



Photo: Hannah Theile (OVGU)

Tseitin or not Tseitin?

The Impact of CNF Transformations on Feature-Model Analyses

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Implementing Configurable Software Systems

A Configurable Graph

```
class Node {  
    #ifdef LABELED  
        std::string label;  
    #endif  
    #ifdef COLORED  
        std::string color;  
    #endif  
};  
  
class Edge {  
    #ifdef DIRECTED  
        Node from, to;  
    #elif UNDIRECTED && HYPER  
        std::set<Node> nodes;  
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```

Product Line Implementation

(here: C++ with C preprocessor)

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Configuration

A Labeled Directed Graph

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A Colored Undirected Hypergraph

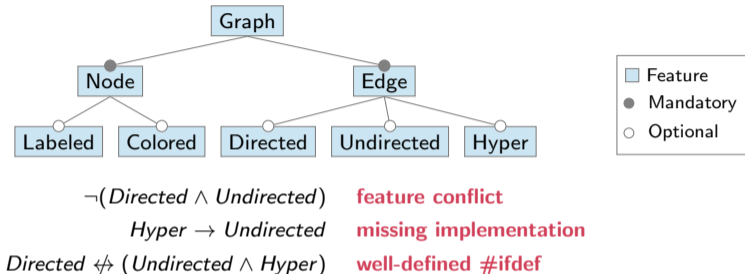
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Product Implementation

Modeling Features and their Dependencies

Feature Models

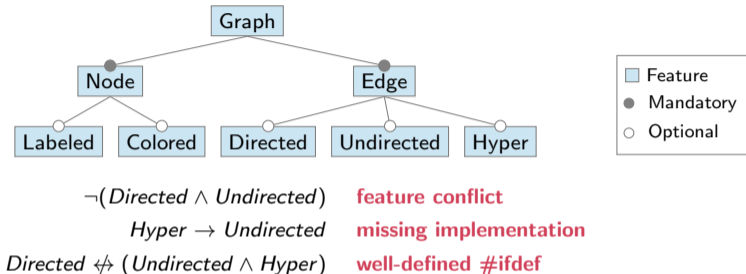
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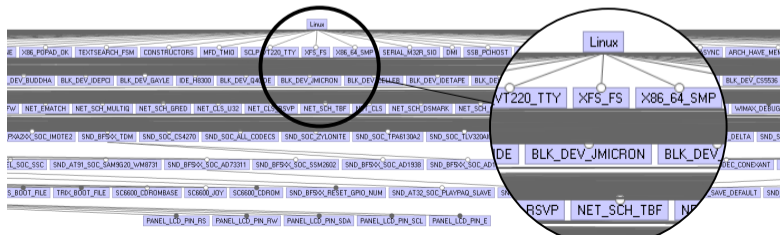
Feature Models

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The Linux Kernel

- > 12000 features [2016]
- > 10⁵⁰⁰⁰ products [2016]
- 114 dead features [2013]
- 151 reverse dependency bugs [2019]



Analyzing Feature Models with SAT and #SAT Solvers

Feature-Model Analysis

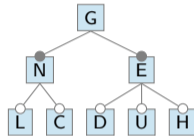


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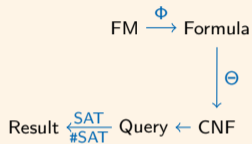
A Feature Model FM



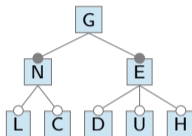
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$\xrightarrow{\Phi}$

As a Formula $\Phi(FM)$

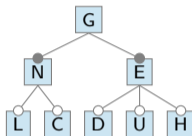
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$\downarrow \Theta$

As a CNF $\Theta(\Phi(FM))$

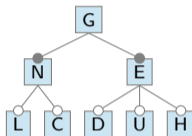
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Core Features

$\{G, N, E\}$

\leftarrow

Core Feature $F?$

$\text{SAT}(\Theta(\Phi(FM)) \wedge \neg F)$

\leftarrow

As a CNF $\Theta(\Phi(FM))$

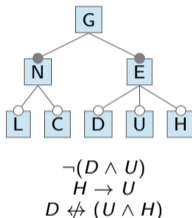
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Φ

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Core Feature F ?

