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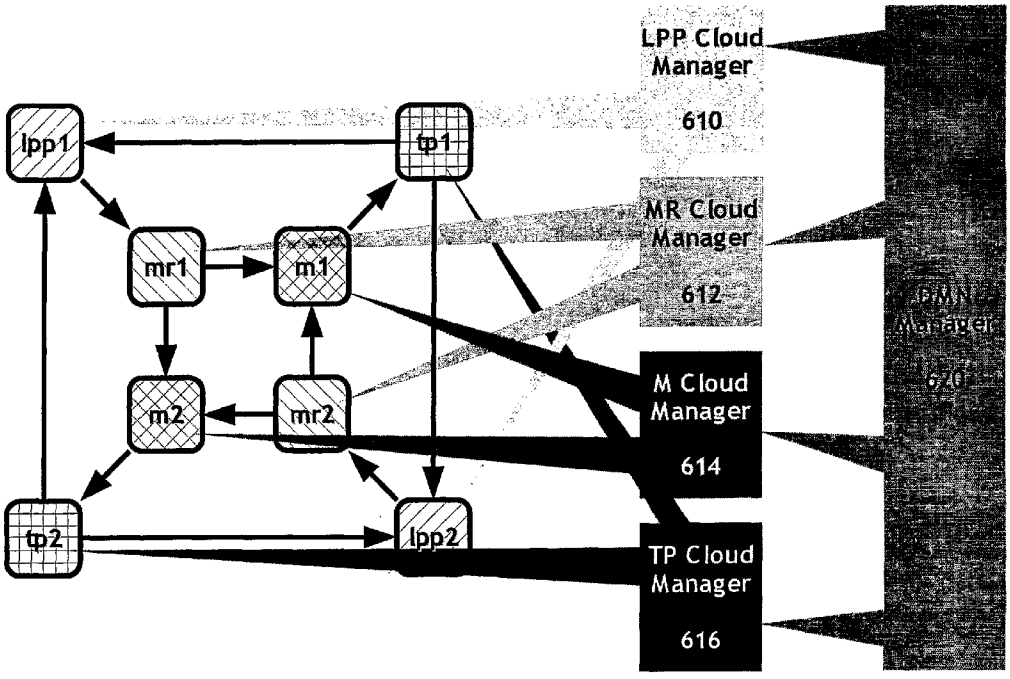

(57) Abstract: A distributed mediation network and method of employing such is provided, having a plurality of different types of network module. Each module has a non-reciprocal path therethrough for network traffic and the distribution of network traffic across the network is managed by an autonomic control plane.

WO 2007/144611 A1

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SELF-MANAGED DISTRIBUTED MEDIATION NETWORKS

Introduction

The present invention relates to self-managed distributed mediation networks having a requirement both for disseminated, peer-to-peer communication, and for a degree of control of information gathering over the sum of such disseminated messages.

Background to the invention

Peer-to-peer (P2P) communication systems allow computational entities (peers) to establish software connections (virtual channels) between one another. P2P systems therefore allow peers to communicate or share computational tasks and resources without the explicit need for centralised control. P2P can operate in a generalised network having one or more servers: a peer may provide information (publish) to at least one service on the network and/or register (subscribe) with services on that network to receive information published by another peer.

Messaging systems that benefit from the provision of centralised control are also known. Here, all messages are directed from publishers to subscribers, via a central locus where some computation (mediation) is performed on the messages. New messages (digests, for example) are generated from the input messages and sent to appropriate subscribers.

In prior art centralised mediation systems, all message traffic is transmitted through a central network point (locus), where the mediation service resides. Viewed in terms of logical elements, such systems are constructed as a star-shaped architectural model with a central point of control, where mediation tasks are executed. This model is shown in Figure 1A: each source (publisher) and sink (subscriber) of information has a line of communication that connects to the central mediation hub. In many cases, the sources and sinks represent the same entities operating in different modes, and may not be architecturally distinguishable.

The problems associated with such an architecture are well known. They are prone to suffer from a lack of bandwidth at the point of mediation. Even though the logical star shape may be superimposed upon a physical network that is highly connected, the essential flow of all information through a central point causes an inherent throughput bottleneck, based upon the bandwidth available between this point and the network

(see Figure 1B). Although advances in networking technologies mean that bandwidth availability continues to improve, increasing bandwidth has an inherent financial cost, and in certain scenarios can cause a real limitation to the throughput of the overall system. This limitation is manifested as a restriction on either the maximum number of users, or the rate at which each user is able to send and receive information.

Indeed there are many examples of systems where neither P2P architectures nor centralised mediation architectures are wholly satisfactory. Often some logical process is required to act over the sum of messages broadcast within the messaging system. Examples of classes of systems where neither architecture is completely suitable include: a trading system where potential buyers and sellers advertise to each other, mediation is required to ensure a transactional matching of requirements; a mediated news or publishing system where a central authority acts as the editorial control, before information is disseminated; a system which is not actively controlled but which requires an ordered log of information flow to be maintained in a central repository; a conversation service which allows a recent context to be presented to a user joining an ongoing conversation; distributing cryptographic keys (the so-called key distribution problem); systems for finding the location of data (state) and services on a distributed network; and systems for locating and communicating with mobile users.

All the examples above have in common a requirement both for peer-to-peer communication, and for a degree of centralised mediation of the flow of information when the communication is viewed as a whole.

Applicant's co-pending US patent application number 10/903,156, incorporated herein by reference, describes a distributed mediation network that overcomes many of the disadvantages described above. A distributed mediation network of this type overcomes the problems associated with providing a mediation service at a single server by distributing the service among a number of logically discrete entities. In order to do this, a mediation application must be amenable to logical partitioning into discrete mediation application components. This permits the mediation service to be partitioned into a set of mediation segment services distributed across a resource pool of servers, with each server providing the mediation service for one or more of these segment services. Hereinafter, this approach will be termed distributed mediation.

To properly eliminate the bandwidth problems at every mediation point, it must be possible to evenly load balance the mediation service across the available resource pool. In systems that exhibit fluctuating demands over time the load across the pool must be dynamically balanced. As such, it must be possible to dynamically change the way in which the application is partitioned. It is therefore necessary to be able to move a mediation segment service from one server to another. Moreover, the movement of a segment service must preserve externally observed causality. That is, the ordering of the interactions with each segment service must be preserved in the face of changes to how that segment service is implemented. This requirement is vitally important in many systems in which out of order interactions have serious consequences, such as financial systems.

In the distributed mediation network, information is classified by content, and mediation requirements are separately served in different processes according to that content-based classification. As demand varies with time over the classification, the corresponding mediation application components may be physically moved to balance both network and computational loads for the whole system. Such load balancing can satisfy the demands placed upon it up to some threshold governed by the sum of the computational, I/O and memory resources of the available servers offering mediation. Beyond this threshold, the quality of service will degrade as the available resources simply cannot handle the load. In order to address this problem, a distributed mediation network will preferably provide mechanisms for the introduction of additional computer resources to the system. Similarly, when excessive resources are available to a mediated application it is preferably possible to remove deployed computational capacity.

The term 'autonomic' has historically been used to refer to the aspect of the nervous system that acts subconsciously to regulate the body, such as the control of breathing rates or the heartbeat. It has recently been used to refer to computer networks that are capable of analogous self-regulation. An autonomic system may be capable of, amongst other things, self-repair, self-configuration, self-monitoring, and self-optimisation, all without the need for external input. Indeed, in the autonomic paradigm, any changes that occur autonomically are in fact impossible for the user to detect.

An autonomic computing system consists of autonomic elements. These autonomic elements are logical constructs that monitor some aspect of the system, analysing its output and taking specific actions to adjust it, so that the overall system is able to meet specific requirements, often expressed as service level agreements (SLAs). SLAs specify the information technology resources needed by a line of business and the specific applications that they maintain.

Autonomic elements are self-organising and are able to discover each other, operate independently, negotiate or collaborate as required, and organise themselves such that the emergent stratified management of the system as a whole reflects both the bottom up demand for resources and the top down business-directed application of those resources to achieve specific goals.

It is an object of the present invention to provide autonomic functionality to a distributed mediation network.

Throughout this document, only the term physical node refers to physical machines. The terms "node" and "logical node" are used interchangeably to refer to the locus having state properties. In terms of the logical topology of the mediation network, a module provides the functionality of an associated logical node.

Hereinafter, the term "high watermark" is used to indicate a maximum threshold level of traffic that may be handled by a single element or node within the network, while the term "low watermark" indicates the minimum threshold, each element being configured to handle traffic levels in a range between the high watermark and the low watermark.

Summary of the Invention

According to a first aspect of the present invention, there is provided a distributed mediation network, comprising:

a plurality of types of network module, including:

local points of presence (LPP) modules for receiving and transmitting network traffic between the mediation network and client programs;

mediator (M) modules for hosting mediation tasks;

mediator router (MR) modules for analyzing the content of incoming messages, each MR module routing the incoming messages to a predetermined mediation task in dependence upon said content; and,

transmission proxy (TP) modules for forwarding messages to at least one of said LPP modules, wherein each of the MR, M and TP modules are adapted such that all paths for network traffic therethrough are non-reciprocal; and,

an autonomic control plane for managing the distribution of network traffic amongst the modules.

Preferably, the network couples network traffic along a unidirectional mediation cycle, in which: LPP modules address MR modules, MR modules in turn address M modules, M modules in turn address TP modules, and TP modules in turn address LPP modules.

According to a second aspect of the present invention, there is provided method for mediating the flow of network traffic in a computer network, wherein the computer network has a mediation network that comprises: a plurality of types of network module, including:

local points of presence (LPP) modules for receiving and transmitting network traffic between the mediation network and client programs;

mediator (M) modules for hosting mediation tasks;

mediator router (MR) modules for analyzing the content of incoming messages, each MR module routing the incoming messages to a predetermined mediation task in dependence upon said content; and,

transmission proxy (TP) modules for forwarding messages to at least one of said LPP modules, wherein each of the MR, M and TP modules are adapted such that all paths for network traffic therethrough are non-reciprocal; and,

an autonomic control plane for managing the distribution of network traffic amongst the modules,

wherein, in the method, incoming messages are propagated along a mediation cycle that comprises the steps of:

an LPP module addressing incoming messages to a respective one of said at least one mediator router (MR) modules;

at said addressed MR module, analyzing the content of incoming messages and routing said messages to a predetermined mediator module in dependence upon said analyzed content;

at said predetermined mediator module, applying the mediation task to said analyzed messages and directing mediated messages to a respective one of said TP modules; and

at said TP module that receives said mediated messages, forwarding said mediated messages to at least one of said LPP modules.

The present invention provides a distributed mediation network in which the traffic load on each node or module is autonomically managed. The autonomic control plane adjusts the distribution of network traffic amongst the modules/nodes according to the circumstances in which the network finds itself.

The paths for the network traffic passing through the modules are non-reciprocal. That is, while the network paths between modules are unidirectional, the modules may direct network traffic in more than one direction. For instance, a given MR module may receive traffic from a plurality of LPP modules and may direct traffic to a plurality of M modules, but may not direct traffic to LPP modules or receive traffic from M modules. The result is a cyclic network, in which network traffic passes from LPP to MR to M to TP to LPP.

As well as network traffic or messages, control messages/signals may be passed through the network in order to effect commands given by the autonomic control plane. Control messages need not necessarily follow the non-reciprocal paths through the modules that are followed by network messages. Nevertheless, in some embodiments control messages cannot flow against the unidirectional paths followed by network traffic between nodes. For example, in such embodiments, a control message may not be directed against the cycle followed by network traffic (for example, a control message cannot pass from an M module to an MR module, or from an LPP module to a TP module).

In a preferred embodiment, the autonomic control plane receives information regarding the status of nodes within the distributed network from sensor interfaces incorporated therein, and may instruct action by the nodes through effector interfaces.

In a preferred embodiment, the control plane is itself distributed across a number of resources. This allows the control plane to continue functioning should one or more of these resources malfunction, and, moreover, allows the spread of the computational

load due to autonomic functions to be optimally distributed at any given time. For example, should one resource be in heavy demand to perform an unrelated task, the remaining resources may assume the burden of autonomic control.

The autonomic control plane is therefore able to monitor the overall status of the network and adjust aspects of the architecture in order to optimise use of the available resources. In particular, the mediation tasks represent segments of an offered mediation service, and the autonomic control plain is therefore effective to distribute the overall mediation service across the available M modules such that no one M module is overloaded. The autonomic control plane may also manage traffic levels on the MR and TPP modules by altering the destination of traffic from the M and LPP modules respectively.

The autonomic control plane may also be capable of increasing and decreasing the number of at least the cross-stream modules (the Ms and MRs), and preferably those of the TPs and LPPs, in order to ensure that there is sufficient capacity in the system. Preferably, the autonomic control plane may cause traffic to be directed through the LPP, MR, and TP modules in such a way as to ensure that no one module is overloaded. As stated above, traffic passes through the M modules in dependence on the mediation segment service required.

Preferably, the autonomic control plane has a hierarchical structure, with a single autonomic manager for each type of network module, known as cloud managers, and an overall distributed mediation network manager that acts as an autonomic manager to the cloud managers. In some examples, the cloud managers may contain further divisions, for example into geographic regions. One advantage of such a structure is that the cloud manager optimising the use of the M modules may act independently other types of network module. For instance, it may be that a first business entity responsible for an LPP employs a second business entity, responsible for the Ms, to provide a mediation service, and the two business entities are hesitant to allow access to each other's resources for security reasons. The preferred hierarchical structure of the autonomic control plane is sufficient to securely separate these two tasks.

The distributed mediation network manager is responsible for allocating resources to the cloud managers but need not be aware of the purpose for which these resources are intended.

The distributed mediation architecture of the present invention is capable of autonomically load balancing a mediation service across a number of servers, thereby enabling the resources consumed by a mediated application to be dynamically adjusted. This is achieved without any breaks in service while maintaining causal delivery of messages to and from the mediated application.

In accordance with the above, a number of key benefits are offered by the present invention:

- services are independent of location
- the architecture is scalable in that no node requires global knowledge of the system
- the architecture is dynamic – services may be redeployed, and new nodes may be transparently added or removed as required, and
- the architecture is autonomic – operational policies automatically optimise the size of the resource pool, the distributed mediation network topology, and the distribution of services.

Brief Description of the Drawings

Examples of the present invention will now be described in detail with reference to the accompanying drawings, in which:

Figure 1A is a node diagram showing a prior art mediated information flow system with a star-shaped logical architecture;

Figure 1B is a schematic node diagram showing the physical architecture of the system in Figure 1A;

Figure 2A is a node diagram showing a prior art mediated information flow system with a central network logical architecture;

Figure 2B is a schematic node diagram showing the physical architecture of the system in Figure 2A;

Figure 3 shows a minimal node diagram showing the fundamental cycle (LPP → MR → M → TP → LPP) of which every effective distributed mediation network node is a part;

Figure 4 is a node diagram showing a "cubic" distributed mediation model in accordance with the present invention;

Figure 5 illustrates the responsibilities of an autonomic control plane in accordance with the present invention;

Figure 6 illustrates the hierarchical structure of the autonomic control plane of the preferred embodiment of the present invention;

Figures 7A to 7F show the steps of handing a mediation task over from a sending mediator module *m2* to a recipient mediator node *m1* in the "cubic" distributed mediation architecture;

Figure 8 illustrates the addition of mediator node to the "cubic" distributed mediation model in Figure 4;

Figure 9 is a node diagram that is topologically equivalent to the "cubic" distributed mediation model in Figure 4. Using this so-called "cylindrical" layout it is easier to discuss other changes to an arbitrary distributed mediation network;

Figures 10A to 10E show the steps of switching an LPP node so that it is associated with a new MR node, given a starting point where at least two LPP nodes are sharing an MR node;

Figures 11A to 11D show the steps of switching a mediator node so that it is associated with a new TP node; and,

Figures 12A to 12D illustrate the addition of MR and TP nodes.

Detailed Description

The P2P and centrally mediated messaging models that provide the background to the present invention are first explained. Throughout this discussion, the term *source* designates a client that generates new messages to send into a network service and the term *sink* designates a client that receives messages from a network service. Each client of a network service may be a source, a sink, or both. In an alternative terminology, sources of information are referred to as *publishers* and sinks for information are referred to as *subscribers*.

In peer-to-peer content-based routing, a network is configured to allow the efficient transmission of messages from source to sink, based upon the establishment of virtual channels between the appropriate sources and sinks. Efficiency is typically achieved by the detection and removal of unnecessary edges (lines of communication connecting nodes) from a fully connected graph, with the resulting optimised graph then tailored to available network infrastructure. To establish a P2P virtual channel requires an expression of interest from one peer and an acceptance of that interest by the other peer.

On the other hand, in centrally mediated models, all messages are transmitted via a central mediation node (see Figures 1A and 1B). In the parlance of mediation networks, a mediation service is the general term applied to some computation that is applied to all incoming messages; the mediation requirement for any particular instance of a mediated architecture refers to the collation of all mediation services provided therein; a mediation authority is a person or persons providing such a mediation service; a mediation network is the network of physical computational entities (machines) under the control of the mediation authority; and a mediation server, the physical machine hosting one or more mediation services.

In a simplified model of a generalised mediated information flow system, messages sent to the mediation authority may belong to one of the following types: new information, emanating from a process acting as an information source; queries about state held by the mediation authority, requesting an immediate reply; and expressions of interest, essentially persistent queries requiring ongoing replies whenever pertinent new information is received by the mediation authority.

It is worth remarking that even in a fully mediated model, expressions of interest may still be significant, especially in the delivery of this network service to sinks, where they can reduce the bandwidth requirement on each virtual channel.

In the light of the above definitions, a mediated information flow system is one that consists of messages, containing information, being sent to and from a central authority. Actions taken by this authority may include time-ordered logging of received messages, computation over the sum of messages received up to a point, and dissemination of incoming messages among other clients, possibly after performing some processing based on the message content.

The present invention represents a hybrid of the content-based, decentralised P2P model and the simple, centrally-mediated network model. Rather than provide a single, central mediator, the various mediation services are dispersed across a mediation network comprising a number of separate functional components. In the hybrid model, expressions of intent are used to open virtual channels between source nodes and mediator nodes, and expressions of interest are used to open virtual channels between the mediator nodes and sink nodes. Messages received by sink nodes are therefore governed by expressions of interest registered with the mediated service. The latency

between source and sink nodes is necessarily greater than in simple content-based routing, as there are two or more logical hops involved. Provided the context allows it, the latency in each logical hop can be successively reduced as more static information becomes available. Relative to the simple mediated model, the mediation task is more complex, because it is spread over multiple nodes (see Figure 2A). However, the inherent central bottleneck of the centrally mediated model has been removed and the resulting architecture is scalable.

Figure 3 shows a minimal topology that illustrates the functional components of a (hybrid) distributed mediation model that is in a "quiescent" or "steady state", i.e. a state in which there is no provision for changes in the logical topology. The Figure shows how data flows around the system among the various component nodes.

