

ALLEGREX™ Instruction Manual

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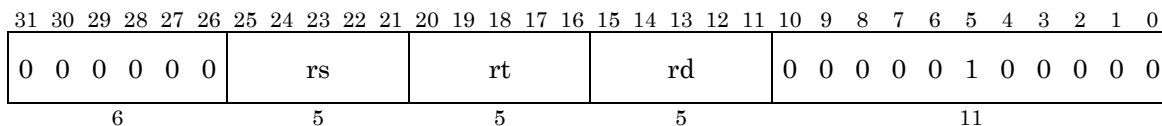
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MIPS Instructions

add

Add



MIPS I

Syntax:

```
add rd, rs, rt
```

Description:

The contents of registers *rs* and *rt* are added together and the result is stored in register *rd*. If a two's complement overflow occurs, an exception is generated.

Operation:

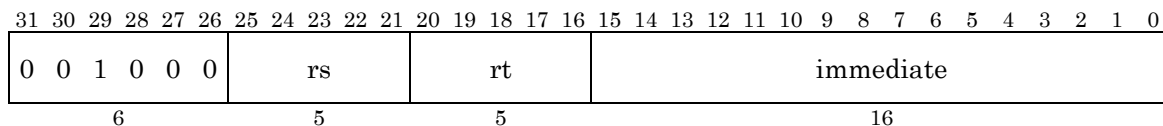
```
temp ← (GPR[rs]31 || GPR[rs]31..0) + (GPR[rt]31 || GPR[rt]31..0)
if (temp32 ≠ temp31) then
    SignalException(IntegerOverflow)
else
    GPR[rd] ← temp
endif
```

Exceptions:

Integer Overflow exception

addi

Add Immediate



MIPS I

Syntax:

```
addi rt,rs,immediate
```

Description:

The 16-bit immediate field is sign-extended to 32 bits and added to register rs. The 32-bit result is stored in register rt.

If a two's complement overflow occurs, an exception is generated.

Operation:

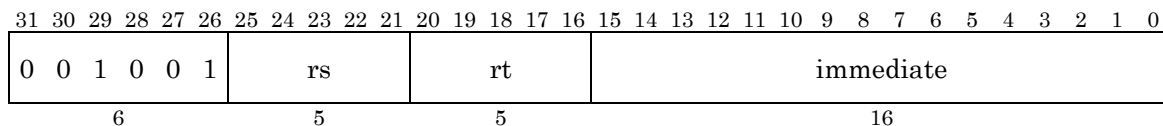
```
temp ← (GPR[rs]31 || GPR[rs]31..0) + sign_extend(immediate)
if (temp32 ≠ temp31) then
    SignalException(IntegerOverflow)
else
    GPR[rt] ← temp
endif
```

Exceptions:

Integer Overflow exception

addiu

Add Immediate Unsigned



MIPS I

Syntax:

```
addiu rt,rs,immediate
```

Description:

The 16-bit immediate field is sign-extended to 32 bits and added to register rs. The 32-bit result is stored in register rt.

No exception is generated even if an overflow occurs.

Operation:

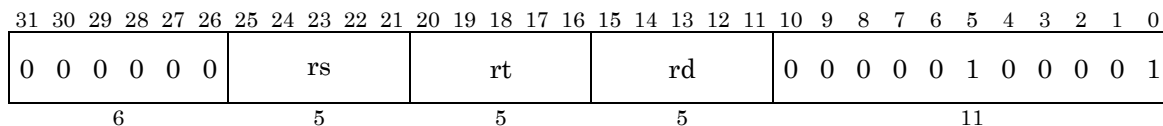
```
temp ← GPR[rs] + sign_extend(immediate)
GPR[rt] ← temp
```

Exceptions:

None

addu

Add Unsigned



MIPS I

Syntax:

```
addu rd,rs,rt
```

Description:

The contents of registers *rs* and *rt* are added together and the result is stored in register *rd*.

No exception is generated even if a two's complement overflow occurs.

Operation:

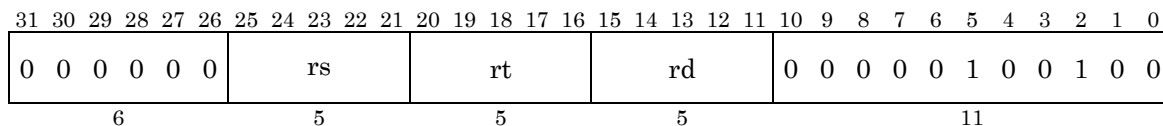
```
temp ← GPR[rs] + GPR[rt]
GPR[rd] ← temp
```

Exceptions:

None

and

AND



MIPS I

Syntax:

and rd, rs, rt

Description:

A bitwise AND (logical product) is performed for the contents of registers rs and rt and the result is stored in register rd.

X	Y	X AND Y
0	0	0
0	1	0
1	0	0
1	1	1

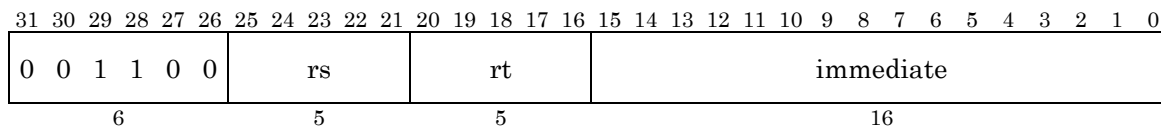
Operation:

$$\text{GPR}[\text{rd}] \leftarrow \text{GPR}[\text{rs}] \text{ and } \text{GPR}[\text{rt}]$$
Exceptions:

None

andi

AND Immediate



MIPS I

Syntax:

```
andi rt,rs,immediate
```

Description:

The 16-bit immediate field is zero-extended, and a bitwise AND (logical product) is performed between the zero-extended value and the contents of register rs. The result is stored in register rt.

X	Y	X AND Y
0	0	0
0	1	0
1	0	0
1	1	1

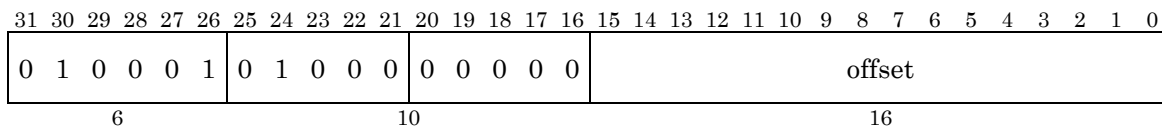
Operation:

$$\text{GPR}[rt] \leftarrow \text{GPR}[rs] \text{ and } \text{zero_extend}(\text{immediate})$$
Exceptions:

None

bc1f

Branch on FPU False

MIPS I
FPU**Syntax:**`bc1f offset`**Description:**

When the FPU condition is false, the program branches with a one instruction delay to the branch target address.

The branch target address is the sum of the PC and the 16-bit offset after it is shifted left two bits and sign-extended to a 32 bit value.

Restrictions:

To change the FPU condition, at least one nop is required between the FPU instruction that changes the condition and the bc1f instruction.

A nop is automatically inserted by the compiler.

The bc1f instruction cannot be placed in the delay slot of a branch/jump instruction.

The ctc1 instruction cannot be placed in the delay slot of the bc1f instruction.

Operation:

```

I - 1:
    condition ← notCOC[1]
I :
    target_offset ← sign_extend(offset || 02)
I + 1:
    if (condition) then
        PC ← PC + target_offset
    endif

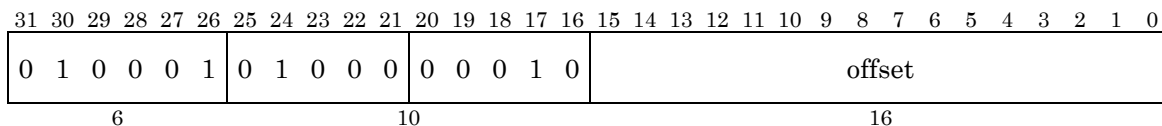
```

Exceptions:

Coprocessor Unusable exception

bc1fl

Branch on FPU False Likely

MIPS II
FPU**Syntax:**`bc1fl offset`**Description:**

When the FPU condition is false, the program branches with a one instruction delay to the branch target address.

The branch target address is the sum of the PC and the 16-bit offset after it is shifted left two bits and sign-extended to a 32 bit value.

If the branch is not taken, the instruction in the branch delay slot is discarded.

Restrictions:

To change the FPU condition, at least one nop is required between the FPU instruction that changes the condition and the bc1fl instruction.

A nop is automatically inserted by the compiler.

The bc1fl instruction cannot be placed in the delay slot of a branch/jump instruction.

The etc1 instruction cannot be placed in the delay slot of the bc1fl instruction.

Operation:

```

I - 1:
    condition ← notCOC[1]
I :
    target_offset ← sign_extend(offset || 02)
I + 1:
    if (condition) then
        PC ← PC + target_offset
    else
        NullifyCurrentInstruction()
    endif

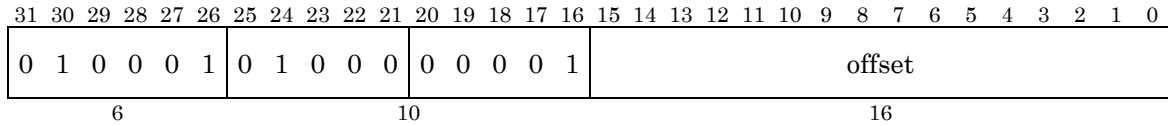
```

Exceptions:

Coprocessor Unusable exception

bc1t

Branch on FPU True

MIPS I
FPU**Syntax:**`bc1t offset`**Description:**

When the FPU condition is true, the program branches with a one instruction delay to the branch target address.

The branch target address is the sum of the PC and the 16-bit offset after it is shifted left two bits and sign-extended to a 32 bit value.

Restrictions:

To change the FPU condition, at least one nop is required between the FPU instruction that changes the condition and the bc1t instruction.

A nop is automatically inserted by the compiler.

The bc1t instruction cannot be placed in the delay slot of a branch/jump instruction.

The ctc1 instruction cannot be placed in the delay slot of the bc1t instruction.

Operation:

```

I - 1:
    condition ← COC[1]
I :
    target_offset ← sign_extend(offset || 02)
I + 1:
    if (condition) then
        PC ← PC + target_offset
    endif

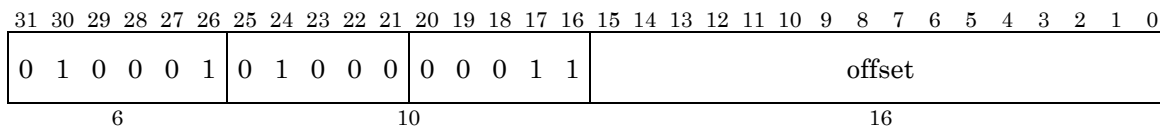
```

Exceptions:

Coprocessor Unusable exception

bc1tl

Branch on FPU True Likely

MIPS II
FPU**Syntax:**`bc1tl offset`**Description:**

When the FPU condition is true, the program branches with a one instruction delay to the branch target address.

The branch target address is the sum of the PC and the 16-bit offset after it is shifted left two bits and sign-extended to a 32 bit value.

If the branch is not taken, the instruction in the branch delay slot is discarded.

Restrictions:

To change the FPU condition, at least one nop is required between the FPU instruction that changes the condition and the bc1tl instruction.

A nop is automatically inserted by the compiler.

The bc1tl instruction cannot be placed in the delay slot of a branch/jump instruction.

The etc1 instruction cannot be placed in the delay slot of the bc1tl instruction.

Operation:

```

I - 1:
    condition ← COC[1]
I :
    target_offset ← sign_extend(offset || 02)
I + 1:
    if (condition) then
        PC ← PC + target_offset
    else
        NullifyCurrentInstruction()
    endif

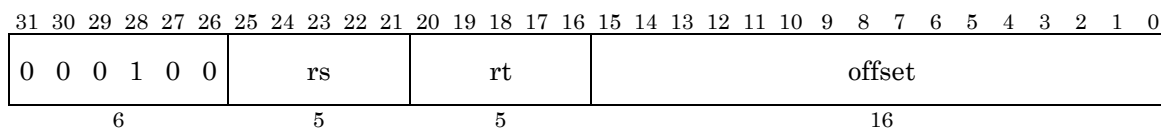
```

Exceptions:

Coprocessor Unusable exception

beq

Branch on Equal



MIPS I

Syntax:

```
beq rs,rt,offset
```

Description:

When register *rs* is equal to register *rt*, the program branches with a one instruction delay to the branch target address.

The branch target address is the sum of the PC and the 16-bit offset after it is shifted left two bits and sign-extended to a 32 bit value.

Restrictions:

The *beq* instruction cannot be placed in the delay slot of a branch/jump instruction.

Operation:

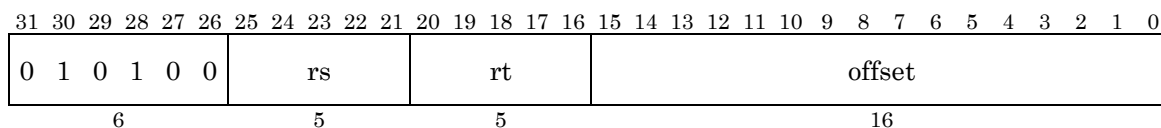
```
I :
    target_offset ← sign_extend(offset || 02)
    condition ← (GPR[rs]=GPR[rt])
I + 1:
    if (condition) then
        PC ← PC + target_offset
    endif
```

Exceptions:

None

beql

Branch on Equal Likely



MIPS II

Syntax:

```
beql rs,rt,offset
```

Description:

When register *rs* is equal to register *rt*, the program branches with a one instruction delay to the branch target address.

The branch target address is the sum of the PC and the 16-bit offset after it is shifted left two bits and sign-extended to a 32 bit value. If the branch is not taken, the instruction in the branch delay slot is discarded.

Restrictions:

The *beql* instruction cannot be placed in the delay slot of a branch/jump instruction.

Operation:

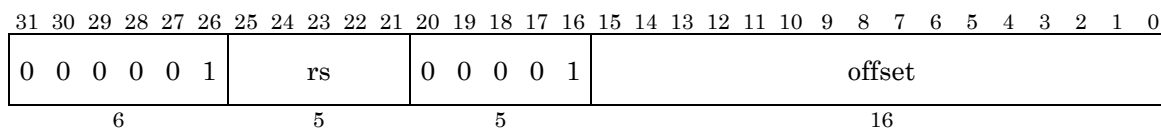
```
I :
    target_offset ← sign_extend(offset || 02)
    condition ← (GPR[rs]=GPR[rt])
I + 1:
    if (condition) then
        PC ← PC + target_offset
    else
        NullifyCurrentInstruction()
    endif
```

Exceptions:

None

bgez

Branch on Greater than or Equal to Zero



MIPS I

Syntax:

```
bgez rs,offset
```

Description:

When register *rs* is greater than or equal to zero, the program branches with a one instruction delay to the branch target address.

The branch target address is the sum of the PC and the 16-bit offset after it is shifted left two bits and sign-extended to a 32 bit value.

Restrictions:

The *bgez* instruction cannot be placed in the delay slot of a branch/jump instruction.

Operation:

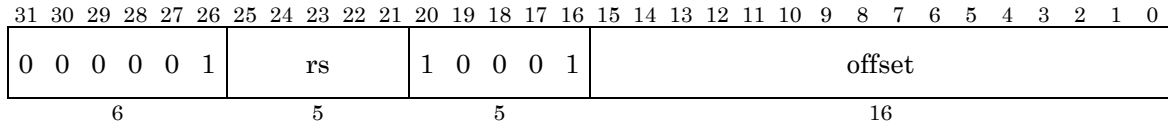
```
I :
    target_offset ← sign_extend(offset || 02)
    condition ← (GPR[rs] ≥ 032)
I + 1:
    if (condition) then
        PC ← PC + target_offset
    endif
```

Exceptions:

None

bgezal

Branch on Greater than or Equal to Zero And Link



MIPS I

Syntax:

```
bgezal rs,offset
```

Description:

When register *rs* is greater than or equal to zero, the program branches with a one instruction delay to the branch target address.

The branch target address is the sum of the PC and the 16-bit offset after it is shifted left two bits and sign-extended to a 32 bit value.

The address of the instruction following the delay slot is stored in register *r31* (the link register).

Restrictions:

The *bgezal* instruction cannot be placed in the delay slot of a branch/jump instruction.

The register *r31* (link register) cannot be specified for the register *rs*.

The instruction for changing the register *r31* (link register) cannot be placed in the delay slot of the *bgezal* instruction.

Operation:

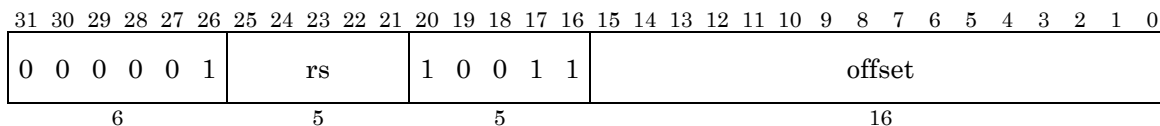
```
I :
  target_offset ← sign_extend(offset || 02)
  condition ← (GPR[rs] ≥ 032)
  GPR[31] ← PC + 8
I + 1 :
  if (condition) then
    PC ← PC + target_offset
  endif
```

Exceptions:

None

bgezall

Branch on Greater than or Equal to Zero And Link Likely



MIPS II

Syntax:

```
bgezall rs,offset
```

Description:

When register *rs* is greater than or equal to zero, the program branches with a one instruction delay to the branch target address. The branch target address is the sum of the PC and the 16-bit offset after it is shifted left two bits and sign-extended to a 32 bit value. The address of the instruction following the delay slot is stored in register *r31* (the link register). If the branch is not taken, the instruction in the branch delay slot is discarded.

Restrictions:

The *bgezall* instruction cannot be placed in the delay slot of a branch/jump instruction. The register *r31* (link register) cannot be specified for the register *rs*. The instruction for changing the register *r31* (link register) cannot be placed in the delay slot of the *bgezall* instruction.

