numeric type](#) and has the value ANY (MATRIX\_A  $(i, :)$  [.AND.](#) MATRIX\_B  $(:, j)$ ) if the arguments are of logical type.

30 *Case (ii):* Element  $(j)$  of the result has the value SUM (MATRIX\_A  $(:)$  \* MATRIX\_B  $(:, j)$ ) if the arguments  
31 are of [numeric type](#) and has the value ANY (MATRIX\_A  $(:)$  [.AND.](#) MATRIX\_B  $(:, j)$ ) if the arguments are of logical type.

32 *Case (iii):* Element  $(i)$  of the result has the value SUM (MATRIX\_A  $(i, :)$  \* MATRIX\_B  $(:)$ ) if the arguments  
33 are of [numeric type](#) and has the value ANY (MATRIX\_A  $(i, :)$  [.AND.](#) MATRIX\_B  $(:)$ ) if the arguments are of logical type.

34 **Examples.** Let A and B be the matrices  $\begin{bmatrix} 1 & 2 & 3 \\ 2 & 3 & 4 \end{bmatrix}$  and  $\begin{bmatrix} 1 & 2 \\ 2 & 3 \\ 3 & 4 \end{bmatrix}$ ; let X and Y be the vectors  $[1, 2]$  and  
35  $[1, 2, 3]$ .



- 1 *Case (i):* The result of MATMUL (A, B) is the matrix-matrix product AB with the value  $\begin{bmatrix} 14 & 20 \\ 20 & 29 \end{bmatrix}$ .
- 2 *Case (ii):* The result of MATMUL (X, A) is the vector-matrix product XA with the value [5, 8, 11].
- 3 *Case (iii):* The result of MATMUL (A, Y) is the matrix-vector product AY with the value [14, 20].

#### 4 **16.9.125 MAX (A1, A2 [, A3, ...])**

- 5 1 **Description.** Maximum value.
- 6 2 **Class.** [Elemental function](#).
- 7 3 **Arguments.** The arguments shall all have the same type which shall be integer, real, or character and they shall  
8 all have the same kind type parameter.
- 9 4 **Result Characteristics.** The type and kind type parameter of the result are the same as those of the arguments.  
10 For arguments of character type, the length of the result is the length of the longest argument.
- 11 5 **Result Value.** The value of the result is that of the largest argument. For arguments of character type, the  
12 result is the value that would be selected by application of intrinsic relational operators; that is, the [collating](#)  
13 [sequence](#) for characters with the kind type parameter of the arguments is applied. If the selected argument is  
14 shorter than the longest argument, the result is extended with blanks on the right to the length of the longest  
15 argument.
- 16 6 **Examples.** MAX (-9.0, 7.0, 2.0) has the value 7.0, MAX ('Z', 'BB') has the value 'Z ', and MAX (['A', 'Z'],  
17 ['BB', 'Y ']) has the value ['BB', 'Z '].

#### 18 **16.9.126 MAXEXPONENT (X)**

- 19 1 **Description.** Maximum exponent of a real model.
- 20 2 **Class.** [Inquiry function](#).
- 21 3 **Argument.** X shall be of type real. It may be a scalar or an array.
- 22 4 **Result Characteristics.** Default integer scalar.
- 23 5 **Result Value.** The result has the value  $e_{\max}$ , as defined in [16.4](#) for the model representing numbers of the same  
24 type and kind type parameter as X.
- 25 6 **Example.** MAXEXPONENT (X) has the value 127 for real X whose model is as in Note [16.4](#).

#### 26 **16.9.127 MAXLOC (ARRAY, DIM [, MASK, KIND, BACK]) or MAXLOC (ARRAY [, MASK, KIND, BACK])**

- 27 1 **Description.** Location(s) of maximum value.
- 28 2 **Class.** [Transformational function](#).
- 29 3 **Arguments.**
- 30 ARRAY shall be an array of type integer, real, or character.
- 31 DIM shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the [rank](#) of ARRAY.
- 32 MASK (optional) shall be of type logical and shall be [conformable](#) with ARRAY.
- 33 KIND (optional) shall be a scalar integer [constant expression](#).
- 34 BACK (optional) shall be a logical scalar.
- 35 4 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of  
36 KIND; otherwise the kind type parameter is that of default integer type. If DIM does not appear, the result is

1 an array of **rank** one and of size equal to the **rank** of ARRAY; otherwise, the result is of **rank**  $n - 1$  and shape  
2  $[d_1, d_2, \dots, d_{\text{DIM}-1}, d_{\text{DIM}+1}, \dots, d_n]$ , where  $[d_1, d_2, \dots, d_n]$  is the shape of ARRAY.

### 3 5 Result Value.

4 *Case (i):* If DIM does not appear and MASK is absent, the result is a rank-one array whose element values  
5 are the values of the subscripts of an element of ARRAY whose value equals the maximum value of  
6 all of the elements of ARRAY. The  $i^{\text{th}}$  subscript returned lies in the range 1 to  $e_i$ , where  $e_i$  is the  
7 extent of the  $i^{\text{th}}$  dimension of ARRAY. If ARRAY has size zero, all elements of the result are zero.

8 *Case (ii):* If DIM does not appear and MASK is present, the result is a rank-one array whose element values  
9 are the values of the subscripts of an element of ARRAY, corresponding to a true element of MASK,  
10 whose value equals the maximum value of all such elements of ARRAY. The  $i^{\text{th}}$  subscript returned  
11 lies in the range 1 to  $e_i$ , where  $e_i$  is the extent of the  $i^{\text{th}}$  dimension of ARRAY. If ARRAY has size  
12 zero or every element of MASK has the value false, all elements of the result are zero.

13 *Case (iii):* If ARRAY has **rank** one and DIM is specified, the result has a value equal to that of the first element  
14 of MAXLOC (ARRAY [, MASK = MASK, KIND = KIND, BACK = BACK]). Otherwise, if DIM  
15 is specified, the value of element  $(s_1, s_2, \dots, s_{\text{DIM}-1}, s_{\text{DIM}+1}, \dots, s_n)$  of the result is equal to  
MAXLOC (ARRAY ( $s_1, s_2, \dots, s_{\text{DIM}-1}, \dots, s_{\text{DIM}+1}, \dots, s_n$ ),  
DIM = 1  
[, MASK = MASK ( $s_1, s_2, \dots, s_{\text{DIM}-1}, \dots, s_{\text{DIM}+1}, \dots, s_n$ ),  
KIND = KIND,  
BACK = BACK] ).

16 6 If only one element has the maximum value, that element's subscripts are returned. Otherwise, if more than  
17 one element has the maximum value and BACK is absent or present with the value false, the element whose  
18 subscripts are returned is the first such element, taken in array element order. If BACK is present with the value  
19 true, the element whose subscripts are returned is the last such element, taken in array element order.

20 7 If ARRAY has type character, the result is the value that would be selected by application of intrinsic relational  
21 operators; that is, the **collating sequence** for characters with the kind type parameter of the arguments is applied.

### 23 8 Examples.

24 *Case (i):* The value of MAXLOC ([2, 6, 4, 6]) is [2] and the value of MAXLOC ([2, 6, 4, 6], BACK=.TRUE.)  
25 is [4].

26 *Case (ii):* If A has the value  $\begin{bmatrix} 0 & -5 & 8 & -3 \\ 3 & 4 & -1 & 2 \\ 1 & 5 & 6 & -4 \end{bmatrix}$ , MAXLOC (A, MASK = A < 6) has the value [3, 2]. This  
27 is independent of the declared lower bounds for A.

28 *Case (iii):* The value of MAXLOC ([5, -9, 3], DIM = 1) is 1. If B has the value  $\begin{bmatrix} 1 & 3 & -9 \\ 2 & 2 & 6 \end{bmatrix}$ , MAXLOC  
29 (B, DIM = 1) is [2, 1, 2] and MAXLOC (B, DIM = 2) is [2, 3]. This is independent of the declared  
30 lower bounds for B.

## 31 16.9.128 MAXVAL (ARRAY, DIM [, MASK]) or MAXVAL (ARRAY [, MASK])

32 1 **Description.** Maximum value(s) of array.

33 2 **Class.** Transformational function.

### 34 3 Arguments.

35 ARRAY shall be an array of type integer, real, or character.

36 DIM shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the **rank** of ARRAY.

37 MASK (optional) shall be of type logical and shall be **conformable** with ARRAY.

38 4 **Result Characteristics.** The result is of the same type and type parameters as ARRAY. It is scalar if DIM  
39 does not appear; otherwise, the result has **rank**  $n - 1$  and shape  $[d_1, d_2, \dots, d_{\text{DIM}-1}, d_{\text{DIM}+1}, \dots, d_n]$  where  
40  $[d_1, d_2, \dots, d_n]$  is the shape of ARRAY.

1 5 **Result Value.**

2 *Case (i):* The result of MAXVAL (ARRAY) has a value equal to the maximum value of all the elements of  
 3 ARRAY if the size of ARRAY is not zero. If ARRAY has size zero and type integer or real, the  
 4 result has the value of the negative number of the largest magnitude supported by the processor  
 5 for numbers of the type and kind type parameter of ARRAY. If ARRAY has size zero and type  
 6 character, the result has the value of a string of characters of length LEN (ARRAY), with each  
 7 character equal to CHAR (0, KIND (ARRAY)).

8 *Case (ii):* The result of MAXVAL (ARRAY, MASK = MASK) has a value equal to that of MAXVAL (PACK  
 9 (ARRAY, MASK)).

10 *Case (iii):* The result of MAXVAL (ARRAY, DIM = DIM [,MASK = MASK]) has a value equal to that of  
 11 MAXVAL (ARRAY [,MASK = MASK]) if ARRAY has **rank** one. Otherwise, the value of element  
 12 ( $s_1, s_2, \dots, s_{\text{DIM}-1}, s_{\text{DIM}+1}, \dots, s_n$ ) of the result is equal to

$$13 \quad \text{MAXVAL (ARRAY (} s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n) \\ 14 \quad [ , \text{MASK = MASK (} s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n) ] ).$$

15 6 If ARRAY is of type character, the result is the value that would be selected by application of intrinsic relational  
 16 operators; that is, the **collating sequence** for characters with the kind type parameter of the arguments is applied.

17 7 **Examples.**

18 *Case (i):* The value of MAXVAL ([1, 2, 3]) is 3.

19 *Case (ii):* MAXVAL (C, MASK = C < 0.0) is the maximum of the negative elements of C.

20 *Case (iii):* If B is the array  $\begin{bmatrix} 1 & 3 & 5 \\ 2 & 7 & 6 \end{bmatrix}$ , MAXVAL (B, DIM = 1) is [2, 7, 6] and MAXVAL (B, DIM = 2) is  
 21 [5, 7].

22 **16.9.129 MERGE (TSOURCE, FSOURCE, MASK)**

23 1 **Description.** Expression value selection.

24 2 **Class.** **Elemental** function.

25 3 **Arguments.**

26 TSOURCE may be of any type.

27 FSOURCE shall be of the same type and type parameters as TSOURCE.

28 MASK shall be of type logical.

29 4 **Result Characteristics.** Same type and **type parameters** as TSOURCE. Because TSOURCE and FSOURCE  
 30 are required to have the same type and **type parameters** (for both the **declared** and **dynamic** types), the result is  
 31 **polymorphic** if and only if both TSOURCE and FSOURCE are **polymorphic**.

32 5 **Result Value.** The result is TSOURCE if MASK is true and FSOURCE otherwise.

33 6 **Examples.** If TSOURCE is the array  $\begin{bmatrix} 1 & 6 & 5 \\ 2 & 4 & 6 \end{bmatrix}$ , FSOURCE is the array  $\begin{bmatrix} 0 & 3 & 2 \\ 7 & 4 & 8 \end{bmatrix}$  and MASK is the  
 34 array  $\begin{bmatrix} \text{T} & \cdot & \text{T} \\ \cdot & \cdot & \text{T} \end{bmatrix}$ , where “T” represents true and “.” represents false, then MERGE (TSOURCE, FSOURCE,  
 35 MASK) is  $\begin{bmatrix} 1 & 3 & 5 \\ 7 & 4 & 6 \end{bmatrix}$ . The value of MERGE (1.0, 0.0, K > 0) is 1.0 for K = 5 and 0.0 for K = -2.

36 **16.9.130 MERGE\_BITS (I, J, MASK)**

37 1 **Description.** Merge of bits under mask.

38 2 **Class.** **Elemental** function.

1 3 **Arguments.**2 I shall be of type integer or a *boz-literal-constant*.3 J shall be of type integer or a *boz-literal-constant*. If both I and J are of type integer they shall have  
4 the same kind type parameter. I and J shall not both be *boz-literal-constants*.5 MASK shall be of type integer or a *boz-literal-constant*. If MASK is of type integer, it shall have the same  
6 kind type parameter as each other argument of type integer.7 4 **Result Characteristics.** Same as I if I is of type integer; otherwise, same as J.8 5 **Result Value.** If any argument is a *boz-literal-constant*, it is first converted as if by the intrinsic function  
9 **INT** to the type and kind type parameter of the result. The result has the value of IOR (IAND (I, MASK),  
10 IAND (J, NOT (MASK))).11 6 **Example.** MERGE\_BITS (13, 18, 22) has the value 4.12 **16.9.131 MIN (A1, A2 [, A3, ...])**13 1 **Description.** Minimum value.14 2 **Class.** **Elemental** function.15 3 **Arguments.** The arguments shall all be of the same type which shall be integer, real, or character and they  
16 shall all have the same kind type parameter.17 4 **Result Characteristics.** The type and kind type parameter of the result are the same as those of the arguments.  
18 For arguments of character type, the length of the result is the length of the longest argument.19 5 **Result Value.** The value of the result is that of the smallest argument. For arguments of character type, the  
20 result is the value that would be selected by application of intrinsic relational operators; that is, the **collating**  
21 **sequence** for characters with the kind type parameter of the arguments is applied. If the selected argument is  
22 shorter than the longest argument, the result is extended with blanks on the right to the length of the longest  
23 argument.24 6 **Examples.** MIN (-9.0, 7.0, 2.0) has the value -9.0, MIN ('A', 'YY') has the value 'A ', and  
25 MIN (['Z', 'A'], ['YY', 'B ']) has the value ['YY', 'A '].26 **16.9.132 MINEXPONENT (X)**27 1 **Description.** Minimum exponent of a real model.28 2 **Class.** **Inquiry** function.29 3 **Argument.** X shall be of type real. It may be a scalar or an array.30 4 **Result Characteristics.** Default integer scalar.31 5 **Result Value.** The result has the value  $e_{\min}$ , as defined in 16.4 for the model representing numbers of the same  
32 type and kind type parameter as X.33 6 **Example.** MINEXPONENT (X) has the value -126 for real X whose model is as in Note 16.4.34 **16.9.133 MINLOC (ARRAY, DIM [, MASK, KIND, BACK]) or  
MINLOC (ARRAY [, MASK, KIND, BACK])**35 1 **Description.** Location(s) of minimum value.36 2 **Class.** **Transformational** function.

1 3 **Arguments.**

- 2 ARRAY shall be an array of type integer, real, or character.  
 3 DIM shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the **rank** of ARRAY.  
 4 MASK (optional) shall be of type logical and shall be **conformable** with ARRAY.  
 5 KIND (optional) shall be a scalar integer **constant expression**.  
 6 BACK (optional) shall be a logical scalar.

7 4 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of  
 8 KIND; otherwise the kind type parameter is that of default integer type. If DIM does not appear, the result is  
 9 an array of **rank** one and of size equal to the **rank** of ARRAY; otherwise, the result is of **rank**  $n - 1$  and shape  
 10  $[d_1, d_2, \dots, d_{\text{DIM}-1}, d_{\text{DIM}+1}, \dots, d_n]$ , where  $[d_1, d_2, \dots, d_n]$  is the shape of ARRAY.

11 5 **Result Value.**

12 *Case (i):* If DIM does not appear and MASK is absent the result is a rank-one array whose element values  
 13 are the values of the subscripts of an element of ARRAY whose value equals the minimum value  
 14 of all the elements of ARRAY. The  $i^{\text{th}}$  subscript returned lies in the range 1 to  $e_i$ , where  $e_i$  is the  
 15 extent of the  $i^{\text{th}}$  dimension of ARRAY. If ARRAY has size zero, all elements of the result are zero.

16 *Case (ii):* If DIM does not appear and MASK is present, the result is a rank-one array whose element values  
 17 are the values of the subscripts of an element of ARRAY, corresponding to a true element of MASK,  
 18 whose value equals the minimum value of all such elements of ARRAY. The  $i^{\text{th}}$  subscript returned  
 19 lies in the range 1 to  $e_i$ , where  $e_i$  is the extent of the  $i^{\text{th}}$  dimension of ARRAY. If ARRAY has size  
 20 zero or every element of MASK has the value false, all elements of the result are zero.

21 *Case (iii):* If ARRAY has **rank** one and DIM is specified, the result has a value equal to that of the first element  
 22 of MINLOC (ARRAY [, MASK = MASK, KIND = KIND, BACK = BACK]). Otherwise, if DIM  
 23 is specified, the value of element  $(s_1, s_2, \dots, s_{\text{DIM}-1}, s_{\text{DIM}+1}, \dots, s_n)$  of the result is equal to

24 
$$\begin{aligned} & \text{MINLOC (ARRAY (} s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n), \\ & \quad \text{DIM = 1} \\ & \quad [, \text{MASK = MASK (} s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n), \\ & \quad \text{KIND = KIND,} \\ & \quad \text{BACK = BACK] ).} \end{aligned}$$

25 6 If only one element has the minimum value, that element's subscripts are returned. Otherwise, if more than one  
 26 element has the minimum value and BACK is absent or present with the value false, the element whose subscripts  
 27 are returned is the first such element, taken in array element order. If BACK is present with the value true, the  
 28 element whose subscripts are returned is the last such element, taken in array element order.

29 7 If ARRAY is of type character, the result is the value that would be selected by application of intrinsic relational  
 30 operators; that is, the **collating sequence** for characters with the kind type parameter of the arguments is applied.

31 8 **Examples.**

32 *Case (i):* The value of MINLOC ([4, 3, 6, 3]) is [2] and the value of MINLOC ([4, 3, 6, 3], BACK = **.TRUE.**)  
 33 is [4].

34 *Case (ii):* If A has the value  $\begin{bmatrix} 0 & -5 & 8 & -3 \\ 3 & 4 & -1 & 2 \\ 1 & 5 & 6 & -4 \end{bmatrix}$ , MINLOC (A, MASK = A > -4) has the value [1, 4].  
 35 This is independent of the declared lower bounds for A.

36 *Case (iii):* The value of MINLOC ([5, -9, 3], DIM = 1) is 2. If B has the value  $\begin{bmatrix} 1 & 3 & -9 \\ 2 & 2 & 6 \end{bmatrix}$ , MIN-  
 37 LOC (B, DIM = 1) is [1, 2, 1] and MINLOC (B, DIM = 2) is [3, 1]. This is independent of  
 38 the declared lower bounds for B.

### 16.9.134 MINVAL (ARRAY, DIM [, MASK]) or MINVAL (ARRAY [, MASK])

1 **Description.** Minimum value(s) of array.

2 **Class.** Transformational function.

3 **Arguments.**

4 ARRAY shall be an array of type integer, real, or character.

5 DIM shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the rank of ARRAY.

6 MASK (optional) shall be of type logical and shall be conformable with ARRAY.

7 **Result Characteristics.** The result is of the same type and type parameters as ARRAY. It is scalar if DIM does not appear; otherwise, the result has rank  $n - 1$  and shape  $[d_1, d_2, \dots, d_{\text{DIM}-1}, d_{\text{DIM}+1}, \dots, d_n]$  where  $[d_1, d_2, \dots, d_n]$  is the shape of ARRAY.

8 **Result Value.**

9 *Case (i):* The result of MINVAL (ARRAY) has a value equal to the minimum value of all the elements of ARRAY if the size of ARRAY is not zero. If ARRAY has size zero and type integer or real, the result has the value of the positive number of the largest magnitude supported by the processor for numbers of the type and kind type parameter of ARRAY. If ARRAY has size zero and type character, the result has the value of a string of characters of length LEN (ARRAY), with each character equal to CHAR ( $n - 1$ , KIND (ARRAY)), where  $n$  is the number of characters in the collating sequence for characters with the kind type parameter of ARRAY.

10 *Case (ii):* The result of MINVAL (ARRAY, MASK = MASK) has a value equal to that of MINVAL (PACK (ARRAY, MASK)).

11 *Case (iii):* The result of MINVAL (ARRAY, DIM = DIM [, MASK = MASK]) has a value equal to that of MINVAL (ARRAY [, MASK = MASK]) if ARRAY has rank one. Otherwise, the value of element  $(s_1, s_2, \dots, s_{\text{DIM}-1}, s_{\text{DIM}+1}, \dots, s_n)$  of the result is equal to

$$\text{MINVAL} (\text{ARRAY} (s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n) \\ [, \text{MASK} = \text{MASK} (s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n) ] ).$$

12 If ARRAY is of type character, the result is the value that would be selected by application of intrinsic relational operators; that is, the collating sequence for characters with the kind type parameter of the arguments is applied.

13 **Examples.**

14 *Case (i):* The value of MINVAL ([1, 2, 3]) is 1.

15 *Case (ii):* MINVAL (C, MASK = C > 0.0) is the minimum of the positive elements of C.

16 *Case (iii):* If B is the array  $\begin{bmatrix} 1 & 3 & 5 \\ 2 & 4 & 6 \end{bmatrix}$ , MINVAL (B, DIM = 1) is [1, 3, 5] and MINVAL (B, DIM = 2) is [1, 2].

### 16.9.135 MOD (A, P)

17 **Description.** Remainder function.

18 **Class.** Elemental function.

19 **Arguments.**

20 A shall be of type integer or real.

21 P shall be of the same type and kind type parameter as A. P shall not be zero.

22 **Result Characteristics.** Same as A.

23 **Result Value.** The value of the result is  $A - \text{INT} (A/P) * P$ .

1 6 **Examples.** MOD (3.0, 2.0) has the value 1.0 (approximately). MOD (8, 5) has the value 3. MOD (-8, 5) has  
 2 the value -3. MOD (8, -5) has the value 3. MOD (-8, -5) has the value -3.

### 3 16.9.136 MODULO (A, P)

4 1 **Description.** Modulo function.

5 2 **Class.** [Elemental](#) function.

6 3 **Arguments.**

7 A shall be of type integer or real.

8 P shall be of the same type and kind type parameter as A. P shall not be zero.

9 4 **Result Characteristics.** Same as A.

10 5 **Result Value.**

11 *Case (i):* A is of type integer. MODULO (A, P) has the value R such that  $A = Q \times P + R$ , where Q is an  
 12 integer, the inequalities  $0 \leq R < P$  hold if  $P > 0$ , and  $P < R \leq 0$  hold if  $P < 0$ .

13 *Case (ii):* A is of type real. The value of the result is  $A - \text{FLOOR}(A / P) * P$ .

14 6 **Examples.** MODULO (8, 5) has the value 3. MODULO (-8, 5) has the value 2. MODULO (8, -5) has the  
 15 value -2. MODULO (-8, -5) has the value -3.

### 16 16.9.137 MOVE\_ALLOC (FROM, TO [, STAT, ERRMSG])

17 1 **Description.** Move an allocation.

18 2 **Class.** Subroutine, [pure](#) if and only if FROM is not a [coarray](#).

19 3 **Arguments.**

20 FROM may be of any type, [rank](#), and [corank](#). It shall be [allocatable](#) and shall not be a [coindexed object](#).  
 21 It is an [INTENT \(INOUT\)](#) argument.

22 TO shall be [type compatible \(7.3.2.3\)](#) with FROM and have the same [rank](#) and [corank](#). It shall be  
 23 [allocatable](#) and shall not be a [coindexed object](#). It shall be [polymorphic](#) if FROM is [polymorphic](#).  
 24 It is an [INTENT \(OUT\)](#) argument. Each nondeferred parameter of the [declared type](#) of TO shall  
 25 have the same value as the corresponding parameter of the [declared type](#) of FROM.

26 STAT (optional) shall be a noncoindexed default integer scalar. It is an [INTENT \(OUT\)](#) argument.

27 ERRMSG (optional) shall be a noncoindexed default character scalar. It is an [INTENT \(INOUT\)](#) argument.

28 4 If execution of MOVE\_ALLOC is successful, or if [STAT\\_FAILED\\_IMAGE](#) is assigned to STAT,

- 29 • On invocation of MOVE\_ALLOC, if the allocation status of TO is allocated, it is deallocated. Then,  
 30 if FROM has an allocation status of allocated on entry to MOVE\_ALLOC, TO becomes allocated with  
 31 [dynamic type](#), [type parameters](#), [bounds](#), [cobounds](#), and value identical to those that FROM had on entry  
 32 to MOVE\_ALLOC. Note that if FROM and TO are the same variable, it shall be unallocated when  
 33 MOVE\_ALLOC is invoked.
- 34 • If TO has the [TARGET attribute](#), any pointer associated with FROM on entry to MOVE\_ALLOC becomes  
 35 correspondingly associated with TO. If TO does not have the [TARGET attribute](#), the [pointer association](#)  
 36 status of any pointer associated with FROM on entry becomes undefined.
- 37 • The allocation status of FROM becomes unallocated.

38 5 When a reference to MOVE\_ALLOC is executed for which the FROM argument is a [coarray](#), there is an implicit  
 39 synchronization of all [active images](#) of the [current team](#). On those [images](#), execution of the segment (11.6.2)  
 40 following the [CALL statement](#) is delayed until all other [active images](#) of the [current team](#) have executed the same  
 41 statement the same number of times. When such a reference is executed, if any [image](#) of the [current team](#) has  
 42 [stopped](#) or [failed](#), an error condition occurs.

- 1 6 If STAT is present and execution is successful, it is assigned the value zero.
- 2 7 If an error condition occurs,
- 3 • if STAT is absent, **error termination** is initiated;
  - 4 • otherwise, if FROM is a **coarray** and the **current team** contains a **stopped image**, STAT is assigned the value **STAT\_STOPPED\_IMAGE** from the intrinsic module **ISO\_FORTRAN\_ENV**;
  - 5 • otherwise, if FROM is a **coarray** and the **current team** contains a **failed image**, and no other error condition
  - 6 occurs, STAT is assigned the value **STAT\_FAILED\_IMAGE** from the intrinsic module **ISO\_FORTRAN\_**
  - 7 **ENV**;
  - 8 • otherwise, STAT is assigned a processor-dependent positive value that differs from that of **STAT\_STOPPED\_**
  - 9 **IMAGE** or **STAT\_FAILED\_IMAGE**.
- 11 8 If the ERRMSG argument is present and an error condition occurs, it is assigned an explanatory message. If no
- 12 error condition occurs, the definition status and value of ERRMSG are unchanged.

- 13 9 **Example.** The example below demonstrates reallocation of GRID to twice its previous size, with its previous
- 14 contents evenly distributed over the new elements so that intermediate points can be inserted.

```

15     REAL,ALLOCATABLE :: GRID(:),TEMPGRID(:)
16     ...
17     ALLOCATE(GRID(-N:N))      ! initial allocation of GRID
18     ...
19     ALLOCATE(TEMPGRID(-2*N:2*N)) ! allocate bigger grid
20     TEMPGRID(::2)=GRID ! distribute values to new locations
21     CALL MOVE_ALLOC(TO=GRID,FROM=TEMPGRID)

```

22 The old grid is deallocated because TO is **INTENT (OUT)**, and GRID then takes over the new grid allocation.

#### NOTE 16.22

It is expected that the implementation of **allocatable** objects will typically involve descriptors to locate the allocated storage; MOVE\_ALLOC could then be implemented by transferring the contents of the descriptor for FROM to the descriptor for TO and clearing the descriptor for FROM.

### 23 16.9.138 MVBITS (FROM, FROMPOS, LEN, TO, TOPOS)

24 1 **Description.** Copy a sequence of bits.

25 2 **Class.** **Elemental** subroutine.

26 3 **Arguments.**

27 FROM shall be of type integer. It is an **INTENT (IN)** argument.

28 FROMPOS shall be of type integer and nonnegative. It is an **INTENT (IN)** argument. FROMPOS + LEN

29 shall be less than or equal to **BIT\_SIZE** (FROM). The model for the interpretation of an integer

30 value as a sequence of bits is in 16.3.

31 LEN shall be of type integer and nonnegative. It is an **INTENT (IN)** argument.

32 TO shall be a variable of the same type and kind type parameter value as FROM and may be associated

33 with FROM (15.8.3). It is an **INTENT (INOUT)** argument. TO is defined by copying the sequence

34 of bits of length LEN, starting at position FROMPOS of FROM to position TOPOS of TO. No

35 other bits of TO are altered. On return, the LEN bits of TO starting at TOPOS are equal to

36 the value that the LEN bits of FROM starting at FROMPOS had on entry. The model for the

37 interpretation of an integer value as a sequence of bits is in 16.3.

38 TOPOS shall be of type integer and nonnegative. It is an **INTENT (IN)** argument. TOPOS + LEN shall

39 be less than or equal to **BIT\_SIZE** (TO).



1 4 **Example.** If TO has the initial value 6, its value after the statement CALL MVBITS (7, 2, 2, TO, 0) is 5.

## 2 16.9.139 NEAREST (X, S)

3 1 **Description.** Adjacent machine number.

4 2 **Class.** [Elemental](#) function.

5 3 **Arguments.**

6 X shall be of type real.

7 S shall be of type real and not equal to zero.

8 4 **Result Characteristics.** Same as X.

9 5 **Result Value.** The result has a value equal to the machine-representable number distinct from X and nearest  
10 to it in the direction of the  $\infty$  with the same sign as S.

11 6 **Example.** NEAREST (3.0, 2.0) has the value  $3 + 2^{-22}$  on a machine whose representation is that of the model  
12 in Note 16.4.

### NOTE 16.23

Unlike other floating-point manipulation functions, NEAREST operates on machine-representable numbers rather than model numbers. On many systems there are machine-representable numbers that lie between adjacent model numbers.

## 13 16.9.140 NEW\_LINE (A)

14 1 **Description.** Newline character.

15 2 **Class.** [Inquiry function](#).

16 3 **Argument.** A shall be of type character. It may be a scalar or an array.

17 4 **Result Characteristics.** Character scalar of length one with the same kind type parameter as A.

18 5 **Result Value.**

19 *Case (i):* If A is default character and the character in position 10 of the [ASCII collating sequence](#) is repres-  
20 entable in the default character set, then the result is [ACHAR](#) (10).

21 *Case (ii):* If A is [ASCII character](#) or [ISO 10646 character](#), then the result is [CHAR](#) (10, [KIND](#) (A)).

22 *Case (iii):* Otherwise, the result is a processor-dependent character that represents a newline in output to files  
23 connected for formatted stream output if there is such a character.

24 *Case (iv):* Otherwise, the result is the blank character.

25 6 **Example.** If there is a suitable newline character, and unit 10 is connected for formatted stream output, the  
26 statement

```
27 WRITE (10, '(A)') 'New'//NEW_LINE('a')//'Line'
```

28 will write a record containing “New” and then a record containing “Line”.

## 29 16.9.141 NINT (A [, KIND])

30 1 **Description.** Nearest integer.

31 2 **Class.** [Elemental](#) function.

1 3 **Arguments.**

2 A shall be of type real.

3 KIND (optional) shall be a scalar integer [constant expression](#).4 4 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of  
5 KIND; otherwise, the kind type parameter is that of default integer type.6 5 **Result Value.** The result is the integer nearest A, or if there are two integers equally near A, the result is  
7 whichever such integer has the greater magnitude.8 6 **Example.** NINT (2.783) has the value 3.9 **16.9.142 NORM2 (X) or NORM2 (X, DIM)**10 1 **Description.**  $L_2$  norm of an array.11 2 **Class.** [Transformational function](#).12 3 **Arguments.**

13 X shall be a real array.

14 DIM shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the [rank](#) of X.15 4 **Result Characteristics.** The result is of the same type and type parameters as X. It is scalar if DIM does not  
16 appear; otherwise the result has [rank](#)  $n - 1$  and shape  $[d_1, d_2, \dots, d_{\text{DIM}-1}, d_{\text{DIM}+1}, \dots, d_n]$ , where  $n$  is the [rank](#)  
17 of X and  $[d_1, d_2, \dots, d_n]$  is the shape of X.18 5 **Result Value.**19 *Case (i):* The result of NORM2 (X) has a value equal to a processor-dependent approximation to the gener-  
20 alized  $L_2$  norm of X, which is the square root of the sum of the squares of the elements of X. If X  
21 has size zero, the result has the value zero.22 *Case (ii):* The result of NORM2 (X, DIM=DIM) has a value equal to that of NORM2 (X) if X has [rank](#)  
23 one. Otherwise, the value of element  $(s_1, s_2, \dots, s_{\text{DIM}-1}, s_{\text{DIM}+1}, \dots, s_n)$  of the result is equal to  
24  $\text{NORM2}(X(s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n))$ .

25 6 It is recommended that the processor compute the result without undue overflow or underflow.

26 7 **Example.** The value of NORM2 ([3.0, 4.0]) is 5.0 (approximately). If X has the value  $\begin{bmatrix} 1.0 & 2.0 \\ 3.0 & 4.0 \end{bmatrix}$  then the  
27 value of NORM2 (X, DIM=1) is [3.162, 4.472] (approximately) and the value of NORM2 (X, DIM=2) is [2.236,  
28 5.0] (approximately).29 **16.9.143 NOT (I)**30 1 **Description.** Bitwise complement.31 2 **Class.** [Elemental function](#).32 3 **Argument.** I shall be of type integer.33 4 **Result Characteristics.** Same as I.34 5 **Result Value.** The result has the value obtained by complementing I bit-by-bit according to the following table:

I	NOT (I)
1	0
0	1

- 1 6 The model for the interpretation of an integer value as a sequence of bits is in 16.3.
- 2 7 **Example.** If I is represented by the string of bits 01010101, NOT (I) has the binary value 10101010.
- 3 **16.9.144 NULL ([MOLD])**
- 4 1 **Description.** *Disassociated* pointer or unallocated *allocatable* entity.
- 5 2 **Class.** *Transformational function*.
- 6 3 **Argument.** MOLD shall be a pointer or *allocatable*. It may be of any type or may be a procedure pointer. If MOLD is a pointer its *pointer association* status may be undefined, *disassociated*, or associated. If MOLD is *allocatable* its allocation status may be allocated or unallocated. It need not be defined with a value.
- 7 4 **Result Characteristics.** If MOLD is present, the *characteristics* are the same as MOLD. If MOLD has *deferred type parameters*, those type parameters of the result are *deferred*.
- 8 5 If MOLD is absent, the *characteristics* of the result are determined by the entity with which the reference is associated. See Table 16.5. MOLD shall not be absent in any other context. If any type parameters of the contextual entity are *deferred*, those type parameters of the result are *deferred*. If any type parameters of the contextual entity are assumed, MOLD shall be present.
- 9 6 If the context of the reference to NULL is an *actual argument* in a generic procedure reference, MOLD shall be present if the type, type parameters, or *rank* are required to resolve the generic reference.

Table 16.5: Characteristics of the result of NULL ( )

Appearance of NULL ( )	Type, type parameters, and rank of result:
right side of a <i>pointer assignment</i>	pointer on the left side
initialization for an object in a declaration	the object
default initialization for a component	the component
in a <i>structure constructor</i>	the corresponding component
as an <i>actual argument</i>	the corresponding dummy argument
in a <i>DATA statement</i>	the corresponding pointer object

- 17 7 **Result.** The result is a *disassociated* pointer or an unallocated *allocatable* entity.
- 18 8 **Examples.**
- 19 *Case (i):* REAL, POINTER, DIMENSION (:) :: VEC => NULL ( ) defines the initial association status of  
20 VEC to be *disassociated*.
- 21 *Case (ii):* The MOLD argument is required in the following:
- ```

22     INTERFACE GEN
23         SUBROUTINE S1 (J, PI)
24             INTEGER J
25             INTEGER, POINTER :: PI
26         END SUBROUTINE S1
27         SUBROUTINE S2 (K, PR)
28             INTEGER K
29             REAL, POINTER :: PR
30         END SUBROUTINE S2
31     END INTERFACE
32     REAL, POINTER :: REAL_PTR
33     CALL GEN (7, NULL (REAL_PTR) )      ! Invokes S2

```

### 16.9.145 NUM\_IMAGES ( ) or NUM\_IMAGES (TEAM) or NUM\_IMAGES (TEAM\_NUMBER)

1 **Description.** Number of *images*.

2 **Class.** *Transformational function*.

3 **Arguments.**

4 TEAM shall be a scalar of type *TEAM\_TYPE* from the intrinsic module *ISO\_FORTRAN\_ENV*, with a  
5 value that identifies the *current* or an *ancestor team*.

6 TEAM\_NUMBER shall be an integer scalar. It shall identify the *initial team* or a *team* whose *parent* is the  
7 same as that of the *current team*.

8 **Result Characteristics.** Default integer scalar.

9 **Result Value.** The number of *images* in the specified *team*, or in the *current team* if no team is specified.

10 **Example.** The following code uses *image 1* to read data and broadcast it to other *images*.

```
11      REAL :: P[*]
12      IF (THIS_IMAGE()==1) THEN
13          READ (6,*) P
14          DO I = 2, NUM_IMAGES()
15              P[I] = P
16          END DO
17      END IF
18      SYNC ALL
```

### 16.9.146 OUT\_OF\_RANGE (X, MOLD [, ROUND])

1 **Description.** Whether a value cannot be converted safely.

2 **Class.** *Elemental function*.

3 **Arguments.**

4 X shall be of type integer or real.

5 MOLD shall be an integer or real scalar. If it is a variable, it need not be defined.

6 ROUND (optional) shall be a logical scalar. ROUND shall be present only if X is of type real and MOLD is of  
7 type integer.

8 **Result Characteristics.** Default logical.

9 **Result Value.**

10 *Case (i):* If MOLD is of type integer, and ROUND is absent or present with the value false, the result is true  
11 if and only if the value of X is an IEEE infinity or NaN, or if the integer with largest magnitude  
12 that lies between zero and X inclusive is not representable by objects with the type and kind of  
13 MOLD.

14 *Case (ii):* If MOLD is of type integer, and ROUND is present with the value true, the result is true if and only  
15 if the value of X is an IEEE infinity or NaN, or if the integer nearest X, or the integer of greater  
16 magnitude if two integers are equally near to X, is not representable by objects with the type and  
17 kind of MOLD.

18 *Case (iii):* Otherwise, the result is true if and only if the value of X is an IEEE infinity or NaN that is not  
19 supported by objects of the type and kind of MOLD, or if X is a finite number and the result of

1 rounding the value of X (according to the IEEE rounding mode if appropriate) to the extended  
 2 model for the kind of MOLD has magnitude larger than that of the largest finite number with the  
 3 same sign as X that is representable by objects with the type and kind of MOLD.

- 4 6 **Examples.** If INT8 is the kind value for an 8-bit binary integer type, OUT\_OF\_RANGE (−128.5, 0\_INT8)  
 5 will have the value false and OUT\_OF\_RANGE (−128.5, 0\_INT8, .TRUE.) will have the value true.

**NOTE 16.24**

MOLD is required to be a scalar because the only information taken from it is its type and kind. Allowing an array MOLD would require that it be conformable with X. ROUND is scalar because allowing an array rounding mode would have severe performance difficulties on many processors.

## 6 16.9.147 PACK (ARRAY, MASK [, VECTOR])

7 1 **Description.** Array packed into a vector.

8 2 **Class.** Transformational function.

9 3 **Arguments.**

10 ARRAY shall be an array of any type.

11 MASK shall be of type logical and shall be conformable with ARRAY.

12 VECTOR (optional) shall be of the same type and type parameters as ARRAY and shall have rank one. VEC-  
 13 TOR shall have at least as many elements as there are true elements in MASK. If MASK is scalar  
 14 with the value true, VECTOR shall have at least as many elements as there are in ARRAY.

15 4 **Result Characteristics.** The result is an array of rank one with the same type and type parameters as  
 16 ARRAY. If VECTOR is present, the result size is that of VECTOR; otherwise, the result size is the number  $t$   
 17 of true elements in MASK unless MASK is scalar with the value true, in which case the result size is the size of  
 18 ARRAY.

19 5 **Result Value.** Element  $i$  of the result is the element of ARRAY that corresponds to the  $i^{th}$  true element of  
 20 MASK, taking elements in array element order, for  $i = 1, 2, \dots, t$ . If VECTOR is present and has size  $n > t$ ,  
 21 element  $i$  of the result has the value VECTOR ( $i$ ), for  $i = t + 1, \dots, n$ .

22 6 **Examples.** The nonzero elements of an array M with the value  $\begin{bmatrix} 0 & 0 & 0 \\ 9 & 0 & 0 \\ 0 & 0 & 7 \end{bmatrix}$  can be “gathered” by the func-  
 23 tion PACK. The result of PACK (M, MASK = M /= 0) is [9, 7] and the result of PACK (M, M /= 0, VEC-  
 24 TOR = [2, 4, 6, 8, 10, 12]) is [9, 7, 6, 8, 10, 12].

## 25 16.9.148 PARITY (MASK) or PARITY (MASK, DIM)

26 1 **Description.** Array reduced by .NEQV. operation.

27 2 **Class.** Transformational function.

28 3 **Arguments.**

29 MASK shall be a logical array.

30 DIM shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the rank of MASK.

31 4 **Result Characteristics.** The result is of type logical with the same kind type parameter as MASK. It is scalar  
 32 if DIM does not appear; otherwise, the result has rank  $n - 1$  and shape  $[d_1, d_2, \dots, d_{\text{DIM}-1}, d_{\text{DIM}+1}, \dots, d_n]$   
 33 where  $[d_1, d_2, \dots, d_n]$  is the shape of MASK.

34 5 **Result Value.**

1 *Case (i):* The result of PARITY (MASK) has the value true if an odd number of the elements of MASK are  
2 true, and false otherwise.

3 *Case (ii):* If MASK has [rank](#) one, PARITY (MASK, DIM) is equal to PARITY (MASK). Otherwise, the  
4 value of element  $(s_1, s_2, \dots, s_{DIM-1}, s_{DIM+1}, \dots, s_n)$  of PARITY (MASK, DIM) is equal to  
5 PARITY (MASK  $(s_1, s_2, \dots, s_{DIM-1}, :, s_{DIM+1}, \dots, s_n)$ ).

## 6 6 Examples.

7 *Case (i):* The value of PARITY ([T, T, T, F]) is true if T has the value true and F has the value false.

8 *Case (ii):* If B is the array  $\begin{bmatrix} T & T & F \\ T & T & T \end{bmatrix}$ , where T has the value true and F has the value false, then  
9 PARITY (B, DIM=1) has the value [F, F, T] and PARITY (B, DIM=2) has the value [F, T].

## 10 16.9.149 POPCNT (I)

11 1 **Description.** Number of one bits.

12 2 **Class.** [Elemental](#) function.

13 3 **Argument.** I shall be of type integer.

14 4 **Result Characteristics.** Default integer.

15 5 **Result Value.** The result value is equal to the number of one bits in the sequence of bits of I. The model for  
16 the interpretation of an integer value as a sequence of bits is in [16.3](#).

17 6 **Examples.** POPCNT ([1, 2, 3, 4, 5, 6]) has the value [1, 1, 2, 1, 2, 2].

## 18 16.9.150 POPPAR (I)

19 1 **Description.** Parity expressed as 0 or 1.

20 2 **Class.** [Elemental](#) function.

21 3 **Argument.** I shall be of type integer.

22 4 **Result Characteristics.** Default integer.

23 5 **Result Value.** POPPAR (I) has the value 1 if POPCNT (I) is odd, and 0 if POPCNT (I) is even.

24 6 **Examples.** POPPAR ([1, 2, 3, 4, 5, 6]) has the value [1, 1, 0, 1, 0, 0].

## 25 16.9.151 PRECISION (X)

26 1 **Description.** Decimal precision of a real model.

27 2 **Class.** [Inquiry](#) function.

28 3 **Argument.** X shall be of type real or complex. It may be a scalar or an array.

29 4 **Result Characteristics.** Default integer scalar.

30 5 **Result Value.** The result has the value  $\text{INT}((p-1) * \text{LOG}_{10}(b)) + k$ , where  $b$  and  $p$  are as defined in [16.4](#)  
31 for the model representing real numbers with the same value for the kind type parameter as X, and where  $k$  is 1  
32 if  $b$  is an integral power of 10 and 0 otherwise.

33 6 **Example.** PRECISION (X) has the value  $\text{INT}(23 * \text{LOG}_{10}(2.)) = \text{INT}(6.92\dots) = 6$  for real X whose model  
34 is as in Note [16.4](#).

## 1 16.9.152 PRESENT (A)

2 1 **Description.** Presence of optional argument.

3 2 **Class.** [Inquiry function](#).

4 3 **Argument.** A shall be the name of an optional dummy argument that is accessible in the subprogram in which  
5 the PRESENT function reference appears. There are no other requirements on A.

6 4 **Result Characteristics.** Default logical scalar.

7 5 **Result Value.** The result has the value true if A is present ([15.5.2.12](#)) and otherwise has the value false.

## 8 16.9.153 PRODUCT (ARRAY, DIM [, MASK]) or 9 PRODUCT (ARRAY [, MASK])

10 1 **Description.** Array reduced by multiplication.

11 2 **Class.** [Transformational function](#).

12 3 **Arguments.**

13 ARRAY shall be an array of [numeric type](#).

14 DIM shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the [rank](#) of ARRAY.

15 MASK (optional) shall be of type logical and shall be [conformable](#) with ARRAY.

16 4 **Result Characteristics.** The result is of the same type and kind type parameter as ARRAY. It is scalar if  
17 DIM does not appear; otherwise, the result has [rank](#)  $n - 1$  and shape  $[d_1, d_2, \dots, d_{\text{DIM}-1}, d_{\text{DIM}+1}, \dots, d_n]$  where  
18  $[d_1, d_2, \dots, d_n]$  is the shape of ARRAY.

19 5 **Result Value.**

20 *Case (i):* The result of PRODUCT (ARRAY) has a value equal to a processor-dependent approximation to  
21 the product of all the elements of ARRAY or has the value one if ARRAY has size zero.

22 *Case (ii):* The result of PRODUCT (ARRAY, MASK = MASK) has a value equal to a processor-dependent  
23 approximation to the product of the elements of ARRAY corresponding to the true elements of  
24 MASK or has the value one if there are no true elements.

25 *Case (iii):* If ARRAY has [rank](#) one, PRODUCT (ARRAY, DIM = DIM [, MASK = MASK]) has a value equal  
26 to that of PRODUCT (ARRAY [, MASK = MASK]). Otherwise, the value of element  $(s_1, s_2, \dots,$   
27  $s_{\text{DIM}-1}, s_{\text{DIM}+1}, \dots, s_n)$  of PRODUCT (ARRAY, DIM = DIM [, MASK = MASK]) is equal to

28  $\text{PRODUCT}(\text{ARRAY}(s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n) [, \text{MASK} = \text{MASK}(s_1, s_2, \dots,$   
 $s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n) ])$ .

29 6 **Examples.**

30 *Case (i):* The value of PRODUCT ([1, 2, 3]) is 6.

31 *Case (ii):* PRODUCT (C, MASK = C > 0.0) forms the product of the positive elements of C.

32 *Case (iii):* If B is the array  $\begin{bmatrix} 1 & 3 & 5 \\ 2 & 4 & 6 \end{bmatrix}$ , PRODUCT (B, DIM = 1) is [2, 12, 30] and PRODUCT (B, DIM = 2)  
33 is [15, 48].

## 34 16.9.154 RADIX (X)

35 1 **Description.** Base of a numeric model.

36 2 **Class.** [Inquiry function](#).

37 3 **Argument.** X shall be of type integer or real. It may be a scalar or an array.

1 4 **Result Characteristics.** Default integer scalar.

2 5 **Result Value.** The result has the value  $r$  if  $X$  is of type integer and the value  $b$  if  $X$  is of type real, where  $r$  and  
3  $b$  are as defined in 16.4 for the model representing numbers of the same type and kind type parameter as  $X$ .

4 6 **Example.** RADIX ( $X$ ) has the value 2 for real  $X$  whose model is as in Note 16.4.

### 5 16.9.155 RANDOM\_INIT (REPEATABLE, IMAGE\_DISTINCT)

6 1 **Description.** Initialize the pseudorandom number generator.

7 2 **Class.** Subroutine.

8 3 **Arguments.**

9 REPEATABLE shall be a logical scalar. It is an **INTENT (IN)** argument. If it has the value true, the seed  
10 accessed by the pseudorandom number generator is set to a processor-dependent value that is the  
11 same each time RANDOM\_INIT is called from the same **image**. If it has the value false, the seed  
12 is set to a processor-dependent, unpredictably different value on each call.

13 IMAGE\_DISTINCT shall be a logical scalar. It is an **INTENT (IN)** argument. If it has the value true, the  
14 seed accessed by the pseudorandom number generator is set to a processor-dependent value that  
15 is distinct from the value that would be set by a call to RANDOM\_INIT by another **image**. If  
16 it has the value false, the value to which the seed is set does not depend on which **image** calls  
17 RANDOM\_INIT.

18 4 **Example.** The following statement initializes the pseudorandom number generator so that the seed is different  
19 on each call and that the sequence generated will differ from that of another **image**:

20 CALL RANDOM\_INIT (REPEATABLE=.FALSE., IMAGE\_DISTINCT=.TRUE.)

### 21 16.9.156 RANDOM\_NUMBER (HARVEST)

22 1 **Description.** Generate pseudorandom number(s).

23 2 **Class.** Subroutine.

24 3 **Argument.** HARVEST shall be of type real. It is an **INTENT (OUT)** argument. It may be a scalar or an array.  
25 It is assigned pseudorandom numbers from the uniform distribution in the interval  $0 \leq x < 1$ . If **images** use a  
26 common generator, the interleaving of values assigned in unordered segments is processor dependent.

27 4 **Example.**

28 REAL X, Y (10, 10)  
29 ! Initialize X with a pseudorandom number  
30 CALL RANDOM\_NUMBER (HARVEST = X)  
31 CALL RANDOM\_NUMBER (Y)  
32 ! X and Y contain uniformly distributed random numbers

### 33 16.9.157 RANDOM\_SEED ([SIZE, PUT, GET])

34 1 **Description.** Restart or query the pseudorandom number generator.

35 2 **Class.** Subroutine.

36 3 **Arguments.** There shall either be exactly one or no arguments present.

37 SIZE (optional) shall be a default integer scalar. It is an **INTENT (OUT)** argument. It is assigned the number  
38  $N$  of integers that the processor uses to hold the value of the seed.



1 PUT (optional) shall be a default integer array of **rank** one and size  $\geq N$ . It is an **INTENT (IN)** argument. It  
 2 is used in a processor-dependent manner to compute the seed value accessed by the pseudorandom  
 3 number generator.

4 GET (optional) shall be a default integer array of **rank** one and size  $\geq N$ . It is an **INTENT (OUT)** argument.  
 5 It is assigned the value of the seed.

6 4 If no argument is present, the processor assigns a processor-dependent value to the seed.

7 5 The pseudorandom number generator used by **RANDOM\_NUMBER** maintains a seed that is updated during the  
 8 execution of **RANDOM\_NUMBER** and that can be retrieved or changed by **RANDOM\_SEED**. Computation  
 9 of the seed from the argument PUT is performed in a processor-dependent manner. The value assigned to GET  
 10 need not be the same as the value of PUT in an immediately preceding reference to **RANDOM\_SEED**. For  
 11 example, following execution of the statements

```
12     CALL RANDOM_SEED (PUT=SEED1)
13     CALL RANDOM_SEED (GET=SEED2)
```

14 SEED2 need not equal SEED1. When the values differ, the use of either value as the PUT argument in a  
 15 subsequent call to **RANDOM\_SEED** shall result in the same sequence of pseudorandom numbers being generated.  
 16 For example, after execution of the statements

```
17     CALL RANDOM_SEED (PUT=SEED1)
18     CALL RANDOM_SEED (GET=SEED2)
19     CALL RANDOM_NUMBER (X1)
20     CALL RANDOM_SEED (PUT=SEED2)
21     CALL RANDOM_NUMBER (X2)
```

22 X2 equals X1.

## 23 6 Examples.

```
24     CALL RANDOM_SEED                ! Processor initialization
25     CALL RANDOM_SEED (SIZE = K)      ! Puts size of seed in K
26     CALL RANDOM_SEED (PUT = SEED (1 : K)) ! Define seed
27     CALL RANDOM_SEED (GET = OLD (1 : K)) ! Read current seed
```

## 28 16.9.158 RANGE (X)

29 1 **Description.** Decimal exponent range of a numeric model (16.4).

30 2 **Class.** [Inquiry function](#).

31 3 **Argument.** X shall be of type integer, real, or complex. It may be a scalar or an array.

32 4 **Result Characteristics.** Default integer scalar.

33 5 **Result Value.**

34 *Case (i):* If X is of type integer, the result has the value **INT (LOG10 (HUGE (X)))**.

35 *Case (ii):* If X is of type real, the result has the value **INT (MIN (LOG10 (HUGE (X)), -LOG10 (TINY (X))))**.

36 *Case (iii):* If X is of type complex, the result has the value **RANGE (REAL (X))**.

37 6 **Examples.** RANGE (X) has the value 38 for real X whose model is as in Note 16.4, because in this case  
 38 **HUGE (X)** =  $(1 - 2^{-24}) \times 2^{127}$  and **TINY (X)** =  $2^{-127}$ .

### 16.9.159 RANK (A)

1 **Description.** Rank of a data object.

2 **Class.** [Inquiry function](#).

3 **Argument.** A shall be a data object of any type.

4 **Result Characteristics.** Default integer scalar.

5 **Result Value.** The value of the result is the rank of A.

6 **Example.** If X is an [assumed-rank](#) dummy argument and its associated [effective argument](#) is an array of rank 3, RANK(X) has the value 3.

### 16.9.160 REAL (A [, KIND])

1 **Description.** Conversion to real type.

2 **Class.** [Elemental function](#).

3 **Arguments.**

A shall be of type integer, real, or complex, or a [boz-literal-constant](#).

KIND (optional) shall be a scalar integer [constant expression](#).

4 **Result Characteristics.** Real.

*Case (i):* If A is of type integer or real and KIND is present, the kind type parameter is that specified by the value of KIND. If A is of type integer or real and KIND is not present, the kind type parameter is that of default real kind.

*Case (ii):* If A is of type complex and KIND is present, the kind type parameter is that specified by the value of KIND. If A is of type complex and KIND is not present, the kind type parameter is the kind type parameter of A.

*Case (iii):* If A is a [boz-literal-constant](#) and KIND is present, the kind type parameter is that specified by the value of KIND. If A is a [boz-literal-constant](#) and KIND is not present, the kind type parameter is that of default real kind.

5 **Result Value.**

*Case (i):* If A is of type integer or real, the result is equal to a processor-dependent approximation to A.

*Case (ii):* If A is of type complex, the result is equal to a processor-dependent approximation to the real part of A.

*Case (iii):* If A is a [boz-literal-constant](#), the value of the result is the value whose internal representation as a bit sequence is the same as that of A as modified by padding or truncation according to [16.3.3](#). The interpretation of the bit sequence is processor dependent.

6 **Examples.** REAL (-3) has the value -3.0. REAL (Z) has the same kind type parameter and the same value as the real part of the complex variable Z.

### 16.9.161 REDUCE (ARRAY, OPERATION [, MASK, IDENTITY, ORDERED]) or REDUCE (ARRAY, OPERATION, DIM [, MASK, IDENTITY, ORDERED])

1 **Description.** General reduction of array.

2 **Class.** [Transformational function](#).

3 **Arguments.**

- 1 ARRAY shall be an array of any type.
- 2 OPERATION shall be a **pure function** with exactly two arguments; each argument shall be a scalar, non-  
3 locatable, nonpointer, nonpolymorphic, nonoptional dummy data object with the same type and  
4 **type parameters** as ARRAY. If one argument has the **ASYNCHRONOUS**, **TARGET**, or **VALUE**  
5 attribute, the other shall have that attribute. Its result shall be a nonpolymorphic scalar and have  
6 the same type and **type parameters** as ARRAY. OPERATION should implement a mathematically  
7 associative operation. It need not be commutative.
- 8 DIM shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the rank of ARRAY.
- 9 MASK (optional) shall be of type logical and shall be **conformable** with ARRAY.
- 10 IDENTITY (optional) shall be scalar with the same type and **type parameters** as ARRAY.
- 11 ORDERED (optional) shall be a logical scalar.

12 4 **Result Characteristics.** The result is of the same type and **type parameters** as ARRAY. It is scalar if DIM  
13 does not appear; otherwise, the result has **rank**  $n - 1$  and shape  $[d_1, d_2, \dots, d_{\text{DIM}-1}, d_{\text{DIM}+1}, \dots, d_n]$  where  
14  $[d_1, d_2, \dots, d_n]$  is the shape of ARRAY.

15 5 **Result Value.**

16 *Case (i):* The result of REDUCE (ARRAY, OPERATION [, IDENTITY = IDENTITY, ORDERED =  
17 ORDERED]) over the sequence of values in ARRAY is the result of an iterative process. The  
18 initial order of the sequence is array element order. While the sequence has more than one element,  
19 each iteration involves the execution of  $r = \text{OPERATION}(x, y)$  for adjacent  $x$  and  $y$  in the sequence,  
20 with  $x$  immediately preceding  $y$ , and the subsequent replacement of  $x$  and  $y$  with  $r$ ; if ORDERED  
21 is present with the value true,  $x$  and  $y$  shall be the first two elements of the sequence. The process  
22 continues until the sequence has only one element which is the value of the reduction. If the initial  
23 sequence is empty, the result has the value IDENTITY if IDENTITY is present, and otherwise,  
24 **error termination** is initiated.

25 *Case (ii):* The result of REDUCE (ARRAY, OPERATION, MASK = MASK [, IDENTITY = IDENTITY,  
26 ORDERED = ORDERED]) is as for Case (i) except that the initial sequence is only those elements  
27 of ARRAY for which the corresponding elements of MASK are true.

28 *Case (iii):* If ARRAY has **rank** one, REDUCE (ARRAY, OPERATION, DIM = DIM [, MASK = MASK,  
29 IDENTITY = IDENTITY, ORDERED = ORDERED]) has a value equal to that of REDUCE (AR-  
30 RAY, OPERATION [, MASK = MASK, IDENTITY = IDENTITY, ORDERED = ORDERED]).  
31 Otherwise, the value of element  $(s_1, s_2, \dots, s_{\text{DIM}-1}, s_{\text{DIM}+1}, \dots, s_n)$  of REDUCE (ARRAY,  
32 DIM = DIM [, MASK = MASK, IDENTITY = IDENTITY]) is equal to

REDUCE (ARRAY  $(s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n)$ ,  
OPERATION = OPERATION,  
DIM=1  
[, MASK = MASK  $(s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n)$ ,  
IDENTITY = IDENTITY,  
ORDERED = ORDERED] ).

34 6 **Examples.** The following examples all use the function MY\_MULT, which returns the product of its two integer  
35 arguments.

36 *Case (i):* The value of REDUCE ([1, 2, 3], MY\_MULT) is 6.

37 *Case (ii):* REDUCE (C, MY\_MULT, MASK= C > 0, IDENTITY=1) forms the product of the positive  
38 elements of C.

39 *Case (iii):* If B is the array  $\begin{bmatrix} 1 & 3 & 5 \\ 2 & 4 & 6 \end{bmatrix}$ , REDUCE (B, MY\_MULT, DIM = 1) is [2, 12, 30] and REDUCE (B,  
40 MY\_MULT, DIM = 2) is [15, 48].

**NOTE 16.25**

If OPERATION is not computationally associative, REDUCE without ORDERED=.TRUE. with the same argument values might not always produce the same result, as the processor can apply the associative law to the evaluation.

### 16.9.162 REPEAT (STRING, NCOPIES)

1 **Description.** Repetitive string concatenation.

2 **Class.** [Transformational function](#).

3 **Arguments.**

4 STRING shall be a character scalar.

5 NCOPIES shall be an integer scalar. Its value shall not be negative.

6 **Result Characteristics.** Character scalar of length NCOPIES times that of STRING, with the same kind type parameter as STRING.

7 **Result Value.** The value of the result is the concatenation of NCOPIES copies of STRING.

8 **Examples.** REPEAT ('H', 2) has the value HH. REPEAT ('XYZ', 0) has the value of a zero-length string.

### 16.9.163 RESHAPE (SOURCE, SHAPE [, PAD, ORDER])

1 **Description.** Arbitrary shape array construction.

2 **Class.** [Transformational function](#).

3 **Arguments.**

4 SOURCE shall be an array of any type. If PAD is absent or of size zero, the size of SOURCE shall be greater than or equal to PRODUCT (SHAPE). The size of the result is the product of the values of the elements of SHAPE.

5 SHAPE shall be a rank-one integer array. [SIZE](#) ( $x$ ), where  $x$  is the [actual argument](#) corresponding to [SHAPE](#), shall be a [constant expression](#) whose value is positive and less than 16. It shall not have an element whose value is negative.

6 PAD (optional) shall be an array of the same type and type parameters as SOURCE.

7 ORDER (optional) shall be of type integer, shall have the same shape as SHAPE, and its value shall be a permutation of  $(1, 2, \dots, n)$ , where  $n$  is the size of SHAPE. If absent, it is as if it were present with value  $(1, 2, \dots, n)$ .

8 **Result Characteristics.** The result is an array of shape SHAPE (that is, SHAPE (RESHAPE (SOURCE, SHAPE, PAD, ORDER)) is equal to SHAPE) with the same type and type parameters as SOURCE.

9 **Result Value.** The elements of the result, taken in permuted subscript order ORDER  $(1), \dots, \text{ORDER}(n)$ , are those of SOURCE in normal array element order followed if necessary by those of PAD in array element order, followed if necessary by additional copies of PAD in array element order.

10 **Examples.** RESHAPE  $([1, 2, 3, 4, 5, 6], [2, 3])$  has the value  $\begin{bmatrix} 1 & 3 & 5 \\ 2 & 4 & 6 \end{bmatrix}$ .

11 RESHAPE  $([1, 2, 3, 4, 5, 6], [2, 4], [0, 0], [2, 1])$  has the value  $\begin{bmatrix} 1 & 2 & 3 & 4 \\ 5 & 6 & 0 & 0 \end{bmatrix}$ .

### 16.9.164 RRSPACING (X)

1 **Description.** Reciprocal of relative spacing of model numbers.

2 **Class.** [Elemental function](#).

3 **Argument.** X shall be of type real.

4 **Result Characteristics.** Same as X.

1 5 **Result Value.** The result has the value  $|Y \times b^{-e}| \times b^p = \text{ABS}(\text{FRACTION}(Y)) * \text{RADIX}(X) / \text{EPSILON}(X)$ ,  
 2 where  $b$ ,  $e$ , and  $p$  are as defined in 16.4 for  $Y$ , the value nearest to  $X$  in the model for real values whose kind type  
 3 parameter is that of  $X$ ; if there are two such values, the value of greater absolute value is taken. If  $X$  is an IEEE  
 4 infinity, the result is an IEEE NaN. If  $X$  is an IEEE NaN, the result is that NaN.

5 6 **Example.** RRSPACING ( $-3.0$ ) has the value  $0.75 \times 2^{24}$  for reals whose model is as in Note 16.4.

## 6 16.9.165 SAME\_TYPE\_AS (A, B)

7 1 **Description.** Dynamic type equality test.

8 2 **Class.** Inquiry function.

9 3 **Arguments.**

10 A shall be an object of extensible declared type or unlimited polymorphic. If it is a polymorphic  
 11 pointer, it shall not have an undefined association status.

12 B shall be an object of extensible declared type or unlimited polymorphic. If it is a polymorphic  
 13 pointer, it shall not have an undefined association status.

14 4 **Result Characteristics.** Default logical scalar.

15 5 **Result Value.** If the dynamic type of A or B is extensible, the result is true if and only if the dynamic type of  
 16 A is the same as the dynamic type of B. If neither A nor B has extensible dynamic type, the result is processor  
 17 dependent.

### NOTE 16.26

The dynamic type of a disassociated pointer or unallocated allocatable variable is its declared type. An unlimited polymorphic entity has no declared type.

18 6 **Example.** Given the declarations and assignments

```
19     TYPE T1
20         REAL C
21     END TYPE
22     TYPE, EXTENDS(T1) :: T2
23     END TYPE
24     CLASS(T1), POINTER :: P, Q, R
25     ALLOCATE(P, Q)
26     ALLOCATE(T2 :: R)
```

27 the value of SAME\_TYPE\_AS (P, Q) will be true, and the value of SAME\_TYPE\_AS (P, R) will be false.

## 28 16.9.166 SCALE (X, I)

29 1 **Description.** Real number scaled by radix power.

30 2 **Class.** Elemental function.

31 3 **Arguments.**

32 X shall be of type real.

33 I shall be of type integer.

34 4 **Result Characteristics.** Same as X.

35 5 **Result Value.** The result has the value  $X \times b^I$ , where  $b$  is defined in 16.4 for model numbers representing values  
 36 of  $X$ , provided this result is representable; if not, the result is processor dependent.

1 6 **Example.** SCALE (3.0, 2) has the value 12.0 for reals whose model is as in Note 16.4.

## 2 16.9.167 SCAN (STRING, SET [, BACK, KIND])

3 1 **Description.** Character set membership search.

4 2 **Class.** [Elemental](#) function.

5 3 **Arguments.**

6 STRING shall be of type character.

7 SET shall be of type character with the same kind type parameter as STRING.

8 BACK (optional) shall be of type logical.

9 KIND (optional) shall be a scalar integer [constant expression](#).

10 4 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of  
11 KIND; otherwise the kind type parameter is that of default integer type.

12 5 **Result Value.**

13 *Case (i):* If BACK is absent or is present with the value false and if STRING contains at least one character  
14 that is in SET, the value of the result is the position of the leftmost character of STRING that is  
15 in SET.

16 *Case (ii):* If BACK is present with the value true and if STRING contains at least one character that is in  
17 SET, the value of the result is the position of the rightmost character of STRING that is in SET.

18 *Case (iii):* The value of the result is zero if no character of STRING is in SET or if the length of STRING or  
19 SET is zero.

20 6 **Examples.**

21 *Case (i):* SCAN ('FORTRAN', 'TR') has the value 3.

22 *Case (ii):* SCAN ('FORTRAN', 'TR', BACK = [.TRUE.](#)) has the value 5.

23 *Case (iii):* SCAN ('FORTRAN', 'BCD') has the value 0.

## 24 16.9.168 SELECTED\_CHAR\_KIND (NAME)

25 1 **Description.** Character kind selection.

26 2 **Class.** [Transformational](#) function.

27 3 **Argument.** NAME shall be default character scalar.

28 4 **Result Characteristics.** Default integer scalar.

29 5 **Result Value.** If NAME has the value DEFAULT, then the result has a value equal to that of the kind type  
30 parameter of default character. If NAME has the value ASCII, then the result has a value equal to that of the  
31 kind type parameter of [ASCII character](#) if the processor supports such a kind; otherwise the result has the value  
32  $-1$ . If NAME has the value ISO\_10646, then the result has a value equal to that of the kind type parameter of  
33 the [ISO 10646 character](#) kind (corresponding to UCS-4 as specified in ISO/IEC 10646) if the processor supports  
34 such a kind; otherwise the result has the value  $-1$ . If NAME is a processor-defined name of some other character  
35 kind supported by the processor, then the result has a value equal to that kind type parameter value. If NAME is  
36 not the name of a supported character type, then the result has the value  $-1$ . The NAME is interpreted without  
37 respect to case or trailing blanks.

38 6 **Examples.** SELECTED\_CHAR\_KIND ('ASCII') has the value 1 on a processor that uses 1 as the kind type  
39 parameter for the [ASCII character](#) set. The following subroutine produces a Japanese date stamp.

40 SUBROUTINE create\_date\_string(string)

```

1      INTRINSIC date_and_time,selected_char_kind
2      INTEGER,PARAMETER :: ucs4 = selected_char_kind("ISO_10646")
3      CHARACTER(1,UCS4),PARAMETER :: nen=CHAR(INT(Z'5e74'),UCS4), & !year
4      gatsu=CHAR(INT(Z'6708'),UCS4), & !month
5      nichi=CHAR(INT(Z'65e5'),UCS4) !day
6      CHARACTER(len= *, kind= ucs4) string
7      INTEGER values(8)
8      CALL date_and_time(values=values)
9      WRITE(string,1) values(1),nen,values(2),gatsu,values(3),nichi
10     1 FORMAT(IO,A,IO,A,IO,A)
11     END SUBROUTINE

```

## 12 16.9.169 SELECTED\_INT\_KIND (R)

13 1 **Description.** Integer kind selection.

14 2 **Class.** [Transformational function](#).

15 3 **Argument.** R shall be an integer scalar.

16 4 **Result Characteristics.** Default integer scalar.

17 5 **Result Value.** The result has a value equal to the value of the kind type parameter of an integer type that represents all values  $n$  in the range  $-10^R < n < 10^R$ , or if no such kind type parameter is available on the processor, the result is  $-1$ . If more than one kind type parameter meets the criterion, the value returned is the one with the smallest decimal exponent range, unless there are several such values, in which case the smallest of these kind values is returned.

22 6 **Example.** Assume a processor supports two integer kinds, 32 with representation method  $r = 2$  and  $q = 31$ , and 64 with representation method  $r = 2$  and  $q = 63$ . On this processor SELECTED\_INT\_KIND (9) has the value 32 and SELECTED\_INT\_KIND (10) has the value 64.

## 25 16.9.170 SELECTED\_REAL\_KIND ([P, R, RADIX])

26 1 **Description.** Real kind selection.

27 2 **Class.** [Transformational function](#).

28 3 **Arguments.** At least one argument shall be present.

29 P (optional) shall be an integer scalar.

30 R (optional) shall be an integer scalar.

31 RADIX (optional) shall be an integer scalar.

32 4 **Result Characteristics.** Default integer scalar.

33 5 **Result Value.** If P or R is absent, the result value is the same as if it were present with the value zero. If RADIX is absent, there is no requirement on the radix of the selected kind.

35 6 The result has a value equal to a value of the kind type parameter of a real type with decimal precision, as returned by the function PRECISION, of at least P digits, a decimal exponent range, as returned by the function RANGE, of at least R, and a radix, as returned by the function RADIX, of RADIX, if such a kind type parameter is available on the processor.