All the component nodes necessary for a distributed mediation network of unlimited scale are present in this minimal topology, including: sources, sinks, local points of presence, mediator routers, mediators and transmission proxies. Throughout the following discussion, these terms and others listed below take a definition as set out below:

Local Point of Presence (LPP): a local point of presence acts as an intermediary between clients (sources and sinks) and the rest of the distributed mediation network. Local points of presence are network nodes that provide proxies for mediation services for a particular geographical region. Each client of a mediation service will thus communicate with only a single local point of presence, and with no other nodes within the mediation architecture. There may be an arbitrary number of LPPs in a system, each serving a respective number of clients.

Mediation Router (MR): a mediation router is a network node, incorporating a mediation router module that analyses the content of incoming messages and routes them to one of a number of cross-stream mediator nodes. Each mediation router sits at the head of an upstream network and receives messages from a number of LPPs. A mediation router may also log incoming messages to allow, for example, a local service within the geographical area it serves.

Mediator (M): a mediator is a network node incorporating a mediator module, which services the mediation requirement. Each mediator has an associated downstream

distribution, which is used to pass relevant messages to LPPs, and therefore, ultimately, to the sinks. Each mediator module implements one or more mediation tasks, each task representing a single mediation segment service to be applied to a particular type of message segment. Mediator modules may be configured to log all incoming messages they receive and to forward these messages to the associated downstream transmission network. Mediation tasks may then include servicing queries over message logs thus generated.

Transmission Proxy (TP): a transmission proxy is a network node, incorporating a transmission proxy module, that analyses messages output by one or more mediator nodes; determines, from registered expressions of interest, to which sink(s) the outgoing message is directed; and forwards messages on the downstream network associated with each mediator.

As explained in detail later in this description, the preferred embodiment of the present invention further comprises an autonomic control plane in addition to the network modules (nodes), thereby enabling autonomic control of the system.

The upstream network (from source to mediator router) is seen to be mediated but not content-based. Routing between mediator routers and mediators, in the so-called *cross-stream*, is content-based. The downstream network, too, requires content-based routing: indeed message routing between mediator routers and LPPs can be regarded as a hybrid content-based delivery mechanism in its own right. The partitioning of the message space as a part of this hybrid allows the introduction of a mid-stream mediation service to the publish and subscribe model without introducing a non-scalable central bottleneck.

For a "quiescent" or "steady state" system, the following statements relating to the distributed mediation network are always true. These statements may be considered "global invariants" of the distributed mediation architecture.

- Every node is part of a cycle LPP → MR → M → TP → LPP
- Every LPP addresses a single MR
- Every MR may address any arbitrary M
- Every M addresses a single TP
- Every TP may address any arbitrary LPP

The symbol "→" used above represents a unidirectional connection (a directed edge).

From inspection of the "global invariants" above, it is clear why the network illustrated in Figure 3 is considered to be the minimal distributed mediation topology: it consists of one LPP, one MR, and one TP, with two Ms over which the mediation is distributed. The nodes are configured in a simple cycle, with unidirectional connections arranged therebetween: LPP → MR → M → TP → LPP.

A more realistic, and complex, distributed mediation network is illustrated in Figure 4. Here, two of each type of node are present in a configuration known hereafter as a "cubic" network. The cubic network illustrates further properties of general distributed mediation networks.

As in the cyclic network, the message flow between connected nodes is unidirectional. Every node in the cubic network is a component of at least one cycle, LPP → MR → M → TP → LPP. The cubic network exhibits a "fan in/fan out" topology: while every LPP sends each message to precisely one MR, each MR may be addressed by a plurality of LPPs (two in Figure 4) – *fan-in*; every MR is capable of sending a message to any mediator – *fan-out*; every mediator sends any given message to precisely one TP, while each TP may be addressed by a plurality of mediators – *fan-in*; and finally, every TP is capable of sending a message to any LPP – *fan-out*.

Distributed mediation networks such as the minimal cyclic and the cubic networks also display another important property: a directed path exists from every node to every other node in the network. In graph theoretical terms, every node is in the transitive closure of every other node. This property holds trivially for cyclic networks: it does however hold in more complex distributed mediation networks, as consequence of the global invariant properties. The directed path can always be considered as a directed cyclic graph. Thus for any two nodes A and B within a generalised distributed mediation network, a cycle exists from A to A which contains B. It is worth noting that in a cubic (single-level) network the maximum path length of such a cycle is 8, rather than 4.

Again, each node has no global dependency or any detailed knowledge of the network beyond its immediate neighbours. Each node stores identities only of those nodes it directly addresses. Nodes may also store information about the totality of nodes which address each node within the network; this may be stored either as a reference count

in the node itself, or as a credit balance value in all nodes that directly address it. In either case the identities of the addressing nodes need not be stored. Indeed, no node need store any further information about the rest of the global system.

A further important property of the distributed mediation network of the present invention it that behaviour is deterministic and ordered within each node; that is, it is possible to ensure that a message B arriving after a message A is also dispatched after message A. Likewise, messages can not overtake on direct links between nodes; thus, for any two nodes N1 and N2 such that N2 is directly addressed by N1, then if a message A is sent to N2 by N1, and subsequently a message B is sent to N2 from N1, then N2 will receive message A before receiving message B.

Distributed Mediation Network Applications

An application *A* is structured according to a message oriented paradigm. *A* receives a partially ordered sequence of messages *m1*, *m2*, *m3*, ... drawn from a set *M*, and responds by generating further messages. The observable behaviour of application *A* (denoted by *obs(A)*) is defined as the set of all possible output sequences derived from a particular input sequence.

For such an application to benefit from being hosted by the distributed mediation network, it must follow the following principles:

- the application must parallelise into a set of application components $\{A_i\}$, each of which operate independently, but where all of the resulting output messages may be interleaved arbitrarily such that the observable behaviour of $\{A_i\}$ is equal to the observable behaviour of *A*,
- the set *M* of all possible input messages must be partitioned into a set of message segments, $\{S_j\}$, according to a function *segment* : ($M \rightarrow S$),
- a relation *mediates* : ($S \rightarrow A$) exists between message segments and application components such that each segment maps to precisely one application component, and
- the observable behaviour of the sum of all application components ΣA_i in response to a sequence of messages *m1*, *m2*, *m3*, ... where each component A_i is passed only those messages filtered from the input stream according to the functions *mediates* and *segment*, is precisely the same as the observable behaviour of *A* on receipt of the entire input stream.

Moreover, given the above constraints, to register with the load-balancing functionality offered by the distributed mediation network, two further side-effecting methods require to be added to instances of application components A_i:

- *gainSegment*(*segDescriptor*, *data*), and
- *loseSegment*(*segDescriptor* → *data*)

such that, for any two application components A_1 and A_2, and any segment *s*, *obs*(A_1 + A_2) in a context where *mediates*(*s*) = A_2 is equal to *obs*(A_1. *gainSegment*(*s*, *d*) + A_2.*loseSegment*(*s*))in a context where *mediates*(*s*) = A_1, where *d* is the result returned by A_2.*loseSegment*(*s*).

The Autonomic Control Plane

The distributed mediation architecture of the present invention is designed to function autonomically. That is, the system will optimise itself towards current requirements without the need for user input or knowledge of this optimisation. In particular, an autonomic system will take steps to ensure that the load on the available components is distributed effectively, thereby avoiding unnecessary bottlenecks in the handling of data in the system.

In the present invention, autonomic functionality is provided by the presence of an autonomic control plane. As illustrated in Figure 5, the autonomic control plane manages the mediation services, the distributed mediation network and the underlying resources (be they virtual or actual). Each network node presents itself as a managed element to the autonomic control plane. A managed element supports both a sensor and an effector interface. The sensor interface emits specified metrics, allowing the autonomic manager to monitor certain attributes of the managed element. The effector interface receives, from the autonomic manager, specified operations to change the behaviour of the managed element. As such, each of the sensor interface and the effector interface allows a unidirectional flow of information, the former functioning from managed element to autonomic manager, and the latter functioning from autonomic manager to managed element. In the preferred embodiment, the autonomic control plane consists of a hierarchical set of autonomic managers, as shown in Figure 6.

In the example shown in Figure 6, there is a single autonomic manager for each type of network node (LPP, MR, M, and TP). Managers at this hierarchical level are referred to as cloud managers. Each cloud manager is responsible for a particular load balancing function. In a preferred embodiment, each cloud manager is itself distributed

across a peer-to-peer network, and receives the sensor events from each network module of the type for which it is responsible.

At the next hierarchical level, an overall distributed mediation network (DMN) manager 620 retains control of the cloud managers, ensuring that they act as a coherent unit. As such, the cloud managers present themselves as managed elements of the DMN manager 620. The DMN manager 620 is responsible for ensuring that the resources available to each cloud manager are sufficient to perform the relevant load-balancing task, while the cloud managers are responsible for relinquishing control of any resources that they do not currently require. The DMN manager 620 resolves any resource conflicts between the cloud managers.

In the example shown in Figure 6, the MR cloud autonomic manger 612 is responsible for upstream load-balancing, that is for ensuring that no MR is overloaded with traffic. The MR cloud manager 612 is capable of increasing and decreasing the number of MRs in the network, and will do so such that the average throughput through the MRs is within a specified optimum range. Moreover, the MR cloud manager 612 will optimise the distribution of LPPs across the MRs such that no individual MR is overloaded. In the preferred embodiment, the actual instruction of a LPP to transfer is output from one MR to another is performed by the LPP cloud manager 610.

The M cloud autonomic manager 614 is responsible for cross-stream load-balancing. As such, the M cloud manager 614 ensures that there are sufficient Ms to handle the throughput routed through the MRs. To achieve this, the M cloud manager 614 will adjust the number of Ms such that the average load on each M module is within a specified range. The M cloud manager 614 is also responsible for distributing the mediation segment services amongst the Ms such that no individual M is overloaded. In order to do this under a range of conditions, the M cloud manager 614 is capable of transferring mediation segments or tasks between the Ms. The algorithm by which this is achieved is discussed in greater detail with respect to the example given below.

The TP cloud autonomic manager 616 is responsible for downstream load balancing. As such the TP cloud manager 616 ensures that the average throughput for each TP lies within a specified range, and that the Ms are distributed across the TPs such that no individual TP is overloaded. In the preferred embodiment, the actual switching of an M from one TP to another is delegated to the M cloud manager 614.

Although the load balancing policies enacted by the cloud managers follow the same basic pattern, it should be noted that switching a mediation segment service is a compound operation consisting of migrating the processing of the segment, including the transfer of any associated state; and updating the routing function used by each MR to determine the target of a given message, rerouting messages that are received by the old mediator in the meantime. It is for this reason that it is the actions of the M cloud manager that are discussed in more detail below.

In general, the load balancing policies should be allowed to operate independently with any resource conflicts resolved by a higher level service (the distributed mediation network manager). However, any switching initiated must be choreographed such that the causal delivery of messages from each client to the appropriate mediation segment service and the causal delivery of messages generated by each mediation service to interested clients are both guaranteed at all times.

Cross Stream Load-balancing – Mediation Segment Service Handover

The distributed mediation architecture of the present invention is based on an arbitrary topology of nodes: LPP, MR, M and TP. This topology of nodes has the properties described above in relation to the "steady state". This same topology is eminently suitable for the dynamic balancing of loads amongst existing functional components.

As described above, whenever a particular segment is deemed to be under heavy load, the hosting of the associated mediation task (mediation segment service) may be autonomically moved (by the M cloud manager) to a machine within the network that has spare capacity. The handover of the mediation task also requires dynamic adjustment of the mediation network to ensure that any messages, either currently within the network, or to be received in the future, are diverted to the new mediator node. This may be achieved by the propagation of special messages around the new logical network topology, i.e. a "causal rippling" through the appropriate machines. Mediation change can thereby occur within a live system without affecting the observable behaviour of that system in terms of message input and output. Recall that the mediation segments or tasks together form a single mediation applications, and that the various segments are distributed across the Ms.

Mediation change is ultimately possible by virtue of the global invariant property whereby an incoming message will be routed to the same mediator regardless of the LPP from which it emanates. Mediation change poses the problem of changing from one consistent state to another, within a live system, and without adversely affecting the system's correctness or performance. Two main functions of the distributed mediation network must be considered: message propagation and querying, and in particular start-up queries.

As mentioned previously, the M cloud manager is responsible for the autonomic load balancing of mediation requirements. As such, it is capable of both introducing new Ms to the system (should sufficient resources be available) and distributing the segments amongst the Ms such that no individual M is overloaded.

Moreover, the M cloud manager is capable of handling the addition or subtraction of mediation segments from a given mediation service. For example, the M cloud manager is capable of identifying a suitable M on which to host a new mediation segment.

For example, consider a system in which ten clients, each generating 15 requests/second, are attached to each of two LPPs (hereinafter referred to as *lpp1 and lpp2*). Moreover, the requests are associated with three distinct segments (hereinafter *s1*, *s2*, and *s3*) in equal proportion. There is a total upstream traffic of 300 requests/second shared equally between *lpp1* and *lpp2*. In turn this generates cross-stream traffic of 300 requests/second that is shared equally (in this simple example) between two MRs (*mr1* and *mr2*) and shared equally between segments *s1*, *s2*, and *s3* (at a rate of 100 requests/second for each segment).

As mentioned previously, the mediators handle traffic according to its segment. In this example, there are two mediators (hereinafter *m1* and *m2*), each capable of handling 200 requests/second. As such, one possible scenario is that *m1* mediates traffic of segment *s1*, while *m2* mediates the remaining traffic (segments *s2* and *s3*). Clearly, the resultant load on *m1* will be 100 requests/second while that on *m2* will be 200 requests/second.

In the simple example given here, the mediation service transmits the current state of the traffic on to the downstream network, and TP *tp1* is associated with *m1* while TP

tp2 is associated with *m2*. As such, *tp1* will be supplied with 100 updates/second while *tp2* will be supplied with 200 updates/second. These will in turn be passed on to *lpp1* and *lpp2*, each of which will receive 300 updates/second from *tp1* and *tp2* combined, representing the sum of the updates received by the clients in the system.

Consider a change in the conditions under which the system of the example above operates. For example, segment *s2* becomes of more interest to the clients while interest in segment *s1* declines, and this is reflected in a change in the number of requests associated with these segments from each client. For the purposes of illustration, assume that each client now transmits 3 requests/second associated with *s1* and 7 requests/second associated with *s2*. The combined number of requests/second for s1 and s2 in the system remains the same (at 200) but there are now only 60 requests/second associated with *s1*, while there are 140 requests/second associated with *s2*. Accordingly, the load on *m1* is now only 60 requests/second, while that on *m2* has risen to 240 requests/second. However, as mentioned previously, *m2* has a limited capacity of 200 requests/second.

Accordingly, the M cloud manager is required to autonomically act to rectify the situation so that neither *m1* nor *m2* is under a load greater than that which it is capable of handling. In this example, the M cloud manager may act to switch the handling of segment *s3* from *m2* to *m1*. Once this has been done the overall load on *m1* and *m2* will be 160 request/second and 140 requests/second respectively. Note that a transfer of *s2* to *m1* would also have left the load on the Ms within acceptable limits, though *m1* would have had to function at maximum capacity.

The migration of a segment from one M to another must be handled in such a way as to maintain causal delivery and without the changes to the distributed mediation network being externally visible.

Consider the progression of states from a time in which the system is in a first consistent state PS1 to a new consistent state PS2. At time t_0, the process of changing to the new state PS2 commences. At some unknown time after this, t_1, the system is known to have changed to the new consistent state PS2. The time t_b when the actual change occurs is unknown but bounded by t_0 and t_1.

Between time t_0 and t_1 we define the system as being *unstable*, meaning that the currently operative apportionment of mediation tasks is not known globally. Each of the system functions however is unaffected, as at each point sufficient local knowledge is available to correctly handle the information flow.

Consider the example above, in which mediator *m1* initially mediates segment *s1* and mediator *m2* while initially mediates segments *s2* and *s3*. We now describe in detail the algorithm which permits the hot migration or handover of mediation services for segment *s3* from *m2* to *m1* in order to load balance the mediation of all three segments (*s1*, *s2*, and *s3*), without interrupting the mediation service from the user perspective.

Figures 7A to 7F illustrate the mediation change cycle. As shown in Figure 7A, the mediation change cycle is initiated by the M Cloud manager calling a **HANDOVER_SEGMENT(s3)** effector method on mediator *m2* instructing it to hand over segment *s3* to *m1*. As a result, *m2* enters the ***HANDOVER_SENDER(s3)*** state. In this state *m2* processes any buffered messages for *s3*. As shown in Figure 7B, once the *s3* buffer has been flushed it sends a downstream **MEDIATION_CHANGE(s3)** control message to all LPPs via its TP; and a serialized snapshot of its current state to *m1* in a **MEDIATION_HANDOVER(s3-state)** control message. From this point, *m2* ceases to mediate segment *s3* and any subsequent messages for *s3* are forwarded to mediator *m1* by *m2*.

When mediator *m1* receives the **MEDIATION_HANDOVER(s3-state)** control message, it enters the ***HANDOVER_RECEIVER(s3)*** state and initialises a mediation service for segment s3 using the state information received. Next, as shown in Figure 7C, *m1* sends a downstream **NEW_MEDIATOR(s3)** control message to all LPPs via its TP.

The **NEW_MEDIATOR(s3)** and **MEDIATION_CHANGE(s3)** control messages are used to ensure the causal delivery to clients of the output of the mediation task associated with segment *s3*. In particular, downstream messages from the old mediator *(m2)* must be delivered to clients before those from the new mediator (*m1*). To ensure this, *s3* messages received by an LPP from the new mediator are buffered if the **NEW_MEDIATOR(s3)** control message is received before its corresponding **MEDIATION_CHANGE(s3)** control message.

Similarly, on entering the ***HANDOVER_RECEIVER(s3)*** state, mediator *m1* immediately begins buffering any *s3* messages received direct from each MR until it is certain that there are no outstanding *s3* messages originating from that particular MR being re-routed from *m2*.

To establish this requires further interaction with the M Cloud manager via the control plane. Therefore, on entering the ***HANDOVER_RECEIVER(s3)*** state mediator *m1* signals this state change to the M Cloud manager by emitting a sensor event.

As shown in Figure 7D, when the M Cloud manager receives this event it calls an effector on the MR Cloud manager to instruct it to update the routing table for each of its MRs. The MR Cloud manager sends each MR an **UPDATE_ROUTING_TABLE (s3, m2 → m1)** control message. Although in this case the M Cloud manager communicates with the MR Cloud manager directly, in other embodiments the communication may happen in other ways. For example, in a more strictly hierarchical embodiment, the communication may be passed through the DMN manager (for example, the M cloud manager may issue a sensor event picked up by the DMN manager which then calls an effector method on the MR Cloud manager). In some embodiments, it may be impossible to communicate directly between Cloud managers. In general, it is contemplated that communication between cloud managers may occur both directly and through the DMN manager in all cases in which such communication is discussed hereinafter.

Once an MR has updated its routing table it sends a **RT_CHANGED(MR-id, s3)** control message to *m2* and routes all subsequent *s3* messages to *m1*. As shown in Figure 7E, this **RT_CHANGED** control message is forwarded from *m2* to *m1* which is the trigger for *m1* to flush any buffered messages for *s3* received direct from this particular MR. Once any buffered messages have been flushed, m1 also emits a **RT_CHANGED** sensor event which alerts the MR cloud manager via the control plane that this particular MR has been updated successfully.