$SAT(\Theta(\Phi(FM)) \wedge \neg F)$

Feature Model Cardinality

8

Products in FM ?

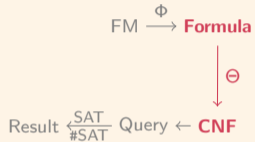
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Often Overlooked: Conjunctive Normal Form (CNF)

Feature-Model Analysis



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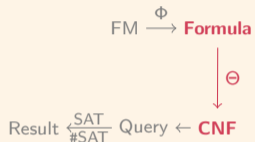
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Often Overlooked: Conjunctive Normal Form (CNF)

Feature-Model Analysis



Conjunctive Normal Form

- **conjunction** \wedge of **disjunctions** \vee of **literals** $X, \neg X$
- here: a set of **clauses**, which are sets of literals
- used by almost all solvers

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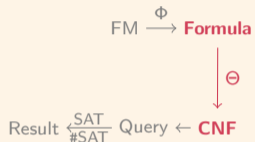
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Our Goal: Raise Awareness for CNF Transformations

- how to **transform** feature-model formulas **into CNF**?
 \Rightarrow describe and classify CNF transformations
- does this **impact** the work of **practitioners and researchers**?
 \Rightarrow evaluate efficiency and correctness on feature models

CNF Transformations

Distributive $\Theta = D$

apply laws of logic (**De Morgan's laws** and **distributivity**)

$$D \not\leftrightarrow (U \wedge H)$$

$$\xrightarrow{D} (D \vee (U \wedge H)) \wedge (\neg D \vee \neg(U \wedge H))$$

$$\xrightarrow{D} \{\{D, U\}, \{D, H\}, \{\neg D, \neg U, \neg H\}\}$$