39 7 Otherwise, the result is  $-1$  if the processor supports a real type with radix RADIX and exponent range of at least R but not with precision of at least P,  $-2$  if the processor supports a real type with radix RADIX and precision of

1 at least P but not with exponent range of at least R,  $-3$  if the processor supports a real type with radix RADIX  
 2 but with neither precision of at least P nor exponent range of at least R,  $-4$  if the processor supports a real type  
 3 with radix RADIX and either precision of at least P or exponent range of at least R but not both together, and  
 4  $-5$  if the processor supports no real type with radix RADIX.

5 8 If more than one kind type parameter value meets the criteria, the value returned is the one with the smallest  
 6 decimal precision, unless there are several such values, in which case the smallest of these kind values is returned.

7 9 **Example.** `SELECTED_REAL_KIND (6, 70)` has the value `KIND (0.0)` on a machine that supports a default  
 8 real approximation method with  $b = 16$ ,  $p = 6$ ,  $e_{\min} = -64$ , and  $e_{\max} = 63$  and does not have a less precise  
 9 approximation method.

### 10 16.9.171 SET\_EXPONENT (X, I)

11 1 **Description.** Real value with specified exponent.

12 2 **Class.** [Elemental function](#).

13 3 **Arguments.**

14 X shall be of type real.

15 I shall be of type integer.

16 4 **Result Characteristics.** Same as X.

17 5 **Result Value.** If X has the value zero, the result has the same value as X. If X is an IEEE infinity, the result is  
 18 an [IEEE NaN](#). If X is an [IEEE NaN](#), the result is the same [NaN](#). Otherwise, the result has the value  $X \times b^{I-e}$ ,  
 19 where  $b$  and  $e$  are as defined in [16.4](#) for the representation for the value of X in the extended real model for the  
 20 kind of X.

21 6 **Example.** `SET_EXPONENT (3.0, 1)` has the value 1.5 for reals whose model is as in Note [16.4](#).

### 22 16.9.172 SHAPE (SOURCE [, KIND])

23 1 **Description.** Shape of an array or a scalar.

24 2 **Class.** [Inquiry function](#).

25 3 **Arguments.**

26 SOURCE shall be a scalar or array of any type. It shall not be an unallocated [allocatable](#) variable or a pointer  
 27 that is not associated. It shall not be an [assumed-size array](#).

28 KIND (optional) shall be a scalar integer [constant expression](#).

29 4 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value  
 30 of KIND; otherwise the kind type parameter is that of default integer type. The result is an array of [rank](#) one  
 31 whose size is equal to the [rank](#) of SOURCE.

32 5 **Result Value.** The result has a value whose  $i^{th}$  element is equal to the extent of dimension  $i$  of SOURCE,  
 33 except that if SOURCE is [assumed-rank](#), and associated with an [assumed-size array](#), the last element is equal to  
 34  $-1$ .

35 6 **Examples.** The value of `SHAPE (A (2:5, -1:1))` is `[4, 3]`. The value of `SHAPE (3)` is the rank-one array of size  
 36 zero.

### 37 16.9.173 SHIFTA (I, SHIFT)

38 1 **Description.** Right shift with fill.

39 2 **Class.** [Elemental function](#).



1 3 **Arguments.**

2 I shall be of type integer.

3 SHIFT shall be of type integer. It shall be nonnegative and less than or equal to BIT\_SIZE (I).

4 4 **Result Characteristics.** Same as I.

5 5 **Result Value.** The result has the value obtained by shifting the bits of I to the right SHIFT bits and replicating  
6 the leftmost bit of I in the left SHIFT bits.

7 6 If SHIFT is zero the result is I. Bits shifted out from the right are lost. The model for the interpretation of an  
8 integer value as a sequence of bits is in 16.3.

9 7 **Example.** SHIFTA (IBSET (0, BIT\_SIZE (0) - 1), 2) is equal to SHIFTL (7, BIT\_SIZE (0) - 3).

10 **16.9.174 SHIFTL (I, SHIFT)**

11 1 **Description.** Left shift.

12 2 **Class.** [Elemental](#) function.

13 3 **Arguments.**

14 I shall be of type integer.

15 SHIFT shall be of type integer. It shall be nonnegative and less than or equal to BIT\_SIZE (I).

16 4 **Result Characteristics.** Same as I.

17 5 **Result Value.** The value of the result is ISHFT (I, SHIFT).

18 6 **Examples.** SHIFTL (3, 1) has the value 6.

19 **16.9.175 SHIFTR (I, SHIFT)**

20 1 **Description.** Right shift.

21 2 **Class.** [Elemental](#) function.

22 3 **Arguments.**

23 I shall be of type integer.

24 SHIFT shall be of type integer. It shall be nonnegative and less than or equal to BIT\_SIZE (I).

25 4 **Result Characteristics.** Same as I.

26 5 **Result Value.** The value of the result is ISHFT (I, -SHIFT).

27 6 **Examples.** SHIFTR (3, 1) has the value 1.

28 **16.9.176 SIGN (A, B)**

29 1 **Description.** Magnitude of A with the sign of B.

30 2 **Class.** [Elemental](#) function.

31 3 **Arguments.**

32 A shall be of type integer or real.

33 B shall be of the same type as A.

34 4 **Result Characteristics.** Same as A.

1 5 **Result Value.**2 *Case (i):* If  $B > 0$ , the value of the result is  $|A|$ .3 *Case (ii):* If  $B < 0$ , the value of the result is  $-|A|$ .4 *Case (iii):* If B is of type integer and  $B=0$ , the value of the result is  $|A|$ .5 *Case (iv):* If B is of type real and is zero, then:6 

- if the processor does not distinguish between positive and negative real zero, or if B is positive real zero, the value of the result is  $|A|$ ;

7 

- if the processor distinguishes between positive and negative real zero, and B is negative real zero, the value of the result is  $-|A|$ .

10 6 **Example.** SIGN ( $-3.0, 2.0$ ) has the value 3.0.11 **16.9.177 SIN (X)**12 1 **Description.** Sine function.13 2 **Class.** [Elemental](#) function.14 3 **Argument.** X shall be of type real or complex.15 4 **Result Characteristics.** Same as X.16 5 **Result Value.** The result has a value equal to a processor-dependent approximation to  $\sin(X)$ . If X is of type real, it is regarded as a value in radians. If X is of type complex, its real part is regarded as a value in radians.18 6 **Example.** SIN (1.0) has the value 0.84147098 (approximately).19 **16.9.178 SINH (X)**20 1 **Description.** Hyperbolic sine function.21 2 **Class.** [Elemental](#) function.22 3 **Argument.** X shall be of type real or complex.23 4 **Result Characteristics.** Same as X.24 5 **Result Value.** The result has a value equal to a processor-dependent approximation to  $\sinh(X)$ . If X is of type complex its imaginary part is regarded as a value in radians.26 6 **Example.** SINH (1.0) has the value 1.1752012 (approximately).27 **16.9.179 SIZE (ARRAY [, DIM, KIND])**28 1 **Description.** Size of an array or one extent.29 2 **Class.** [Inquiry function](#).30 3 **Arguments.**31 ARRAY shall be [assumed-rank](#) or an array. It shall not be an unallocated [allocatable](#) variable or a pointer that is not associated. If ARRAY is an [assumed-size array](#), DIM shall be present with a value less than the [rank](#) of ARRAY.34 DIM (optional) shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the [rank](#) of ARRAY.35 KIND (optional) shall be a scalar integer [constant expression](#).36 4 **Result Characteristics.** Integer scalar. If KIND is present, the kind type parameter is that specified by the value of KIND; otherwise the kind type parameter is that of default integer type.

- 1 5 **Result Value.** If DIM is present, the result has a value equal to the extent of dimension DIM of ARRAY, except  
 2 that if ARRAY is **assumed-rank** and associated with an **assumed-size array** and DIM is present with a value equal  
 3 to the rank of ARRAY, the value is  $-1$ .
- 4 6 If DIM is absent and ARRAY is **assumed-rank**, the result has a value equal to **PRODUCT**(**SHAPE**(ARRAY,  
 5 KIND)). Otherwise, the result has a value equal to the total number of elements of ARRAY.
- 6 7 **Examples.** The value of **SIZE** (A (2:5, -1:1), DIM=2) is 3. The value of **SIZE** (A (2:5, -1:1)) is 12.

**NOTE 16.27**

If ARRAY is **assumed-rank** and has rank zero, DIM cannot be present since it cannot satisfy the requirement  $1 \leq \text{DIM} \leq 0$ .

7 **16.9.180 SPACING (X)**

8 1 **Description.** Spacing of model numbers.

9 2 **Class.** **Elemental** function.

10 3 **Argument.** X shall be of type real.

11 4 **Result Characteristics.** Same as X.

12 5 **Result Value.** If X does not have the value zero and is not an IEEE infinity or NaN, the result has the value  
 13  $b^{e-p}$ , where  $b$ ,  $e$ , and  $p$  are as defined in 16.4 for the value nearest to X in the model for real values whose kind  
 14 type parameter is that of X, provided this result is representable; otherwise, the result is the same as that of  
 15 **TINY** (X). If there are two extended model values equally near to X, the value of greater absolute value is taken.  
 16 If X has the value zero, the result is the same as that of **TINY** (X). If X is an IEEE infinity, the result is an IEEE  
 17 NaN. If X is an IEEE NaN, the result is that NaN.

18 6 **Example.** **SPACING** (3.0) has the value  $2^{-22}$  for reals whose model is as in Note 16.4.

19 **16.9.181 SPREAD (SOURCE, DIM, NCOPIES)**

20 1 **Description.** Value replicated in a new dimension.

21 2 **Class.** **Transformational** function.

22 3 **Arguments.**

23 SOURCE shall be a scalar or array of any type. The **rank** of SOURCE shall be less than 15.

24 DIM shall be an integer scalar with value in the range  $1 \leq \text{DIM} \leq n + 1$ , where  $n$  is the **rank** of SOURCE.

25 NCOPIES shall be an integer scalar.

26 4 **Result Characteristics.** The result is an array of the same type and type parameters as SOURCE and of **rank**  
 27  $n + 1$ , where  $n$  is the **rank** of SOURCE.

28 *Case (i):* If SOURCE is scalar, the shape of the result is (MAX (NCOPIES, 0)).

29 *Case (ii):* If SOURCE is an array with shape  $[d_1, d_2, \dots, d_n]$ , the shape of the result is  $[d_1, d_2, \dots, d_{\text{DIM}-1},$   
 30  $\text{MAX} (\text{NCOPIES}, 0), d_{\text{DIM}}, \dots, d_n]$ .

31 5 **Result Value.**

32 *Case (i):* If SOURCE is scalar, each element of the result has a value equal to SOURCE.

33 *Case (ii):* If SOURCE is an array, the element of the result with subscripts  $(r_1, r_2, \dots, r_{n+1})$  has the value  
 34 SOURCE  $(r_1, r_2, \dots, r_{\text{DIM}-1}, r_{\text{DIM}+1}, \dots, r_{n+1})$ .

35 6 **Examples.** If A is the array [2, 3, 4], **SPREAD** (A, DIM=1, NCOPIES=NC) is the array  $\begin{bmatrix} 2 & 3 & 4 \\ 2 & 3 & 4 \\ 2 & 3 & 4 \end{bmatrix}$  if NC

1 has the value 3 and is a zero-sized array if NC has the value 0.

## 2 **16.9.182 SQRT (X)**

3 1 **Description.** Square root.

4 2 **Class.** [Elemental function](#).

5 3 **Argument.** X shall be of type real or complex. If X is real, its value shall be greater than or equal to zero.

6 4 **Result Characteristics.** Same as X.

7 5 **Result Value.** The result has a value equal to a processor-dependent approximation to the square root of X. A  
8 result of type complex is the principal value with the real part greater than or equal to zero. When the real part  
9 of the result is zero, the imaginary part has the same sign as the imaginary part of X.

10 6 **Example.** SQRT (4.0) has the value 2.0 (approximately).

## 11 **16.9.183 STOPPED\_IMAGES([TEAM, KIND])**

12 1 **Description.** Indices of stopped images.

13 2 **Class.** [Transformational function](#).

14 3 **Arguments.**

15 TEAM (optional) shall be a scalar of type [TEAM\\_TYPE](#) from the intrinsic module [ISO\\_FORTRAN\\_ENV](#),  
16 whose value identifies the [current](#) or an [ancestor](#) team. If TEAM is absent the team specified is the  
17 [current team](#).

18 KIND (optional) shall be a scalar integer [constant expression](#).

19 4 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value  
20 of KIND; otherwise, the kind type parameter is that of default integer type. The result is an array of rank one  
21 whose size is equal to the number of images in the specified [team](#) that have initiated [normal termination](#).

22 5 **Result Value.** The elements of the result are the values of the indices of the images that are known to have  
23 initiated [normal termination](#) in the specified team, in numerically increasing order. If the executing image has  
24 previously executed an [image control statement](#) whose STAT= specifier assigned the value [STAT\\_STOPPED\\_-](#)  
25 [IMAGE](#) from the intrinsic module [ISO\\_FORTRAN\\_ENV](#) or invoked a [collective subroutine](#) whose STAT argu-  
26 ment was assigned [STAT\\_STOPPED\\_IMAGE](#), at least one of the images participating in that [image control](#)  
27 [statement](#) or collective invocation shall be known to have initiated [normal termination](#).

28 6 **Examples.** If image 3 is the only image in the [current team](#) that is known to have initiated [normal termination](#),  
29 STOPPED\_IMAGES() will have the value [3]. If there are no images in the [current team](#) that have initiated  
30 [normal termination](#), the value of STOPPED\_IMAGES() will be a zero-sized array.

## 31 **16.9.184 STORAGE\_SIZE (A [, KIND])**

32 1 **Description.** Storage size in bits.

33 2 **Class.** [Inquiry function](#).

34 3 **Arguments.**

35 A shall be a data object of any type. If it is [polymorphic](#) it shall not be an undefined [pointer](#). If  
36 it is [unlimited polymorphic](#) or has any [deferred type parameters](#), it shall not be an unallocated  
37 [allocatable](#) variable or a [disassociated](#) or undefined pointer.

38 KIND (optional) shall be a scalar integer [constant expression](#).

39 4 **Result Characteristics.** Integer scalar. If KIND is present, the kind type parameter is that specified by the

1 value of KIND; otherwise, the kind type parameter is that of default integer type.

2 5 **Result Value.** The result value is the size expressed in bits for an element of an array that has the [dynamic](#)  
3 [type](#) and type parameters of A. If the type and type parameters are such that [storage association \(19.5.3\)](#) applies,  
4 the result is consistent with the [named constants](#) defined in the intrinsic module [ISO\\_FORTRAN\\_ENV](#).

#### NOTE 16.28

An array element might take more bits to store than an isolated scalar, since any hardware-imposed alignment requirements for array elements might not apply to a simple scalar variable.

#### NOTE 16.29

This is intended to be the size in memory that an object takes when it is stored; this might differ from the size it takes during expression handling (which might be the native register size) or when stored in a file. If an object is never stored in memory but only in a register, this function nonetheless returns the size it would take if it were stored in memory.

5 6 **Example.** STORAGE\_SIZE (1.0) has the same value as the [named constant](#) NUMERIC\_STORAGE\_SIZE in  
6 the intrinsic module [ISO\\_FORTRAN\\_ENV](#).

### 7 16.9.185 SUM (ARRAY, DIM [, MASK]) or SUM (ARRAY [, MASK])

8 1 **Description.** Array reduced by addition.

9 2 **Class.** [Transformational function](#).

10 3 **Arguments.**

11 ARRAY shall be an array of [numeric type](#).

12 DIM shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the [rank](#) of ARRAY.

13 MASK (optional) shall be of type logical and shall be [conformable](#) with ARRAY.

14 4 **Result Characteristics.** The result is of the same type and kind type parameter as ARRAY. It is scalar if  
15 DIM does not appear; otherwise, the result has [rank](#)  $n - 1$  and shape  $[d_1, d_2, \dots, d_{\text{DIM}-1}, d_{\text{DIM}+1}, \dots, d_n]$  where  
16  $[d_1, d_2, \dots, d_n]$  is the shape of ARRAY.

17 5 **Result Value.**

18 *Case (i):* The result of SUM (ARRAY) has a value equal to a processor-dependent approximation to the sum  
19 of all the elements of ARRAY or has the value zero if ARRAY has size zero.

20 *Case (ii):* The result of SUM (ARRAY, MASK = MASK) has a value equal to a processor-dependent approx-  
21 imation to the sum of the elements of ARRAY corresponding to the true elements of MASK or has  
22 the value zero if there are no true elements.

23 *Case (iii):* If ARRAY has [rank](#) one, SUM (ARRAY, DIM = DIM [, MASK = MASK]) has a value equal to that  
24 of SUM (ARRAY [, MASK = MASK]). Otherwise, the value of element  $(s_1, s_2, \dots, s_{\text{DIM}-1}, s_{\text{DIM}+1},$   
25  $\dots, s_n)$  of SUM (ARRAY, DIM = DIM [, MASK = MASK]) is equal to

26  $\text{SUM}(\text{ARRAY}(s_1, s_2, \dots, s_{\text{DIM}-1}, :, s_{\text{DIM}+1}, \dots, s_n) [, \text{MASK} = \text{MASK}(s_1, s_2, \dots, s_{\text{DIM}-1},$   
27  $:, s_{\text{DIM}+1}, \dots, s_n) ])$ .

28 6 **Examples.**

29 *Case (i):* The value of SUM ([1, 2, 3]) is 6.

30 *Case (ii):* SUM (C, MASK = C > 0.0) forms the sum of the positive elements of C.

31 *Case (iii):* If B is the array  $\begin{bmatrix} 1 & 3 & 5 \\ 2 & 4 & 6 \end{bmatrix}$ , SUM (B, DIM = 1) is [3, 7, 11] and SUM (B, DIM = 2) is [9, 12].

## 16.9.186 SYSTEM\_CLOCK ([COUNT, COUNT\_RATE, COUNT\_MAX])

1 **Description.** Query system clock.

2 **Class.** Subroutine.

3 **Arguments.**

4 COUNT (optional) shall be an integer scalar. It is an **INTENT (OUT)** argument. It is assigned a processor-dependent value based on the value of a processor clock, or **–HUGE (COUNT)** if there is no clock for the invoking **image**. The processor-dependent value is incremented by one for each clock count until the value COUNT\_MAX is reached and is reset to zero at the next count. It lies in the range 0 to COUNT\_MAX if there is a clock.

5 COUNT\_RATE (optional) shall be an integer or real scalar. It is an **INTENT (OUT)** argument. It is assigned a processor-dependent approximation to the number of processor clock counts per second, or zero if there is no clock for the invoking **image**.

6 COUNT\_MAX (optional) shall be an integer scalar. It is an **INTENT (OUT)** argument. It is assigned the maximum value that COUNT can have, or zero if there is no clock for the invoking **image**.

7 Whether an **image** has no clock, has a single clock of its own, or shares a clock with another **image**, is processor dependent.

8 **Example.** If the processor clock is a 24-hour clock that registers time at approximately 18.20648193 ticks per second, at 11:30 A.M. the reference

9 CALL SYSTEM\_CLOCK (COUNT = C, COUNT\_RATE = R, COUNT\_MAX = M)

10 defines  $C = (11 \times 3600 + 30 \times 60) \times 18.20648193 = 753748$ ,  $R = 18.20648193$ , and  $M = 24 \times 3600 \times 18.20648193 - 1 = 1573039$ .

## 16.9.187 TAN (X)

1 **Description.** Tangent function.

2 **Class.** **Elemental** function.

3 **Argument.** X shall be of type real or complex.

4 **Result Characteristics.** Same as X.

5 **Result Value.** The result has a value equal to a processor-dependent approximation to  $\tan(X)$ . If X is of type real, it is regarded as a value in radians. If X is of type complex, its real part is regarded as a value in radians.

6 **Example.** TAN (1.0) has the value 1.5574077 (approximately).

## 16.9.188 TANH (X)

1 **Description.** Hyperbolic tangent function.

2 **Class.** **Elemental** function.

3 **Argument.** X shall be of type real or complex.

4 **Result Characteristics.** Same as X.

5 **Result Value.** The result has a value equal to a processor-dependent approximation to  $\tanh(X)$ . If X is of type complex its imaginary part is regarded as a value in radians.

6 **Example.** TANH (1.0) has the value 0.76159416 (approximately).

### 16.9.189 TEAM\_NUMBER([TEAM])

1 **Description.** Team number.

2 **Class.** Transformational function.

3 **Argument.** TEAM (optional) shall be a scalar of type `TEAM_TYPE` from the intrinsic module `ISO_FOR-`  
 4 `TRAN_ENV`, whose value identifies the `current` or an `ancestor` team. If TEAM is absent, the team specified is  
 5 the `current team`.  
 6

7 **Result Characteristics.** Default integer scalar.

8 **Result Value.** The result has the value  $-1$  if the specified team is the `initial team`; otherwise, the result value  
 9 is equal to the positive integer that identifies the specified team within its parent team.

10 **Example.** The team number can be used to control which statements get executed, for example:

```

11     TYPE(Team_Type) :: ODD_EVEN
12     ...
13     FORM TEAM (2-MOD(ME,2), ODD_EVEN)
14     ...
15     CHANGE TEAM (ODD_EVEN)
16         SELECT CASE (TEAM_NUMBER())
17             CASE (1)
18                 ! Case for images with odd image indices in the parent team.
19             CASE (2)
20                 ! Case for images with even image indices in the parent team.
21         END SELECT
22     END TEAM
  
```

### 16.9.190 THIS\_IMAGE ([TEAM]) or THIS\_IMAGE (COARRAY [, TEAM]) or THIS\_IMAGE (COARRAY, DIM [, TEAM])

24 1 **Description.** `Cosubscript(s)` for this `image`.

25 2 **Class.** Transformational function.

26 3 **Arguments.**

27 `COARRAY` shall be a `coarray` of any type. If it is `allocatable` it shall be allocated. If it is of type `TEAM_TYPE`  
 28 from the intrinsic module `ISO_FORTRAN_ENV`, the `TEAM` argument shall appear.

29 `DIM` shall be an integer scalar. Its value shall be in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the `corank` of  
 30 `COARRAY`.

31 `TEAM` (optional) shall be a scalar of type `TEAM_TYPE` from the intrinsic module `ISO_FORTRAN_ENV`,  
 32 whose value identifies the `current` or an `ancestor team`. If `COARRAY` appears, it shall be established  
 33 in that `team`.

34 4 **Result Characteristics.** Default integer. It is scalar if `COARRAY` does not appear or `DIM` appears; otherwise,  
 35 the result has `rank` one and its size is equal to the `corank` of `COARRAY`.

36 5 **Result Value.**

37 *Case (i):* The result of `THIS_IMAGE` ([TEAM]) is a scalar with a value equal to the `index` of the invoking  
 38 `image` in the `team` specified by `TEAM`, if present, or in the `current team` if absent.

1 *Case (ii):* The result of THIS\_IMAGE (COARRAY [, TEAM = TEAM]) is the sequence of **cosubscript** values  
 2 for COARRAY that would specify the invoking **image** in the **team** specified by TEAM, if present,  
 3 or in the **current team** if absent.

4 *Case (iii):* The result of THIS\_IMAGE (COARRAY, DIM [, TEAM = TEAM]) is the value of **cosubscript**  
 5 DIM in the sequence of **cosubscript** values for COARRAY that would specify the invoking **image** in  
 6 the **team** specified by TEAM, if present, or in the **current team** if absent.

7 **6 Examples.** If A is declared by the statement  
 8 REAL A (10, 20) [10, 0:9, 0:\*]  
 9 then on **image** 5, THIS\_IMAGE ( ) has the value 5 and THIS\_IMAGE (A) has the value [5, 0, 0]. For the same  
 10 **coarray** on **image** 213, THIS\_IMAGE (A) has the value [3, 1, 2].

11 7 The following code uses **image** 1 to read data. The other **images** then copy the data.

```
12     IF (THIS_IMAGE()==1) READ (*,*) P
13     SYNC ALL
14     P = P[1]
```

### 15 16.9.191 TINY (X)

16 1 **Description.** Smallest positive model number.

17 2 **Class.** **Inquiry function.**

18 3 **Argument.** X shall be a real scalar or array.

19 4 **Result Characteristics.** Scalar with the same type and kind type parameter as X.

20 5 **Result Value.** The result has the value  $b^{e_{\min}-1}$  where  $b$  and  $e_{\min}$  are as defined in 16.4 for the model representing  
 21 numbers of the same type and kind type parameter as X.

22 6 **Example.** TINY (X) has the value  $2^{-127}$  for real X whose model is as in Note 16.4.

### 23 16.9.192 TRAILZ (I)

24 1 **Description.** Number of trailing zero bits.

25 2 **Class.** **Elemental function.**

26 3 **Argument.** I shall be of type integer.

27 4 **Result Characteristics.** Default integer.

28 5 **Result Value.** If all of the bits of I are zero, the result value is BIT\_SIZE (I). Otherwise, the result value is the  
 29 position of the rightmost 1 bit in I. The model for the interpretation of an integer value as a sequence of bits is  
 30 in 16.3.

31 6 **Examples.** TRAILZ (8) has the value 3.

### 32 16.9.193 TRANSFER (SOURCE, MOLD [, SIZE])

33 1 **Description.** Transfer physical representation.

34 2 **Class.** **Transformational function.**

35 3 **Arguments.**

36 SOURCE shall be a scalar or array of any type.



1 MOLD shall be a scalar or array of any type. If it is a variable, it need not be defined. If the storage size of  
 2 SOURCE is greater than zero and MOLD is an array, a scalar with the type and [type parameters](#)  
 3 of MOLD shall not have a storage size equal to zero.

4 SIZE (optional) shall be an integer scalar. The corresponding [actual argument](#) shall not be an optional dummy  
 5 argument.

6 **4 Result Characteristics.** The result is of the same type and type parameters as MOLD.

7 *Case (i):* If MOLD is a scalar and SIZE is absent, the result is a scalar.

8 *Case (ii):* If MOLD is an array and SIZE is absent, the result is an array and of [rank](#) one. Its size is as small  
 9 as possible such that its physical representation is not shorter than that of SOURCE.

10 *Case (iii):* If SIZE is present, the result is an array of [rank](#) one and size SIZE.

11 **5 Result Value.** If the physical representation of the result has the same length as that of SOURCE, the physical  
 12 representation of the result is that of SOURCE. If the physical representation of the result is longer than that  
 13 of SOURCE, the physical representation of the leading part is that of SOURCE and the remainder is processor  
 14 dependent. If the physical representation of the result is shorter than that of SOURCE, the physical representation  
 15 of the result is the leading part of SOURCE. If D and E are scalar variables such that the physical representation  
 16 of D is as long as or longer than that of E, the value of TRANSFER (TRANSFER (E, D), E) shall be the value  
 17 of E. IF D is an array and E is an array of [rank](#) one, the value of TRANSFER (TRANSFER (E, D), E, SIZE (E))  
 18 shall be the value of E.

19 **6 Examples.**

20 *Case (i):* TRANSFER (1082130432, 0.0) has the value 4.0 on a processor that represents the values 4.0 and  
 21 1082130432 as the string of binary digits 0100 0000 1000 0000 0000 0000 0000 0000.

22 *Case (ii):* TRANSFER ([1.1, 2.2, 3.3], [(0.0, 0.0)]) is a complex rank-one array of length two whose first  
 23 element has the value (1.1, 2.2) and whose second element has a real part with the value 3.3. The  
 24 imaginary part of the second element is processor dependent.

25 *Case (iii):* TRANSFER ([1.1, 2.2, 3.3], [(0.0, 0.0)], 1) is a complex rank-one array of length one whose only  
 26 element has the value (1.1, 2.2).

## 27 16.9.194 TRANSPOSE (MATRIX)

28 **1 Description.** Transpose of an array of [rank](#) two.

29 **2 Class.** [Transformational function](#).

30 **3 Argument.** MATRIX shall be a rank-two array of any type.

31 **4 Result Characteristics.** The result is an array of the same type and type parameters as MATRIX and with  
 32 [rank](#) two and shape  $[n, m]$  where  $[m, n]$  is the shape of MATRIX.

33 **5 Result Value.** Element  $(i, j)$  of the result has the value  $MATRIX(j + LBOUND(MATRIX, 1) - 1, i +$   
 34  $LBOUND(MATRIX, 2) - 1)$ .

35 **6 Example.** If A is the array  $\begin{bmatrix} 1 & 2 & 3 \\ 4 & 5 & 6 \\ 7 & 8 & 9 \end{bmatrix}$ , then TRANSPOSE (A) has the value  $\begin{bmatrix} 1 & 4 & 7 \\ 2 & 5 & 8 \\ 3 & 6 & 9 \end{bmatrix}$ .

## 36 16.9.195 TRIM (STRING)

37 **1 Description.** String without trailing blanks.

38 **2 Class.** [Transformational function](#).

39 **3 Argument.** STRING shall be a character scalar.

1 4 **Result Characteristics.** Character with the same kind type parameter value as STRING and with a length  
 2 that is the length of STRING less the number of trailing blanks in STRING. If STRING contains no nonblank  
 3 characters, the result has zero length.

4 5 **Result Value.** The value of the result is the same as STRING except any trailing blanks are removed.

5 6 **Example.** TRIM (' A B ') has the value ' A B'.

## 6 16.9.196 UBOUND (ARRAY [, DIM, KIND])

7 1 **Description.** Upper bound(s).

8 2 **Class.** Inquiry function.

9 3 **Arguments.**

10 ARRAY shall be **assumed-rank** or an array. It shall not be an unallocated **allocatable** array or a pointer that  
 11 is not associated. If ARRAY is an **assumed-size array**, DIM shall be present with a value less than  
 12 the **rank** of ARRAY.

13 DIM (optional) shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the **rank** of ARRAY.  
 14 The corresponding **actual argument** shall not be an optional dummy argument, a disassociated  
 15 pointer, or an unallocated allocatable.

16 KIND (optional) shall be a scalar integer **constant expression**.

17 4 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of  
 18 KIND; otherwise the kind type parameter is that of default integer type. The result is scalar if DIM is present;  
 19 otherwise, the result is an array of **rank** one and size  $n$ , where  $n$  is the **rank** of ARRAY.

20 5 **Result Value.**

21 *Case (i):* If DIM is present, ARRAY is a **whole array**, and dimension DIM of ARRAY has nonzero extent,  
 22 the result has a value equal to the upper bound for subscript DIM of ARRAY. Otherwise, if DIM  
 23 is present and ARRAY is **assumed-rank**, the value of the result is as if ARRAY were a **whole array**,  
 24 with the extent of the final dimension of ARRAY when ARRAY is associated with an **assumed-size**  
 25 **array** being considered to be  $-1$ . Otherwise, if DIM is present, the result has a value equal to the  
 26 number of elements in dimension DIM of ARRAY.

27 *Case (ii):* If ARRAY has rank zero, UBOUND (ARRAY) has a value that is a zero-sized array. Otherwise,  
 28 UBOUND (ARRAY) has a value whose  $i^{\text{th}}$  element is equal to UBOUND (ARRAY,  $i$ ), for  $i = 1, 2,$   
 29  $\dots, n$ , where  $n$  is the **rank** of ARRAY. UBOUND (ARRAY, KIND=KIND) has a value whose  $i^{\text{th}}$   
 30 element is equal to UBOUND (ARRAY,  $i$ , KIND=KIND), for  $i = 1, 2, \dots, n$ , where  $n$  is the  
 31 **rank** of ARRAY.

32 6 **Examples.** If A is declared by the statement

```
33 REAL A (2:3, 7:10)
```

34 then UBOUND (A) is [3, 10] and UBOUND (A, DIM = 2) is 10.

### NOTE 16.30

If ARRAY is **assumed-rank** and has rank zero, DIM cannot be present since it cannot satisfy the requirement  $1 \leq \text{DIM} \leq 0$ .

## 35 16.9.197 UCBOUND (COARRAY [, DIM, KIND])

36 1 **Description.** Upper **cobound**(s) of a **coarray**.

37 2 **Class.** Inquiry function.

1 3 **Arguments.**

2 COARRAY shall be a **coarray** of any type. It may be a scalar or an array. If it is **allocatable** it shall be allocated.  
 3 DIM (optional) shall be an integer scalar with a value in the range  $1 \leq \text{DIM} \leq n$ , where  $n$  is the **corank**  
 4 of COARRAY. The corresponding **actual argument** shall not be an optional dummy argument, a  
 5 disassociated pointer, or an unallocated allocatable.  
 6 KIND (optional) shall be a scalar integer **constant expression**.

7 4 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of  
 8 KIND; otherwise, the kind type parameter is that of default integer type. The result is scalar if DIM is present;  
 9 otherwise, the result is an array of **rank** one and size  $n$ , where  $n$  is the **corank** of COARRAY.

10 5 **Result Value.** The final upper **cobound** is the final **cosubscript** in the **cosubscript** list for the **coarray** that selects  
 11 the **image** whose **index** is equal to the number of **images** in the **team current** when COARRAY was established.

12 *Case (i):* If DIM is present, the result has a value equal to the upper **cobound** for **codimension** DIM of  
 13 COARRAY.

14 *Case (ii):* If DIM is absent, the result has a value whose  $i^{\text{th}}$  element is equal to the upper **cobound** for  
 15 **codimension**  $i$  of COARRAY, for  $i = 1, 2, \dots, n$ , where  $n$  is the **corank** of COARRAY.

16 6 **Examples.** If NUM\_IMAGES( ) has the value 30 and A is allocated by the statement

17 ALLOCATE (A [2:3, 0:7, \*])

18 then UCOBOUND (A) is [3, 7, 2] and UCOBOUND (A, DIM=2) is 7. Note that the **cosubscripts** [3, 7, 2] do  
 19 not correspond to an actual **image**.

20 **16.9.198 UNPACK (VECTOR, MASK, FIELD)**

21 1 **Description.** Vector unpacked into an array.

22 2 **Class.** **Transformational function.**

23 3 **Arguments.**

24 VECTOR shall be a rank-one array of any type. Its size shall be at least  $t$  where  $t$  is the number of true  
 25 elements in MASK.

26 MASK shall be a logical array.

27 FIELD shall be of the same type and type parameters as VECTOR and shall be **conformable** with MASK.

28 4 **Result Characteristics.** The result is an array of the same type and type parameters as VECTOR and the  
 29 same shape as MASK.

30 5 **Result Value.** The element of the result that corresponds to the  $i^{\text{th}}$  true element of MASK, in array element  
 31 order, has the value VECTOR ( $i$ ) for  $i = 1, 2, \dots, t$ , where  $t$  is the number of true values in MASK. Each other  
 32 element has a value equal to FIELD if FIELD is scalar or to the corresponding element of FIELD if it is an array.

33 6 **Examples.** Particular values can be “scattered” to particular positions in an array by using UNPACK. If M is the

34 array  $\begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$ , V is the array [1, 2, 3], and Q is the logical mask  $\begin{bmatrix} . & \text{T} & . \\ \text{T} & . & . \\ . & . & \text{T} \end{bmatrix}$ , where “T” represents true

35 and “.” represents false, then the result of UNPACK (V, MASK = Q, FIELD = M) has the value  $\begin{bmatrix} 1 & 2 & 0 \\ 1 & 1 & 0 \\ 0 & 0 & 3 \end{bmatrix}$

36 and the result of UNPACK (V, MASK = Q, FIELD = 0) has the value  $\begin{bmatrix} 0 & 2 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 3 \end{bmatrix}$ .

## 16.9.199 VERIFY (STRING, SET [, BACK, KIND])

1 **Description.** Character set non-membership search.

2 **Class.** [Elemental](#) function.

3 **Arguments.**

4 STRING shall be of type character.

5 SET shall be of type character with the same kind type parameter as STRING.

6 BACK (optional) shall be of type logical.

7 KIND (optional) shall be a scalar integer [constant expression](#).

8 **4 Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of KIND; otherwise the kind type parameter is that of default integer type.

9 **5 Result Value.**

10 *Case (i):* If BACK is absent or has the value false and if STRING contains at least one character that is not in SET, the value of the result is the position of the leftmost character of STRING that is not in SET.

11 *Case (ii):* If BACK is present with the value true and if STRING contains at least one character that is not in SET, the value of the result is the position of the rightmost character of STRING that is not in SET.

12 *Case (iii):* The value of the result is zero if each character in STRING is in SET or if STRING has zero length.

13 **6 Examples.**

14 *Case (i):* VERIFY ('ABBA', 'A') has the value 2.

15 *Case (ii):* VERIFY ('ABBA', 'A', BACK = [.TRUE.](#)) has the value 3.

16 *Case (iii):* VERIFY ('ABBA', 'AB') has the value 0.

## 16.10 Standard intrinsic modules

### 16.10.1 General

17 This document defines five standard intrinsic modules: a Fortran environment module, a set of three modules to support floating-point exceptions and IEEE arithmetic, and a module to support interoperability with the C programming language.

18 The intrinsic modules [IEEE\\_EXCEPTIONS](#), [IEEE\\_ARITHMETIC](#), and [IEEE\\_FEATURES](#) are described in [Clause 17](#). The intrinsic module [ISO\\_C\\_BINDING](#) is described in [Clause 18](#). The [module procedures](#) described in [16.10.2](#) are [pure](#).

#### NOTE 16.31

The types and procedures defined in [standard intrinsic](#) modules are not themselves [intrinsic](#).

19 A processor may extend the [standard intrinsic](#) modules to provide public entities in them in addition to those specified in this document.

### 16.10.2 The ISO\_FORTRAN\_ENV intrinsic module

#### 16.10.2.1 General

20 The intrinsic module [ISO\\_FORTRAN\\_ENV](#) provides public entities relating to the Fortran environment.

1 2 The processor shall provide the [named constants](#), derived types, and procedures described in subclause [16.10.2](#).  
2 In the detailed descriptions below, procedure names are generic and not specific.

### 3 16.10.2.2 ATOMIC\_INT\_KIND

4 1 The value of the default integer scalar constant ATOMIC\_INT\_KIND is the kind type parameter value of type  
5 integer variables for which the processor supports atomic operations specified by [atomic subroutines](#).

### 6 16.10.2.3 ATOMIC\_LOGICAL\_KIND

7 1 The value of the default integer scalar constant ATOMIC\_LOGICAL\_KIND is the kind type parameter value  
8 of type logical variables for which the processor supports atomic operations specified by [atomic subroutines](#).

### 9 16.10.2.4 CHARACTER\_KINDS

10 1 The values of the elements of the default integer array constant CHARACTER\_KINDS are the kind values  
11 supported by the processor for variables of type character. The order of the values is processor dependent. The  
12 rank of the array is one, its lower bound is one, and its size is the number of character kinds supported.

### 13 16.10.2.5 CHARACTER\_STORAGE\_SIZE

14 1 The value of the default integer scalar constant CHARACTER\_STORAGE\_SIZE is the size expressed in bits  
15 of the [character storage unit](#) ([19.5.3.2](#)).

### 16 16.10.2.6 COMPILER\_OPTIONS ( )

17 1 **Description.** Processor-dependent string describing the options that controlled the program translation phase.

18 2 **Class.** [Transformational function](#).

19 3 **Argument.** None.

20 4 **Result Characteristics.** Default character scalar with processor-dependent length.

21 5 **Result Value.** A processor-dependent value which describes the options that controlled the translation phase of  
22 program execution. This value should include relevant information that could be useful for diagnosing problems  
23 at a later date.

24 6 **Example.** COMPILER\_OPTIONS ( ) might have the value '/OPTIMIZE /FLOAT=IEEE'.

### 25 16.10.2.7 COMPILER\_VERSION ( )

26 1 **Description.** Processor-dependent string identifying the program translation phase.

27 2 **Class.** [Transformational function](#).

28 3 **Argument.** None.

29 4 **Result Characteristics.** Default character scalar with processor-dependent length.

30 5 **Result Value.** A processor-dependent value that identifies the name and version of the program translation  
31 phase of the processor. This value should include relevant information that could be useful for diagnosing problems  
32 at a later date.

33 6 **Example.** COMPILER\_VERSION ( ) might have the value 'Fast KL-10 Compiler Version 7'.

**NOTE 16.32**

Relevant information that could be useful for diagnosing problems at a later date might include compiler release and patch level, default compiler arguments, environment variable values, and run time library requirements. A processor might include this information in an object file automatically, without the user needing to save the result of this function in a variable.

1     **16.10.2.8 CURRENT\_TEAM**

2     1 The value of the default integer scalar constant CURRENT\_TEAM identifies the [current team](#) when it is used  
3     as the LEVEL argument to [GET\\_TEAM](#).

4     **16.10.2.9 ERROR\_UNIT**

5     1 The value of the default integer scalar constant ERROR\_UNIT identifies the processor-dependent [preconnected](#)  
6     [external unit](#) used for the purpose of error reporting (12.5). This [unit](#) may be the same as OUTPUT\_UNIT.  
7     The value shall not be  $-1$ .

8     **16.10.2.10 EVENT\_TYPE**

9     1 EVENT\_TYPE is a derived type with private components. It is an extensible type with no type parameters.  
10    Each nonallocatable component is fully default-initialized.

11   2 A scalar variable of type EVENT\_TYPE is an [event variable](#). The value of an event variable includes its event  
12   count, which is updated by execution of a sequence of [EVENT POST](#) or [EVENT WAIT](#) statements. The effect  
13   of each change is as if the intrinsic subroutine [ATOMIC\\_ADD](#) were executed with a variable that stores the  
14   event count as its ATOM argument. A coarray that is of type EVENT\_TYPE may be referenced or defined  
15   during execution of a segment that is unordered relative to the execution of another segment in which that  
16   coarray is defined. The event count is of type integer with kind [ATOMIC\\_INT\\_KIND](#) from the intrinsic module  
17   [ISO\\_FORTRAN\\_ENV](#). The initial value of the event count of an event variable is zero.

18   C1603 A named entity with declared type EVENT\_TYPE, or which has a noncoarray [potential subobject](#)  
19   [component](#) with declared type EVENT\_TYPE, shall be a variable. A component that is of such a type  
20   shall be a data component.

21   C1604 A named variable with [declared type](#) EVENT\_TYPE shall be a coarray. A named variable with a  
22   noncoarray [potential subobject component](#) of type EVENT\_TYPE shall be a coarray.

23   C1605 An [event variable](#) shall not appear in a variable definition context except as the *event-variable* in an  
24   [EVENT POST](#) or [EVENT WAIT](#) statement, as an *allocate-object*, or as an [actual argument](#) in a reference  
25   to a procedure with an [explicit interface](#) if the corresponding [dummy argument](#) has [INTENT \(INOUT\)](#).

26   C1606 A variable with a nonpointer subobject of type EVENT\_TYPE shall not appear in a variable definition  
27   context except as an *allocate-object* in an [ALLOCATE statement](#) without a [SOURCE= specifier](#), as an  
28   *allocate-object* in a [DEALLOCATE statement](#), or as an [actual argument](#) in a reference to a procedure  
29   with an [explicit interface](#) if the corresponding [dummy argument](#) has [INTENT \(INOUT\)](#).

**NOTE 16.33**

The restrictions against changing an [event variable](#) except via [EVENT POST](#) and [EVENT WAIT](#) statements ensure the integrity of its value and facilitate efficient implementation, particularly when special synchronization is needed for correct event handling.

**NOTE 16.34**

Updates to variables via atomic subroutines are coherent but not necessarily consistent, so a processor might have to use extra synchronization to obtain the consistency required for the segments ordered by [EVENT POST](#) and [EVENT WAIT](#) statements.

**1 16.10.2.11 FILE\_STORAGE\_SIZE**

2 1 The value of the default integer scalar constant FILE\_STORAGE\_SIZE is the size expressed in bits of the [file](#)  
3 [storage unit](#) (12.3.5).

**4 16.10.2.12 INITIAL\_TEAM**

5 1 The value of the default integer scalar constant INITIAL\_TEAM identifies the [initial team](#) when it is used as  
6 the LEVEL argument to [GET\\_TEAM](#).

**7 16.10.2.13 INPUT\_UNIT**

8 1 The value of the default integer scalar constant INPUT\_UNIT identifies the same processor-dependent [external](#)  
9 [unit preconnected](#) for sequential formatted input as the one identified by an asterisk in a [READ statement](#); this  
10 [unit](#) is the one used for a [READ statement](#) that does not contain an input/output control list (12.6.4.3). The  
11 value shall not be  $-1$ .

**12 16.10.2.14 INT8, INT16, INT32, and INT64**

13 1 The values of these default integer scalar constants shall be those of the [kind type parameters](#) that specify an  
14 INTEGER type whose storage size expressed in bits is 8, 16, 32, and 64 respectively. If, for any of these constants,  
15 the processor supports more than one kind of that size, it is processor dependent which kind value is provided. If  
16 the processor supports no kind of a particular size, that constant shall be equal to  $-2$  if the processor supports  
17 a kind with larger size and  $-1$  otherwise.

**18 16.10.2.15 INTEGER\_KINDS**

19 1 The values of the elements of the default integer array constant INTEGER\_KINDS are the kind values supported  
20 by the processor for variables of type integer. The order of the values is processor dependent. The rank of the  
21 array is one, its lower bound is one, and its size is the number of integer kinds supported.

**22 16.10.2.16 IOSTAT\_END**

23 1 The value of the default integer scalar constant IOSTAT\_END is assigned to the variable specified in an [IOSTAT=](#)  
24 [specifier](#) (12.11.5) if an end-of-file condition occurs during execution of an [input statement](#) and no error condition  
25 occurs. This value shall be negative.

**26 16.10.2.17 IOSTAT\_EOR**

27 1 The value of the default integer scalar constant IOSTAT\_EOR is assigned to the variable specified in an [IOSTAT=](#)  
28 [specifier](#) (12.11.5) if an end-of-record condition occurs during execution of an [input statement](#) and no end-of-file  
29 or error condition occurs. This value shall be negative and different from the value of IOSTAT\_END.

**30 16.10.2.18 IOSTAT\_INQUIRE\_INTERNAL\_UNIT**

31 1 The value of the default integer scalar constant IOSTAT\_INQUIRE\_INTERNAL\_UNIT is assigned to the  
32 variable specified in an [IOSTAT= specifier](#) in an [INQUIRE statement](#) (12.10) if a [file-unit-number](#) identifies an  
33 [internal unit](#) in that statement.

**NOTE 16.35**

This can only occur when a [defined input/output](#) procedure is called by the processor as the result of executing a parent [data transfer statement](#) (12.6.4.8.3) for an [internal unit](#).

### 1 16.10.2.19 LOCK\_TYPE

2 1 LOCK\_TYPE is a [derived type](#) with private components; no component is allocatable or a pointer. It is an  
3 [extensible type](#) with no [type parameters](#). All components have default initialization.

4 2 A scalar variable of type LOCK\_TYPE is a [lock variable](#). A lock variable can have one of two states: locked and  
5 unlocked. The unlocked state is represented by the one value that is the default value of a LOCK\_TYPE variable;  
6 this is the value specified by the [structure constructor](#) LOCK\_TYPE ( ). The locked state is represented by all  
7 other values. The value of a lock variable can be changed with the [LOCK](#) and [UNLOCK](#) statements ([11.6.10](#)).

8 C1607 A named entity with declared type LOCK\_TYPE, or which has a noncoarray [potential subobject com-](#)  
9 [ponent](#) with declared type LOCK\_TYPE, shall be a variable. A component that is of such a type shall  
10 be a data component.

11 C1608 A named variable with [declared type](#) LOCK\_TYPE shall be a [coarray](#). A named variable with a  
12 noncoarray [potential subobject component](#) of type LOCK\_TYPE shall be a [coarray](#).

13 C1609 A lock variable shall not appear in a variable definition context except as the [lock-variable](#) in a [LOCK](#) or  
14 [UNLOCK](#) statement, as an [allocate-object](#), or as an [actual argument](#) in a reference to a procedure with  
15 an [explicit interface](#) where the corresponding [dummy argument](#) has [INTENT \(INOUT\)](#).

16 C1610 A variable with a subobject of type LOCK\_TYPE shall not appear in a variable definition context except  
17 as an [allocate-object](#) or as an [actual argument](#) in a reference to a procedure with an [explicit interface](#)  
18 where the corresponding [dummy argument](#) has [INTENT \(INOUT\)](#).

#### NOTE 16.36

|                                                                                                                                                                                                                                                                                        |
|----------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|
| The restrictions against changing a lock variable except via the <a href="#">LOCK</a> and <a href="#">UNLOCK</a> statements ensure the integrity of its value and facilitate efficient implementation, particularly when special synchronization is needed for correct lock operation. |
|----------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|

### 19 16.10.2.20 LOGICAL\_KINDS

20 1 The values of the elements of the default integer array constant LOGICAL\_KINDS are the kind values supported  
21 by the processor for variables of type logical. The order of the values is processor dependent. The rank of the  
22 array is one, its lower bound is one, and its size is the number of logical kinds supported.

### 23 16.10.2.21 NUMERIC\_STORAGE\_SIZE

24 1 The value of the default integer scalar constant NUMERIC\_STORAGE\_SIZE is the size expressed in bits of the  
25 [numeric storage unit](#) ([19.5.3.2](#)).

### 26 16.10.2.22 OUTPUT\_UNIT

27 1 The value of the default integer scalar constant OUTPUT\_UNIT identifies the same processor-dependent [external](#)  
28 [unit preconnected](#) for sequential formatted output as the one identified by an asterisk in a [WRITE statement](#)  
29 ([12.6.4.3](#)). The value shall not be  $-1$ .

### 30 16.10.2.23 PARENT\_TEAM

31 1 The value of the default integer scalar constant PARENT\_TEAM identifies the [parent team](#) when it is used as  
32 the LEVEL argument to [GET\\_TEAM](#).

### 33 16.10.2.24 REAL\_KINDS

34 1 The values of the elements of the default integer array constant REAL\_KINDS are the kind values supported by  
35 the processor for variables of type real. The order of the values is processor dependent. The rank of the array is  
36 one, its lower bound is one, and its size is the number of real kinds supported.



### 1 16.10.2.25 REAL32, REAL64, and REAL128

2 1 The values of these default integer scalar [named constants](#) shall be those of the kind type parameters that specify  
3 a REAL type whose storage size expressed in bits is 32, 64, and 128 respectively. If, for any of these constants,  
4 the processor supports more than one kind of that size, it is processor dependent which kind value is provided. If  
5 the processor supports no kind of a particular size, that constant shall be equal to  $-2$  if the processor supports  
6 kinds of a larger size and  $-1$  otherwise.

### 7 16.10.2.26 STAT\_FAILED\_IMAGE

8 1 If the processor has the ability to detect that an image has failed, the value of the default integer scalar constant  
9 STAT\_FAILED\_IMAGE is positive; otherwise, the value of STAT\_FAILED\_IMAGE is negative. If an image  
10 involved in execution of an [image control statement](#), a reference to a [coindexed object](#), or execution of a [collective](#)  
11 or [atomic](#) subroutine has failed, and no other error condition occurs, the value of STAT\_FAILED\_IMAGE is  
12 assigned to the variable specified in a [STAT= specifier](#) in the execution of an [image control statement](#) or reference  
13 to a [coindexed object](#), or to the STAT argument in an invocation of a [collective](#) or [atomic](#) subroutine.

### 14 16.10.2.27 STAT\_LOCKED

15 1 The value of the default integer scalar constant STAT\_LOCKED is assigned to the variable specified in a [STAT=](#)  
16 [specifier](#) (11.6.11) of a [LOCK statement](#) if the [lock variable](#) is locked by the executing [image](#).

### 17 16.10.2.28 STAT\_LOCKED\_OTHER\_IMAGE

18 1 The value of the default integer scalar constant STAT\_LOCKED\_OTHER\_IMAGE is assigned to the variable  
19 specified in a [STAT= specifier](#) (11.6.11) of an [UNLOCK statement](#) if the [lock variable](#) is locked by another [image](#).

### 20 16.10.2.29 STAT\_STOPPED\_IMAGE

21 1 The value of the default integer scalar constant STAT\_STOPPED\_IMAGE is assigned to the variable specified  
22 in a [STAT= specifier](#) (9.7.4, 11.6.11), if execution of the statement with that specifier requires synchronization  
23 with an image that has initiated [normal termination](#). It is assigned to a STAT argument in a reference to a  
24 [collective subroutine](#) if any [image](#) of the [current team](#) has initiated [normal termination](#). This value shall be  
25 positive.

### 26 16.10.2.30 STAT\_UNLOCKED

27 1 The value of the default integer scalar constant STAT\_UNLOCKED is assigned to the variable specified in a  
28 [STAT= specifier](#) (11.6.11) of an [UNLOCK statement](#) if the [lock variable](#) is unlocked.

### 29 16.10.2.31 STAT\_UNLOCKED\_FAILED\_IMAGE

30 1 The value of the default integer scalar constant STAT\_UNLOCKED\_FAILED\_IMAGE is assigned to the vari-  
31 able specified in a [STAT= specifier](#) (11.6.11) of a [LOCK statement](#) if the [lock variable](#) is unlocked because of the  
32 failure of the [image](#) that locked it.

### 33 16.10.2.32 TEAM\_TYPE

34 1 TEAM\_TYPE is a [derived type](#) with private components; no component is a pointer. It is an [extensible type](#)  
35 with no [type parameters](#). Each nonallocatable component is fully default-initialized.

36 2 A scalar variable of type TEAM\_TYPE is a [team variable](#), and can identify a team. The default initial value of  
37 a team variable does not identify any team.

1     **16.10.2.33 Uniqueness of named constant values**

- 2     1 The values of the named constants `IOSTAT_INQUIRE_INTERNAL_UNIT`, `STAT_FAILED_IMAGE`, `STAT_-`  
3     `LOCKED`, `STAT_LOCKED_OTHER_IMAGE`, `STAT_STOPPED_IMAGE`, `STAT_UNLOCKED`, and `STAT_-`  
4     `UNLOCKED_FAILED_IMAGE` shall be distinct.

# 17 Exceptions and IEEE arithmetic

## 17.1 Overview of IEEE arithmetic support

- 1 The intrinsic modules IEEE\_EXCEPTIONS, IEEE\_ARITHMETIC, and IEEE\_FEATURES provide support for the facilities defined by ISO/IEC/IEEE 60559:2011\*. Whether the modules are provided is processor dependent. If the module IEEE\_FEATURES is provided, which of the named constants defined in this document are included is processor dependent. The module IEEE\_ARITHMETIC behaves as if it contained a [USE statement](#) for IEEE\_EXCEPTIONS; everything that is public in IEEE\_EXCEPTIONS is public in IEEE\_ARITHMETIC.

### NOTE 17.1

The types and procedures defined in these modules are not themselves intrinsic.

- 2 If IEEE\_EXCEPTIONS or IEEE\_ARITHMETIC is accessible in a [scoping unit](#), the exceptions IEEE\_OVERFLOW and IEEE\_DIVIDE\_BY\_ZERO are supported in the [scoping unit](#) for all kinds of real and complex IEEE floating-point data. Which other exceptions are supported can be determined by the [inquiry function](#) IEEE\_SUPPORT\_FLAG ([17.11.51](#)); whether control of [halting](#) is supported can be determined by the [inquiry function](#) IEEE\_SUPPORT\_HALTING. The extent of support of the other exceptions may be influenced by the accessibility of the [named constants](#) IEEE\_INEXACT\_FLAG, IEEE\_INVALID\_FLAG, and IEEE\_UNDERFLOW\_FLAG of the module IEEE\_FEATURES. If a [scoping unit](#) has access to IEEE\_UNDERFLOW\_FLAG of IEEE\_FEATURES, within the [scoping unit](#) the processor shall support underflow and return true from IEEE\_SUPPORT\_FLAG (IEEE\_UNDERFLOW, X) for at least one kind of real. Similarly, if IEEE\_INEXACT\_FLAG or IEEE\_INVALID\_FLAG is accessible, within the [scoping unit](#) the processor shall support the exception and return true from the corresponding [inquiry function](#) for at least one kind of real. If IEEE\_HALTING is accessible, within the [scoping unit](#) the processor shall support control of [halting](#) and return true from IEEE\_SUPPORT\_HALTING (FLAG) for the flag.

### NOTE 17.2

IEEE\_INVALID is not required to be supported whenever IEEE\_EXCEPTIONS is accessed. This is to allow a processor whose arithmetic does not conform to ISO/IEC/IEEE 60559:2011 to provide support for overflow and divide\_by\_zero. On a processor which does support ISO/IEC/IEEE 60559:2011, invalid is an equally serious condition.

- 3 If a [scoping unit](#) does not access IEEE\_FEATURES, IEEE\_EXCEPTIONS, or IEEE\_ARITHMETIC, the level of support is processor dependent, and need not include support for any exceptions. If a flag is signaling on entry to such a [scoping unit](#), the processor ensures that it is signaling on exit. If a flag is quiet on entry to such a [scoping unit](#), whether it is signaling on exit is processor dependent.
- 4 Additional ISO/IEC/IEEE 60559:2011 facilities are available from the module IEEE\_ARITHMETIC. The extent of support may be influenced by the accessibility of the [named constants](#) of the module IEEE\_FEATURES. If a [scoping unit](#) has access to IEEE\_DATATYPE of IEEE\_FEATURES, within the [scoping unit](#) the processor shall support IEEE arithmetic and return true from IEEE\_SUPPORT\_DATATYPE (X) ([17.11.48](#)) for at least one kind of real. Similarly, if IEEE\_DENORMAL, IEEE\_DIVIDE, IEEE\_INF, IEEE\_NAN, IEEE\_ROUNDING, IEEE\_SQRT, or IEEE\_SUBNORMAL is accessible, within the [scoping unit](#) the processor shall support the feature and return true from the corresponding [inquiry function](#) for at least one kind of real. In the case of IEEE\_ROUNDING, it shall return true for the rounding modes IEEE\_NEAREST, IEEE\_TO\_ZERO, IEEE\_UP, and IEEE\_DOWN; support for [IEEE\\_AWAY](#) is also required if there is at least one kind of real X for

\* Because ISO/IEC/IEEE 60559:2011 was originally an IEEE standard, its facilities are widely known as “IEEE arithmetic”, and this terminology is used by this document.

1 which `IEEE_SUPPORT_DATATYPE` (X) is true and `RADIX` (X) is equal to ten. Note that the effect of  
 2 `IEEE_DENORMAL` is the same as that of `IEEE_SUBNORMAL`.

3 5 Execution might be slowed on some processors by the support of some features. If `IEEE_EXCEPTIONS` or  
 4 `IEEE_ARITHMETIC` is accessed but `IEEE_FEATURES` is not accessed, the supported subset of features is  
 5 processor dependent. The processor's fullest support is provided when all of `IEEE_FEATURES` is accessed as in

```
6 USE, INTRINSIC :: IEEE_ARITHMETIC; USE, INTRINSIC :: IEEE_FEATURES
```

7 but execution might then be slowed by the presence of a feature that is not needed. In all cases, the extent of  
 8 support can be determined by the [inquiry functions](#).

## 9 17.2 Derived types, constants, and operators defined in the modules

10 1 The modules `IEEE_EXCEPTIONS`, `IEEE_ARITHMETIC`, and `IEEE_FEATURES` define five derived types,  
 11 whose components are all private. No [direct component](#) of any of these types is [allocatable](#) or a pointer.

12 2 The module `IEEE_EXCEPTIONS` defines the following types.

- 13 • `IEEE_FLAG_TYPE` is for identifying a particular exception flag. Its only possible values are those of  
 14 [named constants](#) defined in the module: `IEEE_INVALID`, `IEEE_OVERFLOW`, `IEEE_DIVIDE_BY_`  
 15 `ZERO`, `IEEE_UNDERFLOW`, and `IEEE_INEXACT`. The module also defines the array [named constants](#)  
 16 `IEEE_USUAL` = [ `IEEE_OVERFLOW`, `IEEE_DIVIDE_BY_ZERO`, `IEEE_INVALID` ] and `IEEE_`  
 17 `ALL` = [ `IEEE_USUAL`, `IEEE_UNDERFLOW`, `IEEE_INEXACT` ].
- 18 • `IEEE_MODES_TYPE` is for representing the floating-point modes.
- 19 • `IEEE_STATUS_TYPE` is for representing the floating-point status.

20 3 The module `IEEE_ARITHMETIC` defines the following types, constants, and operators.

- 21 • The type `IEEE_CLASS_TYPE`, for identifying a class of floating-point values. Its only possible  
 22 values are those of [named constants](#) defined in the module: `IEEE_SIGNALING_NAN`, `IEEE_QUI-`  
 23 `ET_NAN`, `IEEE_NEGATIVE_INF`, `IEEE_NEGATIVE_NORMAL`, `IEEE_NEGATIVE_DENORMAL`,  
 24 `IEEE_NEGATIVE_ZERO`, `IEEE_POSITIVE_ZERO`, `IEEE_POSITIVE_SUBNORMAL`, `IEEE_POS-`  
 25 `ITIVE_NORMAL`, `IEEE_POSITIVE_INF`, and `IEEE_OTHER_VALUE`. The named constants `IEEE_`  
 26 `NEGATIVE_DENORMAL` and `IEEE_POSITIVE_DENORMAL` are defined with the same value as  
 27 [IEEE\\_NEGATIVE\\_SUBNORMAL](#) and [IEEE\\_POSITIVE\\_SUBNORMAL](#) respectively.
- 28 • The type `IEEE_ROUND_TYPE`, for identifying a particular rounding mode. Its only possible values  
 29 are those of [named constants](#) defined in the module: `IEEE_NEAREST`, `IEEE_TO_ZERO`, `IEEE_UP`,  
 30 `IEEE_DOWN`, [IEEE\\_AWAY](#) and `IEEE_OTHER` for the rounding modes specified in this document.
- 31 • The [pure elemental operator](#) `==` for two values of one of these types to return true if the values are the  
 32 same and false otherwise.
- 33 • The [pure elemental operator](#) `/=` for two values of one of these types to return true if the values differ and  
 34 false otherwise.

35 4 The module `IEEE_FEATURES` defines the type `IEEE_FEATURES_TYPE`, for expressing the need for particu-  
 36 lar ISO/IEC/IEEE 60559:2011 features. Its only possible values are those of [named constants](#) defined in the mod-  
 37 ule: `IEEE_DATATYPE`, `IEEE_DENORMAL`, `IEEE_DIVIDE`, `IEEE_HALTING`, `IEEE_INEXACT_FLAG`,  
 38 `IEEE_INF`, `IEEE_INVALID_FLAG`, `IEEE_NAN`, `IEEE_ROUNDING`, `IEEE_SQRT`, `IEEE_SUBNORMAL`,  
 39 and `IEEE_UNDERFLOW_FLAG`.

## 40 17.3 The exceptions

41 1 The exceptions are the following.

- 1       • IEEE\_OVERFLOW occurs in an intrinsic real addition, subtraction, multiplication, division, or conversion  
 2       by the intrinsic function REAL, as specified by ISO/IEC/IEEE 60559:2011 if IEEE\_SUPPORT\_DATA-  
 3       TYPE is true for the operands of the operation or conversion, and as determined by the processor otherwise.  
 4       It occurs in an intrinsic real exponentiation as determined by the processor. It occurs in a complex op-  
 5       eration, or conversion by the intrinsic function CMPLX, if it is caused by the calculation of the real or  
 6       imaginary part of the result.
- 7       • IEEE\_DIVIDE\_BY\_ZERO occurs in a real division as specified by ISO/IEC/IEEE 60559:2011 if IEEE\_-  
 8       SUPPORT\_DATATYPE is true for the operands of the division, and as determined by the processor  
 9       otherwise. It is processor-dependent whether it occurs in a real exponentiation with a negative exponent.  
 10       It occurs in a complex division if it is caused by the calculation of the real or imaginary part of the result.
- 11       • IEEE\_INVALID occurs when a real or complex operation or assignment is invalid; possible examples are  
 12       SQRT (X) when X is real and has a nonzero negative value, and conversion to an integer (by assignment,  
 13       an intrinsic procedure, or a procedure defined in an intrinsic module) when the result is too large to be  
 14       representable. IEEE\_INVALID occurs for numeric relational intrinsic operations as specified below.
- 15       • IEEE\_UNDERFLOW occurs when the result for an intrinsic real operation or assignment has an absolute  
 16       value less than a processor-dependent limit and loss of accuracy is detected, or the real or imaginary part  
 17       of the result for an intrinsic complex operation or assignment has an absolute value less than a processor-  
 18       dependent limit and loss of accuracy is detected.
- 19       • IEEE\_INEXACT occurs when the result of a real or complex operation or assignment is not exact.
- 20   2   Each exception has a flag whose value is either quiet or signaling. The value can be determined by the subroutine  
 21    IEEE\_GET\_FLAG. Its initial value is quiet and it signals when the associated exception occurs. Its status can  
 22    also be changed by the subroutine IEEE\_SET\_FLAG or the subroutine IEEE\_SET\_STATUS. Once signaling  
 23    within a procedure, it remains signaling unless set quiet by an invocation of the subroutine IEEE\_SET\_FLAG  
 24    or the subroutine IEEE\_SET\_STATUS.
- 25   3   If a flag is signaling on entry to a procedure other than IEEE\_GET\_FLAG or IEEE\_GET\_STATUS, the  
 26    processor will set it to quiet on entry and restore it to signaling on return. If a flag signals during execution of a  
 27    procedure, the processor shall not set it to quiet on return.
- 28   4   Evaluation of a [specification expression](#) might cause an exception to signal.
- 29   5   In a [scoping unit](#) that has access to IEEE\_EXCEPTIONS or IEEE\_ARITHMETIC, if an intrinsic procedure  
 30    or a procedure defined in an intrinsic module executes normally, the values of the flags IEEE\_OVERFLOW,  
 31    IEEE\_DIVIDE\_BY\_ZERO, and IEEE\_INVALID shall be as on entry to the procedure, even if one or more of  
 32    them signals during the calculation. If a real or complex result is too large for the procedure to handle, IEEE\_-  
 33    OVERFLOW may signal. If a real or complex result is a NaN because of an invalid operation (for example,  
 34    LOG (-1.0)), IEEE\_INVALID may signal. Similar rules apply to format processing and to intrinsic operations:  
 35    no signaling flag shall be set quiet and no quiet flag shall be set signaling because of an intermediate calculation  
 36    that does not affect the result.
- 37   6   In a scoping unit that has access to IEEE\_EXCEPTIONS or IEEE\_ARITHMETIC, if  $x_1$  and  $x_2$  are numeric  
 38    entities, the type of  $x_1 + x_2$  is real, and IEEE\_SUPPORT\_NAN ( $x_1 + x_2$ ) is true, the relational intrinsic  
 39    operation  $x_1$  *rel-op*  $x_2$  shall signal IEEE\_INVALID as specified for the conditional predicate of ISO/IEC/IEEE  
 40    60559:2011 corresponding to *rel-op* indicated by Table 17.1. If the types or kind type parameters of  $x_1$  or  $x_2$  differ,  
 41    the conversions (10.1.5.5.1) might signal exceptions instead of or in addition to an IEEE\_INVALID exception  
 42    signaled by the comparison.

**NOTE 17.3**

Each comparison predicate defined by ISO/IEC/IEEE 60559:2011 is either unordered signaling or unordered quiet. An unordered signaling predicate signals an invalid operation exception if and only if one of the values being compared is a NaN. An unordered quiet predicate signals an invalid operation exception if and only if one of the values being compared is a signaling NaN. The comparison predicates do not signal any other exceptions.

Table 17.1: IEEE relational operator correspondence

| Operator   | ISO/IEC/IEEE 60559:2011 comparison predicate |
|------------|----------------------------------------------|
| .LT. or <  | compareSignalingLess                         |
| .LE. or <= | compareSignalingLessEqual                    |
| .GT. or >  | compareSignalingGreater                      |
| .GE. or >= | compareSignalingGreaterEqual                 |
| .EQ. or == | compareQuietEqual                            |
| .NE. or /= | compareQuietNotEqual                         |

1 7 In a scoping unit that has access to IEEE\_EXCEPTIONS or IEEE\_ARITHMETIC, if  $x_1$  or  $x_2$  are numeric  
 2 entities, the type of  $x_1 + x_2$  is complex, and IEEE\_SUPPORT\_NAN (REAL ( $x_1 + x_2$ )) is true, the intrinsic  
 3 equality or inequality operation between  $x_1$  and  $x_2$  may signal IEEE\_INVALID if the value of the real or  
 4 imaginary part of either operand is a signaling NaN. If any conversions are done before the values are compared,  
 5 those conversions might signal exceptions instead of or in addition to an IEEE\_INVALID exception signaled by  
 6 the comparison.

7 8 In a sequence of statements that has no invocations of IEEE\_GET\_FLAG, IEEE\_SET\_FLAG, IEEE\_GET\_  
 8 STATUS, IEEE\_SET\_HALTING\_MODE, or IEEE\_SET\_STATUS, if the execution of an operation would  
 9 cause an exception to signal but after execution of the sequence no value of a variable depends on the operation,  
 10 whether the exception is signaling is processor dependent. For example, when Y has the value zero, whether the  
 11 code

```
12 X = 1.0/Y
13 X = 3.0
```

14 signals IEEE\_DIVIDE\_BY\_ZERO is processor dependent. Another example is the following:

```
15 REAL, PARAMETER :: X=0.0, Y=6.0
16 IF (1.0/X == Y) PRINT *, 'Hello world'
```

17 where the processor is permitted to discard the IF statement because the logical expression can never be true  
 18 and no value of a variable depends on it.

19 9 An exception shall not signal if this could arise only during execution of an operation beyond those required or  
 20 permitted by the standard. For example, the statement

```
21 IF (F (X) > 0.0) Y = 1.0/Z
```

22 shall not signal IEEE\_DIVIDE\_BY\_ZERO when both F (X) and Z are zero and the statement

```
23 WHERE (A > 0.0) A = 1.0/A
```

24 shall not signal IEEE\_DIVIDE\_BY\_ZERO. On the other hand, when X has the value 1.0 and Y has the value  
 25 0.0, the expression

```
26 X>0.00001 .OR. X/Y>0.00001
```

27 is permitted to cause the signaling of IEEE\_DIVIDE\_BY\_ZERO.

28 10 The processor need not support IEEE\_INVALID, IEEE\_UNDERFLOW, and IEEE\_INEXACT. If an exception  
 29 is not supported, its flag is always quiet. The inquiry function IEEE\_SUPPORT\_FLAG can be used to inquire  
 30 whether a particular flag is supported.