Operation:

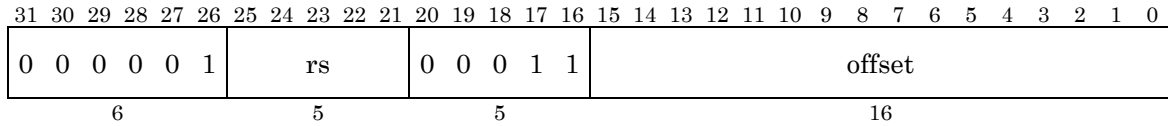
```
I :
  target_offset ← sign_extend(offset || 02)
  condition ← (GPR[rs] ≥ 032)
  GPR[31] ← PC + 8
I + 1 :
  if (condition) then
    PC ← PC + target_offset
  else
    NullifyCurrentInstruction()
  endif
```

Exceptions:

None

bgezl

Branch on Greater than or Equal to Zero Likely



MIPS II

Syntax:

```
bgezl rs,offset
```

Description:

When register *rs* is greater than or equal to zero, the program branches with a one instruction delay to the branch target address. The branch target address is the sum of the PC and the 16-bit offset after it is shifted left two bits and sign-extended to a 32 bit value. If the branch is not taken, the instruction in the branch delay slot is discarded.

Restrictions:

The *bgezl* instruction cannot be placed in the delay slot of a branch/jump instruction.

Operation:

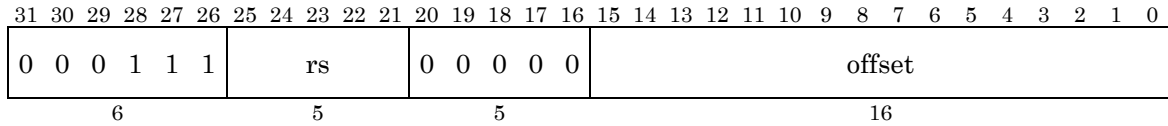
```
I :
  target_offset ← sign_extend(offset || 02)
  condition ← (GPR[rs] ≥ 032)
I + 1:
  if (condition) then
    PC ← PC + target_offset
  else
    NullifyCurrentInstruction()
  endif
```

Exceptions:

None

bgtz

Branch on Greater Than Zero



MIPS I

Syntax:

```
bgtz rs,offset
```

Description:

When register *rs* is greater than zero, the program branches with a one instruction delay to the branch target address.

The branch target address is the sum of the PC and the 16-bit offset after it is shifted left two bits and sign-extended to a 32 bit value.

Restrictions:

The *bgtz* instruction cannot be placed in the delay slot of a branch/jump instruction.

Operation:

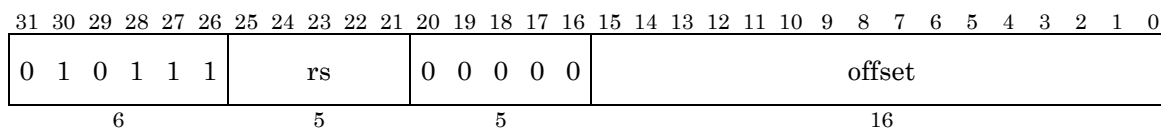
```
I :
    target_offset ← sign_extend(offset || 02)
    condition ← (GPR[rs] > 032)
I + 1:
    if (condition) then
        PC ← PC + target_offset
    endif
```

Exceptions:

None

bgtzl

Branch on Greater Than Zero Likely



MIPS II

Syntax:

```
bgtzl rs,offset
```

Description:

When register *rs* is greater than zero, the program branches with a one instruction delay to the branch target address. The branch target address is the sum of the PC and the 16-bit offset after it is shifted left two bits and sign-extended to a 32 bit value. If the branch is not taken, the instruction in the branch delay slot is discarded.

Restrictions:

The `bgtzl` instruction cannot be placed in the delay slot of a branch/jump instruction.

Operation:

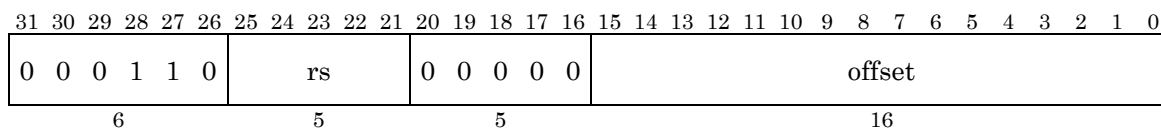
```
I :
    target_offset ← sign_extend(offset || 02)
    condition ← (GPR[rs] > 032)
I + 1:
    if (condition) then
        PC ← PC + target_offset
    else
        NullifyCurrentInstruction()
    endif
```

Exceptions:

None

blez

Branch on Less than or Equal to Zero



MIPS I

Syntax:

```
blez rs,offset
```

Description:

When register *rs* is less than or equal to zero, the program branches with a one instruction delay to the branch target address.

The branch target address is the sum of the PC and the 16-bit offset after it is shifted left two bits and sign-extended to a 32 bit value.

Restrictions:

The *blez* instruction cannot be placed in the delay slot of a branch/jump instruction.

Operation:

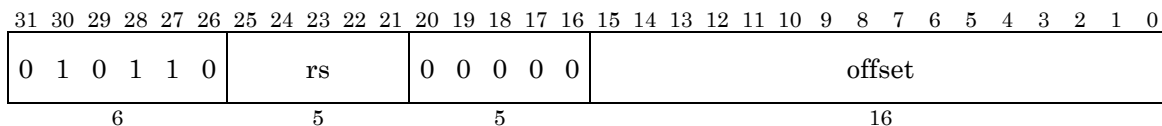
```
I :
  target_offset ← sign_extend(offset || 02)
  condition ← (GPR[rs] ≤ 032)
I + 1:
  if (condition) then
    PC ← PC + target_offset
  endif
```

Exceptions:

None

blezl

Branch on Less than or Equal to Zero Likely



MIPS II

Syntax:

```
blezl rs,offset
```

Description:

When register *rs* is less than or equal to zero, the program branches with a one instruction delay to the branch target address. The branch target address is the sum of the PC and the 16-bit offset after it is shifted left two bits and sign-extended to a 32 bit value. If the branch is not taken, the instruction in the branch delay slot is discarded.

Restrictions:

The `blezl` instruction cannot be placed in the delay slot of a branch/jump instruction.

Operation:

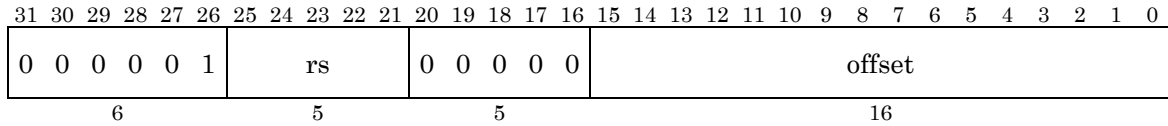
```
I :
    target_offset ← sign_extend(offset || 02)
    condition ← (GPR[rs] ≤ 032)
I + 1:
    if (condition) then
        PC ← PC + target_offset
    else
        NullifyCurrentInstruction()
    endif
```

Exceptions:

None

bltz

Branch on Less Than Zero



MIPS I

Syntax:

```
bltz rs,offset
```

Description:

When register *rs* is less than zero, the program branches with a one instruction delay to the branch target address.

The branch target address is the sum of the PC and the 16-bit offset after it is shifted left two bits and sign-extended to a 32 bit value.

Restrictions:

The *bltz* instruction cannot be placed in the delay slot of a branch/jump instruction.

Operation:

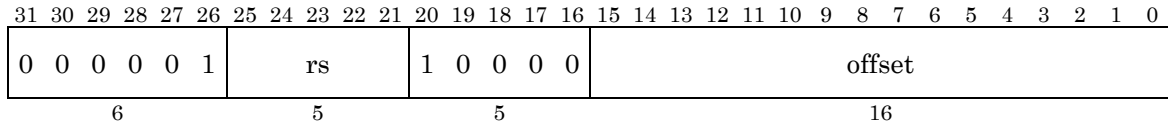
```
I :
    target_offset ← sign_extend(offset || 02)
    condition ← (GPR[rs] < 032)
I + 1 :
    if (condition) then
        PC ← PC + target_offset
    endif
```

Exceptions:

None

bltzal

Branch on Less Than Zero And Link



MIPS I

Syntax:

```
bltzal rs,offset
```

Description:

When register *rs* is less than zero, the program branches with a one instruction delay to the branch target address.

The branch target address is the sum of the PC and the 16-bit offset after it is shifted left two bits and sign-extended to a 32 bit value.

The address of the instruction following the delay slot is stored in register *r31* (the link register).

Restrictions:

The *bltzal* instruction cannot be placed in the delay slot of a branch/jump instruction.

The register *r31* (link register) cannot be specified for the register *rs*.

The instruction for changing the register *r31* (link register) cannot be placed in the delay slot of the *bltzal* instruction.

Operation:

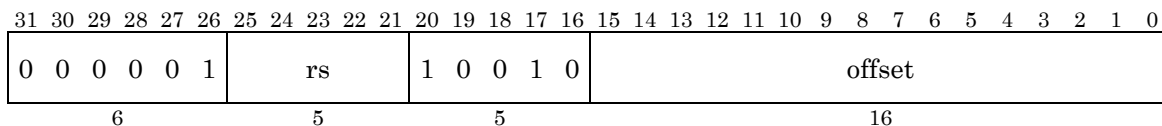
```
I :
  target_offset ← sign_extend(offset || 02)
  condition ← (GPR[rs] < 032)
  GPR[31] ← PC + 8
I + 1 :
  if (condition) then
    PC ← PC + target_offset
  endif
```

Exceptions:

None

bltzall

Branch on Less Than Zero And Link Likely



MIPS II

Syntax:

```
bltzall rs,offset
```

Description:

When register *rs* is less than zero, the program branches with a one instruction delay to the branch target address. The branch target address is the sum of the PC and the 16-bit offset after it is shifted left two bits and sign-extended to a 32 bit value. The address of the instruction following the delay slot is stored in register *r31* (the link register). If the branch is not taken, the instruction in the branch delay slot is discarded.

Restrictions:

The *bltzall* instruction cannot be placed in the delay slot of a branch/jump instruction.

The register *r31* (link register) cannot be specified for the register *rs*.

The instruction for changing the register *r31* (link register) cannot be placed in the delay slot of the *bltzall* instruction.

Operation:

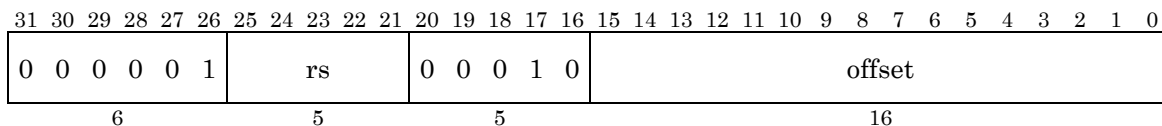
```
I :
  target_offset ← sign_extend(offset || 02)
  condition ← (GPR[rs] < 032)
  GPR[31] ← PC + 8
I + 1 :
  if (condition) then
    PC ← PC + target_offset
  else
    NullifyCurrentInstruction()
  endif
```

Exceptions:

None

bltzl

Branch on Less Than Zero Likely



MIPS II

Syntax:

```
bltzl rs,offset
```

Description:

When register *rs* is less than zero, the program branches with a one instruction delay to the branch target address. The branch target address is the sum of the PC and the 16-bit offset after it is shifted left two bits and sign-extended to a 32 bit value. If the branch is not taken, the instruction in the branch delay slot is discarded.

Restrictions:

The `bltzl` instruction cannot be placed in the delay slot of a branch/jump instruction.

Operation:

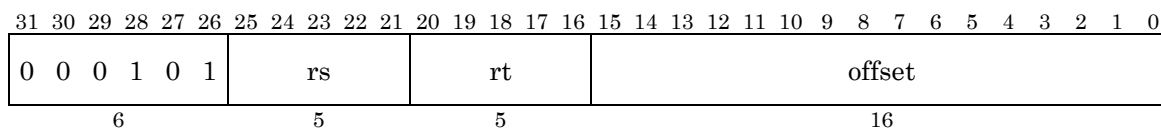
```
I :
  target_offset ← sign_extend(offset || 02)
  condition ← (GPR[rs] < 032)
I + 1:
  if (condition) then
    PC ← PC + target_offset
  else
    NullifyCurrentInstruction()
  endif
```

Exceptions:

None

bne

Branch on Not Equal



MIPS I

Syntax:

```
bne rs,rt,offset
```

Description:

When register *rs* is not equal to register *rt*, the program branches with a one instruction delay to the branch target address.

The branch target address is the sum of the PC and the 16-bit offset after it is shifted left two bits and sign-extended to a 32 bit value.

Restrictions:

The BNE instruction cannot be placed in the delay slot of a branch/jump instruction.

Operation:

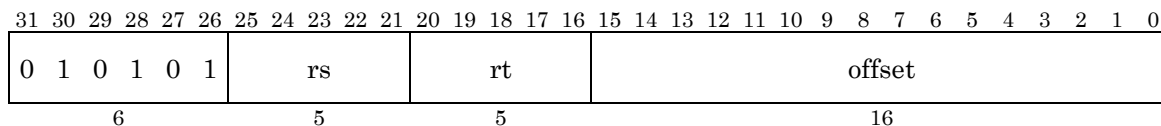
```
I :
    target_offset ← sign_extend(offset || 02)
    condition ← (GPR[rs] ≠ GPR[rt])
I + 1:
    if (condition) then
        PC ← PC + target_offset
    endif
```

Exceptions:

None

bnel

Branch on Not Equal Likely



MIPS II

Syntax:

```
bnel rs,rt,offset
```

Description:

When register *rs* is not equal to register *rt*, the program branches with a one instruction delay to the branch target address.

The branch target address is the sum of the PC and the 16-bit offset after it is shifted left two bits and sign-extended to a 32 bit value.

If the branch is not taken, the instruction in the branch delay slot is discarded.

Restrictions:

The BNEL instruction cannot be placed in the delay slot of a branch/jump instruction.

Operation:

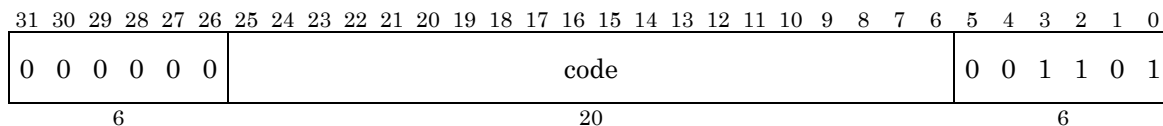
```
I :
    target_offset ← sign_extend(offset || 02)
    condition ← (GPR[rs] ≠ GPR[rt])
I + 1 :
    if (condition) then
        PC ← PC + target_offset
    else
        NullifyCurrentInstruction()
    endif
```

Exceptions:

None

break

Break Point



MIPS I

Syntax:

`break code`

Description:

A breakpoint exception is generated and control is immediately passed to the exception handler.

The code field can contain any arbitrary value.

Operation:

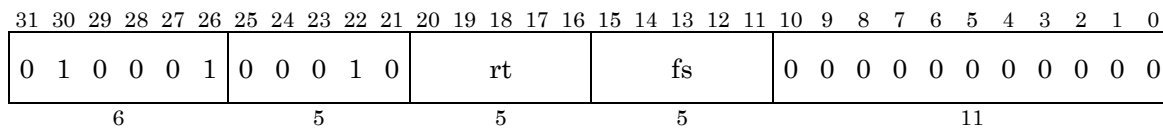
`SignalException(Breakpoint)`

Exceptions:

Break Point exception

cfc1

Move Control from FPU

MIPS I
FPU**Syntax:**`cfc1 rt, fs`**Description:**

The contents of FPU control register fs are transferred to CPU register rt. This instruction is only defined when fs is 0 or 31.

Restrictions:

Only 0 or 31 is valid for fs.

Operation:

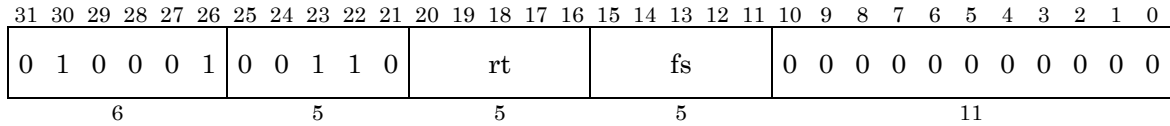
I:
temp ← FCR[fs]
I + 1:
GPR[rt] ← temp

Exceptions:

Coprocessor Unusable exception

ctc1

Move Control to FPU

MIPS I
FPU**Syntax:**`ctc1 rt, fs`**Description:**

The contents of CPU register `rt` are transferred to FPU control register `fs`.

This instruction is only defined when `fs` is 31.

A Floating-Point exception will occur if any Cause bit and its corresponding Enable bit of FCR31 get set to 1 as a result of executing this instruction.

After this instruction completes, the FPU pipeline will stall until the updated control register settings are reflected.

Restrictions:

Only 31 is valid for `fs`.

The `ctc1` instruction cannot be executed continuously after the `ctc1` instruction which causes an exception.

After executing the `ctc1` instruction, at least one `nop` is required until executing an instruction for referring to FCR31.

Operation:

```

I:
    temp ← GPR[rt]
I + 1:
    FCR[fs] ← temp
    COC[1] ← FCR[fs]23

```

Exceptions:

Coprocessor Unusable exception

Floating-Point exception

FPU Exceptions:

Unimplemented Operation exception

Invalid Operation exception

Division By Zero exception

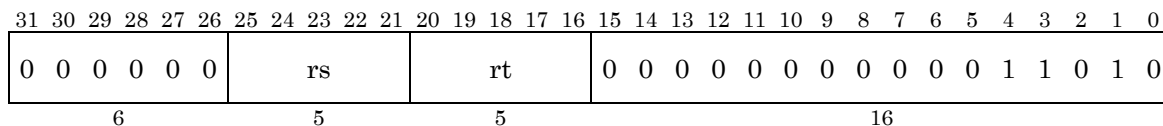
Inexact exception

Overflow exception

Underflow exception

div

Divide



MIPS I

Syntax:

```
div rs,rt
```

Description:

The contents of register *rs* are divided by the contents of register *rt*. The 32-bit quotient is stored in special register *LO*, and the 32-bit remainder is stored in special register *HI*. The operands are treated as signed integers.