- ✓ **equivalence**
SAT ✓, #SAT = 4
- ✓ easy to implement
- ✗ **exponential** complexity

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Tseitin $\Theta = T$ [83]

abbreviate a subformula ϕ with an auxiliary variable $x_\phi \leftrightarrow \phi$

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Plaisted-Greenbaum $\Theta = PG$ [86]

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- ✓ **equi-assignability** SAT ✓
- ✓ **linear** complexity $< T$
- ✗ **equi-countability** #SAT = 5

Evaluation

Research Questions

- RQ 1 **efficiency** of CNF transformations?
- RQ 2 CNF transformation → **efficiency** of analyses?
- RQ 3 CNF transformation → **correctness** of analyses?

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Experimental Setup

- 22 configurable software systems
- 3 CNF transformation tools
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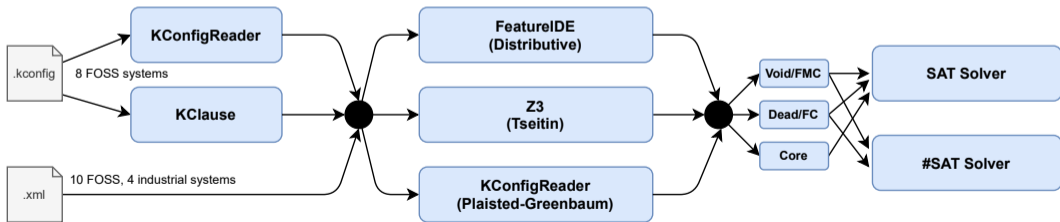
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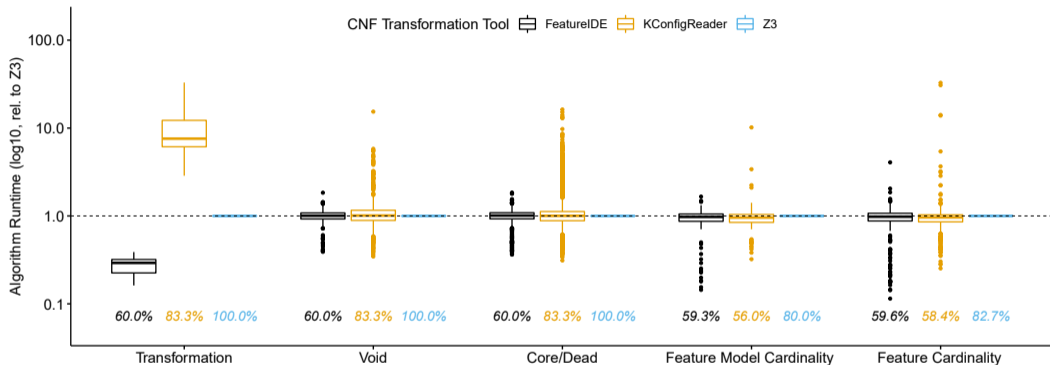
Stage 1: Formula Extraction

Stage 2: CNF Transformation

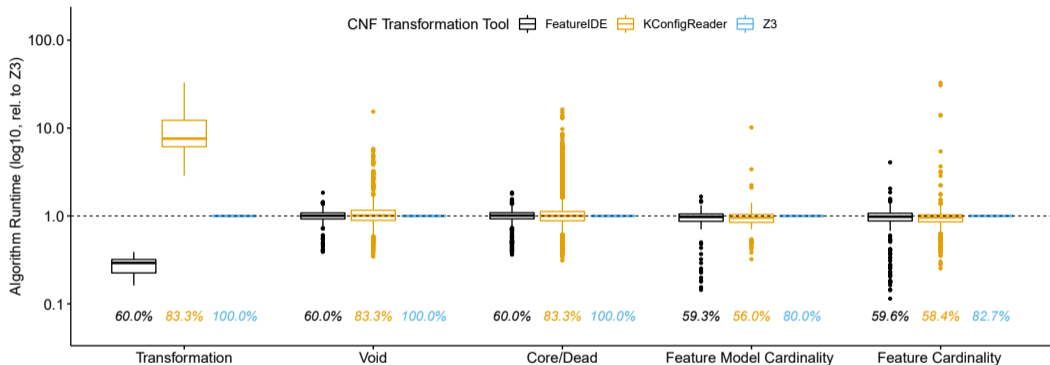
Stage 3: Automated Analysis



Efficiency of CNF Transformations (RQ 1) and Analyses (RQ 2)

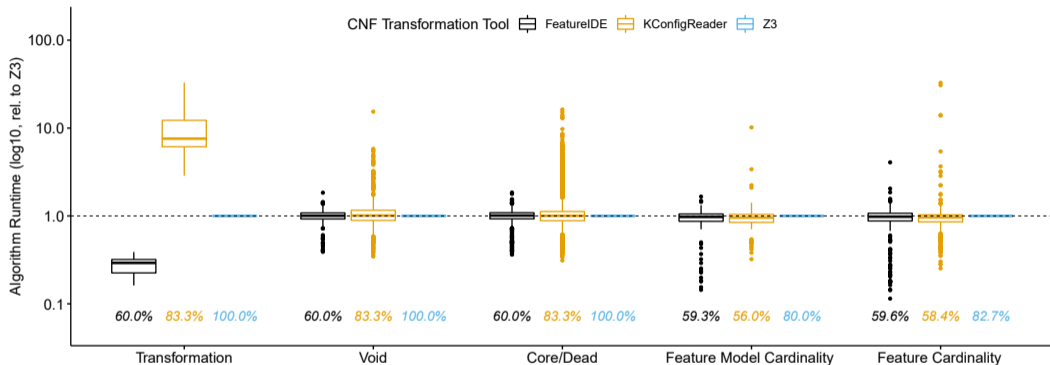


Efficiency of CNF Transformations (RQ 1) and Analyses (RQ 2)



RQ 1: *D* often fails,
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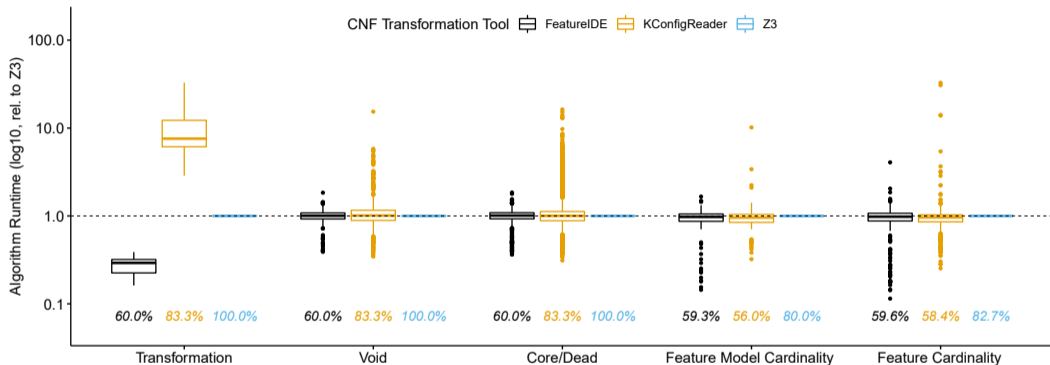