## 17.4 The rounding modes

- 1 This document specifies a binary rounding mode that affects floating-point arithmetic with radix two, and a decimal rounding mode that affects floating-point arithmetic with radix ten. Unqualified references to the rounding mode with respect to a particular arithmetic operation or operands refers to the mode for the radix of the operation or operands, and other unqualified references to the rounding mode refers to both binary and decimal rounding modes.
- 2 ISO/IEC/IEEE 60559:2011 specifies five possible rounding-direction attributes: `roundTiesToEven`, `roundTowardZero`, `roundTowardPositive`, `roundTowardNegative`, and `roundTiesToAway`. These correspond to the rounding modes `IEEE_NEAREST`, `IEEE_TO_ZERO`, `IEEE_UP`, `IEEE_DOWN`, and `IEEE_AWAY` respectively. The rounding mode `IEEE_OTHER` does not correspond to any ISO/IEC/IEEE 60559:2011 rounding-direction attribute; if supported, the effect of this rounding mode is processor dependent.
- 3 The subroutine `IEEE_GET_ROUNDING_MODE` can be used to get the rounding modes. The initial rounding modes are processor dependent.
- 4 If the processor supports the alteration of the rounding modes during execution, the subroutine `IEEE_SET_ROUNDING_MODE` can be used to alter them. The inquiry function `IEEE_SUPPORT_ROUNDING` can be used to inquire whether this facility is available for a particular mode. The inquiry function `IEEE_SUPPORT_IO` can be used to inquire whether rounding for base conversion in formatted input/output ([12.5.6.16](#), [12.6.2.13](#), [13.7.2.3.8](#)) is as specified in ISO/IEC/IEEE 60559:2011.
- 5 In a procedure other than `IEEE_SET_ROUNDING_MODE` or `IEEE_SET_STATUS`, the processor shall not change the rounding modes on entry, and on return shall ensure that the rounding modes are the same as they were on entry.

### NOTE 17.4

ISO/IEC/IEEE 60559:2011 requires support for `roundTiesToAway` only for decimal floating-point.

### NOTE 17.5

ISO/IEC/IEEE 60559:2011 requires that there is a language-defined means to specify a constant value for the rounding-direction attribute for all standard operations in a block. The means provided by this document are a `CALL` to `IEEE_GET_ROUNDING_MODE` at the beginning of the block followed by a `CALL` to `IEEE_SET_ROUNDING_MODE` with constant arguments, together with another `CALL` to `IEEE_SET_ROUNDING_MODE` at the end of the block to restore the rounding mode.

### NOTE 17.6

Within a program, all literal constants that have the same form have the same value ([7.1.4](#)). Therefore, the value of a literal constant is not affected by the rounding modes.

## 17.5 Underflow mode

- 1 Some processors allow control during program execution of whether underflow produces a subnormal number in conformance with ISO/IEC/IEEE 60559:2011 (gradual underflow) or produces zero instead (abrupt underflow). On some processors, floating-point performance is typically better in abrupt underflow mode than in gradual underflow mode.
- 2 Control over the underflow mode is exercised by invocation of `IEEE_SET_UNDERFLOW_MODE`. The subroutine `IEEE_GET_UNDERFLOW_MODE` can be used to get the underflow mode. The inquiry function `IEEE_SUPPORT_UNDERFLOW_CONTROL` can be used to inquire whether this facility is available. The initial underflow mode is processor dependent. In a procedure other than `IEEE_SET_UNDERFLOW_MODE` or `IEEE_SET_STATUS`, the processor shall not change the underflow mode on entry, and on return shall ensure that the underflow mode is the same as it was on entry.

1 3 The underflow mode affects only floating-point calculations whose type is that of an X for which IEEE\_SUP-  
2 PORT\_UNDERFLOW\_CONTROL returns true.

## 3 17.6 Halting

4 1 Some processors allow control during program execution of whether to abort or continue execution after an  
5 exception. Such control is exercised by invocation of the subroutine IEEE\_SET\_HALTING\_MODE. Halting  
6 is not precise and may occur any time after the exception has occurred. The [inquiry function](#) IEEE\_SUP-  
7 PORT\_HALTING can be used to inquire whether this facility is available. The initial halting mode is processor  
8 dependent. In a procedure other than IEEE\_SET\_HALTING\_MODE or IEEE\_SET\_STATUS, the processor  
9 shall not change the halting mode on entry, and on return shall ensure that the halting mode is the same as it  
10 was on entry.

## 11 17.7 The floating-point modes and status

12 1 The values of the rounding modes, [underflow mode](#), and [halting mode](#) are collectively called the floating-point  
13 modes. The values of all the supported flags for exceptions and the floating-point modes are collectively called the  
14 floating-point status. The floating-point modes can be stored in a scalar variable of type [IEEE\\_MODES\\_TYPE](#)  
15 with the subroutine [IEEE\\_GET\\_MODES](#) and restored with the subroutine [IEEE\\_SET\\_MODES](#). The floating-  
16 point status can be stored in a scalar variable of type [IEEE\\_STATUS\\_TYPE](#) with the subroutine [IEEE\\_GET\\_-](#)  
17 STATUS and restored with the subroutine [IEEE\\_SET\\_STATUS](#). There are no facilities for finding the values of  
18 particular flags represented by such a variable.

### NOTE 17.7

Each [image](#) has its own floating-point status ([5.3.4](#)).

### NOTE 17.8

Some processors hold all these flags and modes in one or two status registers that can be obtained and set as a whole faster than all individual flags and modes can be obtained and set. These procedures are provided to exploit this feature.

### NOTE 17.9

The processor is required to ensure that a call to a Fortran procedure does not change the floating-point status other than by setting exception flags to signaling.

## 19 17.8 Exceptional values

20 1 ISO/IEC/IEEE 60559:2011 specifies the following exceptional floating-point values.

- 21 • Subnormal values have very small absolute values and reduced precision.
- 22 • Infinite values (+infinity and -infinity) are created by overflow or division by zero.
- 23 • Not-a-Number ( [NaN](#) ) values are undefined values or values created by an invalid operation.

24 2 A value that does not fall into the above classes is called a normal number.

25 3 The functions [IEEE\\_IS\\_FINITE](#), [IEEE\\_IS\\_NAN](#), [IEEE\\_IS\\_NEGATIVE](#), and [IEEE\\_IS\\_NORMAL](#) are  
26 provided to test whether a value is finite, [NaN](#), negative, or normal. The function [IEEE\\_VALUE](#) is provided to  
27 generate an IEEE number of any class, including an infinity or a [NaN](#). The [inquiry functions](#) [IEEE\\_SUPPORT\\_-](#)  
28 [SUBNORMAL](#), [IEEE\\_SUPPORT\\_INF](#), and [IEEE\\_SUPPORT\\_NAN](#) are provided to determine whether these  
29 facilities are available for a particular kind of real.



## 17.9 IEEE arithmetic

- 1
- 2 1 The [inquiry function](#) `IEEE_SUPPORT_DATATYPE` can be used to inquire whether IEEE arithmetic is sup-  
3 ported for a particular kind of real. Complete conformance with ISO/IEC/IEEE 60559:2011 is not required,  
4 but
- 5 • the normal numbers shall be exactly those of an ISO/IEC/IEEE 60559:2011 floating-point format,
  - 6 • for at least one rounding mode, the intrinsic operations of addition, subtraction and multiplication shall  
7 conform whenever the operands and result specified by ISO/IEC/IEEE 60559:2011 are normal numbers,
  - 8 • the IEEE function `abs` shall be provided by the intrinsic function [ABS](#),
  - 9 • the IEEE operation `rem` shall be provided by the function [IEEE\\_REM](#), and
  - 10 • the IEEE functions `copysign`, `scalb`, `logb`, `nextafter`, and `unordered` shall be provided by the functions  
11 [IEEE\\_COPY\\_SIGN](#), [IEEE\\_SCALB](#), [IEEE\\_LOGB](#), [IEEE\\_NEXT\\_AFTER](#), and [IEEE\\_UNORDERED](#),  
12 respectively,
- 13 for that kind of real.
- 14 2 The [inquiry function](#) `IEEE_SUPPORT_NAN` is provided to inquire whether the processor supports [IEEE NaNs](#).  
15 Where these are supported, the result of the intrinsic operations `+`, `-`, and `*`, and the functions [IEEE\\_REM](#)  
16 and [IEEE\\_RINT](#) from the intrinsic module `IEEE_ARITHMETIC`, shall conform to ISO/IEC/IEEE 60559:2011  
17 when the result is an [IEEE NaN](#).
- 18 3 The [inquiry function](#) `IEEE_SUPPORT_INF` is provided to inquire whether the processor supports IEEE infinities.  
19 Where these are supported, the result of the intrinsic operations `+`, `-`, and `*`, and the functions [IEEE\\_REM](#)  
20 and [IEEE\\_RINT](#) from the intrinsic module `IEEE_ARITHMETIC`, shall conform to ISO/IEC/IEEE 60559:2011  
21 when exactly one operand or the result specified by ISO/IEC/IEEE 60559:2011 is an IEEE infinity.
- 22 4 The [inquiry function](#) `IEEE_SUPPORT_SUBNORMAL` is provided to inquire whether the processor supports  
23 subnormal numbers. Where these are supported, the result of the intrinsic operations `+`, `-`, and `*`, and the  
24 functions [IEEE\\_REM](#) and [IEEE\\_RINT](#) from the intrinsic module `IEEE_ARITHMETIC`, shall conform to  
25 ISO/IEC/IEEE 60559:2011 when the result specified by ISO/IEC/IEEE 60559:2011 is subnormal, or any op-  
26 erand is subnormal and either the result is not an IEEE infinity or [IEEE\\_SUPPORT\\_INF](#) is true.
- 27 5 The [inquiry function](#) `IEEE_SUPPORT_DIVIDE` is provided to inquire whether, on kinds of real for which  
28 [IEEE\\_SUPPORT\\_DATATYPE](#) returns true, the intrinsic division operation conforms to ISO/IEC/IEEE  
29 60559:2011 when both operands and the result specified by ISO/IEC/IEEE 60559:2011 are normal numbers.  
30 If [IEEE\\_SUPPORT\\_NAN](#) is also true for a particular kind of real, the intrinsic division operation on that kind  
31 conforms to ISO/IEC/IEEE 60559:2011 when the result specified by ISO/IEC/IEEE 60559:2011 is a NaN. If  
32 [IEEE\\_SUPPORT\\_INF](#) is also true for a particular kind of real, the intrinsic division operation on that kind  
33 conforms to ISO/IEC/IEEE 60559:2011 when one operand or the result specified by ISO/IEC/IEEE 60559:2011 is  
34 an IEEE infinity. If [IEEE\\_SUPPORT\\_SUBNORMAL](#) is also true for a particular kind of real, the intrinsic divi-  
35 sion operation on that kind conforms to ISO/IEC/IEEE 60559:2011 when the result specified by ISO/IEC/IEEE  
36 60559:2011 is subnormal, or when any operand is subnormal and either the result specified by ISO/IEC/IEEE  
37 60559:2011 is not an infinity or [IEEE\\_SUPPORT\\_INF](#) is true.
- 38 6 ISO/IEC/IEEE 60559:2011 specifies a square root function that returns negative real zero for the square root of  
39 negative real zero and has certain accuracy requirements. The [inquiry function](#) `IEEE_SUPPORT_SQRT` can  
40 be used to inquire whether the intrinsic function [SQRT](#) conforms to ISO/IEC/IEEE 60559:2011 for a particular  
41 kind of real. If [IEEE\\_SUPPORT\\_NAN](#) is also true for a particular kind of real, the intrinsic function [SQRT](#)  
42 on that kind conforms to ISO/IEC/IEEE 60559:2011 when the result specified by ISO/IEC/IEEE 60559:2011  
43 is a NaN. If [IEEE\\_SUPPORT\\_INF](#) is also true for a particular kind of real, the intrinsic function [SQRT](#)  
44 on that kind conforms to ISO/IEC/IEEE 60559:2011 when the result specified by ISO/IEC/IEEE 60559:2011 is an  
45 IEEE infinity. If [IEEE\\_SUPPORT\\_SUBNORMAL](#) is also true for a particular kind of real, the intrinsic function  
46 [SQRT](#) on that kind conforms to ISO/IEC/IEEE 60559:2011 when the argument is subnormal.
- 47 7 The [inquiry function](#) `IEEE_SUPPORT_STANDARD` is provided to inquire whether the processor supports all  
48 the ISO/IEC/IEEE 60559:2011 facilities defined in this document for a particular kind of real.

## 17.10 Summary of the procedures

- 1 For all of the procedures defined in the modules, the arguments shown are the names that shall be used for **argument keywords** if the keyword form is used for the **actual arguments**.
- 2 A procedure classified in 17.10 as an **inquiry function** depends on the properties of one or more of its arguments instead of their values; in fact, these argument values may be undefined. Unless the description of one of these **inquiry functions** states otherwise, these arguments are permitted to be unallocated **allocatable** variables or pointers that are undefined or **disassociated**. A procedure that is classified as a **transformational function** is neither an **inquiry function** nor **elemental**.
- 3 In the Class column of Tables 17.2 and 17.3,
- E indicates that the procedure is an **elemental** function,
  - ES indicates that the procedure is an **elemental** subroutine,
  - I indicates that the procedure is an **inquiry function**,
  - PS indicates that the procedure is a **pure** subroutine,
  - S indicates that the procedure is an impure subroutine, and
  - T indicates that the procedure is a **transformational function**.

Table 17.2: **IEEE\_ARITHMETIC** module procedure summary

| Procedure               | Arguments                  | Class | Description                             |
|-------------------------|----------------------------|-------|-----------------------------------------|
| IEEE_CLASS              | (X)                        | E     | Classify number.                        |
| IEEE_COPY_SIGN          | (X, Y)                     | E     | Copy sign.                              |
| IEEE_FMA                | (A, B, C)                  | E     | Fused multiply-add operation.           |
| IEEE_GET_ROUNDING_MODE  | (ROUND_VALUE<br>[, RADIX]) | S     | Get rounding mode.                      |
| IEEE_GET_UNDERFLOW_MODE | (GRADUAL)                  | S     | Get <b>underflow mode</b> .             |
| IEEE_INT                | (A, ROUND [, KIND])        | E     | Conversion to integer type.             |
| IEEE_IS_FINITE          | (X)                        | E     | Whether a value is finite.              |
| IEEE_IS_NAN             | (X)                        | E     | Whether a value is an <b>IEEE NaN</b> . |
| IEEE_IS_NEGATIVE        | (X)                        | E     | Whether a value is negative.            |
| IEEE_IS_NORMAL          | (X)                        | E     | Whether a value is a normal number.     |
| IEEE_LOGB               | (X)                        | E     | Exponent.                               |
| IEEE_MAX_NUM            | (X, Y)                     | E     | Maximum numeric value.                  |
| IEEE_MAX_NUM_MAG        | (X, Y)                     | E     | Maximum magnitude numeric value.        |
| IEEE_MIN_NUM            | (X, Y)                     | E     | Minimum numeric value.                  |
| IEEE_MIN_NUM_MAG        | (X, Y)                     | E     | Minimum magnitude numeric value.        |
| IEEE_NEXT_AFTER         | (X, Y)                     | E     | Adjacent machine number.                |
| IEEE_NEXT_DOWN          | (X)                        | E     | Adjacent lower machine number.          |
| IEEE_NEXT_UP            | (X)                        | E     | Adjacent higher machine number.         |
| IEEE_QUIET_EQ           | (A, B)                     | E     | Quiet compares equal.                   |
| IEEE_QUIET_GE           | (A, B)                     | E     | Quiet compares greater than or equal.   |
| IEEE_QUIET_GT           | (A, B)                     | E     | Quiet compares greater than.            |
| IEEE_QUIET_LE           | (A, B)                     | E     | Quiet compares less than or equal.      |
| IEEE_QUIET_LT           | (A, B)                     | E     | Quiet compares less than.               |
| IEEE_QUIET_NE           | (A, B)                     | E     | Quiet compares not equal.               |
| IEEE_REAL               | (A [, KIND])               | E     | Conversion to real type.                |
| IEEE_REM                | (X, Y)                     | E     | Exact remainder.                        |
| IEEE_RINT               | (X)                        | E     | Round to integer.                       |
| IEEE_SCALB              | (X, I)                     | E     | $X \times 2^I$ .                        |
| IEEE_SELECTED_REAL_KIND | ([P, R, RADIX])            | T     | IEEE kind type parameter value.         |
| IEEE_SET_ROUNDING_MODE  | (ROUND_VALUE<br>[, RADIX]) | S     | Set rounding mode.                      |

Table 17.2: IEEE\_ARITHMETIC module procedure summary (cont.)

| Procedure                      | Arguments              | Class | Description                                      |
|--------------------------------|------------------------|-------|--------------------------------------------------|
| IEEE_SET_UNDERFLOW_MODE        | (GRADUAL)              | S     | Set <a href="#">underflow mode</a> .             |
| IEEE_SIGNALING_EQ              | (A, B)                 | E     | Signaling compares equal.                        |
| IEEE_SIGNALING_GE              | (A, B)                 | E     | Signaling compares greater than or equal.        |
| IEEE_SIGNALING_GT              | (A, B)                 | E     | Signaling compares greater than.                 |
| IEEE_SIGNALING_LE              | (A, B)                 | E     | Signaling compares less than or equal.           |
| IEEE_SIGNALING_LT              | (A, B)                 | E     | Signaling compares less than.                    |
| IEEE_SIGNALING_NE              | (A, B)                 | E     | Signaling compares not equal.                    |
| IEEE_SIGNBIT                   | (X)                    | E     | Test sign bit.                                   |
| IEEE_SUPPORT_DATATYPE          | ([X])                  | I     | Query IEEE arithmetic support.                   |
| IEEE_SUPPORT_DENORMAL          | ([X])                  | I     | Query subnormal number support.                  |
| IEEE_SUPPORT_DIVIDE            | ([X])                  | I     | Query IEEE division support.                     |
| IEEE_SUPPORT_INF               | ([X])                  | I     | Query IEEE infinity support.                     |
| IEEE_SUPPORT_IO                | ([X])                  | I     | Query IEEE formatting support.                   |
| IEEE_SUPPORT_NAN               | ([X])                  | I     | Query <a href="#">IEEE NaN</a> support.          |
| IEEE_SUPPORT_ROUNDING          | (ROUND_VALUE<br>[, X]) | T     | Query IEEE rounding support.                     |
| IEEE_SUPPORT_SQRT              | ([X])                  | I     | Query IEEE square root support.                  |
| IEEE_SUPPORT_SUBNORMAL         | ([X])                  | I     | Query subnormal number support.                  |
| IEEE_SUPPORT_STANDARD          | ([X])                  | I     | Query IEEE standard support.                     |
| IEEE_SUPPORT_UNDERFLOW_CONTROL | ([X])                  | I     | Query <a href="#">underflow</a> control support. |
| IEEE_UNORDERED                 | (X, Y)                 | E     | Whether two values are unordered.                |
| IEEE_VALUE                     | (X, CLASS)             | E     | Return number in a class.                        |

Table 17.3: IEEE\_EXCEPTIONS module procedure summary

| Procedure             | Arguments          | Class | Description                                 |
|-----------------------|--------------------|-------|---------------------------------------------|
| IEEE_GET_FLAG         | (FLAG, FLAG_VALUE) | ES    | Get an exception flag.                      |
| IEEE_GET_HALTING_MODE | (FLAG, HALTING)    | ES    | Get a <a href="#">halting mode</a> .        |
| IEEE_GET_MODES        | (MODES)            | S     | Get floating-point modes.                   |
| IEEE_GET_STATUS       | (STATUS_VALUE)     | S     | Get floating-point status.                  |
| IEEE_SET_FLAG         | (FLAG, FLAG_VALUE) | PS    | Set an exception flag.                      |
| IEEE_SET_HALTING_MODE | (FLAG, HALTING)    | PS    | Set a <a href="#">halting mode</a> .        |
| IEEE_SET_MODES        | (MODES)            | S     | Set floating-point modes.                   |
| IEEE_SET_STATUS       | (STATUS_VALUE)     | S     | Restore floating-point status.              |
| IEEE_SUPPORT_FLAG     | (FLAG [, X])       | T     | Query exception support.                    |
| IEEE_SUPPORT_HALTING  | (FLAG)             | T     | Query <a href="#">halting mode</a> support. |

1 4 In the intrinsic module IEEE\_ARITHMETIC, the [elemental](#) functions listed are provided for all reals X and Y.

## 2 17.11 Specifications of the procedures

### 3 17.11.1 General

4 1 In the detailed descriptions in 17.11, procedure names are generic and are not specific. All the functions are pure  
5 and all the subroutines are impure unless otherwise stated. All dummy arguments have [INTENT \(IN\)](#) if the  
6 intent is not stated explicitly. In the examples, it is assumed that the processor supports IEEE arithmetic for  
7 default real.

8 2 For the [elemental](#) functions of IEEE\_ARITHMETIC that return a floating-point result, if X or Y has a value

1 that is an infinity or a NaN, the result shall be consistent with the general rules in 6.1 and 6.2 of ISO/IEC/IEEE  
 2 60559:2011. For example, the result for an infinity shall be constructed as the limiting case of the result with a  
 3 value of arbitrarily large magnitude, if such a limit exists.

4 3 A program may contain statements that, if executed, would violate the requirements listed in a **Restriction**  
 5 paragraph.

#### NOTE 17.10

A program can avoid violating those requirements by using **IF constructs** to check whether particular features are supported. For example,

```
IF (IEEE_SUPPORT_DATATYPE (X)) THEN
  C = IEEE_CLASS (X)
ELSE
  ...
END IF
```

avoids invoking **IEEE\_CLASS** except on a processor which supports that facility.

### 6 17.11.2 IEEE\_CLASS (X)

7 1 **Description.** Classify number.

8 2 **Class.** **Elemental** function.

9 3 **Argument.** X shall be of type real.

10 4 **Restriction.** IEEE\_CLASS (X) shall not be invoked if IEEE\_SUPPORT\_DATATYPE (X) has the value false.

11 5 **Result Characteristics.** IEEE\_CLASS\_TYPE.

12 6 **Result Value.** The result value shall be IEEE\_SIGNALING\_NAN or IEEE\_QUIET\_NAN if IEEE\_SUPPORT\_NAN (X) has the value true and the value of X is a signaling or quiet NaN, respectively. The result value shall be IEEE\_NEGATIVE\_INF or IEEE\_POSITIVE\_INF if IEEE\_SUPPORT\_INF (X) has the value true and the value of X is negative or positive infinity, respectively. The result value shall be IEEE\_NEGATIVE\_SUBNORMAL or IEEE\_POSITIVE\_SUBNORMAL if IEEE\_SUPPORT\_SUBNORMAL (X) has the value true and the value of X is a negative or positive subnormal value, respectively. The result value shall be IEEE\_NEGATIVE\_NORMAL, IEEE\_NEGATIVE\_ZERO, IEEE\_POSITIVE\_ZERO, or IEEE\_POSITIVE\_NORMAL if the value of X is negative normal, negative zero, positive zero, or positive normal, respectively. Otherwise, the result value shall be IEEE\_OTHER\_VALUE.

21 7 **Example.** IEEE\_CLASS (−1.0) has the value IEEE\_NEGATIVE\_NORMAL.

#### NOTE 17.11

The result value IEEE\_OTHER\_VALUE is useful on systems that are almost IEEE-compatible, but do not implement all of it. For example, if a subnormal value is encountered on a system that does not support them.

### 22 17.11.3 IEEE\_COPY\_SIGN (X, Y)

23 1 **Description.** Copy sign.

24 2 **Class.** **Elemental** function.

25 3 **Arguments.** The arguments shall be of type real.

1 4 **Restriction.** IEEE\_COPY\_SIGN (X, Y) shall not be invoked if IEEE\_SUPPORT\_DATATYPE (X) or  
2 IEEE\_SUPPORT\_DATATYPE (Y) has the value false.

3 5 **Result Characteristics.** Same as X.

4 6 **Result Value.** The result has the absolute value of X with the sign of Y. This is true even for IEEE special  
5 values, such as a NaN or an infinity (on processors supporting such values).

6 7 **Example.** The value of IEEE\_COPY\_SIGN (X, 1.0) is ABS (X) even when X is a NaN.

#### 7 17.11.4 IEEE\_FMA (A, B, C)

8 1 **Description.** Fused multiply-add operation.

9 2 **Class.** Elemental function.

10 3 **Arguments.**

11 A shall be of type real.

12 B shall be of the same type and kind type parameter as A.

13 C shall be of the same type and kind type parameter as A.

14 4 **Restriction.** IEEE\_FMA (A, B, C) shall not be invoked if IEEE\_SUPPORT\_DATATYPE (A) has the value  
15 false.

16 5 **Result Characteristics.** Same as A.

17 6 **Result Value.** The result has the value specified by ISO/IEC/IEEE 60559:2011 for the fusedMultiplyAdd  
18 operation; that is, when the result is in range, its value is equal to the mathematical value of  $(A \times B) + C$  rounded  
19 to the representation method of A according to the rounding mode. IEEE\_OVERFLOW, IEEE\_UNDERFLOW,  
20 and IEEE\_INEXACT shall be signaled according to the final step in the calculation and not by any intermediate  
21 calculation.

22 7 **Example.** The value of IEEE\_FMA (TINY (0.0), TINY (0.0), 1.0), when the rounding mode is IEEE\_  
23 NEAREST, is equal to 1.0; only the IEEE\_INEXACT exception is signaled.

#### 24 17.11.5 IEEE\_GET\_FLAG (FLAG, FLAG\_VALUE)

25 1 **Description.** Get an exception flag.

26 2 **Class.** Elemental subroutine.

27 3 **Arguments.**

28 FLAG shall be of type IEEE\_FLAG\_TYPE. It specifies the exception flag to be obtained.

29 FLAG\_VALUE shall be of type logical. It is an INTENT (OUT) argument. If the value of FLAG is IEEE\_  
30 INVALID, IEEE\_OVERFLOW, IEEE\_DIVIDE\_BY\_ZERO, IEEE\_UNDERFLOW, or IEEE\_  
31 INEXACT, FLAG\_VALUE is assigned the value true if the corresponding exception flag is signaling  
32 and is assigned the value false otherwise.

33 4 **Example.** Following CALL IEEE\_GET\_FLAG (IEEE\_OVERFLOW, FLAG\_VALUE), FLAG\_VALUE is  
34 true if the IEEE\_OVERFLOW flag is signaling and is false if it is quiet.

#### 35 17.11.6 IEEE\_GET\_HALTING\_MODE (FLAG, HALTING)

36 1 **Description.** Get a halting mode.

37 2 **Class.** Elemental subroutine.

38 3 **Arguments.**

1 FLAG shall be of type `IEEE_FLAG_TYPE`. It specifies the exception flag. It shall have one of the values `IEEE_INVALID`, `IEEE_OVERFLOW`, `IEEE_DIVIDE_BY_ZERO`, `IEEE_UNDERFLOW`,  
 2 or `IEEE_INEXACT`.  
 3  
 4 HALTING shall be of type logical. It is an `INTENT (OUT)` argument. It is assigned the value true if the  
 5 exception specified by FLAG will cause halting. Otherwise, it is assigned the value false.

6 4 **Example.** To store the `halting mode` for `IEEE_OVERFLOW`, do a calculation without halting, and restore the  
 7 `halting mode` later:

```
8     USE, INTRINSIC :: IEEE_ARITHMETIC
9     LOGICAL HALTING
10    ...
11    CALL IEEE_GET_HALTING_MODE (IEEE_OVERFLOW, HALTING) ! Store halting mode
12    CALL IEEE_SET_HALTING_MODE (IEEE_OVERFLOW, .FALSE.) ! No halting
13    ... ! calculation without halting
14    CALL IEEE_SET_HALTING_MODE (IEEE_OVERFLOW, HALTING) ! Restore halting mode
```

### 15 17.11.7 IEEE\_GET\_MODES (MODES)

16 1 **Description.** Get floating-point modes.

17 2 **Class.** Subroutine.

18 3 **Argument.** MODES shall be a scalar of type `IEEE_MODES_TYPE`. It is an `INTENT (OUT)` argument that  
 19 is assigned the value of the floating-point modes.

20 4 **Example.** To save the floating-point modes, do a calculation with specific rounding and `underflow` modes, and  
 21 restore them later:

```
22     USE, INTRINSIC :: IEEE_ARITHMETIC
23     TYPE (IEEE_MODES_TYPE) SAVE_MODES
24     ...
25     CALL IEEE_GET_MODES (SAVE_MODES) ! Save all modes.
26     CALL IEEE_SET_ROUNDING_MODE (IEEE_TO_ZERO)
27     CALL IEEE_SET_UNDERFLOW_MODE (GRADUAL=.FALSE.)
28     ... ! calculation with abrupt round-to-zero.
29     CALL IEEE_SET_MODES (SAVE_MODES) ! Restore all modes.
```

### 30 17.11.8 IEEE\_GET\_ROUNDING\_MODE (ROUND\_VALUE [, RADIX])

31 1 **Description.** Get rounding mode.

32 2 **Class.** Subroutine.

33 3 **Arguments.**

34 `ROUND_VALUE` shall be a scalar of type `IEEE_ROUND_TYPE`. It is an `INTENT (OUT)` argument. It is  
 35 assigned the value `IEEE_NEAREST`, `IEEE_TO_ZERO`, `IEEE_UP`, `IEEE_DOWN`, or `IEEE_-`  
 36 `AWAY` if the corresponding rounding mode is in operation and `IEEE_OTHER` otherwise.

37 `RADIX` (optional) shall be an integer scalar with the value two or ten. If `RADIX` is present with the value ten,  
 38 the rounding mode queried is the decimal rounding mode, otherwise it is the binary rounding mode.

39 4 **Example.** To save the binary rounding mode, do a calculation with round to nearest, and restore the rounding  
 40 mode later:

```

1      USE, INTRINSIC :: IEEE_ARITHMETIC
2      TYPE (IEEE_ROUND_TYPE) ROUND_VALUE
3      ...
4      CALL IEEE_GET_ROUNDING_MODE (ROUND_VALUE) ! Store the rounding mode
5      CALL IEEE_SET_ROUNDING_MODE (IEEE_NEAREST)
6      ... ! calculation with round to nearest
7      CALL IEEE_SET_ROUNDING_MODE (ROUND_VALUE) ! Restore the rounding mode

```

### 8 17.11.9 IEEE\_GET\_STATUS (STATUS\_VALUE)

9 1 **Description.** Get floating-point status.

10 2 **Class.** Subroutine.

11 3 **Argument.** STATUS\_VALUE shall be a scalar of type IEEE\_STATUS\_TYPE. It is an INTENT (OUT)  
12 argument. It is assigned the value of the floating-point status.

13 4 **Example.** To store all the exception flags, do a calculation involving exception handling, and restore them later:

14

```

15      USE, INTRINSIC :: IEEE_ARITHMETIC
16      TYPE (IEEE_STATUS_TYPE) STATUS_VALUE
17      ...
18      CALL IEEE_GET_STATUS (STATUS_VALUE) ! Get the flags
19      CALL IEEE_SET_FLAG (IEEE_ALL, .FALSE.) ! Set the flags quiet.
20      ... ! calculation involving exception handling
21      CALL IEEE_SET_STATUS (STATUS_VALUE) ! Restore the flags

```

### 22 17.11.10 IEEE\_GET\_UNDERFLOW\_MODE (GRADUAL)

23 1 **Description.** Get underflow mode.

24 2 **Class.** Subroutine.

25 3 **Argument.** GRADUAL shall be a logical scalar. It is an INTENT (OUT) argument. It is assigned the value  
26 true if the underflow mode is gradual underflow, and false if the underflow mode is abrupt underflow.

27 4 **Restriction.** IEEE\_GET\_UNDERFLOW\_MODE shall not be invoked unless IEEE\_SUPPORT\_UNDER-  
28 FLOW\_CONTROL (X) is true for some X.

29 5 **Example.** After CALL IEEE\_SET\_UNDERFLOW\_MODE (.FALSE.), a subsequent CALL IEEE\_GET\_-  
30 UNDERFLOW\_MODE (GRADUAL) will set GRADUAL to false.

### 31 17.11.11 IEEE\_INT (A, ROUND [, KIND])

32 1 **Description.** Conversion to integer type.

33 2 **Class.** Elemental function.

34 3 **Arguments.**

35 A shall be of type real.

36 ROUND shall be of type IEEE\_ROUND\_TYPE.

37 KIND (optional) shall be a scalar integer constant expression.

1 4 **Restriction.** IEEE\_INT (A, ROUND, KIND) shall not be invoked if IEEE\_SUPPORT\_DATATYPE (A) has  
2 the value false.

3 5 **Result Characteristics.** Integer. If KIND is present, the kind type parameter is that specified by the value of  
4 KIND; otherwise, the kind type parameter is that of default integer.

5 6 **Result Value.** The result has the value specified by ISO/IEC/IEEE 60559:2011 for the convertToInteger{round}  
6 or the convertToIntegerExact{round} operation; the processor shall consistently choose which operation it  
7 provides. That is, the value of A is converted to an integer according to the rounding mode specified by ROUND;  
8 if this value is representable in the representation method of the result, the result has this value, otherwise IEEE\_  
9 INVALID is signaled and the result is processor dependent. If the processor provides the convertToIntegerExact  
10 operation, IEEE\_INVALID did not signal, and the value of the result differs from that of A, IEEE\_INEXACT  
11 will be signaled.

12 7 **Example.** The value of IEEE\_INT (12.5, IEEE\_UP) is 13; IEEE\_INEXACT will be signaled if the processor  
13 provides the convertToIntegerExact operation.

#### 14 17.11.12 IEEE\_IS\_FINITE (X)

15 1 **Description.** Whether a value is finite.

16 2 **Class.** Elemental function.

17 3 **Argument.** X shall be of type real.

18 4 **Restriction.** IEEE\_IS\_FINITE (X) shall not be invoked if IEEE\_SUPPORT\_DATATYPE (X) has the value  
19 false.

20 5 **Result Characteristics.** Default logical.

21 6 **Result Value.** The result has the value true if the value of X is finite, that is, IEEE\_CLASS (X) has one  
22 of the values IEEE\_NEGATIVE\_NORMAL, IEEE\_NEGATIVE\_SUBNORMAL, IEEE\_NEGATIVE\_ZERO,  
23 IEEE\_POSITIVE\_ZERO, IEEE\_POSITIVE\_SUBNORMAL, or IEEE\_POSITIVE\_NORMAL; otherwise, the  
24 result has the value false.

25 7 **Example.** IEEE\_IS\_FINITE (1.0) has the value true.

#### 26 17.11.13 IEEE\_IS\_NAN (X)

27 1 **Description.** Whether a value is an IEEE NaN.

28 2 **Class.** Elemental function.

29 3 **Argument.** X shall be of type real.

30 4 **Restriction.** IEEE\_IS\_NAN (X) shall not be invoked if IEEE\_SUPPORT\_NAN (X) has the value false.

31 5 **Result Characteristics.** Default logical.

32 6 **Result Value.** The result has the value true if the value of X is an IEEE NaN; otherwise, it has the value false.

33 7 **Example.** IEEE\_IS\_NAN (SQRT (-1.0)) has the value true if IEEE\_SUPPORT\_SQRT (1.0) has the value  
34 true.

#### 35 17.11.14 IEEE\_IS\_NEGATIVE (X)

36 1 **Description.** Whether a value is negative.

37 2 **Class.** Elemental function.



1 3 **Argument.** X shall be of type real.

2 4 **Restriction.** IEEE\_IS\_NEGATIVE (X) shall not be invoked if IEEE\_SUPPORT\_DATATYPE (X) has the  
3 value false.

4 5 **Result Characteristics.** Default logical.

5 6 **Result Value.** The result has the value true if IEEE\_CLASS (X) has one of the values IEEE\_NEGATIVE\_-  
6 NORMAL, IEEE\_NEGATIVE\_SUBNORMAL, IEEE\_NEGATIVE\_ZERO or IEEE\_NEGATIVE\_INF; oth-  
7 erwise, the result has the value false.

8 7 **Example.** IEEE\_IS\_NEGATIVE (0.0) has the value false.

### 9 17.11.15 IEEE\_IS\_NORMAL (X)

10 1 **Description.** Whether a value is a normal number.

11 2 **Class.** Elemental function.

12 3 **Argument.** X shall be of type real.

13 4 **Restriction.** IEEE\_IS\_NORMAL (X) shall not be invoked if IEEE\_SUPPORT\_DATATYPE (X) has the  
14 value false.

15 5 **Result Characteristics.** Default logical.

16 6 **Result Value.** The result has the value true if IEEE\_CLASS (X) has one of the values IEEE\_NEGATIVE\_-  
17 NORMAL, IEEE\_NEGATIVE\_ZERO, IEEE\_POSITIVE\_ZERO or IEEE\_POSITIVE\_NORMAL; otherwise,  
18 the result has the value false.

19 7 **Example.** IEEE\_IS\_NORMAL (SQRT (-1.0)) has the value false if IEEE\_SUPPORT\_SQRT (1.0) has the  
20 value true.

### 21 17.11.16 IEEE\_LOGB (X)

22 1 **Description.** Exponent.

23 2 **Class.** Elemental function.

24 3 **Argument.** X shall be of type real.

25 4 **Restriction.** IEEE\_LOGB (X) shall not be invoked if IEEE\_SUPPORT\_DATATYPE (X) has the value false.

26 5 **Result Characteristics.** Same as X.

27 6 **Result Value.**

28 *Case (i):* If the value of X is neither zero, infinity, nor NaN, the result has the value of the unbiased exponent  
29 of X. Note: this value is equal to EXPONENT (X) - 1.

30 *Case (ii):* If X==0, the result is -infinity if IEEE\_SUPPORT\_INF (X) is true and -HUGE (X) otherwise;  
31 IEEE\_DIVIDE\_BY\_ZERO signals.

32 *Case (iii):* If IEEE\_SUPPORT\_INF (X) is true and X is infinite, the result is +infinity.

33 *Case (iv):* If IEEE\_SUPPORT\_NAN (X) is true and X is a NaN, the result is a NaN.

34 7 **Example.** IEEE\_LOGB (-1.1) has the value 0.0.

### 17.11.17 IEEE\_MAX\_NUM (X, Y)

1 **Description.** Maximum numeric value.

2 **Class.** [Elemental](#) function.

3 **Arguments.**

4 X shall be of type real.

5 Y shall be of the same type and kind type parameter as X.

6 **Restriction.** IEEE\_MAX\_NUM shall not be invoked if [IEEE\\_SUPPORT\\_DATATYPE](#) (X) has the value false.

7 **Result Characteristics.** Same as X.

8 **Result Value.** The result has the value specified for the maxNum operation in ISO/IEC/IEEE 60559:2011; that is,

- 9 • if  $X < Y$  the result has the value of Y;
- 10 • if  $Y < X$  the result has the value of X;
- 11 • if exactly one of X and Y is a quiet NaN the result has the value of the other argument;
- 12 • if one or both of X and Y are signaling NaNs, IEEE\_INVALID signals and the result is a NaN;
- 13 • otherwise, the result is either X or Y (processor dependent).

14 Except when X or Y is a signaling NaN, no exception is signaled.

15 **Example.** The value of IEEE\_MAX\_NUM (1.5, [IEEE\\_VALUE](#) (IEEE\_QUIET\_NAN)) is 1.5.

### 17.11.18 IEEE\_MAX\_NUM\_MAG (X, Y)

1 **Description.** Maximum magnitude numeric value.

2 **Class.** [Elemental](#) function.

3 **Arguments.**

4 X shall be of type real.

5 Y shall be of the same type and kind type parameter as X.

6 **Restriction.** IEEE\_MAX\_NUM\_MAG shall not be invoked if [IEEE\\_SUPPORT\\_DATATYPE](#) (X) has the value false.

7 **Result Characteristics.** Same as X.

8 **Result Value.** The result has the value specified for the maxNumMag operation in ISO/IEC/IEEE 60559:2011; that is,

- 9 • if  $ABS(X) < ABS(Y)$  the result has the value of Y;
- 10 • if  $ABS(Y) < ABS(X)$  the result has the value of X;
- 11 • otherwise, the result has the value of [IEEE\\_MAX\\_NUM](#) (X, Y).

12 Except when X or Y is a signaling NaN, no exception is signaled.

13 **Example.** The value of IEEE\_MAX\_NUM\_MAG (1.5, -2.5) is -2.5.

### 17.11.19 IEEE\_MIN\_NUM (X, Y)

1 **Description.** Minimum numeric value.

2 **Class.** [Elemental](#) function.

3 **Arguments.**

4 X shall be of type real.

5 Y shall be of the same type and kind type parameter as X.

6 **Restriction.** IEEE\_MIN\_NUM shall not be invoked if [IEEE\\_SUPPORT\\_DATATYPE](#) (X) has the value false.

7 **Result Characteristics.** Same as X.

8 **Result Value.** The result has the value specified for the minNum operation in ISO/IEC/IEEE 60559:2011; that is,

- 9 • if  $X < Y$  the result has the value of X;
- 10 • if  $Y < X$  the result has the value of Y;
- 11 • if exactly one of X and Y is a quiet NaN the result has the value of the other argument;
- 12 • if one or both of X and Y are signaling NaNs, [IEEE\\_INVALID](#) signals and the result is a NaN;
- 13 • otherwise, the result is either X or Y (processor dependent).

14 Except when X or Y is a signaling NaN, no exception is signaled.

15 **Example.** The value of IEEE\_MIN\_NUM (1.5, [IEEE\\_VALUE](#) (IEEE\_QUIET\_NAN)) is 1.5.

### 17.11.20 IEEE\_MIN\_NUM\_MAG (X, Y)

1 **Description.** Minimum magnitude numeric value.

2 **Class.** [Elemental](#) function.

3 **Arguments.**

4 X shall be of type real.

5 Y shall be of the same type and kind type parameter as X.

6 **Restriction.** IEEE\_MIN\_NUM\_MAG shall not be invoked if [IEEE\\_SUPPORT\\_DATATYPE](#) (X) has the value false.

7 **Result Characteristics.** Same as X.

8 **Result Value.** The result has the value specified for the minNumMag operation in ISO/IEC/IEEE 60559:2011; that is,

- 9 • if  $ABS(X) < ABS(Y)$  the result has the value of X;
- 10 • if  $ABS(Y) < ABS(X)$  the result has the value of Y;
- 11 • otherwise, the result has the value of [IEEE\\_MIN\\_NUM](#) (X, Y).

12 Except when X or Y is a signaling NaN, no exception is signaled.

13 **Example.** The value of IEEE\_MIN\_NUM\_MAG (1.5, -2.5) is 1.5.

### 17.11.21 IEEE\_NEXT\_AFTER (X, Y)

1 **Description.** Adjacent machine number.

2 **Class.** [Elemental](#) function.

1 3 **Arguments.** The arguments shall be of type real.

2 4 **Restriction.** IEEE\_NEXT\_AFTER (X, Y) shall not be invoked if IEEE\_SUPPORT\_DATATYPE (X) or  
3 IEEE\_SUPPORT\_DATATYPE (Y) has the value false.

4 5 **Result Characteristics.** Same as X.

5 6 **Result Value.**

6 *Case (i):* If  $X == Y$ , the result is X and no exception is signaled.

7 *Case (ii):* If  $X \neq Y$ , the result has the value of the next representable neighbor of X in the direction of Y.  
8 The neighbors of zero (of either sign) are both nonzero. IEEE\_OVERFLOW is signaled when  
9 X is finite but IEEE\_NEXT\_AFTER (X, Y) is infinite; IEEE\_UNDERFLOW is signaled when  
10 IEEE\_NEXT\_AFTER (X, Y) is subnormal; in both cases, IEEE\_INEXACT signals.

11 7 **Example.** The value of IEEE\_NEXT\_AFTER (1.0, 2.0) is  $1.0 + \text{EPSILON}(X)$ .

### 12 17.11.22 IEEE\_NEXT\_DOWN (X)

13 1 **Description.** Adjacent lower machine number.

14 2 **Class.** Elemental function.

15 3 **Argument.** X shall be of type real.

16 4 **Restriction.** IEEE\_NEXT\_DOWN (X) shall not be invoked if IEEE\_SUPPORT\_DATATYPE (X) has the  
17 value false. IEEE\_NEXT\_DOWN ( $-\text{HUGE}(X)$ ) shall not be invoked if IEEE\_SUPPORT\_INF (X) has the  
18 value false.

19 5 **Result Characteristics.** Same as X.

20 6 **Result Value.** The result has the value specified for the nextDown operation in ISO/IEC/IEEE 60559:2011;  
21 that is, it is the greatest value in the representation method of X that compares less than X, except when X is  
22 equal to  $-\infty$  the result has the value  $-\infty$ , and when X is a NaN the result is a NaN. If X is a signaling NaN,  
23 IEEE\_INVALID signals; otherwise, no exception is signaled.

24 7 **Example.** If IEEE\_SUPPORT\_SUBNORMAL (0.0) is true, the value of IEEE\_NEXT\_DOWN (+0.0) is the  
25 negative subnormal number with least magnitude.

### 26 17.11.23 IEEE\_NEXT\_UP (X)

27 1 **Description.** Adjacent higher machine number.

28 2 **Class.** Elemental function.

29 3 **Argument.** X shall be of type real.

30 4 **Restriction.** IEEE\_NEXT\_UP (X) shall not be invoked if IEEE\_SUPPORT\_DATATYPE (X) has the value  
31 false. IEEE\_NEXT\_UP ( $\text{HUGE}(X)$ ) shall not be invoked if IEEE\_SUPPORT\_INF (X) has the value false.

32 5 **Result Characteristics.** Same as X.

33 6 **Result Value.** The result has the value specified for the nextUp operation in ISO/IEC/IEEE 60559:2011; that  
34 is, it is the least value in the representation method of X that compares greater than X, except when X is  
35 equal to  $+\infty$  the result has the value  $+\infty$ , and when X is a NaN the result is a NaN. If X is a signaling NaN,  
36 IEEE\_INVALID signals; otherwise, no exception is signaled.

37 7 **Example.** If IEEE\_SUPPORT\_INF (X) is true, the value of IEEE\_NEXT\_UP ( $\text{HUGE}(X)$ ) is  $+\infty$ .

### 17.11.24 IEEE\_QUIET\_EQ (A, B)

1 **Description.** Quiet compares equal.

2 **Class.** [Elemental](#) function.

3 **Arguments.**

4 A shall be of type real.

5 B shall have the same type and kind type parameter as A.

6 **Restriction.** IEEE\_QUIET\_EQ (A, B) shall not be invoked if [IEEE\\_SUPPORT\\_DATATYPE](#) (A) has the value false.

7 **Result Characteristics.** Default logical.

8 **Result Value.** The result has the value specified for the compareQuietEqual operation in ISO/IEC/IEEE 60559:2011; that is, it is true if and only if A compares equal to B. If A or B is a NaN, the result will be false. If A or B is a signaling NaN, IEEE\_INVALID signals; otherwise, no exception is signaled.

9 **Example.** IEEE\_QUIET\_EQ (1.0, [IEEE\\_VALUE](#) (IEEE\_QUIET\_NAN)) has the value false and no exception is signaled.

### 17.11.25 IEEE\_QUIET\_GE (A, B)

10 **Description.** Quiet compares greater than or equal.

11 **Class.** [Elemental](#) function.

12 **Arguments.**

13 A shall be of type real.

14 B shall have the same type and kind type parameter as A.

15 **Restriction.** IEEE\_QUIET\_GE (A, B) shall not be invoked if [IEEE\\_SUPPORT\\_DATATYPE](#) (A) has the value false.

16 **Result Characteristics.** Default logical.

17 **Result Value.** The result has the value specified for the compareQuietGreaterEqual operation in ISO/IEC/IEEE 60559:2011; that is, it is true if and only if A compares greater than or equal to B. If A or B is a NaN, the result will be false. If A or B is a signaling NaN, IEEE\_INVALID signals; otherwise, no exception is signaled.

18 **Example.** IEEE\_QUIET\_GE (1.0, [IEEE\\_VALUE](#) (IEEE\_QUIET\_NAN)) has the value false and no exception is signaled.

### 17.11.26 IEEE\_QUIET\_GT (A, B)

19 **Description.** Quiet compares greater than.

20 **Class.** [Elemental](#) function.

21 **Arguments.**

22 A shall be of type real.

23 B shall have the same type and kind type parameter as A.

24 **Restriction.** IEEE\_QUIET\_GT (A, B) shall not be invoked if [IEEE\\_SUPPORT\\_DATATYPE](#) (A) has the value false.

25 **Result Characteristics.** Default logical.

1 6 **Result Value.** The result has the value specified for the compareQuietGreater operation in ISO/IEC/IEEE  
2 60559:2011; that is, it is true if and only if A compares greater than B. If A or B is a NaN, the result will be  
3 false. If A or B is a signaling NaN, IEEE\_INVALID signals; otherwise, no exception is signaled.

4 7 **Example.** IEEE\_QUIET\_GT (1.0, IEEE\_VALUE (IEEE\_QUIET\_NAN)) has the value false and no exception  
5 is signaled.

### 6 17.11.27 IEEE\_QUIET\_LE (A, B)

7 1 **Description.** Quiet compares less than or equal.

8 2 **Class.** Elemental function.

9 3 **Arguments.**

10 A shall be of type real.

11 B shall have the same type and kind type parameter as A.

12 4 **Restriction.** IEEE\_QUIET\_LE (A, B) shall not be invoked if IEEE\_SUPPORT\_DATATYPE (A) has the  
13 value false.

14 5 **Result Characteristics.** Default logical.

15 6 **Result Value.** The result has the value specified for the compareQuietLessEqual operation in ISO/IEC/IEEE  
16 60559:2011; that is, it is true if and only if A compares less than or equal to B. If A or B is a NaN, the result will  
17 be false. If A or B is a signaling NaN, IEEE\_INVALID signals; otherwise, no exception is signaled.

18 7 **Example.** IEEE\_QUIET\_LE (1.0, IEEE\_VALUE (IEEE\_QUIET\_NAN)) has the value false and no exception  
19 is signaled.

### 20 17.11.28 IEEE\_QUIET\_LT (A, B)

21 1 **Description.** Quiet compares less than.

22 2 **Class.** Elemental function.

23 3 **Arguments.**

24 A shall be of type real.

25 B shall have the same type and kind type parameter as A.

26 4 **Restriction.** IEEE\_QUIET\_LT (A, B) shall not be invoked if IEEE\_SUPPORT\_DATATYPE (A) has the  
27 value false.

28 5 **Result Characteristics.** Default logical.

29 6 **Result Value.** The result has the value specified for the compareQuietLess operation in ISO/IEC/IEEE  
30 60559:2011; that is, it is true if and only if A compares less than B. If A or B is a NaN, the result will be  
31 false. If A or B is a signaling NaN, IEEE\_INVALID signals; otherwise, no exception is signaled.

32 7 **Example.** IEEE\_QUIET\_LT (1.0, IEEE\_VALUE (IEEE\_QUIET\_NAN)) has the value false and no exception  
33 is signaled.

### 34 17.11.29 IEEE\_QUIET\_NE (A, B)

35 1 **Description.** Quiet compares not equal.

36 2 **Class.** Elemental function.

37 3 **Arguments.**

1 A shall be of type real.

2 B shall have the same type and kind type parameter as A.

3 4 **Restriction.** IEEE\_QUIET\_NE (A, B) shall not be invoked if IEEE\_SUPPORT\_DATATYPE (A) has the  
4 value false.

5 5 **Result Characteristics.** Default logical.

6 6 **Result Value.** The result has the value specified for the compareQuietNotEqual operation in ISO/IEC/IEEE  
7 60559:2011; that is, it is true if and only if A compares not equal to B. If A or B is a NaN, the result will be true.  
8 If A or B is a signaling NaN, IEEE\_INVALID signals; otherwise, no exception is signaled.

9 7 **Example.** IEEE\_QUIET\_NE (1.0, IEEE\_VALUE (IEEE\_QUIET\_NAN)) has the value true and no exception  
10 is signaled.

### 11 17.11.30 IEEE\_REAL (A [, KIND])

12 1 **Description.** Conversion to real type.

13 2 **Class.** Elemental function.

14 3 **Arguments.**

15 A shall be of type integer or real.

16 KIND (optional) shall be a scalar integer constant expression.

17 4 **Restriction.** IEEE\_REAL shall not be invoked if A is of type real and IEEE\_SUPPORT\_DATATYPE (A)  
18 has the value false, or if IEEE\_SUPPORT\_DATATYPE (IEEE\_REAL (A, KIND)) has the value false.

19 5 **Result Characteristics.** Real. If KIND is present, the kind type parameter is that specified by the value of  
20 KIND; otherwise, the kind type parameter is that of default real.

21 6 **Result Value.** The result has the same value as A if that value is representable in the representation method  
22 of the result, and is rounded according to the rounding mode otherwise. This shall be consistent with the  
23 specification of ISO/IEC/IEEE 60559:2011 for the convertFromInt operation when A is of type integer, and with  
24 the convertFormat operation otherwise.

25 7 **Example.** The value of IEEE\_REAL (123) is 123.0.

### 26 17.11.31 IEEE\_REM (X, Y)

27 1 **Description.** Exact remainder.

28 2 **Class.** Elemental function.

29 3 **Arguments.** The arguments shall be of type real and have the same radix.

30 4 **Restriction.** IEEE\_REM (X, Y) shall not be invoked if IEEE\_SUPPORT\_DATATYPE (X) or IEEE\_SUP-  
31 PORT\_DATATYPE (Y) has the value false.

32 5 **Result Characteristics.** Real with the kind type parameter of whichever argument has the greater precision.

33 6 **Result Value.** The result value, regardless of the rounding mode, shall be exactly  $X - Y*N$ , where N is the  
34 integer nearest to the exact value  $X/Y$ ; whenever  $|N - X/Y| = \frac{1}{2}$ , N shall be even. If the result value is zero, the  
35 sign shall be that of X. This function computes the remainder operation specified in ISO/IEC/IEEE 60559:2011.

36 7 **Examples.** The value of IEEE\_REM (4.0, 3.0) is 1.0, the value of IEEE\_REM (3.0, 2.0) is -1.0, and the value  
37 of IEEE\_REM (5.0, 2.0) is 1.0.

### 17.11.32 IEEE\_RINT (X [, ROUND])

1 **Description.** Round to integer.

2 **Class.** [Elemental](#) function.

3 **Arguments.**

4 X shall be of type real.

5 ROUND (optional) shall be of type [IEEE\\_ROUND\\_TYPE](#).

6 **Restriction.** IEEE\_RINT (X) shall not be invoked if [IEEE\\_SUPPORT\\_DATATYPE](#) (X) has the value false.

7 **Result Characteristics.** Same as X.

8 **Result Value.** If ROUND is present, the value of the result is the value of X rounded to an integer according to the mode specified by ROUND; this is the ISO/IEC/IEEE 60559:2011 operation roundToInteger{rounding}. Otherwise, the value of the result is that specified for the operation roundIntegralToExact in ISO/IEC/IEEE 60559:2011; this is the value of X rounded to an integer according to the rounding mode. If the result has the value zero, the sign is that of X.

9 **Examples.** If the rounding mode is round to nearest, the value of IEEE\_RINT (1.1) is 1.0. The value of  
10 IEEE\_RINT (1.1, IEEE\_UP) is 2.0.

### 17.11.33 IEEE\_SCALB (X, I)

11 **Description.**  $X \times 2^I$ .

12 **Class.** [Elemental](#) function.

13 **Arguments.**

14 X shall be of type real.

15 I shall be of type integer.

16 **Restriction.** IEEE\_SCALB (X) shall not be invoked if [IEEE\\_SUPPORT\\_DATATYPE](#) (X) has the value false.

17 **Result Characteristics.** Same as X.

18 **Result Value.**

19 *Case (i):* If  $X \times 2^I$  is representable as a normal number, the result has this value.

20 *Case (ii):* If X is finite and  $X \times 2^I$  is too large, the IEEE\_OVERFLOW exception shall occur. If IEEE\_SUPPORT\_INF (X) is true, the result value is infinity with the sign of X; otherwise, the result value is [SIGN \(HUGE \(X\), X\)](#).

21 *Case (iii):* If  $X \times 2^I$  is too small and there is loss of accuracy, the IEEE\_UNDERFLOW exception shall occur. The result is the representable number having a magnitude nearest to  $|2^I|$  and the same sign as X.

22 *Case (iv):* If X is infinite, the result is the same as X; no exception signals.

23 **Example.** The value of IEEE\_SCALB (1.0, 2) is 4.0.

### 17.11.34 IEEE\_SELECTED\_REAL\_KIND ([P, R, RADIX])

24 **Description.** IEEE kind type parameter value.

25 **Class.** [Transformational](#) function.

26 **Arguments.** At least one argument shall be present.

27 P (optional) shall be an integer scalar.

28 R (optional) shall be an integer scalar.



1 RADIX (optional) shall be an integer scalar.

2 4 **Result Characteristics.** Default integer scalar.

3 5 **Result Value.** If P or R is absent, the result value is the same as if it were present with the value zero. If  
 4 RADIX is absent, there is no requirement on the radix of the selected kind. The result has a value equal to a  
 5 value of the kind type parameter of an ISO/IEC/IEEE 60559:2011 floating-point format with decimal precision,  
 6 as returned by the intrinsic function [PRECISION](#), of at least P digits, a decimal exponent range, as returned  
 7 by the intrinsic function [RANGE](#), of at least R, and a radix, as returned by the intrinsic function [RADIX](#), of  
 8 RADIX, if such a kind type parameter is available on the processor.

9 6 Otherwise, the result is  $-1$  if the processor supports an IEEE real type with radix RADIX and exponent range  
 10 of at least R but not with precision of at least P,  $-2$  if the processor supports an IEEE real type with radix  
 11 RADIX and precision of at least P but not with exponent range of at least R,  $-3$  if the processor supports an  
 12 IEEE real type with radix RADIX but with neither precision of at least P nor exponent range of at least R,  $-4$  if  
 13 the processor supports an IEEE real type with radix RADIX and either precision of at least P or exponent range  
 14 of at least R but not both together, and  $-5$  if the processor supports no IEEE real type with radix RADIX.

15 7 If more than one kind type parameter value meets the criteria, the value returned is the one with the smallest  
 16 decimal precision, unless there are several such values, in which case the smallest of these kind values is returned.

17 8 **Example.** IEEE\_SELECTED\_REAL\_KIND (6, 30) has the value [KIND](#) (0.0) on a machine that supports  
 18 ISO/IEC/IEEE 60559:2011 single precision arithmetic for its default real approximation method.

### 19 17.11.35 IEEE\_SET\_FLAG (FLAG, FLAG\_VALUE)

20 1 **Description.** Set an exception flag.

21 2 **Class.** Pure subroutine.

22 3 **Arguments.**

23 FLAG shall be a scalar or array of type [IEEE\\_FLAG\\_TYPE](#). If a value of FLAG is IEEE\_INVALID,  
 24 IEEE\_OVERFLOW, IEEE\_DIVIDE\_BY\_ZERO, IEEE\_UNDERFLOW, or IEEE\_INEXACT,  
 25 the corresponding exception flag is assigned a value. No two elements of FLAG shall have the same  
 26 value.

27 FLAG\_VALUE shall be a logical scalar or array. It shall be [conformable](#) with FLAG. If an element has the value  
 28 true, the corresponding flag is set to be signaling; otherwise, the flag is set to be quiet.

29 4 **Example.** CALL IEEE\_SET\_FLAG (IEEE\_OVERFLOW, [.TRUE.](#)) sets the IEEE\_OVERFLOW flag to be  
 30 signaling.

### 31 17.11.36 IEEE\_SET\_HALTING\_MODE (FLAG, HALTING)

32 1 **Description.** Set a [halting mode](#).

33 2 **Class.** Pure subroutine.

34 3 **Arguments.**

35 FLAG shall be a scalar or array of type [IEEE\\_FLAG\\_TYPE](#). It shall have only the values IEEE\_  
 36 INVALID, IEEE\_OVERFLOW, IEEE\_DIVIDE\_BY\_ZERO, IEEE\_UNDERFLOW, or IEEE\_  
 37 INEXACT. No two elements of FLAG shall have the same value.

38 HALTING shall be a logical scalar or array. It shall be [conformable](#) with FLAG. If an element has the value  
 39 true, the corresponding exception specified by FLAG will cause halting. Otherwise, execution will  
 40 continue after this exception.

41 4 **Restriction.** IEEE\_SET\_HALTING\_MODE (FLAG, HALTING) shall not be invoked if IEEE\_SUPPORT\_  
 42 HALTING (FLAG) has the value false.

1 5 **Example.** CALL IEEE\_SET\_HALTING\_MODE (IEEE\_DIVIDE\_BY\_ZERO, .TRUE.) causes halting after  
2 a divide\_by\_zero exception.

### 3 17.11.37 IEEE\_SET\_MODES (MODES)

4 1 **Description.** Set floating-point modes.

5 2 **Class.** Subroutine.

6 3 **Argument.** MODES shall be a scalar of type IEEE\_MODES\_TYPE. Its value shall be one that was assigned  
7 by a previous invocation of IEEE\_GET\_MODES to its MODES argument. The floating-point modes (17.7) are  
8 restored to the state at that invocation.

9 4 **Example.**

10 To save the floating-point modes, do a calculation with specific rounding and underflow modes, and restore them  
11 later:

```
12         USE, INTRINSIC :: IEEE_ARITHMETIC
13         TYPE (IEEE_MODES_TYPE) SAVE_MODES
14         ...
15         CALL IEEE_GET_MODES (SAVE_MODES) ! Save all modes.
16         CALL IEEE_SET_ROUNDING_MODE (IEEE_TO_ZERO)
17         CALL IEEE_SET_UNDERFLOW_MODE (GRADUAL=.FALSE.)
18         ... ! calculation with abrupt round-to-zero.
19         CALL IEEE_SET_MODES (SAVE_MODES) ! Restore all modes.
```

### 20 17.11.38 IEEE\_SET\_ROUNDING\_MODE (ROUND\_VALUE [, RADIX])

21 1 **Description.** Set rounding mode.

22 2 **Class.** Subroutine.

23 3 **Arguments.**

24 ROUND\_VALUE shall be a scalar of type IEEE\_ROUND\_TYPE. It specifies the rounding mode to be set.

25 RADIX (optional) shall be an integer scalar with the value two or ten. If RADIX is present with the value ten,  
26 the rounding mode set is the decimal rounding mode; otherwise it is the binary rounding mode.

27 4 **Restriction.** IEEE\_SET\_ROUNDING\_MODE (ROUND\_VALUE) shall not be invoked unless IEEE\_SUPPORT\_ROUNDING  
28 (ROUND\_VALUE, X) is true for some X such that IEEE\_SUPPORT\_DATATYPE (X)  
29 is true. IEEE\_SET\_ROUNDING\_MODE (ROUND\_VALUE, RADIX) shall not be invoked unless IEEE\_SUPPORT\_ROUNDING  
30 (ROUND\_VALUE, X) is true for some X with radix RADIX such that IEEE\_SUPPORT\_DATATYPE (X) is true.  
31

32 5 **Example.** To save the binary rounding mode, do a calculation with round to nearest, and restore the rounding  
33 mode later:

```
34         USE, INTRINSIC :: IEEE_ARITHMETIC
35         TYPE (IEEE_ROUND_TYPE) ROUND_VALUE
36         ...
37         CALL IEEE_GET_ROUNDING_MODE (ROUND_VALUE) ! Store the rounding mode
38         CALL IEEE_SET_ROUNDING_MODE (IEEE_NEAREST)
39         ... ! calculation with round to nearest
40         CALL IEEE_SET_ROUNDING_MODE (ROUND_VALUE) ! Restore the rounding mode
```

### 17.11.39 IEEE\_SET\_STATUS (STATUS\_VALUE)

1 **Description.** Restore floating-point status.

2 **Class.** Subroutine.

3 **Argument.** STATUS\_VALUE shall be a scalar of type `IEEE_STATUS_TYPE`. Its value shall be one that was assigned by a previous invocation of `IEEE_GET_STATUS` to its STATUS\_VALUE argument. The floating-point status (17.7 is restored to the state at that invocation).

4 **Example.** To store all the exceptions flags, do a calculation involving exception handling, and restore them later:

```

9      USE, INTRINSIC :: IEEE_EXCEPTIONS
10     TYPE (IEEE_STATUS_TYPE) STATUS_VALUE
11     ...
12     CALL IEEE_GET_STATUS (STATUS_VALUE) ! Store the flags
13     CALL IEEE_SET_FLAG (IEEE_ALL, .FALSE.) ! Set them quiet
14     ... ! calculation involving exception handling
15     CALL IEEE_SET_STATUS (STATUS_VALUE) ! Restore the flags

```

### 17.11.40 IEEE\_SET\_UNDERFLOW\_MODE (GRADUAL)

1 **Description.** Set `underflow mode`.

2 **Class.** Subroutine.

3 **Argument.** GRADUAL shall be a logical scalar. If it is true, the `underflow mode` is set to gradual underflow. If it is false, the `underflow mode` is set to abrupt underflow.

4 **Restriction.** `IEEE_SET_UNDERFLOW_MODE` shall not be invoked unless `IEEE_SUPPORT_UNDERFLOW_CONTROL (X)` is true for some X.

5 **Example.** To perform some calculations with abrupt underflow and then restore the previous mode:

```

24     USE, INTRINSIC :: IEEE_ARITHMETIC
25     LOGICAL SAVE_UNDERFLOW_MODE
26     ...
27     CALL IEEE_GET_UNDERFLOW_MODE (SAVE_UNDERFLOW_MODE)
28     CALL IEEE_SET_UNDERFLOW_MODE (GRADUAL=.FALSE.)
29     ... ! Perform some calculations with abrupt underflow
30     CALL IEEE_SET_UNDERFLOW_MODE (SAVE_UNDERFLOW_MODE)

```

### 17.11.41 IEEE\_SIGNALING\_EQ (A, B)

1 **Description.** Signaling compares equal.

2 **Class.** `Elemental` function.

3 **Arguments.**

A shall be of type real.

B shall be of the same type and kind type parameter as A.

4 **Restriction.** `IEEE_SIGNALING_EQ (A, B)` shall not be invoked if `IEEE_SUPPORT_DATATYPE (A)` has the value false.

1 5 **Result Characteristics.** Default logical.

2 6 **Result Value.** The result has the value specified for the compareSignalingEqual operation in ISO/IEC/IEEE  
3 60559:2011; that is, it is true if and only if A compares equal to B. If A or B is a NaN, the result will be false  
4 and IEEE\_INVALID signals; otherwise, no exception is signaled.

5 7 **Example.** IEEE\_SIGNALING\_EQ (1.0, IEEE\_VALUE (IEEE\_QUIET\_NAN)) has the value false and signals  
6 IEEE\_INVALID.

#### 7 **17.11.42 IEEE\_SIGNALING\_GE (A, B)**

8 1 **Description.** Signaling compares greater than or equal.

9 2 **Class.** Elemental function.

10 3 **Arguments.**

11 A shall be of type real.

12 B shall be of the same type and kind type parameter as A.

13 4 **Restriction.** IEEE\_SIGNALING\_GE (A, B) shall not be invoked if IEEE\_SUPPORT\_DATATYPE (A) has  
14 the value false.

15 5 **Result Characteristics.** Default logical.

16 6 **Result Value.** The result has the value specified for the compareSignalingGreaterEqual operation in  
17 ISO/IEC/IEEE 60559:2011; that is, it is true if and only if A compares greater than or equal to B. If A or  
18 B is a NaN, the result will be false and IEEE\_INVALID signals; otherwise, no exception is signaled.

19 7 **Example.** IEEE\_SIGNALING\_GE (1.0, IEEE\_VALUE (IEEE\_QUIET\_NAN)) has the value false and signals  
20 IEEE\_INVALID.

#### 21 **17.11.43 IEEE\_SIGNALING\_GT (A, B)**

22 1 **Description.** Signaling compares greater than.

23 2 **Class.** Elemental function.

24 3 **Arguments.**

25 A shall be of type real.

26 B shall be of the same type and kind type parameter as A.

27 4 **Restriction.** IEEE\_SIGNALING\_GT (A, B) shall not be invoked if IEEE\_SUPPORT\_DATATYPE (A) has  
28 the value false.

29 5 **Result Characteristics.** Default logical.

30 6 **Result Value.** The result has the value specified for the compareSignalingGreater operation in ISO/IEC/IEEE  
31 60559:2011; that is, it is true if and only if A compares greater than B. If A or B is a NaN, the result will be false  
32 and IEEE\_INVALID signals; otherwise, no exception is signaled.

33 7 **Example.** IEEE\_SIGNALING\_GT (1.0, IEEE\_VALUE (IEEE\_QUIET\_NAN)) has the value false and signals  
34 IEEE\_INVALID.

#### 35 **17.11.44 IEEE\_SIGNALING\_LE (A, B)**

36 1 **Description.** Signaling compares less than or equal.

37 2 **Class.** Elemental function.

1 3 **Arguments.**

2 A shall be of type real.

3 B shall be of the same type and kind type parameter as A.

4 4 **Restriction.** IEEE\_SIGNALING\_LE (A, B) shall not be invoked if IEEE\_SUPPORT\_DATATYPE (A) has  
5 the value false.

6 5 **Result Characteristics.** Default logical.

7 6 **Result Value.** The result has the value specified for the compareSignalingLessEqual operation in ISO/IEC/IEEE  
8 60559:2011; that is, it is true if and only if A compares less than or equal to B. If A or B is a NaN, the result will  
9 be false and IEEE\_INVALID signals; otherwise, no exception is signaled.

10 7 **Example.** IEEE\_SIGNALING\_LE (1.0, IEEE\_VALUE (IEEE\_QUIET\_NAN)) has the value false and signals  
11 IEEE\_INVALID.

12 **17.11.45 IEEE\_SIGNALING\_LT (A, B)**

13 1 **Description.** Signaling compares less than.

14 2 **Class.** Elemental function.

15 3 **Arguments.**

16 A shall be of type real.

17 B shall be of the same type and kind type parameter as A.

18 4 **Restriction.** IEEE\_SIGNALING\_LT (A, B) shall not be invoked if IEEE\_SUPPORT\_DATATYPE (A) has  
19 the value false.

20 5 **Result Characteristics.** Default logical.

21 6 **Result Value.** The result has the value specified for the compareSignalingLess operation in ISO/IEC/IEEE  
22 60559:2011; that is, it is true if and only if A compares less than B. If A or B is a NaN, the result will be false  
23 and IEEE\_INVALID signals; otherwise, no exception is signaled.

24 7 **Example.** IEEE\_SIGNALING\_LT (1.0, IEEE\_VALUE (IEEE\_QUIET\_NAN)) has the value false and signals  
25 IEEE\_INVALID.

26 **17.11.46 IEEE\_SIGNALING\_NE (A, B)**

27 1 **Description.** Signaling compares not equal.

28 2 **Class.** Elemental function.