When the MR cloud manager has received acknowledgements via the new mediator *m1* from all of the MRs, the mediation change process is effectively complete. As shown in Figure 7F, at this time the MR cloud manager can inform the M Cloud manager via the control plane that it has updated all its MRs and the M Cloud manager

can inform m1 and m2 that mediation change for *s3* is complete at which point they can both revert to **STABLE(s3)** state.

Cross Stream Scale-Out and Scale-Back

As would be clear to one skilled in the art, similar techniques can be utilised to add or remove M nodes from the system on demand. In the case of addition this is achieved simply by obtaining the resources required; creating an empty M node and connecting it to a TP; then adding segments to it over time. The result of adding two M modes to the "cubic" distributed mediation model is illustrated in Figure 8. In the case of removal, this is achieved by first migrating all segments to other M nodes; then disconnecting it from its TP; shutting it down and releasing resources associated with it.

Upstream Load-balancing – Switching Mediation Router

Figure 9 shows an alternative illustration of a distributed mediation network. The network shown in Figure 9 is topologically equivalent to the "cubic" network shown in Figure 4, though the layout adopted in this case is referred to hereinafter as a "cylindrical" layout. Each LPP is illustrated twice in this type of diagram, both at the start and end of the illustrated network transmission paths. The cylindrical layout provides an effective illustration of the opportunities for load balancing that exist within the distributed mediation network.

As described above, cross-stream load balancing ensures the load on each mediator node stays within an acceptable range. As such, the goal is to ensure that the traffic through any given mediator node is within appropriate high/low watermarks and the overall workload handled by the set of mediator nodes is also within acceptable high/low watermarks. Since the level of traffic is determined by the distribution of mediation segment services across the mediator nodes, a technique is provided to migrate or handover a mediation segment service from one mediator node to another.

The upstream load-balancing case can be summarized as follows: the goal is to ensure that the traffic through any given MR node is within appropriate high/low watermarks and the overall workload handled by the set of MR nodes is also within acceptable high/low watermarks. Since the traffic through any given MR node is determined by the throughput generated by the LPP nodes attached to it, the ability to switch the output of an LPP node from one MR node to another is key to the management of load on any given MR. The ability to switch LPPs between MRs allows

an MR to service multiple LPPs when traffic is light. This provides significant efficiency advantages over rigid systems where, for example, a dedicated MR is always allocated to each LPP.

Figure 10A shows an expanded cylindrical layout, consisting of 4 LPPs serviced by 2 MRs and 4 Ms serviced by 2 TPs. The arrows between the nodes in Figure 10A represent the initial flow of network traffic. Taking this as the starting point, we consider the case where *lpp1* starts generating significant throughput such that *mr1* risks being overloaded. We now describe in detail the algorithm which permits the hot switching of *lpp2* to *mr2* to load balance the upstream traffic through the two MR nodes.

Figures 10B to 10E illustrate the LPP switchover cycle. As shown in Figure 10B, the switchover cycle is initiated by the LPP Cloud manager calling a **SWITCH_MR(mr2)** effector method on node *lpp2* instructing it to switch from its current MR node *mr1* to *mr2*. At this point LPP node *lpp2* enters into the ***SWITCHOVER(mr2)*** state.

As shown in Figure 10C, *lpp2* responds by sending **MR_CHANGE(lpp2)** control message to *mr1* which responds by sending an **MR_CHANGE(lpp2, s<i>)** control message to the appropriate mediator for each segment *s<i>* in its routing table and simultaneously emitting an **MR_CHANGE(lpp2, s<i>)** sensor event which is detected by the MR Cloud manager and stored in a shared data space accessible to the M Cloud manager. In some alternative embodiments, each LPP may maintain a list of "active" segments and generate **MR_CHANGE(lpp2, s<i>)** control messages for these which are routed by *mr1*. However, such a variation requires the maintenance of state information at the LPPs that is not otherwise required.

As shown in Figure 10D, *lpp2* then switches to the new MR node *mr2* and sends it a **NEW_MR(lpp2)** control message. As above this results in the generation of **NEW_MR(lpp2, s<i>)** control messages for each segment *s<i>* in its routing table which are transmitted to the appropriate M nodes.

The **NEW_MR(lpp2, s<i>)** and **MR_CHANGE(lpp2, s<i>)** control messages received by the M nodes are used to ensure the causal processing of cross stream messages by the appropriate segment service. In particular, upstream messages sent

by *lpp2* via the old mediation router (*mr1*) must be processed before those sent via the new mediation router (*mr2*).

To ensure this, any *s<i>* messages originating at *lpp2* received by a mediator from the new mediation router are buffered if the **NEW_MR(lpp2, s<i>)** control message is received before its corresponding **MR_CHANGE(lpp2, s<i>)** control message. This buffering process in analogous to that undergone by the new mediator (*m2*) in the cross-stream load balancing described above, though in the case of upstream load balancing the buffering is keyed off the originating LPP whereas for cross-stream load balancing the buffering is keyed off the originating MR.

As shown in Figure 10E, once both control messages have been received and any buffer flushed an **NEW_MR(lpp2, s<i>)** sensor event is emitted by the M nodes involved and detected by the M Cloud manager. Once all the **NEW_MR** sensor events corresponding to the previously stored **MR_CHANGE** sensor events have been detected, the M Cloud manager informs the LPP Cloud manager that the switchover cycle is complete and the LPP Cloud manager in turn invokes an effector method on lpp2 switching it to ***STABLE(mr2)*** state.

In order to ensure optimum stability, it is envisaged that some embodiments of the present invention would enable communication between the LPP Cloud manager and the M Cloud manager to ensure that the switchover of LPPs between MRs would only be initiated when all Ms were in a ***STABLE*** state and that no segment handovers between Ms would be initiated while an LPP is in a ***SWITCHOVER*** state.

Downstream Load-balancing – Switching Transmission Proxy

The downstream load-balancing case mirrors the upstream case above and can be summarized as follows: the goal is to ensure that the traffic through any given TP node is within appropriate high/low watermarks and the overall workload handled by the set of TP nodes is also within acceptable high/low watermarks. Since the traffic through any given TP node is determined by the throughput generated by the M nodes attached to it, the ability to switch an M node from one TP node to another is key. The ability to switch Ms between TPs allows a TP to service multiple Ms when traffic is light. This provides significant efficiency advantages over rigid systems where, for example, a dedicated TP is always allocated to each M.

Taking as our starting point the distributed mediation network described in relation to upstream load balancing above (shown in Figure 10A), we consider the case where *m1* starts generating significant throughput such that *tp1* risks being overloaded. We now describe in detail the algorithm which permits the hot switching of *m2* to *tp2* to load balance the downstream traffic through the two TP nodes.

Figures 11A to 11D illustrate the mediator node switchover cycle. As shown in Figure 11A, this switchover cycle is initiated by the M Cloud manager calling a **SWITCH_TP(tp2)** effector method on mediator *m2* instructing it to switch from its current TP node *tp1* to *tp2*. At this point *m2* enters into the ***SWITCHOVER(tp2)*** state.

As shown in Figure 11B, *m2* responds by sending a **TP_CHANGE(m2, s<i>)** control message to *tp1* for each active segment *s<i>* it is hosting. Node *tp1* multicasts this message to all LPPs simultaneously emitting a **TP_CHANGE(m2, lpp<j>, s<i>)** sensor event corresponding to each LPP which is detected by the TP Cloud manager and stored in a shared data space accessible to the LPP Cloud manager.

As shown in Figure 11C, *m2* then switches to the new TP node *tp2* and sends it a **NEW_TP(m2, s<i>)** control message for each active segment *s<i>* it is hosting. Node *tp2* multicasts this message to all LPPs.

The **NEW_TP(m2, s<i>)** and **TP_CHANGE(m2, s<i>)** control messages are used to ensure the causal delivery of downstream messages generated by the appropriate segment service. In particular downstream messages sent by *m2* via the old transmission proxy (*tp1*) must be delivered before those sent via the new transmission proxy (*tp2*).

To ensure this, any *s<i>* messages originating at *m2* received by an LPP from the new transmission proxy are buffered if the **NEW_TP(m2, s<i>)** control message is received before its corresponding **TP_CHANGE(m2, s<i>)** control message. [Note: This buffering is similar to that done by the LPP during segment handover except that in the former case the buffering is keyed off the originating mediation node rather than the originating transmission proxy.]

As shown in Figure 11D, once both control messages have been received and any buffer flushed an **NEW_TP(m2, lpp<j>, s<i>)** sensor event is emitted by the LPP and

detected by the M Cloud manager. Once all the **NEW_TP** sensor events corresponding to the previously stored **TP_CHANGE** sensor events have been detected, the LPP Cloud manager informs the M Cloud manager that the switchover cycle is complete and the M Cloud manager in turn invokes an effector method on *m2* switching it to ***STABLE(tp2)*** state.

In some embodiments, the switchover of Ms between TPs would only be initiated by the M Cloud manager when all the Ms were in a ***STABLE*** state. Morevover, M Cloud manager may be designed not to initiate the transfer of mediation segments between while any M is in a ***SWITCHOVER*** state.

Upstream and Downstream Scale-Out and Scale-Back

As would be clear to one skilled in the art, similar techniques can be utilised to add or remove MR and TP nodes from the system on demand. In the case of addition this is achieved simply by obtaining the resources required; creating an MR or TP node and switching one or more LPP or M nodes to it respectively.

The net result of adding an MR node (*mr3*) to the network topology shown in Figures 10 and 11 and switching an LPP (*lpp4*) to it in response to the overall increase in workload across the existing MR nodes exceeding a threshold is illustrated in Figure 12A. It follows from this that in order to remove an MR node if the overall reduction in workload across the existing MR nodes drops below a threshold all LPPs pointing at it must first be switched to other MR nodes.

The net result of adding a TP node (*tp3*) to the resultant network topology and switching a mediator node (*m4*) to it in response to the overall increase in workload across the existing TP nodes exceeding a threshold is illustrated in Figures 12B and 12C. It follows from this that in order to remove a TP node if the overall reduction in workload across the existing TP nodes drops below a threshold all Ms pointing at it must first be switched to other TP nodes.

Finally, Figure 12D shows a maximal configuration for the 4 LPP/4 M node network shown in figures 10 and 11. In this case, each LPP has a dedicated MR and each M a dedicated TP. Given the constraints of the distributed mediation network it is not possible for the number of MRs to exceed the number of LPPs or the number of TPs to exceed the number of Ms in the network.

Exception Handling

Even with upstream, cross-stream and downstream load-balancing in place, there are some aspects of network load that cannot be dealt with by the methods described above. These exceptions are referred by the relevant cloud managers to the DMN for action.

M Cloud manager

Exception: An M node hosting single segment *s<i>* becomes overloaded i.e. exceeds the high watermark for an individual M node. The M Cloud manager should raise this exception to the DMN manager.

Action: None possible unless we can differentiate between M nodes i.e. have a mix of M node capabilities with some hosted on more powerful machines (e.g. the latest multi-core chipset versus last year's model) in which case the load balancing algorithm can be refined to move the segment to a more powerful node although we will still hit this edge condition at some point.

Observation: What this exception highlights is a potential bottleneck where weakest link is an M node.

TP Cloud Manager

Exception: A TP node supporting a single *m<j>* node becomes overloaded i.e. exceeds the high water mark for an individual M node. The TP Cloud manager should raise this exception to the DMN Manager.

Action: If we can differentiate between TP nodes i.e. have a mix of TP node capabilities then the load balancing algorithm can be refined although we will still hit this edge condition at some point. Otherwise the DMN manager can inform the M Cloud manager that *m<j>* node is overloading network. If *m<j>* node is hosting multiple segments then in theory these could be redistributed.

Observation: What this highlights is a potential bottleneck where the weakest link is a TP node.

MR Cloud Manager

Exception: An MR node supporting a single *lpp<k>* node becomes overloaded i.e. exceeds the high watermark for an individual MR node. The MR Cloud manager should raise this exception to the DMN Manager.

Action: None possible unless we can differentiate between MR nodes i.e. have a mix of MR node capabilities then the load balancing algorithm can be refined although we will still hit this edge condition at some point.
Observation: What this exception highlights is a potential bottleneck where weakest link is an MR node.

In some embodiments, the observation at the DMN manager that an individual LPP is overloading an MR may lead to a readjustment of users amongst LPPs. The technique by which this is achieved will depend on the nature of the mediation service being provided to the users.

Figure 13 shows Table 1, which provides a brief summary of some of the actions taken by the autonomic control plane to distribute network traffic amongst the nodes within the distributed mediation network. It shows the metrics or variables upon which the autonomic control plan makes decisions. The "X_NODE_HIGH_WATERMARK" represents the maximum threshold value that may be handled by a single type X node (where X may be LPP, TP, MR, or M), while the "X_NODE_POOL_HIGH_WATERMARK" represents the maximum average value of traffic throughput over a group of nodes of type X.

Claims

1. A distributed mediation network, comprising:
a plurality of types of network module, including:
local points of presence (LPP) modules for receiving and transmitting network traffic between the mediation network and client programs;
mediator (M) modules for hosting mediation tasks;
mediator router (MR) modules for analyzing the content of incoming messages, each MR module routing the incoming messages to a predetermined mediation task in dependence upon said content; and,
transmission proxy (TP) modules for forwarding messages to at least one of said LPP modules, wherein each of the MR, M and TP modules are adapted such that all paths for network traffic therethrough are non-reciprocal; and,
an autonomic control plane for managing the distribution of network traffic amongst the modules.

2. A network according to claim 1, wherein the network couples network traffic along a unidirectional mediation cycle, in which: LPP modules address MR modules, MR modules in turn address M modules, M modules in turn address TP modules, and TP modules in turn address LPP modules.

3. A network according to claim 1 or claim 2, wherein the autonomic control plane is adapted to effect the distribution of mediation tasks amongst the mediation modules.

4. A network according to any preceding claim, wherein the autonomic control plane is adapted to regulate the numbers of each type of module.

5. A network according to any preceding claim, wherein the autonomic control plane is hosted across a number of resources.

6. A network according to any preceding claim, wherein the modules contain sensor interfaces for communicating the status of the module to the autonomic control plane.

7. A network according to any preceding claim, wherein the modules contain effector interfaces for receiving commands from the autonomic control plane.

8. A network according to any preceding claim, wherein the autonomic control plane has a hierarchical structure comprising cloud managers for managing the distribution of network traffic amongst modules of a given type and a distributed mediation network (DMN) manager for managing the distribution of resources between the module types.

9. A network according to any preceding claim, wherein the distribution of mediation tasks comprises the transfer of a particular mediation task from a first M module to a second M module, the transfer comprising:

the autonomic control plane calling a HANDOVER_SEGMENT effector method on the first M module,

the first M module, on receipt of the HANDOVER_SEGMENT effector method, changing state to a HANDOVER_SENDER state, the first M module then processing content currently stored at the first M module relating to the particular mediation task and subsequently sending a MEDIATION_CHANGE control signal to all LPP modules,

sending a MEDIATION_CHANGE control signal from the first M module to the second M module, and forwarding content relating to the particular mediation task subsequently received by the first M module to the second M module

the second M module, on receipt of the MEDIATION_CHANGE control signal, changing state to a HANDOVER_RECEIVER state and sending a sensor signal to the autonomic control plane indicating this change of state, the second M module, then sending a NEW_MEDIATOR control signal to all LPP modules, wherein the LPP modules buffer content related to the particular mediation task that is received after the NEW_MEDIATOR control signal but prior to the MEDIATION CHANGE control signal;

the autonomic control plane, on receipt of the sensor signal indicating that the second M module is in a HANDOVER_RECEIVER state, calling an effector method on the MR modules to instruct MR modules to forward content relating to the particular mediation task to the second M module rather than the first M module;

each MR module, on changing the destination of forwarded content to the second M module, emitting a sensor signal to the autonomic control plane indicating the change and forwarding a RT_CHANGED control message for that MR module to the first M module, the first M module subsequently forwarding the RT_CHANGED control message to the second M module, wherein the second M module buffers

content relating to the particular mediation task received directly from each MR module until it receives the RT_CHANGED control message for that MR module,

the second M module emitting a sensor signal to the autonomic control plane when it receives each RT_CHANGED signal originating at each MR module

the autonomic control plane calling effectors methods on the first and second M modules to return them to a stable state once it has received sensor signals from the second M module indicating that RT_CHANGED control signals have been received from all MR modules.

10. A network according to any preceding claim, wherein the management of load on across the network comprises the alteration of the destination of network traffic from an LPP module from a first MR module to a second MR module, the alteration comprising:

the autonomic control plane calling a SWITCH_MR effector method on the LPP module;

the LPP module, on receipt of the SWITCH_MR effector method, changing state to a SWITCHOVER state, sending an MR_CHANGE control signal to the first MR module, and sending a NEW_MR control signal to a second MR module,

the first MR module, on receipt of the MR_CHANGE control signal, forwarding the MR_CHANGE signal to each M module and emitting an MR_CHANGE sensor event to the autonomic control plane for each M module to which the MR_CHANGE control signal has been sent

the second MR module, on receipt of the NEW_MR control signal, forwarding the NEW_MR control signal to each M module, wherein the M modules are adapted to buffer content received after the NEW_MR control signal but prior to the MR_CHANGE control signal, thereby ensuring that content is processed in the correct order

each M module sending a sensor event to the autonomic control plane on receipt of the NEW_MR control signal,

the autonomic control plane, upon receipt of a sensor event from all M modules for which an MR_CHANGE sensor event has been received from the first MR module, calling an effector method to return the LPP module to a STABLE state.

11. A network according to any preceding claim, wherein the management of load on across the network comprises the alteration of the destination of network traffic from an M module from a first TP module to a second TP module, the alteration comprising:

the autonomic control plane calling a SWITCH_TP effector method on the M module;

the M module, on receipt of the SWITCH_TP effector method, changing state to a SWITCHOVER state, sending a TP_CHANGE control signal to the first TP module, and sending a NEW_TP control signal to the second TP module,

the first TP module, on receipt of the TP_CHANGE control signal, forwarding the TP_CHANGE signal to each LPP module and emitting an TP_CHANGE sensor event to the autonomic control plane for each LPP to which the TP_CHANGE control signal has been sent

the second TP module, on receipt of the NEW_TP control signal, forwarding the NEW_TP control signal to each LPP module, wherein the LPP modules are adapted to buffer content received after the NEW_TP control signal but prior to the TP_CHANGE control signal, thereby ensuring that content is processed in the correct order

each LPP module sending a sensor event to the autonomic control plane on receipt of the NEW_TP control signal,

the autonomic control plane, upon receipt of a sensor event from all LPP modules for which a TP_CHANGE sensor event has been received from the first TP, calling an effector method to return the M module to a STABLE state.

12. A network according to any preceding claim, wherein the incoming network traffic to the mediation network belongs to one of the group of message types including:

new information, which emanates from a process acting as an information source;

queries about state of nodes in the mediation network, which require a reply; and

expressions of interest, which require ongoing replies whenever pertinent new information is received by the mediation network.

13. A network according to any preceding claim, wherein the cloud managers are able to communicate directly.

14. A network according to any preceding claim, wherein all communication between cloud managers must pass through the DMN manager.