Operation:

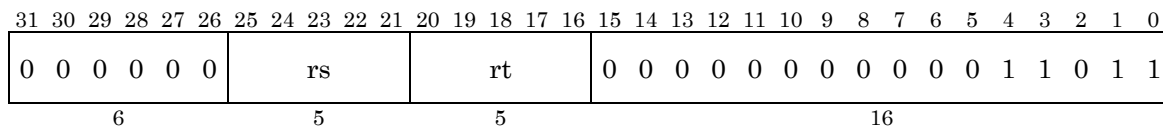
```
q ← GPR[rs] / GPR[rt]
LO ← q
r ← GPR[rs] % GPR[rt]
HI ← r
```

Exceptions:

None

divu

Divide Unsigned



MIPS I

Syntax:

```
divu rs,rt
```

Description:

The contents of register *rs* are divided by the contents of register *rt*.

The 32-bit quotient is stored in special register *LO*, and the 32-bit remainder is stored in special register *HI*. The operands are treated as unsigned integers.

Operation:

$$q \leftarrow (0 \parallel \text{GPR}[rs]) / (0 \parallel \text{GPR}[rt])$$

$$LO \leftarrow q$$

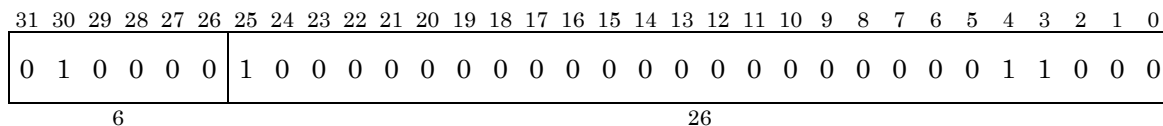
$$r \leftarrow (0 \parallel \text{GPR}[rs]) \% (0 \parallel \text{GPR}[rt])$$

$$HI \leftarrow r$$
Exceptions:

None

eret

Exception Return



MIPS III

Syntax:

eret

Description:

Return from interrupt, exception, or error. The LLbit is cleared to 0.

There is no delay slot of a branch/jump instruction, and the instruction immediately following the eret instruction is discarded.

Restrictions:

Do not place this instruction in the delay slot of a branch/jump instruction.

Two nops are required when an eret instruction follows an mtc0 instruction.

Operation:

```

if (StatusERL = 1) then
    PC ← ErrorEPC
    StatusERL ← 0
else
    PC ← EPC
    StatusEXL ← 0
endif
LLbit ← 0

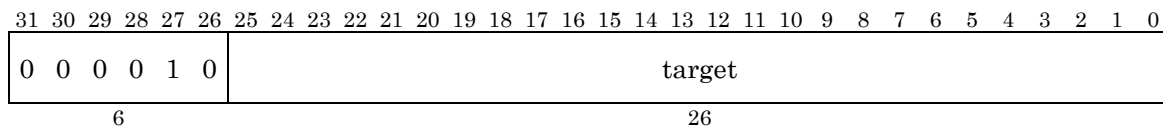
```

Exceptions:

Coprocessor Unusable exception

j

Jump



MIPS I

Syntax:

j target

Description:

The program jumps with a one instruction delay to the destination address. The destination address is formed by shifting the 26-bit target left two bits and concatenating it with the high-order 4 bits of the PC.

Operation:

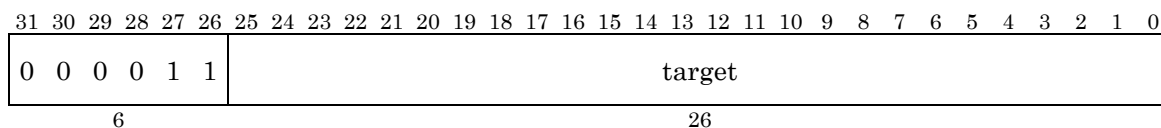
I :
temp \leftarrow target
I + 1:
PC \leftarrow PC_{31..28} || temp || 0²

Exceptions:

None

jal

Jump And Link



MIPS I

Syntax:

jal target

Description:

The program jumps with a one instruction delay to the destination address. The destination address is formed by shifting the 26-bit target left two bits and concatenating it with the high-order 4 bits of the PC. The address of the instruction following the delay slot is stored in r31 (the link register).

Restrictions:

The instruction for changing the register r31 (link register) cannot be placed in the delay slot of the jal instruction.

Operation:

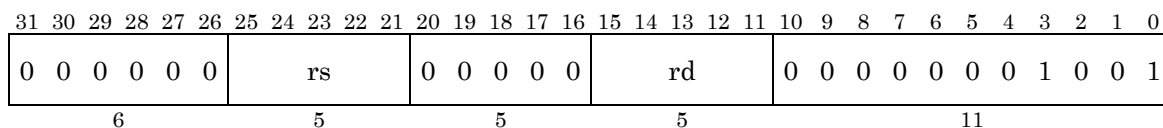
I :
 $GPR[31] \leftarrow PC + 8$
 $temp \leftarrow target$
 I + 1 :
 $PC \leftarrow PC_{31..28} || temp || 0^2$

Exceptions:

None

jalr

Jump And Link Register



MIPS I

Syntax:

```
jalr rd,rs
```

Description:

The program jumps with a one instruction delay to the address in the rs register. The address of the instruction following the delay slot is stored in register rd. If rd is not specified, it is assumed to be r31 by default.

Restrictions:

The same register cannot be specified for rd and rs.

The instruction for changing the register rd or the register rs cannot be placed in the delay slot of the jalr instruction.

Operation:

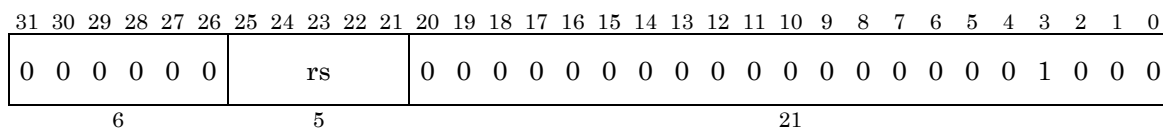
```
I :
    temp ← GPR[rs]
    GPR[rd] ← PC + 8
I + 1:
    PC ← temp
```

Exceptions:

None

jr

Jump Register



MIPS I

Syntax:

jr rs

Description:

The program jumps with a one instruction delay to the address in the rs register.

Operation:

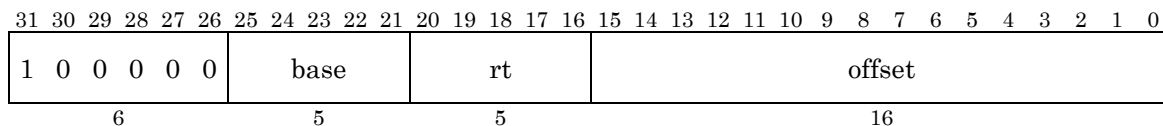
I:
temp ← GPR[rs]
I + 1:
PC ← temp

Exceptions:

None

lb

Load Byte



MIPS I

Syntax:

```
lb rt,offset(base)
```

Description:

The 16-bit offset is sign-extended and added to the contents of the base register to generate an address. The byte in memory at that address is sign-extended and loaded into register `rt`.

Operation:

```
vAddr ← sign_extend(offset) + GPR[base]
(pAddr, CCA) ← AddressTranslation(vAddr, DATA, LOAD)
pAddr ← pAddr_PSIZE-1.2 || (pAddr_1.0 xor ReverseEndian2)
memword ← LoadMemory(CCA, BYTE, pAddr, vAddr, DATA)
byte ← vAddr_1.0 xor ReverseEndian2
GPR[rt] ← sign_extend(memword7+8*byte..8*byte)
```

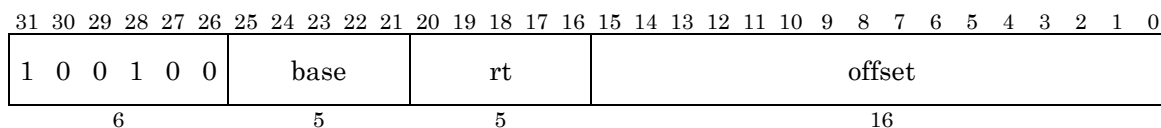
Exceptions:

Address Error exception

Bus Error exception

lbu

Load Byte Unsigned



MIPS I

Syntax:

```
lbu rt, offset(base)
```

Description:

The 16-bit offset is sign-extended and added to the contents of the base register to generate an address.

The byte in memory at that address is zero-extended and loaded into register rt.

Operation:

```
vAddr ← sign_extend(offset) + GPR[base]
(pAddr, CCA) ← AddressTranslation(vAddr, DATA, LOAD)
pAddr ← pAddrPSize-1..2 || (pAddr1..0 xor ReverseEndian2)
memword ← LoadMemory(CCA, BYTE, pAddr, vAddr, DATA)
byte ← vAddr1..0 xor ReverseEndian2
GPR[rt] ← zero_extend(memword7+8*byte..8*byte)
```

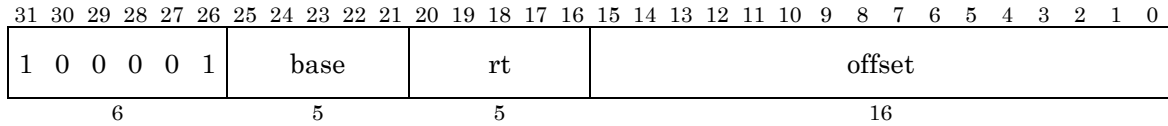
Exceptions:

Address Error exception

Bus Error exception

lh

Load Halfword



MIPS I

Syntax:

```
lh rt, offset(base)
```

Description:

The 16-bit offset is sign-extended and added to the contents of the base register to generate an address.

The halfword in memory at that address is sign-extended and loaded into register rt.

Operation:

```
vAddr ← sign_extend(offset) + GPR[base]
if (vAddr0 ≠ 0) then
    SignalException(AddressError)
endif
(pAddr, CCA) . ← AddressTranslation(vAddr, DATA, LOAD)
pAddr ← pAddrPSize-1..2 || (pAddr1..0 xor (ReverseEndian || 0))
memword ← LoadMemory(CCA, HALFWORD, pAddr, vAddr, DATA)
byte ← vAddr1..0 xor (ReverseEndian || 0)
GPR[rt] ← sign_extend(memword15+8*byte..8*byte)
```

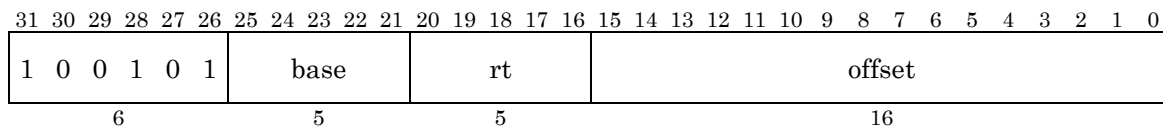
Exceptions:

Address Error exception

Bus Error exception

lhu

Load Halfword Unsigned



MIPS I

Syntax:

```
lhu  rt, offset(base)
```

Description:

The 16-bit offset is sign-extended and added to the contents of the base register to generate an address.

The halfword in memory at that address is zero-extended and loaded into register rt.

Operation:

```
vAddr ← sign_extend(offset) + GPR[base]
if (vAddr0 ≠ 0) then
    SignalException(AddressError)
endif
(pAddr, CCA) ← AddressTranslation(vAddr, DATA, LOAD)
pAddr ← pAddrPSize-1..2 || (pAddr1..0 xor (ReverseEndian || 0))
memword ← LoadMemory(CCA, HALFWORD, pAddr, vAddr, DATA)
byte ← vAddr1..0 xor (ReverseEndian || 0)
GPR[rt] ← zero_extend(memword15+8*byte..8*byte)
```

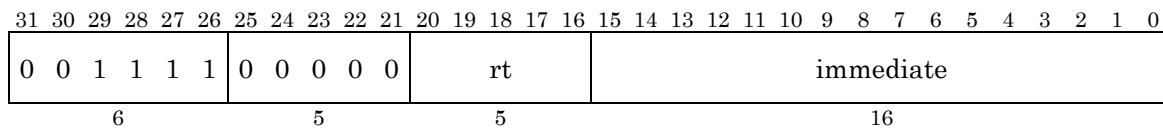
Exceptions:

Address Error exception

Bus Error exception

lui

Load Upper Immediate



MIPS I

Syntax:

```
lui rt,immediate
```

Description:

The 16-bit immediate value is shifted left by 16 bits to produce a 32-bit word. The low-order 16 bits are set to 0 and the result is stored in register rt.

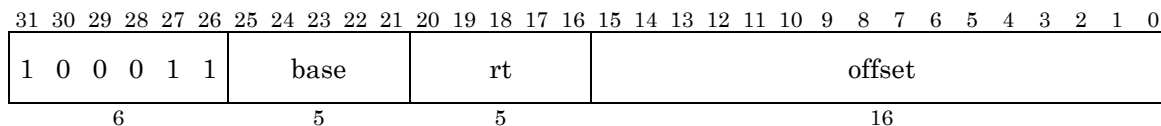
Operation:

$$\text{GPR}[\text{rt}] \leftarrow \text{immediate} \ll 16$$
Exceptions:

None

lw

Load Word



MIPS I

Syntax:

```
lw rt, offset(base)
```

Description:

The 16-bit offset is sign-extended and added to the contents of the base register to generate an address.

The word in memory at that address is loaded into register *rt*.

Operation:

```
vAddr ← sign_extend(offset) + GPR[base]
if (vAddr1:0 ≠ 02) then
    SignalException(AddressError)
endif
(pAddr, CCA) ← AddressTranslation(vAddr, DATA, LOAD)
memword ← LoadMemory(CCA, WORD, pAddr, vAddr, DATA)
GPR[rt] ← memword
```

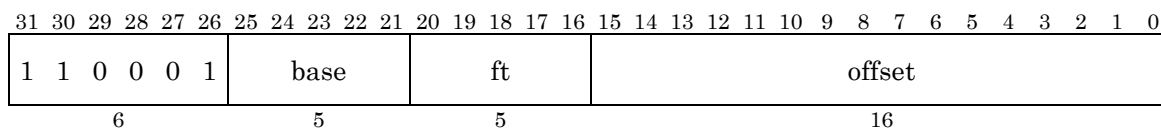
Exceptions:

Address Error exception

Bus Error exception

lwc1

Load Word to FPU

MIPS I
FPU**Syntax:**

```
lwc1 ft,offset(base)
```

Description:

The 16-bit offset is sign-extended and added to the contents of the base register to generate an address.

The word in memory at that address is loaded into FPU register ft.

Operation:

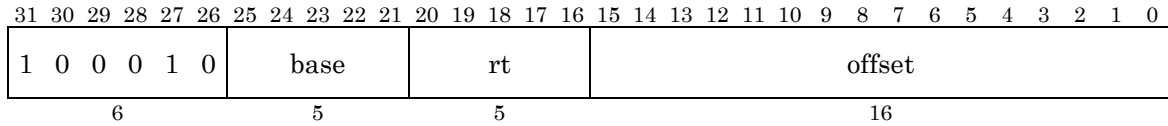
```
vAddr ← sign_extend(offset) + GPR[base];
if (vAddr1:0 ≠ 02) then
    SignalException(AddressError)
endif
(pAddr, CCA) ← AddressTranslation(vAddr, DATA, LOAD)
memword ← LoadMemory(CCA, WORD, pAddr, vAddr, DATA)
StoreFPR(ft, UNINTERPRETED_WORD, memword)
```

Exceptions:

```
Coprocessor Unusable exception
Address Error exception
Bus Error exception
```

lwl

Load Word Left



MIPS I

Syntax:

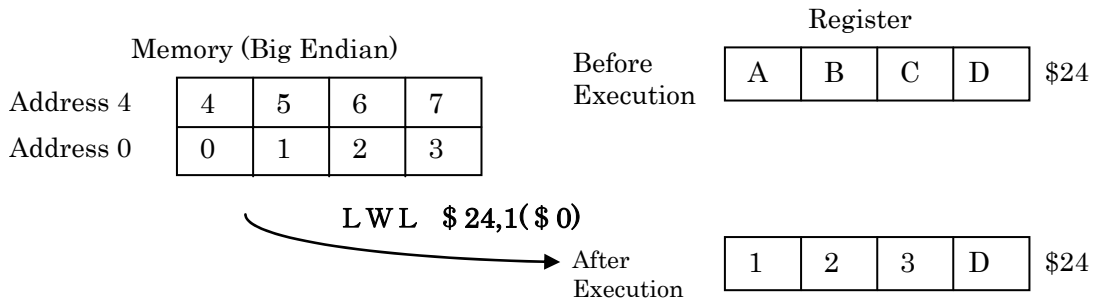
```
lwl rt,offset(base)
```

Description:

The 16-bit offset is sign-extended and added to the contents of the base register to generate an address.

The word in memory at that address is shifted left so that the byte at that address is at the leftmost end of the word.

The result of the shift operation is merged into register rt.

**Operation:**

```
vAddr ← sign_extend(offset) + GPR[base]
(pAddr, CCA) ← AddressTranslation(vAddr, DATA, LOAD)
if (BigEndianMem=0) then
    pAddr ← pAddr_PSIZE - 1.2 || 02
endif
byte ← vAddr1..0 xor ReverseEndian2
memword ← LoadMemory(CCA, byte, pAddr, vAddr, DATA)
temp ← memword7 + 8*byte..0 || GPR[rt]23 - 8*byte..0
GPR[rt] ← temp
```

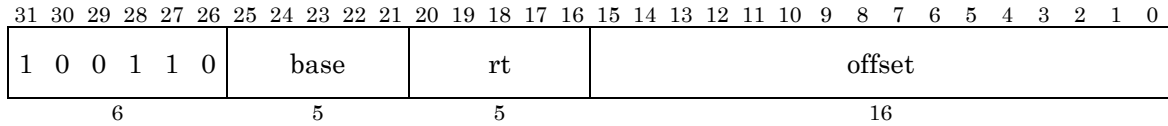
Exceptions:

Address Error exception

Bus Error exception

lwr

Load Word Right



MIPS I

Syntax:

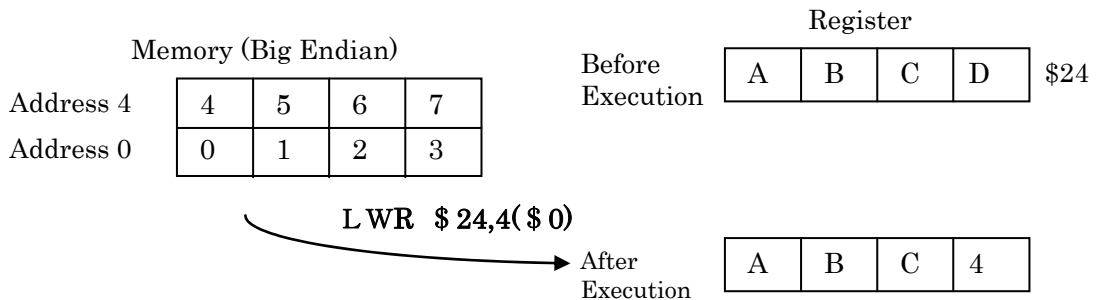
```
lwr rt, offset(base)
```

Description:

The 16-bit offset is sign-extended and added to the contents of the base register to generate an address.