29 3 **Arguments.**

30 A shall be of type real.

31 B shall be of the same type and kind type parameter as A.

32 4 **Restriction.** IEEE\_SIGNALING\_NE (A, B) shall not be invoked if IEEE\_SUPPORT\_DATATYPE (A) has  
33 the value false.

34 5 **Result Characteristics.** Default logical.

35 6 **Result Value.** The result has the value specified for the compareSignalingNotEqual operation in ISO/IEC/IEEE  
36 60559:2011; that is, it is true if and only if A compares not equal to B. If A or B is a NaN, the result will be true  
37 and IEEE\_INVALID signals; otherwise, no exception is signaled.

38 7 **Example.** IEEE\_SIGNALING\_NE (1.0, IEEE\_VALUE (IEEE\_QUIET\_NAN)) has the value true and signals

1 IEEE\_INVALID.

## 2 17.11.47 IEEE\_SIGNBIT (X)

3 1 **Description.** Test sign bit.

4 2 **Class.** [Elemental](#) function.

5 3 **Argument.** X shall be of type real.

6 4 **Restriction.** IEEE\_SIGNBIT (X) shall not be invoked if [IEEE\\_SUPPORT\\_DATATYPE](#) (X) has the value  
7 false.

8 5 **Result Characteristics.** Default logical.

9 6 **Result Value.** The result has the value specified for the isSignMinus operation in ISO/IEC/IEEE 60559:2011;  
10 that is, it is true if and only if the sign bit of X is nonzero. No exception is signaled even if X is a signaling NaN.

11 7 **Example.** IEEE\_SIGNBIT (−1.0) has the value true.

## 12 17.11.48 IEEE\_SUPPORT\_DATATYPE () or 13 IEEE\_SUPPORT\_DATATYPE (X)

13 1 **Description.** Query IEEE arithmetic support.

14 2 **Class.** [Inquiry](#) function.

15 3 **Argument.** X shall be of type real. It may be a scalar or an array.

16 4 **Result Characteristics.** Default logical scalar.

17 5 **Result Value.** The result has the value true if the processor supports IEEE arithmetic for all reals (X does  
18 not appear) or for real variables of the same kind type parameter as X; otherwise, it has the value false. Here,  
19 support is as defined in the first paragraph of [17.9](#).

20 6 **Example.** If default real kind conforms to ISO/IEC/IEEE 60559:2011 except that underflow values flush to zero  
21 instead of being subnormal, IEEE\_SUPPORT\_DATATYPE (1.0) has the value true.

## 22 17.11.49 IEEE\_SUPPORT\_DENORMAL () or 23 IEEE\_SUPPORT\_DENORMAL (X)

23 1 **Description.** Query subnormal number support.

24 2 **Class.** [Inquiry](#) function.

25 3 **Argument.** X shall be of type real. It may be a scalar or an array.

26 4 **Result Characteristics.** Default logical scalar.

27 5 **Result Value.**

28 *Case (i):* IEEE\_SUPPORT\_DENORMAL (X) has the value true if [IEEE\\_SUPPORT\\_DATATYPE](#) (X) has  
29 the value true and the processor supports arithmetic operations and assignments with subnormal  
30 numbers (biased exponent  $e = 0$  and fraction  $f \neq 0$ , see subclause 3.2 of ISO/IEC/IEEE 60559:2011)  
31 for real variables of the same kind type parameter as X; otherwise, it has the value false.

32 *Case (ii):* IEEE\_SUPPORT\_DENORMAL () has the value true if IEEE\_SUPPORT\_DENORMAL (X) has  
33 the value true for all real X; otherwise, it has the value false.

34 6 **Example.** IEEE\_SUPPORT\_DENORMAL (X) has the value true if the processor supports subnormal values  
35 for X.

**NOTE 17.12**

A reference to IEEE\_SUPPORT\_DENORMAL will have the same result value as a reference to IEEE\_SUPPORT\_SUBNORMAL with the same argument list.

**17.11.50 IEEE\_SUPPORT\_DIVIDE () or IEEE\_SUPPORT\_DIVIDE (X)**

1 **Description.** Query IEEE division support.

2 **Class.** [Inquiry function](#).

3 **Argument.** X shall be of type real. It may be a scalar or an array.

4 **Result Characteristics.** Default logical scalar.

5 **Result Value.**

6 *Case (i):* IEEE\_SUPPORT\_DIVIDE (X) has the value true if the processor supports division with the accuracy specified by ISO/IEC/IEEE 60559:2011 for real variables of the same kind type parameter as X; otherwise, it has the value false.

7 *Case (ii):* IEEE\_SUPPORT\_DIVIDE () has the value true if IEEE\_SUPPORT\_DIVIDE (X) has the value true for all real X; otherwise, it has the value false.

8 **Example.** IEEE\_SUPPORT\_DIVIDE (X) has the value true if division of operands with the same kind as X conforms to ISO/IEC/IEEE 60559:2011.

**17.11.51 IEEE\_SUPPORT\_FLAG (FLAG) or IEEE\_SUPPORT\_FLAG (FLAG, X)**

1 **Description.** Query exception support.

2 **Class.** [Transformational function](#).

3 **Arguments.**

4 FLAG shall be a scalar of type [IEEE\\_FLAG\\_TYPE](#). Its value shall be one of IEEE\_INVALID, IEEE\_OVERFLOW, IEEE\_DIVIDE\_BY\_ZERO, IEEE\_UNDERFLOW, or IEEE\_INEXACT.

5 X shall be of type real. It may be a scalar or an array.

6 **Result Characteristics.** Default logical scalar.

7 **Result Value.**

8 *Case (i):* IEEE\_SUPPORT\_FLAG (FLAG, X) has the value true if the processor supports detection of the specified exception for real variables of the same kind type parameter as X; otherwise, it has the value false.

9 *Case (ii):* IEEE\_SUPPORT\_FLAG (FLAG) has the value true if IEEE\_SUPPORT\_FLAG (FLAG, X) has the value true for all real X; otherwise, it has the value false.

10 **Example.** IEEE\_SUPPORT\_FLAG (IEEE\_INEXACT) has the value true if the processor supports the inexact exception.

**17.11.52 IEEE\_SUPPORT\_HALTING (FLAG)**

1 **Description.** Query [halting mode](#) support.

2 **Class.** [Transformational function](#).

3 **Argument.** FLAG shall be a scalar of type [IEEE\\_FLAG\\_TYPE](#). Its value shall be one of IEEE\_INVALID, IEEE\_OVERFLOW, IEEE\_DIVIDE\_BY\_ZERO, IEEE\_UNDERFLOW, or IEEE\_INEXACT.

1 4 **Result Characteristics.** Default logical scalar.

2 5 **Result Value.** The result has the value true if the processor supports the ability to control during program  
3 execution whether to abort or continue execution after the exception specified by FLAG; otherwise, it has the  
4 value false. Support includes the ability to change the mode by CALL [IEEE\\_SET\\_HALTING\\_MODE](#) (FLAG).

5 6 **Example.** IEEE\_SUPPORT\_HALTING (IEEE\_OVERFLOW) has the value true if the processor supports  
6 control of [halting](#) after an overflow.

### 7 **17.11.53 IEEE\_SUPPORT\_INF () or IEEE\_SUPPORT\_INF (X)**

8 1 **Description.** Query IEEE infinity support.

9 2 **Class.** [Inquiry function](#).

10 3 **Argument.** X shall be of type real. It may be a scalar or an array.

11 4 **Result Characteristics.** Default logical scalar.

12 5 **Result Value.**

13 *Case (i):* IEEE\_SUPPORT\_INF (X) has the value true if the processor supports IEEE infinities (positive  
14 and negative) for real variables of the same kind type parameter as X; otherwise, it has the value  
15 false.

16 *Case (ii):* IEEE\_SUPPORT\_INF () has the value true if IEEE\_SUPPORT\_INF (X) has the value true for  
17 all real X; otherwise, it has the value false.

18 6 **Example.** IEEE\_SUPPORT\_INF (X) has the value true if the processor supports IEEE infinities for X.

### 19 **17.11.54 IEEE\_SUPPORT\_IO () or IEEE\_SUPPORT\_IO (X)**

20 1 **Description.** Query IEEE formatting support.

21 2 **Class.** [Inquiry function](#).

22 3 **Argument.** X shall be of type real. It may be a scalar or an array.

23 4 **Result Characteristics.** Default logical scalar.

24 5 **Result Value.**

25 *Case (i):* IEEE\_SUPPORT\_IO (X) has the value true if base conversion during formatted input/output  
26 ([12.5.6.16](#), [12.6.2.13](#), [13.7.2.3.8](#)) conforms to ISO/IEC/IEEE 60559:2011 for the modes UP, DOWN,  
27 ZERO, and NEAREST for real variables of the same kind type parameter as X; otherwise, it has  
28 the value false.

29 *Case (ii):* IEEE\_SUPPORT\_IO () has the value true if IEEE\_SUPPORT\_IO (X) has the value true for all  
30 real X; otherwise, it has the value false.

31 6 **Example.** IEEE\_SUPPORT\_IO (X) has the value true if formatted input/output base conversions conform to  
32 ISO/IEC/IEEE 60559:2011.

### 33 **17.11.55 IEEE\_SUPPORT\_NAN () or IEEE\_SUPPORT\_NAN (X)**

34 1 **Description.** Query [IEEE NaN](#) support.

35 2 **Class.** [Inquiry function](#).

36 3 **Argument.** X shall be of type real. It may be a scalar or an array.

37 4 **Result Characteristics.** Default logical scalar.

38 5 **Result Value.**



1 *Case (i):* IEEE\_SUPPORT\_NAN (X) has the value true if the processor supports IEEE NaNs for real  
2 variables of the same kind type parameter as X; otherwise, it has the value false.

3 *Case (ii):* IEEE\_SUPPORT\_NAN () has the value true if IEEE\_SUPPORT\_NAN (X) has the value true  
4 for all real X; otherwise, it has the value false.

5 6 **Example.** IEEE\_SUPPORT\_NAN (X) has the value true if the processor supports IEEE NaNs for X.

### 6 17.11.56 IEEE\_SUPPORT\_ROUNDING (ROUND\_VALUE) or IEEE\_SUPPORT\_ROUNDING (ROUND\_VALUE, X)

7 1 **Description.** Query IEEE rounding support.

8 2 **Class.** Transformational function.

9 3 **Arguments.**

10 ROUND\_VALUE shall be of type IEEE\_ROUND\_TYPE.

11 X shall be of type real. It may be a scalar or an array.

12 4 **Result Characteristics.** Default logical scalar.

13 5 **Result Value.**

14 *Case (i):* IEEE\_SUPPORT\_ROUNDING (ROUND\_VALUE, X) has the value true if the processor supports  
15 the rounding mode defined by ROUND\_VALUE for real variables of the same kind type parameter  
16 as X; otherwise, it has the value false. Support includes the ability to change the mode by CALL  
17 IEEE\_SET\_ROUNDING\_MODE (ROUND\_VALUE).

18 *Case (ii):* IEEE\_SUPPORT\_ROUNDING (ROUND\_VALUE) has the value true if IEEE\_SUPPORT\_  
19 ROUNDING (ROUND\_VALUE, X) has the value true for all real X; otherwise, it has the value  
20 false.

21 6 **Example.** IEEE\_SUPPORT\_ROUNDING (IEEE\_TO\_ZERO) has the value true if the processor supports  
22 rounding to zero for all reals.

### 23 17.11.57 IEEE\_SUPPORT\_SQRT () or IEEE\_SUPPORT\_SQRT (X)

24 1 **Description.** Query IEEE square root support.

25 2 **Class.** Inquiry function.

26 3 **Argument.** X shall be of type real. It may be a scalar or an array.

27 4 **Result Characteristics.** Default logical scalar.

28 5 **Result Value.**

29 *Case (i):* IEEE\_SUPPORT\_SQRT (X) has the value true if the intrinsic function SQRT conforms to  
30 ISO/IEC/IEEE 60559:2011 for real variables of the same kind type parameter as X; otherwise,  
31 it has the value false.

32 *Case (ii):* IEEE\_SUPPORT\_SQRT () has the value true if IEEE\_SUPPORT\_SQRT (X) has the value true  
33 for all real X; otherwise, it has the value false.

34 6 **Example.** If IEEE\_SUPPORT\_SQRT (1.0) has the value true, SQRT (−0.0) will have the value −0.0.

### 35 17.11.58 IEEE\_SUPPORT\_STANDARD () or IEEE\_SUPPORT\_STANDARD (X)

36 1 **Description.** Query IEEE standard support.

37 2 **Class.** Inquiry function.

1 3 **Argument.** X shall be of type real. It may be a scalar or an array.

2 4 **Result Characteristics.** Default logical scalar.

3 5 **Result Value.**

4 *Case (i):* IEEE\_SUPPORT\_STANDARD (X) has the value true if the results of all the functions  
 5 IEEE\_SUPPORT\_DATATYPE (X), IEEE\_SUPPORT\_DIVIDE (X), IEEE\_SUPPORT\_  
 6 FLAG (FLAG, X) for valid FLAG, IEEE\_SUPPORT\_HALTING (FLAG) for valid  
 7 FLAG, IEEE\_SUPPORT\_INF (X), IEEE\_SUPPORT\_NAN (X), IEEE\_SUPPORT\_  
 8 ROUNDING (ROUND\_VALUE, X) for valid ROUND\_VALUE, IEEE\_SUPPORT\_SQRT  
 9 (X), and IEEE\_SUPPORT\_SUBNORMAL (X) are all true; otherwise, it has the value false.

10 *Case (ii):* IEEE\_SUPPORT\_STANDARD () has the value true if IEEE\_SUPPORT\_STANDARD (X) has  
 11 the value true for all real X; otherwise, it has the value false.

12 6 **Example.** IEEE\_SUPPORT\_STANDARD () has the value false if some but not all kinds of reals conform to  
 13 ISO/IEC/IEEE 60559:2011.

### 14 17.11.59 IEEE\_SUPPORT\_SUBNORMAL () or IEEE\_SUPPORT\_SUBNORMAL (X)

15 1 **Description.** Query subnormal number support.

16 2 **Class.** Inquiry function.

17 3 **Argument.** X shall be of type real. It may be a scalar or an array.

18 4 **Result Characteristics.** Default logical scalar.

19 5 **Result Value.**

20 *Case (i):* IEEE\_SUPPORT\_SUBNORMAL (X) has the value true if IEEE\_SUPPORT\_DATATYPE (X)  
 21 has the value true and the processor supports arithmetic operations and assignments with subnormal  
 22 numbers (biased exponent  $e = 0$  and fraction  $f \neq 0$ , see subclause 3.2 of ISO/IEC/IEEE 60559:2011)  
 23 for real variables of the same kind type parameter as X; otherwise, it has the value false.

24 *Case (ii):* IEEE\_SUPPORT\_SUBNORMAL () has the value true if IEEE\_SUPPORT\_SUBNORMAL (X)  
 25 has the value true for all real X; otherwise, it has the value false.

26 6 **Example.** IEEE\_SUPPORT\_SUBNORMAL (X) has the value true if the processor supports subnormal values  
 27 for X.

#### NOTE 17.13

The subnormal numbers are not included in the 16.4 model for real numbers; they satisfy the inequality  $ABS(X) < TINY(X)$ . They usually occur as a result of an arithmetic operation whose exact result is less than  $TINY(X)$ . Such an operation causes IEEE\_UNDERFLOW to signal unless the result is exact. IEEE\_SUPPORT\_SUBNORMAL (X) is false if the processor never returns a subnormal number as the result of an arithmetic operation.

### 28 17.11.60 IEEE\_SUPPORT\_UNDERFLOW\_CONTROL () or IEEE\_SUPPORT\_UNDERFLOW\_CONTROL (X)

29 1 **Description.** Query underflow control support.

30 2 **Class.** Inquiry function.

31 3 **Argument.** X shall be of type real. It may be a scalar or an array.

32 4 **Result Characteristics.** Default logical scalar.

1 5 **Result Value.**

2 *Case (i):* IEEE\_SUPPORT\_UNDERFLOW\_CONTROL (X) has the value true if the processor supports  
3 control of the [underflow mode](#) for floating-point calculations with the same type as X, and false  
4 otherwise.

5 *Case (ii):* IEEE\_SUPPORT\_UNDERFLOW\_CONTROL () has the value true if the processor supports  
6 control of the [underflow mode](#) for all floating-point calculations, and false otherwise.

7 6 **Example.** IEEE\_SUPPORT\_UNDERFLOW\_CONTROL (2.5) has the value true if the processor supports  
8 [underflow mode](#) control for default real calculations.

9 **17.11.61 IEEE\_UNORDERED (X, Y)**

10 1 **Description.** Whether two values are unordered.

11 2 **Class.** [Elemental](#) function.

12 3 **Arguments.** The arguments shall be of type real.

13 4 **Restriction.** IEEE\_UNORDERED (X, Y) shall not be invoked if IEEE\_SUPPORT\_DATATYPE (X) or  
14 IEEE\_SUPPORT\_DATATYPE (Y) has the value false.

15 5 **Result Characteristics.** Default logical.

16 6 **Result Value.** The result has the value true if X or Y is a NaN or both are NaNs; otherwise, it has the value  
17 false. If X or Y is a signaling NaN, IEEE\_INVALID may signal.

18 7 **Example.** IEEE\_UNORDERED (0.0, [SQRT](#) (−1.0)) has the value true if IEEE\_SUPPORT\_SQRT (1.0) has  
19 the value true.

20 **17.11.62 IEEE\_VALUE (X, CLASS)**

21 1 **Description.** Return number in a class.

22 2 **Class.** [Elemental](#) function.

23 3 **Arguments.**

24 X shall be of type real.

25 CLASS shall be of type [IEEE\\_CLASS\\_TYPE](#). The value is permitted to be: IEEE\_SIGNALING\_NAN or  
26 IEEE\_QUIET\_NAN if [IEEE\\_SUPPORT\\_NAN](#) (X) has the value true, IEEE\_NEGATIVE\_INF  
27 or IEEE\_POSITIVE\_INF if [IEEE\\_SUPPORT\\_INF](#) (X) has the value true, IEEE\_NEGATIVE\_-  
28 SUBNORMAL or IEEE\_POSITIVE\_SUBNORMAL if [IEEE\\_SUPPORT\\_SUBNORMAL](#) (X) has  
29 the value true, IEEE\_NEGATIVE\_NORMAL, IEEE\_NEGATIVE\_ZERO, IEEE\_POSITIVE\_-  
30 ZERO or IEEE\_POSITIVE\_NORMAL.

31 4 **Restriction.** IEEE\_VALUE (X, CLASS) shall not be invoked if [IEEE\\_SUPPORT\\_DATATYPE](#) (X) has the  
32 value false.

33 5 **Result Characteristics.** Same as X.

34 6 **Result Value.** The result value is an IEEE value as specified by CLASS. Although in most cases the value is  
35 processor dependent, the value shall not vary between invocations for any particular X kind type parameter and  
36 CLASS value.

37 7 **Example.** IEEE\_VALUE (1.0, IEEE\_NEGATIVE\_INF) has the value −infinity.

38 8 Whenever IEEE\_VALUE returns a signaling NaN, it is processor dependent whether or not invalid is raised and  
39 processor dependent whether or not the signaling NaN is converted into a quiet NaN.

**NOTE 17.14**

If the *expr* in an **assignment statement** is a reference to the IEEE\_VALUE function that returns a signaling NaN and the *variable* is of the same type and kind as the function result, it is recommended that the signaling NaN be preserved.

1 **17.12 Examples****NOTE 17.15**

```

MODULE DOT
  ! Module for dot product of two real arrays of rank 1.
  ! The caller needs to ensure that exceptions do not cause halting.
  USE, INTRINSIC :: IEEE_EXCEPTIONS
  LOGICAL :: MATRIX_ERROR = .FALSE.
  INTERFACE OPERATOR(.dot.)
    MODULE PROCEDURE MULT
  END INTERFACE
CONTAINS
  REAL FUNCTION MULT (A, B)
    REAL, INTENT (IN) :: A(:), B(:)
    INTEGER I
    LOGICAL OVERFLOW
    IF (SIZE(A) /= SIZE(B)) THEN
      MATRIX_ERROR = .TRUE.
      RETURN
    END IF
    ! The processor ensures that IEEE_OVERFLOW is quiet.
    MULT = 0.0
    DO I = 1, SIZE (A)
      MULT = MULT + A(I)*B(I)
    END DO
    CALL IEEE_GET_FLAG (IEEE_OVERFLOW, OVERFLOW)
    IF (OVERFLOW) MATRIX_ERROR = .TRUE.
  END FUNCTION MULT
END MODULE DOT

```

This module provides a function that computes the dot product of two real arrays of **rank 1**. If the sizes of the arrays are different, an immediate return occurs with **MATRIX\_ERROR** true. If overflow occurs during the actual calculation, the **IEEE\_OVERFLOW** flag will signal and **MATRIX\_ERROR** will be true.

**NOTE 17.16**

```

USE, INTRINSIC :: IEEE_EXCEPTIONS
USE, INTRINSIC :: IEEE_FEATURES, ONLY: IEEE_INVALID_FLAG
! The other exceptions of IEEE_USUAL (IEEE_OVERFLOW and
! IEEE_DIVIDE_BY_ZERO) are always available with IEEE_EXCEPTIONS
TYPE (IEEE_STATUS_TYPE) STATUS_VALUE

```

## NOTE 17.16 (cont.)

```

LOGICAL, DIMENSION(3) :: FLAG_VALUE
...
CALL IEEE_GET_STATUS (STATUS_VALUE)
CALL IEEE_SET_HALTING_MODE (IEEE_USUAL, .FALSE.) ! Needed in case the
!
!           default on the processor is to halt on exceptions
CALL IEEE_SET_FLAG (IEEE_USUAL, .FALSE.)
! First try the "fast" algorithm for inverting a matrix:
MATRIX1 = FAST_INV (MATRIX) ! This shall not alter MATRIX.
CALL IEEE_GET_FLAG (IEEE_USUAL, FLAG_VALUE)
IF (ANY(FLAG_VALUE)) THEN
  ! "Fast" algorithm failed; try "slow" one:
  CALL IEEE_SET_FLAG (IEEE_USUAL, .FALSE.)
  MATRIX1 = SLOW_INV (MATRIX)
  CALL IEEE_GET_FLAG (IEEE_USUAL, FLAG_VALUE)
  IF (ANY (FLAG_VALUE)) THEN
    WRITE (*, *) 'Cannot invert matrix'
    STOP
  END IF
END IF
CALL IEEE_SET_STATUS (STATUS_VALUE)

```

In this example, the function FAST\_INV might cause a condition to signal. If it does, another try is made with SLOW\_INV. If this still fails, a message is printed and the program stops. Note, also, that it is important to set the flags quiet before the second try. The state of all the flags is stored and restored.

## NOTE 17.17

```

USE, INTRINSIC :: IEEE_EXCEPTIONS
LOGICAL FLAG_VALUE
...
CALL IEEE_SET_HALTING_MODE (IEEE_OVERFLOW, .FALSE.)
! First try a fast algorithm for inverting a matrix.
CALL IEEE_SET_FLAG (IEEE_OVERFLOW, .FALSE.)
DO K = 1, N
  ...
  CALL IEEE_GET_FLAG (IEEE_OVERFLOW, FLAG_VALUE)
  IF (FLAG_VALUE) EXIT
END DO
IF (FLAG_VALUE) THEN
  ! Alternative code which knows that K-1 steps have executed normally.
  ...
END IF

```

Here the code for matrix inversion is in line and the transfer is made more precise by adding extra tests of the flag.

1

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# 18 Interoperability with C

## 18.1 General

- 1 Fortran provides a means of referencing procedures that are defined by means of the C programming language or procedures that can be described by C prototypes as defined in 6.7.6.3 of ISO/IEC 9899:2011, even if they are not actually defined by means of C. Conversely, there is a means of specifying that a procedure defined by a Fortran subprogram can be referenced from a function defined by means of C. In addition, there is a means of declaring global variables that are associated with C variables whose names have external linkage as defined in 6.2.2 of ISO/IEC 9899:2011.
- 2 The ISO\_C\_BINDING module provides access to [named constants](#) that represent [kind type parameters](#) of data representations compatible with C types. Fortran also provides facilities for defining derived types (7.5) and enumerations (7.6) that correspond to C types.
- 3 The source file ISO\_Fortran\_binding.h provides definitions and prototypes to enable a C function to interoperate with a Fortran procedure that has a dummy data object that is [allocatable](#), [assumed-shape](#), [assumed-rank](#), [pointer](#), or is of type character with an assumed length.

## 18.2 The ISO\_C\_BINDING intrinsic module

### 18.2.1 Summary of contents

- 1 The processor shall provide the intrinsic module ISO\_C\_BINDING. This module shall make accessible the following entities: the [named constants](#) C\_NULL\_PTR, C\_NULL\_FUNPTR, and those with names listed in the first column of Table 18.1 and the second column of Table 18.2, the types C\_PTR and C\_FUNPTR, and the procedures in 18.2.3. A processor may provide other public entities in the ISO\_C\_BINDING intrinsic module in addition to those listed here.

### 18.2.2 Named constants and derived types in the module

- 1 The entities listed in the second column of Table 18.2 shall be default integer [named constants](#).
- 2 A Fortran [intrinsic type](#) whose [kind type parameter](#) is one of the values in the module shall have the same representation as the C type with which it interoperates, for each value that a variable of that type can have. For C\_BOOL, the internal representation of [.TRUE.](#)\_C\_BOOL and [.FALSE.](#)\_C\_BOOL shall be the same as those of the C values (\_Bool)1 and (\_Bool)0 respectively.
- 3 The value of C\_INT shall be a valid value for an integer kind parameter on the processor. The values of C\_SHORT, C\_LONG, C\_LONG\_LONG, C\_SIGNED\_CHAR, C\_SIZE\_T, C\_INT8\_T, C\_INT16\_T, C\_INT32\_T, C\_INT64\_T, C\_INT\_LEAST8\_T, C\_INT\_LEAST16\_T, C\_INT\_LEAST32\_T, C\_INT\_LEAST64\_T, C\_INT\_FAST8\_T, C\_INT\_FAST16\_T, C\_INT\_FAST32\_T, C\_INT\_FAST64\_T, C\_INTMAX\_T, and C\_INTPTR\_T shall each be a valid value for an integer [kind type parameter](#) on the processor or shall be  $-1$  if the [companion processor](#) (5.5.7) defines the corresponding C type and there is no interoperating Fortran processor kind, or  $-2$  if the [companion processor](#) does not define the corresponding C type.
- 4 The values of C\_FLOAT, C\_DOUBLE, and C\_LONG\_DOUBLE shall each be a valid value for a real [kind type parameter](#) on the processor or shall be  $-1$  if the [companion processor](#)'s type does not have a precision equal to the precision of any of the Fortran processor's real kinds,  $-2$  if the [companion processor](#)'s type does not have a range equal to the range of any of the Fortran processor's real kinds,  $-3$  if the [companion processor](#)'s type

- 1 has neither the precision nor range of any of the Fortran processor's real kinds, and equal to  $-4$  if there is no  
 2 interoperating Fortran processor kind for other reasons. The values of `C_FLOAT_COMPLEX`, `C_DOUBLE_`  
 3 `COMPLEX`, and `C_LONG_DOUBLE_COMPLEX` shall be the same as those of `C_FLOAT`, `C_DOUBLE`, and  
 4 `C_LONG_DOUBLE`, respectively.
- 5 5 The value of `C_BOOL` shall be a valid value for a logical kind parameter on the processor or shall be  $-1$ .
- 6 6 The value of `C_CHAR` shall be a valid value for a character [kind type parameter](#) on the processor or shall be  $-1$ .  
 7 If the value of `C_CHAR` is nonnegative, the character kind specified is the C character kind; otherwise, there is  
 8 no C character kind.
- 9 7 The following entities shall be [named constants](#) of type character with a length parameter of one. The kind  
 10 parameter value shall be equal to the value of `C_CHAR` unless `C_CHAR` =  $-1$ , in which case the kind parameter  
 11 value shall be the same as for default kind. The values of these constants are specified in Table 18.1. In the case  
 12 that `C_CHAR`  $\neq -1$  the value is specified using C syntax. The semantics of these values are explained in 5.2.1  
 13 and 5.2.2 of ISO/IEC 9899:2011.

Table 18.1: Names of C characters with special semantics

| Name                           | C definition    | Value                      |                               |
|--------------------------------|-----------------|----------------------------|-------------------------------|
|                                |                 | <code>C_CHAR</code> = $-1$ | <code>C_CHAR</code> $\neq -1$ |
| <code>C_NULL_CHAR</code>       | null character  | <code>CHAR(0)</code>       | '\0'                          |
| <code>C_ALERT</code>           | alert           | <code>ACHAR(7)</code>      | '\a'                          |
| <code>C_BACKSPACE</code>       | backspace       | <code>ACHAR(8)</code>      | '\b'                          |
| <code>C_FORM_FEED</code>       | form feed       | <code>ACHAR(12)</code>     | '\f'                          |
| <code>C_NEW_LINE</code>        | new line        | <code>ACHAR(10)</code>     | '\n'                          |
| <code>C_CARRIAGE_RETURN</code> | carriage return | <code>ACHAR(13)</code>     | '\r'                          |
| <code>C_HORIZONTAL_TAB</code>  | horizontal tab  | <code>ACHAR(9)</code>      | '\t'                          |
| <code>C_VERTICAL_TAB</code>    | vertical tab    | <code>ACHAR(11)</code>     | '\v'                          |

- 14 8 The entities `C_PTR` and `C_FUNPTR` are described in 18.3.3.
- 15 9 The entity `C_NULL_PTR` shall be a [named constant](#) of type `C_PTR`. The value of `C_NULL_PTR` shall be the  
 16 same as the value `NULL` in C. The entity `C_NULL_FUNPTR` shall be a [named constant](#) of type `C_FUNPTR`.  
 17 The value of `C_NULL_FUNPTR` shall be that of a null pointer to a function in C.

**NOTE 18.1**

|                                                                                       |
|---------------------------------------------------------------------------------------|
| The value of <code>NEW_LINE(C_NEW_LINE)</code> is <code>C_NEW_LINE</code> (16.9.140). |
|---------------------------------------------------------------------------------------|

18 **18.2.3 Procedures in the module**19 **18.2.3.1 General**

- 20 1 In the detailed descriptions below, procedure names are generic and not specific. The `C_F_POINTER` subroutine  
 21 is impure; all other procedures in the module are pure.

22 **18.2.3.2 C\_ASSOCIATED (C\_PTR\_1 [, C\_PTR\_2])**

23 1 **Description.** Query C pointer status.

24 2 **Class.** [Transformational function](#).

25 3 **Arguments.**

26 `C_PTR_1` shall be a scalar of type `C_PTR` or `C_FUNPTR`.

27 `C_PTR_2` (optional) shall be a scalar of the same type as `C_PTR_1`.



1 4 **Result Characteristics.** Default logical scalar.

2 5 **Result Value.**

3 *Case (i):* If C\_PTR\_2 is absent, the result is false if C\_PTR\_1 is a C null pointer and true otherwise.

4 *Case (ii):* If C\_PTR\_2 is present, the result is false if C\_PTR\_1 is a C null pointer. If C\_PTR\_1 is not a  
5 C null pointer, the result is true if C\_PTR\_1 compares equal to C\_PTR\_2 in the sense of 6.3.2.3  
6 and 6.5.9 of ISO/IEC 9899:2011, and false otherwise.

#### NOTE 18.2

The following example illustrates the use of C\_LOC and C\_ASSOCIATED.

```

USE, INTRINSIC :: ISO_C_BINDING, ONLY: C_PTR, C_FLOAT, C_ASSOCIATED, C_LOC
INTERFACE
  SUBROUTINE FOO(GAMMA) BIND(C)
    IMPORT C_PTR
    TYPE(C_PTR), VALUE :: GAMMA
  END SUBROUTINE FOO
END INTERFACE
REAL(C_FLOAT), TARGET, DIMENSION(100) :: ALPHA
TYPE(C_PTR) :: BETA
...
IF (.NOT. C_ASSOCIATED(BETA)) THEN
  BETA = C_LOC(ALPHA)
ENDIF
CALL FOO(BETA)

```

#### 7 18.2.3.3 C\_F\_POINTER (CPTR, FPTR [, SHAPE])

8 1 **Description.** Associate a data pointer with the [target](#) of a C pointer and specify its shape.

9 2 **Class.** Subroutine.

10 3 **Arguments.**

11 CPTR shall be a scalar of type [C\\_PTR](#). It is an **INTENT (IN)** argument. Its value shall be

- 12 • the [C address](#) of an [interoperable](#) data entity,
- 13 • the result of a reference to [C\\_LOC](#) with a noninteroperable argument, or
- 14 • the [C address](#) of a storage sequence that is not in use by any other Fortran entity.

15 The value of CPTR shall not be the [C address](#) of a Fortran variable that does not have the **TARGET**  
16 **attribute**.

17 FPTR shall be a pointer, shall not have a [deferred type parameter](#), and shall not be a [coindexed object](#).  
18 It is an **INTENT (OUT)** argument. If FPTR is an array, its shape is specified by SHAPE and each  
19 lower bound is equal to 1.

20 *Case (i):* If the value of CPTR is the [C address](#) of an [interoperable](#) data entity, FPTR  
21 shall be a [data pointer](#) with type and [type parameter](#) values [interoperable](#)  
22 with the type of the entity. If the target T of CPTR is scalar, FPTR becomes pointer  
23 associated with T; if FPTR is an array, SHAPE shall specify a size of 1. If T is  
24 an array, and FPTR is scalar, FPTR becomes associated with the first element of  
25 T. If both T and FPTR are arrays, SHAPE shall specify a size that is less than or  
26 equal to the size of T, and FPTR becomes associated with the first **PRODUCT**  
27 (SHAPE) elements of T (this could be the entirety of T).

1           *Case (ii):* If the value of CPTR is the result of a reference to C\_LOC with a noninter-  
 2 operable **effective argument** X, FPTR shall be a nonpolymorphic pointer with  
 3 the same type and type parameters as X. In this case, X shall not have been  
 4 deallocated or have become undefined due to execution of a **RETURN** or **END**  
 5 statement since the reference. If X is scalar, FPTR becomes **pointer associated**  
 6 with X; if FPTR is an array, SHAPE shall specify a size of 1. If X is an array and  
 7 FPTR is scalar, FPTR becomes associated with the first element of X. If both X  
 8 and FPTR are arrays, SHAPE shall specify a size that is less than or equal to  
 9 the size of X, and FPTR becomes associated with the first **PRODUCT** (SHAPE)  
 10 elements of X (this could be the entirety of X).

11           *Case (iii):* If the value of CPTR is the **C address** of a storage sequence that is not in use by  
 12 any other Fortran entity, FPTR becomes associated with that storage sequence.  
 13 If FPTR is an array, its shape is specified by SHAPE and each lower bound  
 14 is 1. The storage sequence shall be large enough to contain the target object  
 15 described by FPTR and shall satisfy any other processor-dependent requirement  
 16 for association.

17 SHAPE (optional) shall be a rank-one integer array. It is an **INTENT (IN)** argument. SHAPE shall be present  
 18 if and only if FPTR is an array; its size shall be equal to the **rank** of FPTR.

#### 19 4 Examples.

20   *Case (i):*   extern double c\_x;  
 21               void \*address\_of\_x (void)  
 22               {  
 23                return &c\_x;  
 24               }  
 25  
 26               ! Assume interface to "address\_of\_x" is available.  
 27               Real (C\_double), Pointer :: xp  
 28               Call C\_F\_Pointer (address\_of\_x (), xp)

29   *Case (ii):* Type t  
 30               Real, Allocatable :: v(:, :)  
 31               End Type  
 32               Type(t), Target :: x  
 33               Type(C\_ptr) :: xloc  
 34               xloc = C\_Loc (x)  
 35               ...  
 36               Type(t), Pointer :: y  
 37               Call C\_F\_Pointer (xloc, y)

38   *Case (iii):* void \*getmem (int nbits)  
 39               {  
 40                return malloc ((nbits+CHAR\_BIT-1)/CHAR\_BIT);  
 41               }  
 42  
 43               ! Assume interface to "getmem" is available,  
 44               ! and there is a derived type "mytype" accessible.  
 45               Type(mytype), Pointer :: x  
 46               Call C\_F\_Pointer (getmem (**Storage\_Size** (x)), x)

47 The following statements illustrate the use of C\_F\_POINTER when the pointer to be set has a **deferred type**

```

1 parameter:
2     Character(42), Pointer :: C1
3     Character(:), Pointer :: C2
4     Call C_F_Pointer (CPTR, C1)
5     C2 => C1

```

6 This will associate C2 with the entity at the [C address](#) specified by CPTR, and specify its length to be the same  
7 as that of C1.

#### NOTE 18.3

In the case of associating FPTR with a storage sequence, there might be processor-dependent requirements such as alignment of the memory address or placement in memory.

### 8 18.2.3.4 C\_F\_PROCPOINTER (CPTR, FPTR)

9 1 **Description.** Associate a procedure pointer with the [target](#) of a C function pointer.

10 2 **Class.** [Pure subroutine](#).

11 3 **Arguments.**

12 CPTR shall be a scalar of type [C\\_FUNPTR](#). It is an [INTENT \(IN\)](#) argument. Its value shall be the [C address](#) of a procedure that is [interoperable](#), or the result of a reference to the function [C\\_FUNLOC](#)  
13 from the intrinsic module [ISO\\_C\\_BINDING](#).  
14

15 FPTR shall be a [procedure pointer](#), and shall not be a component of a [coindexed object](#). It is an [INTENT \(OUT\)](#) argument. If the target of CPTR is [interoperable](#), the [interface](#) for FPTR shall be [interoperable](#)  
16 with the prototype that describes the [target](#) of CPTR; otherwise, the [interface](#) for FPTR shall  
17 have the same characteristics as that target. FPTR becomes [pointer associated](#) with the [target](#) of  
18 CPTR.  
19

#### NOTE 18.4

The term “target” in the descriptions of [C\\_F\\_POINTER](#) and [C\\_F\\_PROCPOINTER](#) denotes the entity referenced by a C pointer, as described in 6.2.5 of ISO/IEC 9899:2011.

### 20 18.2.3.5 C\_FUNLOC (X)

21 1 **Description.** [C address](#) of the argument.

22 2 **Class.** [Transformational function](#).

23 3 **Argument.** X shall be a procedure; if it is a procedure pointer it shall be associated. It shall not be a [coindexed](#)  
24 [object](#).

25 4 **Result Characteristics.** Scalar of type [C\\_FUNPTR](#).

26 5 **Result Value.** The result value is described using the result name FUNPTR. The result is determined as if  
27 [C\\_FUNPTR](#) were a derived type containing a procedure pointer component PX with an [implicit interface](#) and  
28 the [pointer assignment](#) FUNPTR%PX => X were executed.

29 6 The result is a value that can be used as an [actual](#) CPTR argument in a call to [C\\_F\\_PROCPOINTER](#) where  
30 the FPTR argument has attributes that would allow the [pointer assignment](#) FPTR => X. Such a call to [C\\_F\\_-](#)  
31 [PROCPOINTER](#) shall have the effect of the pointer assignment FPTR => X.

### 1 18.2.3.6 C\_LOC (X)

2 1 **Description.** C address of the argument.

3 2 **Class.** Transformational function.

4 3 **Argument.** X shall have either the POINTER or TARGET attribute. It shall not be a coindexed object. It shall either be a variable with interoperable type and kind type parameters, or be a nonpolymorphic variable with no length type parameters. If it is allocatable, it shall be allocated. If it is a pointer, it shall be associated. If it is an array, it shall be contiguous and have nonzero size. It shall not be a zero-length string.

8 4 **Result Characteristics.** Scalar of type C\_PTR.

9 5 **Result Value.** The result value is described using the result name CPTR.

10 6 If X is a scalar data entity, the result is determined as if C\_PTR were a derived type containing a scalar pointer component PX of the type and type parameters of X and the pointer assignment CPTR%PX => X were executed.

12 7 If X is an array data entity, the result is determined as if C\_PTR were a derived type containing a scalar pointer component PX of the type and type parameters of X and the pointer assignment of CPTR%PX to the first element of X were executed.

15 8 If X is a data entity that is interoperable or has interoperable type and type parameters, the result is the value that the C processor returns as the result of applying the unary "&" operator (as defined in ISO/IEC 9899:2011, 6.5.3.2) to the target of CPTR%PX.

18 9 The result is a value that can be used as an actual CPTR argument in a call to C\_F\_POINTER where FPTR has attributes that would allow the pointer assignment FPTR => X. Such a call to C\_F\_POINTER shall have the effect of the pointer assignment FPTR => X.

#### NOTE 18.5

Where the actual argument is of noninteroperable type or type parameters, the result of C\_LOC provides an opaque "handle" for it. In an actual implementation, this handle might be the C address of the argument; however, only a C function that treats it as a void (generic) C pointer that cannot be dereferenced (6.5.3.2 in ISO/IEC 9899:2011) is likely to be portable.

### 21 18.2.3.7 C\_SIZEOF (X)

22 1 **Description.** Size of X in bytes.

23 2 **Class.** Inquiry function.

24 3 **Argument.** X shall be an interoperable data entity that is not an assumed-size array or an assumed-rank array that is associated with an assumed-size array.

26 4 **Result Characteristics.** Scalar integer of kind C\_SIZE\_T (18.3.2).

27 5 **Result Value.** If X is scalar, the result value is the value that the companion processor returns as the result of applying the sizeof operator (ISO/IEC 9899:2011, subclause 6.5.3.4) to an object of a type that interoperates with the type and type parameters of X.

30 6 If X is an array, the result value is the value that the companion processor returns as the result of applying the sizeof operator to an object of a type that interoperates with the type and type parameters of X, multiplied by the number of elements in X.

## 18.3 Interoperability between Fortran and C entities

### 18.3.1 General

- 1 Subclause 18.3 defines the conditions under which a Fortran entity is [interoperable](#). If a Fortran entity is [interoperable](#), an equivalent entity could be defined by means of C and the Fortran entity would interoperate with the C entity. There does not have to be such an interoperating C entity.

#### NOTE 18.6

A Fortran entity can be [interoperable](#) with more than one C entity.

### 18.3.2 Interoperability of intrinsic types

- 1 Table 18.2 shows the interoperability between Fortran intrinsic types and C types. A Fortran intrinsic type with particular type parameter values is [interoperable](#) with a C type if the type and [kind type parameter](#) value are listed in the table on the same row as that C type. If the type is character, the [length type parameter](#) is [interoperable](#) if and only if its value is one. A combination of Fortran type and type parameters that is [interoperable](#) with a C type listed in the table is also [interoperable](#) with any unqualified C type that is compatible with the listed C type.
- 2 The second column of the table refers to the [named constants](#) made accessible by the ISO\_C\_BINDING intrinsic module. If the value of any of these [named constants](#) is negative, there is no combination of Fortran type and type parameters [interoperable](#) with the C type shown in that row.
- 3 A combination of intrinsic type and type parameters is [interoperable](#) if it is [interoperable](#) with a C type. The C types mentioned in table 18.2 are defined in subclauses 6.2.5, 7.19, and 7.20.1 of ISO/IEC 9899:2011.

Table 18.2: Interoperability between Fortran and C types

| Fortran type | Named constant from the ISO_C_BINDING module<br>(kind type parameter if value is positive) | C type                       |
|--------------|--------------------------------------------------------------------------------------------|------------------------------|
| INTEGER      | C_INT                                                                                      | int                          |
|              | C_SHORT                                                                                    | short int                    |
|              | C_LONG                                                                                     | long int                     |
|              | C_LONG_LONG                                                                                | long long int                |
|              | C_SIGNED_CHAR                                                                              | signed char<br>unsigned char |
|              | C_SIZE_T                                                                                   | size_t                       |
|              | C_INT8_T                                                                                   | int8_t                       |
|              | C_INT16_T                                                                                  | int16_t                      |
|              | C_INT32_T                                                                                  | int32_t                      |
|              | C_INT64_T                                                                                  | int64_t                      |
|              | C_INT_LEAST8_T                                                                             | int_least8_t                 |
|              | C_INT_LEAST16_T                                                                            | int_least16_t                |
|              | C_INT_LEAST32_T                                                                            | int_least32_t                |
|              | C_INT_LEAST64_T                                                                            | int_least64_t                |
|              | C_INT_FAST8_T                                                                              | int_fast8_t                  |
|              | C_INT_FAST16_T                                                                             | int_fast16_t                 |
|              | C_INT_FAST32_T                                                                             | int_fast32_t                 |
|              | C_INT_FAST64_T                                                                             | int_fast64_t                 |
|              | C_INTMAX_T                                                                                 | intmax_t                     |

## Interoperability between Fortran and C types

(cont.)

| Fortran type | Named constant from the ISO_C_BINDING module<br>(kind type parameter if value is positive) | C type               |
|--------------|--------------------------------------------------------------------------------------------|----------------------|
|              | C_INTPTR_T                                                                                 | intptr_t             |
|              | C_PTRDIFF_T                                                                                | ptrdiff_t            |
| REAL         | C_FLOAT                                                                                    | float                |
|              | C_DOUBLE                                                                                   | double               |
|              | C_LONG_DOUBLE                                                                              | long double          |
| COMPLEX      | C_FLOAT_COMPLEX                                                                            | float _Complex       |
|              | C_DOUBLE_COMPLEX                                                                           | double _Complex      |
|              | C_LONG_DOUBLE_COMPLEX                                                                      | long double _Complex |
| LOGICAL      | C_BOOL                                                                                     | _Bool                |
| CHARACTER    | C_CHAR                                                                                     | char                 |

**NOTE 18.7**

ISO/IEC 9899:2011 specifies that the representations for nonnegative signed integers are the same as the corresponding values of unsigned integers. Because Fortran does not provide direct support for unsigned kinds of integers, the ISO\_C\_BINDING module does not make accessible [named constants](#) for their [kind type parameter](#) values. A user can use the signed kinds of integers to interoperate with the unsigned types and all their qualified versions as well. This has the potentially surprising side effect that the C type unsigned char is [interoperable](#) with the type integer with a [kind type parameter](#) of C\_SIGNED\_CHAR.

1 **18.3.3 Interoperability with C pointer types**

- 2 1 C\_PTR and C\_FUNPTR shall be derived types with only private components. No [direct component](#) of either  
3 of these types is [allocatable](#) or a pointer. C\_PTR is [interoperable](#) with any C object pointer type. C\_FUNPTR  
4 is [interoperable](#) with any C function pointer type.

**NOTE 18.8**

This means that only a C processor with the same representation method for all C object pointer types, and the same representation method for all C function pointer types, can be the target of interoperability of a Fortran processor. ISO/IEC 9899:2011 does not require this to be the case.

**NOTE 18.9**

The function C\_LOC can be used to return a value of type C\_PTR that is the [C address](#) of an allocated [allocatable](#) variable. The function C\_FUNLOC can be used to return a value of type C\_FUNPTR that is the [C address](#) of a procedure. For C\_LOC and C\_FUNLOC the returned value is of an [interoperable](#) type and thus can be used in contexts where the procedure or [allocatable](#) variable is not directly allowed. For example, it could be passed as an [actual argument](#) to a C function.

Similarly, type C\_FUNPTR or C\_PTR can be used in a dummy argument or [structure component](#) and can have a value that is the [C address](#) of a procedure or [allocatable](#) variable, even in contexts where a procedure or [allocatable](#) variable is not directly allowed.

5 **18.3.4 Interoperability of derived types and C struct types**

- 6 1 Interoperability between a [derived type](#) in Fortran and a struct type in C is provided by the [BIND attribute](#) on  
7 the Fortran type.  
8 C1801 (R726) A [derived type](#) with the [BIND attribute](#) shall not have the [SEQUENCE attribute](#).

- 1 C1802 (R726) A *derived type* with the *BIND attribute* shall not have *type parameters*.
- 2 C1803 (R726) A *derived type* with the *BIND attribute* shall not have the *EXTENDS attribute*.
- 3 C1804 (R726) A *derived-type-def* that defines a *derived type* with the *BIND attribute* shall not have a *type-bound-procedure-part*.
- 4
- 5 C1805 (R726) A *derived type* with the *BIND attribute* shall have at least one component.
- 6 C1806 (R726) Each component of a *derived type* with the *BIND attribute* shall be a nonpointer, nonallocatable data component with *interoperable type* and *type parameters*.
- 7

**NOTE 18.10**

The syntax rules and their constraints require that a *derived type* that is *interoperable* with a C struct type have components that are all data entities that are *interoperable*. No component is permitted to be *allocatable* or a pointer, but the value of a component of type *C\_FUNPTR* or *C\_PTR* can be the *C address* of such an entity.

- 8 2 A *derived type* is *interoperable* with a C struct type if and only if the *derived type* has the *BIND attribute* (7.5.2), the *derived type* and the C struct type have the same number of components, and the components of the *derived type* would interoperate with corresponding components of the C struct type as described in 18.3.5 and 18.3.6 if the components were variables. A component of a *derived type* and a component of a C struct type correspond if they are declared in the same relative position in their respective type definitions.
- 9
- 10
- 11
- 12

**NOTE 18.11**

The names of the corresponding components of the *derived type* and the C struct type need not be the same.

- 13 3 There is no Fortran type that is *interoperable* with a C struct type that contains a bit field or that contains a flexible array member. There is no Fortran type that is *interoperable* with a C union type.
- 14

**NOTE 18.12**

For example, the C type *myctype*, declared below, is *interoperable* with the Fortran type *myftype*, declared below.

```
typedef struct {
    int m, n;
    float r;
} myctype;

USE, INTRINSIC :: ISO_C_BINDING
TYPE, BIND(C) :: MYFCTYPE
    INTEGER(C_INT) :: I, J
    REAL(C_FLOAT) :: S
END TYPE MYFCTYPE
```

The names of the types and the names of the components are not significant for the purposes of determining whether a Fortran derived type is *interoperable* with a C struct type.

**NOTE 18.13**

ISO/IEC 9899:2011 requires the names and component names to be the same in order for the types to be compatible (ISO/IEC 9899:2011, subclause 6.2.7). This is similar to Fortran's rule describing when different derived type definitions describe the same *sequence type*. This rule was not extended to determine whether

**NOTE 18.13 (cont.)**

a Fortran derived type is [interoperable](#) with a C struct type because the case of identifiers is significant in C but not in Fortran.

**1 18.3.5 Interoperability of scalar variables**

- 2 1 A named scalar Fortran variable is [interoperable](#) if and only if its type and type parameters are [interoperable](#), it  
3 is not a [coarray](#), it has neither the [ALLOCATABLE](#) nor the [POINTER](#) attribute, and if it is of type character  
4 its length is not assumed or declared by an expression that is not a [constant expression](#).
- 5 2 An [interoperable](#) scalar Fortran variable is [interoperable](#) with a scalar C entity if their types and type parameters  
6 are [interoperable](#).

**7 18.3.6 Interoperability of array variables**

- 8 1 A Fortran variable that is a named array is [interoperable](#) if and only if its type and type parameters are [interop-  
9 erable](#), it is not a [coarray](#), it is of explicit shape or assumed size, and if it is of type character its length is not  
10 assumed or declared by an expression that is not a [constant expression](#).
- 11 2 An [explicit-shape](#) or [assumed-size](#) array of [rank](#)  $r$ , with a shape of  $[ e_1 \dots e_r ]$  is [interoperable](#) with a C array  
12 if its size is nonzero and
- 13 (1) either
- 14 (a) the array is [assumed-size](#), and the C array does not specify a size, or
- 15 (b) the array is an [explicit-shape array](#), and the extent of the last dimension ( $e_r$ ) is the same as  
16 the size of the C array, and
- 17 (2) either
- 18 (a)  $r$  is equal to one, and an element of the array is [interoperable](#) with an element of the C array,  
19 or
- 20 (b)  $r$  is greater than one, and an [explicit-shape array](#) with shape of  $[ e_1 \dots e_{r-1} ]$ , with the  
21 same type and type parameters as the original array, is [interoperable](#) with a C array of a type  
22 equal to the element type of the original C array.

**NOTE 18.14**

An element of a multi-dimensional C array is an array type, so a Fortran array of [rank](#) one is not [interop-  
erable](#) with a multidimensional C array.

**NOTE 18.15**

An [allocatable](#) array or [array pointer](#) is never [interoperable](#). Such an array does not meet the requirement  
of being an [explicit-shape](#) or [assumed-size](#) array.

**NOTE 18.16**

For example, a Fortran array declared as

```
INTEGER(C_INT) :: A(18, 3:7, *)
```

is [interoperable](#) with a C array declared as

```
int b[][5][18];
```



**NOTE 18.17**

The C programming language defines null-terminated strings, which are actually arrays of the C type char that have a C null character in them to indicate the last valid element. A Fortran array of type character with a **kind type parameter** equal to C\_CHAR is **interoperable** with a C string.

Fortran's rules of sequence association (15.5.2.11) permit a character scalar **actual argument** to correspond to a dummy argument array. This makes it possible to argument associate a Fortran character string with a C string.

Note 18.21 has an example of interoperation between Fortran and C strings.

**18.3.7 Interoperability of procedures and procedure interfaces**

1 A Fortran procedure is **interoperable** if and only if it has the **BIND attribute**, that is, if its **interface** is specified with a *proc-language-binding-spec*.

2 A Fortran procedure **interface** is **interoperable** with a C function prototype if

- (1) the **interface** has the **BIND attribute**,
- (2) either
  - (a) the **interface** describes a function whose **result** is a scalar variable that is **interoperable** with the result of the prototype or
  - (b) the **interface** describes a subroutine and the prototype has a result type of void,
- (3) the number of **dummy arguments** of the **interface** is equal to the number of formal parameters of the prototype,
- (4) any scalar dummy argument with the **VALUE attribute** is **interoperable** with the corresponding formal parameter of the prototype,
- (5) any dummy argument without the **VALUE attribute** corresponds to a formal parameter of the prototype that is of a pointer type, and either
  - the dummy argument is **interoperable** with an entity of the referenced type (ISO/IEC 9899:2011, 6.2.5, 7.19, and 7.20.1) of the formal parameter,
  - the dummy argument is a nonallocatable nonpointer variable of type CHARACTER with assumed character length and the formal parameter is a pointer to CFI\_cdasc\_t,
  - the dummy argument is **allocatable**, **assumed-shape**, **assumed-rank**, or a **pointer** without the **CONTIGUOUS attribute**, and the formal parameter is a pointer to CFI\_cdasc\_t, or
  - the dummy argument is **assumed-type** and not **allocatable**, **assumed-shape**, **assumed-rank**, or a **pointer**, and the formal parameter is a pointer to void,
- (6) each **allocatable** or **pointer** dummy argument of type CHARACTER has **deferred character length**, and
- (7) the prototype does not have variable arguments as denoted by the ellipsis (...).

**NOTE 18.18**

The **referenced type** of a C pointer type is the C type of the object that the C pointer type points to. For example, the referenced type of the pointer type `int *` is `int`.

**NOTE 18.19**

The C language allows specification of a C function that can take a variable number of arguments (ISO/IEC 9899:2011, 7.16). This document does not provide a mechanism for Fortran procedures to interoperate with such C functions.

3 A formal parameter of a C function prototype corresponds to a **dummy argument** of a Fortran **interface** if they are in the same relative positions in the C parameter list and the **dummy argument** list, respectively.

- 1 4 In a reference from C to a Fortran procedure with an interoperable [interface](#), a C actual argument shall be the  
 2 address of a [C descriptor](#) for the intended [effective argument](#) if the corresponding [dummy argument](#) interoperates  
 3 with a C formal parameter that is a pointer to `CFI_cdesc_t`. In this [C descriptor](#), the members other than  
 4 `attribute` and `type` shall describe an object with the same characteristics as the intended [effective argument](#).  
 5 The value of the `attribute` member of the [C descriptor](#) shall be compatible with the characteristics of the [dummy](#)  
 6 [argument](#). The `type` member shall have a value that depends on the intended [effective argument](#) as follows:
- 7 • if the [dynamic type](#) of the intended [effective argument](#) is an interoperable type listed in Table 18.4, the  
 8 corresponding value for that type;
  - 9 • if the [dynamic type](#) of the intended [effective argument](#) is an [intrinsic type](#) for which the processor defines  
 10 a nonnegative type specifier value not listed in Table 18.4, that type specifier value;
  - 11 • otherwise, `CFI_type_other`.
- 12 5 When an interoperable Fortran procedure with a [simply contiguous dummy argument](#) is invoked from C and the  
 13 actual argument is the address of a [C descriptor](#) for a discontinuous object, the Fortran processor shall handle  
 14 the difference in contiguity.
- 15 6 When an interoperable C procedure whose Fortran interface has a [simply contiguous dummy argument](#) is invoked  
 16 from Fortran and the [effective argument](#) is discontinuous, the Fortran processor shall ensure that the C procedure  
 17 receives a descriptor for a contiguous object.
- 18 7 When an interoperable C procedure whose Fortran interface has a [simply contiguous dummy argument](#) is invoked  
 19 from C, and the actual argument is the address of a [C descriptor](#) for a discontinuous object, the C code within  
 20 the procedure shall be prepared to handle the discontinuous argument.
- 21 8 If an interoperable procedure defined by means other than Fortran has an optional [dummy argument](#), and the  
 22 corresponding [actual argument](#) in a reference from Fortran is absent, the procedure is invoked with a null pointer  
 23 for that argument. If an interoperable procedure defined by means of Fortran is invoked by a C function, an  
 24 optional [dummy argument](#) is absent if and only if the corresponding argument in the invocation is a null pointer.

**NOTE 18.20**

For example, a Fortran procedure [interface](#) described by

```

INTERFACE
  FUNCTION FUNC(I, J, K, L, M) BIND(C)
    USE, INTRINSIC :: ISO_C_BINDING
    INTEGER(C_SHORT) :: FUNC
    INTEGER(C_INT), VALUE :: I
    REAL(C_DOUBLE) :: J
    INTEGER(C_INT) :: K, L(10)
    TYPE(C_PTR), VALUE :: M
  END FUNCTION FUNC
END INTERFACE

```

is [interoperable](#) with the C function prototype

```
short func(int i, double *j, int *k, int l[10], void *m);
```

A C pointer can correspond to a Fortran dummy argument of type `C_PTR` with the [VALUE attribute](#) or to a Fortran scalar that does not have the [VALUE attribute](#). In the above example, the C pointers `j` and `k` correspond to the Fortran scalars `J` and `K`, respectively, and the C pointer `m` corresponds to the Fortran dummy argument `M` of type `C_PTR`.

**NOTE 18.21**

The interoperability of Fortran procedure [interfaces](#) with C function prototypes is only one part of invocation of a C function from Fortran. There are four pieces to consider in such an invocation: the procedure reference, the Fortran procedure [interface](#), the C function prototype, and the C function. Conversely, the invocation of a Fortran procedure from C involves the function reference, the C function prototype, the Fortran procedure [interface](#), and the Fortran procedure. In order to determine whether a reference is allowed, it is necessary to consider all four pieces.

For example, consider a C function that can be described by the C function prototype

```
void copy(char in[], char out[]);
```

Such a function can be invoked from Fortran as follows:

```
USE, INTRINSIC :: ISO_C_BINDING, ONLY: C_CHAR, C_NULL_CHAR
INTERFACE
  SUBROUTINE COPY(IN, OUT) BIND(C)
    IMPORT C_CHAR
    CHARACTER(KIND=C_CHAR), DIMENSION(*) :: IN, OUT
  END SUBROUTINE COPY
END INTERFACE

CHARACTER(LEN=10, KIND=C_CHAR) :: &
&   DIGIT_STRING = C_CHAR_'123456789' // C_NULL_CHAR
CHARACTER(KIND=C_CHAR) :: DIGIT_ARR(10)

CALL COPY(DIGIT_STRING, DIGIT_ARR)
PRINT '(1X, A1)', DIGIT_ARR(1:9)
END
```

The procedure reference has character string [actual arguments](#). These correspond to character array [dummy arguments](#) in the procedure [interface body](#) as allowed by Fortran's rules of sequence association ([15.5.2.11](#)). Those array [dummy arguments](#) in the procedure [interface](#) are [interoperable](#) with the formal parameters of the C function prototype. The C function is not shown here, but is assumed to be compatible with the C function prototype.

**NOTE 18.22**

The requirements on the Fortran processor when a discontinuous [actual argument](#) is being passed to a [simply contiguous dummy argument](#) are such that the C programmer has the same functionality (and safety) as the Fortran programmer, when either the calling or called procedure are Fortran procedures. The requirements on the C procedure, in the case when both the called and calling procedure are C procedures, effectively require that the C procedure not access memory through the descriptor that is not described by the descriptor; otherwise the program will not conform to this document. A [dummy argument](#) is [simply contiguous](#) if it has the [CONTIGUOUS attribute](#) or is an array that is not [assumed-shape](#), [assumed-rank](#), or a [pointer](#).

## 1 18.4 C descriptors

- 2 1 A [C descriptor](#) is a C structure of type [CFI\\_cdesc\\_t](#). Together with library functions that have standard  
3 prototypes, it provides a means for describing and manipulating Fortran data objects from within a C function.  
4 This C structure is defined in the source file [ISO\\_Fortran\\_binding.h](#).

## 18.5 The source file ISO\_Fortran\_binding.h

### 18.5.1 Summary of contents

- 1 The source file ISO\_Fortran\_binding.h shall contain the C structure definitions, typedef declarations, macro definitions, and function prototypes specified in subclauses 18.5.2 to 18.5.5. The definitions and declarations in ISO\_Fortran\_binding.h can be used by a C function to interpret and manipulate a C descriptor. These provide a means to specify a C prototype that interoperates with a Fortran interface that has a non-interoperable dummy variable (18.3.7).
- 2 The source file ISO\_Fortran\_binding.h may be included in any order relative to the standard C headers, and may be included more than once in a given scope, with no effect different from being included only once, other than the effect on line numbers.
- 3 A C source file that includes the ISO\_Fortran\_binding.h header file shall not use any names starting with CFI\_ that are not defined in the header, and shall not define any of the structure names defined in the header as macro names. All names other than structure member names defined in the header begin with CFI\_ or an underscore character, or are defined by a standard C header that it includes.

### 18.5.2 The CFI\_dim\_t structure type

- 1 CFI\_dim\_t is a typedef name for a C structure. It is used to represent lower bound, extent, and memory stride information for one dimension of an array. The type CFI\_index\_t is described in 18.5.4. CFI\_dim\_t contains at least the following members in any order.

**CFI\_index\_t lower\_bound;** The value is equal to the value of the lower bound for the dimension being described.

**CFI\_index\_t extent;** The value is equal to the number of elements in the dimension being described, or -1 for the final dimension of an assumed-size array.

**CFI\_index\_t sm;** The value is equal to the memory stride for a dimension; this is the difference in bytes between the addresses of successive elements in the dimension being described.

### 18.5.3 The CFI\_cdesc\_t structure type

- 1 CFI\_cdesc\_t is a typedef name for a C structure, which contains a flexible array member. It shall contain at least the members described in this subclause. The values of these members of a structure of type CFI\_cdesc\_t that is produced by the functions and macros specified in this document, or received by a C function when invoked by a Fortran procedure, shall have the properties described in this subclause.

- 2 The first three members of the structure shall be **base\_addr**, **elem\_len**, and **version** in that order. The final member shall be **dim**. All other members shall be between **version** and **dim**, in any order. The types CFI\_attribute\_t, CFI\_rank\_t, and CFI\_type\_t are described in 18.5.4. The type CFI\_dim\_t is described in 18.5.2.

**void \* base\_addr;** If the object is an unallocated allocatable variable or a pointer that is disassociated, the value is a null pointer. If the object has zero size, the value is not a null pointer but is otherwise processor-dependent. Otherwise, the value is the base address of the object being described. The base address of a scalar is its C address. The base address of an array is the C address of the first element in Fortran array element order.

**size\_t elem\_len;** If the object is scalar, the value is the storage size in bytes of the object; otherwise, the value is the storage size in bytes of an element of the object.

**int version;** The value is equal to the value of CFI\_VERSION in the source file ISO\_Fortran\_binding.h that defined the format and meaning of this C descriptor when the descriptor was established.

- 1 **CFI\_rank\_t rank**; The value is equal to the number of dimensions of the Fortran object being described; if  
2 the object is scalar, the value is zero.
- 3 **CFI\_type\_t type**; The value is equal to the specifier for the type of the object. Each interoperable intrinsic C  
4 type has a specifier. Specifiers are also provided to indicate that the type of the object is an interoperable  
5 structure, or is unknown. The macros listed in Table 18.4 provide values that correspond to each specifier.
- 6 **CFI\_attribute\_t attribute**; The value is equal to the value of an attribute code that indicates whether the  
7 object described is [allocatable](#), a [data pointer](#), or a nonallocatable nonpointer data object. The macros  
8 listed in Table 18.3 provide values that correspond to each code.
- 9 **CFI\_dim\_t dim**; The number of elements in the `dim` array is equal to the rank of the object. Each element of  
10 the array contains the lower bound, extent, and memory stride information for the corresponding dimension  
11 of the Fortran object.
- 12 3 For a [C descriptor](#) of an array pointer or [allocatable](#) array, the value of the `lower_bound` member of each element  
13 of the `dim` member of the descriptor is determined by argument association, allocation, or pointer association.  
14 For a [C descriptor](#) of a nonallocatable nonpointer object, the value of the `lower_bound` member of each element  
15 of the `dim` member of the descriptor is zero.
- 16 4 There shall be an ordering of the dimensions such that the absolute value of the `sm` member of the first dimension  
17 is not less than the `elem_len` member of the [C descriptor](#) and the absolute value of the `sm` member of each  
18 subsequent dimension is not less than the absolute value of the `sm` member of the previous dimension multiplied  
19 by the extent of the previous dimension.
- 20 5 In a [C descriptor](#) of an [assumed-size array](#), the `extent` member of the last element of the `dim` member has the  
21 value `-1`.

**NOTE 18.23**

The reason for the restriction on the absolute values of the `sm` members is to ensure that there is no overlap between the elements of the array that is being described, while allowing for the reordering of subscripts. Within Fortran, such a reordering can be achieved with the intrinsic function `TRANSPOSE` or the intrinsic function `RESHAPE` with the optional argument `ORDER`, and an optimizing compiler can accommodate it without making a copy by constructing the appropriate descriptor whenever it can determine that a copy is not needed.

**NOTE 18.24**

The value of `elem_len` for a Fortran `CHARACTER` object is equal to the character length times the number of bytes of a single character of that kind. If the kind is `C_CHAR`, this value will be equal to the character length.

22 **18.5.4 Macros and typedefs in ISO\_Fortran\_binding.h**

- 23 1 Except for `CFI_CDESC_T`, each macro defined in `ISO_Fortran_binding.h` expands to an integer constant  
24 expression that is either a single token or a parenthesized expression that is suitable for use in `#if` preprocessing  
25 directives.
- 26 2 `CFI_CDESC_T` is a function-like macro that takes one argument, which is the rank of the [C descriptor](#) to create,  
27 and evaluates to an unqualified type of suitable size and alignment for defining a variable to use as a [C descriptor](#)  
28 of that rank. The argument shall be an integer constant expression with a value that is greater than or equal to  
29 zero and less than or equal to `CFI_MAX_RANK`. A pointer to a variable declared using `CFI_CDESC_T` can  
30 be cast to `CFI_cdesc_t*`. A variable declared using `CFI_CDESC_T` shall not have an initializer.

**NOTE 18.25**

The `CFI_CDESC_T` macro provides the memory for a **C descriptor**. The address of an entity declared using the macro is not usable as an actual argument corresponding to a formal parameter of type `CFI_cdesc_t *` without an explicit cast. For example, the following code uses `CFI_CDESC_T` to declare a **C descriptor** of rank 5 and pass it to `CFI_deallocate` (18.5.5.4).

```
CFI_CDESC_T(5) object;
int ind;
... code to define and use C descriptor ...
ind = CFI_deallocate((CFI_cdesc_t *)&object);
```

- 1 3 `CFI_index_t` is a typedef name for a standard signed integer type capable of representing the result of subtracting  
2 two pointers.
- 3 4 The `CFI_MAX_RANK` macro has a processor-dependent value equal to the largest rank supported. The value  
4 shall be greater than or equal to 15. `CFI_rank_t` is a typedef name for a standard integer type capable of  
5 representing the largest supported rank.
- 6 5 The `CFI_VERSION` macro has a processor-dependent value that encodes the version of the `ISO_Fortran_`  
7 `binding.h` source file containing this macro. This value should be increased if a new version of the source file is  
8 incompatible with the previous version.
- 9 6 The macros in Table 18.3 are for use as attribute codes. The values shall be nonnegative and distinct. `CFI_`  
10 `attribute_t` is a typedef name for a standard integer type capable of representing the values of the attribute  
11 codes.

Table 18.3: `ISO_Fortran_binding.h` macros for attribute codes

| Macro name                             | Attribute                 |
|----------------------------------------|---------------------------|
| <code>CFI_attribute_pointer</code>     | data pointer              |
| <code>CFI_attribute_allocatable</code> | allocatable               |
| <code>CFI_attribute_other</code>       | nonallocatable nonpointer |

- 12 7 `CFI_attribute_pointer` specifies a data object with the Fortran `POINTER` attribute. `CFI_attribute_allocatable`  
13 specifies an object with the Fortran `ALLOCATABLE` attribute. `CFI_attribute_other` specifies a nonallocatable  
14 nonpointer object.
- 15 8 The macros in Table 18.4 are for use as type specifiers. The value for `CFI_type_other` shall be negative and  
16 distinct from all other type specifiers. `CFI_type_struct` specifies a C structure that is interoperable with a  
17 Fortran derived type; its value shall be positive and distinct from all other type specifiers. If a C type is not  
18 interoperable with a Fortran type and kind supported by the Fortran processor, its macro shall evaluate to a  
19 negative value. Otherwise, the value for an intrinsic type shall be positive.
- 20 9 If the processor supports interoperability of a Fortran **intrinsic type** with a C type not listed in Table 18.4,  
21 the processor shall define a type specifier value for that type which is positive and distinct from all other type  
22 specifiers.
- 23 10 `CFI_type_t` is a typedef name for a standard integer type capable of representing the values for the supported  
24 type specifiers.