15. A method for mediating the flow of network traffic in a computer network, wherein the computer network has a mediation network that comprises: a plurality of types of network module, including:

local points of presence (LPP) modules for receiving and transmitting network traffic between the mediation network and client programs;

mediator (M) modules for hosting mediation tasks;

mediator router (MR) modules for analyzing the content of incoming messages, each MR module routing the incoming messages to a predetermined mediation task in dependence upon said content; and,

transmission proxy (TP) modules for forwarding messages to at least one of said LPP modules, wherein each of the MR, M and TP modules are adapted such that all paths for network traffic therethrough are non-reciprocal; and,

an autonomic control plane for managing the distribution of network traffic amongst the modules,

wherein, in the method, incoming messages are propagated along a mediation cycle that comprises the steps of:

an LPP module addressing incoming messages to a respective one of said at least one mediator router (MR) modules;

at said addressed MR module, analyzing the content of incoming messages and routing said messages to a predetermined mediator module in dependence upon said analyzed content;

at said predetermined mediator module, applying the mediation task to said analyzed messages and directing mediated messages to a respective one of said TP modules; and

at said TP module that receives said mediated messages, forwarding said mediated messages to at least one of said LPP modules.

16. A method according to claim 15, wherein the autonomic control plane is adapted to effect the distribution of mediation tasks amongst the mediation modules.

17. A method according to claim 15 or 16, wherein the autonomic control plane is adapted to regulate the numbers of each type of module.

18. A method according to any of claim 15 to 17, wherein the autonomic control plane is hosted across a number of resources.

19. A method according to any of claim 15 to 18, wherein the modules contain sensor interfaces for communicating the status of the module to the autonomic control plane.

20. A method according to any of claim 15 to 19, wherein the modules contain effector interfaces for receiving commands from the autonomic control plane.

21. A method according to any of claim 15 to 20, wherein the autonomic control plane has a hierarchical structure comprising cloud managers for managing the distribution of network traffic amongst modules of a given type and a distributed mediation network (DMN) manager for managing the distribution of resources between the module types.

22. A method according to any of claim 15 to 21, wherein the distribution of mediation tasks comprises the transfer of a particular mediation task from a first M module to a second M module, the transfer comprising:

the autonomic control plane calling a HANDOVER_SEGMENT effector method on the first M module,

the first M module, on receipt of the HANDOVER_SEGMENT effector method, changing state to a HANDOVER_SENDER state, the first M module then processing content currently stored at the first M module relating to the particular mediation task and subsequently sending a MEDIATION_CHANGE control signal to all LPP modules,

sending a MEDIATION_CHANGE control signal from the first M module to the second M module, and forwarding content relating to the particular mediation task subsequently received by the first M module to the second M module

the second M module, on receipt of the MEDIATION_CHANGE control signal, changing state to a HANDOVER_RECEIVER state and sending a sensor signal to the autonomic control plane indicating this change of state, the second M module, then sending a NEW_MEDIATOR control signal to all LPP modules, wherein the LPP modules buffer content related to the particular mediation task that is received after the NEW_MEDIATOR control signal but prior to the MEDIATION CHANGE control signal;

the autonomic control plane, on receipt of the sensor signal indicating that the second M module is in a HANDOVER_RECEIVER state, calling an effector method on the MR modules to instruct MR modules to forward content relating to the particular mediation task to the second M module rather than the first M module;

each MR module, on changing the destination of forwarded content to the second M module, emitting a sensor signal to the autonomic control plane indicating the change and forwarding a RT_CHANGED control message for that MR module to the first M module, the first M module subsequently forwarding the RT_CHANGED control message to the second M module, wherein the second M module buffers content relating to the particular mediation task received directly from each MR module until it receives the RT_CHANGED control message for that MR module,

the second M module emitting a sensor signal to the autonomic control plane when it receives each RT_CHANGED signal originating at each MR module

the autonomic control plane calling effectors methods on the first and second M modules to return them to a stable state once it has received sensor signals from the second M module indicating that RT_CHANGED control signals have been received from all MR modules.

23. A method according to any of claim 15 to 22, wherein the management of load on across the network comprises the alteration of the destination of network traffic from an LPP module from a first MR module to a second MR module, the alteration comprising:

the autonomic control plane calling a SWITCH_MR effector method on the LPP module;

the LPP module, on receipt of the SWITCH_MR effector method, changing state to a SWITCHOVER state, sending an MR_CHANGE control signal to the first MR module, and sending a NEW_MR control signal to a second MR module,

the first MR module, on receipt of the MR_CHANGE control signal, forwarding the MR_CHANGE signal to each M module and emitting an MR_CHANGE sensor event to the autonomic control plane for each M module to which the MR_CHANGE control signal has been sent

the second MR module, on receipt of the NEW_MR control signal, forwarding the NEW_MR control signal to each M module, wherein the M modules are adapted to buffer content received after the NEW_MR control signal but prior to the MR_CHANGE control signal, thereby ensuring that content is processed in the correct order

each M module sending a sensor event to the autonomic control plane on receipt of the NEW_MR control signal,

the autonomic control plane, upon receipt of a sensor event from all M modules for which an MR_CHANGE sensor event has been received from the first MR module, calling an effector method to return the LPP module to a STABLE state.

24. A method according to any of claim 15 to 23, wherein the management of load on across the network comprises the alteration of the destination of network traffic from an M module from a first TP module to a second TP module, the alteration comprising:

the autonomic control plane calling a SWITCH_TP effector method on the M module;

the M module, on receipt of the SWITCH_TP effector method, changing state to a SWITCHOVER state, sending a TP_CHANGE control signal to the first TP module, and sending a NEW_TP control signal to the second TP module,

the first TP module, on receipt of the TP_CHANGE control signal, forwarding the TP_CHANGE signal to each LPP module and emitting an TP_CHANGE sensor event to the autonomic control plane for each LPP to which the TP_CHANGE control signal has been sent

the second TP module, on receipt of the NEW_TP control signal, forwarding the NEW_TP control signal to each LPP module, wherein the LPP modules are adapted to buffer content received after the NEW_TP control signal but prior to the TP_CHANGE control signal, thereby ensuring that content is processed in the correct order

each LPP module sending a sensor event to the autonomic control plane on receipt of the NEW_TP control signal,

the autonomic control plane, upon receipt of a sensor event from all LPP modules for which a TP_CHANGE sensor event has been received from the first TP, calling an effector method to return the M module to a STABLE state.

25. A network according to any of claim 15 to 24, wherein the incoming network traffic to the mediation network belongs to one of the group of message types including:

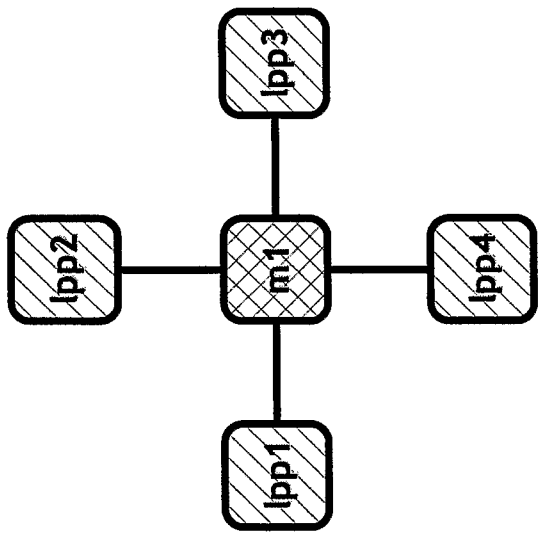
new information, which emanates from a process acting as an information source;

queries about state of nodes in the mediation network, which require a reply; and

expressions of interest, which require ongoing replies whenever pertinent new information is received by the mediation network.

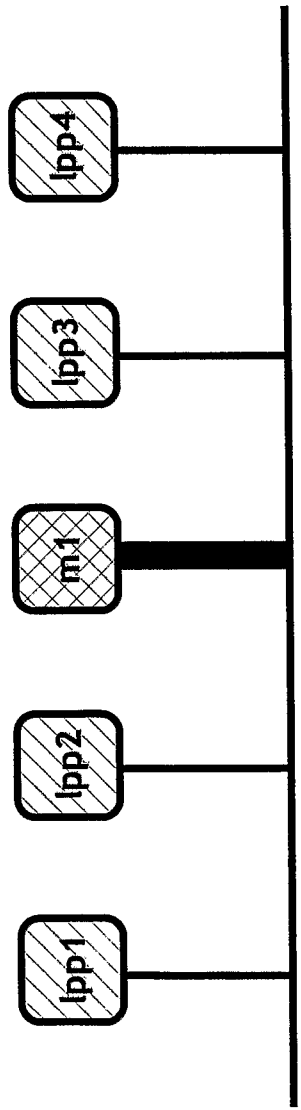
26. A network according to any of claim 15 to 25, wherein the cloud managers are able to communicate directly.

27. A network according to any of claim 15 to 26, wherein all communication between cloud managers must pass through the DMN manager.



Logical grouping with central mediation node

Figure 1A



Physical manifestation with heavy bandwidth requirement on central node

Figure 1B

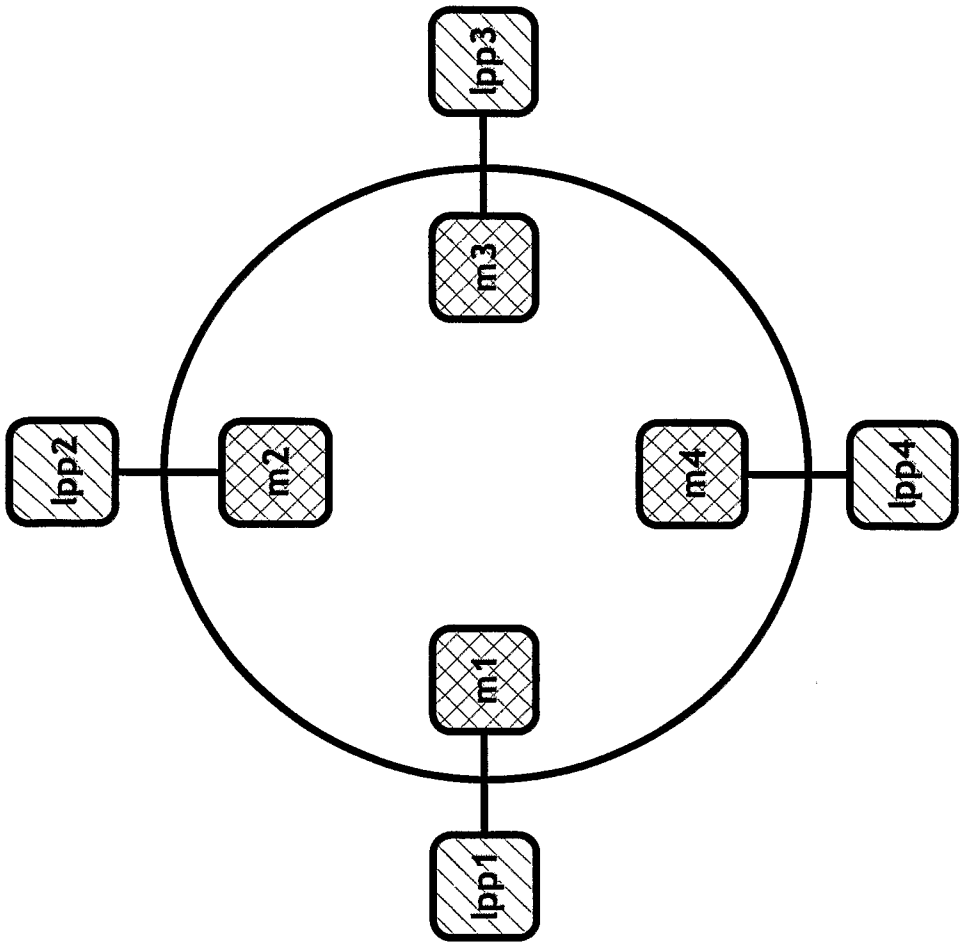
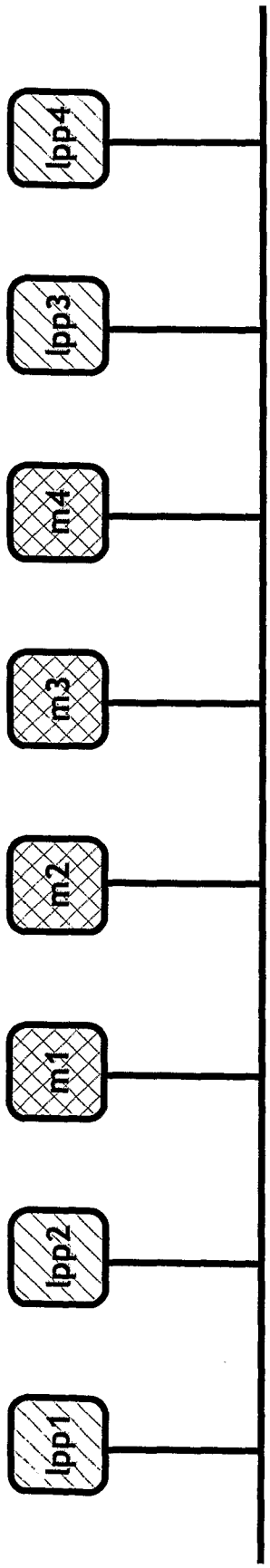