The word in memory at that address is shifted right so that the byte at that address is at the rightmost end of the word.

The result of the shift operation is merged into register rt.

**Operation:**

```
vAddr ← sign_extend(offset) + GPR[base]
(pAddr, CCA) ← AddressTranslation(vAddr, DATA, LOAD)
if (BigEndianMem=0) then
    pAddr ← pAddr_PSIZE - 1..2 || 02
endif
byte ← vAddr1..0 xor ReverseEndian2
memword ← LoadMemory(CCA, byte, pAddr, vAddr, DATA)
temp ← memword31..32 - 8*byte..0 || GPR[rt]31 - 8*byte..0
GPR[rt] ← temp
```

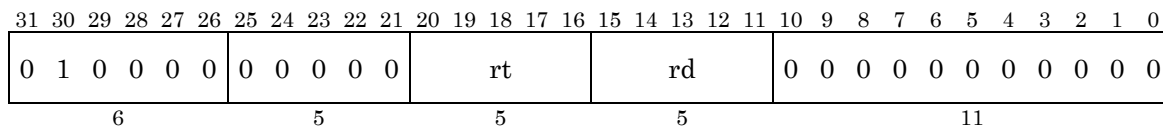
Exceptions:

Address Error exception

Bus Error exception

mfc0

Move From Coprocessor0



MIPS I

Syntax:

mfc0 rt,rd

Description:

The contents of CP0 register rd are transferred to CPU register rt.

Operation:

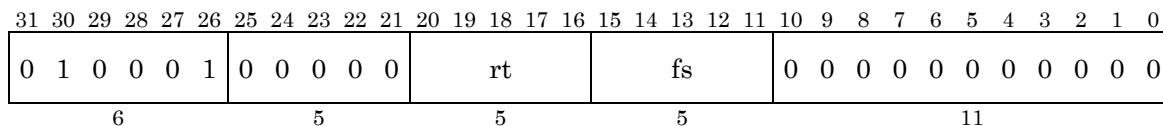
$$\text{data} \leftarrow \text{CP0R}[\text{rd}]$$

$$\text{GPR}[\text{rt}] \leftarrow \text{data}$$
Exceptions:

Coprocessor Unusable exception

mfc1

Move From FPU

MIPS I
FPU**Syntax:**

```
mfc1 rt, fs
```

Description:

The contents of FPU general purpose register fs are transferred to CPU register rt.

Operation:

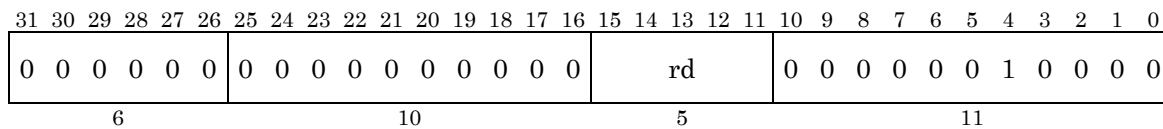
```
data ← ValueFPR(fs, UNINTERPRETED_WORD)
GPR[rt] ← data
```

Exceptions:

Coprocessor Unusable exception

mfhi

Move From HI Register



MIPS I

Syntax:

mfhi rd

Description:

The contents of special register HI are transferred to register rd.

Operation:

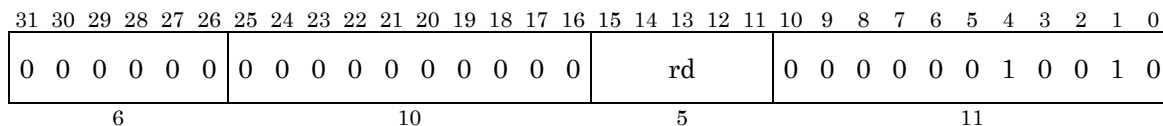
GPR[rd] ← HI

Exceptions:

None

mflo

Move From LO Register



MIPS I

Syntax:

`mflo rd`

Description:

The contents of special register LO are transferred to register rd.

Operation:

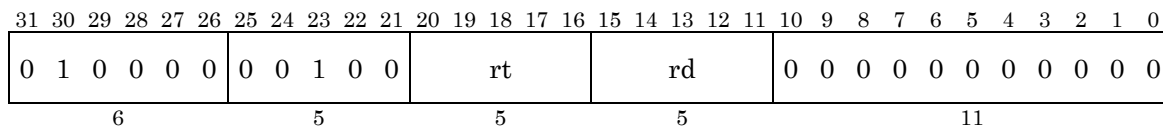
$GPR[rd] \leftarrow LO$

Exceptions:

None

mtc0

Move To Coprocessor0



MIPS I

Syntax:

```
mtc0 rt, rd
```

Description:

The contents of CPU register *rt* are transferred to CP0 register *rd*.

Operation:

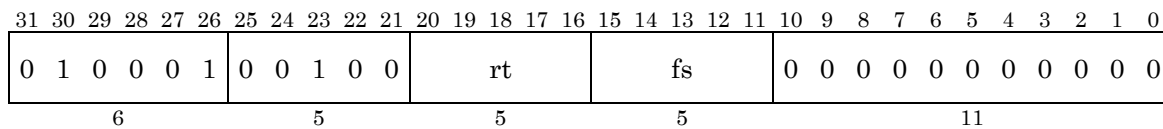
```
DATA ← GPR[rt]
CP0R[rd] ← DATA
```

Exceptions:

Coprocessor Unusable exception

mtc1

Move To FPU

MIPS I
FPU**Syntax:**`mtc1 rt, fs`**Description:**

The contents of CPU register `rt` are transferred to FPU general-purpose register `fs`.

Operation:

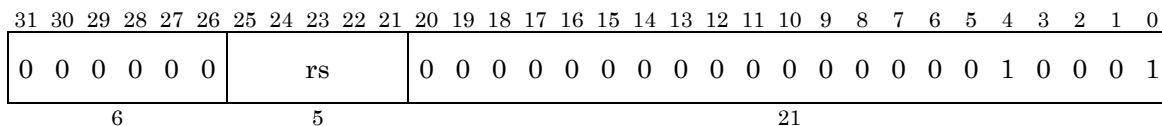
```
data ← GPR[rt]31..0
StoreFPR(fs, UNINTERPRETED_WORD, data)
```

Exceptions:

Coprocessor Unusable exception

mthi

Move To HI Register



MIPS I

Syntax:

mthi rs

Description:

The contents of register rs are transferred to special register HI.

Operation:

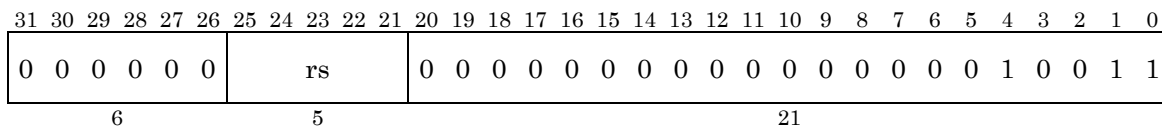
HI ← GPR[rs]

Exceptions:

None

mtlo

Move To LO Register



MIPS I

Syntax:

mtlo rs

Description:

The contents of register rs are transferred to special register LO.

Operation:

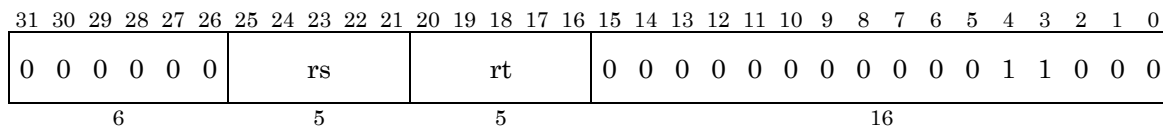
LO ← GPR[rs]

Exceptions:

None

mult

Multiply



MIPS I

Syntax:

```
mult rs,rt
```

Description:

The contents of register *rs* are multiplied by the contents of register *rt* and the 64-bit result is stored in special registers HI and LO. The operands are treated as signed integers.

Operation:

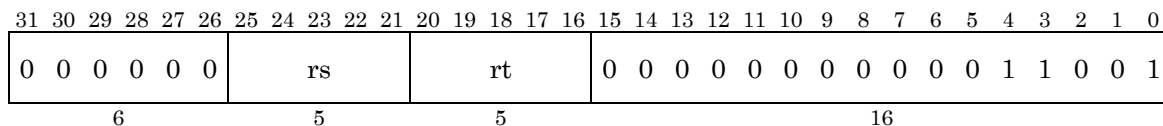
```
temp ← (GPR[rs] × GPR[rt])
HI ← temp63..32
LO ← temp31..0
```

Exceptions:

None

multu

Multiply Unsigned



MIPS I

Syntax:

```
multu rs,rt
```

Description:

The contents of register *rs* are multiplied by the contents of register *rt* and the 64-bit result is stored in special registers HI and LO. The operands are treated as unsigned integers.

Operation:

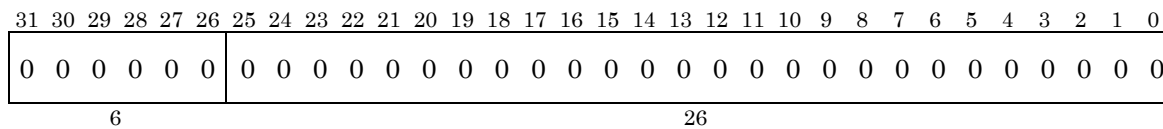
```
temp ← (GPR[rs] × GPR[rt])
HI ← temp63..32
LO ← temp31..0
```

Exceptions:

None

nop

No Operation



MIPS I

Syntax:

`nop`

Description:

This instruction does nothing. It is interpreted by the actual hardware as "SLL r0, r0, 0".

Operation:

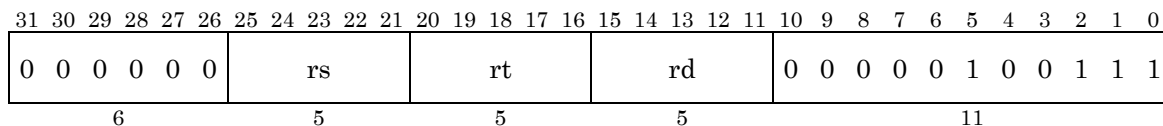
None

Exceptions:

None

nor

NOR



MIPS I

Syntax:

```
nor rd, rs, rt
```

Description:

A bitwise NOR (negative OR) is performed for the contents of registers *rs* and *rt* and the result is stored in register *rd*.

X	Y	X NOR Y
0	0	1
0	1	0
1	0	0
1	1	0

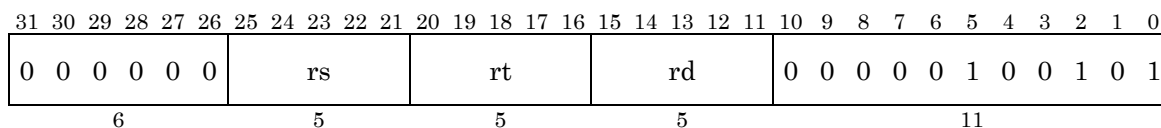
Operation:

$$\text{GPR}[\text{rd}] \leftarrow \text{GPR}[\text{rs}] \text{ nor } \text{GPR}[\text{rt}]$$
Exceptions:

None

or

OR



MIPS I

Syntax:

or rd,rs,rt

Description:

A bitwise OR (logical sum) is performed for the contents of registers rs and rt and the result is stored in register rd.

X	Y	X OR Y
0	0	0
0	1	1
1	0	1
1	1	1

Operation:

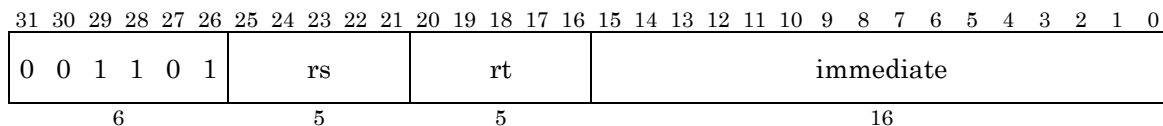
GPR[rd] ← GPR[rs] or GPR[rt]

Exceptions:

None

ori

OR Immediate



MIPS I

Syntax:

```
ori rt,rs,immediate
```

Description:

The 16-bit immediate field is zero-extended and a bitwise OR (logical sum) is performed with the contents of register rs. The result is stored in register rt.

X	Y	X OR Y
0	0	0
0	1	1
1	0	1
1	1	1

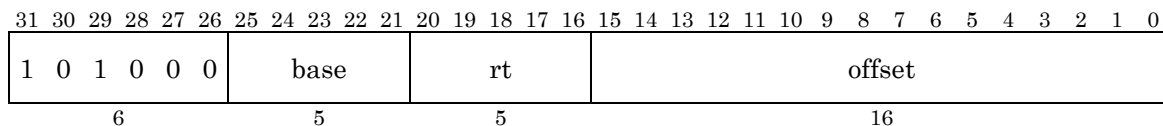
Operation:

$$\text{GPR}[rt] \leftarrow \text{GPR}[rs] \text{ or } \text{zero_extend}(\text{immediate})$$
Exceptions:

None

sb

Store Byte



MIPS I

Syntax:

```
sb rt, offset(base)
```

Description:

The 16-bit offset is sign-extended and added to the contents of the base register to generate an address. The low-order byte of register `rt` is stored in memory at that address.

Bus error exceptions will not be reported because writing is done via a buffer. If a bus error occurs, it is handled by the system as an interrupt.

Operation:

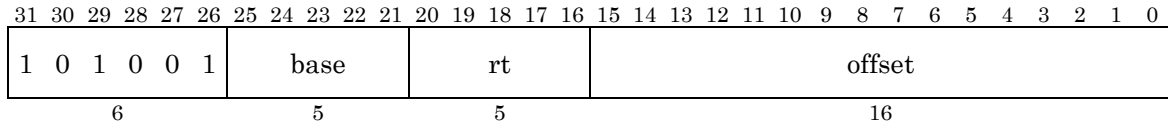
```
vAddr ← sign_extend(offset) + GPR[base]
(pAddr, CCA) ← AddressTranslation(vAddr, DATA, STORE)
pAddr ← pAddrPSIZE - 1..2 || (pAddr1..0 xor ReverseEndian2)
bytesel ← vAddr1..0 xor ReverseEndian2
dataword ← GPR[rt]31 - 8*bytesel..0 || 08*bytesel
StoreMemory(CCA, BYTE, dataword, pAddr, vAddr, DATA)
```

Exceptions:

```
Address Error exception
```


sh

Store HalfWord



MIPS I

Syntax:

```
sh rt,offset(base)
```

Description:

The 16-bit offset is sign-extended and added to the contents of the base register to generate an address. The low-order halfword of register `rt` is stored in memory at that address.

Bus error exceptions will not be reported because writing is done via a buffer. If a bus error occurs, it is handled by the system as an interrupt.

Operation:

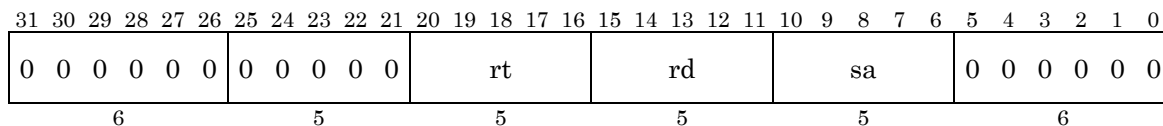
```
vAddr ← sign_extend(offset) + GPR[base]
if (vAddr0 ≠ 0) then
    SignalException(AddressError)
endif
(pAddr, CCA) ← AddressTranslation(vAddr, DATA, STORE)
pAddr ← pAddrPSIZE-1..2 || (pAddr1..0 xor (ReverseEndian || 0))
bytesel ← vAddr1..0 xor (ReverseEndian || 0)
dataword ← GPR[rt]31-8*bytesel..0 || 08*bytesel
StoreMemory(CCA, HALFWORD, dataword, pAddr, vAddr, DATA)
```

Exceptions:

Address Error exception

sll

Shift Left Logical



MIPS I

Syntax:

```
sll rd, rt, sa
```

Description:

The contents of register *rt* are shifted left *sa* bits. Zeroes are inserted into the low-order bit positions from the right.

The 32-bit result is stored in register *rd*.

Operation:

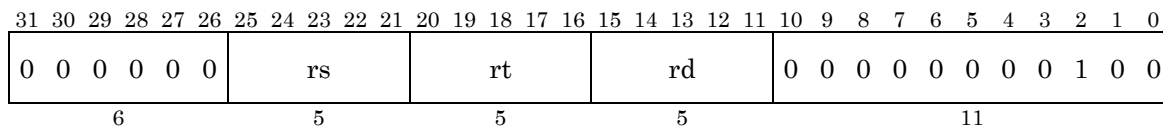
```
s ← sa
temp ← GPR[rt](31-s)..0 || 0s
GPR[rd] ← temp
```

Exceptions:

None

sllv

Shift Left Logical Variable



MIPS I

Syntax:

```
sllv rd,rt,rs
```

Description:

The contents of register *rt* are shifted left. Zeroes are inserted into the low-order bit positions from the right. The shift amount is specified by the low-order 5 bits of register *rs*.

The 32-bit result is stored in register *rd*.

Operation:

$$s \leftarrow \text{GPR}[\text{rs}]_{4..0}$$

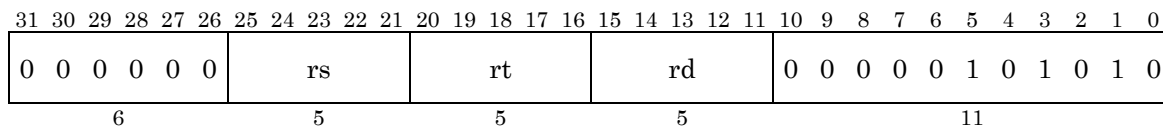
$$\text{temp} \leftarrow \text{GPR}[\text{rt}]_{(31-s)..0} \parallel 0^s$$

$$\text{GPR}[\text{rd}] \leftarrow \text{temp}$$
Exceptions:

None

slt

Set on Less Than



MIPS I

Syntax:

```
slt rd, rs, rt
```

Description:

The contents of registers *rs* and *rt* are compared as 32-bit signed integers. If the comparison result is $rs < rt$, 1 is stored in register *rd*, otherwise 0 is stored.