Table 18.4: `ISO_Fortran_binding.h` macros for type codes

| Macro name                        | C Type      |
|-----------------------------------|-------------|
| <code>CFI_type_signed_char</code> | signed char |
| <code>CFI_type_short</code>       | short int   |

## ISO\_Fortran\_binding.h macros for type codes (cont.)

| Macro name                   | C Type                    |
|------------------------------|---------------------------|
| CFI_type_int                 | int                       |
| CFI_type_long                | long int                  |
| CFI_type_long_long           | long long int             |
| CFI_type_size_t              | size_t                    |
| CFI_type_int8_t              | int8_t                    |
| CFI_type_int16_t             | int16_t                   |
| CFI_type_int32_t             | int32_t                   |
| CFI_type_int64_t             | int64_t                   |
| CFI_type_int_least8_t        | int_least8_t              |
| CFI_type_int_least16_t       | int_least16_t             |
| CFI_type_int_least32_t       | int_least32_t             |
| CFI_type_int_least64_t       | int_least64_t             |
| CFI_type_int_fast8_t         | int_fast8_t               |
| CFI_type_int_fast16_t        | int_fast16_t              |
| CFI_type_int_fast32_t        | int_fast32_t              |
| CFI_type_int_fast64_t        | int_fast64_t              |
| CFI_type_intmax_t            | intmax_t                  |
| CFI_type_intptr_t            | intptr_t                  |
| CFI_type_ptrdiff_t           | ptrdiff_t                 |
| CFI_type_float               | float                     |
| CFI_type_double              | double                    |
| CFI_type_long_double         | long double               |
| CFI_type_float_Complex       | float _Complex            |
| CFI_type_double_Complex      | double _Complex           |
| CFI_type_long_double_Complex | long double _Complex      |
| CFI_type_Bool                | _Bool                     |
| CFI_type_char                | char                      |
| CFI_type_cptr                | void *                    |
| CFI_type_struct              | interoperable C structure |
| CFI_type_other               | Not otherwise specified   |

**NOTE 18.26**

The values for different C types can be the same; for example, CFI\_type\_int and CFI\_type\_int32\_t might have the same value.

- 1 11 The macros in Table 18.5 are for use as error codes. The macro CFI\_SUCCESS shall be defined to be the integer  
 2 constant 0. The value of each macro other than CFI\_SUCCESS shall be nonzero and shall be different from the  
 3 values of the other macros specified in this subclause. Error conditions other than those listed in this subclause  
 4 should be indicated by error codes different from the values of the macros named in this subclause.
- 5 12 The values of the macros in Table 18.5 indicate the error condition described.

Table 18.5: ISO\_Fortran\_binding.h macros for error codes

| Macro name                   | Error condition                                                                                                                  |
|------------------------------|----------------------------------------------------------------------------------------------------------------------------------|
| CFI_SUCCESS                  | No error detected.                                                                                                               |
| CFI_ERROR_BASE_ADDR_NULL     | The base address member of a <a href="#">C descriptor</a> is a null pointer in a context that requires a non-null pointer value. |
| CFI_ERROR_BASE_ADDR_NOT_NULL | The base address member of a <a href="#">C descriptor</a> is not a null pointer in a context that requires a null pointer value. |
| CFI_INVALID_ELEM_LEN         | The value supplied for the element length member of a <a href="#">C descriptor</a> is not valid.                                 |

## ISO\_Fortran\_binding.h macros for error codes

(cont.)

| Macro name               | Error condition                                                                 |
|--------------------------|---------------------------------------------------------------------------------|
| CFI_INVALID_RANK         | The value supplied for the rank member of a C descriptor is not valid.          |
| CFI_INVALID_TYPE         | The value supplied for the type member of a C descriptor is not valid.          |
| CFI_INVALID_ATTRIBUTE    | The value supplied for the attribute member of a C descriptor is not valid.     |
| CFI_INVALID_EXTENT       | The value supplied for the extent member of a CFI_dim_t structure is not valid. |
| CFI_INVALID_DESCRIPTOR   | A C descriptor is invalid in some way.                                          |
| CFI_ERROR_MEM_ALLOCATION | Memory allocation failed.                                                       |
| CFI_ERROR_OUT_OF_BOUNDS  | A reference is out of bounds.                                                   |

## 18.5.5 Functions declared in ISO\_Fortran\_binding.h

### 18.5.5.1 Arguments and results of the functions

1 Some of the functions described in 18.5.5 return an error indicator; this is an integer value that indicates whether an error condition was detected. The value zero indicates that no error condition was detected, and a nonzero value indicates which error condition was detected. Table 18.5 lists standard error conditions and macro names for their corresponding error codes. A processor is permitted to detect other error conditions. If an invocation of a function defined in 18.5.5 could detect more than one error condition and an error condition is detected, which error condition is detected is processor dependent.

2 In function arguments representing subscripts, bounds, extents, or strides, the ordering of the elements is the same as the ordering of the elements of the dim member of a C descriptor.

3 Prototypes for these functions, or equivalent macros, are provided in the ISO\_Fortran\_binding.h file as described in 18.5.5. It is unspecified whether the functions defined by this header are macros or identifiers declared with external linkage. If a macro definition is suppressed in order to access an actual function, the behavior is undefined.

#### NOTE 18.27

These functions are allowed to be macros to provide extra implementation flexibility. For example, CFI\_ establish could include the value of CFI\_VERSION in the header used to compile the call to CFI\_ establish as an extra argument of the actual function used to establish the C descriptor.

### 18.5.5.2 The CFI\_address function

1 **Synopsis.** C address of an object described by a C descriptor.

```
void *CFI_address(const CFI_cdesc_t *dv, const CFI_index_t subscripts[]);
```

2 **Formal Parameters.**

**dv** shall be the address of a C descriptor describing the object. The object shall not be an unallocated allocatable variable or a pointer that is not associated.

**subscripts** shall be a null pointer or the address of an array of type CFI\_index\_t. If the object is an array, **subscripts** shall be the address of an array of CFI\_index\_t with at least *n* elements, where *n* is the rank of the object. The value of **subscripts**[*i*] shall be within the bounds of dimension *i* specified by the **dim** member of the C descriptor.

3 **Result Value.** If the object is an array of rank *n*, the result is the C address of the element of the object that



1 the first  $n$  elements of the `subscripts` argument would specify if used as subscripts. If the object is scalar, the  
2 result is its `C address`.

3 4 **Example.** If `dv` is the address of a `C descriptor` for the Fortran array `A` declared as

```
4 REAL(C_FLOAT) :: A(100, 100)
```

5 the following code calculates the `C address` of `A(5, 10)`:

```
6 CFI_index_t subscripts[2];
7 float *address;
8 subscripts[0] = 4;
9 subscripts[1] = 9;
10 address = (float *) CFI_address(dv, subscripts );
```

### 11 18.5.5.3 The `CFI_allocate` function

12 1 **Synopsis.** Allocate memory for an object described by a `C descriptor`.

```
13 int CFI_allocate(CFI_cdesc_t *dv, const CFI_index_t lower_bounds[],
14                const CFI_index_t upper_bounds[], size_t elem_len);
```

15 2 **Formal Parameters.**

16 `dv` shall be the address of a `C descriptor` specifying the rank and type of the object. The `base_`  
17 `addr` member of the `C descriptor` shall be a null pointer. If the type is not a character type, the  
18 `elem_len` member shall specify the element length. The `attribute` member shall have a value of  
19 `CFI_attribute_allocatable` or `CFI_attribute_pointer`.

20 `lower_bounds` shall be the address of an array with at least `dv->rank` elements, if `dv->rank`>0.

21 `upper_bounds` shall be the address of an array with at least `dv->rank` elements, if `dv->rank`>0.

22 `elem_len` If the type specified in the `C descriptor` type is a Fortran character type, the value of `elem_len`  
23 shall be the storage size in bytes of an element of the object; otherwise, `elem_len` is ignored.

24 3 **Description.** Successful execution of `CFI_allocate` allocates memory for the object described by the `C`  
25 `descriptor` with the address `dv` using the same mechanism as the Fortran `ALLOCATE statement`, and assigns the  
26 address of that memory to `dv->base_addr`. The first `dv->rank` elements of the `lower_bounds` and `upper_bounds`  
27 arguments provide the lower and upper Fortran bounds, respectively, for each corresponding dimension of the  
28 object. The supplied lower and upper bounds override any current dimension information in the `C descriptor`.  
29 If the rank is zero, the `lower_bounds` and `upper_bounds` arguments are ignored. If the type specified in the `C`  
30 `descriptor` is a character type, the supplied element length overrides the current element-length information in  
31 the descriptor.

32 If an error is detected, the `C descriptor` is not modified.

33 4 **Result Value.** The result is an error indicator.

34 5 **Example.** If `dv` is the address of a `C descriptor` for the Fortran array `A` declared as

```
35 REAL, ALLOCATABLE :: A(:, :)
```

36 and the array is not allocated, the following code allocates it to be of shape `[100, 500]`:

```
37 CFI_index_t lower[2], upper[2];
38 int ind;
39 lower[0] = 1; lower[1] = 1;
40 upper[0] = 100; upper[1] = 500;
41 ind = CFI_allocate(dv, lower, upper, 0);
```

#### 1 18.5.5.4 The CFI\_deallocate function

2 1 **Synopsis.** Deallocate memory for an object described by a [C descriptor](#).

```
3 int CFI_deallocate(CFI_cdesc_t *dv);
```

4 2 **Formal Parameter.** `dv` shall be the address of a [C descriptor](#) describing the object. It shall have been allocated using the same mechanism as the Fortran [ALLOCATE statement](#). If the object is a pointer, it shall be associated with a target satisfying the conditions for successful deallocation by the Fortran [DEALLOCATE statement \(9.7.3\)](#).

8 3 **Description.** Successful execution of `CFI_deallocate` deallocates memory for the object using the same mechanism as the Fortran [DEALLOCATE statement](#), and the `base_addr` member of the [C descriptor](#) becomes a null pointer.

11 If an error is detected, the [C descriptor](#) is not modified.

12 4 **Result Value.** The result is an error indicator.

13 5 **Example.** If `dv` is the address of a [C descriptor](#) for the Fortran array `A` declared as

```
14 REAL, ALLOCATABLE :: A(:, :)
```

15 and the array is allocated, the following code deallocates it:

```
16 int ind;
17 ind = CFI_deallocate(dv);
```

#### 18 18.5.5.5 The CFI\_establish function

19 1 **Synopsis.** Establish a [C descriptor](#).

```
20 int CFI_establish(CFI_cdesc_t *dv, void *base_addr, CFI_attribute_t attribute,
21                 CFI_type_t type, size_t elem_len, CFI_rank_t rank,
22                 const CFI_index_t extents[]);
```

23 2 **Formal Parameters.**

24 `dv` shall be the address of a data object large enough to hold a [C descriptor](#) of the rank specified by `rank`. It shall not have the same value as either a C formal parameter that corresponds to a Fortran actual argument or a C actual argument that corresponds to a Fortran dummy argument. It shall not be the address of a [C descriptor](#) that describes an allocated allocatable object.

28 `base_addr` shall be a null pointer or the base address of the object to be described. If it is not a null pointer, it shall be the address of a contiguous storage sequence that is appropriately aligned (ISO/IEC 9899:2011 3.2) for an object of the type specified by `type`.

31 `attribute` shall be one of the attribute codes in Table 18.3. If it is `CFI_attribute_allocatable`, `base_addr` shall be a null pointer.

33 `type` shall be one of the type codes in Table 18.4.

34 `elem_len` If the type is equal to `CFI_type_struct`, `CFI_type_other`, or a Fortran character type code, `elem_len` shall be greater than zero and equal to the storage size in bytes of an element of the object. Otherwise, `type` will be ignored.

37 `rank` shall have a value in the range  $0 \leq \text{rank} \leq \text{CFI\_MAX\_RANK}$ . It specifies the rank of the object.

38 `extents` is ignored if `rank` is equal to zero or if `base_addr` is a null pointer. Otherwise, it shall be the address of an array with `rank` elements; the value of each element shall be nonnegative, and `extents[i]` specifies the extent of dimension  $i$  of the object.

1 3 **Description.** Successful execution of `CFI_establish` updates the object with the address `dv` to be an established  
 2 **C descriptor** for a nonallocatable nonpointer data object of known shape, an unallocated allocatable object, or a  
 3 data pointer. If `base_addr` is not a null pointer, it is for a nonallocatable entity that is a scalar or a contiguous  
 4 array; if the `attribute` argument has the value `CFI_attribute_pointer`, the lower bounds of the object described  
 5 by `dv` are set to zero. If `base_addr` is a null pointer, the established **C descriptor** is for an unallocated allocatable,  
 6 a disassociated pointer, or is a **C descriptor** that has the `attribute` `CFI_attribute_other` but does not describe  
 7 a data object. If `base_addr` is the **C address** of a Fortran data object, the `type` and `elem_len` arguments shall be  
 8 consistent with the type and type parameters of the Fortran data object. The remaining properties of the object  
 9 are given by the other arguments.

10 If an error is detected, the object with the address `dv` is not modified.

11 4 **Result Value.** The result is an error indicator.

#### NOTE 18.28

`CFI_establish` is used to initialize a **C descriptor** declared in C with `CFI_CDESC_T` before passing it to any other functions as an actual argument, in order to set the rank, attribute, type and element length.

#### NOTE 18.29

A **C descriptor** with `attribute` `CFI_attribute_other` and `base_addr` a null pointer can be used as the argument `result` in calls to `CFI_section` or `CFI_select_part`, which will produce a **C descriptor** for a nonallocatable nonpointer data object.

12 5 **Examples.**

13 *Case (i):* The following code fragment establishes a **C descriptor** for an unallocated rank-one allocatable array  
 14 that can be passed to Fortran for allocation there.

```
15     CFI_rank_t rank;
16     CFI_CDESC_T(1) field;
17     int ind;
18     rank = 1;
19     ind = CFI_establish((CFI_cdesc_t *)&field, NULL, CFI_attribute_allocatable,
20                       CFI_type_double, 0, rank, NULL);
```

21 *Case (ii):* Given the Fortran type definition

```
22     TYPE, BIND(C) :: T
23         REAL(C_DOUBLE) :: X
24         COMPLEX(C_DOUBLE_COMPLEX) :: Y
25     END TYPE
```

26 and a Fortran subprogram that has an assumed-shape dummy argument of type T, the following  
 27 code fragment creates a descriptor `a_fortran` for an array of size 100 that can be used as the actual  
 28 argument in an invocation of the subprogram from C:

```
29     typedef struct {double x; double _Complex y;} t;
30     t a_c[100];
31     CFI_CDESC_T(1) a_fortran;
32     int ind;
33     CFI_index_t extent[1];
34
35     extent[0] = 100;
36     ind = CFI_establish((CFI_cdesc_t *)&a_fortran, a_c, CFI_attribute_other,
37                       CFI_type_struct, sizeof(t), 1, extent);
```

### 1 18.5.5.6 The CFI\_is\_contiguous function

2 1 **Synopsis.** Test contiguity of an array.

```
3 int CFI_is_contiguous(const CFI_cdesc_t * dv);
```

4 2 **Formal Parameter.** `dv` shall be the address of a **C descriptor** describing an array. The `base_addr` member of the **C descriptor** shall not be a null pointer.

6 3 **Result Value.** The value of the result is 1 if the array described by `dv` is contiguous, and 0 otherwise.

#### NOTE 18.30

Assumed-size and allocatable arrays are always contiguous, and therefore the result of `CFI_is_contiguous` on a **C descriptor** for such an array will be equal to 1.

### 7 18.5.5.7 The CFI\_section function

8 1 **Synopsis.** Update a **C descriptor** for an array section for which each element is an element of a given array.

```
9 int CFI_section(CFI_cdesc_t *result, const CFI_cdesc_t *source,
10               const CFI_index_t lower_bounds[], const CFI_index_t upper_bounds[],
11               const CFI_index_t strides[]);
```

12 2 **Formal Parameters.**

13 **result** shall be the address of a **C descriptor** with rank equal to the rank of `source` minus the number of zero strides. The `attribute` member shall have the value `CFI_attribute_other` or `CFI_attribute_pointer`. If the value of `result` is the same as either a C formal parameter that corresponds to a Fortran actual argument or a C actual argument that corresponds to a Fortran dummy argument, the `attribute` member shall have the value `CFI_attribute_pointer`.

18 Successful execution of `CFI_section` updates the `base_addr` and `dim` members of the **C descriptor** with the address `result` to describe the array section determined by `source`, `lower_bounds`, `upper_bounds`, and `strides`, as follows.

21 The array section is equivalent to the Fortran array section `SOURCE(sectsub1, sectsub2, ... sectsubn)`, where `SOURCE` is the array described by `source`,  $n$  is the rank of that array, and `sectsubi` is the subscript `loweri` if `stridei` is zero, and the section subscript `loweri : upperi : stridei` otherwise. The value of `loweri` is the lower bound of dimension  $i$  of `SOURCE` if `lower_bounds` is a null pointer and `lower_bounds[i]` otherwise. The value of `upperi` is the upper bound of dimension  $i$  of `SOURCE` if `upper_bounds` is a null pointer and `upper_bounds[i]` otherwise. The value of `stridei` is 1 if `strides` is a null pointer and `strides[i]` otherwise. If `stridei` has the value zero, `loweri` shall have the same value as `upperi`.

29 If an error is detected, the **C descriptor** with the address `result` is not modified.

30 **source** shall be the address of a **C descriptor** that describes a nonallocatable nonpointer array, an allocated allocatable array, or an associated array pointer. The `elem_len` and `type` members of `source` shall have the same values as the corresponding members of `result`.

33 **lower\_bounds** shall be a null pointer or the address of an array with at least `source->rank` elements. If it is not a null pointer, and `stridei` is zero or  $(upper<sub>i</sub> - lower\_bounds[i] + stride<sub>i</sub>)/stride<sub>i</sub> > 0$ , the value of `lower_bounds[i]` shall be within the bounds of dimension  $i$  of `SOURCE`.

36 **upper\_bounds** shall be a null pointer or the address of an array with at least `source->rank` elements. If `source` describes an **assumed-size array**, `upper_bounds` shall not be a null pointer. If it is not a null pointer and `stridei` is zero or  $(upper\_bounds[i] - lower<sub>i</sub> + stride<sub>i</sub>)/stride<sub>i</sub> > 0$ , the value of `upper\_bounds[i]` shall be within the bounds of dimension  $i$  of `SOURCE`.

40 **strides** shall be a null pointer or the address of an array with at least `source->rank` elements.

1 3 **Result Value.** The result is an error indicator.

2 4 **Examples.**

3 *Case (i):* If `source` is already the address of a **C descriptor** for the rank-one Fortran array `A`, the lower  
4 bounds of `A` are equal to 1, and the lower bounds in the **C descriptor** are equal to 0, the following  
5 code fragment establishes a new **C descriptor** section and updates it to describe the array section  
6 `A(3::5)`:

```
7     CFI_index_t lower[1], strides[1];
8     CFI_CDESC_T(1) section;
9     int ind;
10    lower[0] = 2;
11    strides[0] = 5;
12    ind = CFI_establish((CFI_cdesc_t *)&section, NULL, CFI_attribute_other,
13                      CFI_type_float, 0, 1, NULL);
14    ind = CFI_section((CFI_cdesc_t *)&section, source, lower, NULL, strides);
```

15 *Case (ii):* If `source` is already the address of a **C descriptor** for a rank-two Fortran assumed-shape array `A`  
16 with lower bounds equal to 1, the following code fragment establishes a **C descriptor** and updates  
17 it to describe the rank-one array section `A(:, 42)`.

```
18    CFI_index_t lower[2], upper[2], strides[2];
19    CFI_CDESC_T(1) section;
20    int ind;
21    lower[0] = source->dim[0].lower_bound;
22    upper[0] = source->dim[0].lower_bound + source->dim[0].extent - 1;
23    strides[0] = 1;
24    lower[1] = upper[1] = source->dim[1].lower_bound + 41;
25    strides[1] = 0;
26    ind = CFI_establish((CFI_cdesc_t *)&section, NULL, CFI_attribute_other,
27                      CFI_type_float, 0, 1, NULL);
28    ind = CFI_section((CFI_cdesc_t *)&section, source, lower, upper, strides);
```

#### 29 18.5.5.8 The `CFI_select_part` function

30 1 **Synopsis.** Update a **C descriptor** for an array section for which each element is a part of the corresponding  
31 element of an array.

```
32    int CFI_select_part(CFI_cdesc_t *result, const CFI_cdesc_t *source, size_t displacement,
33                      size_t elem_len);
```

#### 34 2 **Formal Parameters.**

35 **result** shall be the address of a **C descriptor**; `result->rank` shall have the same value as `source->rank`  
36 and `result->attribute` shall have the value `CFI_attribute_other` or `CFI_attribute_pointer`. If  
37 the address specified by `result` is the value of a C formal parameter that corresponds to a For-  
38 tran actual argument or of a C actual argument that corresponds to a Fortran dummy argument,  
39 `result->attribute` shall have the value `CFI_attribute_pointer`. The value of `result->type` spe-  
40 cifies the type of the array section.

41 **source** shall be the address of a **C descriptor** for a nonallocatable nonpointer array, an allocated allocatable  
42 array, or an associated array pointer.

43 **displacement** shall have a value  $0 \leq \text{displacement} \leq \text{source->elem\_len} - 1$ , and the sum of the displacement  
44 and the size in bytes of an element of the array section shall be less than or equal to `source->elem_-`  
45 `len`. The address `displacement` bytes greater than the value of `source->base_addr` is the base of

1 the array section and shall be appropriately aligned (ISO/IEC 9899:2011, 3.2) for an object of the  
2 type of the array section.

3 `elem_len` shall have a value equal to the storage size in bytes of an element of the array section if `result->type`  
4 specifies a Fortran character type; otherwise, `elem_len` is ignored.

5 **3 Description.** Successful execution of `CFI_select_part` updates the `base_addr`, `dim`, and `elem_len` members of  
6 the **C descriptor** with the address `result` for an array section for which each element is a part of the corresponding  
7 element of the array described by the **C descriptor** with the address `source`. The part shall be a component of a  
8 structure, a substring, or the real or imaginary part of a complex value.

9 If an error is detected, the **C descriptor** with the address `result` is not modified.

10 **4 Result Value.** The result is an error indicator.

11 **5 Example.** If `source` is already the address of a **C descriptor** for the Fortran array `A` declared with

```
12     TYPE, BIND(C) :: T
13         REAL(C_DOUBLE) :: X
14         COMPLEX(C_DOUBLE_COMPLEX) :: Y
15     END TYPE
16     TYPE(T) A(100)
```

17 the following code fragment establishes a **C descriptor** for the array `A%Y`:

```
18     typedef struct {
19         double x; double _Complex y;
20     } t;
21     CFI_CDESC_T(1) component;
22     CFI_cdesc_t * comp_cdesc = (CFI_cdesc_t *)&component;
23     CFI_index_t extent[] = { 100 };
24     (void)CFI_establish(comp_cdesc, NULL, CFI_attribute_other, CFI_type_double_Complex,
25                         sizeof(double _Complex), 1, extent);
26     (void)CFI_select_part(comp_cdesc, source, offsetof(t,y), 0);
```

### 27 **18.5.5.9 The CFI\_setpointer function**

28 **1 Synopsis.** Update a **C descriptor** for a Fortran pointer to be associated with the whole of a given object or to  
29 be disassociated.

```
30     int CFI_setpointer(CFI_cdesc_t *result, CFI_cdesc_t *source,
31                       const CFI_index_t lower_bounds[]);
```

### 32 **2 Formal Parameters.**

33 `result` shall be the address of a **C descriptor** for a Fortran pointer. It is updated using information from  
34 the `source` and `lower_bounds` arguments.

35 `source` shall be a null pointer or the address of a **C descriptor** for a nonallocatable nonpointer data object,  
36 an allocated allocatable object, or a data pointer object. If `source` is not a null pointer, the corres-  
37 ponding values of the `elem_len`, `rank`, and `type` members shall be the same in the **C descriptors**  
38 with the addresses `source` and `result`.

39 `lower_bounds` If `source` is not a null pointer and `source->rank` is nonzero, `lower_bounds` shall be a null pointer  
40 or the address of an array with at least `source->rank` elements.

1 3 **Result Value.** The result is an error indicator.

2 4 **Description.** Successful execution of `CFI_setpointer` updates the `base_addr` and `dim` members of the **C descriptor** with the address `result` as follows:

- 4 • if `source` is a null pointer or the address of a **C descriptor** for a disassociated pointer, the updated **C descriptor** describes a disassociated pointer;
- 5 • otherwise, the **C descriptor** with the address `result` becomes a **C descriptor** for the object described by the **C descriptor** with the address `source`, except that if `source->rank` is nonzero and `lower_bounds` is not a null pointer, the lower bounds are replaced by the values of the first `source->rank` elements of the `lower_bounds` array.

10 If an error is detected, the **C descriptor** with the address `result` is not modified.

11 5 **Example.** If `ptr` is already the address of a **C descriptor** for an array pointer of rank 1, the following code updates it to be a **C descriptor** for a pointer to the same array with lower bound 0.

```
13     CFI_index_t lower_bounds[1];
14     int ind;
15     lower_bounds[0] = 0;
16     ind = CFI_setpointer(ptr, ptr, lower_bounds);
```

## 17 18.6 Restrictions on C descriptors

18 1 A **C descriptor** shall not be initialized, updated, or copied other than by calling the functions specified in 18.5.5.

19 2 If the address of a **C descriptor** is a formal parameter that corresponds to a Fortran actual argument or a C actual argument that corresponds to a Fortran dummy argument,

- 21 • the **C descriptor** shall not be modified if either the corresponding dummy argument in the Fortran interface has the **INTENT (IN) attribute** or the **C descriptor** is for a nonallocatable nonpointer object, and
- 22 • the `base_addr` member of the **C descriptor** shall not be accessed before it is given a value if the corresponding dummy argument in the Fortran interface has the **POINTER** and **INTENT (OUT)** attributes.

### NOTE 18.31

In this context, modification refers to any change to the location or contents of the **C descriptor**, including establishment and update. The intent of these restrictions is that **C descriptors** remain intact at all times they are accessible to an active Fortran procedure, so that the Fortran code is not required to copy them.

## 25 18.7 Restrictions on formal parameters

26 1 Within a C function, an **allocatable** object shall be allocated or deallocated only by execution of the **CFI\_allocate** and **CFI\_deallocate** functions. A Fortran pointer can become associated with a target by execution of the **CFI\_allocate** function.

29 2 Calling **CFI\_allocate** or **CFI\_deallocate** for a **C descriptor** changes the allocation status of the Fortran variable it describes.

31 3 If the address of an object is the value of a formal parameter that corresponds to a nonpointer dummy argument in an interface with the **BIND attribute**, then

- 33 • if the dummy argument has the **INTENT (IN) attribute**, the object shall not be defined or become undefined, and

- 1       • if the dummy argument has the **INTENT (OUT) attribute**, the object shall not be referenced before it is  
2       defined.

- 3 4 If a formal parameter that is a pointer to **CFI\_cdesc\_t** corresponds to a dummy argument in an interoperable  
4 procedure interface, a pointer based on the **base\_addr** in that **C descriptor** shall not be used to access memory  
5 that is not part of the object described by the **C descriptor**.

## 6       **18.8 Restrictions on lifetimes**

- 7 1 A **C descriptor** of, or C pointer to, any part of a Fortran object becomes undefined under the same conditions  
8 that the association status of a Fortran pointer associated with that object would become undefined, and any  
9 further use of it is undefined behavior (ISO/IEC 9899:2011, 3.4.3).

- 10 2 A **C descriptor** whose address is a formal parameter that corresponds to a Fortran dummy argument becomes  
11 undefined on return from a call to the function from Fortran. If the dummy argument does not have either the  
12 **TARGET** or **ASYNCHRONOUS** attribute, all C pointers to any part of the object described by the **C descriptor**  
13 become undefined on return from the call, and any further use of them is undefined behavior.

- 14 3 If the address of a **C descriptor** is passed as an actual argument to a Fortran procedure, the lifetime (ISO/IEC  
15 9899:2011, 6.2.4) of the **C descriptor** shall not end before the return from the procedure call. If an object is passed  
16 to a Fortran procedure as a nonallocatable, nonpointer dummy argument, its lifetime shall not end before the  
17 return from the procedure call.

- 18 4 If the lifetime of a **C descriptor** for an allocatable object that was established by C ends before the program exits,  
19 the object shall be unallocated at that time.

- 20 5 If a Fortran pointer becomes associated with a data object defined by the companion processor, the association  
21 status of the Fortran pointer becomes undefined when the lifetime of that data object ends.

### **NOTE 18.32**

The following example illustrates how a **C descriptor** becomes undefined upon returning from a call to a C function.

```
REAL, TARGET :: X(1000), B
INTERFACE
  REAL FUNCTION CFUN(ARRAY) BIND(C, NAME="Cfun")
    REAL ARRAY(:)
  END FUNCTION
END INTERFACE
B = CFUN(X)
```

**Cfun** is a C function. Before or during the invocation of **Cfun**, the processor will create a **C descriptor** for the array **x**. On return from **Cfun**, that **C descriptor** will become undefined. In addition, because the dummy argument **ARRAY** does not have the **TARGET** or **ASYNCHRONOUS** attribute, a C pointer whose value was set during execution of **Cfun** to be the address of any part of **X** will become undefined.

## 22       **18.9 Interoperation with C global variables**

### 23       **18.9.1 General**

- 24 1 A C variable whose name has external linkage may interoperate with a **common block** or with a variable declared in  
25 the scope of a module. The **common block** or variable shall be specified to have the **BIND attribute**.



- 1 2 At most one variable that is associated with a particular C variable whose name has external linkage is permitted  
 2 to be declared within all the Fortran **program units** of a program. A variable shall not be initially defined by  
 3 more than one processor.
- 4 3 If a **common block** is specified in a **BIND statement**, it shall be specified in a **BIND statement** with the same **binding label** in each  
 5 **scoping unit** in which it is declared. A C variable whose name has external linkage interoperates with a **common block** that has been  
 6 specified in a **BIND statement** if
- 7 • the C variable is of a struct type and the variables that are members of the **common block** are **interoperable** with corresponding  
 8 components of the struct type, or
  - 9 • the **common block** contains a single variable, and the variable is **interoperable** with the C variable.
- 10 4 There does not have to be an associated C entity for a Fortran entity with the **BIND attribute**.

**NOTE 18.33**

The following are examples of the usage of the **BIND attribute** for variables and for a **common block**. The Fortran variables, C\_EXTERN and C2, interoperate with the C variables, c\_extern and myVariable, respectively. The Fortran **common blocks**, COM and SINGLE, interoperate with the C variables, com and single, respectively.

```

MODULE LINK_TO_C_VARS
  USE, INTRINSIC :: ISO_C_BINDING
  INTEGER(C_INT), BIND(C) :: C_EXTERN
  INTEGER(C_LONG) :: C2
  BIND(C, NAME='myVariable') :: C2

  COMMON /COM/ R, S
  REAL(C_FLOAT) :: R, S, T
  BIND(C) :: /COM/, /SINGLE/
  COMMON /SINGLE/ T
END MODULE LINK_TO_C_VARS

/* Global variables. */
int c_extern;
long myVariable;
struct { float r, s; } com;
float single;

```

## 11 18.9.2 Binding labels for common blocks and variables

- 12 1 The **binding label** of a variable or **common block** is a default character value that specifies the name by which the  
 13 variable or **common block** is known to the **companion processor**.
- 14 2 If a variable or **common block** has the **BIND attribute** with the NAME= specifier and the value of its expression,  
 15 after discarding leading and trailing blanks, has nonzero length, the variable or **common block** has this as its **binding**  
 16 **label**. The case of letters in the **binding label** is significant. If a variable or **common block** has the **BIND attribute**  
 17 specified without a NAME= specifier, the **binding label** is the same as the name of the entity using lower case  
 18 letters. Otherwise, the variable or **common block** has no **binding label**.
- 19 3 The **binding label** of a C variable whose name has external linkage is the same as the name of the C variable. A  
 20 Fortran variable or **common block** with the **BIND attribute** that has the same **binding label** as a C variable whose  
 21 name has external linkage is **linkage associated** (19.5.1.5) with that variable.

## 18.10 Interoperation with C functions

### 18.10.1 Definition and reference of interoperable procedures

- 1 A procedure that is [interoperable](#) may be defined either by means other than Fortran or by means of a Fortran subprogram, but not both. A C function that has an inline definition and no external definition is not considered to be defined in this sense.
- 2 If the procedure is defined by means other than Fortran,
  - it shall be describable by a C prototype that is [interoperable](#) with the [interface](#), and
  - if it is accessed using its [binding label](#), it shall
    - have a name that has external linkage as defined by 6.2.2 of ISO/IEC 9899:2011, and
    - have the same [binding label](#) as the [interface](#).
- 3 A reference to such a procedure causes the function described by the C prototype to be called as specified in ISO/IEC 9899:2011.
- 4 A reference in C to a procedure that has the [BIND attribute](#), has the same [binding label](#), and is defined by means of Fortran, causes the Fortran procedure to be invoked. A C function shall not invoke a function pointer whose value is the result of a reference to [C\\_FUNLOC](#) with a noninteroperable argument.
- 5 A procedure defined by means of Fortran shall not invoke `setjmp` or `longjmp` (ISO/IEC 9899:2011, 7.13). If a procedure defined by means other than Fortran invokes `setjmp` or `longjmp`, that procedure shall not cause any procedure defined by means of Fortran to be invoked. A procedure defined by means of Fortran shall not be invoked as a signal handler (ISO/IEC 9899:2011, 7.14.1).
- 6 If a procedure defined by means of Fortran and a procedure defined by means other than Fortran perform input/output operations on the same [external file](#), the results are processor dependent ([12.5.4](#)).
- 7 If the value of a C function pointer will be the result of a reference to [C\\_FUNLOC](#) with a noninteroperable argument, it is recommended that the C function pointer be declared to have the type `void (*)()`.

### 18.10.2 Binding labels for procedures

- 1 The [binding label](#) of a procedure is a default character value that specifies the name by which a procedure with the [BIND attribute](#) is known to the [companion processor](#).
- 2 If a procedure has the [BIND attribute](#) with the `NAME=` specifier and the value of its expression, after discarding leading and trailing blanks, has nonzero length, the procedure has this as its [binding label](#). The case of letters in the [binding label](#) is significant. If a procedure has the [BIND attribute](#) with no `NAME=` specifier, and the procedure is not a [dummy procedure](#), [internal procedure](#), or procedure pointer, then the [binding label](#) of the procedure is the same as the name of the procedure using lower case letters. Otherwise, the procedure has no [binding label](#).
 

C1807 A procedure defined in a submodule shall not have a [binding label](#) unless its [interface](#) is declared in the ancestor module.
- 3 The [binding label](#) for a C function whose name has external linkage is the same as the C function name.

#### NOTE 18.34

In the following sample, the [binding label](#) of `C_SUB` is `"c_sub"`, and the [binding label](#) of `C_FUNC` is `"C_funC"`.

```
SUBROUTINE C_SUB() BIND(C)
...

```

**NOTE 18.34 (cont.)**

```
END SUBROUTINE C_SUB
```

```
INTEGER(C_INT) FUNCTION C_FUNC() BIND(C, NAME="C_funC")
```

```
  USE, INTRINSIC :: ISO_C_BINDING
```

```
  ...
```

```
END FUNCTION C_FUNC
```

ISO/IEC 9899:2011 permits functions to have names that are not permitted as Fortran names; it also distinguishes between names that would be considered as the same name in Fortran. For example, a C name can begin with an underscore, and C names that differ in case are distinct names.

The specification of a [binding label](#) allows a program to use a Fortran name to refer to a procedure defined by a [companion processor](#).

### 1 18.10.3 Exceptions and IEEE arithmetic procedures

2 1 A procedure defined by means other than Fortran shall not use signal (ISO/IEC 9899:2011, 7.14.1) to change the  
3 [handling of any exception](#) that is being handled by the Fortran processor.

4 2 A procedure defined by means other than Fortran shall not alter the floating-point status ([17.7](#)) other than by  
5 setting an exception flag to signaling.

6 3 The values of the floating-point exception flags on entry to a procedure defined by means other than Fortran are  
7 processor dependent.

### 8 18.10.4 Asynchronous communication

9 1 Asynchronous communication for a Fortran variable with the [ASYNCHRONOUS attribute](#) occurs through the  
10 action of procedures defined by means other than Fortran. It is initiated by execution of an asynchronous  
11 communication initiation procedure and completed by execution of an asynchronous communication completion  
12 procedure. Between the execution of the initiation and completion procedures, any variable of which any part  
13 is associated with any part of the asynchronous communication variable is a pending communication affector.  
14 Whether a procedure is an asynchronous communication initiation or completion procedure is processor depend-  
15 ent.

16 2 Asynchronous communication is either input communication or output communication. For input communication,  
17 a pending communication affector shall not be referenced, become defined, become undefined, become associated  
18 with a dummy argument that has the [VALUE attribute](#), or have its pointer association status changed. For  
19 output communication, a pending communication affector shall not be redefined, become undefined, or have its  
20 pointer association status changed. The restrictions for asynchronous input communication are the same as for  
21 asynchronous input data transfer. The restrictions for asynchronous output communication are the same as for  
22 asynchronous output data transfer.

**NOTE 18.35**

Asynchronous communication can be used for nonblocking MPI calls such as MPI\_IRecv and MPI\_ISEND. For example,

```
REAL :: BUF(100,100)
```

```
... ! Code that involves BUF
```

```
BLOCK
```

```
  ASYNCHRONOUS :: BUF
```

```
  CALL MPI_IRecv(BUF,...REQ,...)
```

```
  ... ! Code that does not involve BUF
```

**NOTE 18.35 (cont.)**

```
CALL MPI_WAIT(REQ,...)
END BLOCK
... ! Code that involves BUF
```

In this example, there is asynchronous input communication and BUF is a pending communication affector between the two calls. MPI\_Irecv can return while the communication (reading values into BUF) is still underway. The intent is that the code between MPI\_Irecv and MPI\_WAIT can execute without waiting for this communication to complete.

Similar code with the call of MPI\_Irecv replaced by a call of MPI\_Isend is asynchronous output communication.

## 19 Scope, association, and definition

### 19.1 Scopes, identifiers, and entities

1 An entity is identified by an identifier.

2 The scope of

- a global identifier is a program (5.2.2),
- a local identifier is an [inclusive scope](#),
- an identifier of a [construct entity](#) is that construct (10.2.4, 11.1), and
- an identifier of a [statement entity](#) is that statement or part of that statement (6.3),

excluding any nested scope where the identifier is treated as the identifier of a different entity (19.3, 19.4).

3 An entity may be identified by

- an [image index](#) (3.85),
- a [name](#) (3.103),
- a [statement label](#) (3.135),
- an [external input/output unit](#) number (12.5),
- an identifier of a pending data transfer operation (12.6.2.9, 12.7),
- a submodule identifier (14.2.3),
- a [generic identifier](#) (3.79), or
- a [binding label](#) (3.15).

4 By means of association, an entity may be referred to by the same identifier or a different identifier in a different scope, or by a different identifier in the same scope.

### 19.2 Global identifiers

1 [Program units](#), [common blocks](#), [external procedures](#), entities with [binding labels](#), [external input/output units](#), pending data transfer operations, and [images](#) are global entities of a program. The name of a [common block](#) with no [binding label](#), [external procedure](#) with no [binding label](#), or [program unit](#) that is not a submodule is a global identifier. The submodule identifier of a submodule is a global identifier. A [binding label](#) of an entity of the program is a global identifier. An entity of the program shall not be identified by more than one [binding label](#).

2 The global identifier of an entity shall not be the same as the global identifier of any other entity. Furthermore, a [binding label](#) shall not be the same as the global identifier of any other global entity, ignoring differences in case. A processor may assign a global identifier to an entity that is not specified by this document to have a global identifier (such as an intrinsic procedure); in such a case, the processor shall ensure that this assigned global identifier differs from all other global identifiers in the program.

#### NOTE 19.1

An intrinsic module is not a [program unit](#), so a global identifier can be the same as the name of an intrinsic module.

**NOTE 19.2**

Submodule identifiers are global identifiers, but because they consist of a module name and a descendant submodule name, the name of a submodule can be the same as the name of another submodule so long as they do not have the same ancestor module.

**19.3 Local identifiers****19.3.1 Classes of local identifiers**

## 1 Identifiers of entities in the classes

- (1) except for statement or **construct** entities (19.4), named variables, **named constants**, named **procedure pointers**, named constructs, statement functions, **internal procedures**, **module procedures**, **dummy procedures**, **intrinsic procedures**, **external procedures** that have **binding labels**, intrinsic modules, **abstract interfaces**, **generic interfaces**, derived types, namelist groups, **external procedures** accessed via **USE**, and statement labels,
- (2) type parameters, components, and **type-bound procedure bindings**, in a separate class for each type,
- (3) **argument keywords**, in a separate class for each procedure with an **explicit interface**, and
- (4) common blocks that have **binding labels**

are local identifiers.

## 2 Within its scope, a local identifier of an entity of class (1) or class (4) shall not be the same as a global identifier used in that scope unless the global identifier

- is used only as the *use-name* of a **rename** in a **USE statement**,
- is a **common block** name (19.3.2),
- is an **external procedure** name that is also a generic name, or
- is an external function name and the **inclusive scope** is its defining subprogram (19.3.3).

## 3 Within its scope, a local identifier of one class shall not be the same as another local identifier of the same class, except that a generic name may be the same as the name of a procedure as explained in 15.4.3.4 or the same as the name of a derived type (7.5.10). A local identifier of one class may be the same as a local identifier of another class.

**NOTE 19.3**

An intrinsic procedure is inaccessible by its own name in a **scoping unit** that uses the same name as a local identifier of class (1) for a different entity. For example, in the program fragment

```

SUBROUTINE SUB
  ...
  A = SIN (K)
  ...
CONTAINS
  FUNCTION SIN (X)
    ...
  END FUNCTION SIN
END SUBROUTINE SUB

```

any reference to function SIN in subroutine SUB refers to the internal function SIN, not to the intrinsic function of the same name.

1 4 A local identifier identifies an entity in a scope and may be used to identify an entity in another scope except in  
2 the following cases.

- 3 • The name that appears as a *subroutine-name* in a *subroutine-stmt* has limited use within the scope estab-  
4 lished by the *subroutine-stmt*. It can be used to identify recursive references of the subroutine or to identify  
5 a **common block** (the latter is possible only for internal and module subroutines).
- 6 • The name that appears as a *function-name* in a *function-stmt* has limited use within the scope established  
7 by that *function-stmt*. It can be used to identify the **function result**, to identify recursive references of the  
8 function, or to identify a **common block** (the latter is possible only for internal and module functions).
- 9 • The name that appears as an *entry-name* in an *entry-stmt* has limited use within the scope of the subprogram in which  
10 the *entry-stmt* appears. It can be used to identify the **function result** if the subprogram is a function, to identify recursive  
11 references, or to identify a **common block** (the latter is possible only if the *entry-stmt* is in a module subprogram).

### 12 19.3.2 Local identifiers that are the same as common block names

13 1 A name that identifies a **common block** in a **scoping unit** shall not be used to identify a constant or an intrinsic procedure in that  
14 **scoping unit**. If a local identifier of class (1) is also the name of a **common block**, the appearance of that name in any context other  
15 than as a **common block** name in a **BIND**, **COMMON**, or **SAVE** statement is an appearance of the local identifier.

### 16 19.3.3 Function results

17 1 For each **FUNCTION statement** or **ENTRY statement** in a function subprogram, there is a **function result**. A **function**  
18 **result** is either a variable or a **procedure pointer**, and thus the name of a **function result** is a local identifier of  
19 class (1).

### 20 19.3.4 Components, type parameters, and bindings

21 1 A component name has the scope of its derived-type definition. Outside the type definition, it may also appear  
22 within a designator of a component of a structure of that type or as a **component keyword** in a **structure**  
23 **constructor** for that type.

24 2 A type parameter name has the scope of its derived-type definition. Outside the derived-type definition, it may  
25 also appear as a type parameter keyword in a *derived-type-spec* for the type or as the *type-param-name* of a  
26 *type-param-inquiry*.

27 3 The **binding name** (7.5.5) of a **type-bound procedure** has the scope of its derived-type definition. Outside of the  
28 derived-type definition, it may also appear as the *binding-name* in a procedure reference.

29 4 A generic **binding** for which the *generic-spec* is not a *generic-name* has a scope that consists of all **scoping units**  
30 in which an entity of the type is accessible.

31 5 A component name or **binding name** may appear only in a scope in which it is accessible.

32 6 The accessibility of components and **bindings** is specified in 7.5.4.8 and 7.5.5.

### 33 19.3.5 Argument keywords

34 1 As an **argument keyword**, a **dummy argument** name in an **internal procedure**, **module procedure**, or an **interface**  
35 **body** has a scope of the **scoping unit** of the **host** of the procedure or **interface body**. As an **argument keyword**,  
36 the name of a **dummy argument** of a procedure declared by a **procedure declaration statement** that specifies an  
37 **explicit interface** has a scope of the **scoping unit** containing the **procedure declaration statement**. It may appear  
38 only in a procedure reference for the procedure of which it is a **dummy argument**. If the procedure is accessible

1 in another [scoping unit](#) by [use](#) or [host](#) association ([19.5.1.3](#), [19.5.1.4](#)), the [argument keyword](#) is accessible for  
2 procedure references for that procedure in that [scoping unit](#).

3 2 A [dummy argument](#) name in an [intrinsic](#) procedure has a scope as an [argument keyword](#) of the [scoping unit](#)  
4 in which the reference to the procedure occurs. As an [argument keyword](#), it may appear only in a procedure  
5 reference for the procedure of which it is a [dummy argument](#).

## 6 19.4 Statement and construct entities

7 1 A variable that appears as a [data-i-do-variable](#) in a [DATA statement](#) or an [ac-do-variable](#) in an array constructor,  
8 as a dummy argument in a [statement function statement](#), or as an [index-name](#) in a [FORALL statement](#) is a [statement entity](#). A  
9 variable that appears as an [index-name](#) in a [FORALL](#) or [DO CONCURRENT](#), as an [associate-name](#) in a [SELECT](#)  
10 [TYPE](#) or [ASSOCIATE](#) construct, or as a [coarray-name](#) in a [codimension-decl](#) in a [CHANGE TEAM](#) construct  
11 is a [construct entity](#). A variable that has [LOCAL](#) or [LOCAL\\_INIT](#) locality in a [DO CONCURRENT](#) construct  
12 is a [construct entity](#). An entity that is explicitly declared in the specification part of a [BLOCK construct](#), other  
13 than only in [ASYNCHRONOUS](#) and [VOLATILE](#) statements, is a [construct entity](#). Two [construct entities](#) of the  
14 same construct shall not have the same identifier.

15 2 Even if the name of a statement entity is the same as another identifier and the statement is in the scope of that  
16 identifier, within the scope of the statement entity the name is interpreted as that of the statement entity.

17 3 The name of a statement entity shall not be the same as an accessible global identifier or local identifier of class  
18 (1) ([19.3.1](#)), except for a [common block](#) name or a scalar variable name. Within the scope of a [statement entity](#),  
19 another [statement entity](#) shall not have the same name.

20 4 The name of a [data-i-do-variable](#) in a [DATA statement](#) or an [ac-do-variable](#) in an array constructor has a scope  
21 of its [data-implied-do](#) or [ac-implied-do](#). It is a scalar variable. If [integer-type-spec](#) appears in [data-implied-do](#) or  
22 [ac-implied-do-control](#) it has the specified type and type parameters; otherwise it has the type and type parameters  
23 that it would have if it were the name of a variable in the innermost executable construct or [scoping unit](#) that  
24 includes the [DATA statement](#) or array constructor, and this type shall be integer type. It has no other attributes.  
25 The appearance of a name as a [data-i-do-variable](#) of an implied DO in a [DATA statement](#) or an [ac-do-variable](#)  
26 in an array constructor is not an implicit declaration of a variable whose scope is the [scoping unit](#) that contains  
27 the statement.

28 5 The name of a variable that appears as an [index-name](#) in a [DO CONCURRENT construct](#), [FORALL statement](#), or  
29 [FORALL construct](#) has a scope of the statement or construct. It is a scalar variable. If [integer-type-spec](#) appears in  
30 [concurrent-header](#) it has the specified type and type parameters; otherwise it has the type and type parameters  
31 that it would have if it were the name of a variable in the innermost executable construct or [scoping unit](#) that  
32 includes the [DO CONCURRENT](#) or [FORALL](#), and this type shall be integer type. It has no other attributes.  
33 The appearance of a name as an [index-name](#) in a [DO CONCURRENT construct](#), [FORALL statement](#), or [FORALL](#)  
34 [construct](#) is not an implicit declaration of a variable whose scope is the [scoping unit](#) that contains the statement or  
35 construct.

36 6 A variable that has [LOCAL](#) or [LOCAL\\_INIT](#) locality in a [DO CONCURRENT construct](#) has the scope of that  
37 construct. Its attributes are specified in [11.1.7.5](#).

38 7 If [integer-type-spec](#) does not appear in a [concurrent-header](#), an [index-name](#) shall not be the same as an accessible  
39 global identifier, local identifier, or identifier of an outer construct entity, except for a [common block](#) name or



1 a scalar variable name. An *index-name* of a contained **DO CONCURRENT** construct, **FORALL** statement, or  
2 **FORALL** construct shall not be the same as an *index-name* of any of its containing **DO CONCURRENT** or **FORALL**  
3 constructs.

4 8 The **associate name** of a **SELECT TYPE** construct has a separate scope for each block of the construct. Within  
5 each block, it has the **declared type**, **dynamic type**, type parameters, **rank**, and bounds specified in 11.1.11.2.

6 9 The **associate names** of an **ASSOCIATE** construct have the scope of the block. They have the **declared type**,  
7 **dynamic type**, type parameters, **rank**, and bounds specified in 11.1.3.2.

8 10 The **associate names** of a **CHANGE TEAM** construct have the scope of the block. They have the **declared type**,  
9 **dynamic type**, type parameters, **rank**, **corank**, **bounds**, and **cobounds** specified in 11.1.5.

10 11 The name of a variable that appears as a dummy argument in a **statement function statement** has a scope of the statement in which  
11 it appears. It is a scalar that has the type and type parameters that it would have if it were the name of a variable in the **scoping**  
12 **unit** that includes the **statement function**; it has no other attributes.

## 13 19.5 Association

### 14 19.5.1 Name association

#### 15 19.5.1.1 Forms of name association

16 1 There are five forms of **name association**: **argument association**, **use association**, **host association**, **linkage asso-**  
17 **ciation**, and **construct association**. **Argument**, **use**, and **host** association provide mechanisms by which entities  
18 known in one scope may be accessed in another scope.

#### 19 19.5.1.2 Argument association

20 1 The rules governing **argument association** are given in Clause 15. As explained in 15.5, execution of a procedure  
21 reference establishes a correspondence between each **actual argument** and a **dummy argument** and thus an associ-  
22 ation between each present **dummy argument** and its **effective argument**. Argument association can be sequence  
23 association (15.5.2.11).

24 2 The name of the **dummy argument** may be different from the name, if any, of its **effective argument**. The **dummy**  
25 **argument** name is the name by which the **effective argument** is known, and by which it may be accessed, in the  
26 referenced procedure.

#### NOTE 19.4

An **effective argument** can be a nameless data entity, such as the result of evaluating an expression that is not simply a variable or constant.

27 3 Upon termination of execution of a procedure reference, all **argument associations** established by that reference  
28 are terminated. A **dummy argument** of that procedure can be associated with an entirely different **effective**  
29 **argument** in a subsequent invocation of the procedure.

### 19.5.1.3 Use association

1 **Use association** is the association of names in different scopes specified by a **USE statement**. The rules governing  
2 **use association** are given in 14.2.2. They allow for renaming of entities being accessed. **Use association** allows  
3 access in one scope to entities defined or declared in another scope; it remains in effect throughout the execution  
4 of the program.

### 19.5.1.4 Host association

1 A nested **scoping unit** has access to named entities from its host as specified in 8.8. In the case of an **internal**  
2 **subprogram**, the access is to the entities in its **host instance**. The accessed entities are identified by the same  
3 identifier and have the same attributes as in the **host**, except that a local entity may have the **ASYNCHRONOUS**  
4 **attribute** even if the host entity does not, and a noncoarray local entity may have the **VOLATILE attribute**  
5 even if the host entity does not. The accessed entities are named data objects, derived types, **abstract interfaces**,  
6 procedures, **generic identifiers**, and namelist groups.

1 If an entity that is accessed by **use association** has the same nongeneric name as a host entity, the host entity is  
2 inaccessible by that name. The name of an **external procedure** that is given the **EXTERNAL attribute** (8.5.9)  
3 within the **scoping unit**, or a name that appears within the **scoping unit** as a *module-name* in a *use-stmt* is a  
4 global identifier; any entity of the **host** that has this as its nongeneric name is inaccessible by that name. A name  
5 that appears in the **scoping unit** as

- (1) a *function-name* in a *stmt-function-stmt* or in an *entity-decl* in a *type-declaration-stmt*, unless it is a global identifier,
- (2) an *object-name* in an *entity-decl* in a *type-declaration-stmt*, in a *pointer-stmt*, in a *save-stmt*, in an *allocatable-stmt*, or in a *target-stmt*,
- (3) a *type-param-name* in a *derived-type-stmt*,
- (4) a *named-constant* in a *named-constant-def* in a *parameter-stmt*,
- (5) a *coarray-name* in a *codimension-stmt*,
- (6) an *array-name* in a *dimension-stmt*,
- (7) a *variable-name* in a *common-block-object* in a *common-stmt*,
- (8) a **procedure pointer** given the **EXTERNAL attribute** in the **scoping unit**,
- (9) the name of a variable that is wholly or partially initialized in a *data-stmt*,
- (10) the name of an object that is wholly or partially equivalenced in an *equivalence-stmt*,
- (11) a *dummy-arg-name* in a *function-stmt*, in a *subroutine-stmt*, in an *entry-stmt*, or in a *stmt-function-stmt*,
- (12) a *result-name* in a *function-stmt* or in an *entry-stmt*,
- (13) the name of an entity declared by an **interface body**, unless it is a global identifier,
- (14) an *intrinsic-procedure-name* in an *intrinsic-stmt*,
- (15) a *namelist-group-name* in a *namelist-stmt*,
- (16) a *generic-name* in a *generic-spec* in an *interface-stmt*, or
- (17) the name of a named construct

is a local identifier in the **scoping unit** and any entity of the **host** that has this as its nongeneric name is inaccessible  
by that name by **host association**. If a **scoping unit** is the **host** of a derived-type definition or a subprogram that  
does not define a separate module procedure, the name of the derived type or of any procedure defined by the  
subprogram is a local identifier in the **scoping unit**; any entity of the **host** that has this as its nongeneric name is

1 inaccessible by that name. Local identifiers of a subprogram are not accessible to its [host](#).

**NOTE 19.5**

A name that appears in an [ASYNCHRONOUS](#) or [VOLATILE](#) statement is not necessarily the name of a local variable. In an [internal](#) or module procedure, if a variable that is accessible via [host association](#) is specified in an [ASYNCHRONOUS](#) or [VOLATILE](#) statement, that host variable is given the [ASYNCHRONOUS](#) or [VOLATILE](#) attribute in the local scope.

2 3 If a host entity is inaccessible only because a local variable with the same name is wholly or partially initialized  
3 in a [DATA statement](#), the local variable shall not be referenced or defined prior to the [DATA statement](#).

4 4 If a derived-type name of a [host](#) is inaccessible, data entities of that type or subobjects of such data entities still  
5 can be accessible.

**NOTE 19.6**

An [interface body](#) that is not a module procedure interface body accesses by [host association](#) only those entities made accessible by [IMPORT statements](#).

6 5 If an [external](#) or [dummy](#) procedure with an [implicit interface](#) is accessed via [host association](#), then it shall have  
7 the [EXTERNAL attribute](#) in the [host scoping unit](#); if it is invoked as a function in the inner [scoping unit](#), its type  
8 and type parameters shall be established in the [host scoping unit](#). The type and type parameters of a function  
9 with the [EXTERNAL attribute](#) are established in a [scoping unit](#) if that [scoping unit](#) explicitly declares them,  
10 invokes the function, accesses the function from a module, or accesses the function from its [host](#) where its type  
11 and type parameters are established.

12 6 If an intrinsic procedure is accessed via [host association](#), then it shall be established to be intrinsic in the [host](#)  
13 [scoping unit](#). An intrinsic procedure is established to be intrinsic in a [scoping unit](#) if that [scoping unit](#) explicitly  
14 gives it the [INTRINSIC attribute](#), invokes it as an intrinsic procedure, accesses it from a module, or accesses it  
15 from its [host](#) where it is established to be intrinsic.

**NOTE 19.7**

A host subprogram and an internal subprogram can contain the same and differing use-associated entities, as illustrated in the following example.

```

MODULE B; REAL BX, Q; INTEGER IX, JX; END MODULE B
MODULE C; REAL CX; END MODULE C
MODULE D; REAL DX, DY, DZ; END MODULE D
MODULE E; REAL EX, EY, EZ; END MODULE E
MODULE F; REAL FX; END MODULE F
MODULE G; USE F; REAL GX; END MODULE G
PROGRAM A
USE B; USE C; USE D
...
CONTAINS
  SUBROUTINE INNER_PROC (Q)
    USE C          ! Not needed
    USE B, ONLY: BX ! Entities accessible are BX, IX, and JX
                  ! if no other IX or JX

```

## NOTE 19.7 (cont.)

```

! is accessible to INNER_PROC
! Q is local to INNER_PROC,
! because Q is a dummy argument
USE D, X => DX ! Entities accessible are DX, DY, and DZ
! X is local name for DX in INNER_PROC
! X and DX denote same entity if no other
! entity DX is local to INNER_PROC
USE E, ONLY: EX ! EX is accessible in INNER_PROC, not in program A
! EY and EZ are not accessible in INNER_PROC
! or in program A
USE G ! FX and GX are accessible in INNER_PROC
...
END SUBROUTINE INNER_PROC
END PROGRAM A

```

Because program A contains the statement

```
USE B
```

all of the entities in module B, except for Q, are accessible in INNER\_PROC, even though INNER\_PROC contains the statement

```
USE B, ONLY: BX
```

The [USE statement](#) with the [ONLY](#) option means that this particular statement brings in only the entity named, not that this is the only variable from the module accessible in this [scoping unit](#).

## NOTE 19.8

For more examples of [host association](#), see subclause [C.13](#).

1 **19.5.1.5 Linkage association**

2 1 [Linkage association](#) occurs between a module variable that has the [BIND attribute](#) and the C variable with which  
3 it interoperates, or between a Fortran [common block](#) and the C variable with which it interoperates ([18.9](#)). Such association  
4 remains in effect throughout the execution of the program.

5 **19.5.1.6 Construct association**

6 1 Execution of a [SELECT RANK](#) or [SELECT TYPE](#) statement establishes an association between the selector and  
7 the [associate name](#) of the construct. Execution of an [ASSOCIATE](#) or [CHANGE TEAM statement](#) statement  
8 establishes an association between each selector and the corresponding [associate name](#) of the construct.

9 2 In an [ASSOCIATE](#) or [SELECT TYPE](#) construct, the following rules apply.

- 10 • If a selector is [allocatable](#), it shall be allocated; the [associate name](#) is associated with the data object and  
11 does not have the [ALLOCATABLE attribute](#).
- 12 • If a selector has the [POINTER attribute](#), it shall be associated; the [associate name](#) is associated with the  
13 [target](#) of the pointer and does not have the [POINTER attribute](#).

- 1 3 If the selector is a variable other than an [array section](#) having a [vector subscript](#), the association is with the data  
2 object specified by the selector; otherwise, the association is with the value of the selector expression, which is  
3 evaluated prior to execution of the block.
- 4 4 Each [associate name](#) remains associated with the corresponding selector throughout the execution of the executed  
5 block. Within the block, each selector is known by and may be accessed by the corresponding [associate name](#).  
6 On completion of execution of the construct, the association is terminated.

**NOTE 19.9**

The association between the [associate name](#) and a data object is established prior to execution of the block and is not affected by subsequent changes to variables that were used in subscripts or substring ranges in the [selector](#).

**19.5.2 Pointer association****19.5.2.1 General**

- 9 1 [Pointer association](#) between a [pointer](#) and a [target](#) allows the [target](#) to be referenced by a reference to the pointer.  
10 At different times during the execution of a program, a pointer may be [undefined](#), associated with different [targets](#)  
11 on its own [image](#), or be [disassociated](#). If a pointer is associated with a [target](#), the definition status of the pointer  
12 is either defined or undefined, depending on the definition status of the [target](#). If the pointer has deferred [type](#)  
13 [parameters](#) or shape, their values are assumed from the [target](#). If the pointer is [polymorphic](#), its [dynamic type](#)  
14 is assumed from the [dynamic type](#) of the [target](#).

**19.5.2.2 Pointer association status**

- 16 1 A pointer may have a pointer association status of associated, [disassociated](#), or [undefined](#). Its association status  
17 may change during execution of a program. Unless a pointer is initialized ([explicitly](#) or by [default](#)), it has an  
18 initial association status of [undefined](#). A pointer may be initialized to have an association status of [disassociated](#)  
19 or associated.

**NOTE 19.10**

A pointer from a module [program unit](#) might be accessible in a subprogram via [use association](#). Such pointers have a lifetime that is greater than [targets](#) that are declared in the subprogram, unless such [targets](#) are [saved](#). Therefore, if such a pointer is associated with a local [target](#), there is the possibility that when a procedure defined by the subprogram completes execution, the [target](#) will cease to exist, leaving the pointer “dangling”. This document considers such pointers to have an undefined association status. They are neither associated nor [disassociated](#). They cannot be used again in the program until their status has been reestablished. A processor is not required to detect when a pointer [target](#) ceases to exist.

**19.5.2.3 Events that cause pointers to become associated**

- 21 1 A pointer becomes associated when any of the following events occur.
- 22 (1) The pointer is allocated ([9.7.1](#)) as the result of the successful execution of an [ALLOCATE statement](#)  
23 referencing the pointer.
  - 24 (2) The pointer is pointer-assigned to a [target](#) ([10.2.2](#)) that is associated or is specified with the [TARGET](#)  
25 [attribute](#) and, if [allocatable](#), is allocated.

- 1 (3) The pointer is a **subobject** of an object that is allocated by an **ALLOCATE statement** in which  
2 **SOURCE=** appears and the corresponding **subobject** of *source-expr* is associated.
- 3 (4) The pointer is a dummy argument and its corresponding **actual argument** is not a pointer.
- 4 (5) The pointer is a **default-initialized subcomponent** of an object, the corresponding initializer is not a  
5 reference to the intrinsic function **NULL**, and
  - 6 (a) a procedure is invoked with this object as an **actual argument** corresponding to a nonpointer  
7 nonallocatable dummy argument with **INTENT (OUT)**,
  - 8 (b) a procedure with this object as an **unsaved** nonpointer nonallocatable local variable is invoked,
  - 9 (c) a **BLOCK construct** is entered and this object is an **unsaved** local nonpointer nonallocatable  
10 local variable of the **BLOCK construct**,
  - 11 or
  - 12 (d) this object is allocated other than by an **ALLOCATE statement** in which **SOURCE=** appears.

#### 13 19.5.2.4 Events that cause pointers to become disassociated

14 1 A pointer becomes **disassociated** when

- 15 (1) the pointer is nullified (9.7.2),
- 16 (2) the pointer is deallocated (9.7.3),
- 17 (3) the pointer is **pointer-assigned** (10.2.2) to a **disassociated** pointer,
- 18 (4) the pointer is a **subobject** of an object that is allocated by an **ALLOCATE statement** in which  
19 **SOURCE=** appears and the corresponding **subobject** of *source-expr* is **disassociated**,  
20 or
- 21 (5) the pointer is a **default-initialized subcomponent** of an object, the corresponding initializer is a  
22 reference to the intrinsic function **NULL**, and
  - 23 (a) a procedure is invoked with this object as an **actual argument** corresponding to a nonpointer  
24 nonallocatable dummy argument with **INTENT (OUT)**,
  - 25 (b) a procedure with this object as an **unsaved** nonpointer nonallocatable local variable is invoked,
  - 26 (c) a **BLOCK construct** is entered and this object is an **unsaved** local nonpointer nonallocatable  
27 local variable of the **BLOCK construct**,
  - 28 or
  - 29 (d) this object is allocated other than by an **ALLOCATE statement** in which **SOURCE=** appears.

#### 30 19.5.2.5 Events that cause the association status of pointers to become undefined

31 1 The association status of a pointer becomes undefined when

- 32 (1) the pointer is pointer-assigned to a **target** that has an undefined association status,
- 33 (2) the pointer is pointer-assigned to a **target** on a different **image**,
- 34 (3) the **target** of the pointer is deallocated other than through the pointer,
- 35 (4) the **target** of the pointer is a data object defined by the companion processor and the lifetime of that  
36 data object ends,
- 37 (5) the allocation transfer procedure (16.9.137) is executed, the pointer is associated with the argument  
38 FROM, and the argument TO does not have the **TARGET attribute**,

- 1 (6) completion of execution of an instance of a subprogram causes the pointer's **target** to become un-  
2 defined (item (3) of 19.6.6),
- 3 (7) completion of execution of a **BLOCK construct** causes the pointer's **target** to become undefined (item  
4 (23) of 19.6.6),
- 5 (8) execution of the **host instance** of a procedure pointer is completed,
- 6 (9) execution of an instance of a subprogram completes and the pointer is declared or accessed in the  
7 subprogram that defines the procedure if the pointer
- 8 (a) does not have the **SAVE attribute**,
- 9 (b) is not in **blank common**,
- 10 (c) is not in a named **common block** that is declared in any other **scoping unit** that is in execution,
- 11 (d) is not accessed by **host association**, and
- 12 (e) is not the result of a function declared to have the **POINTER attribute**,
- 13 (10) execution of an instance of a subprogram completes, the pointer is associated with a **dummy argument**  
14 of the procedure, and
- 15 (a) the **effective argument** does not have the **TARGET attribute** or is an **array section** with a  
16 **vector subscript**, or
- 17 (b) the **dummy argument** has the **VALUE attribute**,
- 18 (11) a **BLOCK construct** completes execution and the pointer is an **unsaved construct entity** of that  
19 **BLOCK construct**,
- 20 (12) a **DO CONCURRENT construct** is terminated and the pointer's association status was changed in  
21 more than one iteration of the construct,
- 22 (13) an iteration of a **DO CONCURRENT construct** completes and the pointer is associated with a  
23 variable of that construct that has **LOCAL** or **LOCAL\_INIT** locality,
- 24 (14) the pointer is a **subcomponent** of an object that is allocated and either
- 25 (a) the pointer is not **default-initialized** and **SOURCE=** does not appear, or
- 26 (b) **SOURCE=** appears and the association status of the corresponding **subcomponent** of *source-*  
27 *expr* is undefined,
- 28 (15) the pointer is a **subcomponent** of an object, the pointer is not **default-initialized**, and a procedure is  
29 invoked with this object as an **actual argument** corresponding to a dummy argument with **INTENT**  
30 **(OUT)**,
- 31 (16) a procedure is invoked with the pointer as an **actual argument** corresponding to a pointer dummy  
32 argument with **INTENT (OUT)**, or
- 33 (17) evaluation of an expression containing a function reference that need not be evaluated completes, if  
34 execution of that function would change the association status of the pointer.

#### 35 19.5.2.6 Other events that change the association status of pointers

- 36 1 When a pointer becomes associated with another pointer by **argument association**, **construct association**, or **host**  
37 **association**, the effects on its association status are specified in 19.5.5.
- 38 2 While two pointers are **name associated**, **storage associated**, or **inheritance associated**, if the association status of  
39 one pointer changes, the association status of the other changes accordingly.

1 3 The association status of a pointer object with the **VOLATILE attribute** might change by means not specified  
2 by the program.

### 3 **19.5.2.7 Pointer definition status**

4 1 The definition status of an associated pointer is that of its **target**. If a pointer is associated with a **definable**  
5 **target**, it may be defined or become undefined according to the rules for a variable (19.6). The definition status  
6 of a pointer that is not associated is undefined.

### 7 **19.5.2.8 Relationship between association status and definition status**

8 1 If the association status of a pointer is **disassociated** or undefined, the pointer shall not be referenced or dealloc-  
9 ated. Whatever its association status, a pointer always may be nullified, allocated, or pointer-assigned. A nullified  
10 pointer is **disassociated**. When a pointer is allocated, it becomes associated but undefined. When a pointer is  
11 pointer-assigned, its association and definition status become those of the specified **data-target** or **proc-target**.

## 12 **19.5.3 Storage association**

### 13 **19.5.3.1 General**

14 1 **Storage sequences** are used to describe relationships that exist among variables and **common blocks**. **Storage asso-**  
15 **ciation** is the association of two or more data objects that occurs when two or more **storage sequences** share or  
16 are aligned with one or more **storage units**.

### 17 **19.5.3.2 Storage sequence**

18 1 A **storage sequence** is a sequence of **storage units**. The size of a **storage sequence** is the number of **storage units**  
19 in the **storage sequence**. A **storage unit** is a **character storage unit**, a **numeric storage unit**, a **file storage unit**  
20 (12.3.5), or an **unspecified storage unit**. The sizes of the **numeric storage unit**, the **character storage unit** and the  
21 **file storage unit** are the values of constants in the ISO\_FORTRAN\_ENV intrinsic module (16.10.2).

22 2 In a storage association context

- 23 (1) a nonpointer scalar object that is default integer, default real, or default logical occupies a single  
24 **numeric storage unit**,
- 25 (2) a nonpointer scalar object that is double precision real or default complex occupies two **contiguous**  
26 **numeric storage units**,
- 27 (3) a default character nonpointer scalar object of character length *len* occupies *len* **contiguous character**  
28 **storage units**,
- 29 (4) if C character kind is not the same as default character kind a nonpointer scalar object of type char-  
30 acter with the C character kind (18.2.2) and character length *len* occupies *len* **contiguous unspecified**  
31 **storage units**,
- 32 (5) a nonpointer scalar object of **sequence type** occupies a sequence of **storage sequences** corresponding  
33 to the sequence of its **ultimate components**,
- 34 (6) a nonpointer scalar object of any type not specified in items (1)-(5) occupies a single **unspecified**  
35 **storage unit** that is different for each case and each set of type parameter values, and that is different  
36 from the **unspecified storage units** of item (4),
- 37 (7) a nonpointer array occupies a sequence of **contiguous storage sequences**, one for each array element,  
38 in array element order (9.5.3.2), and



1 (8) a **data pointer** occupies a single **unspecified storage unit** that is different from that of any nonpointer  
2 object and is different for each combination of type, type parameters, and **rank**. A **data pointer** that  
3 has the **CONTIGUOUS attribute** occupies a **storage unit** that is different from that of a **data pointer**  
4 that does not have the **CONTIGUOUS attribute**.