Figure 2A



Central network arrangement avoids heavy bandwidth requirement on any one node

Figure 2B

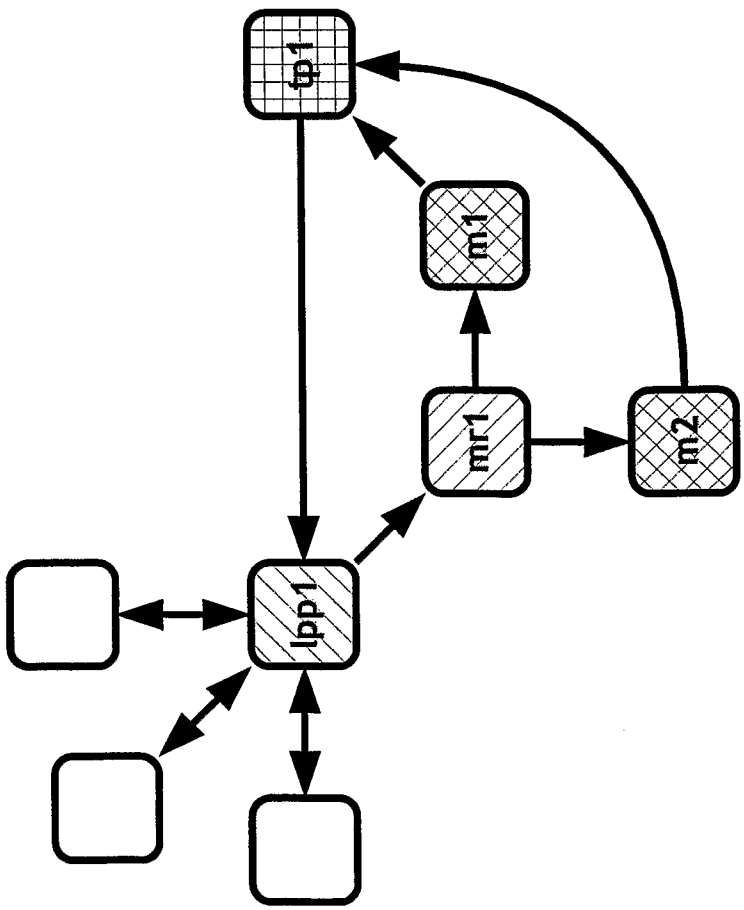


Figure 3

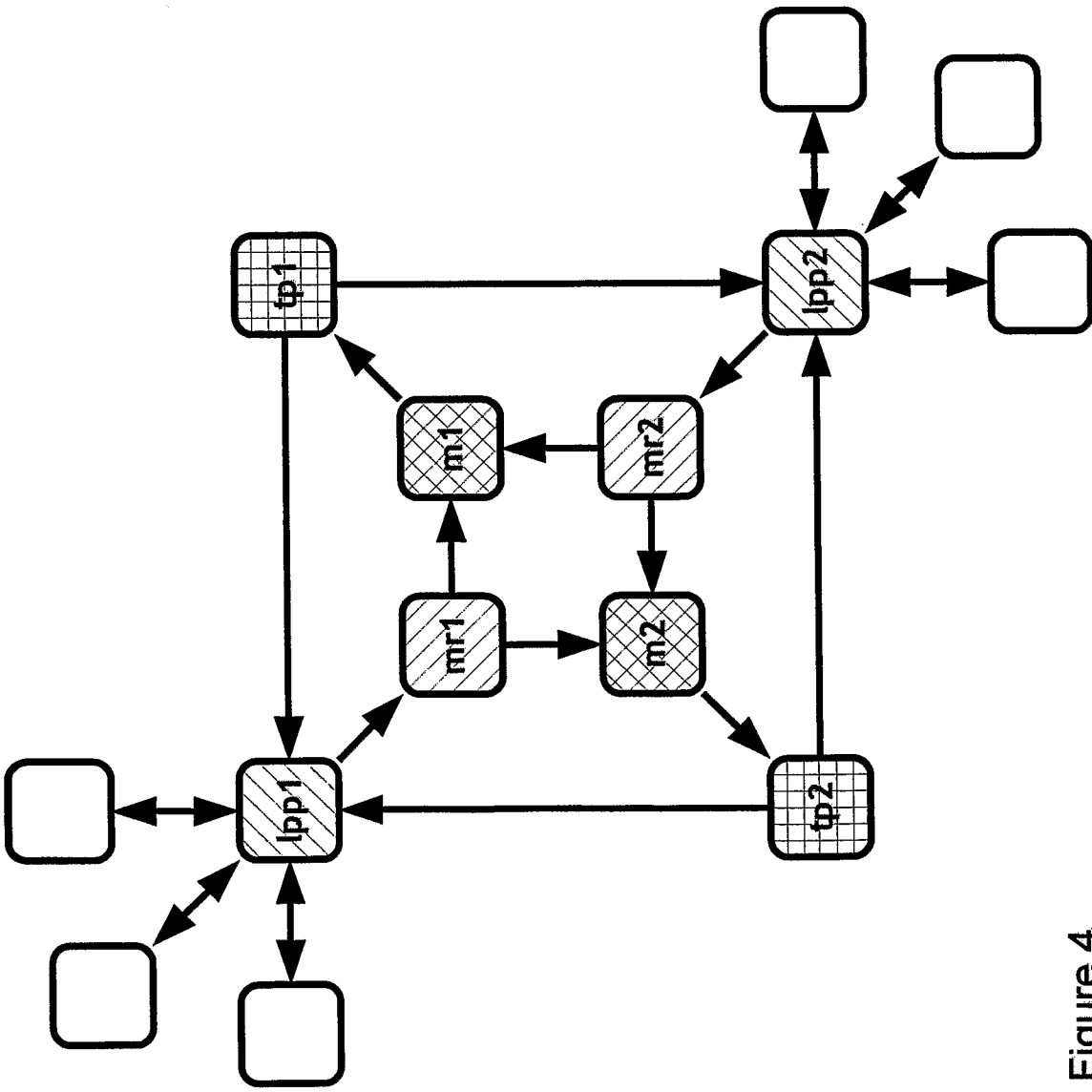


Figure 4

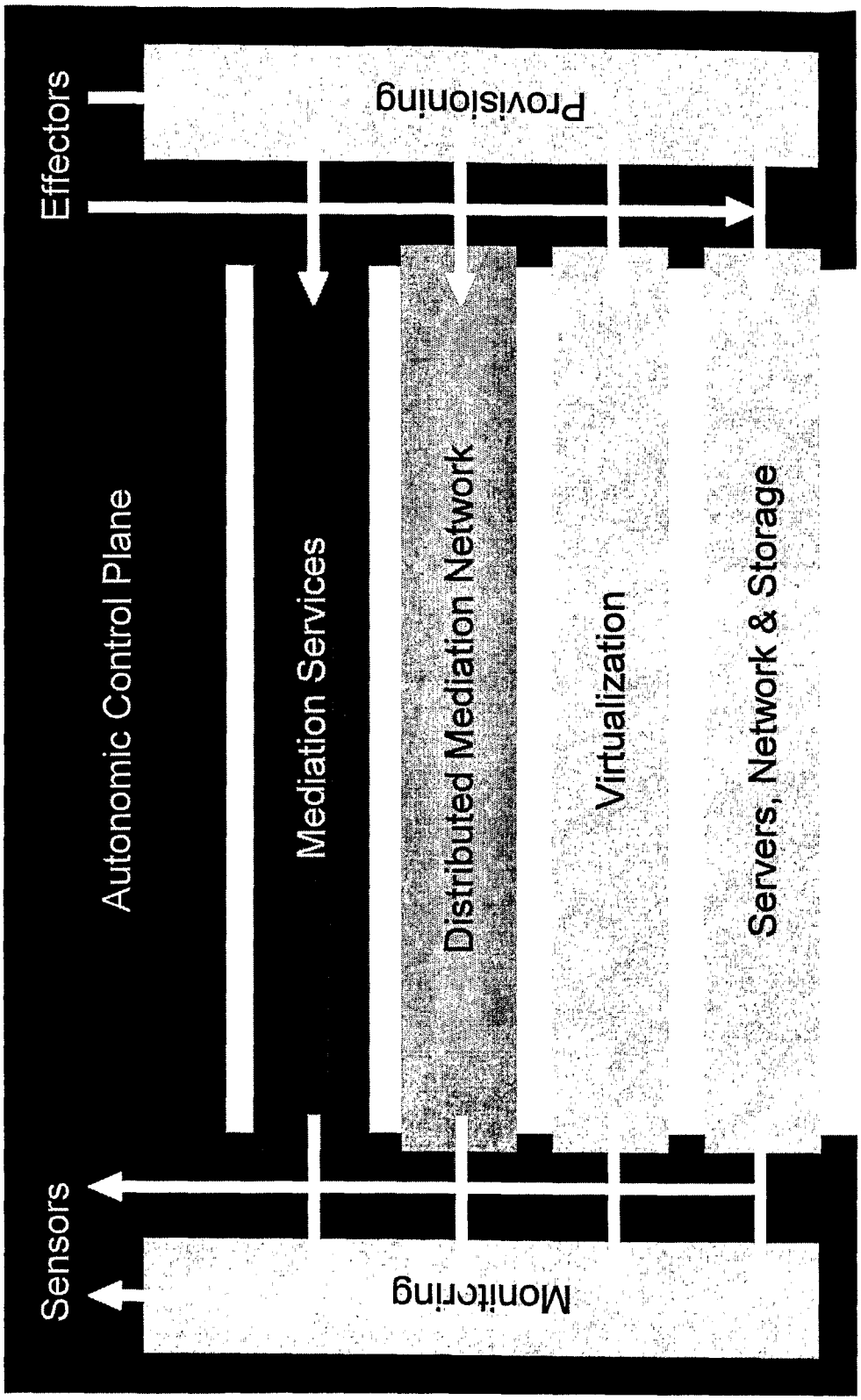

Effectors
Provisioning
Autonomic Control Plane
Mediation Services
Distributed Mediation Network
Virtualization
Servers, Network & Storage
Sensors
Monitoring

Figure 5

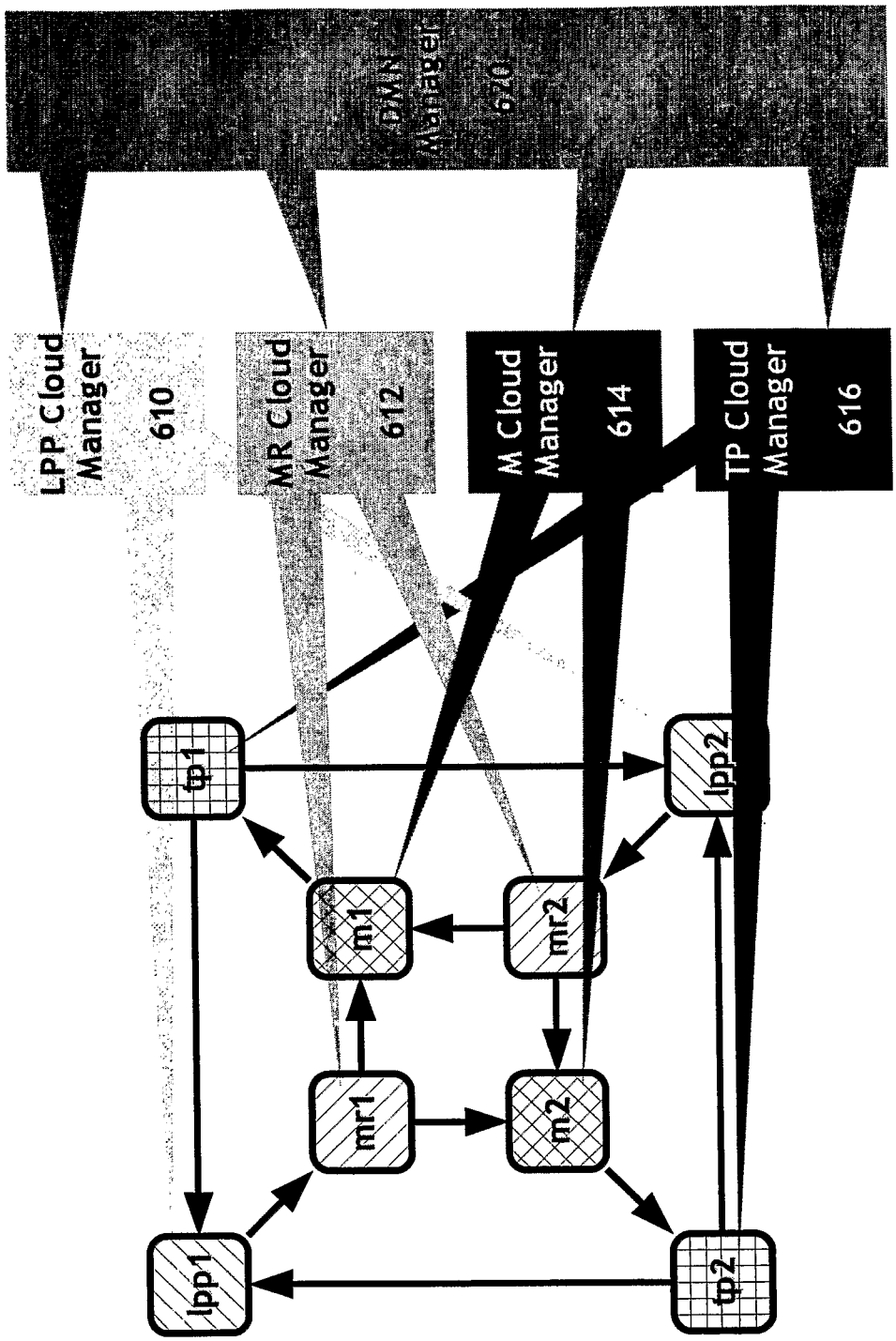

DMN Manager 620
LPP Cloud Manager 610
MR Cloud Manager 612
M Cloud Manager 614
TP Cloud Manager 616
tp1
lpp2
m1
mr2
mr1
m2
lpp1
tp2

Figure 6

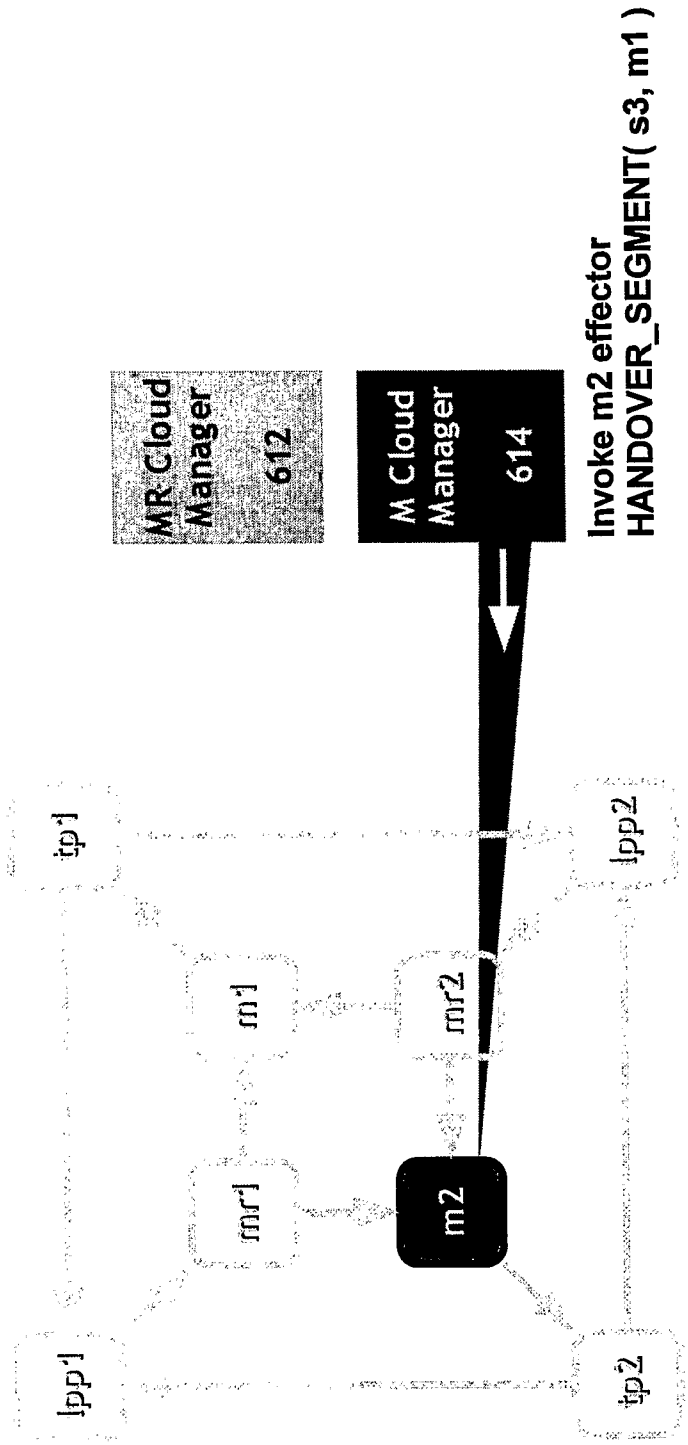

MR Cloud Manager 612
M Cloud Manager 614
Invoke m2 effector HANDOVER_SEGMENT(s3, m1)
tp1
lpp2
m1
mr2
mr1
m2
lpp1
tp2

Figure 7A

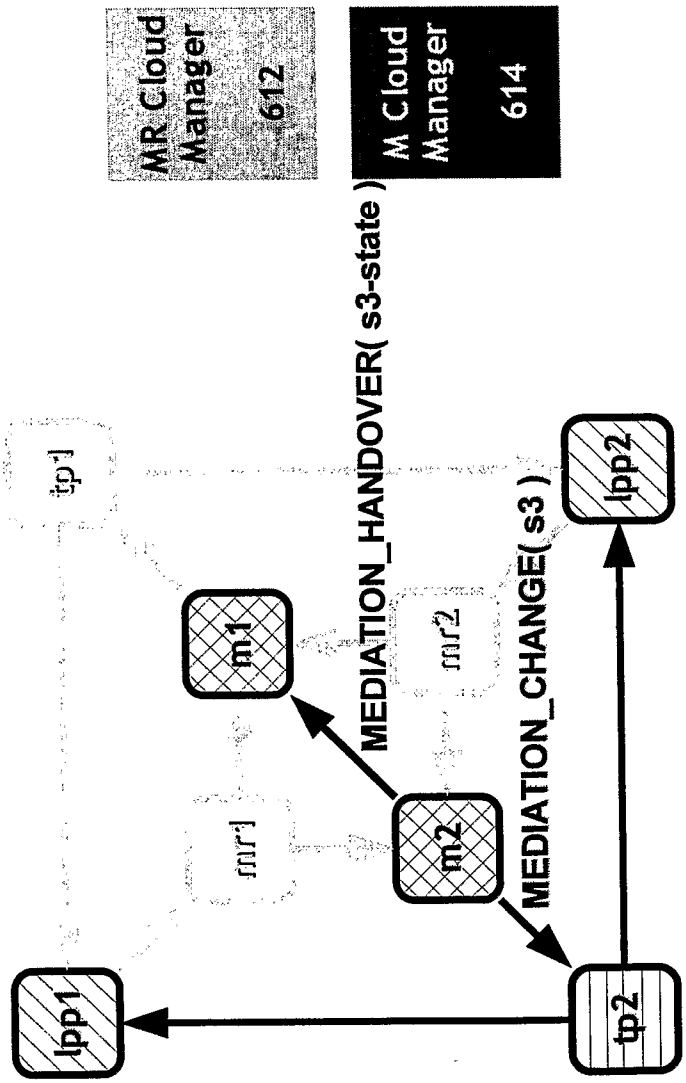

MR Cloud Manager 612
M Cloud Manager 614
MEDIATION_HANDOVER(s3-state)
MEDIATION_CHANGE(s3)
tp1
m1
mr2
mr1
m2
lpp2
lpp1
tp2

Figure 7B

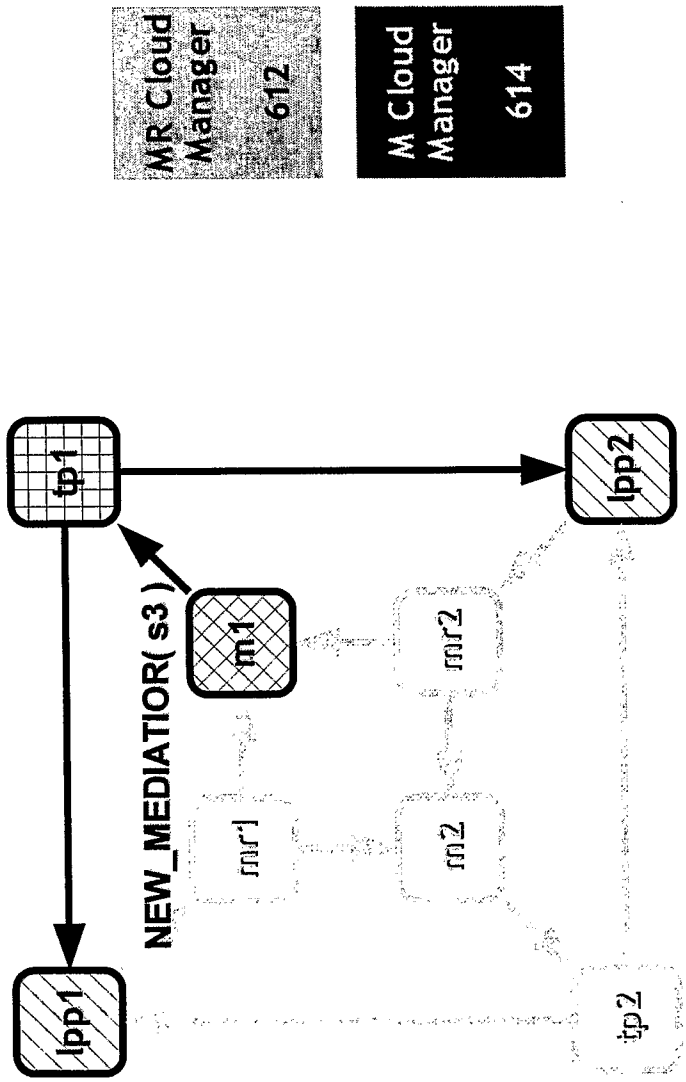

MR Cloud Manager 612
M Cloud Manager 614
tp1
lpp2
NEW_MEDIATIOR(s3)
m1
mr2
mr1
m2
lpp1
tp2

Figure 7C

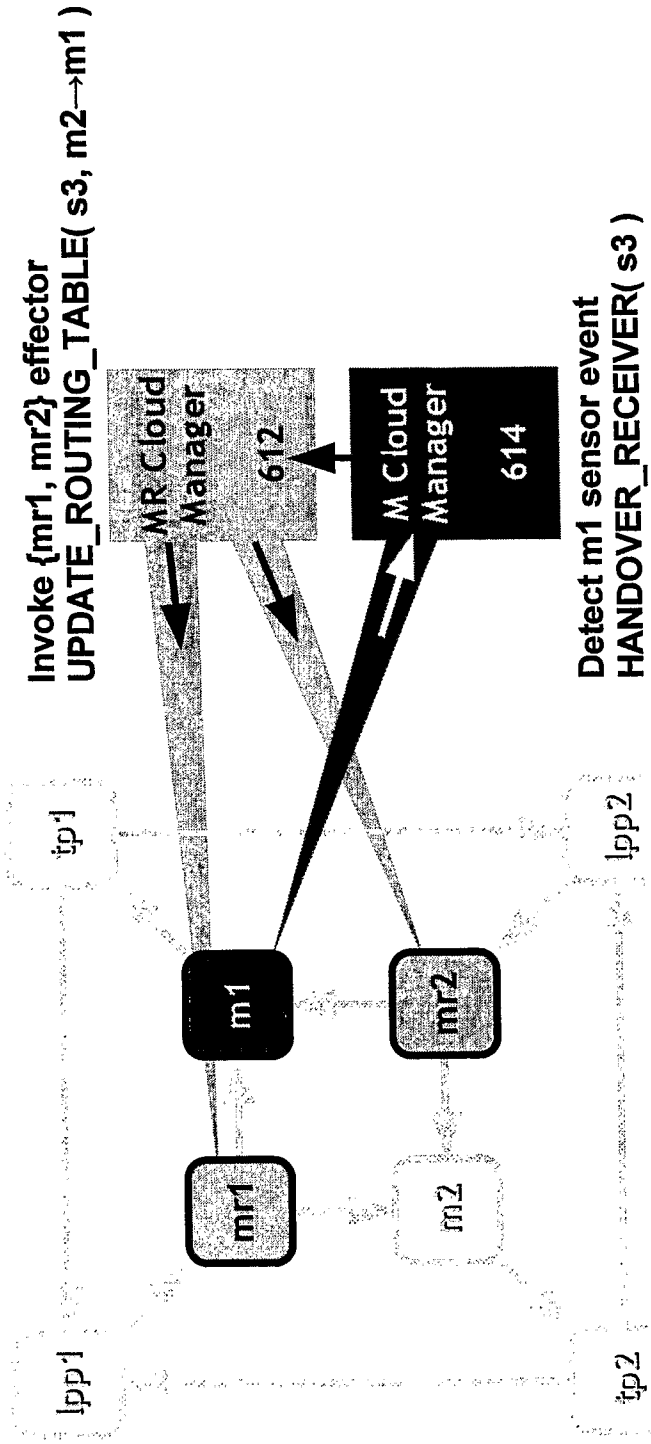

Invoke {mr1, mr2} effector
UPDATE_ROUTING_TABLE(s3, m2→m1)
MR Cloud Manager
612
M Cloud Manager
614
Detect m1 sensor event
HANDOVER_RECEIVER(s3)
tp1
lpp2
m1
mr2
mr1
m2
lpp1
tp2