Operation:

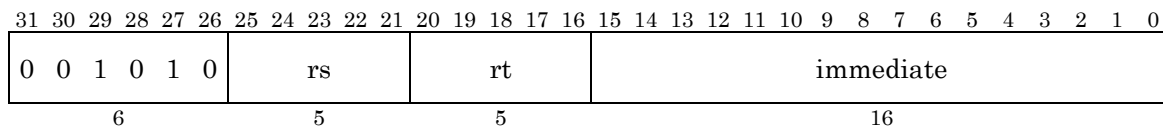
```
if (GPR[rs] < GPR[rt]) then
    GPR[rd] ← 031 || 1
else
    GPR[rd] ← 032
endif
```

Exceptions:

None

slti

Set on Less Than Immediate



MIPS I

Syntax:

```
slti rt,rs,immediate
```

Description:

The contents of register *rs* and the 16-bit immediate value sign-extended to 32 bits are compared as 32-bit signed integers. If the comparison result is $rs < \text{immediate}$, 1 is stored in register *rt*, otherwise 0 is stored.

Operation:

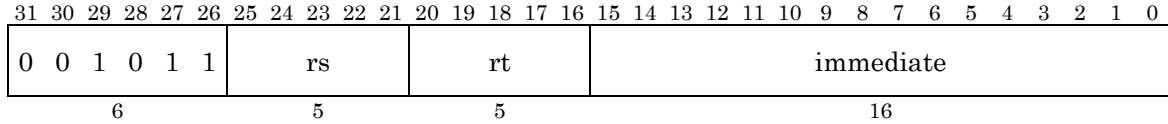
```
if (GPR[rs] < sign_extend(immediate)) then
    GPR[rt] ← 031 || 1
else
    GPR[rt] ← 032
endif
```

Exceptions:

None

sltiu

Set on Less Than Immediate Unsigned



MIPS I

Syntax:

```
sltiu  rt,rs,immediate
```

Description:

The contents of register *rs* and the sign-extended 16-bit immediate value are compared as 32-bit unsigned integers. If the comparison result is $rs < \text{immediate}$, 1 is stored in register *rt*, otherwise 0 is stored.

Because the 16-bit immediate value is sign-extended before comparison, the range of numeric values that the immediate represents is not sequential, but split into two areas: around the smallest and largest 32-bit unsigned integers. That is $[0, 32767]$ and $[\text{max_unsigned} - 32767, \text{max_unsigned}]$, respectively.

Operation:

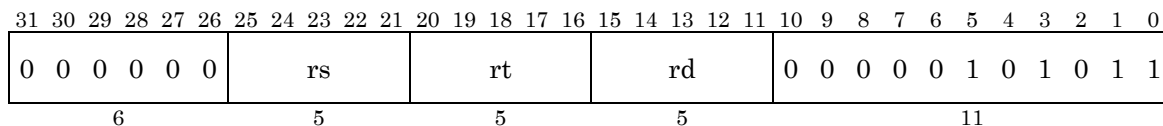
```
if ((0 || GPR[rs]) < (0 || sign_extend(immediate))).then
    GPR[rt] ← 031 || 1
else
    GPR[rt] ← 032
endif
```

Exceptions:

None

sltu

Set on Less Than Unsigned



MIPS I

Syntax:

```
sltu rd,rs,rt
```

Description:

The contents of registers *rs* and *rt* are compared as 32-bit unsigned integers. If the comparison result is $rs < rt$, 1 is stored in register *rd*, otherwise 0 is stored.

Operation:

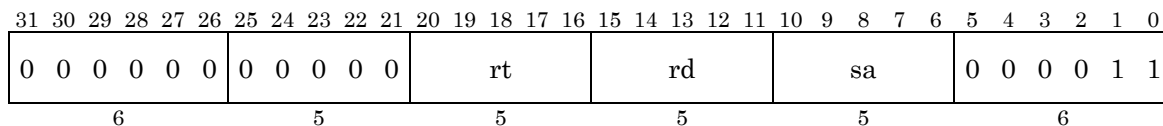
```
if ((0 || GPR[rs]) < (0 || GPR[rt])) then
    GPR[rd] ← 031 || 1
else
    GPR[rd] ← 032
endif
```

Exceptions:

None

sra

Shift Right Arithmetic



MIPS I

Syntax:

```
sra rd, rt, sa
```

Description:

The contents of register *rt* are shifted right *sa* bits. The sign is extended into the high-order bit positions.

The 32-bit result is stored in register *rd*.

Operation:

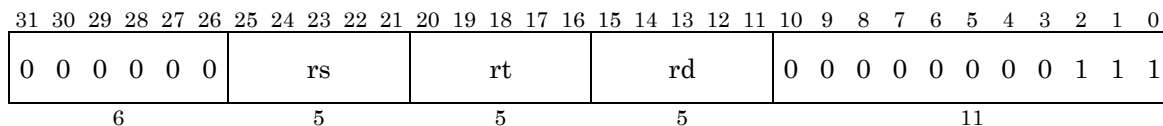
```
s ← sa
temp ← ((GPR[rt]31)s || GPR[rt]31..s)
GPR[rd] ← temp
```

Exceptions:

None

srav

Shift Right Arithmetic Variable



MIPS I

Syntax:

```
srav rd, rt, rs
```

Description:

The contents of register *rt* are shifted right. The sign is extended into the high-order bit positions. The shift amount is specified by the low-order 5 bits of register *rs*.

The 32-bit result is stored in register *rd*.

Operation:

$$s \leftarrow \text{GPR}[\text{rs}]_{4..0}$$

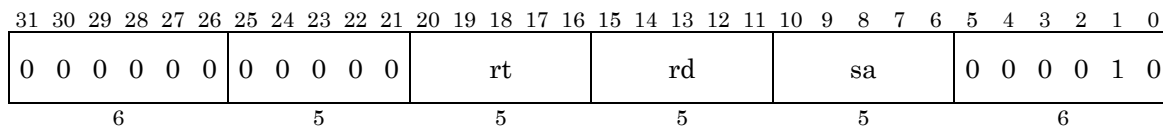
$$\text{temp} \leftarrow (\text{GPR}[\text{rt}]_{31})^s \parallel \text{GPR}[\text{rt}]_{31..s}$$

$$\text{GPR}[\text{rd}] \leftarrow \text{temp}$$
Exceptions:

None

srl

Shift Right Logical



MIPS I

Syntax:

```
srl rd, rt, sa
```

Description:

The contents of register *rt* are shifted right *sa* bits. Zeroes are inserted into the high-order bit positions from the left.

The 32-bit result is stored in register *rd*.

Operation:

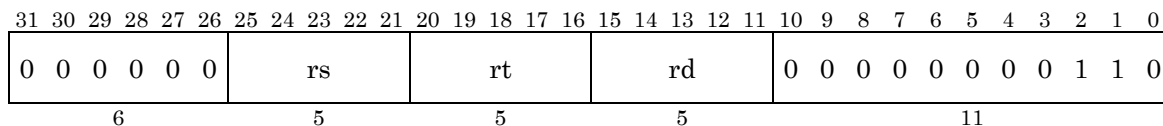
```
s ← sa
temp ← (0s || GPR[rt]31..s)
GPR[rd] ← temp
```

Exceptions:

None

srlv

Shift Right Logical Variable



MIPS I

Syntax:

```
srlv rd,rt,rs
```

Description:

The contents of register *rt* are shifted right. Zeroes are inserted into the high-order bit positions from the left. The shift amount is specified by the low-order 5 bits of register *rs*.

The 32-bit result is stored in register *rd*.

Operation:

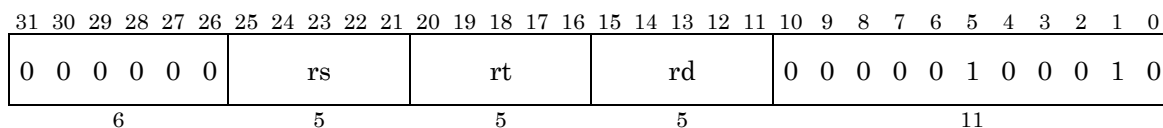
```
s ← GPR[rs]4..0
temp ← (0s || GPR[rt]31..s)
GPR[rd] ← temp
```

Exceptions:

None

sub

Subtract



MIPS I

Syntax:

```
sub rd, rs, rt
```

Description:

The contents of register *rt* are subtracted from the contents of register *rs* and the result is stored in register *rd*.

If a two's complement overflow occurs, an exception is generated.

Operation:

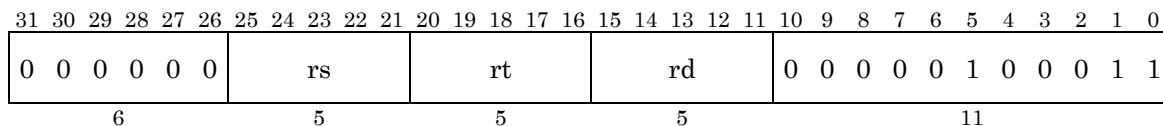
```
temp ← (GPR[rs]31 || GPR[rs]31..0) - (GPR[rt]31 || GPR[rt]31..0)
if (temp32 ≠ temp31) then
    SignalException(IntegerOverflow)
else
    GPR[rd] ← temp
endif
```

Exceptions:

Integer Overflow exception

subu

Subtract Unsigned



MIPS I

Syntax:

```
subu rd,rs,rt
```

Description:

The contents of register *rt* are subtracted from the contents of register *rs* and the result is stored in register *rd*.

No exception is generated even if a two's complement overflow occurs.

Operation:

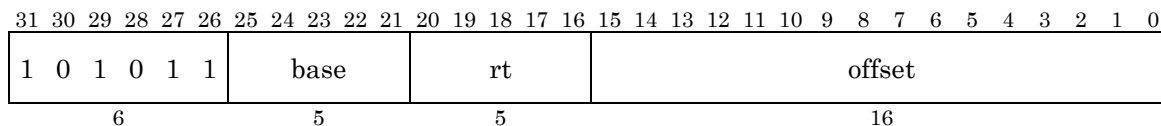
```
temp ← GPR[rs] - GPR[rt]
GPR[rd] ← temp
```

Exceptions:

None

SW

Store Word



MIPS I

Syntax:

```
sw rt, offset(base)
```

Description:

The 16-bit offset is sign-extended and added to the contents of the base register to generate an address. The contents of register *rt* are stored in memory at that address. Bus error exceptions will not be reported because writing is done via a buffer. If a bus error occurs, it is handled by the system as an interrupt.

Operation:

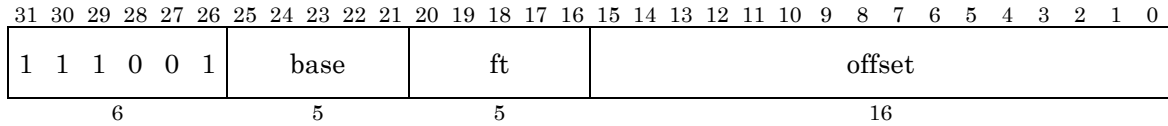
```
vAddr ← sign_extend(offset) + GPR[base]
if (vAddr1:0 ≠ 02) then
    SignalException(AddressError)
endif
(pAddr, CCA) ← AddressTranslation(vAddr, DATA, STORE)
dataword ← GPR[rt]
StoreMemory(CCA, WORD, dataword, pAddr, vAddr, DATA)
```

Exceptions:

```
Address Error exception
```

swc1

Store Word from FPU

MIPS I
FPU**Syntax:**

```
swc1 ft,offset(base)
```

Description:

The 16-bit offset is sign-extended and added to the contents of the base register to generate an address. The contents of FPU register ft are stored in memory at that address.

Bus error exceptions will not be reported because writing is done via a buffer. If a bus error occurs, it is handled by the system as an interrupt.

Operation:

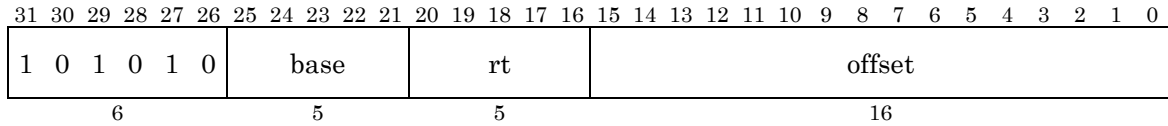
```
vAddr ← sign_extend(offset) + GPR[base]
if (vAddr1:0 ≠ 02) then
    SignalException(AddressError)
endif
(pAddr, CCA) ← AddressTranslation(vAddr, DATA, STORE)
dataword ← ValueFPR(ft, UNINTERPRETED_WORD)
StoreMemory(CCA, WORD, dataword, pAddr, vAddr, DATA)
```

Exceptions:

```
Coprocessor Unusable exception
Address Error exception
```

swl

Store Word Left



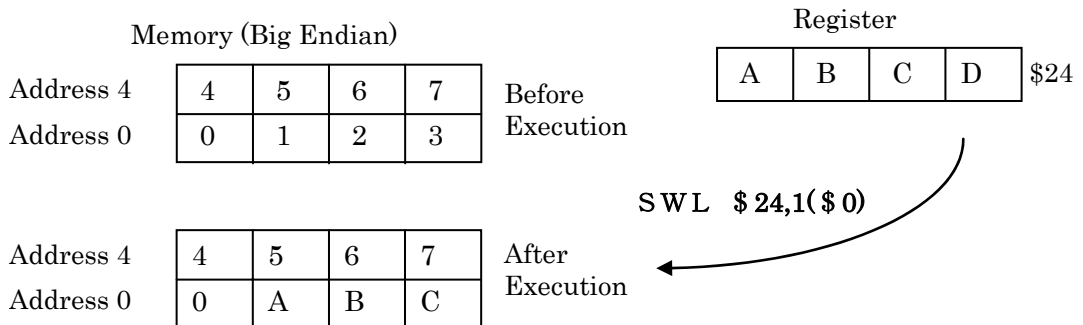
MIPS I

Syntax:

```
swl rt,offset(base)
```

Description:

The 16-bit offset is sign-extended and added to the contents of the base register to generate an address. The contents of register *rt* are shifted right so that the leftmost byte is in the same position in the word as the byte at the generated address. The byte (bytes) that contains the original data of register *rt* is stored in the byte (bytes) from the byte position of the specified address to the word boundary on the right side. Bus error exceptions are not reported because writing is done via a buffer. If a bus error occurs, it is handled by the system as an interrupt.

**Operation:**

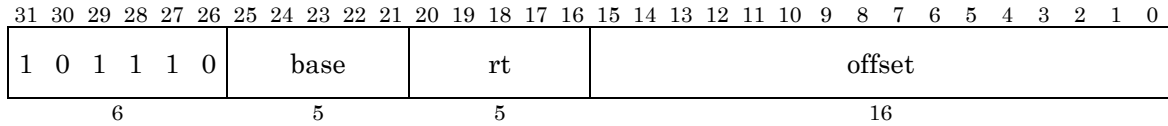
```
vAddr ← sign_extend(offset) + GPR[base]
(pAddr, CCA) ← AddressTranslation(vAddr, DATA, STORE)
if (BigEndianMem=0) then
    pAddr ← pAddr_PSIZE - 1..2 || 02
endif
byte ← vAddr1..0 xor ReverseEndian2
dataword ← 024 - 8*byte || GPR[rt]31..24 - 8*byte
StoreMemory(CCA, byte, dataword, pAddr, vAddr, DATA)
```


Exceptions:

Address Error exception

SWR

Store Word Right



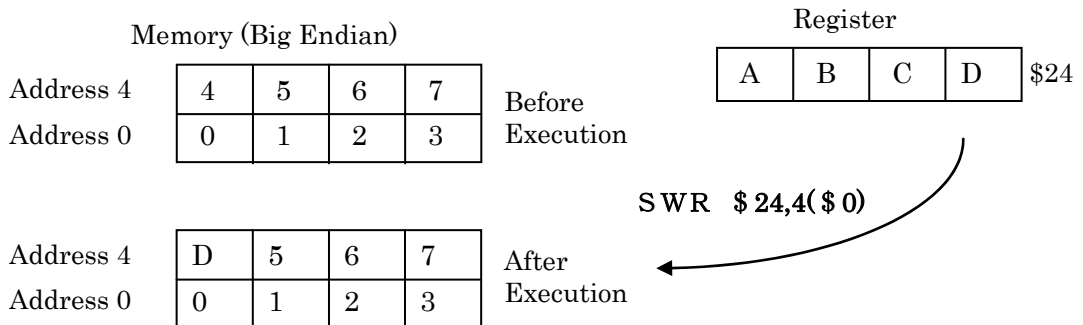
MIPS I

Syntax:

```
swr rt,offset(base)
```

Description:

The 16-bit offset is sign-extended and added to the contents of the base register to generate an address. The contents of register *rt* are shifted left so that the rightmost byte is in the same position in the word as the byte at the generated address. The byte (bytes) that contains the original data of register *rt* is stored in the byte (bytes) from the byte position of the specified address to the word boundary on the left side. Bus error exceptions are not reported because writing is done via a buffer. If a bus error occurs, it is handled by the system as an interrupt.

**Operation:**

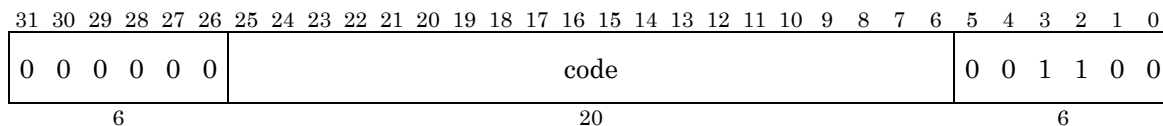
```
vAddr ← sign_extend(offset) + GPR[base]
(pAddr,CCA) ← AddressTranslation(vAddr, DATA, STORE)
if (BigEndianMem=0) then
    pAddr ← pAddrPSIZE-1..2 || 02
endif
byte ← vAddr1..0 xor ReverseEndian2
dataword ← GPR[rt]31-8*byte || 08*byte
StoreMemory(CCA,WORD – byte, dataword, pAddr, vAddr, DATA)
```

Exceptions:

Address Error exception

syscall

System Call



MIPS I

Syntax:

`syscall code`

Description:

A system call exception occurs, immediately transferring control to the exception handling program. The code field is arbitrary.