5 3 A sequence of **storage sequences** forms a **storage sequence**. The order of the **storage units** in such a composite  
6 **storage sequence** is that of the individual **storage units** in each of the constituent **storage sequences** taken in  
7 succession, ignoring any zero-sized constituent sequences.

8 4 Each **common block** has a **storage sequence** (8.10.2.2).

### 9 19.5.3.3 Association of storage sequences

10 1 Two nonzero-sized **storage sequences**  $s_1$  and  $s_2$  are storage associated if the  $i$ th **storage unit** of  $s_1$  is the same as  
11 the  $j$ th **storage unit** of  $s_2$ . This causes the  $(i + k)$ th **storage unit** of  $s_1$  to be the same as the  $(j + k)$ th **storage**  
12 **unit** of  $s_2$ , for each integer  $k$  such that  $1 \leq i + k \leq \text{size of } s_1$  and  $1 \leq j + k \leq \text{size of } s_2$  where *size of* measures  
13 the number of **storage units**.

14 2 Storage association also is defined between two zero-sized **storage sequences**, and between a zero-sized **storage**  
15 **sequence** and a **storage unit**. A zero-sized **storage sequence** in a sequence of **storage sequences** is storage associated  
16 with its successor, if any. If the successor is another zero-sized **storage sequence**, the two **sequences** are storage  
17 associated. If the successor is a nonzero-sized **storage sequence**, the zero-sized **sequence** is storage associated with  
18 the first **storage unit** of the successor. Two **storage units** that are each storage associated with the same zero-sized  
19 **storage sequence** are the same **storage unit**.

### 20 19.5.3.4 Association of scalar data objects

21 1 Two scalar data objects are storage associated if their **storage sequences** are storage associated. Two scalar entities  
22 are totally associated if they have the same **storage sequence**. Two scalar entities are partially associated if they  
23 are associated without being totally associated.

24 2 The definition status and value of a data object affects the definition status and value of any storage associated  
25 entity. An **EQUIVALENCE statement**, a **COMMON statement**, or an **ENTRY statement** can cause storage association of **storage**  
26 **sequences**.

27 3 An **EQUIVALENCE statement** causes storage association of data objects only within one **scoping unit**, unless one of the equivalenced  
28 entities is also in a **common block** (8.10.1.2, 8.10.2.2).

29 4 **COMMON statements** cause data objects in one **scoping unit** to become storage associated with data objects in another **scoping unit**.

30 5 A **common block** is permitted to contain a sequence of differing **storage units**. All **scoping units** that access named **common blocks**  
31 with the same name shall specify an identical sequence of **storage units**. **Blank common blocks** may be declared with differing sizes  
32 in different **scoping units**. For any two **blank common blocks**, the initial sequence of **storage units** of the longer **blank common block**  
33 shall be identical to the sequence of **storage units** of the shorter **common block**. If two **blank common blocks** are the same length,  
34 they shall have the same sequence of **storage units**.

35 6 An **ENTRY statement** in a function subprogram causes storage association of the **function results** that are variables.

36 7 Partial association shall exist only between

- 37 • an object that is default character or of **character sequence type** and an object that is default character or

1 of [character sequence type](#), or  
2 • an object that is default complex, double precision real, or of [numeric sequence type](#) and an object that is  
3 default integer, default real, default logical, double precision real, default complex, or of [numeric sequence](#)  
4 [type](#).

5 8 For noncharacter entities, partial association may occur only through the use of [COMMON](#), [EQUIVALENCE](#), or [ENTRY](#) statements.  
6 For character entities, partial association may occur only through [argument association](#) or the use of [COMMON](#) or  
7 [EQUIVALENCE](#) statements.

8 9 Partial association of character entities occurs when some, but not all, of the [storage units](#) of the entities are the  
9 same.

10 10 A [storage unit](#) shall not be [explicitly initialized](#) more than once in a program. [Explicit initialization](#) overrides  
11 [default initialization](#), and [default initialization](#) for an object of derived type overrides [default initialization](#) for  
12 a component of the object (7.5.4.6). [Default initialization](#) may be specified for a [storage unit](#) that is storage  
13 associated provided the objects supplying the [default initialization](#) are of the same type and type parameters,  
14 and supply the same value for the [storage unit](#).

#### 15 19.5.4 Inheritance association

16 1 [Inheritance association](#) occurs between components of the [parent component](#) and components [inherited](#) by type  
17 extension into an [extended type](#) (7.5.7.2). This association is persistent; it is not affected by the accessibility of  
18 the [inherited](#) components.

#### 19 19.5.5 Establishing associations

20 1 When an association is established between two entities by [argument association](#), [host association](#), or [construct](#)  
21 [association](#), certain properties of the [associating entity](#) become those of the pre-existing entity.

22 2 For [argument association](#), the pre-existing entity is the [effective argument](#) and the [associating entity](#) is the dummy  
23 argument.

24 3 For [host association](#), the [associating entity](#) is the entity in the contained [scoping unit](#). When a procedure is  
25 invoked, the pre-existing entity that participates in the association is the one from its [host instance](#) (15.6.2.4).  
26 Otherwise the pre-existing entity that participates in the association is the entity in the [host scoping unit](#).

27 4 For [construct association](#), the [associating entity](#) is identified by the [associate name](#) and the pre-existing entity is  
28 the selector.

29 5 When an association is established by [argument association](#), [host association](#), or [construct association](#), the fol-  
30 lowing applies.

31 • If the entities have the [POINTER attribute](#), the [pointer association](#) status of the [associating entity](#) becomes  
32 the same as that of the pre-existing entity. If the pre-existing entity has a [pointer association](#) status of  
33 associated, the [associating entity](#) becomes [pointer associated](#) with the same [target](#) and, if they are arrays,  
34 the bounds of the [associating entity](#) become the same as those of the pre-existing entity.

35 • If the [associating entity](#) has the [ALLOCATABLE attribute](#), its allocation status becomes the same as that  
36 of the pre-existing entity. If the pre-existing entity is allocated, the bounds (if it is an array), values of  
37 [deferred type parameters](#), definition status, and value (if it is defined) become the same as those of the

1 pre-existing entity. If the **associating entity** is **polymorphic** and the pre-existing entity is allocated, the  
2 **dynamic type** of the **associating entity** becomes the same as that of the pre-existing entity.

- 3 • If the **associating entity** is neither a pointer nor allocatable, its definition status, value (if it is defined), and  
4 **dynamic type** (if it is **polymorphic**) become the same as those of the pre-existing entity. If the entities are  
5 arrays and the association is not **argument association**, the bounds of the **associating entity** become the  
6 same as those of the pre-existing entity.
- 7 • If the **associating entity** is a pointer dummy argument and the pre-existing entity is a nonpointer **actual**  
8 **argument** the **associating entity** becomes **pointer associated** with the pre-existing entity and, if the entities  
9 are arrays, the bounds of the **associating entity** become the same as those of the pre-existing entity.

## 10 19.6 Definition and undefinition of variables

### 11 19.6.1 Definition of objects and subobjects

12 1 A variable may be defined or may be undefined and its definition status may change during execution of a  
13 program. An action that causes a variable to become undefined does not imply that the variable was previously  
14 defined. An action that causes a variable to become defined does not imply that the variable was previously  
15 undefined.

16 2 Arrays, including sections, and variables of derived, character, or complex type are objects that consist of zero  
17 or more subobjects. Associations may be established between variables and subobjects and between subobjects  
18 of different variables. These subobjects may become defined or undefined.

19 3 An array is defined if and only if all of its elements are defined.

20 4 A derived-type scalar object is defined if and only if all of its nonpointer components are defined.

21 5 A complex or character scalar object is defined if and only if all of its subobjects are defined.

22 6 If an object is undefined, at least one (but not necessarily all) of its subobjects are undefined.

### 23 19.6.2 Variables that are always defined

24 1 Zero-sized arrays and zero-length strings are always defined.

### 25 19.6.3 Variables that are initially defined

26 1 The following variables are initially defined:

- 27 (1) variables specified to have initial values by **DATA statements**;
- 28 (2) variables specified to have initial values by **type declaration statements**;
- 29 (3) nonpointer **default-initialized subcomponents** of **saved** variables that do not have the **ALLOCAT-**  
30 **ABLE** or **POINTER** attribute;
- 31 (4) pointers specified to be initially associated with a variable that is initially defined;
- 32 (5) variables that are always defined;
- 33 (6) variables with the **BIND attribute** that are initialized by means other than Fortran.

**NOTE 19.11**

Fortran code:

```
module mod
  integer, bind(c,name="blivet") :: foo
end module mod
```

C code:

```
int blivet = 123;
```

In the above example, the Fortran variable foo is initially defined to have the value 123 by means other than Fortran.

## 1 19.6.4 Variables that are initially undefined

2 1 Variables that are not initially defined are initially undefined.

## 3 19.6.5 Events that cause variables to become defined

4 1 Variables become defined by the following events.

- 5 (1) Execution of an [intrinsic assignment statement](#) other than a [masked array assignment](#) or [FORALL](#)  
6 [assignment statement](#) causes the variable that precedes the equals to become defined.
- 7 (2) Execution of a [masked array assignment](#) or [FORALL](#) [assignment statement](#) might cause some or all of  
8 the array elements in the [assignment statement](#) to become defined ([10.2.3](#)).
- 9 (3) As execution of an [input statement](#) proceeds, each variable that is assigned a value from the input  
10 file becomes defined at the time that data are transferred to it. (See (4) in [19.6.6](#).) Execution of a  
11 [WRITE statement](#) whose [unit](#) specifier identifies an [internal file](#) causes each record that is written  
12 to become defined.
- 13 (4) Execution of a [DO statement](#) causes the DO variable, if any, to become defined.
- 14 (5) Beginning of execution of the action specified by an [io-implied-do](#) in a synchronous data transfer  
15 statement causes the [do-variable](#) to become defined.
- 16 (6) A reference to a procedure causes an entire dummy data object to become defined if the dummy data  
17 object does not have [INTENT \(OUT\)](#) and the entire [effective argument](#) is defined.  
18 A reference to a procedure causes a subobject of a dummy argument to become defined if the dummy  
19 argument does not have [INTENT \(OUT\)](#) and the corresponding subobject of the [effective argument](#)  
20 is defined.
- 21 (7) Execution of an input/output statement containing an [IOSTAT= specifier](#) causes the specified integer  
22 variable to become defined.
- 23 (8) Execution of a synchronous [input statement](#) containing a [SIZE= specifier](#) causes the specified integer  
24 variable to become defined.
- 25 (9) Execution of a wait operation ([12.7.1](#)) corresponding to an asynchronous [input statement](#) containing  
26 a [SIZE= specifier](#) causes the specified integer variable to become defined.
- 27 (10) Execution of an [INQUIRE statement](#) causes any variable that is assigned a value during the execution  
28 of the statement to become defined if no error condition exists.

- 1 (11) If an error, end-of-file, or end-of-record condition occurs during execution of an input/output state-  
2 ment that has an **IOMSG= specifier**, the *iomsg-variable* becomes defined.
- 3 (12) When a **character storage unit** becomes defined, all associated **character storage units** become defined.  
4 When a **numeric storage unit** becomes defined, all associated **numeric storage units** of the same type  
5 become defined. When an entity of double precision real type becomes defined, all totally associated  
6 entities of double precision real type become defined.  
7 When an **unspecified storage unit** becomes defined, all associated **unspecified storage units** become  
8 defined.
- 9 (13) When a default complex entity becomes defined, all partially associated default real entities become  
10 defined.
- 11 (14) When both parts of a default complex entity become defined as a result of partially associated default  
12 real or default complex entities becoming defined, the default complex entity becomes defined.
- 13 (15) When all components of a structure of a **numeric sequence type** or **character sequence type** become  
14 defined as a result of partially associated objects becoming defined, the structure becomes defined.
- 15 (16) Execution of a statement with a **STAT= specifier** causes the variable specified by the **STAT= specifier**  
16 to become defined.
- 17 (17) If an error condition occurs during execution of a statement that has an **ERRMSG= specifier**, the  
18 variable specified by the **ERRMSG= specifier** becomes defined.
- 19 (18) Allocation of a zero-sized array or zero-length character variable causes the array or variable to  
20 become defined.
- 21 (19) Allocation of an object that has a nonpointer **default-initialized subcomponent**, except by an **AL-**  
22 **LOCATE statement** with a **SOURCE= specifier**, causes that **subcomponent** to become defined.
- 23 (20) Successful execution of an **ALLOCATE statement** with a **SOURCE= specifier** causes a subobject of  
24 the allocated object to become defined if the corresponding subobject of the **SOURCE=** expression  
25 is defined.
- 26 (21) Invocation of a procedure causes any **automatic data object** of zero size or zero character length in  
27 that procedure to become defined.
- 28 (22) When a pointer becomes associated with a **target** that is defined, the pointer becomes defined.
- 29 (23) Invocation of a procedure that contains an **unsaved** nonpointer nonallocatable local variable causes  
30 all nonpointer **default-initialized subcomponents** of the object to become defined.
- 31 (24) Invocation of a procedure that has a nonpointer nonallocatable **INTENT (OUT)** dummy argument  
32 causes all nonpointer **default-initialized subcomponents** of the dummy argument to become defined.
- 33 (25) In a **DO CONCURRENT** or **FORALL** construct, the *index-name* becomes defined when the *index-*  
34 *name* value set is evaluated.
- 35 (26) In a **DO CONCURRENT construct**, a variable with **LOCAL\_INIT** locality becomes defined at the  
36 beginning of each iteration.
- 37 (27) An object with the **VOLATILE attribute** that is changed by a means not specified by the program  
38 might become defined (see 8.5.19).
- 39 (28) Execution of the **BLOCK statement** of a **BLOCK construct** that has an **unsaved** nonpointer non-  
40 allocatable local variable causes all nonpointer **default-initialized subcomponents** of the variable to  
41 become defined.
- 42 (29) Execution of an **OPEN statement** containing a **NEWUNIT= specifier** causes the specified integer  
43 variable to become defined.

- 1 (30) Execution of a **LOCK statement** containing an **ACQUIRED\_LOCK= specifier** causes the specified  
2 logical variable to become defined. If the logical variable becomes defined with the value true, the  
3 **lock variable** in the **LOCK statement** also becomes defined.
- 4 (31) Successful execution of a **LOCK statement** that does not contain an **ACQUIRED\_LOCK= specifier**  
5 causes the **lock variable** to become defined.
- 6 (32) Successful execution of an **UNLOCK statement** causes the **lock variable** to become defined.
- 7 (33) Failure of an **image** that locked a **lock variable** without unlocking it causes the **lock variable** to become  
8 defined.
- 9 (34) Successful execution of an **EVENT POST** or **EVENT WAIT** statement causes the **event variable** to  
10 become defined.

### 11 19.6.6 Events that cause variables to become undefined

12 1 Variables become undefined by the following events.

- 13 (1) When a scalar variable of **intrinsic type** becomes defined, all totally associated variables of different type become  
14 undefined. When a double precision scalar variable becomes defined, all partially associated scalar  
15 variables become undefined. When a scalar variable becomes undefined, all partially associated double  
16 precision scalar variables become undefined.
- 17 (2) If the evaluation of a function would cause a variable to become defined and if a reference to the  
18 function appears in an expression in which the value of the function is not needed to determine the  
19 value of the expression, the variable becomes undefined when the expression is evaluated.
- 20 (3) When execution of an instance of a subprogram completes,  
21 (a) its **unsaved** local variables become undefined,  
22 (b) **unsaved** variables in a named **common block** that appears in the subprogram become undefined if they have been  
23 defined or redefined, unless another active **scoping unit** is referencing the **common block**, and  
24 (c) a variable of type **C\_PTR** from the intrinsic module **ISO\_C\_BINDING** whose value is the **C**  
25 **address** of an **unsaved** local variable of the subprogram becomes undefined.
- 26 (4) When an error condition or end-of-file condition occurs during execution of an **input statement**, all of  
27 the variables specified by the input list or namelist group of the statement become undefined.
- 28 (5) When an error condition occurs during execution of an **output statement** in which the **unit** is an **internal**  
29 **file**, the **internal file** becomes undefined.
- 30 (6) When an error condition, end-of-file condition, or end-of-record condition occurs during execution of  
31 an input/output statement and the statement contains any *io-implied-dos*, all of the *do-variables* in  
32 the statement become undefined (12.11).
- 33 (7) Execution of a direct access **input statement** that specifies a record that has not been written previously  
34 causes all of the variables specified by the input list of the statement to become undefined.
- 35 (8) Execution of an **INQUIRE statement** might cause the **NAME=**, **RECL=**, and **NEXTREC=** variables  
36 to become undefined (12.10).
- 37 (9) When a **character storage unit** becomes undefined, all associated **character storage units** become un-  
38 defined.
- 39 When a **numeric storage unit** becomes undefined, all associated **numeric storage units** become undefined  
40 unless the undefinition is a result of defining an associated **numeric storage unit** of different type (see  
41 (1) above).

- 1           When an entity of double precision real type becomes undefined, all totally associated entities of double  
2           precision real type become undefined.
- 3           When an [unspecified storage unit](#) becomes undefined, all associated [unspecified storage units](#) become  
4           undefined.
- 5           (10) When an [allocatable](#) entity is deallocated, it becomes undefined.
- 6           (11) When the allocation transfer procedure ([16.9.137](#)) causes the allocation status of an [allocatable](#) entity  
7           to become unallocated, the entity becomes undefined.
- 8           (12) Successful execution of an [ALLOCATE statement](#) with no [SOURCE= specifier](#) causes a [subcomponent](#)  
9           of an allocated object to become undefined if [default initialization](#) has not been specified for that  
10           [subcomponent](#).
- 11           (13) Successful execution of an [ALLOCATE statement](#) with a [SOURCE= specifier](#) causes a subobject of  
12           the allocated object to become undefined if the corresponding subobject of the [SOURCE=](#) expression  
13           is undefined.
- 14           (14) Execution of an [INQUIRE statement](#) causes all inquiry specifier variables to become undefined if an  
15           error condition exists, except for any variable in an [IOSTAT=](#) or [IOMSG=](#) specifier.
- 16           (15) When a procedure is invoked
- 17           (a) an optional dummy argument that has no corresponding [actual argument](#) becomes undefined,  
18           (b) a dummy argument with [INTENT \(OUT\)](#) becomes undefined except for any nonpointer [default-](#)  
19           [initialized subcomponents](#) of the argument,  
20           (c) an [actual argument](#) corresponding to a dummy argument with [INTENT \(OUT\)](#) becomes un-  
21           defined except for any nonpointer [default-initialized subcomponents](#) of the argument,  
22           (d) a subobject of a dummy argument that does not have [INTENT \(OUT\)](#) becomes undefined if the  
23           corresponding subobject of the [effective argument](#) is undefined, and  
24           (e) a variable that is the [function result](#) of that procedure becomes undefined except for any of its  
25           nonpointer [default-initialized subcomponents](#).
- 26           (16) When the association status of a pointer becomes undefined or [disassociated](#) ([19.5.2.4](#), [19.5.2.5](#)), the  
27           pointer becomes undefined.
- 28           (17) When a [DO CONCURRENT construct](#) terminates, a variable that is defined or becomes undefined  
29           during more than one iteration of the construct becomes undefined.
- 30           (18) When execution of an iteration of a [DO CONCURRENT construct](#) completes, a construct entity of  
31           that construct which has [LOCAL](#) or [LOCAL\\_INIT](#) locality becomes undefined.
- 32           (19) Execution of an asynchronous [READ statement](#) causes all of the variables specified by the input list or  
33           [SIZE= specifier](#) to become undefined. Execution of an asynchronous namelist [READ statement](#) causes  
34           any variable in the namelist group to become undefined if that variable will subsequently be defined  
35           during the execution of the [READ statement](#) or the corresponding wait operation ([12.7.1](#)).
- 36           (20) When a variable with the [TARGET attribute](#) is deallocated, a variable of type [C\\_PTR](#) from the  
37           intrinsic module [ISO\\_C\\_BINDING](#) becomes undefined if its value is the [C address](#) of any part of the  
38           variable that is deallocated.
- 39           (21) When a pointer is deallocated, a variable of type [C\\_PTR](#) from the intrinsic module [ISO\\_C\\_BINDING](#)  
40           becomes undefined if its value is the [C address](#) of any part of the [target](#) that is deallocated.
- 41           (22) Execution of the allocation transfer procedure ([16.9.137](#)) where the argument TO does not have the  
42           [TARGET attribute](#) causes a variable of type [C\\_PTR](#) from the intrinsic module [ISO\\_C\\_BINDING](#) to  
43           become undefined if its value is the [C address](#) of any part of the argument FROM.

- 1 (23) When a **BLOCK construct** completes execution,
- 2     • its **unsaved** local variables become undefined, and
- 3     • a variable of type **C\_PTR** from the intrinsic module **ISO\_C\_BINDING**, whose value is the **C**
- 4     **address** of an **unsaved** local variable of the **BLOCK construct**, becomes undefined.
- 5 (24) When execution of the **host instance** of the **target** of a variable of type **C\_FUNPTR** from the intrinsic
- 6     module **ISO\_C\_BINDING** is completed by execution of a **RETURN** or **END** statement, the variable
- 7     becomes undefined.
- 8 (25) Execution of an intrinsic assignment of the type **C\_PTR** or **C\_FUNPTR** from the intrinsic module
- 9     **ISO\_C\_BINDING** in which the variable and *expr* are not on the same **image** causes the variable to
- 10     become undefined.
- 11 (26) An object with the **VOLATILE attribute** (8.5.19) might become undefined by means not specified by
- 12     the program.
- 13 (27) When a pointer becomes associated with a target that is undefined, the pointer becomes undefined.
- 14 (28) When an image fails during execution of a segment, a data object on a nonfailed image becomes
- 15     undefined if it is not a **lock variable** and it might become undefined by execution of a statement of
- 16     the segment other than an invocation of an **atomic subroutine** with the object as an **actual argument**
- 17     corresponding to the **ATOM dummy argument**.

**NOTE 19.12**

|                                                                                                       |
|-------------------------------------------------------------------------------------------------------|
| Execution of a <b>defined assignment statement</b> could leave all or part of the variable undefined. |
|-------------------------------------------------------------------------------------------------------|

**19.6.7 Variable definition context**

- 18
- 19 1 Some variables are prohibited from appearing in a syntactic context that would imply definition or undefinition
- 20 of the variable (8.5.10, 8.5.15, 15.7). The following are the contexts in which the appearance of a variable implies
- 21 such definition or undefinition of the variable:

- 22 (1) the *variable* of an *assignment-stmt*;
- 23 (2) a *do-variable* in a *do-stmt* or *io-implied-do*;
- 24 (3) an *input-item* in a *read-stmt*;
- 25 (4) a *variable-name* in a *namelist-stmt* if the *namelist-group-name* appears in a **NML= specifier** in a
- 26 *read-stmt*;
- 27 (5) an *internal-file-variable* in a *write-stmt*;
- 28 (6) an **IOSTAT=**, **SIZE=**, or **IOMSG=** specifier in an input/output statement;
- 29 (7) a specifier in an **INQUIRE statement** other than **FILE=**, **ID=**, and **UNIT=**;
- 30 (8) a **NEWUNIT=** specifier in an **OPEN statement**;
- 31 (9) a *stat-variable*, *allocate-object*, or *errmsg-variable*;
- 32 (10) an **actual argument** in a reference to a procedure with an **explicit interface** if the corresponding
- 33 **dummy argument** is not a pointer and has **INTENT (OUT)** or **INTENT (INOUT)**;
- 34 (11) a *variable* that is a *selector* in an **ASSOCIATE**, **CHANGE TEAM**, **SELECT RANK**, or **SELECT**
- 35 **TYPE** construct if the corresponding **associate name** or any subobject thereof appears in a variable
- 36 definition context;
- 37 (12) an *event-variable* in an **EVENT POST** or **EVENT WAIT** statement;
- 38 (13) a *lock-variable* in a **LOCK** or **UNLOCK** statement;



1 (14) a *scalar-logical-variable* in an ACQUIRED\_LOCK= specifier;

2 (15) a *team-variable* in a FORM TEAM statement.

3 2 If a reference to a function appears in a variable definition context the result of the function reference shall be a  
4 pointer that is associated with a *definable target*. That *target* is the variable that becomes defined or undefined.

### 5 **19.6.8 Pointer association context**

6 1 Some pointers are prohibited from appearing in a syntactic context that would imply alteration of the *pointer*  
7 *association* status (19.5.2.2, 8.5.10, 8.5.15, 15.7). The following are the contexts in which the appearance of a  
8 pointer implies such alteration of its *pointer association* status:

9 • a *pointer-object* in a *nullify-stmt*;

10 • a *data-pointer-object* or *proc-pointer-object* in a *pointer-assignment-stmt*;

11 • an *allocate-object* in an *allocate-stmt* or *deallocate-stmt*;

12 • an *actual argument* in a reference to a procedure if the corresponding *dummy argument* is a pointer with  
13 the *INTENT (OUT)* or *INTENT (INOUT)* attribute.

1

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# Annex A

(Informative)

## Processor Dependencies

### A.1 Unspecified Items

1 This document does not specify the following:

- the properties excluded in 1;
- a processor's error detection capabilities beyond those listed in 4.2;
- which additional intrinsic procedures or modules a processor provides (4.2);
- the number and kind of **companion processors** (5.5.7);
- the number of representation methods and associated kind type parameter values of the intrinsic types (7.4), except that there shall be at least two representation methods for type real, and a representation method of type complex that corresponds to each representation method for type real.

### A.2 Processor Dependencies

1 According to this document, the following are processor dependent:

- the order of evaluation of the **specification expressions** within the specification part of an invoked Fortran procedure (5.3.5);
- how soon an **image** terminates if another **image** initiates **error termination** (5.3.5);
- the value of a reference to to a **coindexed object** on a **failed image** (5.3.6);
- the conditions that cause an image to fail (5.3.6);
- whether the processor supports a concept of process exit status, and if so, the process exit status on program termination (5.3.7);
- the mechanism of a **companion processor**, and the means of selecting between multiple **companion processors** (5.5.7);
- the processor character set (6.1);
- the means for specifying the source form of a **program unit** (6.3);
- the maximum number of characters allowed on a source line containing characters not of default kind (6.3.2, 6.3.3);
- the maximum depth of nesting of include lines (6.4);
- the interpretation of the **char-literal-constant** in the include line (6.4);
- the set of values supported by an intrinsic type, other than logical (7.1.3);
- the **kind type parameter** value of a complex literal constant, if both the real part and imaginary part are of type real with the same precision, but have different **kind type parameter** values (7.4.3.3);
- the kind of a character length type parameter (7.4.4.1);
- the blank padding character for nondefault character kind (7.4.4.2)

- 1 • whether particular control characters can appear within a character literal constant in fixed source form  
2 (7.4.4.3);
- 3 • the *collating sequence* for each character set (7.4.4.4);
- 4 • the order of *finalization* of components of objects of derived type (7.5.6.2);
- 5 • the order of *finalization* when several objects are finalized as the consequence of a single event (7.5.6.2);
- 6 • whether and when an object is finalized if it is allocated by pointer allocation and it later becomes un-  
7 reachable due to all pointers associated with the object having their pointer association status changed  
8 (7.5.6.3);
- 9 • whether an object is finalized by a deallocation in which an error condition occurs (7.5.6.3);
- 10 • the kind type parameter of each enumeration and its enumerators (7.6);
- 11 • whether an array is contiguous, except as specified in 8.5.7;
- 12 • the set of error conditions that can occur in *ALLOCATE* and *DEALLOCATE* statements (9.7.1, 9.7.3);
- 13 • the allocation status of a variable after evaluation of an expression if the evaluation of a function would  
14 change the allocation status of the variable and if a reference to the function appears in the expression in  
15 which the value of the function is not needed to determine the value of the expression (9.7.1.3);
- 16 • the order of deallocation when several objects are deallocated by a *DEALLOCATE statement* (9.7.3);
- 17 • the order of deallocation when several objects are deallocated due to the occurrence of an event described  
18 in 9.7.3.2;
- 19 • whether an allocated allocatable subobject is deallocated when an error condition occurs in the deallocation  
20 of an object (9.7.3.2);
- 21 • the positive integer values assigned to the *stat-variable* in a *STAT= specifier* as the result of an error  
22 condition (9.7.4, 11.6.11);
- 23 • the allocation status or pointer association status of an *allocate-object* if an error occurs during execution  
24 of an *ALLOCATE* or *DEALLOCATE* statement (9.7.4);
- 25 • the value assigned to the *errmsg-variable* in an *ERRMSG= specifier* as the result of an error condition  
26 (9.7.5, 11.6.11);
- 27 • the kind type parameter value of the result of a numeric intrinsic binary operation where
  - 28 – both operands are of type integer but with different kind type parameters, and the decimal exponent  
29 ranges are the same,
  - 30 – one operand is of type real or complex and the other is of type real or complex with a different kind  
31 type parameter, and the decimal precisions are the same,
- 32 and for a logical intrinsic binary operation where the operands have different kind type parameters (10.1.9.3);
- 33 • the character assigned to the variable in an *intrinsic assignment statement* if the kind of the expression is  
34 different and the character is not representable in the kind of the variable (10.2.1.3);
- 35 • the order of evaluation of the *specification expressions* within the specification part of a *BLOCK construct*  
36 when the construct is executed (11.1.4);
- 37 • the pointer association status of a pointer that has its pointer association changed in more than one iteration  
38 of a *DO CONCURRENT construct*, on termination of the construct (11.1.7);
- 39 • the ordering between records written by different iterations of a *DO CONCURRENT construct* if the records  
40 are written to a file connected for sequential access by more than one iteration (11.1.7);
- 41 • the manner in which the stop code of a *STOP* or *ERROR STOP* statement is made available (11.4);
- 42 • the mechanisms available for creating dependencies for cooperative synchronization (11.6.5);

- 1 • the value of the count of the event variable in an `EVENT POST` or `EVENT WAIT` statement if an error  
2 condition occurs (11.6.7, 11.6.8);
- 3 • the `image index` value established for each image in a `team` by a `FORM TEAM` statement without a  
4 `NEW_INDEX= specifier` (11.6.9);
- 5 • the set of error conditions that can occur in `image control statements` (11.6.11);
- 6 • the relationship between the `file storage units` when viewing a file as a `stream file`, and the records when  
7 viewing that file as a `record file` (12);
- 8 • whether particular control characters can appear in a formatted `record` or a formatted `stream file` (12.2.2);
- 9 • the form of values in an unformatted record (12.2.3);
- 10 • at any time, the set of allowed access methods, set of allowed forms, set of allowed actions, and set of  
11 allowed record lengths for a `file` (12.3);
- 12 • the set of allowable names for a `file` (12.3);
- 13 • whether a named file on one `image` is the same as a file with the same name on another `image` (12.3.1);
- 14 • the set of `external files` that exist for a program (12.3.2);
- 15 • the relationship between positions of successive `file storage units` in an `external file` that is connected for  
16 formatted stream access (12.3.3.4);
- 17 • the `external unit` preconnected for sequential formatted input and identified by an asterisk or the `named`  
18 `constant` `INPUT_UNIT` of the `ISO_FORTRAN_ENV` intrinsic module (12.5);
- 19 • the `external unit` preconnected for sequential formatted output and identified by an asterisk or the `named`  
20 `constant` `OUTPUT_UNIT` of the `ISO_FORTRAN_ENV` intrinsic module (12.5);
- 21 • the `external unit` preconnected for sequential formatted output and identified by the `named constant` `ER-`  
22 `ROR_UNIT` of the `ISO_FORTRAN_ENV` intrinsic module, and whether this unit is the same as `OUT-`  
23 `PUT_UNIT` (12.5);
- 24 • at any time, the set of `external units` that exist for an `image` (12.5.3);
- 25 • whether a `unit` can be connected to a file that is also connected to a C stream (12.5.4);
- 26 • whether a file can be connected to more than one `unit` at the same time (12.5.4);
- 27 • the effect of performing input/output operations on multiple units while they are connected to the same  
28 `external file` (12.5.4);
- 29 • the result of performing input/output operations on a `unit` connected to a file that is also connected to a C  
30 stream (12.5.4);
- 31 • whether the files connected to the `units` `INPUT_UNIT`, `OUTPUT_UNIT`, and `ERROR_UNIT` correspond  
32 to the predefined C text streams standard input, standard output, and standard error, respectively (12.5.4);
- 33 • the results of performing input/output operations on an `external file` both from Fortran and from a procedure  
34 defined by means other than Fortran (12.5.4);
- 35 • the default value for the `ACTION= specifier` in an `OPEN` statement (12.5.6.4);
- 36 • the encoding of a file opened with `ENCODING='DEFAULT'` (12.5.6.9);
- 37 • the file connected by an `OPEN` statement with `STATUS='SCRATCH'` (12.5.6.10);
- 38 • the interpretation of case in a file name (12.5.6.10, 12.10.2.2);
- 39 • the position of a file after executing an `OPEN` statement with a `POSITION= specifier` of ASIS, when the  
40 file previously existed but was not connected (12.5.6.14);
- 41 • the default value for the `RECL= specifier` in an `OPEN` statement (12.5.6.15);
- 42 • the effect of `RECL=` on a record containing any nondefault characters (12.5.6.15);

- 1       • the default input/output rounding mode (12.5.6.16);
- 2       • the default sign mode (12.5.6.17);
- 3       • the file status when `STATUS='UNKNOWN'` is specified in an `OPEN` statement (12.5.6.18);
- 4       • the value assigned to the variable in the `ID= specifier` in an asynchronous `data transfer statement` when
- 5       execution of the statement is successfully completed (12.6.2.9);
- 6       • whether `POS=` is permitted with particular files, and whether `POS=` can position a particular file to a
- 7       position prior to its current position (12.6.2.11);
- 8       • the form in which a single value of derived type is treated in an unformatted input/output statement if the
- 9       `effective item` is not processed by a `defined input/output procedure` (12.6.3);
- 10       • the result of unformatted input when the value stored in the file has a different type or `type parameters`
- 11       from that of the input list item (12.6.4.5.2);
- 12       • the negative value of the `unit` argument to a `defined input/output procedure` if the parent `data transfer`
- 13       `statement` accesses an `internal file` (12.6.4.8.2);
- 14       • the manner in which the processor makes the value of the `iomsg` argument of a `defined input/output`
- 15       procedure available if the procedure assigns a nonzero value to the `iostat` argument and the processor
- 16       therefore `terminates execution` of the program (12.6.4.8.2);
- 17       • the action caused by the flush operation, whether the processor supports the flush operation for the specified
- 18       `unit`, and the negative value assigned to the `IOSTAT=` variable if the processor does not support the flush
- 19       operation for the specified `unit` (12.9);
- 20       • the case of characters assigned to the variable in a `NAME= specifier` in an `INQUIRE statement` (12.10.2.15);
- 21       • which of the connected external unit numbers is assigned to the `scalar-int-variable` in the `NUMBER=`
- 22       `specifier` in an `INQUIRE by file statement`, if more than one unit on an `image` is connected to the file
- 23       (12.10.2.18);
- 24       • the value of the variable in a `POSITION= specifier` in an `INQUIRE statement` if the file has been repositioned
- 25       since connection (12.10.2.23);
- 26       • the relationship between file size and the data stored in records in a sequential or direct access file
- 27       (12.10.2.30);
- 28       • the number of `file storage units` needed to store data in an unformatted file (12.10.3);
- 29       • the set of error conditions that can occur in input/output statements (12.11.1);
- 30       • when an input/output error condition occurs or is detected (12.11.1);
- 31       • the positive integer value assigned to the variable in an `IOSTAT= specifier` as the result of an error condition
- 32       (12.11.5);
- 33       • the value assigned to the variable in an `IOMSG= specifier` as the result of an error condition (12.11.6);
- 34       • the result of output of non-representable characters to a Unicode file (13.7.1);
- 35       • the interpretation of the optional non-blank characters within the parentheses of a real NaN input field
- 36       (13.7.2.3.2);
- 37       • the interpretation of a sign in a NaN input field (13.7.2.3.2);
- 38       • for output of an IEEE NaN, whether after the letters 'NaN', the processor produces additional alphanumeric
- 39       characters enclosed in parentheses (13.7.2.3.2);
- 40       • the choice of binary exponent in EX output editing (13.7.2.3.6);
- 41       • the effect of the input/output rounding mode `PROCESSOR_DEFINED` (13.7.2.3.8);
- 42       • which value is chosen if the input/output rounding mode is `NEAREST` and the value to be converted is
- 43       exactly halfway between the two nearest representable values in the result format (13.7.2.3.8);

- 1 • the field width, decimal part width, and exponent width used for the G0 edit descriptor (13.7.5);
- 2 • the file position when position editing skips a character of nondefault kind in an [internal file](#) of default
- 3 character kind or an [external unit](#) that is not connected to a Unicode file (13.8.1.1);
- 4 • when the sign mode is PROCESSOR\_DEFINED, whether a plus sign appears in a numeric output field
- 5 for a nonnegative value (13.8.4);
- 6 • the results of list-directed output (13.10.4);
- 7 • the results of namelist output (13.11.4);
- 8 • the interaction between [argument association](#) and pointer association, (15.5.2.4);
- 9 • the values returned by some intrinsic functions (16);
- 10 • how the sequences of atomic actions in unordered segments interleave (16.5);
- 11 • the value assigned to a STAT argument in a reference to an atomic subroutine when an error condition
- 12 occurs (16.5);
- 13 • the effect of calling EXECUTE\_COMMAND\_LINE on any [image](#) other than [image 1](#) in the [initial team](#)
- 14 (16.7);
- 15 • whether each [image](#) uses a separate random number generator, or if some or all [images](#) use common random
- 16 number generators (16.7);
- 17 • whether the results returned from CPU\_TIME, DATE\_AND\_TIME and SYSTEM\_CLOCK are depend-
- 18 ent on which [image](#) calls them (16.7);
- 19 • the set of error conditions that can occur in some intrinsic subroutines (16.9);
- 20 • the value assigned to a CMDSTAT, ERRMSG, EXITSTAT, STAT, or STATUS argument to indicate a
- 21 processor-dependent error condition (16.9);
- 22 • the computed value of the the intrinsic subroutine CO\_REDUCE (16.9.49) and the intrinsic subroutine
- 23 CO\_SUM (16.9.50);
- 24 • the value assigned to the TIME argument by the intrinsic subroutine CPU\_TIME (16.9.57);
- 25 • whether date, clock, and time zone information is available (16.9.59);
- 26 • whether date, clock, and time zone information on one [image](#) is the same as that on another [image](#) (16.9.59);
- 27 • the value of command argument zero, if the processor does not support the concept of a command name
- 28 (16.9.83);
- 29 • the order of command arguments (16.9.83);
- 30 • whether the significant length of a command argument includes trailing blanks (16.9.83);
- 31 • the interpretation of case for the NAME argument of the intrinsic subroutine GET\_ENVIRONMENT\_-
- 32 VARIABLE (16.9.84);
- 33 • whether an environment variable that exists on an [image](#) also exists on another [image](#), and if it does exist
- 34 on both [images](#), whether the values are the same or different (16.9.84);
- 35 • the value assigned to the pseudorandom number generator by the intrinsic subroutine RANDOM\_INIT
- 36 (16.9.155);
- 37 • the computation of the seed value used by the pseudorandom number generator (16.9.157);
- 38 • on [images](#) that use a common random number generator, the interleaving of values assigned by RANDOM\_-
- 39 NUMBER in unordered segments(16.7);
- 40 • the value assigned to the seed by the intrinsic subroutine RANDOM\_SEED when no argument is present
- 41 (16.9.157);
- 42 • the values assigned to its arguments by the intrinsic subroutine SYSTEM\_CLOCK (16.9.186);

- 1 • the values of the named constants in the intrinsic module `ISO_FORTRAN_ENV` (16.10.2);
- 2 • the values returned by the functions `COMPILER_OPTIONS` and `COMPILER_VERSION` in the intrinsic
- 3 module `ISO_FORTRAN_ENV` (16.10.2);
- 4 • the extent to which a processor supports IEEE arithmetic (17);
- 5 • whether a flag that is quiet on entry to a scoping unit that does not access `IEEE_FEATURES`, `IEEE_`
- 6 `EXCEPTIONS`, or `IEEE_ARITHMETIC` is signaling on exit (17.1);
- 7 • the conditions under which `IEEE_OVERFLOW` is raised in a calculation involving non-ISO/IEC/IEEE
- 8 60559:2011 floating-point data (17.3);
- 9 • the conditions under which `IEEE_OVERFLOW` and `IEEE_DIVIDE_BY_ZERO` are raised in a floating-
- 10 point exponentiation operation (17.3);
- 11 • the conditions under which `IEEE_DIVIDE_BY_ZERO` is raised in a calculation involving non-ISO/IEC/IEEE
- 12 60559:2011 floating-point data (17.3);
- 13 • whether an exception signals at the end of a sequence of statements that has no invocations of `IEEE_GET_`
- 14 `FLAG`, `IEEE_SET_FLAG`, `IEEE_GET_STATUS`, `IEEE_SET_STATUS`, or `IEEE_SET_HALTING_`
- 15 `MODE`, in which execution of an operation would cause it to signal, if no value of a variable depends upon
- 16 the result of the operation (17.3);
- 17 • the initial rounding modes (17.4);
- 18 • whether the processor supports a particular rounding mode (17.4);
- 19 • the effect of the rounding mode `IEEE_OTHER`, if supported (17.4);
- 20 • the initial underflow mode (17.5);
- 21 • the initial halting mode (17.6);
- 22 • whether `IEEE_INT` implements the `convertToInteger{round}` or `convertToIntegerExact{round}` operation
- 23 specified by ISO/IEC/IEEE 60559:2011 (17.11.11);
- 24 • which argument is the result value of `IEEE_MAX_NUM`, `IEEE_MAX_NUM_MAG`, `IEEE_MIN_NUM`,
- 25 or `IEEE_MIN_NUM_MAG` when both arguments are quiet NaNs or are zeros (17.11.17, 17.11.18, 17.11.19,
- 26 17.11.20);
- 27 • the requirements on the storage sequence to be associated with the pointer `FPTR` by the `C_F_POINTER`
- 28 subroutine (18.2.3.4);
- 29 • the order of the members of the `CFI_dim_t` structure defined in the source file `CFI_Fortran_binding.h`
- 30 (18.5.2);
- 31 • members of the `CFI_cdesc_t` structure defined in the source file `CFI_Fortran_binding.h` beyond the re-
- 32 quirements of 18.5.3;
- 33 • the value of `CFI_MAX_RANK` in the source file `CFI_Fortran_binding.h` (18.5.4);
- 34 • the value of `CFI_VERSION` in the source file `CFI_Fortran_binding.h` (18.5.4);
- 35 • which error condition is detected if more than one error condition could be detected for an invocation of
- 36 one of the functions declared in the source file `CFI_Fortran_binding.h` (18.5.5.1);
- 37 • the values of the attribute specifier macros defined in the source file `CFI_Fortran_binding.h` (18.5.4);
- 38 • the values of the type specifier macros defined in the source file `CFI_Fortran_binding.h`;
- 39 • which additional type specifier values are defined in the source file `CFI_Fortran_binding.h` (18.5.4);
- 40 • the values of the error code macros other than `CFI_SUCCESS` that are defined in the source file `CFI_`
- 41 `Fortran_binding.h` (18.5.4);
- 42 • the base address of a zero-sized array (18.5.3);
- 43 • the values of the floating-point exception flags on entry to a procedure defined by means other than Fortran



- 1           (18.10.3);
- 2           • whether a procedure defined by means other than Fortran is an asynchronous communication initiation or
- 3           completion procedure (18.10.4).

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## Annex B

(Informative)

### Deleted and obsolescent features

#### B.1 Deleted features from Fortran 90

1 These deleted features are those features of Fortran 90 that were redundant and considered largely unused.

2 The following Fortran 90 features are not required.

3 (1) Real and double precision DO variables.

4 In FORTRAN 77 and Fortran 90, a DO variable was allowed to be of type real or double precision  
5 in addition to type integer; this has been deleted. A similar result can be achieved by using a **DO**  
6 **construct** with no loop control and the appropriate exit test.

7 (2) Branching to an **END IF statement** from outside its block.

8 In FORTRAN 77 and Fortran 90, it was possible to branch to an **END IF statement** from outside the  
9 **IF construct**; this has been deleted. A similar result can be achieved by branching to a **CONTINUE**  
10 **statement** that is immediately after the **END IF statement**.

11 (3) PAUSE statement.

12 The PAUSE statement, provided in FORTRAN 66, FORTRAN 77, and Fortran 90, has been deleted.  
13 A similar result can be achieved by writing a message to the appropriate **unit**, followed by reading  
14 from the appropriate **unit**.

15 (4) ASSIGN and assigned GO TO statements, and assigned format specifiers.

16 The ASSIGN statement and the related assigned GO TO statement, provided in FORTRAN 66,  
17 FORTRAN 77, and Fortran 90, have been deleted. Further, the ability to use an assigned integer as a  
18 format, provided in FORTRAN 77 and Fortran 90, has been deleted. A similar result can be achieved  
19 by using other control constructs instead of the assigned GO TO statement and by using a default  
20 character variable to hold a format specification instead of using an assigned integer.

21 (5) H edit descriptor.

22 In FORTRAN 77 and Fortran 90, there was an alternative form of character string edit descriptor,  
23 which had been the only such form in FORTRAN 66; this has been deleted. A similar result can be  
24 achieved by using a character string edit descriptor.

25 (6) Vertical format control.

26 In FORTRAN 66, FORTRAN 77, Fortran 90, and Fortran 95 formatted output to certain **units** resulted  
27 in the first character of each record being interpreted as controlling vertical spacing. There was no  
28 standard way to detect whether output to a **unit** resulted in this vertical format control, and no  
29 way to specify that it should be applied; this has been deleted. The effect can be achieved by  
30 post-processing a formatted file.

31 3 See ISO/IEC 1539:1991 for detailed rules of how these deleted features worked.

## B.2 Deleted features from Fortran 2008

1 These deleted features are those features of Fortran 2008 that were redundant and considered largely unused.

2 The following Fortran 2008 features are not required.

(1) Arithmetic IF statement.

The arithmetic IF statement is incompatible with ISO/IEC/IEEE 60559:2011 and necessarily involves the use of [statement labels](#); [statement labels](#) can hinder optimization, and make code hard to read and maintain. Similar logic can be more clearly encoded using other conditional statements.

(2) Nonblock DO construct

The nonblock forms of the DO loop were confusing and hard to maintain. Shared termination and dual use of labeled action statements as do termination and branch targets were especially error-prone.

## B.3 Obsolescent features

### B.3.1 General

1 The obsolescent features are those features of Fortran 90 that were redundant and for which better methods were available in Fortran 90. Subclause [4.4.3](#) describes the nature of the obsolescent features. The obsolescent features in this document are the following.

(1) Alternate return — see [B.3.2](#).

(2) [Computed GO TO](#) — see [B.3.3](#).

(3) Statement functions — see [B.3.4](#).

(4) [DATA statements](#) amongst executable statements — see [B.3.5](#).

(5) Assumed length character functions — see [B.3.6](#).

(6) Fixed form source — see [B.3.7](#).

(7) CHARACTER\* form of CHARACTER declaration — see [B.3.8](#).

(8) [ENTRY statements](#) — see [B.3.9](#).

(9) Label form of [DO statement](#) — see [B.3.10](#).

(10) [COMMON](#) and [EQUIVALENCE](#) statements, and the [block data program unit](#) — see [B.3.11](#).

(11) Specific names for intrinsic functions — see [B.3.12](#).

(12) [FORALL](#) construct and statement — see [B.3.13](#)

### B.3.2 Alternate return

1 An alternate return introduces labels into an argument list to allow the called procedure to direct the execution of the caller upon return. The same effect can be achieved with a return code that is used in a [SELECT CASE construct](#) on return. This avoids an irregularity in the syntax and semantics of [argument association](#). For example,

```
CALL SUBR_NAME (X, Y, Z, *100, *200, *300)
```

can be replaced by

```
CALL SUBR_NAME (X, Y, Z, RETURN_CODE)
SELECT CASE (RETURN_CODE)
```

```
1      CASE (1)
2          ...
3      CASE (2)
4          ...
5      CASE (3)
6          ...
7      CASE DEFAULT
8          ...
9  END SELECT
```

### 10 **B.3.3 Computed GO TO statement**

11 1 The [computed GO TO statement](#) has been superseded by the [SELECT CASE construct](#), which is a generalized,  
12 easier to use, and clearer means of expressing the same computation.

### 13 **B.3.4 Statement functions**

14 1 [Statement functions](#) are subject to a number of nonintuitive restrictions and are a potential source of error because  
15 their syntax is easily confused with that of an assignment statement.

16 2 The internal function is a more generalized form of the [statement function](#) and completely supersedes it.

### 17 **B.3.5 DATA statements among executables**

18 1 The statement ordering rules allow [DATA statements](#) to appear anywhere in a [program unit](#) after the specific-  
19 ation statements. The ability to position [DATA statements](#) amongst executable statements is very rarely used,  
20 unnecessary, and a potential source of error.

### 21 **B.3.6 Assumed character length functions**

22 1 Assumed character length for functions is an irregularity in the language in that elsewhere in Fortran the philo-  
23 sophy is that the attributes of a function result depend only on the [actual arguments](#) of the invocation and on  
24 any data accessible by the function through host or use association. Some uses of this facility can be replaced  
25 with an [automatic](#) character length function, where the length of the function result is declared in a [specification](#)  
26 [expression](#). Other uses can be replaced by the use of a subroutine whose arguments correspond to the function  
27 result and the function arguments.

28 2 Note that dummy arguments of a function can have assumed character length.

### 29 **B.3.7 Fixed form source**

30 1 Fixed form source was designed when the principal machine-readable input medium for new programs was punched  
31 cards. Now that new and amended programs are generally entered via keyboards with screen displays, it is an  
32 unnecessary overhead, and is potentially error-prone, to have to locate positions 6, 7, or 72 on a line. Free form  
33 source was designed expressly for this more modern technology.

34 2 It is a simple matter for a software tool to convert from fixed to free form source.

### 1 **B.3.8 CHARACTER\* form of CHARACTER declaration**

2 1 In addition to the CHARACTER\**char-length* form introduced in FORTRAN 77, Fortran 90 provided the CHAR-  
3 ACTER([ LEN = ] *type-param-value*) form. The older form (CHARACTER\**char-length*) is redundant.

### 4 **B.3.9 ENTRY statements**

5 1 **ENTRY statements** allow more than one entry point to a subprogram, facilitating sharing of data items and  
6 executable statements local to that subprogram.

7 2 This can be replaced by a module containing the (private) data items, with a module procedure for each entry  
8 point and the shared code in a private module procedure.

### 9 **B.3.10 Label DO statement**

10 1 The label in the **DO statement** is redundant with the construct name. Furthermore, the label allows unrestricted  
11 branches and, for its main purpose (the target of a conditional branch to skip the rest of the current iteration),  
12 is redundant with the **CYCLE statement**, which is clearer.

### 13 **B.3.11 COMMON and EQUIVALENCE statements and the block data program unit**

14 1 **Common blocks** are error-prone and have largely been superseded by **modules**. **EQUIVALENCE** similarly is  
15 error-prone. Whilst use of these statements was invaluable prior to Fortran 90 they are now redundant and  
16 can inhibit performance. The **block data program unit** exists only to serve **common blocks** and hence is also  
17 redundant.

### 18 **B.3.12 Specific names for intrinsic functions**

19 1 The **specific names of the intrinsic functions** are often obscure and hinder portability. They have been redundant  
20 since Fortran 90. Use generic names for references to intrinsic procedures.

### 21 **B.3.13 FORALL construct and statement**

22 1 The **FORALL** construct and statement were added to the language in the expectation that they would enable  
23 highly efficient execution, especially on parallel processors. However, experience indicates that they are too  
24 complex and have too many restrictions for compilers to take advantage of them. They are redundant with the  
25 **DO CONCURRENT construct**, and many of the manipulations for which they might be used can be done more  
26 effectively using pointers, especially using pointer rank remapping.

# Annex C

(Informative)

## Extended notes

### C.1 Fortran 2008 features not mentioned in its Introduction

1 The following features were new in Fortran 2008 but not originally listed in its Introduction as being new features:

- 2 • An array or object with a nonconstant [length type parameter](#) can have the [VALUE attribute](#).
- 3 • Multiple allocations are permitted in a single [ALLOCATE statement](#) with the [SOURCE= specifier](#).
- 4 • A [PROCEDURE statement](#) can have a double colon before the first procedure name.
- 5 • An argument to a [pure procedure](#) can have default [INTENT](#) if it has the [VALUE attribute](#).
- 6 • The [PROTECTED attribute](#) can be specified by the [procedure declaration statement](#).
- 7 • A [defined-operator](#) can be used in a [specification expression](#).
- 8 • All transformational functions from the intrinsic modules [IEEE\\_ARITHMETIC](#) and [IEEE\\_EXCEPTIONS](#) can be used in [constant expressions](#). All transformational functions from the intrinsic modules [IEEE\\_ARITHMETIC](#), [IEEE\\_EXCEPTIONS](#), and [ISO\\_C\\_BINDING](#) can be used in [specification expressions](#).
- 9 • A [contiguous](#) array variable that is not interoperable but which has interoperable type and [kind type parameter](#) (if any), and a scalar character variable with length greater than 1 and [kind C\\_CHAR](#) in the intrinsic module [ISO\\_C\\_BINDING](#), can be used as the argument of the function [C\\_LOC](#) in the intrinsic module [ISO\\_C\\_BINDING](#), provided the variable has the [POINTER](#) or [TARGET](#) attribute.
- 10 • The name of an [external procedure](#) that has a [binding label](#) is a local identifier and not a global identifier.
- 11 • A procedure that is not a [procedure pointer](#) can be an [actual argument](#) that corresponds to a [procedure pointer dummy argument](#) with the [INTENT \(IN\)](#) attribute.
- 12 • An interface body for an [external procedure](#) that does not exist in a program can be used to specify an [explicit specific interface](#).

2 All but the last two of the above list were subsequently added to the Introduction by Technical Corrigenda.

### C.2 Clause 7 notes

#### C.2.1 Selection of the approximation methods (7.4.3.2)

1 One can select the real approximation method for an entire program through the use of a module and the parameterized real type. This is accomplished by defining a named integer constant to have a particular kind type parameter value and using that [named constant](#) in all real, complex, and derived-type declarations. For example, the specification statements

```

21     INTEGER, PARAMETER :: LONG_FLOAT = 8
22     REAL (LONG_FLOAT) X, Y
23     COMPLEX (LONG_FLOAT) Z

```

1 specify that the approximation method corresponding to a kind type parameter value of 8 is supplied for the data  
 2 objects X, Y, and Z in the [program unit](#). The kind type parameter value LONG\_FLOAT can be made available  
 3 to an entire program by placing the INTEGER specification statement in a module and accessing the [named](#)  
 4 [constant](#) LONG\_FLOAT with a [USE statement](#). Note that by changing 8 to 4 once in the module, a different  
 5 approximation method is selected.

6 2 To avoid the use of the processor-dependent values 4 or 8, replace 8 by [KIND](#) (0.0) or [KIND](#) (0.0D0). Another  
 7 way to avoid these processor-dependent values is to select the kind value using the intrinsic function [SELEC-](#)  
 8 [TED\\_REAL\\_KIND](#) (16.9.170). In the above specification statement, the 8 might be replaced by, for instance,  
 9 [SELECTED\\_REAL\\_KIND](#) (10, 50), which requires an approximation method to be selected with at least 10  
 10 decimal digits of precision and a range from  $10^{-50}$  to  $10^{50}$ . There are no magnitude or ordering constraints placed  
 11 on kind values, in order that implementers have flexibility in assigning such values and can add new kinds without  
 12 changing previously assigned kind values.

13 3 As kind values have no portable meaning, a good practice is to use them in programs only through [named](#)  
 14 [constants](#) as described above (for example, SINGLE, IEEE\_SINGLE, DOUBLE, and QUAD), rather than using  
 15 the kind values directly.

## 16 C.2.2 Type extension and component accessibility (7.5.2.2, 7.5.4)

17 1 The default accessibility of the components of an [extended type](#) can be specified in the type definition. The  
 18 accessibility of its components can be specified individually. For example:

```

19     module types
20         type base_type
21             private                !-- Sets default accessibility
22             integer :: i            !-- a private component
23             integer, private :: j  !-- another private component
24             integer, public :: k   !-- a public component
25         end type base_type
26
27         type, extends(base_type) :: my_type
28             private                !-- Sets default for components declared in my_type
29             integer :: l            !-- A private component.
30             integer, public :: m   !-- A public component.
31         end type my_type
32
33     end module types
34
35     subroutine sub
36         use types
37         type (my_type) :: x
38
39         ...
40
41     call another_sub( &
42         x%base_type, & !-- ok because base_type is a public subobject of x
43         x%base_type%k, & !-- ok because x%base_type is ok and has k as a

```



```

1           !-- public component.
2       x%k,           & !-- ok because it is shorthand for x%base_type%k
3       x%base_type%i, & !-- Invalid because i is private.
4       x%i)           !-- Invalid because it is shorthand for x%base_type%i
5   end subroutine sub

```

### 6 C.2.3 Generic type-bound procedures (7.5.5)

7 Example of a derived type with generic type-bound procedures:

8 1 The only difference between this example and the same thing rewritten to use generic [interface blocks](#) is that  
9 with [type-bound procedures](#),

```
10     USE rational_numbers, ONLY: rational
```

11 does not block the [type-bound procedures](#); the user still gets access to the defined assignment and extended  
12 operations.

```

13     MODULE rational_numbers
14     IMPLICIT NONE
15     PRIVATE
16     TYPE,PUBLIC :: rational
17     PRIVATE
18     INTEGER n,d
19     CONTAINS
20     ! ordinary type-bound procedure
21     PROCEDURE :: real => rat_to_real
22     ! specific type-bound procedures for generic support
23     PROCEDURE,PRIVATE :: rat_asgn_i, rat_plus_i, rat_plus_rat => rat_plus
24     PROCEDURE,PRIVATE,PASS(b) :: i_plus_rat
25     ! generic type-bound procedures
26     GENERIC :: ASSIGNMENT(=) => rat_asgn_i
27     GENERIC :: OPERATOR(+) => rat_plus_rat, rat_plus_i, i_plus_rat
28     END TYPE
29     CONTAINS
30     ELEMENTAL REAL FUNCTION rat_to_real(this) RESULT(r)
31     CLASS(rational),INTENT(IN) :: this
32     r = REAL(this%n)/this%d
33     END FUNCTION
34     ELEMENTAL SUBROUTINE rat_asgn_i(a,b)
35     CLASS(rational),INTENT(OUT) :: a
36     INTEGER,INTENT(IN) :: b
37     a%n = b
38     a%d = 1
39     END SUBROUTINE
40     ELEMENTAL TYPE(rational) FUNCTION rat_plus_i(a,b) RESULT(r)
41     CLASS(rational),INTENT(IN) :: a
42     INTEGER,INTENT(IN) :: b

```

```

1         r%n = a%n + b*a%d
2         r%d = a%d
3     END FUNCTION
4     ELEMENTAL TYPE(rational) FUNCTION i_plus_rat(a,b) RESULT(r)
5         INTEGER,INTENT(IN) :: a
6         CLASS(rational),INTENT(IN) :: b
7         r%n = b%n + a*b%d
8         r%d = b%d
9     END FUNCTION
10    ELEMENTAL TYPE(rational) FUNCTION rat_plus(a,b) RESULT(r)
11        CLASS(rational),INTENT(IN) :: a,b
12        r%n = a%n*b%d + b%n*a%d
13        r%d = a%d*b%d
14    END FUNCTION
15    END

```

## 16 C.2.4 Abstract types (7.5.7.1)

17 1 The following illustrates how an abstract type can be used as the basis for a collection of related types, and how  
18 a non-abstract member of that collection can be created by type extension.

```

19    TYPE, ABSTRACT :: DRAWABLE_OBJECT
20        REAL, DIMENSION(3) :: RGB_COLOR = (/1.0,1.0,1.0/) ! White
21        REAL, DIMENSION(2) :: POSITION = (/0.0,0.0/) ! Centroid
22    CONTAINS
23        PROCEDURE(RENDER_X), PASS(OBJECT), DEFERRED :: RENDER
24    END TYPE DRAWABLE_OBJECT
25
26    ABSTRACT INTERFACE
27        SUBROUTINE RENDER_X(OBJECT, WINDOW)
28            IMPORT DRAWABLE_OBJECT, X_WINDOW
29            CLASS(DRAWABLE_OBJECT), INTENT(IN) :: OBJECT
30            CLASS(X_WINDOW), INTENT(INOUT) :: WINDOW
31        END SUBROUTINE RENDER_X
32    END INTERFACE
33
34    ...
35
36    TYPE, EXTENDS(DRAWABLE_OBJECT) :: DRAWABLE_TRIANGLE ! Not ABSTRACT
37        REAL, DIMENSION(2,3) :: VERTICES ! In relation to centroid
38    CONTAINS
39        PROCEDURE, PASS(OBJECT) :: RENDER=>RENDER_TRIANGLE_X
40    END TYPE DRAWABLE_TRIANGLE

```

41 2 The actual drawing procedure will draw a triangle in WINDOW with vertices at x and y coordinates at  
42 OBJECT%POSITION(1)+OBJECT%VERTICES(1,1:3) and OBJECT%POSITION(2)+OBJECT%VERTICES(2,1:3):

```

1      SUBROUTINE RENDER_TRIANGLE_X(OBJECT, WINDOW)
2          CLASS(DRAWABLE_TRIANGLE), INTENT(IN) :: OBJECT
3          CLASS(X_WINDOW), INTENT(INOUT) :: WINDOW
4          ...
5      END SUBROUTINE RENDER_TRIANGLE_X

```

## 6 C.2.5 Structure constructors and generic names (7.5.10)

7 1 A generic name can be the same as a type name. This can be used to emulate user-defined [structure constructors](#)  
8 for that type, even if the type has private components. For example:

```

9      MODULE mytype_module
10         TYPE mytype
11             PRIVATE
12             COMPLEX value
13             LOGICAL exact
14         END TYPE
15         INTERFACE mytype
16             MODULE PROCEDURE int_to_mytype
17         END INTERFACE
18         ! Operator definitions etc.
19         ...
20     CONTAINS
21         TYPE(mytype) FUNCTION int_to_mytype(i)
22             INTEGER, INTENT(IN) :: i
23             int_to_mytype%value = i
24             int_to_mytype%exact = .TRUE.
25         END FUNCTION
26         ! Procedures to support operators etc.
27         ...
28     END
29
30     PROGRAM example
31         USE mytype_module
32         TYPE(mytype) x
33         x = mytype(17)
34     END

```

35 2 The type name can still be used as a generic name if the type has type parameters. For example:

```

36     MODULE m
37         TYPE t(kind)
38             INTEGER, KIND :: kind
39             COMPLEX(kind) value
40         END TYPE
41         INTEGER, PARAMETER :: single = KIND(0.0), double = KIND(0d0)
42         INTERFACE t

```

```

1      MODULE PROCEDURE real_to_t1, dble_to_t2, int_to_t1, int_to_t2
2      END INTERFACE
3      ...
4      CONTAINS
5      TYPE(t(single)) FUNCTION real_to_t1(x)
6          REAL(single) x
7          real_to_t1%value = x
8      END FUNCTION
9      TYPE(t(double)) FUNCTION dble_to_t2(x)
10         REAL(double) x
11         dble_to_t2%value = x
12     END FUNCTION
13     TYPE(t(single)) FUNCTION int_to_t1(x,mold)
14         INTEGER x
15         TYPE(t(single)) mold
16         int_to_t1%value = x
17     END FUNCTION
18     TYPE(t(double)) FUNCTION int_to_t2(x,mold)
19         INTEGER x
20         TYPE(t(double)) mold
21         int_to_t2%value = x
22     END FUNCTION
23     ...
24     END
25
26     PROGRAM example
27         USE m
28         TYPE(t(single)) x
29         TYPE(t(double)) y
30         x = t(1.5)                ! References real_to_t1
31         x = t(17,mold=x)          ! References int_to_t1
32         y = t(1.5d0)              ! References dble_to_t2
33         y = t(42,mold=y)         ! References int_to_t2
34         y = t(kind(0d0)) ((0,1)) ! Uses the structure constructor for type t
35     END

```

## 36 C.2.6 Final subroutines (7.5.6, 7.5.6.2, 7.5.6.3, 7.5.6.4)

37 Example of a parameterized derived type with final subroutines:

```

38     MODULE m
39         TYPE t(k)
40             INTEGER, KIND :: k
41             REAL(k), POINTER :: vector(:) => NULL()
42     CONTAINS
43         FINAL :: finalize_t1s, finalize_t1v, finalize_t2e
44     END TYPE

```

```

1      CONTAINS
2      SUBROUTINE finalize_t1s(x)
3          TYPE(t(KIND(0.0))) x
4          IF (ASSOCIATED(x%vector)) DEALLOCATE(x%vector)
5      END SUBROUTINE
6      SUBROUTINE finalize_t1v(x)
7          TYPE(t(KIND(0.0))) x(:)
8          DO i=LBOUND(x,1),UBOUND(x,1)
9              IF (ASSOCIATED(x(i)%vector)) DEALLOCATE(x(i)%vector)
10         END DO
11     END SUBROUTINE
12     ELEMENTAL SUBROUTINE finalize_t2e(x)
13         TYPE(t(KIND(0.0d0))),INTENT(INOUT) :: x
14         IF (ASSOCIATED(x%vector)) DEALLOCATE(x%vector)
15     END SUBROUTINE
16 END MODULE
17
18 SUBROUTINE example(n)
19     USE m
20     TYPE(t(KIND(0.0))) a,b(10),c(n,2)
21     TYPE(t(KIND(0.0d0))) d(n,n)
22     ...
23     ! Returning from this subroutine will effectively do
24     !   CALL finalize_t1s(a)
25     !   CALL finalize_t1v(b)
26     !   CALL finalize_t2e(d)
27     ! No final subroutine will be called for variable C because the user
28     ! omitted to define a suitable specific procedure for it.
29 END SUBROUTINE

```

30 **Example of extended types with final subroutines:**

```

31     MODULE m
32         TYPE t1
33             REAL a,b
34         END TYPE
35         TYPE,EXTENDS(t1) :: t2
36             REAL,POINTER :: c(:),d(:)
37         CONTAINS
38             FINAL :: t2f
39         END TYPE
40         TYPE,EXTENDS(t2) :: t3
41             REAL,POINTER :: e
42         CONTAINS
43             FINAL :: t3f
44         END TYPE
45         ...

```

```

1      CONTAINS
2          SUBROUTINE t2f(x) ! Finalizer for TYPE(t2)'s extra components
3              TYPE(t2) :: x
4              IF (ASSOCIATED(x%c)) DEALLOCATE(x%c)
5              IF (ASSOCIATED(x%d)) DEALLOCATE(x%d)
6          END SUBROUTINE
7          SUBROUTINE t3f(y) ! Finalizer for TYPE(t3)'s extra components
8              TYPE(t3) :: y
9              IF (ASSOCIATED(y%e)) DEALLOCATE(y%e)
10         END SUBROUTINE
11     END MODULE
12
13     SUBROUTINE example
14         USE m
15         TYPE(t1) x1
16         TYPE(t2) x2
17         TYPE(t3) x3
18         ...
19         ! Returning from this subroutine will effectively do
20         !     ! Nothing to x1; it is not finalizable
21         !     CALL t2f(x2)
22         !     CALL t3f(x3)
23         !     CALL t2f(x3%t2)
24     END SUBROUTINE

```

### 25 C.3 Clause 8 notes: The VOLATILE attribute (8.5.19)

- 26 1 The following example shows the use of a variable with the [VOLATILE attribute](#) to communicate with an  
27 asynchronous process, in this case the operating system. The program detects a user keystroke on the terminal  
28 and reacts at a convenient point in its processing.
- 29 2 The [VOLATILE attribute](#) is necessary to prevent an optimizing compiler from storing the communication variable  
30 in a register or from doing flow analysis and deciding that the [EXIT statement](#) can never be executed.

```

31     SUBROUTINE TERMINATE_ITERATIONS
32         LOGICAL, VOLATILE :: USER_HIT_ANY_KEY
33
34         ! Have the OS start to look for a user keystroke and set the variable
35         ! "USER_HIT_ANY_KEY" to TRUE as soon as it detects a keystroke.
36         ! This call is operating system dependent.
37
38         CALL OS_BEGIN_DETECT_USER_KEYSTROKE( USER_HIT_ANY_KEY )
39         USER_HIT_ANY_KEY = .FALSE.      ! This will ignore any recent keystrokes.
40         PRINT *, " Hit any key to terminate iterations!"
41
42         DO I = 1,100
43             ...                          ! Compute a value for R.

```

```

1         PRINT *, I, R
2         IF (USER_HIT_ANY_KEY)      EXIT
3     ENDDO
4
5         ! Have the OS stop looking for user keystrokes.
6         CALL OS_STOP_DETECT_USER_KEYSTROKE
7
8     END SUBROUTINE TERMINATE_ITERATIONS

```

## 9 C.4 Clause 9 notes

### 10 C.4.1 Structure components (9.4.2)

- 11 1 Components of a structure are referenced by writing the components of successive levels of the structure hierarchy  
12 until the desired component is described. For example,

```

13     TYPE ID_NUMBERS
14         INTEGER SSN
15         INTEGER EMPLOYEE_NUMBER
16     END TYPE ID_NUMBERS
17
18     TYPE PERSON_ID
19         CHARACTER (LEN=30) LAST_NAME
20         CHARACTER (LEN=1) MIDDLE_INITIAL
21         CHARACTER (LEN=30) FIRST_NAME
22         TYPE (ID_NUMBERS) NUMBER
23     END TYPE PERSON_ID
24
25     TYPE PERSON
26         INTEGER AGE
27         TYPE (PERSON_ID) ID
28     END TYPE PERSON
29
30     TYPE (PERSON) GEORGE, MARY
31
32     PRINT *, GEORGE % AGE           ! Print the AGE component
33     PRINT *, MARY % ID % LAST_NAME ! Print LAST_NAME of MARY
34     PRINT *, MARY % ID % NUMBER % SSN ! Print SSN of MARY
35     PRINT *, GEORGE % ID % NUMBER ! Print SSN and EMPLOYEE_NUMBER of GEORGE

```

- 36 2 A [structure component](#) can be a data object of intrinsic type as in the case of GEORGE % AGE or it can be  
37 of derived type as in the case of GEORGE % ID % NUMBER. The resultant component can be a scalar or an  
38 array of intrinsic or derived type.

```

39     TYPE LARGE
40         INTEGER ELT (10)
41         INTEGER VAL

```

```

1      END TYPE LARGE
2
3      TYPE (LARGE) A (5)      ! 5 element array, each of whose elements
4                             ! includes a 10 element array ELT and
5                             ! a scalar VAL.
6      PRINT *, A (1)        ! Prints 10 element array ELT and scalar VAL.
7      PRINT *, A (1) % ELT (3) ! Prints scalar element 3
8                             ! of array element 1 of A.
9      PRINT *, A (2:4) % VAL ! Prints scalar VAL for array elements
10                             ! 2 to 4 of A.