Figure 7D

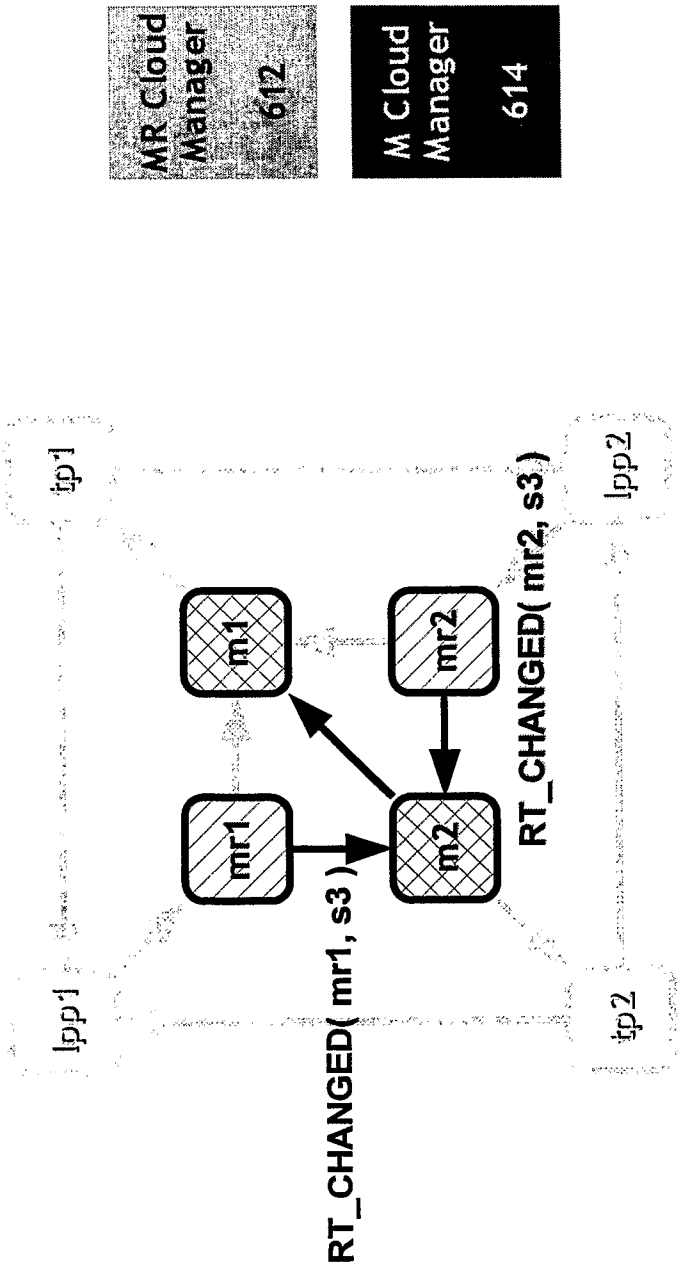

MR Cloud Manager 612
M Cloud Manager 614
tp1
tp2
lpp1
lpp2
m1
mr2
mr1
m2
RT_CHANGED(mr2, s3)
RT_CHANGED(mr1, s3)

Figure 7E

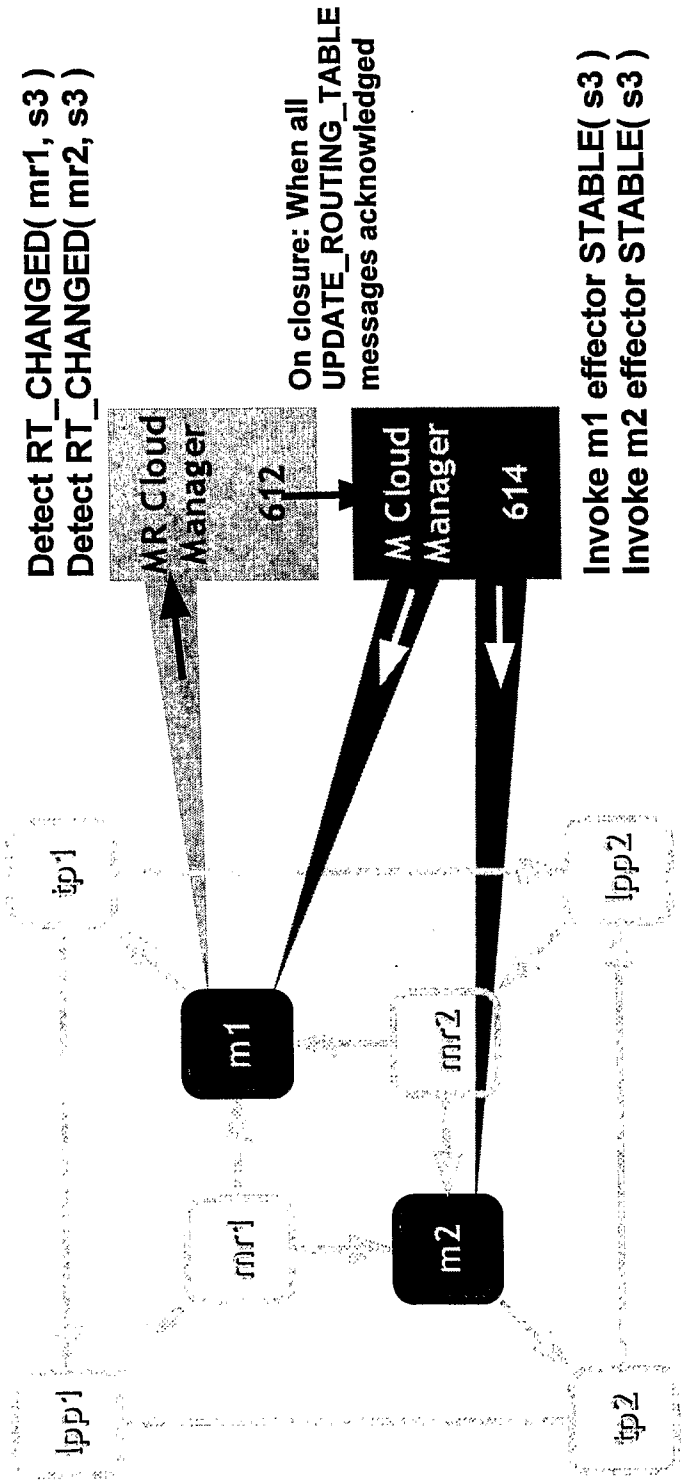

Detect RT_CHANGED(mr1, s3)
Detect RT_CHANGED(mr2, s3)
MR Cloud Manager
612
On closure: When all UPDATE_ROUTING_TABLE messages acknowledged
M Cloud Manager
614
Invoke m1 effector STABLE(s3)
Invoke m2 effector STABLE(s3)
tp1
lpp1
mr1
m1
mr2
m2
lpp2
tp2

Figure 7F

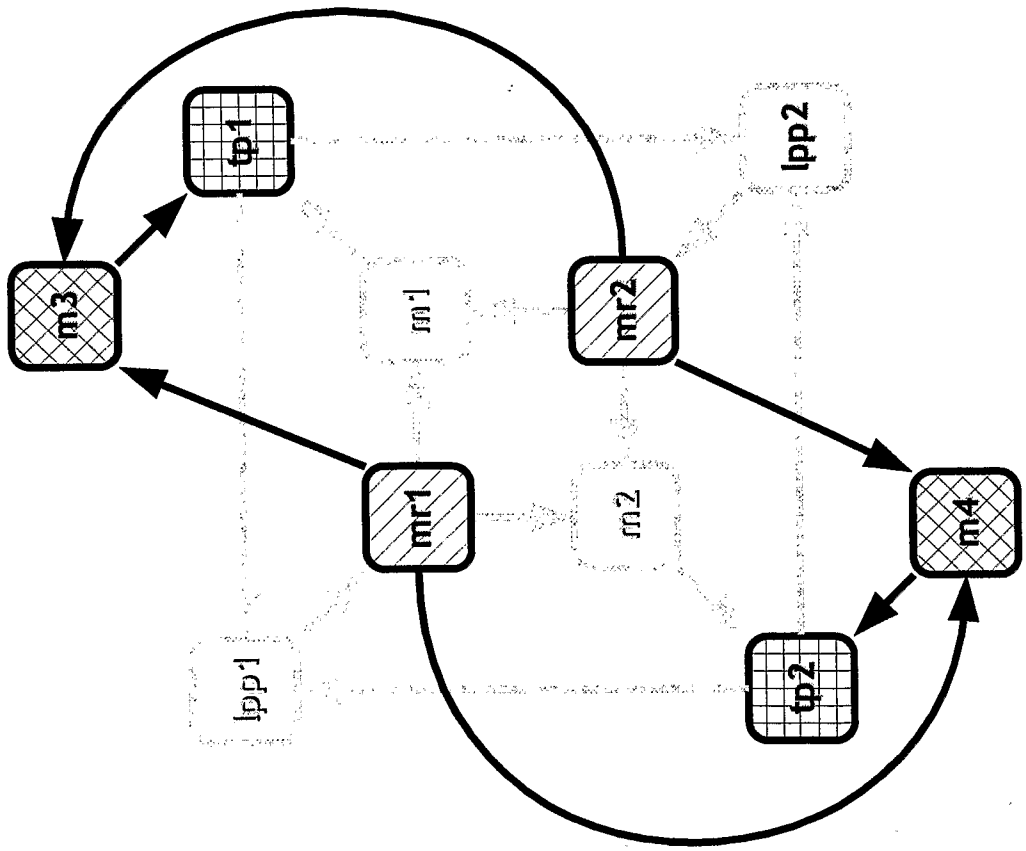


Figure 8

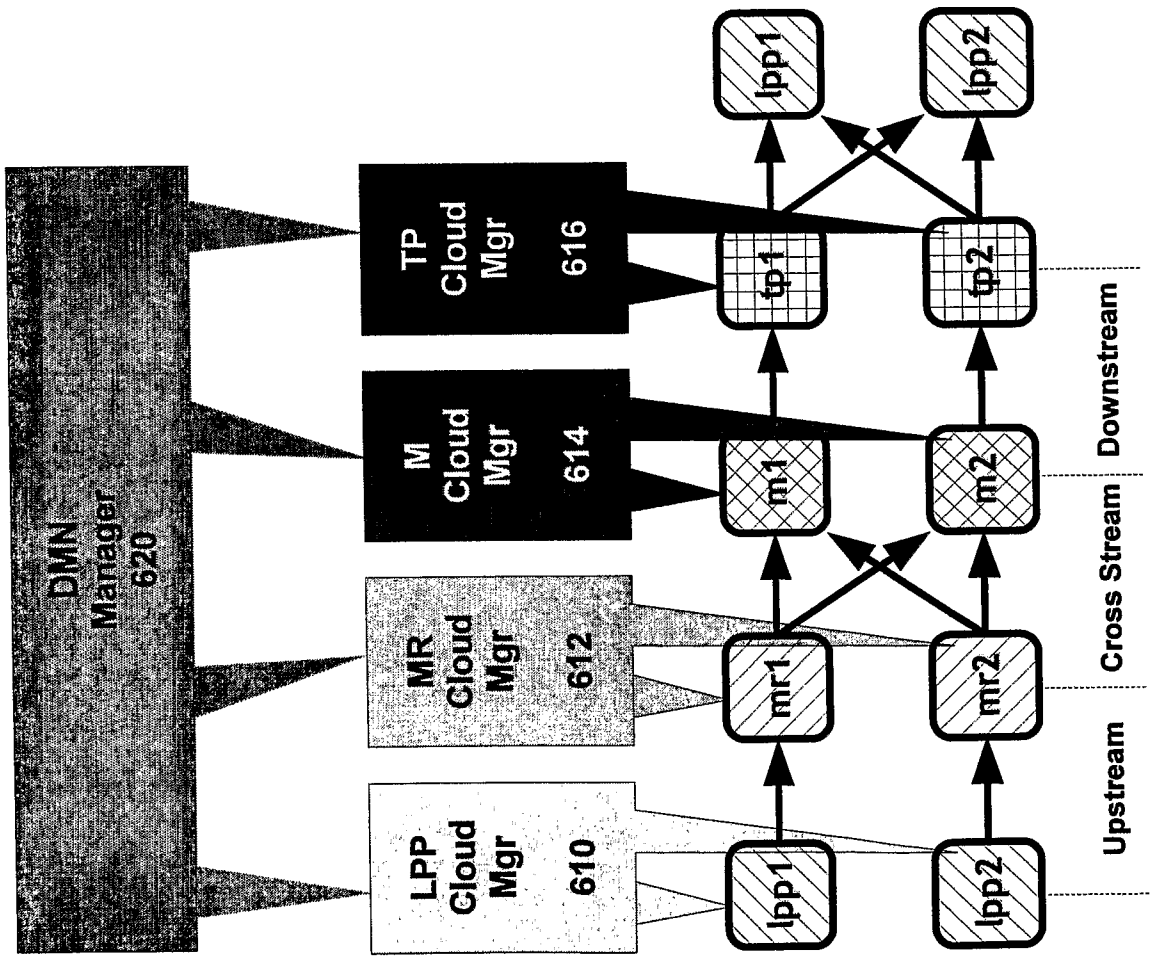

DMN Manager 620
LPP Cloud Mgr 610
MR Cloud Mgr 612
M Cloud Mgr 614
TP Cloud Mgr 616
lpp1
lpp2
mr1
mr2
m1
m2
tp1
tp2
lpp1
lpp2
Upstream
Cross Stream
Downstream

Figure 9

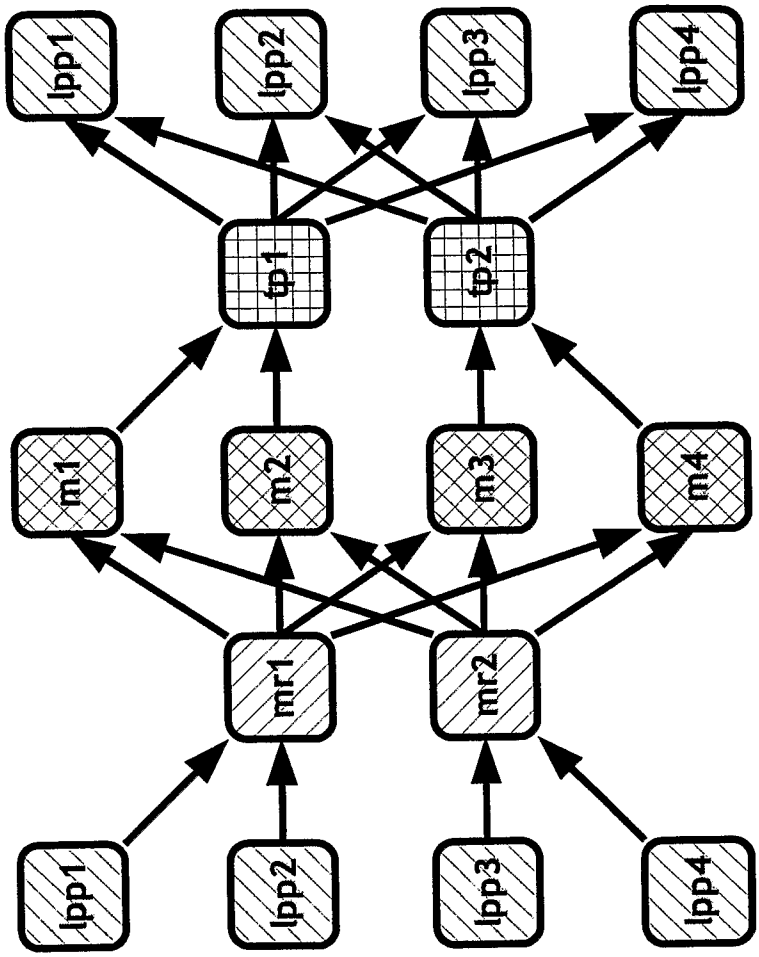

lpp1
lpp2
lpp3
lpp4
tp1
tp2
m1
m2
m3
m4
mr1
mr2
lpp1
lpp2
lpp3
lpp4

Figure 10A

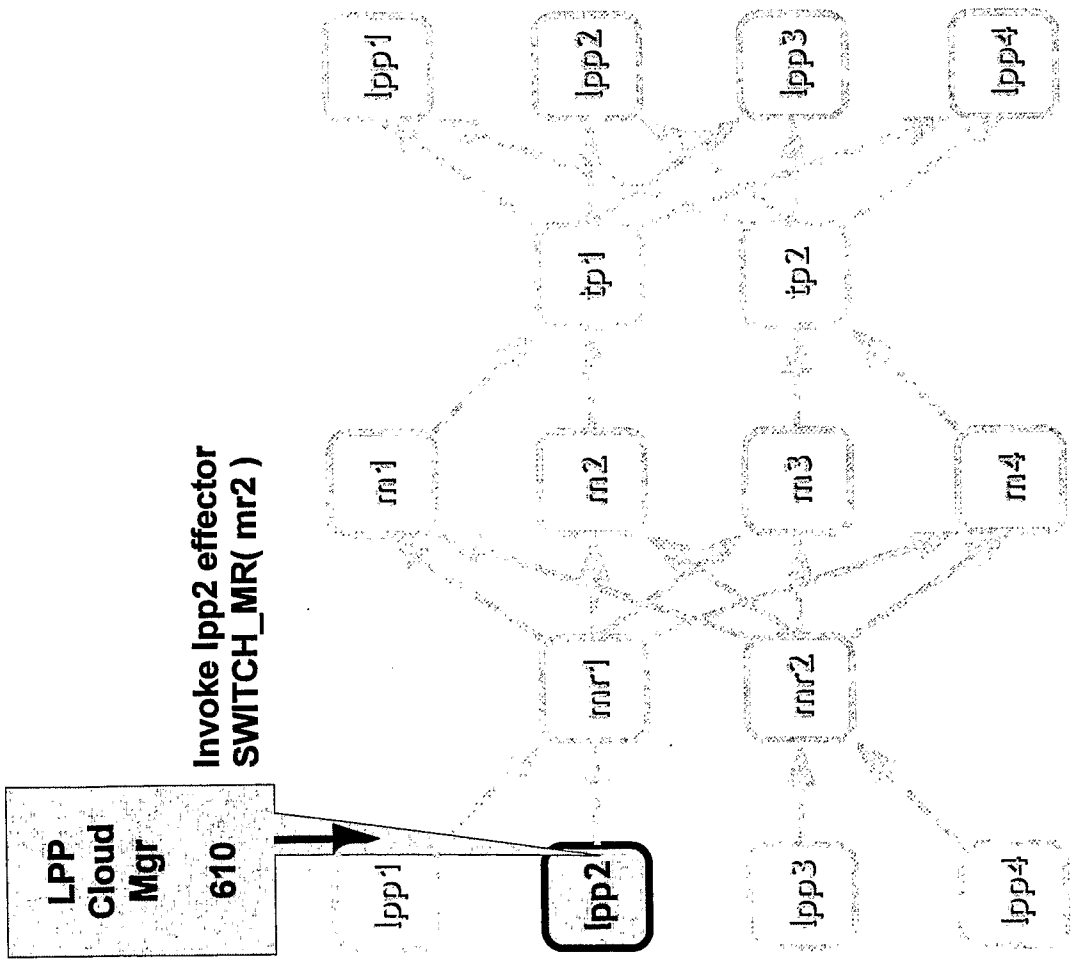

LPP Cloud Mgr 610
Invoke lpp2 effector SWITCH_MR(mr2)
lpp1
lpp2
lpp3
lpp4
tp1
tp2
rn1
rn2
rn3
rn4
mr1
mr2
lpp1
lpp2
lpp3
lpp4