Operation:

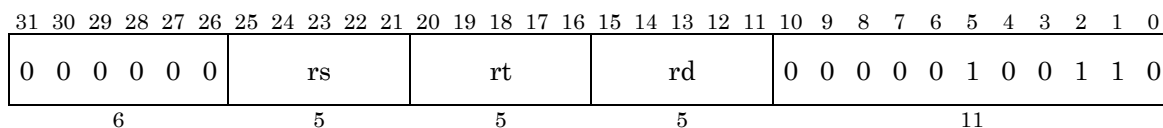
`SignalException(SystemCall)`

Exceptions:

`System Call exception`

xor

XOR



MIPS I

Syntax:

```
xor rd,rs,rt
```

Description:

A bitwise Exclusive OR (exclusive logical sum) is performed for the contents of registers rs and rt and the result is stored in register rd.

X	Y	X XOR Y
0	0	0
0	1	1
1	0	1
1	1	0

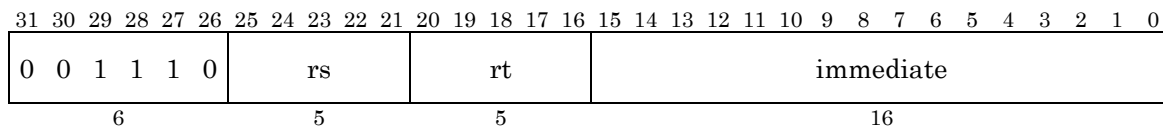
Operation:

$$\text{GPR}[\text{rd}] \leftarrow \text{GPR}[\text{rs}] \text{ xor } \text{GPR}[\text{rt}]$$
Exceptions:

None

xori

Exclusive OR Immediate



MIPS I

Syntax:

```
xori rt,rs,immediate
```

Description:

The 16-bit immediate field is zero-extended and a bitwise Exclusive OR (exclusive logical sum) is performed with the contents of register rs. The result is stored in register rt.

X	Y	X XOR Y
0	0	0
0	1	1
1	0	1
1	1	0

Operation:

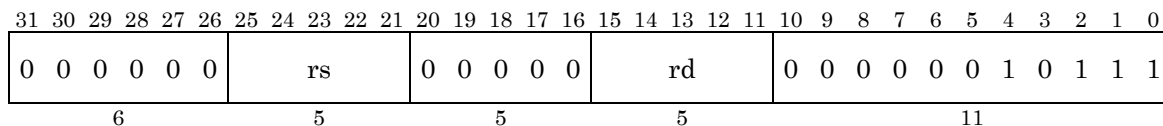
$$\text{GPR}[rt] \leftarrow \text{GPR}[rs] \text{ xor zero_extend}(\text{immediate})$$
Exceptions:

None

ALLEGREX™ Instructions

clo

Count Leading One



ALLEGREX™

Syntax:

```
clo rd, rs
```

Description:

The number of consecutive ones in register *rs* are counted starting from the leftmost bit (MSB). The result (0-32) is stored in register *rd*.

Operation:

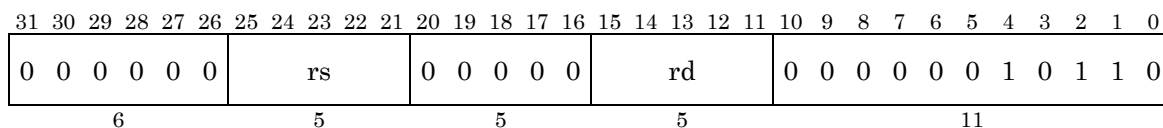
```
temp ← 32
for (i) in(31..0)
  if (GPR[rs]i = 0) then
    temp ← 31 - i
    break
  endif
endfor
GPR[rd] ← temp
```

Exceptions:

None

clz

Count Leading Zero



ALLEGREX™

Syntax:

```
clz rd, rs
```

Description:

The number of consecutive zeroes in register *rs* are counted starting from the leftmost bit (MSB). The result (0-32) is stored in register *rd*.

Operation:

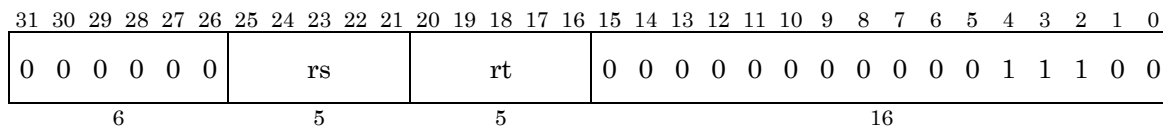
```
temp ← 32
for (i) in(31..0)
  if (GPR[rs]i = 1) then
    temp ← 31 - i
    break
  endif
endfor
GPR[rd] ← temp
```

Exceptions:

None

madd

Multiply Add



ALLEGREX™

Syntax:

```
madd rs,rt
```

Description:

This is a product-sum operation for multiplying the contents of registers *rs* and *rt* as signed integers, adding the 64-bit product to the 64-bit value obtained by concatenating special registers *HI* and *LO*, and updating the special registers *HI* and *LO*.

Operation:

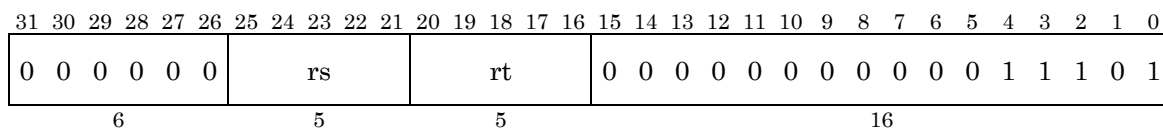
```
temp ← (HI || LO) + (GPR[rs] × GPR[rt])
HI ← temp63..32
LO ← temp31..0
```

Exceptions:

None

maddu

Multiply Add Unsigned



ALLEGREX™

Syntax:

```
maddu rs,rt
```

Description:

This is a product-sum operation for multiplying the contents of registers *rs* and *rt* as unsigned integers, adding the 64-bit product to the 64-bit value obtained by concatenating special registers HI and LO, and updating the special registers HI and LO.

Operation:

$$\text{temp} \leftarrow (\text{HI} \parallel \text{LO}) + (\text{GPR}[\text{rs}] \times \text{GPR}[\text{rt}])$$

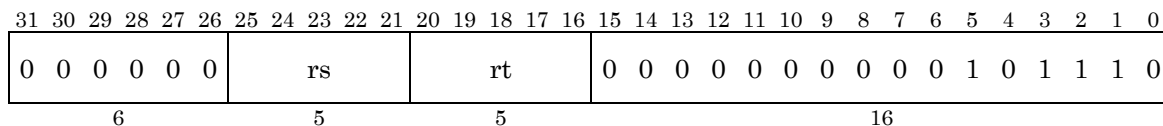
$$\text{HI} \leftarrow \text{temp}_{63..32}$$

$$\text{LO} \leftarrow \text{temp}_{31..0}$$
Exceptions:

None

msub

Multiply Subtract



ALLEGREX™

Syntax:

```
msub rs,rt
```

Description:

This is a product-difference operation for multiplying the contents of registers *rs* and *rt* as signed integers, subtracting the 64-bit product from the 64-bit value obtained by concatenating special registers HI and LO, and updating the special registers HI and LO.

Operation:

$$\text{temp} \leftarrow (\text{HI} \parallel \text{LO}) - (\text{GPR}[\text{rs}] \times \text{GPR}[\text{rt}])$$

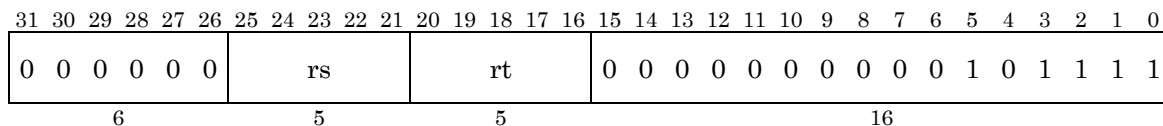
$$\text{HI} \leftarrow \text{temp}_{63..32}$$

$$\text{LO} \leftarrow \text{temp}_{31..0}$$
Exceptions:

None

msubu

Multiply Subtract Unsigned



ALLEGREX™

Syntax:

```
msubu rs,rt
```

Description:

This is a product-difference operation for multiplying the contents of registers *rs* and *rt* as unsigned integers, subtracting the 64-bit product from the 64-bit value obtained by concatenating special registers HI and LO, and updating the special registers HI and LO.

Operation:

$$\text{temp} \leftarrow (\text{HI} \parallel \text{LO}) - (\text{GPR}[\text{rs}] \times \text{GPR}[\text{rt}])$$

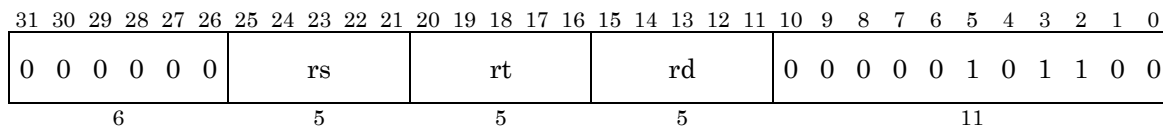
$$\text{HI} \leftarrow \text{temp}_{63..32}$$

$$\text{LO} \leftarrow \text{temp}_{31..0}$$
Exceptions:

None

max

Select Max



ALLEGREX™

Syntax:

```
max rd, rs, rt
```

Description:

The contents of registers *rs* and *rt* are compared as 32-bit signed integers. The contents of the register containing the larger value are stored in register *rd*.

Operation:

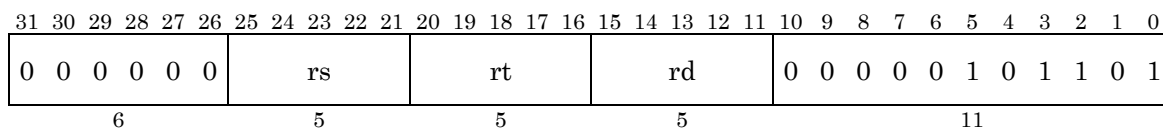
```
if (GPR[rs] < GPR[rt]) then
    GPR[rd] ← GPR[rt]
else
    GPR[rd] ← GPR[rs]
endif
```

Exceptions:

None

min

Select Min



ALLEGREX™

Syntax:

```
min rd, rs, rt
```

Description:

The contents of registers *rs* and *rt* are compared as 32-bit signed integers. The contents of the register containing the smaller value are stored in register *rd*.

Operation:

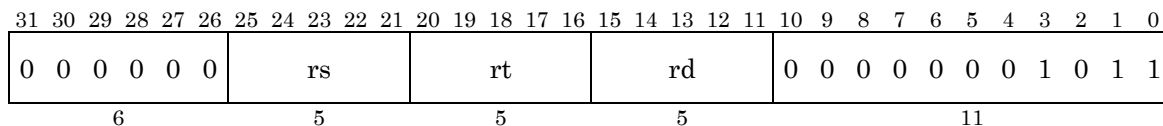
```
if (GPR[rs] < GPR[rt]) then
    GPR[rd] ← GPR[rs]
else
    GPR[rd] ← GPR[rt]
endif
```

Exceptions:

None

movn

Move Conditional on Not Zero



ALLEGREX™

Syntax:

```
movn rd,rs,rt
```

Description:

If the contents of register *rt* are not equal to zero, the contents of register *rs* are copied to register *rd*.

Operation:

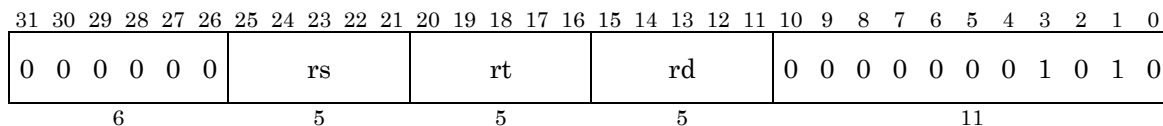
```
if (GPR[rt] ≠ 0) then
    GPR[rd] ← GPR[rs]
endif
```

Exceptions:

None

MOVZ

Move Conditional on Zero



ALLEGREX™

Syntax:

```
movz rd,rs,rt
```

Description:

If the contents of register *rt* are equal to zero, the contents of register *rs* are copied to register *rd*.

Operation:

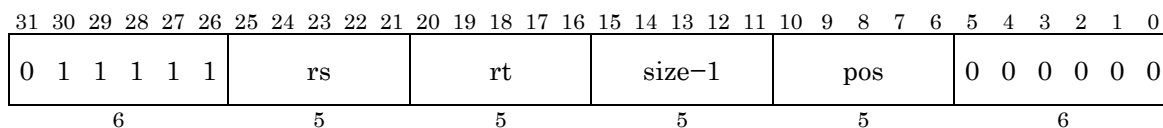
```
if (GPR[rt] = 0) then
    GPR[rd] ← GPR[rs]
endif
```

Exceptions:

None

ext

Extract Bit Field



ALLEGREX™

Syntax:

```
ext rt,rs,pos,size
```

Description:

size bits are extracted from register rs starting at offset bit position pos within the word. The result is stored right-justified in register rt. The high-order bits of register rt are filled with zeros.

Restrictions:

The values of pos and size must be specified so that the sum of pos and size is equal to or less than 32.

Operation:

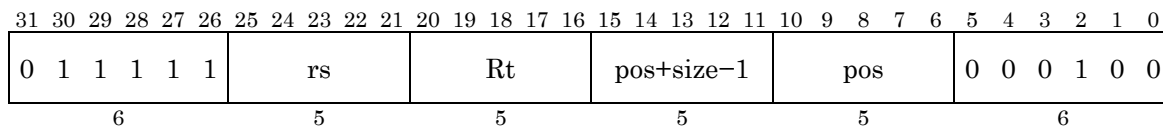
```
if (pos + size > 32) then
    UNPREDICTABLE
endif
temp ← 032-size || GPR[rs](size+pos-1)..pos
GPR[rt] ← temp
```

Exceptions:

None

ins

Insert Bit Field



ALLEGREX™

Syntax:

```
ins rt,rs,pos,size
```

Description:

size bits are extracted from register rs starting from the low-order bit position and inserted into register rt at offset bit position pos within the word.

Restrictions:

The value of size must be specified to be equal to or more than 1.

Operation:

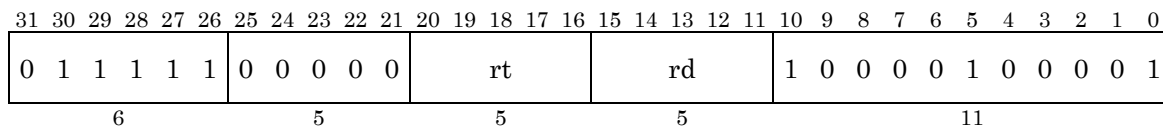
```
if (size < 1) then
    UNPREDICTABLE
endif
GPR[rt] ← GPR[rt]31..(size + pos) || GPR[rs](size-1)..0 || GPR[rt](pos-1)..0
```

Exceptions:

None

seb

Sign-Extend Byte



ALLEGREX™

Syntax:

```
seb rd, rt
```

Description:

The lowest byte of register *rt* is sign-extended to 32 bits. The result is stored in register *rd*.

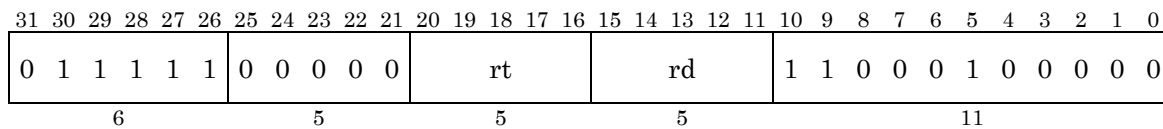
Operation:

$$\text{GPR}[\text{rd}] \leftarrow \text{sign_extend}(\text{GPR}[\text{rt}]_{7..0})$$
Exceptions:

None

seh

Sign-Extend Halfword



ALLEGREX™

Syntax:

seh rd, rt

Description:

The low-order halfword of register *rt* is sign-extended to 32 bits. The result is stored in register *rd*.

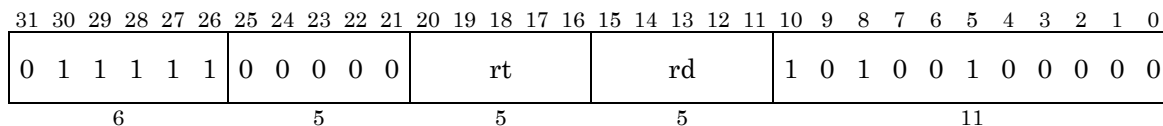
Operation:

$$\text{GPR}[\text{rd}] \leftarrow \text{sign_extend}(\text{GPR}[\text{rt}]_{15..0})$$
Exceptions:

None

bitrev

Bit Reverse



ALLEGREX™

Syntax:

```
bitrev rd, rt
```

Description:

The contents of register `rt` are swapped bit-for-bit within the word. The 32-bit result is stored in register `rd`.

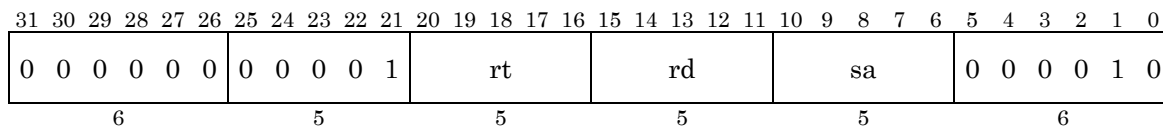
Operation:

$$\text{GPR}[\text{rd}] \leftarrow (\text{GPR}[\text{rt}]_0 \parallel \text{GPR}[\text{rt}]_1 \parallel \dots \parallel \text{GPR}[\text{rt}]_{31})$$
Exceptions:

None

rotr

Rotate Word Right



ALLEGREX™

Syntax:

```
rotr rd, rt, sa
```

Description:

The contents of register *rt* are rotated right by right-shifting them *sa* bits. The 32-bit result is stored in register *rd*.