```

3 Components of an object of [extensible type](#) that are [inherited](#) from the [parent type](#) can be accessed as a whole by using the [parent component](#) name, or individually, either with or without qualifying them by the [parent component](#) name. For example:

```

14     TYPE POINT              ! A base type
15     REAL :: X, Y
16     END TYPE POINT
17     TYPE, EXTENDS(POINT) :: COLOR_POINT ! An extension of TYPE(POINT)
18     ! Components X and Y, and component name POINT, inherited from parent
19     INTEGER :: COLOR
20     END TYPE COLOR_POINT
21
22     TYPE(POINT) :: PV = POINT(1.0, 2.0)
23     TYPE(COLOR_POINT) :: CPV = COLOR_POINT(POINT=PV, COLOR=3)
24
25     PRINT *, CPV%POINT      ! Prints 1.0 and 2.0
26     PRINT *, CPV%POINT%X, CPV%POINT%Y ! And this does, too
27     PRINT *, CPV%X, CPV%Y   ! And this does, too

```

## 28 C.4.2 Allocation with dynamic type (9.7.1)

29 1 The following example illustrates the use of allocation with the value and [dynamic type](#) of the allocated object  
30 given by another object. The example copies a list of objects of any type. It copies the list starting at IN\_LIST.  
31 After copying, each element of the list starting at LIST\_COPY has a polymorphic component, ITEM, for which  
32 both the value and type are taken from the ITEM component of the corresponding element of the list starting at  
33 IN\_LIST.

```

34     TYPE :: LIST ! A list of anything
35     TYPE(LIST), POINTER :: NEXT => NULL()
36     CLASS(*), ALLOCATABLE :: ITEM
37     END TYPE LIST
38     ...
39     TYPE(LIST), POINTER :: IN_LIST, LIST_COPY => NULL()
40     TYPE(LIST), POINTER :: IN_WALK, NEW_TAIL
41     ! Copy IN_LIST to LIST_COPY
42     IF (ASSOCIATED(IN_LIST)) THEN

```



```

1      IN_WALK => IN_LIST
2      ALLOCATE(LIST_COPY)
3      NEW_TAIL => LIST_COPY
4      DO
5          ALLOCATE(NEW_TAIL%ITEM, SOURCE=IN_WALK%ITEM)
6          IN_WALK => IN_WALK%NEXT
7          IF (.NOT. ASSOCIATED(IN_WALK)) EXIT
8          ALLOCATE(NEW_TAIL%NEXT)
9          NEW_TAIL => NEW_TAIL%NEXT
10     END DO
11 END IF

```

## 12 C.5 Clause 10 notes

### 13 C.5.1 Evaluation of function references (10.1.7)

14 1 If more than one function reference appears in a statement, they can be executed in any order (subject to a  
15 function result being evaluated after the evaluation of its arguments) and their values cannot depend on the order  
16 of execution. This lack of dependence on order of evaluation enables parallel execution of the function references.

### 17 C.5.2 Pointers in expressions (10.1.9.2)

18 1 A pointer is considered to be like any other variable when it is used as a primary in an expression. If a pointer  
19 is used as an operand to an operator that expects a value, the pointer will automatically deliver the value stored  
20 in the space described by the pointer, that is, the value of the target object associated with the pointer.

### 21 C.5.3 Pointers in variable definition contexts (10.2.1.3, 19.6.7)

22 1 The appearance of a pointer in a context that requires its value is a reference to its target. Similarly, where a  
23 pointer appears in a variable definition context the variable that is defined is the target of the pointer.

24 2 Executing the program fragment

```

25     REAL, POINTER :: A
26     REAL, TARGET :: B = 10.0
27     A => B
28     A = 42.0
29     PRINT '(F4.1)', B

```

30 produces “42.0” as output.

## 31 C.6 Clause 11 notes

### 32 C.6.1 The SELECT CASE construct (11.1.9)

33 1 At most one case block is selected for execution within a [SELECT CASE construct](#), and there is no fall-through  
34 from one block into another block within a [SELECT CASE construct](#). Thus there is no requirement for the user  
35 to exit explicitly from a block.

## 1 C.6.2 Loop control (11.1.7)

2 1 Fortran provides several forms of loop control:

- 3 (1) With an iteration count and a DO variable. This is the classic Fortran DO loop.
- 4 (2) Test a logical condition before each execution of the loop (DO WHILE).
- 5 (3) DO “forever”.

## 6 C.6.3 Examples of DO constructs (11.1.7)

7 1 The following are all valid examples of DO constructs.

### 8 Example 1:

```

9         SUM = 0.0
10        READ (IUN) N
11        OUTER: DO L = 1, N           ! A DO with a construct name
12            READ (IUN) IQUAL, M, ARRAY (1:M)
13            IF (IQUAL < IQUAL_MIN) CYCLE OUTER ! Skip inner loop
14            INNER: DO 40 I = 1, M     ! A DO with a label and a name
15                CALL CALCULATE (ARRAY (I), RESULT)
16                IF (RESULT < 0.0) CYCLE
17                SUM = SUM + RESULT
18                IF (SUM > SUM_MAX) EXIT OUTER
19        40    END DO INNER
20        END DO OUTER

```

21 2 The outer loop has an iteration count of `MAX (N, 0)`, and will execute that number of times or until SUM exceeds SUM\_MAX, in which case the `EXIT OUTER` statement terminates both loops. The inner loop is skipped by the first `CYCLE statement` if the quality flag, IQUAL, is too low. If CALCULATE returns a negative RESULT, the second `CYCLE statement` prevents it from being summed. Both loops have construct names and the inner loop also has a label. A construct name is required in the `EXIT statement` in order to terminate both loops, but is optional in the `CYCLE statements` because each belongs to its innermost loop.

### 27 Example 2:

```

28        N = 0
29        DO 50, I = 1, 10
30            J = I
31            DO K = 1, 5
32                L = K
33                N = N + 1 ! This statement executes 50 times
34            END DO      ! Nonlabeled DO inside a labeled DO
35        50 CONTINUE

```

36 3 After execution of the above program fragment, I = 11, J = 10, K = 6, L = 5, and N = 50.

### 37 Example 3:

```

1         N = 0
2         DO I = 1, 10
3             J = I
4             DO 60, K = 5, 1 ! This inner loop is never executed
5                 L = K
6                 N = N + 1
7         60     CONTINUE      ! Labeled DO inside a nonlabeled DO
8         END DO

```

9 4 After execution of the above program fragment, I = 11, J = 10, K = 5, N = 0, and L is not defined by these  
10 statements.

### 11 C.6.4 Examples of invalid DO constructs (11.1.7)

12 1 The following are all examples of invalid skeleton DO constructs:

#### 13 Example 1:

```

14         DO I = 1, 10
15             ...
16         END DO LOOP ! No matching construct name

```

#### 17 Example 2:

```

18         LOOP: DO 1000 I = 1, 10 ! No matching construct name
19             ...
20         1000 CONTINUE

```

#### 21 Example 3:

```

22         LOOP1: DO
23             ...
24         END DO LOOP2 ! Construct names don't match

```

#### 25 Example 4:

```

26         DO I = 1, 10 ! Label required or ...
27             ...
28         1010 CONTINUE ! ... END DO required

```

#### 29 Example 5:

```

30         DO 1020 I = 1, 10
31             ...
32         1021 END DO ! Labels don't match

```

#### 33 Example 6:

```

1      FIRST: DO I = 1, 10
2          SECOND: DO J = 1, 5
3              ...
4          END DO FIRST    ! Improperly nested DOs
5      END DO SECOND

```

## 6 C.6.5 Simple example using events

7 1 A tree is a graph in which every node except one has a single “parent” node to which it is connected by an edge.  
8 The node without a parent is the “root” of the tree. The nodes that have a particular node as their parent are  
9 the “children” of that node. The root is at level 1, its children are at level 2, and so on.

10 2 A multifrontal code to solve a sparse set of linear equations involves a tree. Work at a node can start after all of  
11 its children’s work is complete and their data have been passed to it.

12 3 Here we assume that each node has been assigned to an image. Each image has a list of its nodes and these  
13 are ordered in decreasing tree level (all those at level  $L$  preceding those at level  $L - 1$ ). For each node, array  
14 elements hold the number of children, details about the parent, and an event variable. This allows the processing to  
15 proceed asynchronously subject to the rule that a parent has to wait for all its children.

### 16 Outline of example code:

```

17      PROGRAM TREE
18          USE, INTRINSIC :: ISO_FORTRAN_ENV
19          INTEGER, ALLOCATABLE :: NODE (:) ! Tree nodes that this image handles.
20          INTEGER, ALLOCATABLE :: NC (:)  ! NODE(I) has NC(I) children.
21          INTEGER, ALLOCATABLE :: PARENT (:), SUB (:)
22              ! The parent of NODE (I) is NODE (SUB (I)) [PARENT (I)].
23          TYPE (EVENT_TYPE), ALLOCATABLE :: DONE (:) [:]
24          INTEGER :: I, J, STATUS
25          ! Set up the tree, including allocation of all arrays.
26          DO I = 1, SIZE (NODE)
27              ! Wait for children to complete
28              IF (NC (I) > 0) THEN
29                  EVENT WAIT (DONE (I), UNTIL_COUNT=NC (I), STAT=STATUS)
30                  IF (STATUS/=0) EXIT
31              END IF
32
33              ! Process node, using data from children.
34              IF (PARENT (I)>0) THEN
35                  ! Node is not the root.
36                  ! Place result on image PARENT (I) for node NODE (SUB) [PARENT (I)]
37                  ! Tell PARENT (I) that this has been done.
38                  EVENT POST (DONE (SUB (I)) [PARENT (I)], STAT=STATUS)
39                  IF (STATUS/=0) EXIT
40              END IF
41          END DO
42      END PROGRAM TREE

```

## 1 C.6.6 Example using three teams

- 2 1 The following example illustrates the structure of a routine that will compute fluxes based on surface properties  
3 over land, sea, and ice, each in a different team. Each image will deal with areas containing exactly one of the  
4 three surface types.

```

5      SUBROUTINE COMPUTE_FLUXES (FLUX_MOM, FLUX_SENS, FLUX_LAT)
6      USE, INTRINSIC :: ISO_FORTRAN_ENV, ONLY: TEAM_TYPE
7      REAL, INTENT (OUT) :: FLUX_MOM (:,:), FLUX_SENS (:,:), FLUX_LAT (:,:)
8      INTEGER, PARAMETER :: LAND = 1, SEA = 2, ICE = 3
9      CHARACTER (LEN=10) :: SURFACE_TYPE
10     INTEGER :: MY_SURFACE_TYPE, N_IMAGE
11     TYPE (TEAM_TYPE) :: TEAM_SURFACE_TYPE
12
13     CALL GET_SURFACE_TYPE(THIS_IMAGE ( ), SURFACE_TYPE)
14     SELECT CASE (SURFACE_TYPE)
15     CASE ("LAND")
16         MY_SURFACE_TYPE = LAND
17     CASE ("SEA")
18         MY_SURFACE_TYPE = SEA
19     CASE ("ICE")
20         MY_SURFACE_TYPE = ICE
21     CASE DEFAULT
22         ERROR STOP
23     END SELECT
24     FORM TEAM (MY_SURFACE_TYPE, TEAM_SURFACE_TYPE)
25
26     CHANGE TEAM (TEAM_SURFACE_TYPE)
27     SELECT CASE (TEAM_NUMBER ( ))
28     CASE (LAND) ! Compute fluxes over land surface
29         CALL COMPUTE_FLUXES_LAND (FLUX_MOM, FLUX_SENS, FLUX_LAT)
30     CASE (SEA) ! Compute fluxes over sea surface
31         CALL COMPUTE_FLUXES_SEA (FLUX_MOM, FLUX_SENS, FLUX_LAT)
32     CASE (ICE) ! Compute fluxes over ice surface
33         CALL COMPUTE_FLUXES_ICE (FLUX_MOM, FLUX_SENS, FLUX_LAT)
34     CASE DEFAULT
35         ERROR STOP
36     END SELECT
37     END TEAM
38     END SUBROUTINE COMPUTE_FLUXES

```

## 39 C.6.7 Accessing coarrays in sibling teams

- 40 1 The following program illustrates subdividing a  $4 \times 4$  grid into  $2 \times 2$  teams, and the denotation of sibling teams.

```

41     PROGRAM DEMO
42         ! Initial team : 16 images. Algorithm design is a 4 by 4 grid.
43         ! Desire 4 teams, for the upper left (UL), upper right (UR),

```

```

1           !               lower left (LL), lower right (LR)
2 USE,INTRINSIC :: ISO_FORTRAN_ENV, ONLY: TEAM_TYPE
3 TYPE (TEAM_TYPE) :: t
4 INTEGER, PARAMETER :: UL=11, UR=22, LL=33, LR=44
5 REAL    :: A(10,10)[4,*]
6 INTEGER :: MYPE, TEAMNUM, NEWPE
7 INTEGER :: UL_IMAGE_LIST(4) = [1, 2, 5, 6], &
8           LL_IMAGE_LIST(4) = UL_IMAGE_LIST + 2, &
9           UR_IMAGE_LIST(4) = UL_IMAGE_LIST + 8, &
10          LR_IMAGE_LIST(4) = UL_IMAGE_LIST + 10
11
12 MYPE = THIS_IMAGE()
13 IF (ANY (MYPE==UL_IMAGE_LIST)) TEAMNUM = UL
14 IF (ANY (MYPE==LL_IMAGE_LIST)) TEAMNUM = LL
15 IF (ANY (MYPE==UR_IMAGE_LIST)) TEAMNUM = UR
16 IF (ANY (MYPE==LR_IMAGE_LIST)) TEAMNUM = LR
17 FORM TEAM (TEAMNUM, T)
18
19 A = 3.14
20
21 CHANGE TEAM (T, B[2,*] => A)
22     ! Inside change team, image pattern for B is a 2 by 2 grid.
23     B (5, 5) = B (1, 1)[2, 1]
24
25     ! Outside the team addressing:
26
27     NEWPE = THIS_IMAGE()
28     SELECT CASE (TEAM_NUMBER ())
29     CASE (UL)
30         IF (NEWPE==3) THEN
31             ! Right column of UL gets left column of UR.
32             B (:, 10) = B (:, 1)[1, 1, TEAM_NUMBER=UR]
33         ELSE IF (NEWPE==4) THEN
34             B (:, 10) = B (:, 1)[2, 1, TEAM_NUMBER=UR]
35         END IF
36     CASE (LL)
37         ! Similar to complete column exchange across middle of the original grid.
38         ...
39     END SELECT
40 END TEAM
41 END PROGRAM DEMO

```

## 42 C.6.8 Example involving failed images

43 1 Parallel algorithms often use work sharing schemes based on a specific mapping between image indices and global  
44 data addressing. To allow such programs to continue when one or more images fail, spare images can be used

1 to re-establish execution of the algorithm with the failed images replaced by spare images, while retaining the  
2 previous image mapping for nonfailed images.

3 2 The following example illustrates how this might be done. In this example, failure cannot be tolerated for image  
4 1 in the initial team.

```

5     PROGRAM possibly_recoverable_simulation
6         USE, INTRINSIC :: ISO_FORTRAN_ENV, ONLY: TEAM_TYPE, STAT_FAILED_IMAGE
7         IMPLICIT NONE
8         INTEGER, ALLOCATABLE :: failures (:)
9         INTEGER :: images_used, i, images_spare, status
10        INTEGER :: id [*], me [*]
11        TYPE (TEAM_TYPE) :: simulation_team
12        LOGICAL :: read_checkpoint, done [*]
13
14        ! Keep 1% spare images if we have a lot, just 1 if 10-199 images, 0 if <10.
15        images_spare = MAX (INT (0.01*NUM_IMAGES ()), 0, MIN (NUM_IMAGES () - 10, 1)
16        images_used = NUM_IMAGES () - images_spare
17        read_checkpoint = THIS_IMAGE () > images_used
18
19        setup : DO
20            me = THIS_IMAGE ()
21            id = MERGE (1, 2, me<=images_used)
22            !
23            ! Set up spare images as replacement for failed ones.
24            !
25            IF (IMAGE_STATUS (1) == STAT_FAILED_IMAGE) &
26                ERROR STOP "cannot recover"
27            IF (me == 1) THEN
28                failures = FAILED_IMAGES ()
29                k = images_used
30                DO i = 1, SIZE (failures)
31                    DO k = k+1, NUM_IMAGES ()
32                        IF (IMAGE_STATUS (k) == 0) EXIT
33                    END DO
34                    IF (k > NUM_IMAGES ()) ERROR STOP "cannot recover"
35                    me [k] = failures (i)
36                    id [k] = 1
37                END DO
38                images_used = k
39            END IF
40            !
41            ! Set up a simulation team of constant size.
42            ! Team 2 is the set of spares, so does not participate in the simulation.
43            !
44            FORM TEAM (id, simulation_team, NEW_INDEX=me, STAT=status)
45            simulation : CHANGE TEAM (simulation_team, STAT=status)

```

```

1      IF (status == STAT_FAILED_IMAGE) EXIT simulation
2      IF (TEAM_NUMBER () == 1) THEN
3          iter : DO
4              CALL simulation_procedure (read_checkpoint, status, done)
5              ! The simulation_procedure:
6              ! - sets up and performs some part of the simulation;
7              ! - resets to the last checkpoint if requested;
8              ! - sets status from its internal synchronizations;
9              ! - sets done to .TRUE. when the simulation has completed.
10             IF (status == STAT_FAILED_IMAGE) THEN
11                 read_checkpoint = .TRUE.
12                 EXIT simulation
13             ELSE IF (done)
14                 EXIT iter
15             END IF
16             read_checkpoint = .FALSE.
17         END DO iter
18     END IF
19     END TEAM simulation (STAT=status)
20     SYNC ALL (STAT=status)
21     IF (THIS_IMAGE () > images_used) done = done[1]
22     IF (done) EXIT setup
23 END DO setup
24 END PROGRAM possibly_recoverable_simulation

```

- 25 3 Supporting fault-tolerant execution imposes obligations on library writers who use the parallel language facilities.  
 26 Every synchronization statement, allocation or deallocation of coarrays, or invocation of a [collective procedure](#)  
 27 will need to be prepared to handle error conditions, and implicit deallocation of coarrays will need to be avoided.  
 28 Also, coarray module variables that are allocated inside the [team](#) execution context are not persistent.

### 29 C.6.9 EVENT\_QUERY example that tolerates image failure

- 30 1 This example is an adaptation of the later [EVENT\\_QUERY](#) example of [C.11.2](#) to make it able to execute in  
 31 the presence of the failure of one or more of the worker images. The function `create_work_item` now accepts an  
 32 integer argument to indicate which work item is required. It is assumed that the work items are indexed 1, 2,  
 33 ... . It is also assumed that if an image fails while processing a work item, that work item can subsequently be  
 34 processed by another image.

```

35     PROGRAM work_share
36         USE, INTRINSIC :: ISO_FORTRAN_ENV, ONLY: EVENT_TYPE
37         USE :: mod_work, ONLY: & ! Module that creates work items
38             work, & ! Type for holding a work item
39             create_work_item, & ! Function that creates work item
40             process_item, & ! Function that processes an item
41             work_done ! Logical function that returns true
42                 ! if all work done
43

```



```

1      TYPE :: worker_type
2          TYPE (EVENT_TYPE), ALLOCATABLE :: free (:)
3      END TYPE
4      TYPE (EVENT_TYPE) :: submit [*]      ! Whether work ready for a worker
5      TYPE (worker_type) :: worker [*]     ! Whether worker is free
6      TYPE (work) :: work_item [*] ! Holds the data for a work item
7      INTEGER :: count, i, k, kk, nbusy [*], np, status
8      INTEGER, ALLOCATABLE :: working (:) ! Items being worked on
9      INTEGER, ALLOCATABLE :: pending (:) ! Items pending after image failure
10
11     IF (THIS_IMAGE () == 1) THEN
12         ! Get started
13         ALLOCATE (worker%free (2:NUM_IMAGES ()))
14         ALLOCATE (working (2: NUM_IMAGES ()), pending(NUM_IMAGES ()-1))
15         nbusy = 0          ! This holds the number of workers working
16         k = 1             ! Index of next work item
17         np = 0           ! Number of work items in array pending
18         DO i = 2, NUM_IMAGES () ! Start the workers working
19             IF (work_done ()) EXIT
20             working (i) = 0
21             IF (IMAGE_STATUS (i) == STAT_FAILED_IMAGE) CYCLE
22             work_item [i] = create_work_item (k)
23             working (i) = k
24             k = k + 1
25             nbusy = nbusy + 1
26             EVENT POST (submit [i], STAT=status)
27         END DO
28         ! Main work distribution loop
29         master : DO
30             image : DO i = 2, NUM_IMAGES ()
31                 IF (IMAGE_STATUS (i) == STAT_FAILED_IMAGE) THEN
32                     IF (working (i)>0) THEN          ! It failed while working
33                         np = np + 1
34                         pending (np) = working (i)
35                         working (i) = 0
36                     END IF
37                     CYCLE image
38                 END IF
39                 CALL EVENT_QUERY (worker%free (i), count)
40                 IF (count == 0) CYCLE image          ! Worker is not free
41                 EVENT WAIT (worker%free (i))
42                 nbusy = nbusy - 1
43                 IF (np>0) THEN
44                     kk = pending (np)
45                     np = np - 1
46                 ELSE

```

```

1           IF (work_done ()) CYCLE image
2           kk = k
3           k = k + 1
4       END IF
5       nbusy = nbusy + 1
6       working (i) = kk
7       work_item [i] = create_work_item (kk)
8       EVENT POST (submit [i], STAT=status)
9       ! If image i has failed, the failure will be handled on
10      ! the next iteration of the master loop.
11  END DO image
12  IF ( nbusy==0 ) THEN ! All done. Exit on all images.
13      DO i = 2, NUM_IMAGES ()
14          EVENT POST (submit [i], STAT=status)
15          IF (status == STAT_FAILED_IMAGE) CYCLE
16      END DO
17      EXIT master
18  END IF
19  END DO master
20  ELSE
21      ! Work processing loop
22      worker : DO
23          EVENT WAIT (submit)
24          IF (nbusy [1] == 0) EXIT worker
25          CALL process_item(work_item)
26          EVENT POST (worker[1]%free (THIS_IMAGE ()))
27      END DO worker
28  END IF
29  END PROGRAM work_share

```

## 30 C.7 Clause 12 notes

### 31 C.7.1 External files (12.3)

32 1 This document accommodates, but does not require, file cataloging. To do this, several concepts are introduced.

#### 33 C.7.1.1 File existence (12.3.2)

34 1 Totally independent of the connection state is the property of existence, this being a file property. The processor  
35 “knows” of a set of files that exist at a given time for a given program. This set would include tapes ready to  
36 read, files in a catalog, a keyboard, a printer, etc. The set might exclude files inaccessible to the program because  
37 of security, because they are already in use by another program, etc. This document does not specify which  
38 files exist, hence wide latitude is available to a processor to implement security, locks, privilege techniques, etc.  
39 Existence is a convenient concept to designate all of the files that a program can potentially process.

1 2 All four combinations of connection and existence can occur:

| Connect | Exist | Examples                                             |
|---------|-------|------------------------------------------------------|
| Yes     | Yes   | A card reader loaded and ready to be read            |
| Yes     | No    | A printer before the first line is written           |
| No      | Yes   | A file named 'JOAN' in the catalog                   |
| No      | No    | A file on a reel of tape, not known to the processor |

2 3 Means are provided to create, delete, connect, and disconnect files.

### 3 **C.7.1.2 File access (12.3.3)**

4 1 This document does not address problems of security, protection, locking, and many other concepts that might  
5 be part of the concept of “right of access”. Such concepts are considered to be in the province of an operating  
6 system.

7 2 The [OPEN](#) and [INQUIRE](#) statements can be extended naturally to consider these things.

8 3 Possible access methods for a file are: sequential, stream and direct. The processor might implement three  
9 different types of files, each with its own access method. It might instead implement one type of file with three  
10 different access methods.

11 4 Direct access to files is of a simple and commonly available type, that is, fixed-length records. The key is a  
12 positive integer.

### 13 **C.7.1.3 File connection (12.5)**

14 1 Before any input/output can be performed on a file, it needs to be connected to a [unit](#). The [unit](#) then serves as a  
15 designator for that file as long as it is connected. To be connected does not imply that “buffers” have or have not  
16 been allocated, that “file-control tables” have or have not been filled, or that any other method of implementation  
17 has been used. Connection means that (barring some other fault) a [READ](#) or [WRITE](#) statement can be executed  
18 on the [unit](#), hence on the file. Without a connection, a [READ](#) or [WRITE](#) statement cannot be executed.

### 19 **C.7.1.4 File names (12.5.6.10)**

20 1 A file can have a name. The form of a file name is not specified. If a system does not have some form of cataloging  
21 or tape labeling for at least some of its files, all file names disappear at the [termination of execution](#). This is a  
22 valid implementation. Nowhere does this document require names to survive for any period of time longer than  
23 the execution time span of a program. Therefore, this document does not impose cataloging as a prerequisite.  
24 The naming feature is intended to enable use of a cataloging system where one exists.

## 25 **C.7.2 Nonadvancing input/output (12.3.4.2)**

26 1 [Data transfer statements](#) affect the positioning of an [external file](#). In FORTRAN 77, if no error or end-of-file  
27 condition exists, the file is positioned after the record just read or written and that record becomes the preceding  
28 record. This document contains the [ADVANCE= specifier](#) in a [data transfer statement](#) that provides the capab-  
29 ility of maintaining a position within the current record from one formatted [data transfer statement](#) to the next  
30 [data transfer statement](#). The value NO provides this capability. The value YES positions the file after the record  
31 just read or written. The default is YES.

- 1 2 The tab edit descriptor and the slash are still appropriate for use with this type of record access but the tab  
2 cannot reposition before the left tab limit.
- 3 3 A **BACKSPACE** of a file that is positioned within a record causes the specified unit to be positioned before the  
4 current record.
- 5 4 If the next input/output operation on a file after a nonadvancing write is a **rewind**, **backspace**, **end file** or **close**  
6 operation, the file is positioned implicitly after the current record before an **ENDFILE** record is written to the  
7 file, that is, a **REWIND**, **BACKSPACE**, or **ENDFILE** statement following a nonadvancing **WRITE statement**  
8 causes the file to be positioned at the end of the current output record before the endfile record is written to the  
9 file.
- 10 5 This document provides a **SIZE= specifier** to be used with formatted **data transfer statements**. The variable in  
11 the **SIZE= specifier** is assigned the count of the number of characters that make up the sequence of values read  
12 by the data edit descriptors in the **input statement**.
- 13 6 The count is especially helpful if there is only one list item in the input list because it is the number of characters  
14 that appeared for the item.
- 15 7 The **EOR= specifier** is provided to indicate when an EOR condition is encountered during a nonadvancing **data**  
16 **transfer statement**. The EOR condition is not an error condition. If this specifier appears, an input list item that  
17 requires more characters than the record contained is padded with blanks if **PAD= 'YES'** is in effect. This means  
18 that the input list item completed successfully. The file is positioned after the current record. If the **IOSTAT=**  
19 **specifier** appears, the specified variable is defined with the value of the **named constant IOSTAT\_EOR** from  
20 the intrinsic module **ISO\_FORTRAN\_ENV** and the **data transfer statement** is terminated. Program execution  
21 continues with the statement specified in the **EOR= specifier**. The **EOR= specifier** gives the capability of taking  
22 control of execution when the EOR condition is encountered. The *do-variables* in *io-implied-dos* retain their last  
23 defined value and any remaining items in the *input-item-list* retain their definition status when an EOR condition  
24 occurs. If the **SIZE= specifier** appears, the specified variable is assigned the number of characters read with the  
25 data edit descriptors during the **READ statement**.
- 26 8 For nonadvancing input, the processor is not required to read partial records. The processor could read the entire  
27 record into an internal buffer and make successive portions of the record available to successive **input statements**.
- 28 9 In an implementation of nonadvancing input/output in which a nonadvancing write to a terminal device causes  
29 immediate display of the output, such a write can be used as a mechanism to output a prompt. In this case, the  
30 statement

```
31 WRITE (*, FMT='(A)', ADVANCE='NO') 'CONTINUE?(Y/N): '
```

32 would result in the prompt

```
33 CONTINUE?(Y/N):
```

34 being displayed with no subsequent line feed.

- 35 10 The response, which might be read by a statement of the form

```
36 READ (*, FMT='(A)') ANSWER
```

37 can then be entered on the same line as the prompt as in

```
38 CONTINUE?(Y/N): Y
```

1 11 This document does not require that an implementation of nonadvancing input/output operate in this manner.  
 2 For example, an implementation of nonadvancing output in which the display of the output is deferred until  
 3 the current record is complete is also standard-conforming. Such an implementation will not, however, allow a  
 4 prompting mechanism of this kind to operate.

### 5 C.7.3 OPEN statement (12.5.6)

6 1 A file can become connected to a **unit** either by preconnection or by execution of an **OPEN statement**. Precon-  
 7 nection is performed prior to the beginning of execution of a program by means external to Fortran. For example,  
 8 it could be done by job control action or by processor-established defaults. Execution of an **OPEN statement** is  
 9 not required in order to access preconnected files (12.5.5).

10 2 The **OPEN statement** provides a means to access existing files that are not preconnected. An **OPEN statement**  
 11 can be used in either of two ways: with a file name (open-by-name) and without a file name (open-by-unit). A  
 12 **unit** is given in either case. Open-by-name connects the specified file to the specified **unit**. Open-by-unit connects  
 13 a processor-dependent default file to the specified **unit**. (The default file might or might not have a name.)

14 3 Therefore, there are three ways a file can become connected and hence processed: preconnection, open-by-name,  
 15 and open-by-unit. Once a file is connected, there is no means in standard Fortran to determine how it became  
 16 connected.

17 4 An **OPEN statement** can also be used to create a new file. In fact, any of the foregoing three connection methods  
 18 can be performed on a file that does not exist. When a **unit** is preconnected, writing the first record creates the  
 19 file. With the other two methods, execution of the **OPEN statement** creates the file.

20 5 When an **OPEN statement** is executed, the **unit** specified in the **OPEN statement** might or might not already be  
 21 connected to a file. If it is already connected to a file (either through preconnection or by prior execution of an  
 22 **OPEN statement**), then omitting the **FILE= specifier** in the **OPEN statement** implies that the file is to remain  
 23 connected to the **unit**. Such an **OPEN statement** can be used to change the values of the blank interpretation  
 24 mode, decimal edit mode, pad mode, input/output rounding mode, delimiter mode, and sign mode.

25 6 If the value of the **ACTION= specifier** is WRITE, then a **READ statement** cannot refer to the connection.  
 26 **ACTION = 'WRITE'** does not restrict positioning by a **BACKSPACE statement** or positioning specified by the  
 27 **POSITION= specifier** with the value APPEND. However, a **BACKSPACE statement** or an **OPEN statement**  
 28 containing **POSITION = 'APPEND'** might fail if the processor needs to read the file to achieve the positioning.

29 7 The following examples illustrate these rules. In the first example, **unit 10** is preconnected to a SCRATCH file;  
 30 the **OPEN statement** changes the value of **PAD=** to YES.

```
31 CHARACTER (LEN = 20) CH1
32 WRITE (10, '(A)') 'THIS IS RECORD 1'
33 OPEN (UNIT = 10, STATUS = 'OLD', PAD = 'YES')
34 REWIND 10
35 READ (10, '(A20)') CH1 ! CH1 now has the value
36 ! 'THIS IS RECORD 1'
```

37 8 In the next example, **unit 12** is first connected to a file named FRED, with a status of OLD. The second **OPEN**  
 38 **statement** then opens **unit 12** again, retaining the connection to the file FRED, but changing the value of the  
 39 **DELIM= specifier** to QUOTE.

```

1      CHARACTER (LEN = 25) CH2, CH3
2      OPEN (12, FILE = 'FRED', STATUS = 'OLD', DELIM = 'NONE')
3      CH2 = '"THIS STRING HAS QUOTES."'
4              ! Quotes in string CH2
5      WRITE (12, *) CH2          ! Written with no delimiters
6      OPEN (12, DELIM = 'QUOTE') ! Now quote is the delimiter
7      REWIND 12
8      READ (12, *) CH3  ! CH3 now has the value
9              ! 'THIS STRING HAS QUOTES. '

```

10 9 The next example is invalid because it attempts to change the value of the `STATUS=` specifier.

```

11      OPEN (10, FILE = 'FRED', STATUS = 'OLD')
12      WRITE (10, *) A, B, C
13      OPEN (10, STATUS = 'SCRATCH') ! Attempts to make FRED a SCRATCH file

```

14 10 The previous example could be made valid by closing the `unit` first, as in the next example.

```

15      OPEN (10, FILE = 'FRED', STATUS = 'OLD')
16      WRITE (10, *) A, B, C
17      CLOSE (10)
18      OPEN (10, STATUS = 'SCRATCH') ! Opens a different SCRATCH file

```

#### 19 C.7.4 Connection properties (12.5.4)

20 1 When a `unit` becomes connected to a file, either by execution of an `OPEN statement` or by preconnection, the  
 21 following connection properties, among others, are established.

- 22 (1) An access method, which is sequential, direct, or stream, is established for the connection (12.5.6.3).
- 23 (2) A form, which is formatted or unformatted, is established for a connection to a file that exists or  
 24 is created by the connection. For a connection that results from execution of an `OPEN statement`,  
 25 a default form (which depends on the access method, as described in 12.3.3) is established if no  
 26 form is specified. For a preconnected file that exists, a form is established by preconnection. For a  
 27 preconnected file that does not exist, a form might be established, or the establishment of a form  
 28 might be delayed until the file is created (for example, by execution of a formatted or unformatted  
 29 `WRITE statement`) (12.5.6.11).
- 30 (3) A record length might be established. If the access method is direct, the connection establishes a  
 31 record length that specifies the length of each record of the file. A direct access file can only contain  
 32 records that are all of equal length.
- 33 (4) A sequential file can contain records of varying lengths. In this case, the record length established  
 34 specifies the maximum length of a record in the file (12.5.6.15).

35 2 A processor has wide latitude in adapting these concepts and actions to its own cataloging and job control  
 36 conventions. Some processors might need job control action to specify the set of files that exist or that will  
 37 be created by a program. Some processors might not need any job control action prior to execution. This  
 38 document enables processors to perform dynamic open, close, or file creation operations, but it does not require  
 39 such capabilities of the processor.

1 3 The meaning of “open” in contexts other than Fortran might include such things as mounting a tape, console  
2 messages, spooling, label checking, security checking, etc. These actions might occur upon job control action  
3 external to Fortran, upon execution of an **OPEN statement**, or upon execution of the first read or write of the  
4 file. The **OPEN statement** describes properties of the connection to the file and might or might not cause physical  
5 activities to take place.

### 6 **C.7.5 Asynchronous input/output (12.6.2.5)**

7 1 Rather than limit support for asynchronous input/output to what has been traditionally provided by facilities  
8 such as BUFFERIN/BUFFEROUT, this document builds upon existing Fortran syntax. This permits alternative  
9 approaches for implementing asynchronous input/output, and simplifies the task of adapting existing standard-  
10 conforming programs to use asynchronous input/output.

11 2 Not all processors actually perform input/output asynchronously, nor will every processor that does be able to  
12 handle **data transfer statements** with complicated input/output item lists in an asynchronous manner. Such  
13 processors can still be standard-conforming.

14 3 This document allows for at least two different conceptual models for asynchronous input/output.

15 4 Model 1: the processor performs asynchronous input/output when the item list is simple (perhaps one contiguous  
16 named array) and the input/output is unformatted. The implementation cost is reduced, and this is the scenario  
17 most likely to be beneficial on traditional “big-iron” machines.

18 5 Model 2: The processor is free to do any of the following:

19 (1) on output, create a buffer inside the input/output library, completely formatted, and then start an  
20 asynchronous write of the buffer, and immediately return to the next statement in the program. The  
21 processor is free to wait for previously issued WRITES, or not, or

22 (2) pass the input/output list addresses to another processor/process, which processes the list items  
23 independently of the processor that executes the user’s code. The addresses of the list items will  
24 need to be computed before the asynchronous **READ/WRITE** statement completes. There is still  
25 an ordering requirement on list item processing to handle things like **READ (...) N,(a(i),i=1,N)**.

26 6 A program can issue a large number of asynchronous input/output requests, without waiting for any of them to  
27 complete, and then wait for any or all of them. That does not constitute a requirement for the processor to keep  
28 track of each individual request separately.

29 7 It is not necessary for all requests to be tracked by the runtime library. If an **ID= specifier** does not appear in on a  
30 **READ** or **WRITE** statement, the runtime library can forget about this particular request once it has successfully  
31 completed. If an error or end-of-file condition occurs for a request, the processor can report this during any  
32 input/output operation to that **unit**. If an **ID= specifier** appears, the processor’s runtime input/output library  
33 will need to keep track of any end-of-file or error conditions for that particular input/output request. However, if  
34 the input/output request succeeds without any exceptional conditions occurring, then the runtime can forget that  
35 **ID=** value. A runtime library might only keep track of the last request made, or perhaps a very few. Then, when  
36 a user **WAITs** for a particular request, either the library will know about it (and does the right thing with respect  
37 to error handling, etc.), or can assume it is a request that successfully completed and was forgotten about (and  
38 will just return without signaling any end-of-file or error condition). A standard-conforming program can only  
39 pass valid **ID=** values, but there is no requirement on the processor to detect invalid **ID=** values. There might  
40 be a processor dependent limit on how many outstanding input/output requests that generate an end-of-file or

1 error condition can be handled before the processor runs out of memory to keep track of such conditions. The  
2 restrictions on the `SIZE=` variables are designed to enable the processor to update such variables at any time  
3 (after the request has been processed, but before the wait operation), and then forget about them. Only error and  
4 end-of-file conditions are expected to be tracked by individual request by the runtime, and then only if an `ID=`  
5 `specifier` appears. The `END=` and `EOR=` specifiers have not been added to all statements that can perform wait  
6 operations. Instead, the `IOSTAT` variable can be queried after a wait operation to handle this situation. This  
7 choice was made because the `WAIT statement` is expected to be the usual method of waiting for input/output  
8 to complete (and `WAIT` does support the `END=` and `EOR=` specifiers). This particular choice is philosophical,  
9 and was not based on significant technical difficulties.

- 10 8 The requirement to set the `IOSTAT` variable correctly means that a processor will need to remember which  
11 input/output requests encountered an end-of-record condition, so that a subsequent wait operation can return  
12 the correct `IOSTAT` value. Therefore there might be a processor defined limit on the number of outstanding  
13 nonadvancing input/output requests that have encountered an end-of-record condition (constrained by available  
14 memory to keep track of this information, similar to end-of-file and error conditions).

## 15 C.8 Clause 13 notes

### 16 C.8.1 Number of records (13.4, 13.5, 13.8.2)

- 17 1 The number of records read by an explicitly formatted advancing input statement can be determined from the  
18 following rule: a record is read at the beginning of the format scan (even if the input list is empty unless the most  
19 recently previous operation on the `unit` was not a nonadvancing read operation), at each slash edit descriptor  
20 encountered in the format, and when a format rescan occurs at the end of the format.
- 21 2 The number of records written by an explicitly formatted advancing output statement can be determined from  
22 the following rule: a record is written when a slash edit descriptor is encountered in the format, when a format  
23 rescan occurs at the end of the format, and at completion of execution of an advancing `output statement` (even if  
24 the output list is empty). Thus, the occurrence of  $n$  successive slashes between two other edit descriptors causes  
25  $n - 1$  blank lines if the records are printed. The occurrence of  $n$  slashes at the beginning or end of a complete  
26 format specification causes  $n$  blank lines if the records are printed. However, a complete format specification  
27 containing  $n$  slashes ( $n > 0$ ) and no other edit descriptors causes  $n + 1$  blank lines if the records are printed. For  
28 example, the statements

```
29 PRINT 3  
30 3 FORMAT (/)
```

31 will write two records that cause two blank lines if the records are printed.

### 32 C.8.2 List-directed input (13.10.3)

- 33 1 The following examples illustrate list-directed input. A blank character is represented by b.

#### 34 2 Example 1:

##### 35 Program:

```
36 J = 3  
37 READ *, I  
38 READ *, J
```



1 **Sequential input file:**

2           record 1:  b1b,4bbbb  
3           record 2:  ,2bbbbbbb

4   3 Result: I = 1, J = 3.

5   4 Explanation: The second **READ statement** reads the second record. The initial comma in the record designates  
6   a null value; therefore, J is not redefined.

7 **5 Example 2:**

8 **Program:**

9           CHARACTER A \*8, B \*1  
10          READ \*, A, B

11 **Sequential input file:**

12          record 1:  'bbbbbbbb'  
13          record 2:  'QXY'b'Z'

14   6 Result: A = 'bbbbbbbb', B = 'Q'

15   7 Explanation: In the first record, the rightmost apostrophe is interpreted as delimiting the constant (it cannot  
16   be the first of a pair of embedded apostrophes representing a single apostrophe because this would involve  
17   the prohibited “splitting” of the pair by the end of a record); therefore, A is assigned the character constant  
18   'bbbbbbbb'. The end of a record acts as a blank, which in this case is a value separator because it occurs between  
19   two constants.

20 **C.9 Clause 14 notes**

21 **C.9.1 Main program and block data program unit (14.1, 14.3)**

22   1 The name of the main program or of a block data program unit has no explicit use within the Fortran language.  
23   It is available for documentation and for possible use by a processor.

24   2 A processor might implement an unnamed program unit by assigning it a global identifier that is not used  
25   elsewhere in the program. This could be done by using a default name that does not satisfy the rules for Fortran  
26   names.

27 **C.9.2 Dependent compilation (14.2)**

28   1 This document, like its predecessors, is intended to enable the implementation of conforming processors in which  
29   a program can be broken into multiple units, each of which can be separately translated in preparation for  
30   execution. Such processors are commonly described as supporting separate compilation. There is an important  
31   difference between the way separate compilation can be implemented under this document and the way it could be  
32   implemented under the FORTRAN 77 International Standard. Under the FORTRAN 77 standard, any information

1 required to translate a **program unit** was specified in that **program unit**. Each translation was thus totally  
2 independent of all others. Under this document, a **program unit** can use information that was specified in a  
3 separate module and thus can be dependent on that module. The implementation of this dependency in a  
4 processor might be that the translation of a **program unit** depends on the results of translating one or more  
5 modules. Processors implementing the dependency this way are commonly described as supporting dependent  
6 compilation.

7 2 The dependencies involved here are new only in the sense that the Fortran processor is now aware of them. The  
8 same information dependencies existed under the FORTRAN 77 International Standard, but it was the program-  
9 mer's responsibility to transport the information necessary to resolve them by making redundant specifications of  
10 the information in multiple **program units**. The availability of separate but dependent compilation offers several  
11 potential advantages over the redundant textual specification of information.

12 (1) Specifying information at a single place in the program ensures that different **program units** using that  
13 information are translated consistently. Redundant specification leaves the possibility that different  
14 information can be erroneously be specified. Even if an INCLUDE line is used to ensure that the  
15 text of the specifications is identical in all involved **program units**, the presence of other specifications  
16 (for example, an **IMPLICIT statement**) could change the interpretation of that text.

17 (2) During the revision of a program, it is possible for a processor to assist in determining whether differ-  
18 ent **program units** have been translated using different (incompatible) versions of a module, although  
19 there is no requirement that a processor provide such assistance. Inconsistencies in redundant textual  
20 specification of information, on the other hand, tend to be much more difficult to detect.

21 (3) Putting information in a module provides a way of packaging it. Without modules, redundant spe-  
22 cifications frequently are interleaved with other specifications in a **program unit**, making convenient  
23 packaging of such information difficult.

24 (4) Because a processor can be implemented such that the specifications in a module are translated once  
25 and then repeatedly referenced, there is the potential for greater efficiency than when the processor  
26 translates redundant specifications of information in multiple **program units**.

27 3 The exact meaning of the requirement that the public portions of a module be available at the time of reference  
28 is processor dependent. For example, a processor could consider a module to be available only after it has been  
29 compiled and require that if the module has been compiled separately, the result of that compilation be identified  
30 to the compiler when compiling **program units** that use it.

### 31 **C.9.2.1 USE statement and dependent compilation (14.2.2)**

32 1 Another benefit of the **USE statement** is its enhanced facilities for name management. If one needs to use only  
33 selected entities in a module, one can do so without having to worry about the names of all the other entities  
34 in that module. If one needs to use two different modules that happen to contain entities with the same name,  
35 there are several ways to deal with the conflict. If none of the entities with the same name are to be used, they  
36 can simply be ignored. If the name happens to refer to the same entity in both modules (for example, if both  
37 modules obtained it from a third module), then there is no confusion about what the name denotes and the name  
38 can be freely used. If the entities are different and one or both is to be used, the local renaming facility in the  
39 **USE statement** makes it possible to give those entities different names in the **program unit** containing the **USE**  
40 **statements**.

41 2 A benefit of using the **ONLY** option consistently, as compared to **USE** without it, is that the module from which  
42 each accessed entity is accessed is explicitly specified in each **program unit**. This means that one need not search

- 1 other [program units](#) to find where each one is defined. This reduces maintenance costs.
- 2 3 A typical implementation of dependent but separate compilation might involve storing the result of translating a  
3 module in a file whose name is derived from the name of the module. Note, however, that the name of a module  
4 is limited only by the Fortran rules and not by the names allowed in the file system. Thus the processor might  
5 have to provide a mapping between Fortran names and file system names.
- 6 4 The result of translating a module could reasonably either contain only the information textually specified in the  
7 module (with “pointers” to information originally textually specified in other modules) or contain all information  
8 specified in the module (including copies of information originally specified in other modules). Although the former  
9 approach would appear to save on storage space, the latter approach can greatly simplify the logic necessary to  
10 process a [USE statement](#) and can avoid the necessity of imposing a limit on the logical “nesting” of modules via  
11 the [USE statement](#).
- 12 5 There is an increased potential for undetected errors in a [scoping unit](#) that uses both implicit typing and the  
13 [USE statement](#). For example, in the program fragment

```

14         SUBROUTINE SUB
15             USE MY_MODULE
16             IMPLICIT INTEGER (I-N), REAL (A-H, O-Z)
17             X = F (B)
18             A = G (X) + H (X + 1)
19         END SUBROUTINE SUB

```

- 20 X could be either an implicitly typed real variable or a variable obtained from the module MY\_MODULE and  
21 might change from one to the other because of changes in MY\_MODULE unrelated to the action performed by  
22 SUB. Logic errors resulting from this kind of situation can be extremely difficult to locate. Thus, the use of these  
23 features together is discouraged.

#### 24 C.9.2.2 Accessibility attributes (8.5.2)

- 25 1 The [PUBLIC](#) and [PRIVATE](#) attributes, which can be declared only in modules, divide the entities in a module  
26 into those that are actually relevant to a [scoping unit](#) referencing the module and those that are not. This  
27 information might be used to improve the performance of a Fortran processor. For example, it might be possible  
28 to discard much of the information about the private entities once a module has been translated, thus saving on  
29 both storage and the time to search it. Similarly, it might be possible to recognize that two versions of a module  
30 differ only in the private entities they contain and avoid retranslating [program units](#) that use that module when  
31 switching from one version of the module to the other.

### 32 C.9.3 Examples of the use of modules (14.2.1)

#### 33 C.9.3.1 Global data (14.2.1)

- 34 1 A module could contain only data objects, for example:

```

35         MODULE DATA_MODULE
36             SAVE
37             REAL A (10), B, C (20,20)
38             INTEGER :: I=0
39             INTEGER, PARAMETER :: J=10

```

```

1         COMPLEX D (J,J)
2         END MODULE DATA_MODULE

```

3 2 Data objects made global in this manner can have any combination of data types.

4 3 Access to some of these can be made by a [USE statement](#) with the [ONLY](#) option, such as:

```

5         USE DATA_MODULE, ONLY: A, B, D

```

6 and access to all of them can be made by the following [USE statement](#):

```

7         USE DATA_MODULE

```

8 4 Access to all of them with some renaming to avoid name conflicts can be made by, for example:

```

9         USE DATA_MODULE, AMODULE => A, DMODULE => D

```

### 10 **C.9.3.2 Derived types (14.2.1)**

11 1 A derived type can be defined in a module and accessed in a number of [program units](#). For example,

```

12         MODULE SPARSE
13             TYPE NONZERO
14                 REAL A
15                 INTEGER I, J
16             END TYPE NONZERO
17         END MODULE SPARSE

```

18 defines a type consisting of a real component and two integer components for holding the numerical value of a  
19 nonzero matrix element and its row and column indices.

### 20 **C.9.3.3 Global allocatable arrays (14.2.1)**

21 1 Many programs need large global allocatable arrays whose sizes are not known before program execution. A  
22 simple form for such a program is:

```

23         PROGRAM GLOBAL_WORK
24             CALL CONFIGURE_ARRAYS      ! Perform the appropriate allocations
25             CALL COMPUTE              ! Use the arrays in computations
26         END PROGRAM GLOBAL_WORK
27         MODULE WORK_ARRAYS          ! An example set of work arrays
28             INTEGER N
29             REAL, ALLOCATABLE :: A (:), B (:, :), C (:, :, :)
30         END MODULE WORK_ARRAYS
31         SUBROUTINE CONFIGURE_ARRAYS ! Process to set up work arrays
32             USE WORK_ARRAYS
33             READ (*, *) N
34             ALLOCATE (A (N), B (N, N), C (N, N, 2 * N))
35         END SUBROUTINE CONFIGURE_ARRAYS
36         SUBROUTINE COMPUTE

```

```

1      USE WORK_ARRAYS
2      ... ! Computations involving arrays A, B, and C
3      END SUBROUTINE COMPUTE

```

4 2 Typically, many subprograms need access to the work arrays, and all such subprograms would contain the  
5 statement

```

6      USE WORK_ARRAYS

```

#### 7 **C.9.3.4 Procedure libraries (14.2.2)**

8 1 [Interface bodies](#) for [external procedures](#) in a library can be gathered into a module. An [interface body](#) specifies  
9 an [explicit interface](#) (15.4.2.2).

10 2 An example is the following library module:

```

11      MODULE LIBRARY_LLS
12      INTERFACE
13      SUBROUTINE LLS (X, A, F, FLAG)
14      REAL X (:, :)
15      ! The SIZE in the next statement is an intrinsic function
16      REAL, DIMENSION (SIZE (X, 2)) :: A, F
17      INTEGER FLAG
18      END SUBROUTINE LLS
19      ...
20      END INTERFACE
21      ...
22      END MODULE LIBRARY_LLS

```

23 3 This module provides an [explicit interface](#) that is necessary for the subroutine LLS to be invoked. for example:

```

24      USE LIBRARY_LLS
25      ...
26      CALL LLS (X = ABC, A = D, F = XX, FLAG = IFLAG)
27      ...

```

28 4 Because [dummy argument](#) names in an [interface body](#) for an [external procedure](#) are not required to be the same  
29 as in the procedure definition, different versions can be constructed for different applications using [argument](#)  
30 [keywords](#) appropriate to each application.

#### 31 **C.9.3.5 Operator extensions (14.2.2)**

32 1 In order to extend an intrinsic operator symbol to have an additional meaning, an [interface block](#) specifying that  
33 operator symbol in the [OPERATOR](#) option of the [INTERFACE statement](#) could be placed in a module.

34 2 For example, // can be extended to perform concatenation of two derived-type objects serving as varying length  
35 character strings and + can be extended to specify matrix addition for type MATRIX or interval arithmetic  
36 addition for type INTERVAL.

1 3 A module might contain several such [interface blocks](#). An operator can be defined by an external function (either  
2 in Fortran or some other language) and its procedure [interface](#) placed in the module.

### 3 **C.9.3.6 Data abstraction (14.2.2)**

4 1 In addition to providing a portable means of avoiding the redundant specification of information in multiple  
5 [program units](#), a module provides a convenient means of “packaging” related entities, such as the definitions of  
6 the representation and operations of an abstract data type. The following example of a module defines a data  
7 abstraction for a SET type where the elements of each set are of type integer. The usual set operations of UNION,  
8 INTERSECTION, and DIFFERENCE are provided. The CARDINALITY function returns the cardinality of  
9 (number of elements in) its set argument. Two functions returning logical values are included, ELEMENT and  
10 SUBSET. ELEMENT defines the operator .IN. and SUBSET extends the operator <=. ELEMENT determines  
11 if a given scalar integer value is an element of a given set, and SUBSET determines if a given set is a subset of  
12 another given set. (Two sets can be checked for equality by comparing cardinality and checking that one is a  
13 subset of the other, or checking to see if each is a subset of the other.)

14 2 The transfer function SETF converts a vector of integer values to the corresponding set, with duplicate values  
15 removed. Thus, a vector of constant values can be used as set constants. An inverse transfer function VECTOR  
16 returns the elements of a set as a vector of values in ascending order. In this SET implementation, set data  
17 objects have a maximum cardinality of 200.

18 3 Here is the example module:

```
19     MODULE INTEGER_SETS
20     ! This module is intended to illustrate use of the module facility
21     ! to define a new type, along with suitable operators.
22
23     INTEGER, PARAMETER :: MAX_SET_CARD = 200
24
25     TYPE SET                                ! Define SET type
26     PRIVATE
27     INTEGER CARD
28     INTEGER ELEMENT (MAX_SET_CARD)
29     END TYPE SET
30
31     INTERFACE OPERATOR (.IN.)
32     MODULE PROCEDURE ELEMENT
33     END INTERFACE OPERATOR (.IN.)
34
35     INTERFACE OPERATOR (<=)
36     MODULE PROCEDURE SUBSET
37     END INTERFACE OPERATOR (<=)
38
39     INTERFACE OPERATOR (+)
40     MODULE PROCEDURE UNION
41     END INTERFACE OPERATOR (+)
42
43     INTERFACE OPERATOR (-)
```

```

1      MODULE PROCEDURE DIFFERENCE
2      END INTERFACE OPERATOR (-)
3
4      INTERFACE OPERATOR (*)
5      MODULE PROCEDURE INTERSECTION
6      END INTERFACE OPERATOR (*)
7
8      CONTAINS
9
10     INTEGER FUNCTION CARDINALITY (A)      ! Returns cardinality of set A
11     TYPE (SET), INTENT (IN) :: A
12     CARDINALITY = A % CARD
13     END FUNCTION CARDINALITY
14
15     LOGICAL FUNCTION ELEMENT (X, A)      ! Determines if
16     INTEGER, INTENT(IN) :: X           ! element X is in set A
17     TYPE (SET), INTENT(IN) :: A
18     ELEMENT = ANY (A % ELEMENT (1 : A % CARD) == X)
19     END FUNCTION ELEMENT
20
21     FUNCTION UNION (A, B)                ! Union of sets A and B
22     TYPE (SET) UNION
23     TYPE (SET), INTENT(IN) :: A, B
24     INTEGER J
25     UNION = A
26     DO J = 1, B % CARD
27         IF (.NOT. (B % ELEMENT (J) .IN. A)) THEN
28             IF (UNION % CARD < MAX_SET_CARD) THEN
29                 UNION % CARD = UNION % CARD + 1
30                 UNION % ELEMENT (UNION % CARD) = B % ELEMENT (J)
31             ELSE
32                 ! Maximum set size exceeded . . .
33             END IF
34         END IF
35     END DO
36     END FUNCTION UNION
37
38     FUNCTION DIFFERENCE (A, B)           ! Difference of sets A and B
39     TYPE (SET) DIFFERENCE
40     TYPE (SET), INTENT(IN) :: A, B
41     INTEGER J, X
42     DIFFERENCE % CARD = 0                ! The empty set
43     DO J = 1, A % CARD
44         X = A % ELEMENT (J)
45         IF (.NOT. (X .IN. B)) DIFFERENCE = DIFFERENCE + SET (1, X)
46     END DO

```

```

1      END FUNCTION DIFFERENCE
2
3      FUNCTION INTERSECTION (A, B)          ! Intersection of sets A and B
4          TYPE (SET) INTERSECTION
5          TYPE (SET), INTENT(IN) :: A, B
6          INTERSECTION = A - (A - B)
7      END FUNCTION INTERSECTION
8
9      LOGICAL FUNCTION SUBSET (A, B)        ! Determines if set A is
10         TYPE (SET), INTENT(IN) :: A, B    ! a subset of set B
11         INTEGER I
12         SUBSET = A % CARD <= B % CARD
13         IF (.NOT. SUBSET) RETURN          ! For efficiency
14         DO I = 1, A % CARD
15             SUBSET = SUBSET .AND. (A % ELEMENT (I) .IN. B)
16         END DO
17     END FUNCTION SUBSET
18
19     TYPE (SET) FUNCTION SETF (V)          ! Transfer function between a vector
20         INTEGER V (:)                    ! of elements and a set of elements
21         INTEGER J                          ! removing duplicate elements
22         SETF % CARD = 0
23         DO J = 1, SIZE (V)
24             IF (.NOT. (V (J) .IN. SETF)) THEN
25                 IF (SETF % CARD < MAX_SET_CARD) THEN
26                     SETF % CARD = SETF % CARD + 1
27                     SETF % ELEMENT (SETF % CARD) = V (J)
28                 ELSE
29                     ! Maximum set size exceeded . . .
30                 END IF
31             END IF
32         END DO
33     END FUNCTION SETF
34
35     FUNCTION VECTOR (A)                   ! Transfer the values of set A
36         TYPE (SET), INTENT (IN) :: A     ! into a vector in ascending order
37         INTEGER, POINTER :: VECTOR (:)
38         INTEGER I, J, K
39         ALLOCATE (VECTOR (A % CARD))
40         VECTOR = A % ELEMENT (1 : A % CARD)
41         DO I = 1, A % CARD - 1            ! Use a better sort if
42             DO J = I + 1, A % CARD        ! A % CARD is large
43                 IF (VECTOR (I) > VECTOR (J)) THEN
44                     K = VECTOR (J); VECTOR (J) = VECTOR (I); VECTOR (I) = K
45                 END IF
46             END DO

```



```

1           END DO
2       END FUNCTION VECTOR
3
4       END MODULE INTEGER_SETS

5 4 Examples of using INTEGER_SETS (A, B, and C are variables of type SET; X is an integer variable):
6           ! Check to see if A has more than 10 elements
7       IF (CARDINALITY (A) > 10) ...
8
9           ! Check for X an element of A but not of B
10      IF (X .IN. (A - B)) ...
11
12           ! C is the union of A and the result of B intersected
13           ! with the integers 1 to 100
14      C = A + B * SETF ([[I, I = 1, 100]])
15
16           ! Does A have any even numbers in the range 1:100?
17      IF (CARDINALITY (A * SETF ([[I, I = 2, 100, 2]])) > 0) ...
18
19      PRINT *, VECTOR (B) ! Print out the elements of set B, in ascending order

```

### 20 C.9.3.7 Public entities renamed (14.2.2)

- 21 1 At times it might be necessary to rename entities that are accessed with [USE statements](#).
- 22 2 The following example illustrates renaming features of the [USE statement](#).

```

23      MODULE J; REAL JX, JY, JZ; END MODULE J
24      MODULE K
25          USE J, ONLY : KX => JX, KY => JY
26          ! KX and KY are local names to module K
27          REAL KZ          ! KZ is local name to module K
28          REAL JZ          ! JZ is local name to module K
29      END MODULE K
30      PROGRAM RENAME
31          USE J; USE K
32          ! Module J's entity JX is accessible under names JX and KX
33          ! Module J's entity JY is accessible under names JY and KY
34          ! Module K's entity KZ is accessible under name KZ
35          ! Module J's entity JZ and K's entity JZ are different entities
36          ! and cannot be referenced
37          ...
38      END PROGRAM RENAME

```

### 39 C.9.4 Modules with submodules (14.2.3)

- 40 1 Each submodule specifies that it is the child of exactly one parent module or submodule. Therefore, a module
- 41 and all of its descendant submodules stand in a tree-like relationship one to another.

1 2 A separate module procedure that is declared in a module to have public accessibility can be accessed by use  
 2 association even if it is defined in a submodule. No other entity in a submodule can be accessed by use association.  
 3 Each **program unit** that references a module by use association depends on it, and each submodule depends on  
 4 its ancestor module. Therefore, if one changes a separate module procedure body in a submodule but does not  
 5 change its corresponding module procedure interface, a tool for automatic program translation would not need  
 6 to reprocess **program units** that reference the module by use association. This is so even if the tool exploits the  
 7 relative modification times of files as opposed to comparing the result of translating the module to the result of  
 8 a previous translation.

9 3 By constructing taller trees, one can put entities at intermediate levels that are shared by submodules at lower  
 10 levels; changing these entities cannot change the interpretation of anything that is accessible from the module  
 11 by use association. Developers of modules that embody large complicated concepts can exploit this possibility  
 12 to organize components of the concept into submodules, while preserving the privacy of entities that are shared  
 13 by the submodules and that ought not to be exposed to users of the module. Putting these shared entities at an  
 14 intermediate level also prevents cascades of reprocessing and testing if some of them are changed.

15 4 The following example illustrates a module, `color_points`, with a submodule, `color_points_a`, that in turn has  
 16 a submodule, `color_points_b`. Public entities declared within `color_points` can be accessed by use association.  
 17 The submodules `color_points_a` and `color_points_b` can be changed without causing retranslation of **program**  
 18 **units** that reference the module `color_points`.

19 5 The module `color_points` does not have a *module-subprogram-part*, but a *module-subprogram-part* is not pro-  
 20 hibited. The module could be published as definitive specification of the interface, without revealing trade secrets  
 21 contained within `color_points_a` or `color_points_b`. Of course, a similar module without the `module` prefix in  
 22 the interface bodies would serve equally well as documentation – but the procedures would be **external procedures**.  
 23 It would make little difference to the consumer, but the developer would forfeit all of the advantages of modules.

```

24     module color_points
25
26         type color_point
27             private
28                 real :: x, y
29                 integer :: color
30         end type color_point
31
32         interface                ! Interfaces for procedures with separate
33                                 ! bodies in the submodule color_points_a
34             module subroutine color_point_del ( p ) ! Destroy a color_point object
35                 type(color_point), allocatable :: p
36             end subroutine color_point_del
37             ! Distance between two color_point objects
38             real module function color_point_dist ( a, b )
39                 type(color_point), intent(in) :: a, b
40             end function color_point_dist
41             module subroutine color_point_draw ( p ) ! Draw a color_point object
42                 type(color_point), intent(in) :: p
43             end subroutine color_point_draw
44             module subroutine color_point_new ( p ) ! Create a color_point object

```

```

1         type(color_point), allocatable :: p
2     end subroutine color_point_new
3 end interface
4
5     end module color_points

```

6 The only entities within `color_points_a` that can be accessed by use association are the separate module procedures that were declared in `color_points`. If the procedures are changed but their interfaces are not, the interface from `program units` that access them by use association is unchanged. If the module and submodule are in separate files, utilities that examine the time of modification of a file would notice that changes in the module could affect the translation of its submodules or of `program units` that reference the module by use association, but that changes in submodules could not affect the translation of the parent module or `program units` that reference it by use association.

7 The variable `instance_count` in the following example is not accessible by use association of `color_points`, but is accessible within `color_points_a`, and its submodules.

```

15     submodule ( color_points ) color_points_a ! Submodule of color_points
16
17     integer :: instance_count = 0
18
19     interface
20         ! Interface for a procedure with a separate
21         ! body in submodule color_points_b
22     module subroutine inquire_palette ( pt, pal )
23         use palette_stuff      ! palette_stuff, especially submodules
24         ! thereof, can reference color_points by use
25         ! association without causing a circular
26         ! dependence during translation because this
27         ! use is not in the module. Furthermore,
28         ! changes in the module palette_stuff do not
29         ! affect the translation of color_points.
30         type(color_point), intent(in) :: pt
31         type(palette), intent(out) :: pal
32     end subroutine inquire_palette
33
34     end interface
35
36     contains ! Invisible bodies for public separate module procedures
37         ! declared in the module
38
39     module subroutine color_point_del ( p )
40         type(color_point), allocatable :: p
41         instance_count = instance_count - 1
42         deallocate ( p )
43     end subroutine color_point_del
44
45     real module function color_point_dist ( a, b ) result ( dist )
46         type(color_point), intent(in) :: a, b

```

```

1      dist = SQRT ( (b%x - a%x)**2 + (b%y - a%y)**2 )
2  end function color_point_dist
3  module subroutine color_point_new ( p )
4      type(color_point), allocatable :: p
5      instance_count = instance_count + 1
6      allocate ( p )
7  end subroutine color_point_new
8
9  end submodule color_points_a

```

10 8 The subroutine `inquire_palette` is accessible within `color_points_a` because its interface is declared therein. It is not, however, accessible by use association, because its interface is not declared in the module, `color_points`. Since the interface is not declared in the module, changes in the interface cannot affect the translation of [program units](#) that reference the module by use association.

```

14  module palette_stuff
15      type :: palette ; ... ; end type palette
16  contains
17      subroutine test_palette ( p )
18          ! Draw a color wheel using procedures from the color_points module
19          use color_points ! This does not cause a circular dependency because
20                          ! the "use palette_stuff" that is logically within
21                          ! color_points is in the color_points_a submodule.
22          type(palette), intent(in) :: p
23          ...
24      end subroutine test_palette
25  end module palette_stuff
26
27  submodule ( color_points:color_points_a ) color_points_b ! Subsidiary**2 submodule
28
29  contains
30      ! Invisible body for interface declared in the ancestor module
31      module subroutine color_point_draw ( p )
32          use palette_stuff, only: palette
33          type(color_point), intent(in) :: p
34          type(palette) :: MyPalette
35          ...; call inquire_palette ( p, MyPalette ); ...
36      end subroutine color_point_draw
37
38      ! Invisible body for interface declared in the parent submodule
39      module procedure inquire_palette
40          ... implementation of inquire_palette
41      end procedure inquire_palette
42
43      subroutine private_stuff ! not accessible from color_points_a
44          ...
45      end subroutine private_stuff

```

```

1
2     end submodule color_points_b
3
4 9 There is a use palette_stuff in color_points_a, and a use color_points in palette_stuff. The use
5 palette_stuff would cause a circular reference if it appeared in color_points. In this case, it does not cause
6 a circular dependence because it is in a submodule. Submodules cannot be referenced by use association, and
7 therefore what would be a circular appearance of use palette_stuff is not accessed.
8
9     program main
10        use color_points
11        ! "instance_count" and "inquire_palette" are not accessible here
12        ! because they are not declared in the "color_points" module.
13        ! "color_points_a" and "color_points_b" cannot be referenced by
14        ! use association.
15        interface draw                ! just to demonstrate it's possible
16            module procedure color_point_draw
17        end interface
18        type(color_point) :: C_1, C_2
19        real :: RC
20        ...
21        call color_point_new (c_1)    ! body in color_points_a, interface in color_points
22        ...
23        call draw (c_1)               ! body in color_points_b, specific interface
24        ! in color_points, generic interface here.
25        ...
26        rc = color_point_dist (c_1, c_2) ! body in color_points_a, interface in color_points
27        ...
28        call color_point_del (c_1)    ! body in color_points_a, interface in color_points
29        ...
30        end program main

```

10 A multilevel submodule system can be used to package and organize a large and interconnected concept without exposing entities of one subsystem to other subsystems.

11 Consider a **Plasma** module from a Tokamak simulator. A plasma simulation requires attention at least to fluid flow, thermodynamics, and electromagnetism. Fluid flow simulation requires simulation of subsonic, supersonic, and hypersonic flow. This problem decomposition can be reflected in the submodule structure of the **Plasma** module:

| Plasma module      |                      |                      |                   |                            |  |  |
|--------------------|----------------------|----------------------|-------------------|----------------------------|--|--|
| Flow submodule     |                      |                      | Thermal submodule | Electromagnetics submodule |  |  |
| Subsonic submodule | Supersonic submodule | Hypersonic submodule |                   |                            |  |  |
|                    |                      |                      |                   |                            |  |  |

12 Entities can be shared among the **Subsonic**, **Supersonic**, and **Hypersonic** submodules by putting them within the **Flow** submodule. One then need not worry about accidental use of these entities by use association or by the **Thermal** or **Electromagnetics** submodules, or the development of a dependency of correct operation of those

1 subsystems upon the representation of entities of the Flow subsystem as a consequence of maintenance. Since  
2 these entities are not accessible by use association, if any of them are changed, the new values cannot be accessed  
3 in [program units](#) that reference the Plasma module by use association; the answer to the question “where are  
4 these entities used” is therefore confined to the set of descendant submodules of the Flow submodule.

## 5 C.10 Clause 15 notes

### 6 C.10.1 Portability problems with external procedures (15.4.3.5)

7 1 There is a potential portability problem in a [scoping unit](#) that references an [external procedure](#) without explicitly  
8 declaring it to have the [EXTERNAL attribute](#) (8.5.9). On a different processor, the name of that procedure  
9 might be the name of a [nonstandard intrinsic](#) procedure and in such a case the processor would interpret those  
10 procedure references as references to that intrinsic procedure. (On that processor, the program would also be  
11 viewed as not conforming to this document because of the references to the [nonstandard intrinsic](#) procedure.)  
12 Declaration of the [EXTERNAL attribute](#) causes the references to be to the [external procedure](#) regardless of the  
13 availability of an intrinsic procedure with the same name. Note that declaration of the type of a procedure is not  
14 enough to make it [external](#), even if the type is inconsistent with the type of the result of an intrinsic procedure  
15 of the same name.

### 16 C.10.2 Procedures defined by means other than Fortran (15.6.3)

17 1 A processor is not required to provide any means other than Fortran for defining [external procedures](#). Among the  
18 means that might be supported are the machine assembly language, other high level languages, the Fortran lan-  
19 guage extended with nonstandard features, and the Fortran language as supported by another Fortran processor  
20 (for example, a previously existing FORTRAN 77 processor). The means other than Fortran for defining [external](#)  
21 [procedures](#), including any restrictions on the structure for organization of those procedures, are not specified by  
22 this document.