Figure 10B

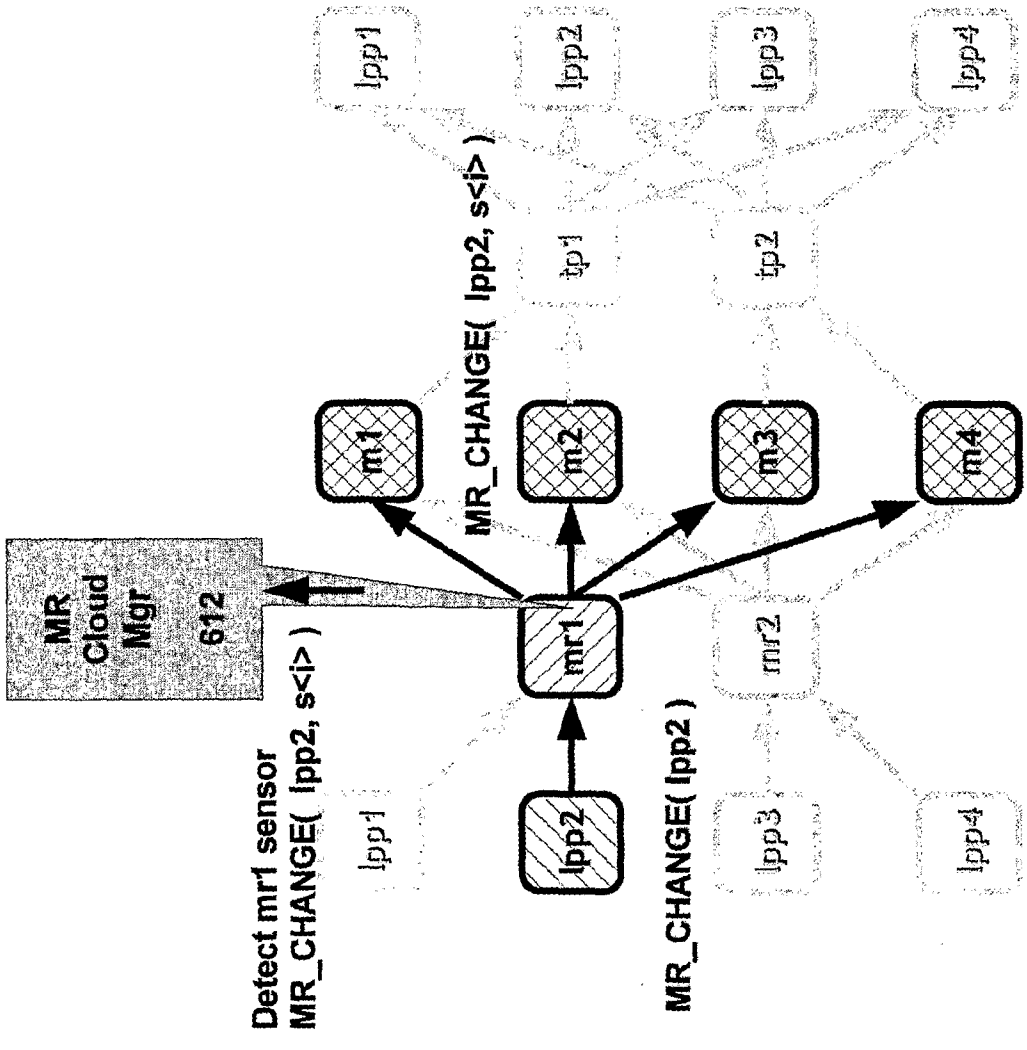

MR Cloud Mgr 612
Detect mr1 sensor MR_CHANGE(lpp2, s<i>)
MR_CHANGE(lpp2, s<i>)
MR_CHANGE(lpp2)
lpp1
lpp2
lpp3
lpp4
tp1
tp2
m1
m2
m3
m4
mr1
mr2
lpp2
lpp1
lpp3
lpp4

Figure 10C

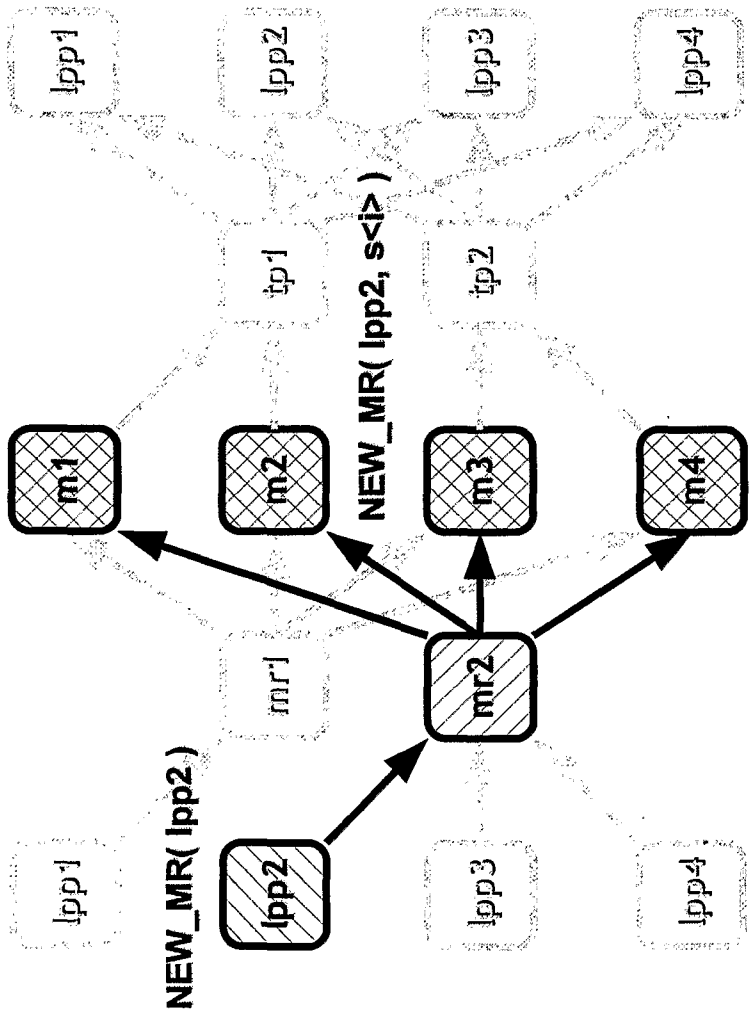

lpp1
lpp2
lpp3
lpp4
tp1
tp2
m1
m2
m3
m4
NEW_MR(lpp2, s<i>)
mr1
mr2
NEW_MR(lpp2)
lpp2
lpp1
lpp3
lpp4

Figure 10D

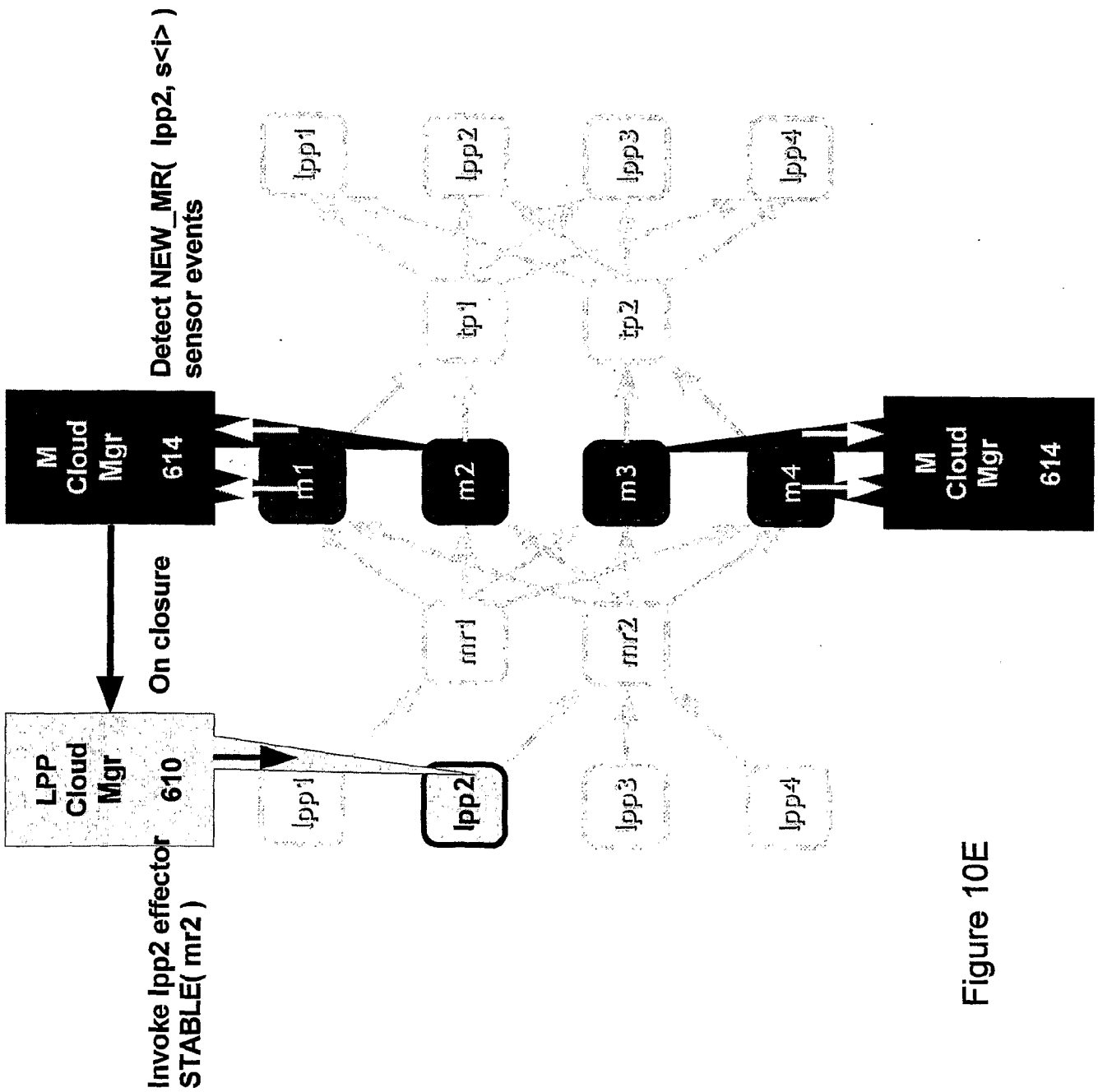

M Cloud Mgr 614
Detect NEW_MR(lpp2, s<i>) sensor events
lpp1
lpp2
lpp3
lpp4
tp1
tp2
m1
m2
m3
m4
M Cloud Mgr 614
On closure
mr1
mr2
LPP Cloud Mgr 610
lpp1
lpp2
lpp3
lpp4
Invoke lpp2 effector STABLE(mr2)

Figure 10E

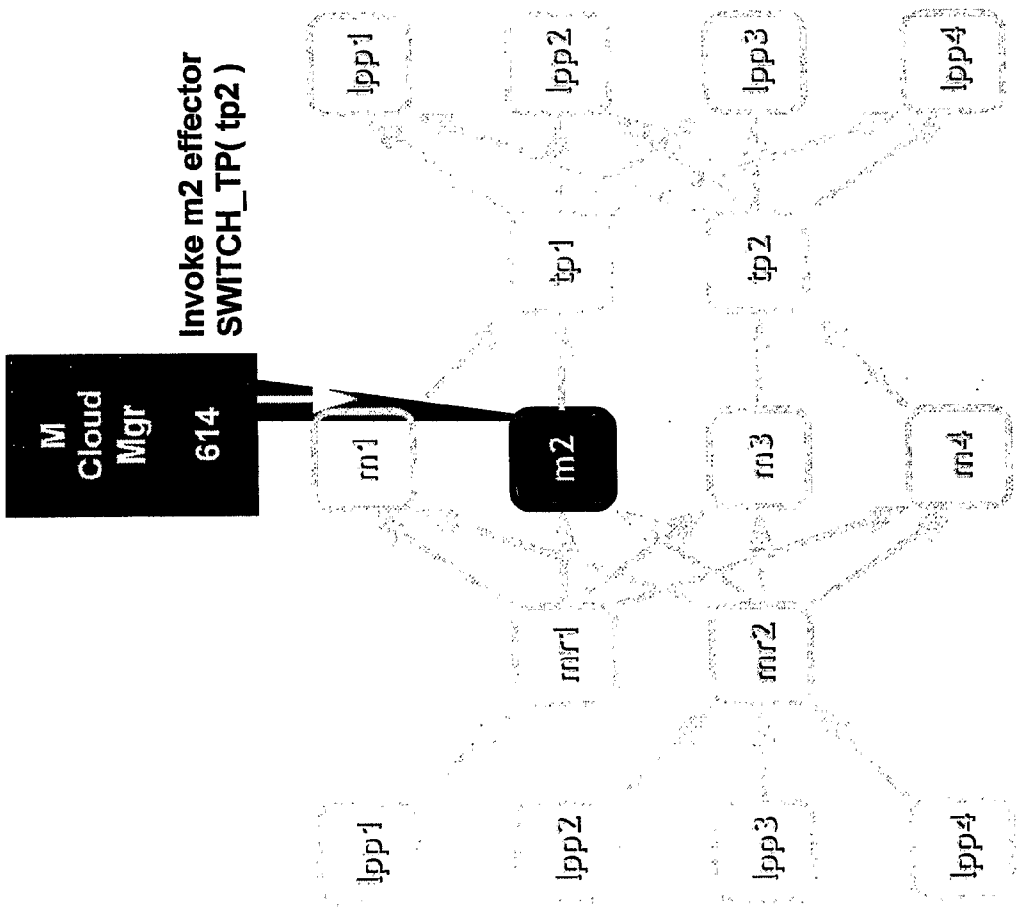

M
Cloud
Mgr
614
Invoke m2 effector
SWITCH_TP(tp2)
lpp1
lpp2
lpp3
lpp4
tp1
tp2
m1
m2
m3
m4
mr1
mr2
lpp1
lpp2
lpp3
lpp4

Figure 11A

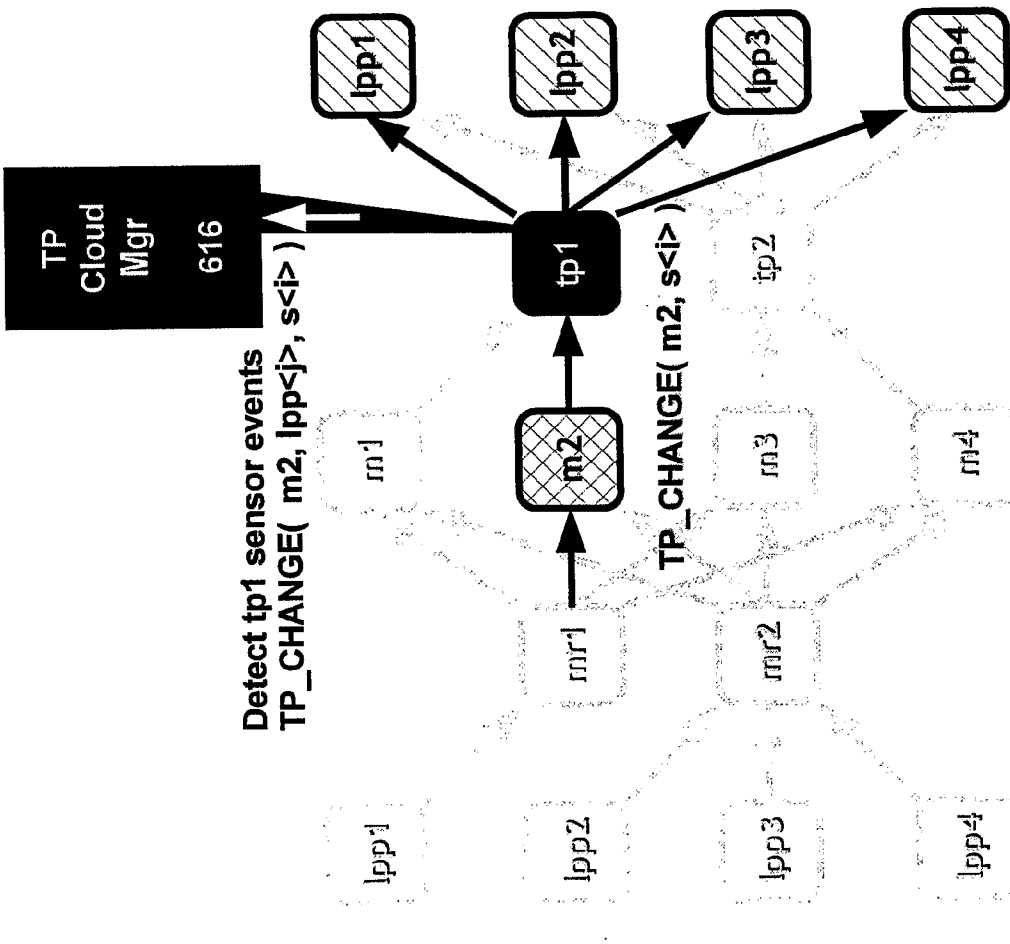

TP Cloud Mgr 616
Detect tp1 sensor events
TP_CHANGE(m2, lpp<j>, s<i>)
TP_CHANGE(m2, s<i>)
lpp1
lpp2
lpp3
lpp4
tp1
tp2
m1
m2
m3
m4
mr1
mr2