Operation:

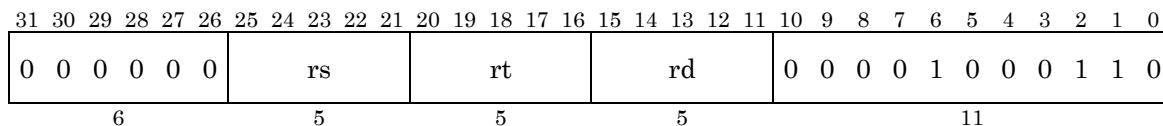
```
s ← sa
temp ← (GPR[rt]s-1..0 || GPR[rt]31..s)
GPR[rd] ← temp
```

Exceptions:

None

rotrv

Rotate Word Right Variable



ALLEGREX™

Syntax:

```
rotrv rd, rt, rs
```

Description:

The contents of register *rt* are rotated right by performing a right shift. The shift amount is specified by the low-order 5 bits of register *rs*.

The 32-bit result is stored in register *rd*.

Operation:

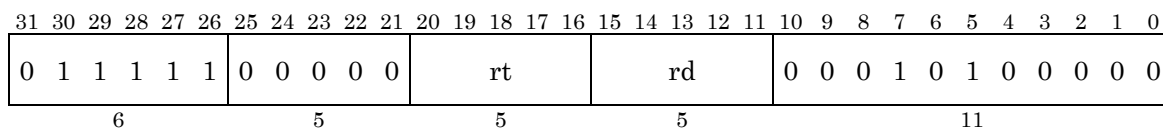
```
s ← GPR[rs]4..0
temp ← (GPR[rt]s-1..0 || GPR[rt]31..s)
GPR[rd] ← temp
```

Exceptions:

None

wsbh

Word Swap Bytes within Halfword



ALLEGREX™

Syntax:

wsbh rd, rt

Description:

The contents of register *rt* are swapped byte-for-byte within each halfword of the register. The 32-bit result is stored in register *rd*.

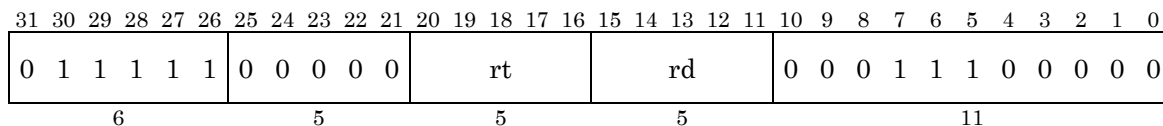
Operation:

$$\text{GPR}[\text{rd}] \leftarrow (\text{GPR}[\text{rt}]_{23..16} \parallel \text{GPR}[\text{rt}]_{31..24} \parallel \text{GPR}[\text{rt}]_{7..0} \parallel \text{GPR}[\text{rt}]_{15..8})$$
Exceptions:

None

wsbw

Word Swap Bytes within Word



ALLEGREX™

Syntax:

wsbw rd, rt

Description:

The contents of register *rt* are swapped byte-for-byte within the word. The 32-bit result is stored in register *rd*.

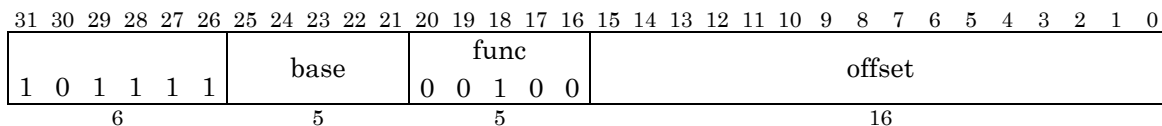
Operation:

$$\text{GPR}[\text{rd}] \leftarrow (\text{GPR}[\text{rt}]_{7..0} \parallel \text{GPR}[\text{rt}]_{15..8} \parallel \text{GPR}[\text{rt}]_{23..16} \parallel \text{GPR}[\text{rt}]_{31..24})$$
Exceptions:

None

cache : Index Invalidate (I)

Index Invalidate (I)



ALLEGREX™

Syntax:

```
cache func(=0x04), offset(base)
```

Description:

This instruction is used for clearing a line in the instruction cache.

The 16-bit offset is sign-extended and added to the contents of the base register to generate the address of a cache line. The cache operation indicated by the 5-bit func code is performed on that cache line.

The generated address is used to obtain an index that points to a pair of entries in the cache. The appropriate entry, either WayA or WayB, is selected and invalidated (the Valid bit is cleared to 0) according to the following rules.

- 1) When both WayA and WayB are valid: Invalidate the oldest entry as determined by the LRU value
- 2) When either WayA or WayB is valid, but not both: Invalidate the valid entry and clear the Lock bit
- 3) When both WayA and WayB are invalid: Do nothing

To clear both entries WayA and WayB, issue this instruction twice in a row to the same address (index).

This instruction can be executed in user mode.

Operation:

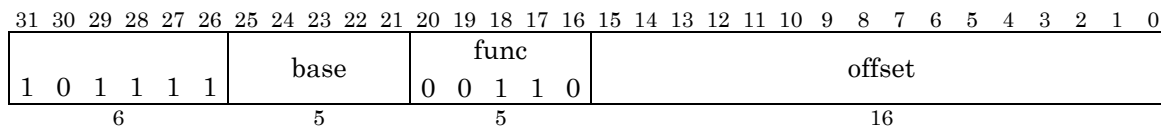
```
vAddr ← GPR[base] + sign_extend(offset)
CacheOp(func, vAddr)
```

Exceptions:

```
Address Error exception
```

cache : Index Unlock (I)

Index Unlock (I)



ALLEGREX™

Syntax:

```
cache func(=0x06), offset(base)
```

Description:

This instruction is used for canceling the lock on an instruction cache line.

The 16-bit offset is sign-extended and added to the contents of the base register to generate the address of a cache line. The cache operation indicated by the 5-bit func code is performed on that cache line.

The generated address is used to obtain an index that points to a cache line. The Lock bit for that cache line is cleared (the lock on the cache line is canceled).

This instruction can be executed in user mode.

Operation:

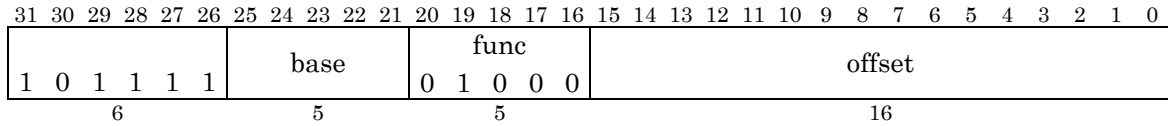
```
vAddr ← GPR[base] + sign_extend(offset)
CacheOp(func, vAddr)
```

Exceptions:

```
Address Error exception
```

cache : Hit Invalidate (I)

Hit Invalidate (I)



ALLEGREX™

Syntax:

```
cache func (=0x08) , offset (base)
```

Description:

This instruction is used to invalidate a specific line in the instruction cache.

The 16-bit offset is sign-extended and added to the contents of the base register to generate the address of a cache line. The cache operation indicated by the 5-bit func code is performed on that cache line.

The tag for the cache index is obtained from the generated address. If the tag indicates that the addressed line is present in the cache (a cache hit), its entry is invalidated (its Valid bit is cleared). If the Lock bit in the tag is 1 and the LRU bit points to the other entry (in other words, the target entry is locked), the Lock bit is cleared (the lock is canceled).

If the addressed line is not present in the cache, no operation is performed.

This instruction can be executed in user mode.

Operation:

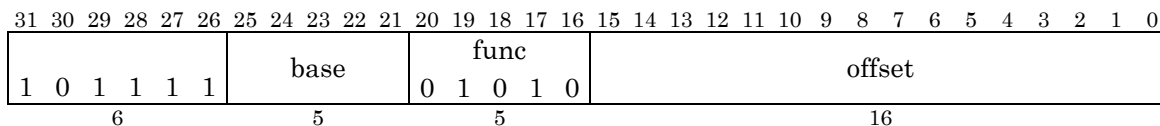
```
vAddr ← GPR[base] + sign_extend(offset)
CacheOp(func, vAddr)
```

Exceptions:

```
Address Error exception
```

cache : Fill (I)

Fill (I)



ALLEGREX™

Syntax:

```
cache func (=0x0a) , offset (base)
```

Description:

This instruction is used for explicitly filling a specific line in the instruction cache, if the line is not in the cache.

The 16-bit offset is sign-extended and added to the contents of the base register to generate the address of a cache line. The cache operation indicated by the 5-bit func code is performed on that cache line.

If the cache line at the generated address is not in the cache (a cache miss), the cache line is filled. If the line is already in the cache (a cache hit), no operation is performed. If the Lock bit = 0 (not locked), the LRU bit is updated to point to the way that is not being filled.

During the instruction cache fill, the pipeline is interlocked so no other processing can be performed in parallel.

This instruction can be executed in user mode.

Operation:

```
vAddr ← GPR[base] + sign_extend(offset)
CacheOp(func, vAddr)
```

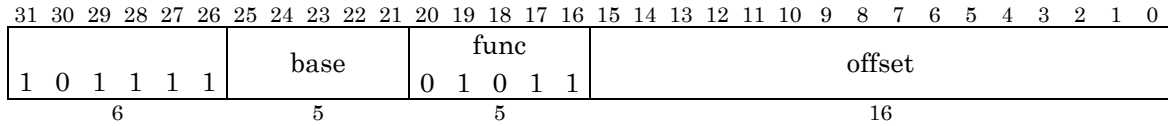
Exceptions:

```
Address Error exception
```

```
Bus Error exception
```

cache : Fill with Lock (I)

Fill with Lock (I)



ALLEGREX™

Syntax:

```
cache func (=0x0b) , offset (base)
```

Description:

This instruction is used for explicitly filling and locking a specific line in the instruction cache, if the line is not in the cache.

The 16-bit offset is sign-extended and added to the contents of the base register to generate the address of a cache line. The cache operation indicated by the 5-bit func code is performed on that cache line.

If the cache line at the generated address is not in the cache (a cache miss), the cache line is filled. If the line is already in the cache (a cache hit), no operation is performed. If the Lock bit = 0 (not locked), the LRU bit is updated to point to the way that is not being filled.

This instruction also sets the Lock bit. As a result, the LRU bit is held, and the target line will be locked and not be subject to replacement.

During the instruction cache fill, the pipeline is interlocked so no other processing can be performed in parallel.

This instruction can be executed in user mode.

Operation:

```
vAddr ← GPR[base] + sign_extend(offset)
CacheOp(func, vAddr)
```

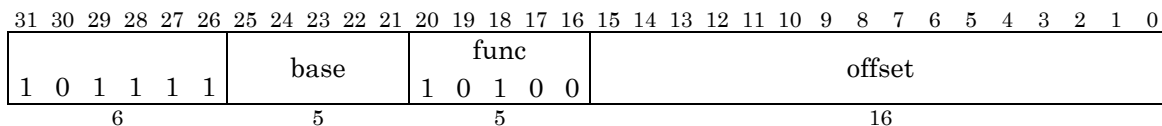
Exceptions:

```
Address Error exception
```

```
Bus Error exception
```

cache : Index Writeback Invalidate (D)

Index Writeback Invalidate (D)



ALLEGREX™

Syntax:

```
cache func (=0x14) , offset (base)
```

Description:

This instruction is used to flush and clear a line in the data cache.

The 16-bit offset is sign-extended and added to the contents of the base register to generate the address of a cache line. The cache operation indicated by the 5-bit func code is performed on that cache line.

The generated address is used to obtain an index that points to a pair of entries in the cache. The appropriate entry, either WayA or WayB, is selected according to the rules listed below, and if the Dirty bit in the entry is set, the line is written back to main memory and the entry is invalidated (the Valid bit is cleared to 0). If the Dirty bit is not set, the entry is only invalidated and no writeback is performed.

- 1) When both WayA and WayB are valid: Select the oldest as determined by the value of the LRU bit
- 2) When either WayA or WayB is valid, but not both: Select the valid entry and clear the Lock bit
- 3) When both WayA and WayB are invalid: Do nothing.

To flush and clear both entries WayA and WayB, issue this instruction twice in a row to the same address (index).

This instruction can be executed in user mode.

Operation:

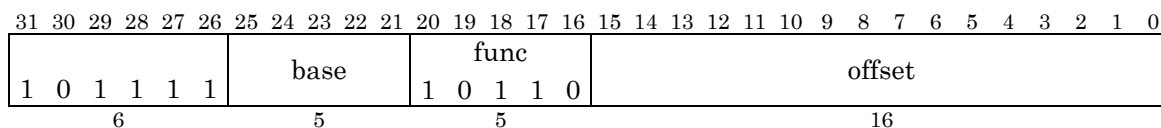
```
vAddr ← GPR[base] + sign_extend(offset)
CacheOp(func, vAddr)
```

Exceptions:

```
Address Error exception
```


cache : Index Unlock (D)

Index Unlock (D)



ALLEGREX™

Syntax:

```
cache func(=0x16), offset(base)
```

Description:

This instruction is used to unlock a line in the data cache.

The 16-bit offset is sign-extended and added to the contents of the base register to generate the address of a cache line. The cache operation indicated by the 5-bit func code is performed on that cache line.

The generated address is used to obtain an index that points to a cache line. The Lock bit for that cache line is cleared (the lock on the cache line is canceled).

This instruction can be executed in user mode.

Operation:

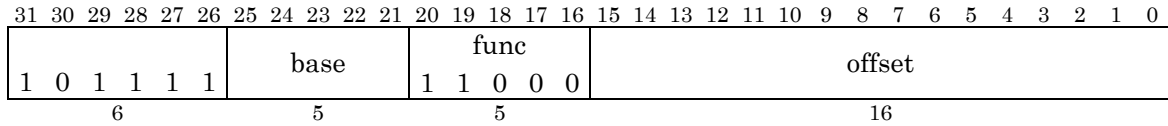
```
vAddr ← GPR[base] + sign_extend(offset)
CacheOp(func, vAddr)
```

Exceptions:

```
Address Error exception
```

cache : Create Dirty Exclusive (D)

Create Dirty Exclusive (D)



ALLEGREX™

Syntax:

```
cache func(=0x18), offset(base)
```

Description:

This instruction is used to create a dirty line in the data cache. It prevents unnecessary cache fills from being performed on lines that are only going to be fully written.

The 16-bit offset is sign-extended and added to the contents of the base register to generate the address of a cache line. The cache operation indicated by the 5-bit func code is performed on that cache line.

An entry in the data cache is created for the line at the generated address, in the data cache tag. At the same time, the Dirty bit is set in the tag. This instruction does not perform a data fill.

If all the words in the cache line are written right after this instruction is executed, the cache line can be completed without having to perform a wasteful cache fill. If the Lock bit = 0 (not locked), the LRU bit will be updated to point to the way that was not created. This instruction can be executed in user mode.

Specifying a non-cache address with the Create Dirty Exclusive (D) instruction will cause indeterminate behavior and is thus prohibited.

Operation:

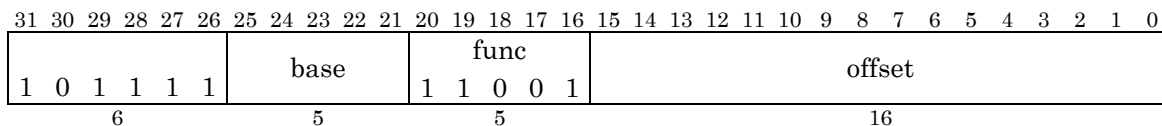
```
vAddr ← GPR[base] + sign_extend(offset)
CacheOp(func, vAddr)
```

Exceptions:

```
Address Error exception
```

cache : Hit Invalidate (D)

Hit Invalidate (D)



ALLEGREX™

Syntax:

```
cache func (=0x19) , offset (base)
```

Description:

This instruction is used to invalidate a specific line in the data cache.

The 16-bit offset is sign-extended and added to the contents of the base register to generate the address of a cache line. The cache operation indicated by the 5-bit func code is performed on that cache line.

The tag for the cache index is obtained from the generated address. If the tag indicates that the addressed line is present in the cache (a cache hit), its entry is invalidated (its Valid bit is cleared). If the Lock bit in the tag is 1 and the LRU bit points to the other entry (in other words, the target entry is locked), the Lock bit is cleared (the lock is canceled).

If the addressed line is not present in the cache, no operation is performed.

This instruction can be executed in user mode.

Operation:

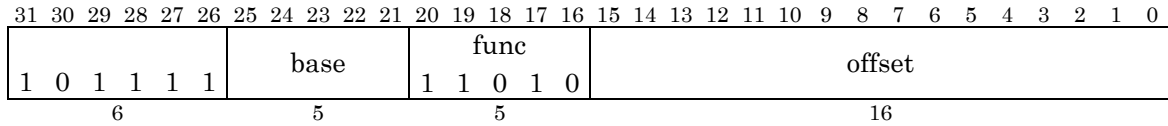
```
vAddr ← GPR[base] + sign_extend(offset)
CacheOp(func, vAddr)
```

Exceptions:

```
Address Error exception
```

cache : Hit WriteBack (D)

Hit WriteBack (D)



ALLEGREX™

Syntax:

```
cache func(=0x1a), offset(base)
```

Description:

This instruction is used to perform an explicit writeback of a dirty line in the data cache. It is used to explicitly synchronize data in the cache with data in main memory.

The 16-bit offset is sign-extended and added to the contents of the base register to generate the address of a cache line. The cache operation indicated by the 5-bit func code is performed on that cache line.

The tag for the cache index is obtained from the generated address. If the tag indicates that the addressed line is present in the cache (a cache hit), and if its Dirty bit is set, a writeback is performed on that cache line and the Dirty bit is cleared.

If the addressed line is not present in the cache, no operation is performed.

This instruction can be executed in user mode.

Operation:

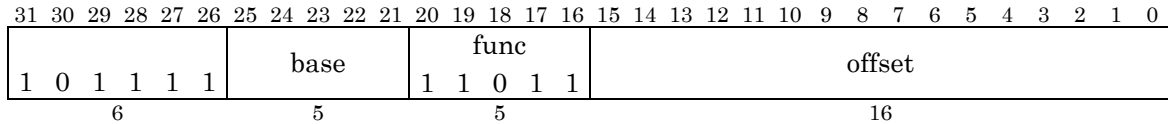
```
vAddr ← GPR[base] + sign_extend(offset)
CacheOp(func, vAddr)
```

Exceptions:

```
Address Error exception
```

cache : Hit WriteBack Invalidate(D)

Hit WriteBack Invalidate(D)



ALLEGREX™

Syntax:

```
cache func(=0x1b), offset(base)
```

Description:

This instruction is used to perform an explicit writeback of a dirty line in the data cache. It is used to explicitly synchronize data in the cache with data in main memory. After the writeback is performed, the line is invalidated.