23 2 A Fortran processor might limit its support of procedures defined by means other than Fortran such that these  
24 procedures can affect entities in the Fortran environment only on the same basis as procedures written in Fortran.  
25 For example, it might not support the value of a local variable from being changed by a procedure reference unless  
26 that variable were one of the arguments to the procedure.

### 27 C.10.3 Abstract interfaces and procedure pointer components (15.4, 7.5)

28 1 This is an example of a library module providing lists of callbacks that the user can register and invoke.

```
29     MODULE callback_list_module
30     !
31     ! Type for users to extend with their own data, if they so desire
32     !
33     TYPE callback_data
34     END TYPE
35     !
36     ! Abstract interface for the callback procedures
37     !
38     ABSTRACT INTERFACE
```

```

1      SUBROUTINE callback_procedure(data)
2          IMPORT callback_data
3          CLASS(callback_data),OPTIONAL :: data
4      END SUBROUTINE
5  END INTERFACE
6      !
7      ! The callback list type.
8      !
9      TYPE callback_list
10         PRIVATE
11         CLASS(callback_record),POINTER :: first => NULL()
12     END TYPE
13     !
14     ! Internal: each callback registration creates one of these
15     !
16     TYPE,PRIVATE :: callback_record
17         PROCEDURE(callback_procedure),POINTER,NOPASS :: proc
18         CLASS(callback_record),POINTER :: next
19         CLASS(callback_data),POINTER :: data => NULL();
20     END TYPE
21     PRIVATE invoke,forward_invoke
22 CONTAINS
23     !
24     ! Register a callback procedure with optional data
25     !
26     SUBROUTINE register_callback(list, entry, data)
27         TYPE(callback_list),INTENT(INOUT) :: list
28         PROCEDURE(callback_procedure) :: entry
29         CLASS(callback_data),OPTIONAL :: data
30         TYPE(callback_record),POINTER :: new,last
31         ALLOCATE(new)
32         new%proc => entry
33         IF (PRESENT(data)) ALLOCATE(new%data,SOURCE=data)
34         new%next => list%first
35         list%first => new
36     END SUBROUTINE
37     !
38     ! Internal: Invoke a single callback and destroy its record
39     !
40     SUBROUTINE invoke(callback)
41         TYPE(callback_record),POINTER :: callback
42         IF (ASSOCIATED(callback%data) THEN
43             CALL callback%proc(list%first%data)
44             DEALLOCATE(callback%data)
45         ELSE
46             CALL callback%proc

```

```
1         END IF
2         DEALLOCATE(callback)
3     END SUBROUTINE
4     !
5     ! Call the procedures in reverse order of registration
6     !
7     SUBROUTINE invoke_callback_reverse(list)
8         TYPE(callback_list),INTENT(INOUT) :: list
9         TYPE(callback_record),POINTER :: next,current
10        current => list%first
11        NULLIFY(list%first)
12        DO WHILE (ASSOCIATED(current))
13            next => current%next
14            CALL invoke(current)
15            current => next
16        END DO
17    END SUBROUTINE
18    !
19    ! Internal: Forward mode invocation
20    !
21    SUBROUTINE forward_invoke(callback)
22        IF (ASSOCIATED(callback%next)) CALL forward_invoke(callback%next)
23        CALL invoke(callback)
24    END SUBROUTINE
25    !
26    ! Call the procedures in forward order of registration
27    !
28    SUBROUTINE invoke_callback_forward(list)
29        TYPE(callback_list),INTENT(INOUT) :: list
30        IF (ASSOCIATED(list%first)) CALL forward_invoke(list%first)
31    END SUBROUTINE
32    END
```

#### 33 C.10.4 Pointers and targets as arguments (15.5.2.4, 15.5.2.6, 15.5.2.7)

- 34 1 If a dummy argument is declared to be a [pointer](#), the corresponding [actual argument](#) could be a [pointer](#) or could  
35 be a nonpointer variable or procedure. Consider the two cases separately.
- 36 *Case (i):* The [actual argument](#) is a pointer. When procedure execution commences the [pointer association](#)  
37 status of the [dummy argument](#) becomes the same as that of the [actual argument](#). If the [pointer](#)  
38 [association](#) status of the [dummy argument](#) is changed, the [pointer association](#) status of the [actual](#)  
39 [argument](#) changes in the same way.
- 40 *Case (ii):* The [actual argument](#) is not a pointer. This only occurs when the [actual argument](#) has the [TARGET](#)  
41 [attribute](#) or is a procedure, and the [dummy argument](#) has the [INTENT \(IN\)](#) attribute. The [dummy](#)  
42 [argument](#) becomes pointer associated with the [actual argument](#).
- 43 2 When execution of a procedure completes, any pointer that remains defined and that is associated with a dummy  
44 argument that has the [TARGET attribute](#) and is either a scalar or an [assumed-shape array](#), remains associated



1 with the corresponding [actual argument](#) if the [actual argument](#) has the [TARGET attribute](#) and is not an array  
2 section with a [vector subscript](#).

3 For example, consider:

```

4     REAL, POINTER      :: PBEST
5     REAL, TARGET      :: B (10000)
6     CALL BEST (PBEST, B)          ! Upon return PBEST is associated
7     ...                          ! with the 'best' element of B
8     CONTAINS
9     SUBROUTINE BEST (P, A)
10    REAL, POINTER, INTENT (OUT)  :: P
11    REAL, TARGET, INTENT (IN)   :: A (:)
12    ...                          ! Find the 'best' element A(I)
13    P => A (I)
14    RETURN
15    END SUBROUTINE BEST
16    END

```

17 When procedure BEST completes, the pointer PBEST is associated with an element of B.

18 4 An [actual argument](#) without the [TARGET attribute](#) can become associated with a dummy argument with the  
19 [TARGET attribute](#). This enables a pointer to become associated with the dummy argument during execution of  
20 the procedure that contains the dummy argument. For example:

```

21    INTEGER LARGE(100,100)
22    CALL SUB (LARGE)
23    ...
24    CALL SUB ()
25    CONTAINS
26    SUBROUTINE SUB(ARG)
27    INTEGER, TARGET, OPTIONAL :: ARG(100,100)
28    INTEGER, POINTER, DIMENSION(:, :) :: PARG
29    IF (PRESENT(ARG)) THEN
30    PARG => ARG
31    ELSE
32    ALLOCATE (PARG(100,100))
33    PARG = 0
34    ENDIF
35    ... ! Code with lots of references to PARG
36    IF (.NOT. PRESENT(ARG)) DEALLOCATE(PARG)
37    END SUBROUTINE SUB
38    END

```

39 Within subroutine SUB the pointer PARG is either associated with the dummy argument ARG or it is associated  
40 with an allocated target. The bulk of the code can reference PARG without further calls to the intrinsic function  
41 [PRESENT](#).

42 5 If a nonpointer dummy argument has the [TARGET attribute](#) and the corresponding [actual argument](#) does not,  
43 any pointers that become associated with the dummy argument, and therefore with the [actual argument](#), during

1 execution of the procedure, become undefined when execution of the procedure completes.

## 2 **C.10.5 Polymorphic Argument Association (15.5.2.9)**

3 1 The following example illustrates the polymorphic [argument association](#) rules using the derived types defined in  
4 Note 7.55.

```

5     TYPE(POINT) :: T2
6     TYPE(COLOR_POINT) :: T3
7     CLASS(POINT) :: P2
8     CLASS(COLOR_POINT) :: P3
9     ! Dummy argument is polymorphic and actual argument is of fixed type
10    SUBROUTINE SUB2 ( X2 ); CLASS(POINT) :: X2; ...
11    SUBROUTINE SUB3 ( X3 ); CLASS(COLOR_POINT) :: X3; ...
12
13    CALL SUB2 ( T2 ) ! Valid -- The declared type of T2 is the same as the
14                    !           declared type of X2.
15    CALL SUB2 ( T3 ) ! Valid -- The declared type of T3 is extended from
16                    !           the declared type of X2.
17    CALL SUB3 ( T2 ) ! Invalid -- The declared type of T2 is neither the
18                    !           same as nor extended from the declared type
19                    !           type of X3.
20    CALL SUB3 ( T3 ) ! Valid -- The declared type of T3 is the same as the
21                    !           declared type of X3.
22    ! Actual argument is polymorphic and dummy argument is of fixed type
23    SUBROUTINE TUB2 ( D2 ); TYPE(POINT) :: D2; ...
24    SUBROUTINE TUB3 ( D3 ); TYPE(COLOR_POINT) :: D3; ...
25
26    CALL TUB2 ( P2 ) ! Valid -- The declared type of P2 is the same as the
27                    !           declared type of D2.
28    CALL TUB2 ( P3 ) ! Invalid -- The declared type of P3 differs from the
29                    !           declared type of D2.
30    CALL TUB2 ( P3%POINT ) ! Valid alternative to the above
31    CALL TUB3 ( P2 ) ! Invalid -- The declared type of P2 differs from the
32                    !           declared type of D3.
33    SELECT TYPE ( P2 ) ! Valid conditional alternative to the above
34    CLASS IS ( COLOR_POINT ) ! Works if the dynamic type of P2 is the same
35        CALL TUB3 ( P2 )      ! as the declared type of D3, or a type
36                                ! extended therefrom.
37    CLASS DEFAULT
38                                ! Cannot work if not.
39    END SELECT
40    CALL TUB3 ( P3 ) ! Valid -- The declared type of P3 is the same as the
41                    !           declared type of D3.
42    ! Both the actual and dummy arguments are of polymorphic type.
43    CALL SUB2 ( P2 ) ! Valid -- The declared type of P2 is the same as the
44                    !           declared type of X2.
```

```

1      CALL SUB2 ( P3 ) ! Valid -- The declared type of P3 is extended from
2          !           the declared type of X2.
3      CALL SUB3 ( P2 ) ! Invalid -- The declared type of P2 is neither the
4          !           same as nor extended from the declared
5          !           type of X3.
6      SELECT TYPE ( P2 ) ! Valid conditional alternative to the above
7      CLASS IS ( COLOR_POINT ) ! Works if the dynamic type of P2 is the
8          CALL SUB3 ( P2 )      ! same as the declared type of X3, or a
9          ! type extended therefrom.
10     CLASS DEFAULT
11         ! Cannot work if not.
12     END SELECT
13     CALL SUB3 ( P3 ) ! Valid -- The declared type of P3 is the same as the
14         !           declared type of X3.

```

## 15 C.10.6 Rules ensuring unambiguous generics (15.4.3.4.5)

- 16 1 The rules in 15.4.3.4.5 are intended to ensure
- 17 • that it is possible to reference each specific procedure or binding in the generic collection,
  - 18 • that for any valid generic procedure reference, the determination of the specific procedure referenced is unambiguous, and
  - 19 • that the determination of the specific procedure or binding referenced can be made before execution of the program begins (during compilation).
- 22 2 Interfaces of specific procedures or bindings are distinguished by fixed properties of their arguments, specifically type, kind type parameters, [rank](#), and whether the dummy argument has the [POINTER](#) or [ALLOCATABLE](#) attribute. A valid reference to one procedure in a generic collection will differ from another because it has an argument that the other cannot accept, because it is missing an argument that the other requires, or because one of these fixed properties is different.
- 27 3 Although the declared type of a data entity is a fixed property, polymorphic variables allow for a limited degree of type mismatch between dummy arguments and [actual arguments](#), so the requirement for distinguishing two dummy arguments is type incompatibility, not merely different types. (This is illustrated in the BAD6 example later in this note.)
- 31 4 That same limited type mismatch means that two dummy arguments that are not type incompatible can be distinguished on the basis of the values of the kind type parameters they have in common; if one of them has a kind type parameter that the other does not, that is irrelevant in distinguishing them.
- 34 5 [Rank](#) is a fixed property, but some forms of array dummy arguments allow [rank](#) mismatches when a procedure is referenced by its specific name. In order to allow [rank](#) to always be usable in distinguishing generics, such [rank](#) mismatches are disallowed for those arguments when the procedure is referenced as part of a generic. Additionally, the fact that elemental procedures can accept array arguments is not taken into account when applying these rules, so apparent ambiguity between elemental and nonelemental procedures is possible; in such cases, the reference is interpreted as being to the nonelemental procedure.
- 40 6 For procedures referenced as operators or defined-assignment, syntactically distinguished arguments are mapped

1 to specific positions in the argument list, so the rule for distinguishing such procedures is that it be possible to  
2 distinguish the arguments at one of the argument positions.

3 7 For [defined input/output](#) procedures, only the `dtv` argument corresponds to something explicitly written in the  
4 program, so it is the `dtv` that is required to be distinguished. Because `dtv` arguments are required to be scalar,  
5 they cannot differ in [rank](#). Thus this rule effectively involves only type and kind type parameters.

6 8 For generic procedure names, the rules are more complicated because optional arguments can be omitted and  
7 because arguments can be specified either positionally or by name.

8 9 In the special case of type-bound procedures with [passed-object dummy arguments](#), the passed-object argument  
9 is syntactically distinguished in the reference, so rule (3) in [15.4.3.4.5](#) can be applied. The type of passed-object  
10 arguments is constrained in ways that prevent passed-object arguments in the same [scoping unit](#) from being type  
11 incompatible. Thus this rule effectively involves only kind type parameters and [rank](#).

12 10 The primary means of distinguishing named generics is rule (4). The most common application of that rule is a  
13 single argument satisfying both (4a) and (4b):

```
14         INTERFACE GOOD1
15             FUNCTION F1A(X)
16                 REAL :: F1A,X
17             END FUNCTION F1A
18             FUNCTION F1B(X)
19                 INTEGER :: F1B,X
20             END FUNCTION F1B
21         END INTERFACE GOOD1
```

22 11 Whether one writes `GOOD1(1.0)` or `GOOD1(X=1.0)`, the reference is to `F1A` because `F1B` would require an integer  
23 argument whereas these references provide the real constant 1.0.

24 12 This example and those that follow are expressed using interface bodies, with type as the distinguishing property.  
25 This was done to make it easier to write and describe the examples. The principles being illustrated are equally  
26 applicable when the procedures get their explicit interfaces in some other way or when kind type parameters or  
27 [rank](#) are the distinguishing property.

28 13 Another common variant is the argument that satisfies (4a) and (4b) by being required in one specific and  
29 completely missing in the other:

```
30         INTERFACE GOOD2
31             FUNCTION F2A(X)
32                 REAL :: F2A,X
33             END FUNCTION F2A
34             FUNCTION F2B(X,Y)
35                 COMPLEX :: F2B
36                 REAL :: X,Y
37             END FUNCTION F2B
38         END INTERFACE GOOD2
```

39 14 Whether one writes `GOOD2(0.0,1.0)`, `GOOD2(0.0,Y=1.0)`, or `GOOD2(Y=1.0,X=0.0)`, the reference is to `F2B`,  
40 because `F2A` has no argument in the second position or with the name `Y`. This approach is used as an alternative

1 to optional arguments when one wants a function to have different result type, kind type parameters, or [rank](#),  
 2 depending on whether the argument is present. In many of the intrinsic functions, the DIM argument works this  
 3 way.

4 15 It is possible to construct cases where different arguments are used to distinguish positionally and by name:

```
5      INTERFACE GOOD3
6          SUBROUTINE S3A(W,X,Y,Z)
7              REAL :: W,Y
8              INTEGER :: X,Z
9          END SUBROUTINE S3A
10         SUBROUTINE S3B(X,W,Z,Y)
11             REAL :: W,Z
12             INTEGER :: X,Y
13         END SUBROUTINE S3B
14     END INTERFACE GOOD3
```

15 16 If one writes GOOD3(1.0,2,3.0,4) to reference S3A, then the third and fourth arguments are consistent with a  
 16 reference to S3B, but the first and second are not. If one switches to writing the first two arguments as keyword  
 17 arguments in order for them to be consistent with a reference to S3B, the latter two arguments will also need  
 18 to be written as keyword arguments, GOOD3(X=2,W=1.0,Z=4,Y=3.0), and the named arguments Y and Z are  
 19 distinguished.

20 17 The ordering requirement in rule (4) is critical:

```
21     INTERFACE BAD4 ! this interface is invalid !
22         SUBROUTINE S4A(W,X,Y,Z)
23             REAL :: W,Y
24             INTEGER :: X,Z
25         END SUBROUTINE S4A
26         SUBROUTINE S4B(X,W,Z,Y)
27             REAL :: X,Y
28             INTEGER :: W,Z
29         END SUBROUTINE S4B
30     END INTERFACE BAD4
```

31 18 In this example, the positionally distinguished arguments are Y and Z, and it is W and X that are distinguished by  
 32 name. In this order it is possible to write BAD4(1.0,2,Y=3.0,Z=4), which is a valid reference for both S4A and  
 33 S4B.

34 19 Rule (1) can be used to distinguish some cases that are not covered by rule (4):

```
35     INTERFACE GOOD5
36         SUBROUTINE S5A(X)
37             REAL :: X
38         END SUBROUTINE S5A
39         SUBROUTINE S5B(Y,X)
40             REAL :: Y,X
41         END SUBROUTINE S5B
```

1           END INTERFACE GOOD5

2 20 In attempting to apply rule (4), position 2 and name Y are distinguished, but they are in the wrong order, just like  
3 the BAD4 example. However, when we try to construct a similarly ambiguous reference, we get GOOD5(1.0,X=2.0),  
4 which can't be a reference to S5A because it would be attempting to associate two different **actual arguments**  
5 with the dummy argument X. Rule (4) catches this case by recognizing that S5B requires two real arguments, and  
6 S5A cannot possibly accept more than one.

7 21 The application of rule (1) becomes more complicated when **extensible types** are involved. If FRUIT is an **extensible**  
8 **type**, PEAR and APPLE are **extensions** of FRUIT, and BOSC is an **extension** of PEAR, then

```
9           INTERFACE BAD6 ! this interface is invalid !
10           SUBROUTINE S6A(X,Y)
11           CLASS(PEAR) :: X,Y
12           END SUBROUTINE S6A
13           SUBROUTINE S6B(X,Y)
14           CLASS(FRUIT) :: X
15           CLASS(BOSC) :: Y
16           END SUBROUTINE S6B
17           END INTERFACE BAD6
```

18 might, at first glance, seem distinguishable this way, but because of the limited type mismatching allowed,  
19 BAD6(A\_PEAR,A\_BOSC) is a valid reference to both S6A and S6B.

20 22 It is important to try rule (1) for each type that appears:

```
21           INTERFACE GOOD7
22           SUBROUTINE S7A(X,Y,Z)
23           CLASS(PEAR) :: X,Y,Z
24           END SUBROUTINE S7A
25           SUBROUTINE S7B(X,Z,W)
26           CLASS(FRUIT) :: X
27           CLASS(BOSC) :: Z
28           CLASS(APPLE),OPTIONAL :: W
29           END SUBROUTINE S7B
30           END INTERFACE GOOD7
```

31 23 Looking at the most general type, S7A has a minimum and maximum of 3 FRUIT arguments, while S7B has a  
32 minimum of 2 and a maximum of three. Looking at the most specific, S7A has a minimum of 0 and a maximum  
33 of 3 BOSC arguments, while S7B has a minimum of 1 and a maximum of 2. However, when we look at the  
34 intermediate, S7A has a minimum and maximum of 3 PEAR arguments, while S7B has a minimum of 1 and a  
35 maximum of 2. Because S7A's minimum exceeds S7B's maximum, they can be distinguished.

36 24 In identifying the minimum number of arguments with a particular set of properties, we exclude optional argu-  
37 ments and test TKR compatibility, so the corresponding **actual arguments** are required to have those properties.  
38 In identifying the maximum number of arguments with those properties, we include the optional arguments and  
39 test not distinguishable, so we include **actual arguments** which could have those properties but are not required  
40 to have them.

1 25 These rules are sufficient to ensure that references to procedures that meet them are unambiguous, but there  
2 remain examples that fail to meet these rules but which can be shown to be unambiguous:

```
3 26     INTERFACE BAD8 ! this interface is invalid !
4         ! despite the fact that it is unambiguous !
5         SUBROUTINE S8A(X,Y,Z)
6             REAL,OPTIONAL :: X
7             INTEGER :: Y
8             REAL :: Z
9         END SUBROUTINE S8A
10        SUBROUTINE S8B(X,Z,Y)
11            INTEGER,OPTIONAL :: X
12            INTEGER :: Z
13            REAL :: Y
14        END SUBROUTINE S8B
15    END INTERFACE BAD8
```

16 27 This interface fails rule (4) because there are no required arguments that can be distinguished from the positionally  
17 corresponding argument, but in order for the mismatch of the optional arguments not to be relevant, the later  
18 arguments need to be specified as keyword arguments, so distinguishing by name does the trick. This interface is  
19 nevertheless invalid so a standard-conforming Fortran processor is not required to do such reasoning. The rules  
20 to cover all cases are too complicated to be useful.

21 28 The real data objects that would be valid arguments for S9A are entirely disjoint from procedures that are valid  
22 arguments to S9B and S9C, and the procedures that valid arguments for S9B are disjoint from the procedures  
23 that are valid arguments to S9C because the former are required to accept real arguments and the latter in-  
24 teger arguments. Again, this interface is invalid, so a standard-conforming Fortran processor need not examine  
25 such properties when deciding whether a generic collection is valid. Again, the rules to cover all cases are too  
26 complicated to be useful.

27 29 If one dummy argument has the [POINTER attribute](#) and a corresponding argument in the other interface body  
28 has the [ALLOCATABLE attribute](#) the [generic interface](#) is not ambiguous. If one dummy argument has either the  
29 [POINTER](#) or [ALLOCATABLE](#) attribute and a corresponding argument in the other interface body has neither  
30 attribute, the [generic interface](#) might be ambiguous.

## 31 C.11 Clause 16 notes

### 32 C.11.1 Atomic memory consistency

#### 33 C.11.1.1 Relaxed memory model

34 1 Parallel programs sometimes have apparently impossible behavior because data transfers and other messages can  
35 be delayed, reordered and even repeated, by hardware, communication software, and caching and other forms  
36 of optimization. Requiring processors to deliver globally consistent behavior is incompatible with performance  
37 on many systems. This document specifies that all ordered actions will be consistent (5.3.5 and 11.6), but all  
38 consistency between unordered segments is deliberately left processor dependent. Depending on the hardware,  
39 this can be observed even when only two images and one mechanism are involved.

### 1 C.11.1.2 Examples with atomic operations

2 1 When variables are being referenced (atomically) from segments that are unordered with respect to the segment  
3 that is atomically defining or redefining the variables, the results are processor dependent. This supports use  
4 of so-called “relaxed memory model” architectures, which can enable more efficient execution on some hardware  
5 implementations.

6 2 The following examples assume these declarations:

```
7     MODULE EXAMPLE
8         USE, INTRINSIC :: ISO_FORTRAN_ENV
9         INTEGER(ATOMIC_INT_KIND) :: X [*] = 0, Y [*] = 0, TMP
```

#### 10 3 Example 1

11 With X [*j*] and Y [*j*] still in their initial state (both zero), image *j* executes the following sequence of statements:

```
12     CALL ATOMIC_DEFINE (X, 1)
13     CALL ATOMIC_DEFINE (Y, 1)
```

14 and a different image, *k*, executes the following sequence of statements:

```
15     DO
16         CALL ATOMIC_REF (TMP, Y [j])
17         IF (TMP==1) EXIT
18     END DO
19     CALL ATOMIC_REF (TMP, X [j])
20     PRINT *, TMP
```

21 4 The final value of TMP on image *k* could be either 0 or 1. That is, even though image *j* thinks that it defined X  
22 [*j*] before it defined Y [*j*], this ordering is not guaranteed to be observed on image *k*. There are many aspects of  
23 hardware and software implementation that can cause this effect, but conceptually this example can be thought  
24 of as the change in the value of Y propagating faster through the inter-image connections than the change in the  
25 value of X.

26 5 Even if image *j* executed the sequence

```
27     CALL ATOMIC_DEFINE (X, 1)
28     SYNC MEMORY
29     CALL ATOMIC_DEFINE (Y, 1)
```

30 the same effect could be seen. That is because even though X and Y are defined in ordered segments, the  
31 references from image *k* are both from a segment that is unordered with respect to image *j*.

32 6 Only if the reference on image *k* to Y [*j*] is in a segment that is ordered after the segment on image *j* that defined  
33 Y, will TMP be guaranteed to have the value 1.

#### 34 7 Example 2:

35 With the initial state of X and Y on image *j* (i.e. X [*j*] and Y [*j*]) still being zero, execution of

```
36     CALL ATOMIC_REF (TMP, X [j])
37     CALL ATOMIC_DEFINE (Y [j], 1)
38     PRINT *, TMP
```



1 on image  $k_1$ , and execution of

```
2     CALL ATOMIC_REF (TMP, Y [j])
3     CALL ATOMIC_DEFINE (X [j], 1)
4     PRINT *, TMP
```

5 on image  $k_2$ , in unordered segments, might print the value 1 both times.

6 8 This can happen by such mechanisms as “load buffering”; one might imagine that what is happening is that  
7 the definitions (`ATOMIC_DEFINE`) are overtaking the references (`ATOMIC_REF`). On some processors it is  
8 possible that insertion of `SYNC MEMORY` statements between the calls to `ATOMIC_REF` and `ATOMIC_`  
9 `DEFINE` might be sufficient to make the output print the value 1 at most one time (or even exactly one time),  
10 but this is still processor dependent unless the `SYNC MEMORY` statement executions cause the relevant segments  
11 on images  $k_1$  and  $k_2$  to be ordered.

### 12 9 Example 3:

13 Because there are no segment boundaries implied by collective subroutines, with the initial state as before,  
14 execution of

```
15     IF (THIS_IMAGE ()==1) THEN
16         CALL ATOMIC_DEFINE (X [3], 23)
17         Y = 42
18     END IF
19     CALL CO_BROADCAST (Y, 1)
20     IF (THIS_IMAGE ()==2) THEN
21         CALL ATOMIC_REF (TMP, X [3])
22         PRINT *, Y, TMP
23     END IF
```

24 could print the values 42 and 0.

### 25 10 Example 4:

26 Assuming the declarations

```
27     INTEGER (ATOMIC_INT_KIND) :: X [*] = 0, Z = 0
```

28 the statements

```
29     CALL ATOMIC_ADD (X [1], 1)           ! (A)
30     IF (THIS_IMAGE() == 2) THEN
31         wait: DO
32             CALL ATOMIC_REF (Z, X [1])   ! (B)
33             IF (Z == NUM_IMAGES ()) EXIT wait
34         END DO wait                       ! (C)
35     END IF
```

36 will execute the “wait” loop on image 2 until all images have completed statement (A). The updates of X [1] are  
37 performed by each image in the same manner, but in an arbitrary order. Because the result from the complete  
38 set of updates will eventually become visible by execution of statement (B) for some loop iteration on image 2,  
39 the termination condition is guaranteed to be eventually fulfilled, provided that no image failure occurs, every

1 image executes the above code, and no other code is executed in an unordered segment that performs an update  
 2 to X [1]. Furthermore, if two **SYNC MEMORY** statements are inserted in the above code before statement (A)  
 3 and after statement (C), respectively, the segment started by the second **SYNC MEMORY** on image 2 is ordered  
 4 after the segments on all images that end with the first **SYNC MEMORY**.

### 5 **C.11.2 EVENT\_QUERY example**

- 6 1 The following example illustrates the use of events via a program in which image 1 acts as the master image,  
 7 distributing work items to the other images. Only one work item at a time can be active on a worker image, and  
 8 each deals with the result (e.g. via input/output) without directly feeding data back to the master image.
- 9 2 Because the work items are not expected to be balanced, the master keeps cycling through all images to find one  
 10 that is waiting for work.
- 11 3 An event is posted by each worker to indicate that it has completed its work item. Since the corresponding  
 12 variables are needed only on the master, we place them in an allocatable array component of a coarray. An event  
 13 on each worker is needed for the master to post the fact that it has made a work item available for it.

#### 14 **Example code:**

```

15     PROGRAM work_share
16         USE, INTRINSIC :: ISO_FORTRAN_ENV, ONLY: EVENT_TYPE
17         USE :: mod_work, ONLY:
18             work,           & ! Module that creates work items
19             create_work_item, & ! Function that creates work item
20             process_item,   & ! Function that processes an item
21             work_done       ! Logical function that returns true
22                             ! if all work has been done.
23
24     TYPE :: worker_type
25         TYPE (EVENT_TYPE), ALLOCATABLE :: free (:)
26     END TYPE
27     TYPE (EVENT_TYPE) :: submit [*] ! Post when work ready for a worker
28     TYPE (worker_type) :: worker [*] ! Post when worker is free
29     TYPE (work) :: work_item [*] ! Holds the data for a work item
30     INTEGER :: count, i, nbusy [*]
31
32     IF (THIS_IMAGE ()==1) THEN
33         ! Get started
34         ALLOCATE (worker%free (2:NUM_IMAGES ()))
35         nbusy = 0 ! This holds the number of workers working
36         DO i = 2, NUM_IMAGES () ! Start the workers working
37             IF (work_done ()) EXIT
38             nbusy = nbusy + 1
39             work_item [i] = create_work_item ()
40             EVENT POST (submit [i])
41         END DO
42         ! Main work distribution loop

```

```

1     master: DO
2         image: DO i = 2, NUM_IMAGES ()
3             CALL EVENT_QUERY (worker%free (i), count)
4             IF (count==0) CYCLE image ! Worker is not free
5             EVENT WAIT (worker%free (i))
6             nbusy = nbusy - 1
7             IF (work_done ()) CYCLE
8             nbusy = nbusy + 1
9             work_item [i] = create_work_item ()
10            EVENT POST (submit [i])
11        END DO image
12        IF (nbusy==0) THEN
13            ! All done. Exit on all images.
14            DO i = 2, NUM_IMAGES ()
15                EVENT POST (submit [i])
16            END DO
17            EXIT master
18        END IF
19    END DO master
20    ELSE
21        ! Work processing loop
22    worker: DO
23        EVENT WAIT (submit)
24        IF (nbusy[1] == 0) EXIT
25        CALL process_item (work_item)
26        EVENT POST (worker [1]%free (THIS_IMAGE ()))
27    END DO worker
28    END IF
29    END PROGRAM work_share

```

### 30 C.11.3 Collective subroutine examples

31 1 The following example computes a dot product of two scalar coarrays using `CO_SUM` to store the result in a  
32 noncoarray scalar variable.

```

33     SUBROUTINE codot (x, y, x_dot_y)
34         REAL :: x [*], y [*], x_dot_y
35         x_dot_y = x*y
36         CALL CO_SUM (x_dot_y)
37     END SUBROUTINE codot

```

38 2 The function below demonstrates passing a noncoarray dummy argument to `CO_MAX`. The function uses `CO_-`  
39 `MAX` to find the maximum value of the dummy argument across all images. Then the function flags all images  
40 that hold values matching the maximum. The function then returns the maximum image index for an image that  
41 holds the maximum value.

```

42     FUNCTION find_max (j) RESULT (j_max_location)

```

```

1      INTEGER, INTENT (IN) :: j
2      INTEGER j_max, j_max_location
3      j_max = j
4      CALL CO_MAX (j_max)
5      ! Flag images that hold the maximum j.
6      IF (j==j_max) THEN
7          j_max_location = THIS_IMAGE ()
8      ELSE
9          j_max_location = 0
10     END IF
11     ! Return highest image index associated with a maximal j.
12     CALL CO_MAX(j_max_location)
13 END FUNCTION find_max

```

## 14 C.12 Clause 18 notes

### 15 C.12.1 Runtime environments (18.1)

- 16 1 This document allows programs to contain procedures defined by means other than Fortran. That raises the  
17 issues of initialization of and interaction between the runtime environments involved.
- 18 2 Implementations are free to solve these issues as they see fit, provided that
- 19 • heap allocation/deallocation (e.g., (DE)ALLOCATE in a Fortran subprogram and malloc/free in a C func-  
20 tion) can be performed without interference,
  - 21 • input/output to and from [external files](#) can be performed without interference, as long as procedures defined  
22 by different means do not do input/output with the same [external file](#),
  - 23 • input/output preconnections exist as required by the respective standards, and
  - 24 • initialized data are initialized according to the respective standards.

### 25 C.12.2 Example of Fortran calling C (18.3)

#### 26 C Function Prototype:

```
27     int C_Library_Function(void* sendbuf, int sendcount, int *recvcounts);
```

#### 28 Fortran Module:

```

29     MODULE CLIBFUN_INTERFACE
30     INTERFACE
31         INTEGER (C_INT) FUNCTION C_LIBRARY_FUNCTION (SENDBUF, SENDCOUNT, RECVCOUNTS) &
32             BIND(C, NAME='C_Library_Function')
33         USE, INTRINSIC :: ISO_C_BINDING
34         IMPLICIT NONE
35         TYPE (C_PTR), VALUE :: SENDBUF
36         INTEGER (C_INT), VALUE :: SENDCOUNT
37         INTEGER (C_INT) :: RECVCOUNTS(*)
38     END FUNCTION C_LIBRARY_FUNCTION

```

```

1           END INTERFACE
2     END MODULE CLIBFUN_INTERFACE

```

3 1 The module CLIBFUN\_INTERFACE contains the declaration of the Fortran dummy arguments, which correspond to the C formal parameters. The **NAME=** is used in the **BIND attribute** in order to handle the case-sensitive name change between Fortran and C from “c\_library\_function” to “C\_Library\_Function”.

6 2 The first C formal parameter is the pointer to void **sendbuf**, which corresponds to the Fortran dummy argument SENDBUF, which has the type **C\_PTR** and the **VALUE attribute**.

8 3 The second C formal parameter is the int **sendcount**, which corresponds to the Fortran dummy argument SENDCOUNT, which has the type **INTEGER (C\_INT)** and the **VALUE attribute**.

10 4 The third C formal parameter is the pointer to int **recvcounts**, which corresponds to the Fortran dummy argument RECVCOUNTS, which is an **assumed-size array** of type **INTEGER (C\_INT)**.

12 5 This example show how **C\_Library\_Function** might be referenced in a Fortran **program unit**:

```

13         USE, INTRINSIC :: ISO_C_BINDING, ONLY: C_INT, C_FLOAT, C_LOC
14         USE CLIBFUN_INTERFACE
15         ...
16         REAL (C_FLOAT), TARGET :: SEND(100)
17         INTEGER (C_INT)          :: SENDCOUNT, RET
18         INTEGER (C_INT), ALLOCATABLE :: RECVCOUNTS(:)
19         ...
20         ALLOCATE( RECVCOUNTS(100) )
21         ...
22         RET = C_LIBRARY_FUNCTION(C_LOC(SEND), SENDCOUNT, RECVCOUNTS)
23         ...

```

24 6 The first Fortran actual argument is a reference to the function **C\_LOC** which returns the value of the C address of its argument, SEND. This value becomes the value of the first formal parameter, the pointer **sendbuf**, in **C\_Library\_Function**.

27 7 The second Fortran actual argument is SENDCOUNT of type **INTEGER (C\_INT)**. Its value becomes the initial value of the second formal parameter, the int **sendcount**, in **C\_Library\_Function**.

29 8 The third Fortran actual argument is the allocatable array RECVCOUNTS of type **INTEGER (C\_INT)**. The base C address of this array becomes the value of the third formal parameter, the pointer **recvcounts**, in **C\_Library\_Function**. Note that interoperability is based on the characteristics of the dummy arguments in the specified interface and not on those of the actual arguments. Thus, the fact that the actual argument is allocatable is not relevant here.

### 34 C.12.3 Example of C calling Fortran (18.3)

35 **Fortran Code:**

```

36         SUBROUTINE SIMULATION(ALPHA, BETA, GAMMA, DELTA, ARRAYS) BIND(C)
37         USE, INTRINSIC :: ISO_C_BINDING

```

```

1      IMPLICIT NONE
2      INTEGER (C_LONG), VALUE           :: ALPHA
3      REAL (C_DOUBLE), INTENT(INOUT)    :: BETA
4      INTEGER (C_LONG), INTENT(OUT)     :: GAMMA
5      REAL (C_DOUBLE), DIMENSION(*), INTENT(IN) :: DELTA
6      TYPE, BIND(C) :: PASS
7          INTEGER (C_INT) :: LENC, LENF
8          TYPE (C_PTR)    :: C, F
9      END TYPE PASS
10     TYPE (PASS), INTENT(INOUT) :: ARRAYS
11     REAL (C_FLOAT), ALLOCATABLE, TARGET, SAVE :: ETA(:)
12     REAL (C_FLOAT), POINTER :: C_ARRAY(:)
13     ...
14     ! Associate C_ARRAY with an array allocated in C
15     CALL C_F_POINTER (ARRAYS%C, C_ARRAY, [ ARRAYS%LENC ])
16     ...
17     ! Allocate an array and make it available in C
18     ARRAYS%LENF = 100
19     ALLOCATE (ETA(ARRAYS%LENF))
20     ARRAYS%F = C_LOC(ETA)
21     ...
22     END SUBROUTINE SIMULATION

```

### 23 C Struct Declaration:

```

24     struct pass {
25         int lenc, lenf;
26         float *c, *f;
27     };

```

### 28 C Function Prototype:

```

29     void simulation(long alpha, double *beta, long *gamma, double delta[],
30                    struct pass *arrays);

```

### 31 C Calling Sequence:

```

32     simulation(alpha, beta, gamma, delta, arrays);

```

33 1 The above-listed Fortran code specifies a subroutine SIMULATION. This subroutine corresponds to the C void  
34 function simulation.

35 2 The Fortran subroutine references the intrinsic module ISO\_C\_BINDING.

36 3 The first Fortran dummy argument of the subroutine is ALPHA, which has the type INTEGER(C\_LONG) and  
37 the [VALUE attribute](#). This dummy argument corresponds to the C formal parameter alpha, which is a long.  
38 The C actual argument is also a long.

- 1 4 The second Fortran dummy argument of the subroutine is BETA, which has the type REAL(C\_DOUBLE) and  
 2 the **INTENT (INOUT) attribute**. This dummy argument corresponds to the C formal parameter `beta`, which is  
 3 a pointer to double. An address is passed as the C actual argument.
- 4 5 The third Fortran dummy argument of the subroutine is GAMMA, which has the type INTEGER(C\_LONG)  
 5 and the **INTENT (OUT) attribute**. This dummy argument corresponds to the C formal parameter `gamma`, which  
 6 is a pointer to long. An address is passed as the C actual argument.
- 7 6 The fourth Fortran dummy argument is the **assumed-size array** DELTA, which has the type REAL (C\_DOUBLE)  
 8 and the **INTENT (IN) attribute**. This dummy argument corresponds to the C formal parameter `delta`, which is  
 9 a double array. The C actual argument is also a double array.
- 10 7 The fifth Fortran dummy argument is ARRAYS, which is a structure for accessing an array allocated in C and  
 11 an array allocated in Fortran. The lengths of these arrays are held in the components LENC and LENF; their **C**  
 12 **addresses** are held in components C and F.

### 13 C.12.4 Example of calling C functions with noninteroperable data (18.10)

- 14 1 Many Fortran processors support 16-byte real numbers, which might not be supported by the C processor.  
 15 Assume a Fortran programmer wants to use a C procedure from a message passing library for an array of these  
 16 reals. The C prototype of this procedure is

```
17     void ProcessBuffer(void *buffer, int n_bytes);
```

18 with the corresponding Fortran interface

```
19     USE, INTRINSIC :: ISO_C_BINDING
20     INTERFACE
21         SUBROUTINE PROCESS_BUFFER(BUFFER,N_BYTES) BIND(C,NAME="ProcessBuffer")
22             IMPORT :: C_PTR, C_INT
23             TYPE(C_PTR), VALUE :: BUFFER ! The 'C address' of the array buffer
24             INTEGER (C_INT), VALUE :: N_BYTES ! Number of bytes in buffer
25         END SUBROUTINE PROCESS_BUFFER
26     END INTERFACE
```

- 27 2 This can be done using **C\_LOC** if the particular Fortran processor specifies that **C\_LOC** returns an appropriate  
 28 address:

```
29     REAL(R_QUAD), DIMENSION(:), ALLOCATABLE, TARGET :: QUAD_ARRAY
30     ...
31     CALL PROCESS_BUFFER(C_LOC(QUAD_ARRAY), INT(16*SIZE(QUAD_ARRAY),C_INT))
32     ! One quad real takes 16 bytes on this processor
```

### 33 C.12.5 Example of opaque communication between C and Fortran (18.3)

- 34 1 The following example demonstrates how a Fortran processor can make a modern object-oriented random number  
 35 generator written in Fortran available to a C program.

```
36     USE, INTRINSIC :: ISO_C_BINDING
37     ! Assume this code is inside a module
```

```

1
2     TYPE RANDOM_STREAM
3         ! A (uniform) random number generator (URNG)
4     CONTAINS
5         PROCEDURE(RANDOM_UNIFORM), DEFERRED, PASS(STREAM) :: NEXT
6         ! Generates the next number from the stream
7     END TYPE RANDOM_STREAM
8
9     ABSTRACT INTERFACE
10        ! Abstract interface of Fortran URNG
11        SUBROUTINE RANDOM_UNIFORM(STREAM, NUMBER)
12            IMPORT :: RANDOM_STREAM, C_DOUBLE
13            CLASS(RANDOM_STREAM), INTENT(INOUT) :: STREAM
14            REAL(C_DOUBLE), INTENT(OUT) :: NUMBER
15        END SUBROUTINE RANDOM_UNIFORM
16    END INTERFACE

```

- 17 2 A polymorphic object with declared type RANDOM\_STREAM is not [interoperable](#) with C. However, we can  
 18 make such a random number generator available to C by packaging it inside another nonpolymorphic, nonpara-  
 19 meterized derived type:

```

20     TYPE :: URNG_STATE ! No BIND(C), as this type is not interoperable
21         CLASS(RANDOM_STREAM), ALLOCATABLE :: STREAM
22     END TYPE URNG_STATE

```

- 23 3 The following two procedures will enable a C program to use our Fortran uniform random number generator:

```

24     ! Initialize a uniform random number generator:
25     SUBROUTINE INITIALIZE_URNG(STATE_HANDLE, METHOD) &
26         BIND(C, NAME="InitializeURNG")
27         TYPE(C_PTR), INTENT(OUT) :: STATE_HANDLE
28         ! An opaque handle for the URNG
29         CHARACTER(C_CHAR), DIMENSION(*), INTENT(IN) :: METHOD
30         ! The algorithm to be used
31
32         TYPE(URNG_STATE), POINTER :: STATE
33         ! An actual URNG object
34
35         ALLOCATE(STATE)
36         ! There needs to be a corresponding finalization
37         ! procedure to avoid memory leaks, not shown in this example
38         ! Allocate STATE%STREAM with a dynamic type depending on METHOD
39         ...
40         STATE_HANDLE=C_LOC(STATE)
41         ! Obtain an opaque handle to return to C
42     END SUBROUTINE INITIALIZE_URNG
43

```



```

1      ! Generate a random number:
2      SUBROUTINE GENERATE_UNIFORM(STATE_HANDLE, NUMBER) &
3          BIND(C, NAME="GenerateUniform")
4          TYPE(C_PTR), INTENT(IN), VALUE :: STATE_HANDLE
5          ! An opaque handle: Obtained via a call to INITIALIZE_URNG
6          REAL(C_DOUBLE), INTENT(OUT) :: NUMBER
7
8          TYPE(URNG_STATE), POINTER :: STATE
9          ! A pointer to the actual URNG
10
11         CALL C_F_POINTER(CPTR=STATE_HANDLE, FPTR=STATE)
12         ! Convert the opaque handle into a usable pointer
13         CALL STATE%STREAM%NEXT(NUMBER)
14         ! Use the type-bound procedure NEXT to generate NUMBER
15     END SUBROUTINE GENERATE_UNIFORM

```

## 16 C.12.6 Using assumed type to interoperate with C

### 17 C.12.6.1 Overview

18 1 The mechanism for handling [unlimited polymorphic](#) entities whose [dynamic type](#) is interoperable with C is  
19 designed to handle the following two situations:

- 20 (1) A formal parameter that is a C pointer to void. This is an address, and no further information  
21 about the entity is provided. The formal parameter corresponds to a dummy argument that is a  
22 nonallocatable nonpointer scalar or is an [assumed-size array](#).
- 23 (2) A formal parameter that is the address of a [C descriptor](#). Additional information on the status, type,  
24 size, and shape is implicitly provided. The formal parameter corresponds to a dummy argument that  
25 is [assumed-shape](#) or [assumed-rank](#).

26 2 In the first situation, it is the programmer's responsibility to explicitly provide any information needed on the  
27 status, type, size, and shape of the entity.

### 28 C.12.6.2 Mapping of interfaces with void \* C parameters to Fortran

29 1 A C interface for message passing or input/output functionality could be provided in the form

```
30     int EXAMPLE_send(const void *buffer, size_t buffer_size, const HANDLE_t *handle);
```

31 where the `buffer_size` argument is given in units of bytes, and the `handle` argument (which is of a type aliased  
32 to `int`) provides information about the target the buffer is to be transferred to. In this example, type resolution  
33 is not required.

34 2 The first method provides a thin binding; a call to `EXAMPLE_send` from Fortran directly invokes the C function.

```

35     INTERFACE
36         INTEGER (C_INT) FUNCTION example_send(buffer, buffer_size, handle) &
37         BIND(C, NAME='EXAMPLE_send')
38         USE, INTRINSIC :: ISO_C_BINDING

```

```

1         TYPE(*), INTENT (IN) :: buffer(*)
2         INTEGER (C_SIZE_T), VALUE :: buffer_size
3         INTEGER (C_INT), INTENT (IN) :: handle
4     END FUNCTION
5 END INTERFACE

```

6 3 It is assumed that this interface is declared in the specification part of the module MOD\_EXAMPLE\_OLD. An  
7 example of its use follows:

```

8     USE, INTRINSIC :: ISO_C_BINDING
9     USE MOD_EXAMPLE_OLD
10
11     REAL(C_FLOAT) :: x(100)
12     INTEGER(C_INT) :: y(10,10)
13     REAL(C_DOUBLE) :: z
14     INTEGER(C_INT) :: status, handle
15     ...
16     ! Assign values to x, y, z and initialize handle.
17     ...
18     ! Send values in x, y, and z using EXAMPLE_send.
19     status = example_send(x, C_SIZEOF(x), handle)
20     status = example_send(y, C_SIZEOF(y), handle)
21     status = example_send([ z ], C_SIZEOF(z), handle)

```

22 4 In those invocations, x and y are passed directly with sequence association, but it is necessary to make an array  
23 expression containing the value of z to pass it.

24 5 The second method provides a Fortran interface which is easier to use, but requires writing a separate C wrapper  
25 routine. With this method, a [C descriptor](#) is created because the buffer is [assumed-rank](#) in the Fortran interface;  
26 the use of an optional argument is also demonstrated.

```

27     INTERFACE
28         SUBROUTINE example_send(buffer, handle, status) BIND(C, NAME="EG_send_fortran")
29             USE, INTRINSIC :: ISO_C_BINDING
30             TYPE(*), CONTIGUOUS, INTENT (IN) :: buffer(..)
31             INTEGER (C_INT), INTENT (IN) :: handle
32             INTEGER (C_INT), INTENT(OUT), OPTIONAL :: status
33     END SUBROUTINE
34 END INTERFACE

```

35 6 It is assumed that this interface is declared in the specification part of a module MOD\_EXAMPLE\_NEW.  
36 Example invocations from Fortran are then

```

37     USE, INTRINSIC :: iso_c_binding
38     USE mod_example_new
39
40     TYPE, BIND(C) :: my_derived

```

```

1      INTEGER(C_INT) :: len_used
2      REAL(C_FLOAT) :: stuff(100)
3  END TYPE
4  TYPE(my_derived) :: w(3)
5  REAL(C_FLOAT) :: x(100)
6  INTEGER(C_INT) :: y(10,10)
7  REAL(C_DOUBLE) :: z
8  INTEGER(C_INT) :: status, handle
9  ...
10 ! Assign values to w, x, y, z and initialize handle.
11 ...
12 ! Send values in w, x, y, and z using example_send.
13 CALL example_send(w, handle, status)
14 CALL example_send(x, handle)
15 CALL example_send(y, handle)
16 CALL example_send(z, handle)
17 CALL example_send(y(:,5), handle) ! Fifth column of y.
18 CALL example_send(y(1,5), handle) ! Scalar y(1,5) passed by descriptor.

```

19 7 The wrapper routine can be written in C as follows.

```

20 #include "ISO_Fortran_binding.h"
21
22 void EXAMPLE_send_fortran(const CFI_cdesc_t *buffer, const HANDLE_t *handle,
23                          int *status)
24 {
25     int status_local;
26     size_t buffer_size;
27     int i;
28
29     buffer_size = buffer->elem_len;
30     for (i=0; i<buffer->rank; i++) {
31         buffer_size *= buffer->dim[i].extent;
32     }
33     status_local = EXAMPLE_send(buffer->base_addr,buffer_size, handle);
34     if (status != NULL) *status = status_local;
35 }

```

### 36 C.12.7 Using assumed-type variables in Fortran

37 1 An [assumed-type](#) dummy argument in a Fortran procedure can be used as an actual argument corresponding to an  
38 [assumed-type](#) dummy in a call to another procedure. In the following example, the Fortran subroutine SIMPLE\_  
39 SEND serves as a wrapper to hide the complications associated with calls to a C function named ACTUAL\_Send.  
40 Module COMM\_INFO contains node and address information for the current data transfer operations.

```

41 SUBROUTINE SIMPLE_SEND(buffer, nbytes)
42     USE comm_info, ONLY: my_node, r_node, r_addr
43     USE, INTRINSIC :: ISO_C_BINDING

```

```

1      IMPLICIT NONE
2
3      TYPE(*), INTENT (IN) :: buffer(*)
4      INTEGER              :: nbytes, ierr
5
6      INTERFACE
7          SUBROUTINE actual_Send(buffer, nbytes, node, addr, ierr) &
8          BIND(C, NAME="ACTUAL_Send")
9              IMPORT :: C_SIZE_T, C_INT, C_INTPTR_T
10             TYPE(*), INTENT (IN)      :: buffer(*)
11             INTEGER(C_SIZE_T), VALUE  :: nbytes
12             INTEGER(C_INT), VALUE     :: node
13             INTEGER(C_INTPTR_T), VALUE :: addr
14             INTEGER(C_INT), INTENT(OUT) :: ierr
15         END SUBROUTINE actual_Send
16     END INTERFACE
17
18     CALL actual_Send(buffer, INT(nbytes, C_SIZE_T), r_node, r_addr, ierr)
19
20     IF (ierr /= 0) THEN
21         PRINT *, "Error sending from node", my_node, "to node", r_node
22         PRINT *, "Program Aborting" ! Or call a recovery procedure
23         ERROR STOP ! Omit in the recovery case
24     END IF
25     END SUBROUTINE simple_Send

```

## 26 C.12.8 Simplifying interfaces for arbitrary rank procedures

- 27 1 There are situations where an [assumed-rank](#) dummy argument can be useful in Fortran, although a Fortran  
 28 procedure cannot itself access its value. For example, the IEEE inquiry functions in Clause 14 could be written  
 29 using an [assumed-rank](#) dummy argument instead of writing 16 separate specific routines, one for each possible  
 30 rank.
- 31 2 In particular, the specific procedures for the IEEE\_SUPPORT\_DIVIDE function could possibly be implemented  
 32 in Fortran as follows:

```

33     INTERFACE ieee_support_divide
34         MODULE PROCEDURE ieee_support_divide_noarg, ieee_support_divide_onearg_r, &
35             ieee_support_divide_onearg_d
36     END INTERFACE ieee_support_divide
37
38     ...
39
40     LOGICAL FUNCTION ieee_support_divide_noarg ()
41         ieee_support_divide_noarg = .TRUE.
42     END FUNCTION ieee_support_divide_noarg
43

```

```

1      LOGICAL FUNCTION ieee_support_divide_onearg_r (x)
2          REAL, INTENT (IN) :: x(..)
3          ieee_support_divide_onearg_r4 = .TRUE.
4      END FUNCTION ieee_support_divide_onearg_r
5
6      LOGICAL FUNCTION ieee_support_divide_onearg_d (x)
7          DOUBLE PRECISION, INTENT (IN) :: x(..)
8          ieee_support_divide_onearg_r8 = .TRUE.
9      END FUNCTION ieee_support_divide_onearg_d

```

## 10 C.12.9 Processing assumed-shape arrays in C

11 1 The example shown below calculates the product of individual elements of arrays A and B and returns the result  
12 in array C. The Fortran interface of `elemental_mult` will accept arguments of any type and rank. However, the  
13 C function will return an error code if any argument is not a two-dimensional `int` array. Note that the arguments  
14 are permitted to be array sections, so the C function does not assume that any argument is contiguous.

15 2 The Fortran interface is:

```

16      INTERFACE
17          FUNCTION elemental_mult(a, b, c) BIND(C, NAME="elemental_mult_c") RESULT(err)
18              USE, INTRINSIC :: ISO_C_BINDING
19              INTEGER(C_INT) :: err
20              TYPE(*), DIMENSION(..) :: a, b, c
21          END FUNCTION elemental_mult
22      END INTERFACE

```

23 3 The definition of the C function is:

```

24      #include "ISO_Fortran_binding.h"
25
26      int elemental_mult_c(CFI_cdesc_t * a_desc, CFI_cdesc_t * b_desc, CFI_cdesc_t * c_desc)
27      {
28          size_t i, j, ni, nj;
29
30          int err = 1; /* this error code represents all errors */
31
32          char * a_col = (char*) a_desc->base_addr;
33          char * b_col = (char*) b_desc->base_addr;
34          char * c_col = (char*) c_desc->base_addr;
35          char *a_elt, *b_elt, *c_elt;
36
37          /* Only support int. */
38          if (a_desc->type != CFI_type_int || b_desc->type != CFI_type_int ||
39              c_desc->type != CFI_type_int) {
40              return err;
41          }
42

```

```

1      /* Only support two dimensions. */
2      if (a_desc->rank != 2 || b_desc->rank != 2 || c_desc->rank != 2) {
3          return err;
4      }
5
6      ni = a_desc->dim[0].extent;
7      nj = a_desc->dim[1].extent;
8
9      /* Ensure the shapes conform. */
10     if (ni != b_desc->dim[0].extent || ni != c_desc->dim[0].extent) return err;
11     if (nj != b_desc->dim[1].extent || nj != c_desc->dim[1].extent) return err;
12
13     /* Multiply the elements of the two arrays. */
14     for (j = 0; j < nj; j++) {
15         a_elt = a_col;
16         b_elt = b_col;
17         c_elt = c_col;
18         for (i = 0; i < ni; i++) {
19             *(int*)a_elt = *(int*)b_elt * *(int*)c_elt;
20             a_elt += a_desc->dim[0].sm;
21             b_elt += b_desc->dim[0].sm;
22             c_elt += c_desc->dim[0].sm;
23         }
24         a_col += a_desc->dim[1].sm;
25         b_col += b_desc->dim[1].sm;
26         c_col += c_desc->dim[1].sm;
27     }
28     return 0;
29 }

```

### 30 C.12.10 Creating a contiguous copy of an array

31 1 A C function might need to create a contiguous copy of an array section, for example, to pass the array section  
32 as an actual argument corresponding to a dummy argument with the CONTIGUOUS attribute. The following  
33 example provides functions that can be used to copy an array described by a `CFI_cdesc_t` descriptor to a  
34 contiguous buffer. The input array need not be contiguous.

35 2 The C functions are:

```

36     #include "ISO_Fortran_binding.h"
37     /* Other necessary includes omitted. */
38
39     /*
40     * Returns the number of elements in the object described by desc.
41     * If it is an array, it need not be contiguous.
42     * (The number of elements could be zero).
43     */
44     size_t numElements(const CFI_cdesc_t * desc)

```

```
1      {
2          CFI_rank_t r;
3          size_t num = 1;
4
5          for (r = 0; r < desc->rank; r++) {
6              num *= desc->dim[r].extent;
7          }
8          return num;
9      }
10
11     /*
12     * Auxiliary recursive function to copy an array of a given rank.
13     * Recursion is useful because an array of rank n is composed of an
14     * ordered set of arrays of rank n-1.
15     */
16     static void _copyToContiguous (const CFI_cdesc_t *vald, void *output,
17                                     const void *input, CFI_rank_t rank)
18     {
19         CFI_index_t e;
20
21         if (rank == 0) {
22             /* Copy scalar element. */
23             memcpy (output, input, vald->elem_len);
24             output = (void *)((char *)output + vald->elem_len);
25         }
26         else {
27             for (e = 0; e < vald->dim[rank-1].extent; e++) {
28                 /* Recurse on subarrays of lesser rank. */
29                 output = _copyToContiguous (vald, output, input, rank-1);
30                 input = (void *) ((char *)input + vald->dim[rank].sm);
31             }
32         }
33         return output;
34     }
35
36     /*
37     * General routine to copy the elements in the array described by vald
38     * to buffer, as done by sequence association. The array itself can
39     * be non-contiguous. This is not the most efficient approach.
40     */
41     void copyToContiguous (void * buffer, const CFI_cdesc_t * vald) {
42         _copyToContiguous (vald, buffer, vald->base_addr, vald->rank);
43     }
```

### 1 C.12.11 Changing the attributes of an array

2 1 A C programmer might want to call more than one Fortran procedure and the attributes of an array involved  
3 might differ between the procedures. In this case, it is necessary to set up more than one C descriptor for the  
4 array. For example, this code fragment initializes the first C descriptor for an allocatable entity of rank 2, calls  
5 a procedure that allocates the array described by the first C descriptor, constructs the second C descriptor by  
6 invoking CFI\_section with the value CFI\_attribute\_other for the attribute parameter, then calls a procedure  
7 that expects an assumed-shape array.

```
8         CFI_CDESC_T(2) loc_alloc, loc_assum;  
9         CFI_cdesc_t * desc_alloc = (CFI_cdesc_t *)&loc_alloc,  
10            * desc_assum = (CFI_cdesc_t *)&loc_assum;  
11         CFI_index_t extents[2];  
12         CFI_rank_t rank = 2;  
13         int flag;  
14  
15         flag = CFI_establish(desc_alloc,  
16            NULL,  
17            CFI_attribute_allocatable,  
18            CFI_type_double,  
19            sizeof(double),  
20            rank,  
21            NULL);  
22  
23         Fortran_factor (desc_alloc, ...); /* Allocates array described by desc_alloc. */  
24  
25         /* Extract extents from descriptor. */  
26         extents[0] = desc_alloc->dim[0].extent;  
27         extents[1] = desc_alloc->dim[1].extent;  
28  
29         flag = CFI_establish(desc_assum,  
30            desc_alloc->base_addr,  
31            CFI_attribute_other,  
32            CFI_type_double,  
33            sizeof(double),  
34            rank,  
35            extents);  
36  
37         Fortran_solve (desc_assum, ...); /* Uses array allocated in Fortran_factor. */
```

38 2 After invocation of the second CFI\_establish, the lower bounds stored in the dim member of desc\_assum will  
39 have the value 0 even if the corresponding entries in desc\_alloc have different values.

### 40 C.12.12 Creating an array section in C using CFI\_section

41 1 The C function set\_odd sets every second element of an array to a specific value, beginning with the first element.  
42 If does this by making an array section descriptor for the elements to be set, and calling a Fortran subroutine



1 SET\_ALL that sets every element of an assumed-shape array to a specific value. An interface block for set\_odd  
2 permits it to be also called from Fortran.

```
3     SUBROUTINE set_all(int_array, val) BIND(C)
4         INTEGER(C_INT) :: int_array(:)
5         INTEGER(C_INT), VALUE :: val
6         int_array = val
7     END SUBROUTINE
8
9     INTERFACE
10        SUBROUTINE set_odd(int_array, val) BIND(C)
11            USE, INTRINSIC :: ISO_C_BINDING, ONLY : C_INT
12            INTEGER(C_INT) :: int_array(:)
13            INTEGER(C_INT), VALUE :: val
14        END SUBROUTINE
15    END INTERFACE
16
17    #include "ISO_Fortran_binding.h"
18
19    void set_odd(CFI_cdesc_t *int_array, int val)
20    {
21        CFI_index_t lower_bound[1], upper_bound[1], stride[1];
22        CFI_CDESC_T(1) array;
23        int status;
24        /* Create a new descriptor which will contain the section. */
25        status = CFI_establish((CFI_cdesc_t *)&array,
26                               NULL,
27                               CFI_attribute_other,
28                               int_array->type,
29                               int_array->elem_len,
30                               /* rank */ 1,
31                               /* extents is ignored */ NULL);
32
33        lower_bound[0] = int_array->dim[0].lower_bound;
34        upper_bound[0] = lower_bound[0] + (int_array->dim[0].extent - 1);
35        stride[0] = 2;
36
37        status = CFI_section((CFI_cdesc_t *)&array,
38                            int_array,
39                            lower_bound,
40                            upper_bound,
41                            stride);
42
43        set_all( (CFI_cdesc_t *) &array, val);
44
45        /* Here one could make use of int_array and access all its data. */
46    }
```

1 2 The `set_odd` procedure can be called from Fortran as follows:

```
2     INTEGER(C_INT) :: d(5)
3     d = (/ 1, 2, 3, 4, 5 /)
4     CALL set_odd(d, -1)
5     PRINT *, d
```

6 3 This program will print something like:

```
7     -1    2    -1    4    -1
```

8 4 During execution of the subroutine `SET_ALL`, its dummy argument `INT_ARRAY` would have size (and upper bound) 3.

10 5 It is also possible to invoke `set_odd()` from C. However, it would be the C programmer's responsibility to make sure that all members of the [C descriptor](#) have the correct value on entry to the function. Inserting additional checking into the function could alleviate this problem.

13 6 Following is an example C function that dynamically generates a [C descriptor](#) for an [assumed-shape array](#) and calls `set_odd`.

```
15     #include <stdio.h>
16     #include <stdlib.h>
17     #include "ISO_Fortran_binding.h"
18
19     #define ARRAY_SIZE 5
20
21     void example_of_calling_set_odd(void)
22     {
23         CFI_CDESC_T(1) d;
24         CFI_index_t extent[1];
25         CFI_index_t subscripts[1];
26         void *base;
27         int i, status;
28
29         base = malloc(ARRAY_SIZE*sizeof(int));
30         extent[0] = ARRAY_SIZE;
31         status = CFI_establish((CFI_cdesc_t *)&d,
32                               base,
33                               CFI_attribute_other,
34                               CFI_type_int,
35                               /* element length is ignored */ 0,
36                               /* rank */ 1,
37                               extent);
38
39         set_odd((CFI_cdesc_t *)&d, -1);
40
41         for (i=0; i<ARRAY_SIZE; i++) {
42             subscripts[0] = i;
```

```

1         printf(" %d",*((int *)CFI_address((CFI_cdesc_t *)&d, subscripts)));
2     }
3     putc(10, stdout);
4     free(base);
5 }

```

6 The above C function will print similar output to that of the preceding Fortran program.

### 7 C.12.13 Use of CFI\_setpointer

8 1 The C function `change_target` modifies a pointer to an integer variable to become associated with a global  
9 variable defined inside C:

```

10     #include "ISO_Fortran_binding.h"
11
12     int y = 2;
13
14     void change_target(CFI_cdesc_t *ip) {
15         CFI_CDESC_T(0) yp;
16         int status;
17         /* Make local yp point at y. */
18         status = CFI_establish((CFI_cdesc_t *)&yp,
19                               &y,
20                               CFI_attribute_pointer,
21                               CFI_type_int,
22                               /* elem_len is ignored */ sizeof(int),
23                               /* rank */ 0,
24                               /* extents are ignored */ NULL);
25         /* Pointer-associate ip with (the target of) yp. */
26         status = CFI_setpointer(ip, (CFI_cdesc_t *)&yp, NULL);
27         if (status != CFI_SUCCESS) {
28             ... report run time error...
29         }
30     }

```

31 2 The restrictions on the use of `CFI_establish` prohibit direct modification of the incoming pointer entity `ip` by  
32 invoking that function on it.

33 3 The following program illustrates the usage of `change_target` from Fortran.

```

34     PROGRAM change_target_example
35         USE, INTRINSIC :: ISO_C_BINDING
36         INTERFACE
37             SUBROUTINE change_target(ip) BIND(C)
38                 IMPORT :: C_INT
39                 INTEGER(C_INT), POINTER :: ip
40             END SUBROUTINE
41         END INTERFACE

```

```

1      INTEGER(C_INT), TARGET :: it = 1
2      INTEGER(C_INT), POINTER :: it_ptr
3      it_ptr => it
4      WRITE (*,*) it_ptr
5      CALL change_target(it_ptr)
6      WRITE (*,*) it_ptr

```

7 4 This will print something similar to

```

8      1
9      2

```

#### 10 C.12.14 Mapping of MPI interfaces to Fortran

11 1 The Message Passing Interface (MPI) specifies procedures for exchanging data between MPI processes. This  
12 example shows the usage of `MPI_Send` and is similar to the second variant of `EXAMPLE_Send` in C.12.6.2. It also  
13 shows the usage of assumed-length character dummy arguments and optional dummy arguments.

14 2 `MPI_Send` has the C prototype:

```

15      int MPI_Send(void *buf, int count, MPI_Datatype datatype, int dest, int tag,
16                  MPI_Comm comm);

```

17 where `MPI_Datatype` and `MPI_Comm` are opaque handles. Most MPI C functions return an error code, which in  
18 Fortran is the last dummy argument to the corresponding subroutine and can be made optional. Thus, the use  
19 of a Fortran subroutine requires a wrapper function, declared as

```

20      void MPI_Send_f(CFI_cdesc_t *buf, int count, MPI_Datatype_f datatype, int dest,
21                    int tag, MPI_Datatype_f comm, int *ierror);

```

22 3 This wrapper function will convert `MPI_Datatype_f` and `MPI_Comm_f` to `MPI_Datatype` and `MPI_Comm`, and pro-  
23 duce a contiguous `void *` buffer from `CFI_cdesc_t *buf` (if necessary).

24 4 Similarly, the wrapper function for `MPI_Comm_set_name` could have the C prototype:

```

25      void MPI_Comm_set_name_f(MPI_Comm comm, CFI_cdesc_t *comm_name, int *ierror);

```

26 5 The Fortran handle types and interfaces are defined in the module `MPI_F08`. For example,

```

27      MODULE mpi_f08
28      ...
29      TYPE, BIND(C) :: mpi_comm
30      PRIVATE
31      INTEGER(C_INT) :: mpi_val
32      END TYPE mpi_comm
33
34      INTERFACE
35      SUBROUTINE MPI_SEND(buf,count,datatype,dest,tag,comm,ierror) &
36      BIND(C, NAME='MPI_Send_f')
37      USE, INTRINSIC :: ISO_C_BINDING

```

```

1      IMPORT :: MPI_Datatype, MPI_Comm
2      TYPE(*), DIMENSION(..), INTENT (IN) :: buf
3      INTEGER(C_INT), VALUE, INTENT (IN) :: count, dest, tag
4      TYPE(mpi_datatype), INTENT (IN) :: datatype
5      TYPE(mpi_comm), INTENT (IN) :: comm
6      INTEGER(C_INT), OPTIONAL, INTENT (OUT) :: ierror
7  END SUBROUTINE mpi_send
8
9      SUBROUTINE mpi_comm_set_name(comm,comm_name,ierror) &
10     BIND(C, NAME='MPI_Comm_set_name_f')
11     USE, INTRINSIC :: ISO_C_BINDING
12     IMPORT :: mpi_comm
13     TYPE(mpi_comm), INTENT (IN) :: comm
14     CHARACTER(KIND=C_CHAR, LEN=*), INTENT (IN) :: comm_name
15     INTEGER(C_INT), OPTIONAL, INTENT (OUT) :: ierror
16     END SUBROUTINE mpi_comm_set_name
17 END INTERFACE
18 ...
19 END MODULE mpi_f08

```

20 6 Some examples of invocation from Fortran are:

```

21     USE, INTRINSIC :: ISO_C_BINDING
22     USE :: MPI_f08
23
24     TYPE(mpi_comm) :: comm
25     REAL :: x(100)
26     INTEGER :: y(10,10)
27     REAL(KIND(1.0d0)) :: z
28     INTEGER :: dest, tag, ierror
29     ...
30     ! Assign values to x, y, z and initialize MPI variables.
31     ...
32
33     ! Set the name of the communicator.
34     CALL mpi_comm_set_name(comm, "Communicator Name", ierror)
35
36     ! Send values in x, y, and z.
37     CALL mpi_send(x, 100, MPI_REAL, dest, tag, comm, ierror)
38     IF (ierror/=0) PRINT *, 'WARNING: X send error', ierror
39     CALL mpi_send(y(3,:), 10, MPI_INTEGER, dest, tag, comm)
40     CALL mpi_send(z, 1, MPI_DOUBLE_PRECISION, dest, tag, comm)

```

41 7 The first example sends the entire array X and includes the optional error argument return value. The second  
42 example sends a noncontiguous subarray (the third row of Y) and the third example sends a scalar Z. Note the  
43 differences between the calls in this example and those in [C.12.6.2](#).

## 1 C.13 Clause 19 notes : Examples of host association (19.5.1.4)

2 1 The first two examples are examples of valid [host association](#). The third example is an example of invalid [host](#)  
3 [association](#).

### 4 Example 1:

```
5 PROGRAM A
6     INTEGER I, J
7     ...
8 CONTAINS
9     SUBROUTINE B
10        INTEGER I ! Declaration of I hides
11                ! program A's declaration of I
12        ...
13        I = J     ! Use of variable J from program A
14                ! through host association
15    END SUBROUTINE B
16 END PROGRAM A
```

### 17 Example 2:

```
18 PROGRAM A
19     TYPE T
20     ...
21 END TYPE T
22 ...
23 CONTAINS
24     SUBROUTINE B
25        IMPLICIT TYPE (T) (C) ! Refers to type T declared below
26                               ! in subroutine B, not type T
27                               ! declared above in program A
28        ...
29        TYPE T
30        ...
31    END TYPE T
32    ...
33 END SUBROUTINE B
34 END PROGRAM A
```

### 35 Example 3:

```
36 PROGRAM Q
37     REAL (KIND = 1) :: C
38     ...
39 CONTAINS
40     SUBROUTINE R
41        REAL (KIND = KIND (C)) :: D ! Invalid declaration
```

```
1                                     ! See below
2         REAL (KIND = 2) :: C
3         ...
4         END SUBROUTINE R
5     END PROGRAM Q
```

6 2 In the declaration of D in subroutine R, the use of C would refer to the declaration of C in subroutine R, not  
7 program Q. However, it is invalid because the declaration of C is required to occur before it is used in the  
8 declaration of D (10.1.12).

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