Figure 11B

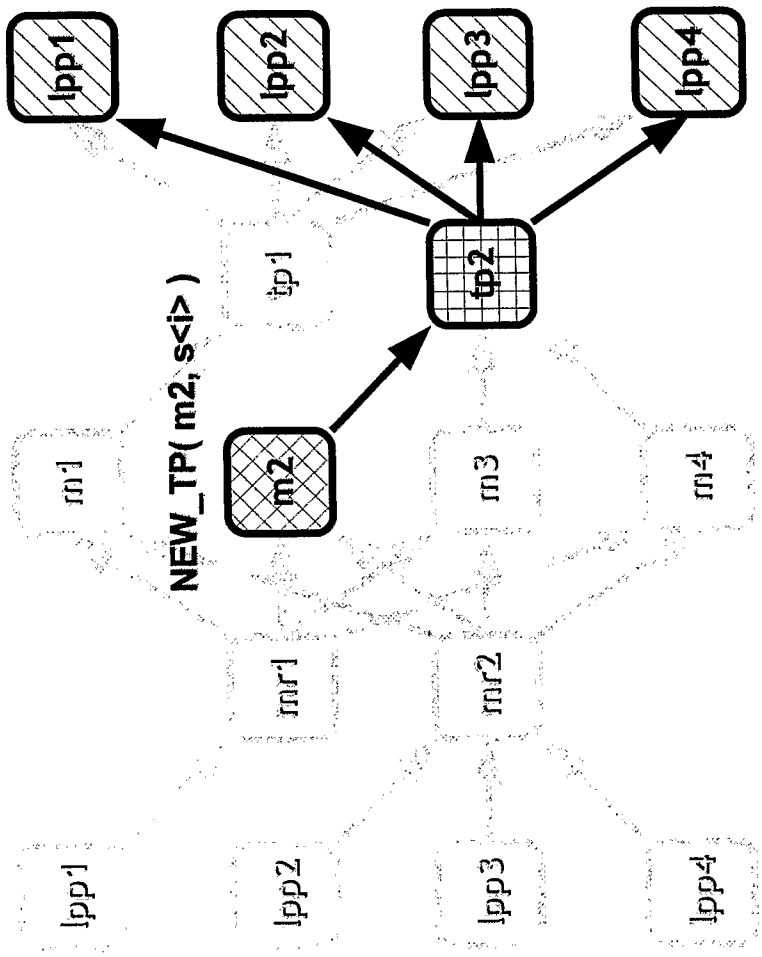

lpp1
lpp2
lpp3
lpp4
tp2
tp1
NEW_TP(m2, s<i>)
m2
m1
m3
m4
mr1
mr2
lpp1
lpp2
lpp3
lpp4

Figure 11C

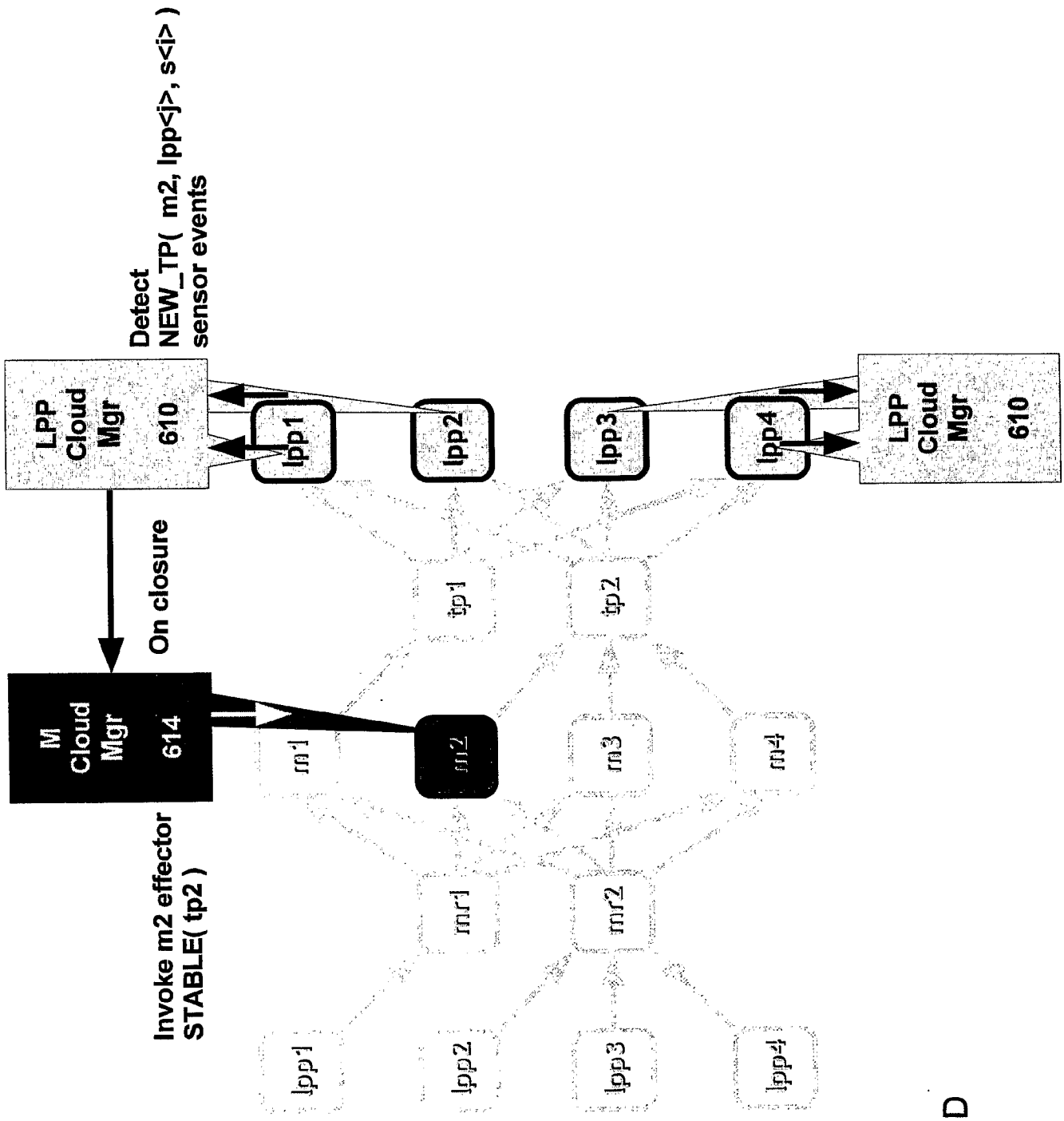

Detect
NEW_TP(m2, lpp<j>, s<i>)
sensor events
LPP
Cloud
Mgr
610
lpp1
lpp2
lpp3
lpp4
LPP
Cloud
Mgr
610
On closure
M
Cloud
Mgr
614
Invoke m2 effector
STABLE(tp2)
tp1
tp2
m1
m2
m3
m4
mr1
mr2
lpp1
lpp2
lpp3
lpp4

Figure 11D

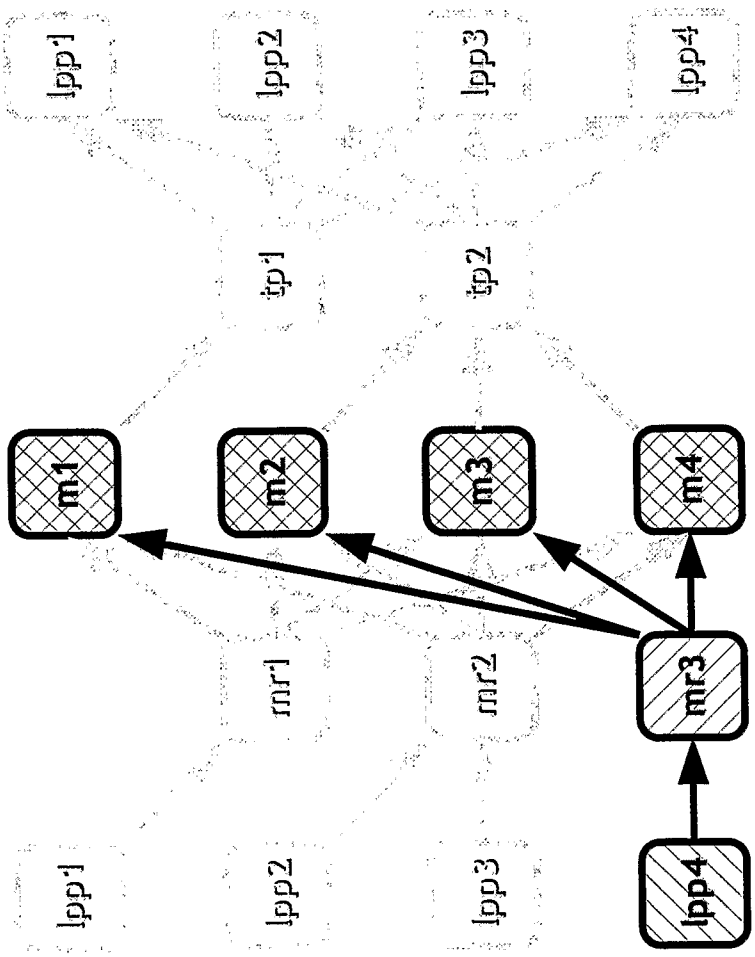

lpp1
lpp2
lpp3
lpp4
tp1
tp2
m1
m2
m3
m4
mr1
mr2
mr3
lpp1
lpp2
lpp3
lpp4

Figure 12A

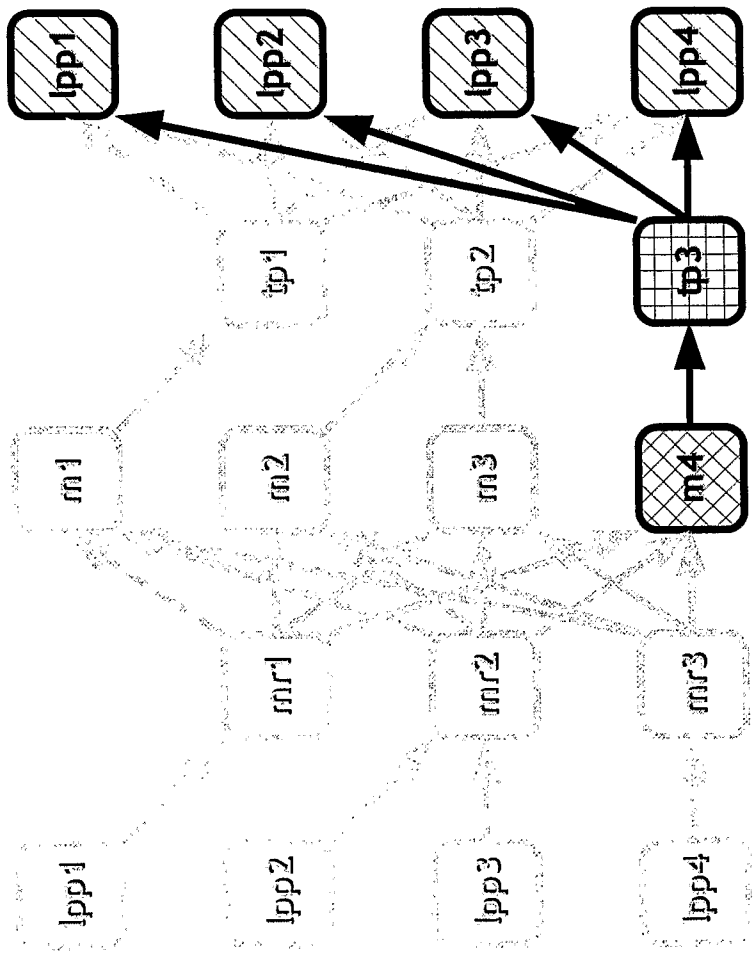

lpp1
lpp2
lpp3
lpp4
tp1
tp2
tp3
m1
m2
m3
m4
mr1
mr2
mr3
lpp1
lpp2
lpp3
lpp4

Figure 12B

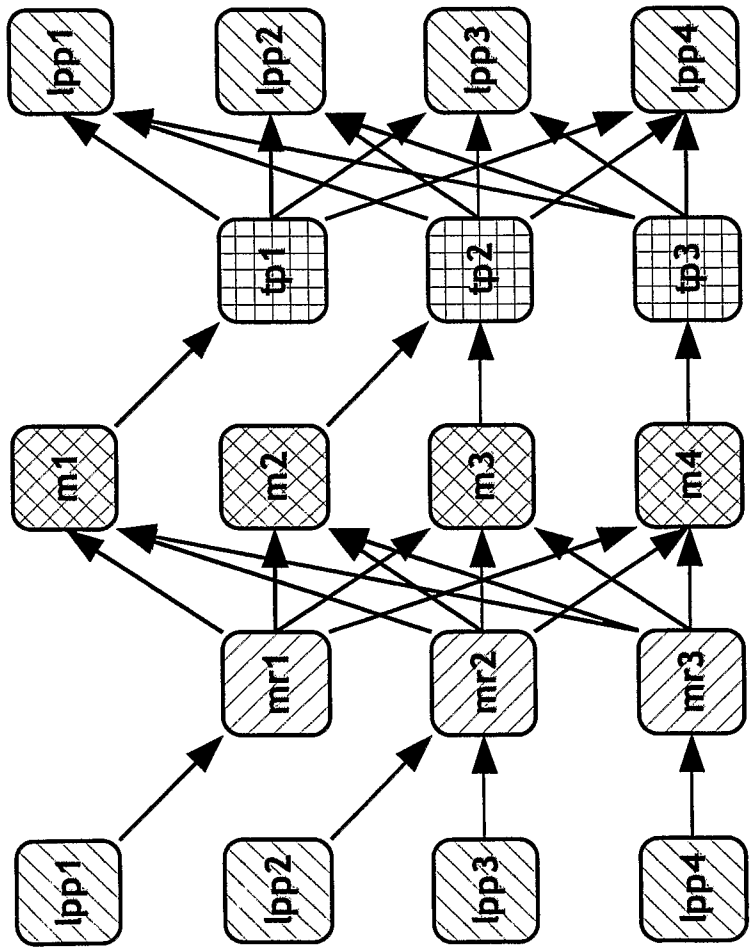

lpp1
lpp2
lpp3
lpp4
tp1
tp2
tp3
m1
m2
m3
m4
mr1
mr2
mr3
lpp1
lpp2
lpp3
lpp4

Figure 12C

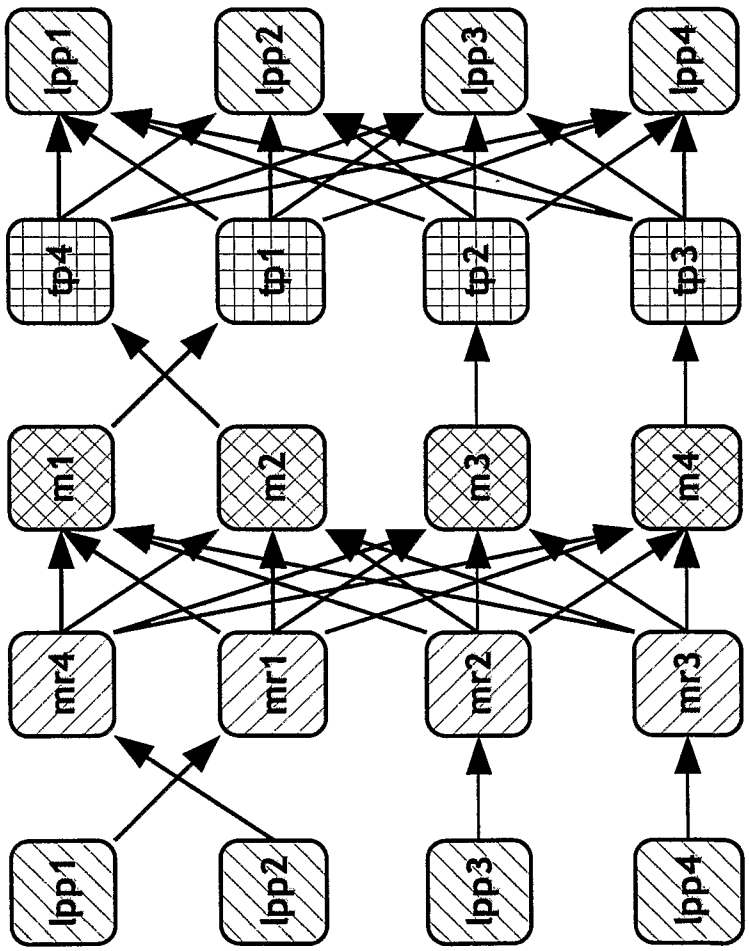

lpp1
lpp2
lpp3
lpp4
φ4
φ1
φ2
φ3
m1
m2
m3
m4
mr4
mr1
mr2
mr3
lpp1
lpp2
lpp3
lpp4

Figure 12D

Fig 13

Management Locus	Workload Metric (msgs/sec)	Condition	Action	Exception
M Cloud Manager	W=throughput, and w(s) = throughput per segment "s"	M node HIGH (W > M_NODE_HIGH_WATERMARK)	MOVE M SEGMENTS, to least-loaded M node	If only one M segment on overloaded M node, then inform DMN manager
		all M nodes HIGH (AVG(W) > M_NODE_POOL_HIGH_WATERMARK)	ADD M NODE (request to DMN Manager); then respond to individual "M node HIGH" events	None
MR Cloud Manager	W=throughput, and w(n) = throughput per upstream LPP node "n"	MR node HIGH (W > MR_NODE_HIGH_WATERMARK)	REASSIGN LPP	If only 1 upstream LPP on overloaded MR node, then inform DMN manager
		all MR nodes HIGH (AVG(W) > MR_NODE_POOL_HIGH_WATERMARK)	ADD MR NODE (request to DMN Manager); then respond to individual "MR node HIGH" events	None
TP Cloud Manager	W=throughput, and w(n) = throughput per upstream M node "n"	TP node HIGH (W > TP_NODE_HIGH_WATERMARK)	MOVE M SEGMENTS or REASSIGN M	If only 1 upstream M on overloaded TP node, then inform DMN manager
		all TP nodes HIGH (AVG(W) > TP_NODE_POOL_HIGH_WATERMARK)	ADD TP NODE (request to DMN Manager); then respond to individual "TP node HIGH" events	None
DMN Manager	n/a	Exception from one of the cloud managers	Various (see text)	
		ADD NODE request, resources available	Allocate resource to cloud and kick off ADD NODE operation	
		ADD NODE request, resources unavailable	Various (see text)	

INTERNATIONAL SEARCH REPORT

International application No
PCT/GB2007/002195

A. CLASSIFICATION OF SUBJECT MATTER
INV. H04L29/08

According to International Patent Classification (IPC) or to both national classification and IPC

B. FIELDS SEARCHED

Minimum documentation searched (classification system followed by classification symbols)
H04L

Documentation searched other than minimum documentation to the extent that such documents are included in the fields searched

Electronic data base consulted during the international search (name of data base and, where practical, search terms used)
EPO-Internal, COMPENDEX, INSPEC, IBM-TDB, WPI Data

C. DOCUMENTS CONSIDERED TO BE RELEVANT

Category*	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
Y	WO 2005/013554 A (ENIGMATEC CORP [GB]; JOHNSTON-WATT DUNCAN [GB]; WEST ANDREW MARTIN [GB) 10 February 2005 (2005-02-10) abstract page 1, line 1 - page 3, line 27 page 4, line 17 - page 5, line 19 page 12, line 16 - page 14, line 11 page 20, line 30 - page 22, line 11 claims 1,23	1-27
Y	WO 99/23784 A2 (ORACLE CORP [US]) 14 May 1999 (1999-05-14) abstract page 1, line 1 - page 4, line 7 page 7, line 17 - page 11, line 5 page 12, line 27 - page 14, line 6	1-27

☐ Further documents are listed in the continuation of Box C.

☒ See patent family annex.

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"&" document member of the same patent family

Date of the actual completion of the international search	Date of mailing of the international search report
1 October 2007	10/10/2007

Name and mailing address of the ISA/	Authorized officer
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INTERNATIONAL SEARCH REPORT

Information on patent family members

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