The 16-bit offset is sign-extended and added to the contents of the base register to generate the address of a cache line. The cache operation indicated by the 5-bit func code is performed on that cache line.

The tag for the cache index is obtained from the generated address. If the tag indicates that the addressed line is present in the cache (a cache hit), and if its Dirty bit is set, a writeback is performed on that cache line and the line is invalidated. If the Dirty bit is not set, the line will only be invalidated. If the Lock bit in the tag is 1 and the LRU bit points to the other entry (in other words, the target entry is locked), the Lock bit is cleared (the lock is canceled).

If the addressed line is not present in the cache, no operation is performed.

This instruction can be executed in user mode.

Operation:

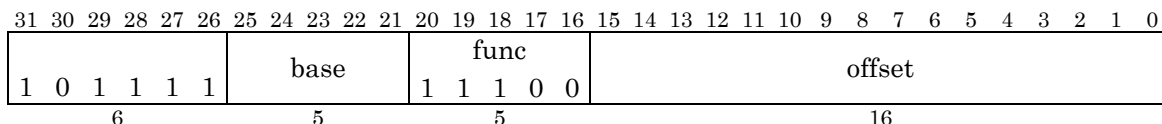
```
vAddr ← GPR[base] + sign_extend(offset)
CacheOp(func, vAddr)
```

Exceptions:

```
Address Error exception
```

cache : Create Dirty Exclusive with Lock (D)

Create Dirty Exclusive with Lock (D)



ALLEGREX™

Syntax:

```
cache func(=0x1c), offset(base)
```

Description:

This instruction is used to create a dirty line in the data cache and lock it. It prevents unnecessary cache fills from being performed on lines that are only going to be fully written.

The 16-bit offset is sign-extended and added to the contents of the base register to generate the address of a cache line. The cache operation indicated by the 5-bit func code is performed on that cache line.

An entry in the data cache is created for the line at the generated address, in the data cache tag. At the same time, the Dirty bit is set in the tag. This instruction does not perform a data fill.

If all the words in the cache line are written right after this instruction is executed, the cache line can be completed without having to perform a wasteful cache fill. If the Lock bit = 0 (not locked), the LRU bit will be updated to point to the way that was not created. This instruction also sets the Lock bit. As a result, the LRU bit is held, and the target line will be locked and not be subject to replacement.

This instruction can be executed in user mode.

Specifying a non-cache address with the Create Dirty Exclusive with Lock(D) instruction will cause indeterminate behavior and is thus prohibited.

Operation:

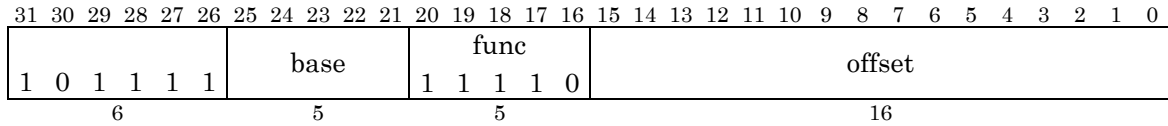
```
vAddr ← GPR[base] + sign_extend(offset)
CacheOp(func, vAddr)
```

Exceptions:

```
Address Error exception
```

cache : Fill (D)

Fill (D)



ALLEGREX™

Syntax:

```
cache func(=0x1e), offset(base)
```

Description:

This instruction is used for explicitly filling a specific line in the data cache, if the line is not in the cache.

The 16-bit offset is sign-extended and added to the contents of the base register to generate the address of a cache line. The cache operation indicated by the 5-bit func code is performed on that cache line.

If the cache line at the generated address is not in the cache (a cache miss), the cache line is filled. If the line is already in the cache (a cache hit), no operation is performed. If the Lock bit = 0 (not locked), the LRU bit is updated to point to the way that is not being filled.

When the instruction cache is being filled, the pipeline is interlocked so no other processing can be performed in parallel. However, when a data cache is being filled, as long as the following instruction is not a load/store or cache instruction, that instruction can be executed in parallel with the data cache fill (the data cache fill is a non-blocking operation).

This instruction can be executed in user mode.

Operation:

```
vAddr ← GPR[base] + sign_extend(offset)
CacheOp(func, vAddr)
```

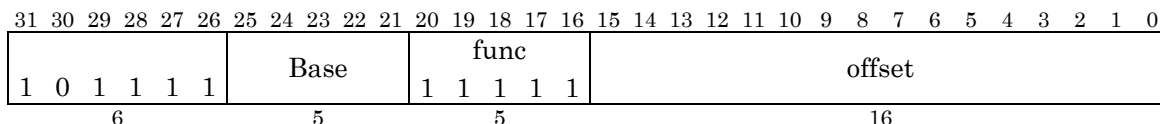
Exceptions:

Address Error exception

Bus Error exception

cache : Fill with Lock (D)

Fill with Lock (D)



ALLEGREX™

Syntax:

```
cache func(=0x1f), offset(base)
```

Description:

This instruction is used for explicitly filling a specific line in the data cache and locking it, if the line is not in the cache.

The 16-bit offset is sign-extended and added to the contents of the base register to generate the address of a cache line. The cache operation indicated by the 5-bit func code is performed on that cache line.

If the cache line at the generated address is not in the cache (a cache miss), the cache line is filled. If the line is already in the cache (a cache hit), no operation is performed. If the Lock bit = 0 (not locked), the LRU bit is updated to point to the way that is not being filled.

This instruction also sets the Lock bit. As a result, the LRU bit is held, and the target line will be locked and not be subject to replacement.

When the instruction cache is being filled, the pipeline is interlocked so no other processing can be performed in parallel. However, when a data cache is being filled, as long as the following instruction is not a load/store or cache instruction, that instruction can be executed in parallel with the data cache fill (the data cache fill is a non-blocking operation).

This instruction can be executed in user mode.

Operation:

```
vAddr ← GPR[base] + sign_extend(offset)
CacheOp(func, vAddr)
```

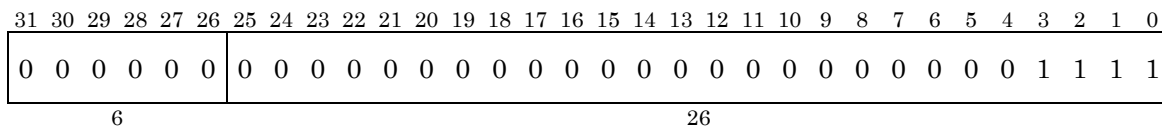
Exceptions:

```
Address Error exception
```


Bus Error exception

sync

Synchronize Shared Memory



ALLEGREX™

Syntax:

sync

Description:

The pipeline is stalled until the CPU's cache writeback buffer and non-cache write buffer have emptied. This is done to block the execution of subsequent instructions.

Operation:

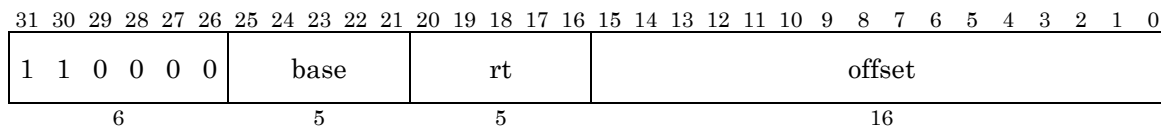
SyncOperation()

Exceptions:

None

II

Load Linked



ALLEGREX™

Syntax:

```
ll rt, offset(base)
```

Description:

The 16-bit offset is sign-extended and added to the contents of the base register to generate an address.

The word in memory at that address is loaded into register *rt* and the LLbit is set to 1.

Operation:

```
vAddr ← sign_extend(offset) + GPR[base]
if (vAddr1..0 ≠ 02) then
    SignalException(AddressError)
endif
(pAddr, CCA) ← AddressTranslation(vAddr, DATA, LOAD)
memword ← LoadMemory(CCA, WORD, pAddr, vAddr, DATA)
GPR[rt] ← memword
LLbit ← 1
```

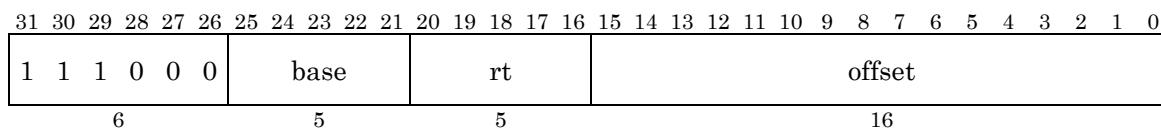
Exceptions:

Address Error exception

Bus Error exception

SC

Store Conditional



ALLEGREX™

Syntax:

```
sc rt, offset(base)
```

Description:

The 16-bit offset is sign-extended and added to the contents of the base register to generate an address.

If the LLbit is 1, the contents of register *rt* are stored in memory at that address and 1 is returned in register *rt*.

If the LLbit is 0, no store is performed and 0 is returned in register *rt*.

Bus error exceptions are not reported because writing is done via a buffer. If a bus error occurs, it is handled by the system as an interrupt.

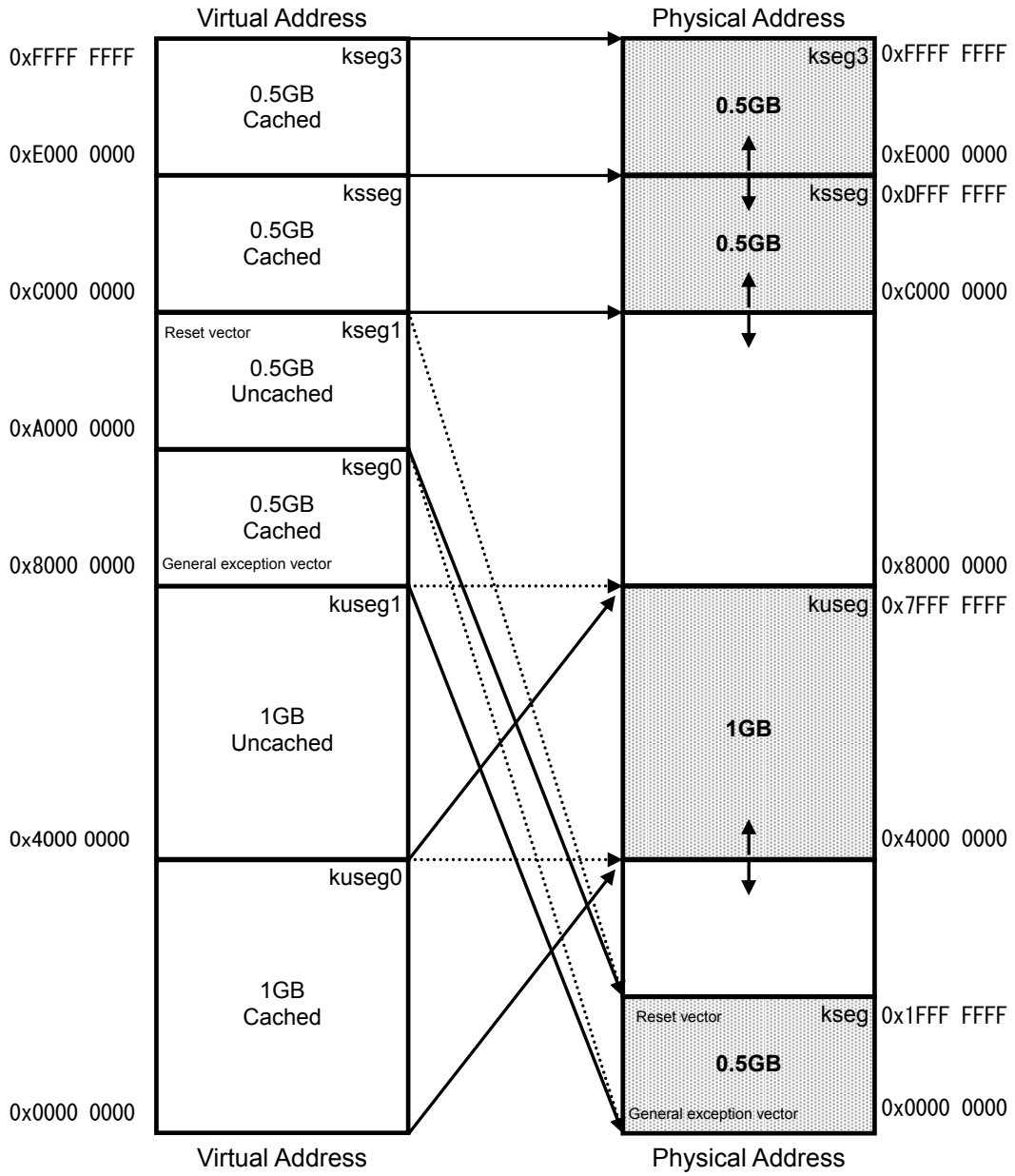
Operation:

```
vAddr ← sign_extend(offset) + GPR[base]
if (vAddr1..0 ≠ 02) then
    SignalException(AddressError)
endif
(pAddr, CCA) ← AddressTranslation(vAddr, DATA, STORE)
dataword ← GPR[rt]
if (LLbit) then
    StoreMemory(CCA, WORD, dataword, pAddr, vAddr, DATA)
endif
GPR[rt] ← (031 || LLbit)
```

Exceptions:

```
Address Error exception
```

Virtual—Physical Address Map



CACHE Instruction func Code Table

func Code	Cache	Instruction	Available Modes
0	I	—	—
1	I	—	—
2	I	—	—
3	I	—	—
4	I	Index Invalidate	u / s/ k
5	I	—	—
6	I	Index Unlock	u / s/ k
7	I	—	—
8	I	Hit Invalidate	u / s/ k
9	I	—	—
10	I	Fill	u / s/ k
11	I	Fill with Lock	u / s/ k
12	I	—	—
13	I	—	—
14	I	—	—
15	I	—	—
16	D	—	—
17	D	—	—
18	D	—	—
19	D	—	—
20	D	Index Writeback Invalidate	u / s/ k
21	D	—	—
22	D	Index Unlock	u / s/ k
23	D	—	—
24	D	Create Dirty Exclusive	u / s/ k
25	D	Hit Invalidate	u / s/ k
26	D	Hit Writeback	u / s/ k
27	D	Hit Writeback Invalidate	u / s/ k
28	D	Create Dirty Exclusive with Lock	u / s/ k
29	D	—	—
30	D	Fill	u / s/ k
31	D	Fill with Lock	u / s/ k

u: User mode, s: Supervisor mode, k: Kernel mode

CPU Instruction Set Code Table

Opcode		Bit 28-26							
Bit 31-29		0	1	2	3	4	5	6	7
0		SPECIAL	REGIMM	J	JAL	BFQ	BNF	BLEZ	BGTZ
1		ADDI	ADDIU	SLTI	SLTIU	ANDI	ORI	XORI	LUI
2		COP0	COP1	COP2	COP3	BRQI	BNEL	BLEZL	BGTZL
3		VFPU0	VFPU1	*	VFPU3	SPECIAL2	*	*	SPECIAL3
4		LB	LH	LWL	LW	LRU	LHU	LWR	*
5		SB	SH	SWL	SW	*	*	SWR	CACHE
6		LL	LWC1	LWC2	*	VFPU4	LQUC2	LQC2	VFPU5
7		SC	SWC1	SWC2	*	VFPU6	SQUC2	SQC2	VFPU7
SPECIAL func		Bit 2-0							
Bit 5-3		0	1	2	3	4	5	6	7
0		SLL		SRL/ROTR	SRA	SLLV		SRLV/ROTRV	SRAV
1		JR	JALR	MOVZ	MOVN	SYSCALL	BREAK	*	SYNC
2		MFHI	MTHI	MFLO	MTLO	*	*	CLZ	CLO
3		MULT	MUL TU	DIV	DIVU	MADD	MADDU	*	*
4		ADD	ADDU	SUB	SUBU	AND	OR	XOR	NOR
5		*	*	SLT	SLTU	MAX	MIN	MSUB	MSUBU
6		*	*	*	*	*	*	*	*
7		*	*	*	*	*	*	*	*
SPECIAL rs		Bit 23-21							
Bit 25-24		0	1	2	3	4	5	6	7
0		SRL	ROTR						
1									
2									
3									
SPECIAL sa		Bit 8-6							
Bit 10-9		0	1	2	3	4	5	6	7
0		SRLV	ROTRV						
1									
2									
3									
REGIMM sa		Bit 18-16							
Bit 20-19		0	1	2	3	4	5	6	7
0		BLTZ	RGEZ	BLTZL	RGEZL	*	*	*	*
1		*	*	*	*	*	*	*	*
2		BLTZAL	RGEZAL	BLTZALL	RGEZALL	*	*	*	*
3		*	*	*	*	*	*	*	*
COPz rs		Bit 23-21							
Bit 25-24		0	1	2	3	4	5	6	7
0		MFCz		GFCz	MFHCz	MTCz		CTCz	MTHCz
1		BCz							
2		COPz							
3		COPz							
COPz rt (BC)		Bit 18-16							
Bit 20-19		0	1	2	3	4	5	6	7
0		BCzF	BCzT	BCzFL	BCzTL				
1									
2									
3									
COP0 rs		Bit 23-21							
Bit 25-24		0	1	2	3	4	5	6	7
0		MFC0				MTC0			
1		BC							
2		COPz							
3		COPz							
COP0 func		Bit 2-0							
bit 4-3		0	1	2	3	4	5	6	7
0		*	*	*				*	
1		*							
2		*							
3		FRET							
SPECIAL2 func		Bit 2-0							
Bit 5-3		0	1	2	3	4	5	6	7
0									
1									
2									
3									
4									
5									
6									
7									
SPECIAL2 rs		Bit 23-21							
Bit 25-24		0	1	2	3	4	5	6	7
0									
1									
2									
3									
SPECIAL3 func		Bit 2-0							
Bit 5-3		0	1	2	3	4	5	6	7
0		EXT				INS			
1									
2									
3		BSHFL							
4									
5									
6									
7									
BSHFL sa		Bit 8-6							
Bit 10-9		0	1	2	3	4	5	6	7
0				WSBH	WSBW				
1									
2		SEB				BITREV			
3		SEH							

MIPS-I FPU
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