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RQ 2 (SAT): almost all calls succeed, solve time varies by factor 0.31–16.27

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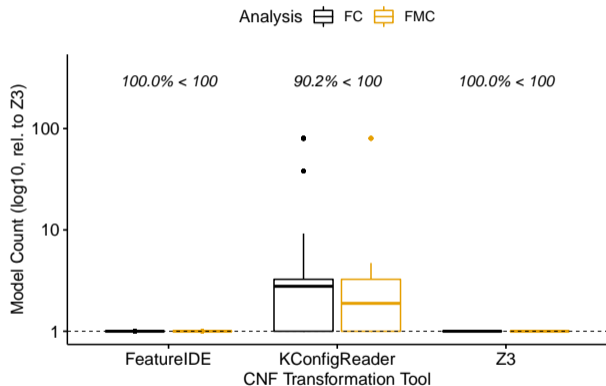


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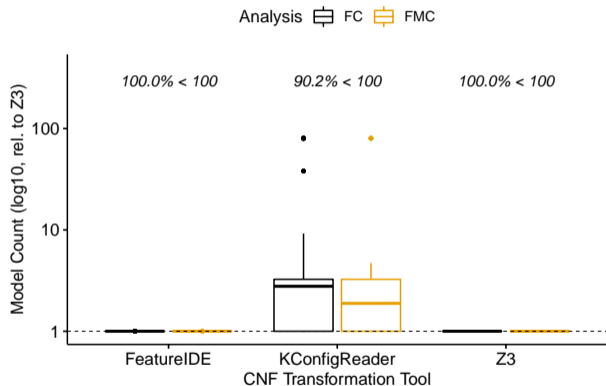
RQ 2 (SAT): almost all calls succeed, solve time varies by factor 0.31–16.27

RQ 2 (#SAT): 81.6% of calls succeed, solve time varies by factor 0.11–32.7

Correctness of #SAT-Based Analyses (RQ 3)



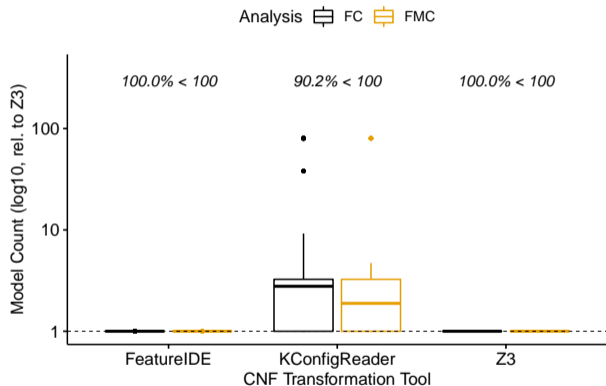
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- incorrect by factor ≈ 3 (median)
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Our Recommendations

- RQ 1** *D* for small, *T* for large models
- RQ 2** largely depends on the model
 \Rightarrow future work
- RQ 3** do not use *PG* for #SAT

Conclusion

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apply laws of logic

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- ✓ easy to implement
- ✗ **exponential** complexity

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often fails on large models

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abbreviate ϕ with $x_\phi \leftrightarrow \phi$

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Tseitin or not Tseitin? \Rightarrow Yes!

find out more:

<https://github.com/ekuitertseitin-or-not-tseitin>

